

Computer Architecture
Machine Representation of Programs: Control

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Conditional Branches

Equivalent to the ‘if’ statement that computers use. There are many types of conditional branches:

Jumps Known as *jX* instructions, these make the program jump to different parts of code depending on condition codes.

jX	Condition	Description
jmp	1	Unconditional
jbe	ZF	Equal / Zero
jnb	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~ (SF^OF) & ~ZF	Greater (Signed)
jge	~ (SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

An example of an ‘Old Sytle’ Conditional Branch

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
absdiff:
    cmpq    %rsi, %rdi    # x:y
    jle     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax
    ret
.L4:      # x <= y
    movq    %rsi, %rax
    subq    %rdi, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

This is old style because The GoTo Function was implemented in C which allows for the following:

C Code

```
val = Test ? Then_Expr : Else_Expr;
```

```
val = x>y ? x-y : y-x;
```

GoTo Version

```
ntest = !Test;
if (ntest) goto Else;
val = Then_Expr;
goto Done;
Else:
val = Else_Expr;
Done:
. . .
```

- Create separate code regions for then & else expressions
- Execute appropriate one

```
long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
long absdiff_j
(long x, long y)
{
    long result;
    int ntest = x <= y;
    if (ntest) goto Else;
    result = x-y;
    goto Done;
Else:
    result = y-x;
Done:
    return result;
}
```

The above programs are the old C code on the left compared to the New GoTo C code on the right. This makes the C code more functionally similar to the Assembly code. But because it is a hardcoded jump rather than a branch (like an if statement) the computer cannot pipeline (predict the program) efficiently which is a drain on performance.

Conditional Move

The following is an example of a Conditional Move (which is preferable to Goto):

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```

long absdiff
(long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}

absdiff:
    movq    %rdi, %rax    # x
    subq    %rsi, %rax    # result = x-y
    movq    %rsi, %rdx
    subq    %rdi, %rdx    # eval = y-x
    cmpq    %rsi, %rdi    # x:y
    cmovle  %rdx, %rax    # if <=, result = eval
    ret
  
```

Note that this is a different representation of the example on page 2 of this note which uses conditional moves rather than jumps.

Bad Cases There are a few cases where a conditional jump is not good to use. First, if the computation for the if statement is very complex the computational benefit for using a conditional statement is outweighed by the computational load used to compute the conditions.

```
val = Test(x) ? HardComp1(x) : HardComp2(x);.
```

When a computation is risky, a conditional branch may cause undefined or undesirable behaviour.

```
val = p ? *p : 0;.
```

Finally, if the computations in each branch have ‘side effects’ then a conditional move should not be used.

```
val = x > 0 ? x*=7 : x+=3;.
```

This is because both of the values get computed and causes more computational resources to be used.

Conditional Branch - if statement summary

we now know that we can use two different methods of implementing *if* statements in code. They are:

- conditional Jumps
- conditional Moves

Loops

Loops can be implemented with simply with an *if* statement. in C code, there are 3 categories of loops.

Do While

The first is the ‘Do While’ loop which does something while it is true. In C code this can be represented with either an *if* statement or using the Goto function:

C Code

```
long pcount_do
(unsigned long x) {
    long result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

Goto Version

```
long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

These are both equivalent; however the goto version makes the code unnecessarily complicated. The assembly code for this loop is:

Goto Version

```

long pcount_goto
(unsigned long x) {
    long result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}

```

Register	Use(s)
%rdi	Argument <i>x</i>
%rax	result

```

    movl    $0, %eax    # result = 0
.L2:
    movq    %rdi, %rdx  # loop:
    andl    $1, %edx    # t = x & 0x1
    addq    %rdx, %rax  # result += t
    shrq    %rdi        # x >>= 1
    jne     .L2         # if (x) goto loop
rep; ret

```

Below is the general form for **do while** loops in C:

C Code

```

do
    Body
while (Test) ;

```

```

■ Body: {
    Statement1;
    Statement2;
    ...
    Statementn;
}

```

Goto Version

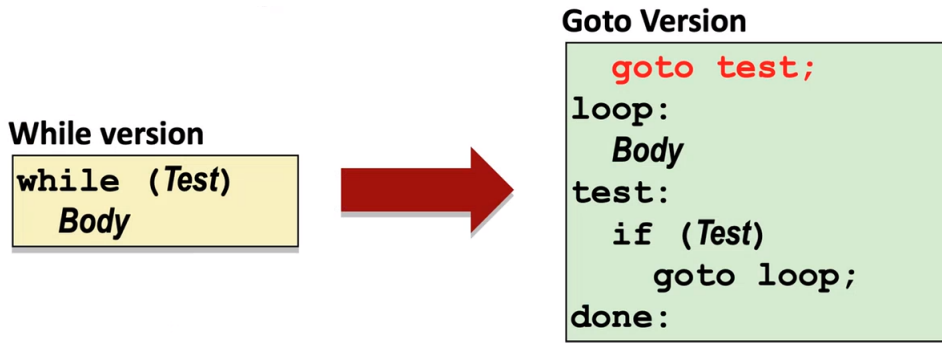
```

