

Notebook - Competitive Programming

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1 Data structures

1.1 Matrix

```
template <typename T>
struct Matrix {
    vector<vector<T>> d;

    Matrix() : Matrix(0) {}
    Matrix(int n) : Matrix(n, n) {}
    Matrix(int n, int m) : Matrix(vector<vector<T>>(n, vector<T>(m))) {}
    Matrix(const vector<vector<T>> &v) : d(v) {}

    constexpr int n() const { return (int)d.size(); }
    constexpr int m() const { return n() ? (int)d[0].size() : 0; }

    void rotate() { *this = rotated(); }

    Matrix<T> rotated() const {
        Matrix<T> res(m(), n());
        for (int i = 0; i < m(); i++) {
            for (int j = 0; j < n(); j++) {
                res[i][j] = d[n() - j - 1][i];
            }
        }
        return res;
    }

    Matrix<T> pow(int power) const {
        assert(n() == m());

        auto res = Matrix<T>::identity(n());
        auto b = *this;
        while (power) {
            if (power & 1) res *= b;
            b *= b;
            power >>= 1;
        }
        return res;
    }

    Matrix<T> submatrix(int start_i, int start_j, int rows = INT_MAX,
                       int cols = INT_MAX) const {
        rows = min(rows, n() - start_i);
        cols = min(cols, m() - start_j);
        if (rows <= 0 or cols <= 0) return {};

        Matrix<T> res(rows, cols);
        for (int i = 0; i < rows; i++)
            for (int j = 0; j < cols; j++) res[i][j] = d[i + start_i][j + start_j];
        return res;
    }

    Matrix<T> translated(int x, int y) const {
        Matrix<T> res(n(), m());
        for (int i = 0; i < n(); i++) {
            for (int j = 0; j < m(); j++) {
                if (i + x < 0 or i + x >= n() or j + y < 0 or j + y >= m()) continue;

```

```
                res[i + x][j + y] = d[i][j];
            }
        }
        return res;
    }

    static Matrix<T> identity(int n) {
        Matrix<T> res(n);
        for (int i = 0; i < n; i++) res[i][i] = 1;
        return res;
    }

    vector<T> &operator[](int i) { return d[i]; }
    const vector<T> &operator[](int i) const { return d[i]; }
    Matrix<T> &operator+=(T value) {
        for (auto &row : d) {
            for (auto &x : row) x += value;
        }
        return *this;
    }
    Matrix<T> operator+(T value) const {
        auto res = *this;
        for (auto &row : res) {
            for (auto &x : row) x = x + value;
        }
        return res;
    }
    Matrix<T> &operator-=(T value) {
        for (auto &row : d) {
            for (auto &x : row) x -= value;
        }
        return *this;
    }
    Matrix<T> operator-(T value) const {
        auto res = *this;
        for (auto &row : res) {
            for (auto &x : row) x = x - value;
        }
        return res;
    }
    Matrix<T> &operator*=(T value) {
        for (auto &row : d) {
            for (auto &x : row) x *= value;
        }
        return *this;
    }
    Matrix<T> operator*(T value) const {
        auto res = *this;
        for (auto &row : res) {
            for (auto &x : row) x = x * value;
        }
        return res;
    }
    Matrix<T> &operator/=(T value) {
        for (auto &row : d) {
            for (auto &x : row) x /= value;
        }
        return *this;
    }
}
```

```

}
Matrix<T> operator/(T value) const {
    auto res = *this;
    for (auto &row : res) {
        for (auto &x : row) x = x / value;
    }
    return res;
}
Matrix<T> &operator+=(const Matrix<T> &o) {
    assert(n() == o.n() and m() == o.m());
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            d[i][j] += o[i][j];
        }
    }
    return *this;
}
Matrix<T> operator+(const Matrix<T> &o) const {
    assert(n() == o.n() and m() == o.m());
    auto res = *this;
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            res[i][j] = res[i][j] + o[i][j];
        }
    }
    return res;
}
Matrix<T> &operator-=(const Matrix<T> &o) {
    assert(n() == o.n() and m() == o.m());
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            d[i][j] -= o[i][j];
        }
    }
    return *this;
}
Matrix<T> operator-(const Matrix<T> &o) const {
    assert(n() == o.n() and m() == o.m());
    auto res = *this;
    for (int i = 0; i < n(); i++) {
        for (int j = 0; j < m(); j++) {
            res[i][j] = res[i][j] - o[i][j];
        }
    }
    return res;
}
Matrix<T> &operator*=(const Matrix<T> &o) {
    *this = *this * o;
    return *this;
}
Matrix<T> operator*(const Matrix<T> &o) const {
    assert(m() == o.n());
    Matrix<T> res(n(), o.m());
    for (int i = 0; i < res.n(); i++) {
        for (int j = 0; j < res.m(); j++) {
            auto &x = res[i][j];
            for (int k = 0; k < m(); k++) {
                x += (d[i][k] * o[k][j]);
            }
        }
    }
}

