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# 1 Data structures

## 1.1 Ufds

```
class UFDS {
private:
    vector<int> size, ps;

public:
    UFDS(int N) : size(N + 1, 1), ps(N + 1) { iota(ps.begin(), ps.end(), 0); }

    int find_set(int x) { return x == ps[x] ? x : (ps[x] = find_set(ps[x])); }

    bool same_set(int x, int y) { return find_set(x) == find_set(y); }

    void union_set(int x, int y) {
        if (same_set(x, y)) return;

        int p = find_set(x);
        int q = find_set(y);

        if (size[p] < size[q]) swap(p, q);

        ps[q] = p;
        size[p] += size[q];
    }
};
```

# 2 Dynamic programming

## 2.1 Kadane

```
int kadane(const vi& xs) {
    vi s(xs.size());
    s[0] = xs[0];

    for (size_t i = 1; i < xs.size(); ++i) s[i] = max(xs[i], s[i - 1] + xs[i]);

    return *max_element(all(s));
}
```

## 2.2 Longest Increasing Subsequence (LIS)

Time:  $O(N \cdot \log N)$ .

```
int lis(vi const& a) {
    int n = a.size();
    const int INF = 1e9;
    vi d(n + 1, INF);
    d[0] = -INF;

    for (int i = 0; i < n; i++) {
        int l = upper_bound(d.begin(), d.end(), a[i]) - d.begin();
        if (d[l - 1] < a[i] && a[i] < d[l]) d[l] = a[i];
    }

    int ans = 0;
```

```
    for (int l = 0; l <= n; l++) {
        if (d[l] < INF) ans = l;
    }

    return ans;
}
```

# 3 Graphs

## 3.1 Dijkstra

```
vector<pll> adj[MAX];
class Graph {
public:
    void add(ll u, ll v, ll w) {
        adj[u].emplace_back(v, w);
        // Undirected Graph
        // adj[u].emplace_back(v, w);
    }

    vl dijkstra(ll src, ll n) {
        vl ds(n, LLONG_MAX);
        ds[src] = 0;
        set<pll> pq;
        pq.emplace(0, src);

        while (!pq.empty()) {
            ll u = pq.begin()->second;
            ll wu = pq.begin()->first;
            pq.erase(pq.begin());

            if (wu != ds[u]) continue;
            for (auto [v, w] : adj[u]) {
                if (ds[v] > ds[u] + w) {
                    ds[v] = ds[u] + w;
                    pq.emplace(ds[v], v);
                }
            }
        }

        return ds;
    }
};
```

# 4 Math

## 4.1 Convex Hull

Given a set of points find the smallest convex polygon that contains all the given points.

Time:  $O(N \cdot \log N)$

By default it removes the collinear points, set the boolean to true if you don't want that

```
struct pt {
    double x, y;
};
```

```

int orientation(pt a, pt b, pt c) {
    double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x * (a.y - b.y);
    if (v < 0) return -1; // clockwise
    if (v > 0) return +1; // counter-clockwise
    return 0;
}

bool cw(pt a, pt b, pt c, bool include_collinear) {
    int o = orientation(a, b, c);
    return o < 0 || (include_collinear && o == 0);
}

bool collinear(pt a, pt b, pt c) { return orientation(a, b, c) == 0; }

void convex_hull(vector<pt>& a, bool include_collinear = false) {
    pt p0 = *min_element(a.begin(), a.end(), [](pt a, pt b) {
        return make_pair(a.y, a.x) < make_pair(b.y, b.x);
    });
    sort(a.begin(), a.end(), [&p0](const pt& a, const pt& b) {
        int o = orientation(p0, a, b);
        if (o == 0)
            return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.y) * (p0.y - a.y) <
                (p0.x - b.x) * (p0.x - b.x) + (p0.y - b.y) * (p0.y - b.y);
        return o < 0;
    });
    if (include_collinear) {
        int i = (int)a.size() - 1;
        while (i >= 0 && collinear(p0, a[i], a.back())) i--;
        reverse(a.begin() + i + 1, a.end());
    }

    vector<pt> st;
    for (int i = 0; i < (int)a.size(); i++) {
        while (st.size() > 1 &&
            !cw(st[st.size() - 2], st.back(), a[i], include_collinear))
            st.pop_back();
        st.push_back(a[i]);
    }

    a = st;
}