loop:
    Body
    if (Test)
        goto loop

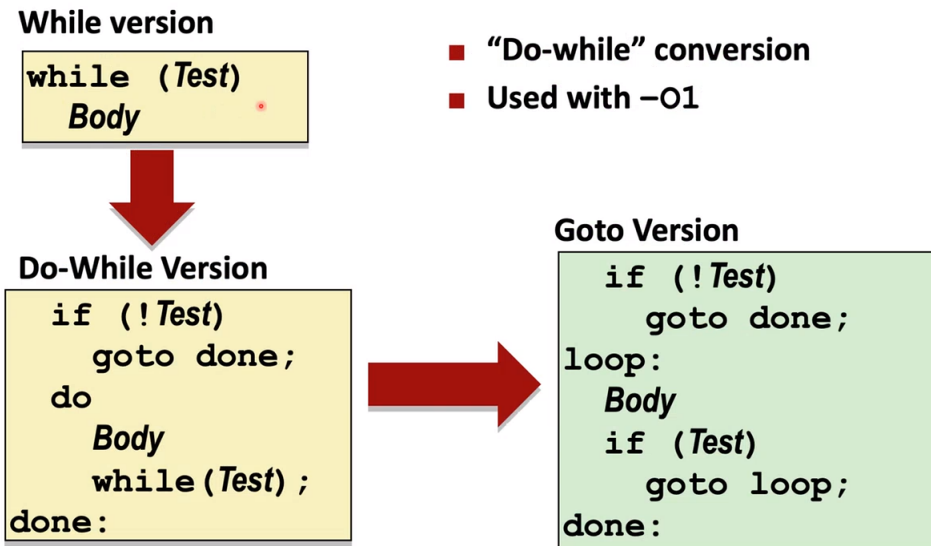
```

While

While loops are similar to do-while for obvious reasons. The general form for a **while** loop in C is:



We can also generalize the while loop in terms of the do while loop like so:



Using this generalization we can reform a while loop to be a ‘do while’ for loop:

C Code

```
long pcount_while
(unsigned long x) {
    long result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Do-While Version

```
long pcount_goto_dw
(unsigned long x) {
    long result = 0;
    if (!x) goto done;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
done:
    return result;
} *
```

For

The for loop has the general form in C (and most other languages) is :

```
for (i = InitVal; i < Test; i++).
```

To convert a for loop to a ‘for while’ loop the general form is:

<p>Init</p> <pre>i = 0</pre> <p>Test</p> <pre>i < WSIZE</pre> <p>Update</p> <pre>i++</pre> <p>Body</p> <pre>{ unsigned bit = (x >> i) & 0x1; result += bit; }</pre>	<pre>long pcount_for_while (unsigned long x) { size_t i; long result = 0; i = 0; while (i < WSIZE) { unsigned bit = (x >> i) & 0x1; result += bit; i++; } return result; }</pre>
--	---

We can also convert from for loop to a do while loop like so:

C Code	Goto Version
<pre>long pcount_for (unsigned long x) { size_t i; long result = 0; for (i = 0; i < WSIZE; i++) { unsigned bit = (x >> i) & 0x1; result += bit; } return result; }</pre>	<pre>long pcount_for_goto_dw (unsigned long x) { size_t i; long result = 0; i = 0; Init if (! (i < WSIZE)) goto done; ! Test loop: { unsigned bit = (x >> i) & 0x1; result += bit; Body } i++; Update if (i < WSIZE) goto loop; Test done: return result; }</pre>
<p>■ Initial test can be optimized away</p>	