```

```

    }
}
return res;
}

friend istream &operator>>(istream &is, Matrix<T> &mat) {
    for (auto &row : mat)
        for (auto &x : row) is >> x;
    return is;
}

friend ostream &operator<<(ostream &os, const Matrix<T> &mat) {
    bool frow = 1;
    for (auto &row : mat) {
        if (not frow) os << '\n';
        bool first = 1;
        for (auto &x : row) {
            if (not first) os << ' ';
            os << x;
            first = 0;
        }

        frow = 0;
    }
    return os;
}

auto begin() { return d.begin(); }
auto end() { return d.end(); }
auto rbegin() { return d.rbegin(); }
auto rend() { return d.rend(); }

auto begin() const { return d.begin(); }
auto end() const { return d.end(); }
auto rbegin() const { return d.rbegin(); }
auto rend() const { return d.rend(); }
};

```

1.2 Merge Sort Tree

Like a segment tree but each node st_i stores a sorted subarray

- $inrange(l, r, a, b)$: counts the number of elements $x \in [l, r]$ such that $a \leq x \leq b$.

Memory: $O(N \log N)$

Build: $O(N \log N)$

$inrange$: $O(\log^2 N)$

```

template <class T>
struct MergeSortTree {
    int n;
    vector<vector<T>> st;
    MergeSortTree(vector<T>& xs) : n(len(xs)), st(n << 1) {
        for (int i = 0; i < n; i++) st[i + n] = vector<T>({xs[i]});

        for (int i = n - 1; i > 0; i--) {
            st[i].resize(len(st[i << 1]) + len(st[i << 1 | 1]));
            merge(all(st[i << 1]), all(st[i << 1 | 1]), st[i].begin());
        }
    }
}

```

```

int count(int i, T a, T b) {
    return upper_bound(all(st[i]), b) - lower_bound(all(st[i]), a);
}

int inrange(int l, int r, T a, T b) {
    int ans = 0;

    for (l += n, r += n + 1; l < r; l >>= 1, r >>= 1) {
        if (l & 1) ans += count(l++, a, b);
        if (r & 1) ans += count(--r, a, b);
    }

    return ans;
}
};

```

1.3 Minimal Excluded With Updates (MEX-U)

In the problem you need to change individual numbers in the array, and compute the new MEX of the array after each such update.

Pre-compute: $O(N \log N)$

Update: $O(\log N)$

Query: $O(1)$

```

class Mex {
private:
    map<ll, ll> frequency;
    set<ll> missing_numbers;
    vl A;

public:
    Mex(vl const& A) : A(A) {
        for (ll i = 0; i <= A.size(); i++) missing_numbers.insert(i);

        for (ll x : A) {
            ++frequency[x];
            missing_numbers.erase(x);
        }

        ll mex() { return *missing_numbers.begin(); }

        void update(ll idx, ll new_value) {
            if (--frequency[A[idx]] == 0) missing_numbers.insert(A[idx]);
            A[idx] = new_value;
            ++frequency[new_value];
            missing_numbers.erase(new_value);
        }
    };
};

```

1.4 Minimal Excluded (MEX)

Given an array A of size N . You have to find the minimal non-negative element that is not present in the array. That number is commonly called the MEX (minimal excluded).

Time: $O(N)$

```

ll mex(vl const& A) {
    static bool used[MAX + 111] = {0};

    for (ll x : A) {
        if (x <= MAX) used[x] = true;
    }

    ll result = 0;
    while (used[result]) ++result;

    for (ll x : A) {
        if (x <= MAX) used[x] = false;
    }

    return result;
}