```

## 4.2 Factorization With Sieve

```

map<ll, ll> factorization_with_sieve(ll n, const vl& primes) {
    map<ll, ll> fact;

    for (ll d : primes) {
        if (d * d > n) break;

        ll k = 0;
        while (n % d == 0) {
            k++;
            n /= d;
        }

        if (k) fact[d] = k;
    }
}

```

```

    if (n > 1) fact[n] = 1;
    return fact;
}

```

## 4.3 Factorization

```

map<ll, ll> factorization(ll n) {
    map<ll, ll> ans;
    for (ll i = 2; i * i <= n; i++) {
        ll count = 0;
        for (; n % i == 0; count++, n /= i)
            ;
        if (count) ans[i] = count;
    }
    if (n > 1) ans[n]++;
    return ans;
}

```

## 4.4 Euler phi $\varphi(n)$

Computes the number of positive integers less than  $n$  that are coprimes with  $n$ , in  $O(\sqrt{N})$ .

```

ll phi(ll n) {
    if (n == 1) return 1;

    auto fs = factorization(n);
    auto res = n;

    for (auto [p, k] : fs) {
        res /= p;
        res *= (p - 1);
    }

    return res;
}

```

## 4.5 Point To Segment

```

typedef pair<double, double> pdb;

#define fst first
#define snd second

double pt2segment(pdb A, pdb B, pdb E) {
    pdb AB = {B.fst - A.fst, B.snd - A.snd};
    pdb BE = {E.fst - B.fst, E.snd - B.snd};
    pdb AE = {E.fst - A.fst, E.snd - A.snd};

    double AB_BE = AB.fst * BE.fst + AB.snd * BE.snd;
    double AB_AE = AB.fst * AE.fst + AB.snd * AE.snd;

    double ans;
    if (AB_BE > 0) {
        double y = E.snd - B.snd;
        double x = E.fst - B.fst;
        ans = sqrt(x * x + y * y);
    }
}

```

```

} else if (AB_AE < 0) {
    double y = E.snd - A.snd;
    double x = E.fst - A.fst;
    ans = sqrt(x * x + y * y);
} else {
    auto [x1, y1] = AB;
    auto [x2, y2] = AE;
    double mod = sqrt(x1 * x1 + y1 * y1);
    ans = abs(x1 * y2 - y1 * x2) / mod;
}

return ans;
}

```

## 4.6 Sieve

```

v1 sieve(ll N) {
    bitset<MAX + 1> sieve;
    v1 ps{2, 3};
    sieve.set();

    for (ll i = 5, step = 2; i <= N; i += step, step = 6 - step) {
        if (sieve[i]) {
            ps.push_back(i);

            for (ll j = i * i; j <= N; j += 2 * i) sieve[j] = false;
        }
    }
    return ps;
}

```

## 5 Problems

### 5.1 Kth Digit String (CSES)

Time:  $O(\log_{10} K)$ .

Space:  $O(1)$ .

```

ll kth_digit_string(ll k) {
    if (k < 10) return k;

    ll c = 180, i = 2, u = 10, r = 0, ans = -1, m;
    for (k -= 9; k > c; i++, u *= 10) {
        k -= c;
        c /= i;
        c *= 10 * (i + 1);
    }

    if ((m = k % i))
        r++;
    else
        m = i;

    ll tmp = (k / i) + r + u - 1;
    for (m = i + 1 - m; m--; tmp /= 10) ans = tmp % 10;

    return ans;
}

```

## 6 Strings

### 6.1 Manacher

Given string  $s$  with length  $n$ . Find all the pairs  $(i, j)$  such that substring  $s[i \dots j]$  is a palindrome. String  $t$  is a palindrome when  $t = t_{rev}$  ( $t_{rev}$  is a reversed string for  $t$ ).