CISC 221: Computer Architecture

Switch Statement

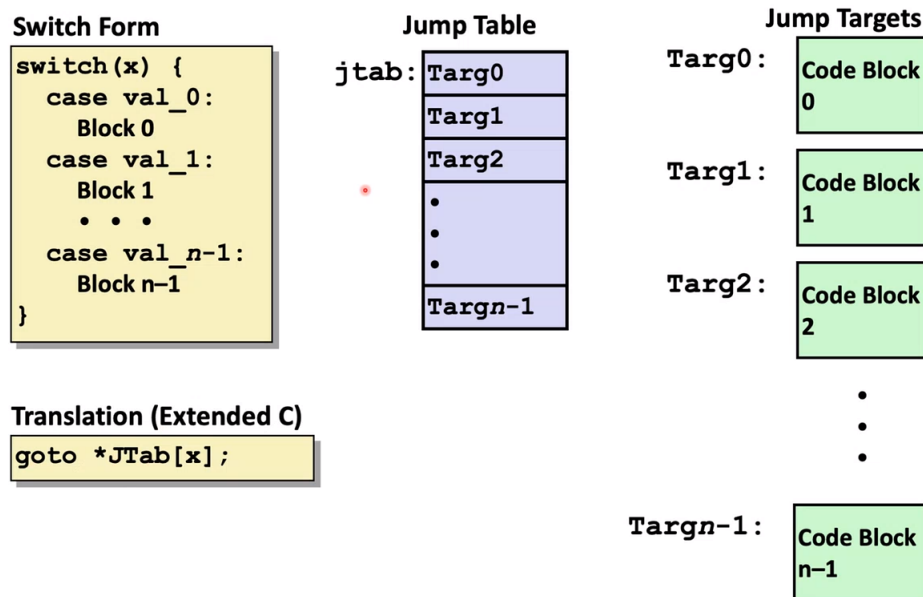
The final control structure we will cover is the switch statement.

```
long switch_eg
(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        case 1:
            w = y*z;
            break;
        case 2:
            w = y/z;
            /* Fall Through */
        case 3:
            w += z;
            break;
        case 5:
        case 6:
            w -= z;
            break;
        default:
            w = 2;
    }
    return w;
}
```

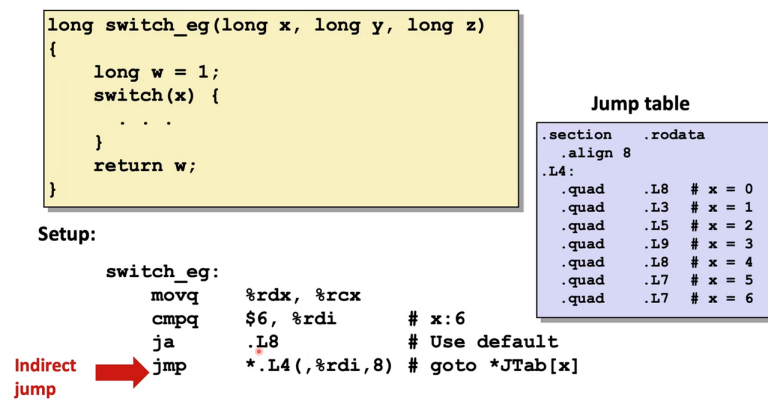
Switch Statement Example

- Multiple case labels
 - Here: 5 & 6
- Fall through cases
 - Here: 2
- Missing cases
 - Here: 4

This structure is achieved by using a Jump Table with different identifiers. This is done like so:



Showing this with the example from before:



And thus, the general setup in assembly for a switch statement is:

■ Table Structure

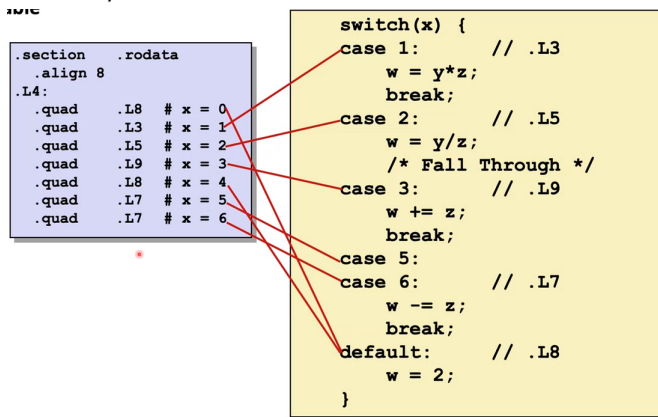
- Each target requires 8 bytes
- Base address at `.L4`

Jump table

```
.section .rodata
.align 8
.L4:
.quad .L8 # x = 0
.quad .L3 # x = 1
.quad .L5 # x = 2
.quad .L9 # x = 3
.quad .L8 # x = 4
.quad .L7 # x = 5
.quad .L7 # x = 6
```

■ Jumping

- **Direct:** `jmp .L8`
- Jump target is denoted by label `.L8`
- **Indirect:** `jmp *.L4(,%rdi,8)`
- Start of jump table: `.L4`
- Must scale by factor of 8 (addresses are 8 bytes)
- Fetch target from effective Address `.L4 + x*8`
- Only for $0 \leq x \leq 6$



Summary

In C we can implement Control by using:

- if-then-else
- do-while
- while, for
- switch

Within the Assembler there are 4 ways of achieving these:

- conditional Jump

- conditional Move
- indirect jump (via jump table)
- compiler generates code sequence to implement more complex control

Some standard techniques for operating on control structures are:

- converting loops to do-while or ‘lump-to-middle’ form
- large switch statements use jump tables
- sparse switch statements may use decision trees (if-elseif-elseif-...)