```

1.5 Range Min Query (RMQ)

Build: $O(N)$

Query: $O(1)$

```

// @brunomaletta
template <typename T>
struct rmq {
    vector<T> v;
    int n;
    static const int b = 30;
    vector<int> mask, t;

    int op(int x, int y) { return v[x] <= v[y] ? x : y; }
    int msb(int x) { return __builtin_clz(1) - __builtin_clz(x); }
    int small(int r, int sz = b) { return r - msb(mask[r] & ((1 << sz) - 1)); }
    rmq() {}
    rmq(const vector<T>& v_) : v(v_), n(v.size()), mask(n), t(n) {
        for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
            at = (at << 1) & ((1 << b) - 1);
            while (at and op(i - msb(at & -at), i) == i) at ^= at & -at;
        }
        for (int i = 0; i < n / b; i++) t[i] = small(b * i + b - 1);
        for (int j = 1; (1 << j) <= n / b; j++)
            for (int i = 0; i + (1 << j) <= n / b; i++)
                t[n / b * j + i] =
                    op(t[n / b * (j - 1) + i], t[n / b * (j - 1) + i + (1 << (j - 1))]);
    }

    int index_query(int l, int r) {
        if (r - l + 1 <= b) return small(r, r - l + 1);
        int x = l / b + 1, y = r / b - 1;
        if (x > y) return op(small(l + b - 1), small(r));
        int j = msb(y - x + 1);
        int ans = op(small(l + b - 1),
                    op(t[n / b * j + x], t[n / b * j + y - (1 << j) + 1]));
        return op(ans, small(r));
    }

    T query(int l, int r) { return v[index_query(l, r)]; }
};

```

1.6 Segment Tree (Parameterized OP)

Query: $O(\log N)$
Update: $O(\log N)$

```
template <typename T, auto op>
class SegTree {
private:
    T e;
    ll N;
    vector<T> seg;

public:
    SegTree(ll N, T e) : e(e), N(N), seg(N + N, e) {}

    void assign(ll i, T v) {
        i += N;
        seg[i] = v;
        for (i >= 1; i; i >= 1) seg[i] = op(seg[2 * i], seg[2 * i + 1]);
    }

    T query(ll l, ll r) {
        T la = e, ra = e;
        l += N;
        r += N;

        while (l <= r) {
            if (l & 1) la = op(la, seg[l++]);
            if (~r & 1) ra = op(seg[r--], ra);
            l >>= 1;
            r >>= 1;
        }

        return op(la, ra);
    }
};
```

1.7 Segment Tree 2D

Query: $O(\log N \cdot \log M)$
Update: $O(\log N \cdot \log M)$

```
template <typename T, auto op>
class SegTree {
private:
    T e;
    ll n, m;
    vector<vector<T>> seg;

public:
    SegTree(ll n, ll m, T e)
        : e(e), n(n), m(m), seg(2 * n, vector<T>(2 * m, e)) {}

    void assign(ll x, ll y, T v) {
        ll ny = y += m;
        for (x += n; x; x >>= 1, y = ny) {
            if (x >= n)
                seg[x][y] = v;
            else
```

```
                seg[x][y] = op(seg[2 * x][y], seg[2 * x + 1][y]);

            while (y >= 1) seg[x][y] = op(seg[x][2 * y], seg[x][2 * y + 1]);
        }

    T query(ll lx, ll rx, ll ly, ll ry) {
        ll ans = e, nx = rx + n, my = ry + m;

        for (lx += n, ly += m; lx <= ly; ++lx >= 1, --ly >= 1)
            for (rx = nx, ry = my; rx <= ry; ++rx >= 1, --ry >= 1) {
                if (lx & 1 and rx & 1) ans = op(ans, seg[lx][rx]);
                if (lx & 1 and !(ry & 1)) ans = op(ans, seg[lx][ry]);
                if (!(ly & 1) and rx & 1) ans = op(ans, seg[ly][rx]);
                if (!(ly & 1) and !(ry & 1)) ans = op(ans, seg[ly][ry]);
            }

        return ans;
    }
};
```