Time:  $O(N)$

```

vi manacher(string s) {
    string t;
    for (auto c : s) t += string("#") + c;
    t = t + '#';

    int n = t.size();
    t = "$" + t + "~";

    vi p(n + 2);
    int l = 1, r = 1;
    for (int i = 1; i <= n; i++) {
        p[i] = max(0, min(r - i, p[l + (r - i)]));
        while (t[i - p[i]] == t[i + p[i]]) p[i]++;
        if (i + p[i] > r) {
            l = i - p[i], r = i + p[i];
        }
        p[i]--;
    }

    return vi(begin(p) + 1, end(p) - 1);
}

```

## 7 Trees

### 7.1 LCA Binary Lifting (CP Algo)

The algorithm described will need  $O(N \cdot \log N)$  for preprocessing the tree, and then  $O(\log N)$  for each LCA query.

```

ll n, l;
vector<ll> adj[MAX];

ll timer;
vector<ll> tin, tout;
vector<vector<ll>> up;

void dfs(ll v, ll p) {
    tin[v] = ++timer;
    up[v][0] = p;
    for (ll i = 1; i <= l; ++i) up[v][i] = up[up[v][i - 1]][i - 1];

    for (ll u : adj[v]) {
        if (u != p) dfs(u, v);
    }

    tout[v] = ++timer;
}

```

```

bool is_ancestor(ll u, ll v) { return tin[u] <= tin[v] && tout[u] >= tout[v];
}

ll lca(ll u, ll v) {
    if (is_ancestor(u, v)) return u;
    if (is_ancestor(v, u)) return v;
    for (ll i = 1; i >= 0; --i) {
        if (!is_ancestor(up[u][i], v)) u = up[u][i];
    }
    return up[u][0];
}

void preprocess(ll root) {
    tin.resize(n);
    tout.resize(n);
    timer = 0;
    l = ceil(log2(n));
    up.assign(n, vector<ll>(l + 1));
    dfs(root, root);
}

```

## 7.2 LCA SegTree (CP Algo)

The algorithm can answer each query in  $O(\log N)$  with preprocessing in  $O(N)$  time.

```

struct LCA {
    vector<ll> height, euler, first, segtree;
    vector<bool> visited;
    ll n;

    LCA(vector<vector<ll>>& adj, ll root = 0) {
        n = adj.size();
        height.resize(n);
        first.resize(n);
        euler.reserve(n * 2);
        visited.assign(n, false);
        dfs(adj, root);
        ll m = euler.size();
        segtree.resize(m * 4);
        build(1, 0, m - 1);
    }

    void dfs(vector<vector<ll>>& adj, ll node, ll h = 0) {
        visited[node] = true;
        height[node] = h;
        first[node] = euler.size();
        euler.push_back(node);
        for (auto to : adj[node]) {
            if (!visited[to]) {
                dfs(adj, to, h + 1);
                euler.push_back(node);
            }
        }
    }

    void build(ll node, ll b, ll e) {
        if (b == e) {
            segtree[node] = euler[b];

```

```

        } else {
            ll mid = (b + e) / 2;
            build(node << 1, b, mid);
            build(node << 1 | 1, mid + 1, e);
            ll l = segtree[node << 1], r = segtree[node << 1 | 1];
            segtree[node] = (height[l] < height[r]) ? l : r;
        }
    }

    ll query(ll node, ll b, ll e, ll L, ll R) {
        if (b > R || e < L) return -1;
        if (b >= L && e <= R) return segtree[node];
        ll mid = (b + e) >> 1;

        ll left = query(node << 1, b, mid, L, R);
        ll right = query(node << 1 | 1, mid + 1, e, L, R);
        if (left == -1) return right;
        if (right == -1) return left;
        return height[left] < height[right] ? left : right;
    }

    ll lca(ll u, ll v) {
        ll left = first[u], right = first[v];
        if (left > right) swap(left, right);
        return query(1, 0, euler.size() - 1, left, right);
    }
};

```

## 7.3 LCA Sparse Table

The algorithm described will need  $O(N)$  for preprocessing, and then  $O(1)$  for each LCA query.