1.8 Segment Tree Lazy

Query (Range Sum): $O(\log N)$
Update (Sum Value): $O(\log N)$

```
template <typename T>
class SegTreeLazy {
private:
    int N;
    vector<T> seg, lzy;

    void down(int k, int l, int r) {
        seg[k] += (r - l + 1) * lzy[k];
        if (l < r) {
            lzy[k << 1] += lzy[k];
            lzy[k << 1 | 1] += lzy[k];
        }
        lzy[k] = 0;
    }

    void update(int i, int j, int k, int l, int r, T v) {
        if (lzy[k]) down(k, l, r);
        if (i > r or j < l) return;
        if (i <= l and j >= r) {
            seg[k] += (r - l + 1) * v;
            if (l < r) {
                lzy[k << 1] += v;
                lzy[k << 1 | 1] += v;
            }
            return;
        }

        update(i, j, k << 1, l, (l + r) / 2, v);
        update(i, j, k << 1 | 1, (l + r) / 2 + 1, r, v);
        seg[k] = seg[k << 1] + seg[k << 1 | 1];
    }
};
```

```

T query(int i, int j, int k, int l, int r) {
    if (lzy[k]) down(k, l, r);
    if (i > r or j < l) return 0;
    if (i <= l and j >= r) return seg[k];

    T lft = query(i, j, k << 1, l, (l + r) / 2);
    T rgt = query(i, j, k << 1 | 1, (l + r) / 2 + 1, r);
    return lft + rgt;
}

public:
    SegTreeLazy(int N) : N(N), seg(N << 2, 0), lzy(N << 2, 0) {}

    void update(int i, int j, T v) { update(i, j, 1, 0, N - 1, v); }

    T query(int i, int j) { return query(i, j, 1, 0, N - 1); }
};

```

1.9 Union Find Disjoint Set (UFDS)

Uncomment the lines to recover which element belong to each set.

Time: $\approx O(1)$ for everything.

```

class UFDS {
public:
    vi ps, size;
    // vector<unordered_set<int>> sts;

    UFDS(int N) : size(N + 1, 1), ps(N + 1), sts(N) {
        iota(ps.begin(), ps.end(), 0);
        // for (int i = 0; i < N; i++) sts[i].insert(i);
    }

    int find_set(int x) { return x == ps[x] ? x : (ps[x] = find_set(ps[x])); }

    bool same_set(int x, int y) { return find_set(x) == find_set(y); }

    void union_set(int x, int y) {
        if (same_set(x, y)) return;

        int px = find_set(x);
        int py = find_set(y);

        if (size[px] < size[py]) swap(px, py);

        ps[py] = px;
        size[px] += size[py];
        // sts[px].merge(sts[py]);
    }
};

```

1.10 Wavelet Tree

Build: $O(N \cdot \log \sigma)$.

Queries: $O(\log \sigma)$.

σ = alphabet length

```

typedef vector<int>::iterator iter;

class WaveletTree {
public:
    int L, H;
    WaveletTree *l, *r;
    vector<int> frq;

    WaveletTree(iter fr, iter to, int x, int y) {
        L = x, H = y;
        if (fr >= to) return;

        int M = L + ((H - L) >> 1);
        auto F = [M](int x) { return x <= M; };

        frq.reserve(to - fr + 1);
        frq.push_back(0);
        for (auto it = fr; it != to; it++) frq.push_back(frq.back() + F(*it));

        if (H == L) return;
        auto pv = stable_partition(fr, to, F);
        l = new WaveletTree(fr, pv, L, M);
        r = new WaveletTree(pv, to, M + 1, H);
    }

    // Find the k-th smallest element in positions [i,j]
    int quantile(int l, int r, int k) {
        if (l > r) return 0;
        if (L == H) return L;
        int inLeft = frq[r] - frq[l - 1];
        int lb = frq[l - 1], rb = frq[r];
        if (k <= inLeft) return this->l->quantile(lb + 1, rb, k);
        return this->r->quantile(l - lb, r - rb, k - inLeft);
    }

    // Count occurrences of number c until position i -> [0, i].
    int rank(int c, int i) { return until(c, min(i + 1, (int)frq.size() - 1)); }

    int until(int c, int i) {
        if (c > H or c < L) return 0;
        if (L == H) return i;

        int M = L + ((H - L) >> 1);
        int r = frq[i];
        if (c <= M)
            return this->l->until(c, r);
        else
            return this->r->until(c, i - r);
    }

    // Count number of occurrences of numbers in the range [a, b]
    int range(int i, int j, int a, int b) const {
        if (b < a or j < i) return 0;
        return range(i, j + 1, L, H, a, b);
    }

    int range(int i, int j, int a, int b, int L, int U) const {
        if (b < L or U < a) return 0;

```

```

    if (L <= a and b <= U) return j - i;
    int M = a + ((b - a) >> 1);
    int ri = frq[i], rj = frq[j];
    return this->l->range(ri, rj, a, M, L, U) +
           this->r->range(i - ri, j - rj, M + 1, b, L, U);
}

// Number of elements greater than or equal to k in [l, r];
// Can count distinct in a range with aux vector of next pos
int greater(int l, int r, int k) { return _greater(l + 1, r + 1, k); }

int _greater(int l, int r, int k) {
    if (l > r or k > H) return 0;
    if (L >= k) return r - l + 1;

    int ri = frq[l - 1], rj = frq[r];
    return this->l->_greater(ri + 1, rj, k) +
           this->r->_greater(l - ri, r - rj, k);
}
};

```

2 Dynamic programming

2.1 Kadane

```

int kadane(const vi& xs) {
    vi s(xs.size());
    s[0] = xs[0];

    for (size_t i = 1; i < xs.size(); ++i) s[i] = max(xs[i], s[i - 1] + xs[i]);

    return *max_element(all(s));
}

```

2.2 Longest Increasing Subsequence (LIS)

Time: $O(N \cdot \log N)$.

```

int lis(vi const& a) {
    int n = a.size();
    const int INF = 1e9;
    vi d(n + 1, INF);
    d[0] = -INF;

    for (int i = 0; i < n; i++) {
        int l = upper_bound(d.begin(), d.end(), a[i]) - d.begin();
        if (d[l - 1] < a[i] && a[i] < d[l]) d[l] = a[i];
    }

    int ans = 0;
    for (int l = 0; l <= n; l++) {
        if (d[l] < INF) ans = l;
    }

    return ans;
}

```

3 Extras

3.1 cin/cout __int128_t

Allows standard reading and writing with cin/cout for 128-bit integers using `__int128_t` type.

```

ostream& operator<<(ostream& dest, __int128_t value) {
    ostream::sentry s(dest);
    if (s) {
        __uint128_t tmp = value < 0 ? -value : value;
        char buffer[128];
        char* d = end(buffer);
        do {
            --d;
            *d = "0123456789"[tmp % 10];
            tmp /= 10;
        } while (tmp != 0);
        if (value < 0) {
            --d;
            *d = '-';
        }
        int len = end(buffer) - d;
        if (dest.rdbuf()->sputn(d, len) != len) dest.setstate(ios_base::badbit);
    }
    return dest;
}

istream& operator>>(istream& is, __int128_t& value) {
    string s;
    is >> s;

    __int128_t res = 0;
    size_t i = 0;

    bool neg = false;
    if (s[i] == '-') neg = 1, i++;
    for (; i < s.size(); ++i) (res *= 10) += (s[i] - '0');

    value = neg ? -res : res;
    return is;
}

```

4 Geometry

4.1 Convex Hull

Given a set of points find the smallest convex polygon that contains all the given points.

Time: $O(N \cdot \log N)$

By default it removes the collinear points, set the boolean to true if you don't want that

```

struct pt {
    double x, y;
};

int orientation(pt a, pt b, pt c) {
    double v = a.x * (b.y - c.y) + b.x * (c.y - a.y) + c.x * (a.y - b.y);
    if (v < 0) return -1; // clockwise
}

```

```

    if (v > 0) return +1; // counter-clockwise
    return 0;
}

bool cw(pt a, pt b, pt c, bool include_collinear) {
    int o = orientation(a, b, c);
    return o < 0 || (include_collinear && o == 0);
}

bool collinear(pt a, pt b, pt c) { return orientation(a, b, c) == 0; }

void convex_hull(vector<pt>& a, bool include_collinear = false) {
    pt p0 = *min_element(a.begin(), a.end(), [](pt a, pt b) {
        return make_pair(a.y, a.x) < make_pair(b.y, b.x);
    });
    sort(a.begin(), a.end(), [&p0](const pt& a, const pt& b) {
        int o = orientation(p0, a, b);
        if (o == 0)
            return (p0.x - a.x) * (p0.x - a.x) + (p0.y - a.y) * (p0.y - a.y) <
                (p0.x - b.x) * (p0.x - b.x) + (p0.y - b.y) * (p0.y - b.y);
        return o < 0;
    });
    if (include_collinear) {
        int i = (int)a.size() - 1;
        while (i >= 0 && collinear(p0, a[i], a.back())) i--;
        reverse(a.begin() + i + 1, a.end());
    }

    vector<pt> st;
    for (int i = 0; i < (int)a.size(); i++) {
        while (st.size() > 1 &&
            !cw(st[st.size() - 2], st.back(), a[i], include_collinear))
            st.pop_back();
        st.push_back(a[i]);
    }

    a = st;
}