0 indexed !

```

#define len(__x) (int)__x.size()

using ll = long long;
using pll = pair<ll, ll>;
using vi = vector<int>;
using vi2d = vector<vi>;
#define all(a) a.begin(), a.end()
#define pb(___x) push_back(___x)
#define mp(____a, ____b) make_pair(____a, ____b)
#define eb(____x) emplace_back(____x)

template <typename T>
struct SparseTable {
    vector<T> v;
    ll n;
    static const ll b = 30;
    vi mask, t;

    ll op(ll x, ll y) { return v[x] < v[y] ? x : y; }
    ll msb(ll x) { return __builtin_clz(1) - __builtin_clz(x); }
    SparseTable() {}
    SparseTable(const vector<T>& v_) : v(v_), n(v.size()), mask(n), t(n) {
        for (ll i = 0, at = 0; i < n; mask[i++] = at |= 1) {
            at = (at << 1) & ((1 << b) - 1);
            while (at and op(i, i - msb(at & -at)) == i) at ^= at & -at;

```

```

}
for (ll i = 0; i < n / b; i++)
    t[i] = b * i + b - 1 - msb(mask[b * i + b - 1]);
for (ll j = 1; (1 << j) <= n / b; j++)
    for (ll i = 0; i + (1 << j) <= n / b; i++)
        t[n / b * j + i] =
            op(t[n / b * (j - 1) + i], t[n / b * (j - 1) + i + (1 << (j - 1))]);
}
ll small(ll r, ll sz = b) { return r - msb(mask[r] & ((1 << sz) - 1)); }
T query(ll l, ll r) {
    if (r - l + 1 <= b) return small(r, r - l + 1);
    ll ans = op(small(l + b - 1), small(r));
    ll x = l / b + 1, y = r / b - 1;
    if (x <= y) {
        ll j = msb(y - x + 1);
        ans = op(ans, op(t[n / b * j + x], t[n / b * j + y - (1 << j) + 1]));
    }
    return ans;
}
};

struct LCA {
    SparseTable<ll> st;
    ll n;
    vi v, pos, dep;

    LCA(const vi2d& g, ll root) : n(len(g)), pos(n) {
        dfs(root, 0, -1, g);
        st = SparseTable<ll>(vector<ll>(all(dep)));
    }

    void dfs(ll i, ll d, ll p, const vi2d& g) {
        v.pb(len(dep)) = i, pos[i] = len(dep), dep.pb(d);
        for (auto j : g[i])
            if (j != p) {
                dfs(j, d + 1, i, g);
                v.pb(len(dep)) = i, dep.pb(d);
            }
    }

    ll lca(ll a, ll b) {
        ll l = min(pos[a], pos[b]);
        ll r = max(pos[a], pos[b]);
        return v[st.query(l, r)];
    }

    ll dist(ll a, ll b) {
        return dep[pos[a]] + dep[pos[b]] - 2 * dep[pos[lca(a, b)]];
    }
};

```

## 8 Settings and macros

### 8.1 macro.cpp

```

#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>

```

```

#include <ext/pb_ds/tree_policy.hpp>

using namespace __gnu_pbds;
#define ordered_set tree<int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update>

using namespace std;

typedef long long ll;
typedef pair<int, int> pii;
typedef pair<ll, ll> pll;
typedef vector<int> vi;
typedef vector<ll> vl;
typedef vector<pii> vii;

#define all(x) x.begin(), x.end()
#define vin(vt) for (auto &e : vt) cin >> e
#define LSONe(S) ((S) & -(S))
#define MSONe(S) (1ull << (63 - __builtin_clzll(S)))

const vii dir4{ {1,0},{-1,0},{0,1},{0,-1} };

auto solve() { }

int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);

    ll t = 1;
    //cin >> t;

    while (t--) solve();

    return 0;
}

```

### 8.2 short-macro.cpp

```

#include <bits/stdc++.h>

using namespace std;

typedef long long ll;
typedef pair<int, int> ii;

#define all(x) x.begin(), x.end()
#define vin(vt) for (auto &e : vt) cin >> e

auto solve() { }

int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);

    ll t = 1;
    //cin >> t;

    while (t--) solve();
}

```

```
return 0;
```

```
    }  
}
```