```

4.2 Point To Segment

```

typedef pair<double, double> pdb;

double pt2segment(pdb A, pdb B, pdb E) {
    pdb AB = {B.fst - A.fst, B.snd - A.snd};
    pdb BE = {E.fst - B.fst, E.snd - B.snd};
    pdb AE = {E.fst - A.fst, E.snd - A.snd};

    double AB_BE = AB.fst * BE.fst + AB.snd * BE.snd;
    double AB_AE = AB.fst * AE.fst + AB.snd * AE.snd;

    double ans;
    if (AB_BE > 0) {
        double y = E.snd - B.snd;
        double x = E.fst - B.fst;
        ans = hypot(x, y);
    } else if (AB_AE < 0) {
        double y = E.snd - A.snd;

```

```

        double x = E.fst - A.fst;
        ans = hypot(x, y);
    } else {
        auto [x1, y1] = AB;
        auto [x2, y2] = AE;
        double mod = hypot(x1, y1);
        ans = abs(x1 * y2 - y1 * x2) / mod;
    }

    return ans;
}

```

4.3 Polynominoes

Geometric figure made by equal squares, connected between themselves in a way that at least one side of each square coincide with a side of another square.

Watch out: the number of polynominoes increases fastly (size 12 has 63.600 figures)

```

// We consider the rotations
// as distinct (0, 10, 10+9, 10+9+8...)
vi pos = {0, 10, 19, 27, 34, 40, 45, 49, 52, 54, 55};

```

```

const int MAXP = 10;

struct Poly {
    ii v[MAXP];
    int64_t id;
    int n;

    Poly() {
        n = 1;
        v[0] = {0, 0};
        normalize();
    }

    Poly(vii &vp) {
        n = vp.size();
        for (int i = 0; i < n; i++) v[i] = vp[i];
        normalize();
    }

    ii &operator[](int i) { return v[i]; }

    bool add(int a, int b) {
        for (int i = 0; i < n; i++) {
            auto [f, s] = v[i];
            if (f == a and s == b) return false;
        }

        v[n++] = ii{a, b};
        normalize();
        return true;
    }

    void normalize() {
        int mx = 100, my = 100;
        for (int i = 0; i < n; i++) {
            auto [f, s] = v[i];
            mx = min(mx, f), my = min(my, s);

```



```

}

id = 0;
for (int i = 0; i < n; i++) {
    auto &[f, s] = v[i];
    f -= mx, s -= my;
    id |= (1LL << (pos[f] + s));
}
}

bool operator<(Poly oth) { return id < oth.id; }
};

vector<Poly> poly[MAXP + 1];

void buildPoly(int mxN) {
    for (int i = 0; i <= mxN; i++) poly[i].clear();

    Poly init;
    queue<Poly> q;
    unordered_set<int64_t> used;
    q.push(init);
    used.insert(init.id);
    while (not q.empty()) {
        Poly u = q.front();
        q.pop();
        poly[u.n].emplace_back(u);

        if (u.n == mxN) continue;

        for (int i = 0; i < u.n; i++) {
            for (auto [dx, dy] : dir4) {
                Poly to = u;
                auto [f, s] = to[i];
                bool ok = to.add(f + dx, s + dy);

                if (ok and not used.count(to.id)) {
                    q.push(to);
                    used.insert(to.id);
                }
            }
        }
    }
}
}

```

5 Graphs

5.1 Articulation Points

```

int dfs_num[MAX], dfs_low[MAX];
vi adj[MAX];

int dfs_articulation_points(int u, int p, int& next, set<int>& points) {
    int children = 0;
    dfs_low[u] = dfs_num[u] = next++;

```

```

    for (auto v : adj[u])
        if (not dfs_num[v]) {
            ++children;

            dfs_articulation_points(v, u, next, points);

            if (dfs_low[v] >= dfs_num[u]) points.insert(u);

            dfs_low[u] = min(dfs_low[u], dfs_low[v]);
        } else if (v != p)
            dfs_low[u] = min(dfs_low[u], dfs_num[v]);

    return children;
}

set<int> articulation_points(int N) {
    memset(dfs_num, 0, (N + 1) * sizeof(int));
    memset(dfs_low, 0, (N + 1) * sizeof(int));

    set<int> points;

    for (int u = 1, next = 1; u <= N; ++u)
        if (not dfs_num[u]) {
            auto children = dfs_articulation_points(u, u, next, points);

            if (children == 1) points.erase(u);
        }

    return points;
}

```

5.2 Bellman Ford

Time: $O(V \cdot E)$. Returns the shortest path from s to all other nodes.

```

using edge = tuple<int, int, int>;

pair<vi, vi> bellman_ford(int s, int N, const vector<edge>& edges) {
    vi dist(N + 1, oo), pred(N + 1, oo);

    dist[s] = 0;
    pred[s] = s;

    for (int i = 1; i <= N - 1; i++)
        for (auto [u, v, w] : edges)
            if (dist[u] < oo and dist[v] > dist[u] + w) {
                dist[v] = dist[u] + w;
                pred[v] = u;
            }

    return {dist, pred};
}

```

5.3 BFS 0/1

Time: $O(V + E)$.

```

vii adj[MAX];

```

```

vi bfs_01(int s, int N) {
    vi dist(N + 1, oo);
    dist[s] = 0;

    deque<int> q;
    q.emplace_back(s);

    while (not q.empty()) {
        auto u = q.front();
        q.pop_front();

        for (auto [v, w] : adj[u])
            if (dist[v] > dist[u] + w) {
                dist[v] = dist[u] + w;
                w == 0 ? q.emplace_front(v) : q.emplace_back(v);
            }
    }

    return dist;
}

```

5.4 Binary Lifting

Time: $O(N \cdot \log_2 K)$

```

const int MAXN = 2e5, MAXLOG2 = 60;
int bl[MAXN][MAXLOG2 + 1];

int jump(int u, ll k) {
    for (int i = 0; i <= MAXLOG2; i++)
        if (k & (1LL << i)) u = bl[u][i];

    return u;
}

void build(int N) {
    for (int i = 1; i <= MAXLOG2; i++)
        for (int j = 0; j < N; j++) bl[j][i] = bl[bl[j][i - 1]][i - 1];
}

```

5.5 Bridges

```

int dfs_num[MAX], dfs_low[MAX];
vi adj[MAX];

void dfs_bridge(int u, int p, int& next, vii& bridges) {
    dfs_low[u] = dfs_num[u] = next++;

    for (auto v : adj[u])
        if (not dfs_num[v]) {
            dfs_bridge(v, u, next, bridges);

            if (dfs_low[v] > dfs_num[u]) bridges.emplace_back(u, v);

            dfs_low[u] = min(dfs_low[u], dfs_low[v]);
        }
}

```

```

    } else if (v != p)
        dfs_low[u] = min(dfs_low[u], dfs_num[v]);
}

vii bridges(int N) {
    memset(dfs_num, 0, (N + 1) * sizeof(int));
    memset(dfs_low, 0, (N + 1) * sizeof(int));

    vii bridges;

    for (int u = 1, next = 1; u <= N; ++u)
        if (not dfs_num[u]) dfs_bridge(u, u, next, bridges);

    return bridges;
}

```

5.6 Negative Cycle Bellman Ford

Time: $O(V \cdot E)$. Detects whether there is a negative cycle in the graph using Bellman Ford.

```

using edge = tuple<int, int, int>;

bool has_negative_cycle(int s, int N, const vector<edge>& edges) {
    vi dist(N + 1, oo);
    dist[s] = 0;

    for (int i = 1; i <= N - 1; i++)
        for (auto [u, v, w] : edges)
            if (dist[u] < oo and dist[v] > dist[u] + w) dist[v] = dist[u] + w;

    for (auto [u, v, w] : edges)
        if (dist[u] < oo and dist[v] > dist[u] + w) return true;

    return false;
}

```

5.7 Negative Cycle Floyd Warshall

Time: $O(n^3)$. Detects whether there is a negative cycle in the graph using Floyd Warshall.

```

int dist[MAX][MAX];
vii adj[MAX];

bool has_negative_cycle(int N) {
    for (int u = 1; u <= N; ++u)
        for (int v = 1; v <= N; ++v) dist[u][v] = u == v ? 0 : oo;

    for (int u = 1; u <= N; ++u)
        for (auto [v, w] : adj[u]) dist[u][v] = w;

    for (int k = 1; k <= N; ++k)
        for (int u = 1; u <= N; ++u)
            for (int v = 1; v <= N; ++v)
                if (dist[u][k] < oo and dist[k][v] < oo)
                    dist[u][v] = min(dist[u][v], dist[u][k] + dist[k][v]);

    for (int i = 1; i <= N; ++i)
        if (dist[i][i] < 0) return true;
}

```

```
    return false;
}
```

5.8 Dijkstra

```
pair<vl, vl> Graph::dijkstra(ll src) {
    vl pd(this->N, LLONG_MAX), ds(this->N, LLONG_MAX);
    pd[src] = src;
    ds[src] = 0;

    set<pll> st;
    st.emplace(0, src);

    while (!st.empty()) {
        ll u = st.begin()->snd;
        ll wu = st.begin()->fst;
        st.erase(st.begin());

        if (wu != ds[u]) continue;
        for (auto& [v, w] : adj[u]) {
            if (ds[v] > ds[u] + w) {
                ds[v] = ds[u] + w;
                pd[v] = u;
                st.emplace(ds[v], v);
            }
        }
    }

    return {ds, pd};
}
```

5.9 Floyd Warshall

```
vii adj[MAX];

pair<vector<vi>, vector<vi>> floyd_warshall(int N) {
    vector<vi> dist(N + 1, vi(N + 1, oo));
    vector<vi> pred(N + 1, vi(N + 1, oo));

    for (int u = 1; u <= N; ++u) {
        dist[u][u] = 0;
        pred[u][u] = u;
    }

    for (int u = 1; u <= N; ++u)
        for (auto [v, w] : adj[u]) {
            dist[u][v] = w;
            pred[u][v] = u;
        }

    for (int k = 1; k <= N; ++k) {
        for (int u = 1; u <= N; ++u) {
            for (int v = 1; v <= N; ++v) {
                if (dist[u][k] < oo and dist[k][v] < oo and
                    dist[u][v] > dist[u][k] + dist[k][v]) {
```

```
                    dist[u][v] = dist[u][k] + dist[k][v];
                    pred[u][v] = pred[k][v];
                }
            }
        }

    return {dist, pred};
}
```

5.10 Graph

```
class Graph {
private:
    ll N;
    bool undirected;
    vector<vll> adj;

public:
    Graph(ll N, bool is_undirected = true) {
        this->N = N;
        adj.resize(N);
        undirected = is_undirected;
    }

    void add(ll u, ll v, ll w) {
        adj[u].emplace_back(v, w);
        if (undirected) adj[v].emplace_back(u, w);
    }
};
```

5.11 TopSort - Kahn

Works only on Directed Acyclic Graphs (DAGs). For each edge (u,v), u comes before v in the ordering. If the task A is a prerequisite for task B, then A comes before B in the ordering. Time: $O(E \cdot \log(v))$

```
unordered_set<int> in[MAX], out[MAX];

vi topological_sort(int N) {
    vi o;
    queue<int> q;

    for (int u = 1; u <= N; ++u)
        if (in[u].empty()) q.push(u);

    while (not q.empty()) {
        auto u = q.front();
        q.pop();

        o.emplace_back(u);

        for (auto v : out[u]) {
            in[v].erase(u);

            if (in[v].empty()) q.push(v);
        }
    }

    return (int)o.size() == N ? o : vi{};
}
```

```
}
```

5.12 Kosaraju

Time: $O(V + E)$. Returns a vector of vectors indicating the directed strongly connected nodes.

```
vi adj[MAX], rev[MAX];
bitset<MAX> visited;

void dfs(int u, vi& order) {
    if (visited[u]) return;

    visited[u] = true;

    for (auto v : adj[u]) dfs(v, order);

    order.emplace_back(u);
}

vi dfs_order(int N) {
    vi order;

    for (int u = 1; u <= N; ++u) dfs(u, order);

    return order;
}

void dfs_cc(int u, vi& cc) {
    if (visited[u]) return;

    visited[u] = true;
    cc.emplace_back(u);

    for (auto v : rev[u]) dfs_cc(v, cc);
}

vector<vi> kosaraju(int N) {
    auto order = dfs_order(N);
    reverse(order.begin(), order.end());

    for (int u = 1; u <= N; ++u)
        for (auto v : adj[u]) rev[v].emplace_back(u);

    vector<vi> cs;
    visited.reset();

    for (auto u : order) {
        if (visited[u]) continue;

        cs.emplace_back(vi());
        dfs_cc(u, cs.back());
    }

    return cs;
}
```

5.13 Kruskal

Time: $O(e \cdot \log(v))$

```
using edge = tuple<int, int, int>;

int kruskal(int N, vector<edge>& es) {
    sort(es.begin(), es.end());

    int cost = 0;
    UnionFind udfs(N);

    for (auto [w, u, v] : es) {
        if (not udfs.same_set(u, v)) {
            cost += w;
            udfs.union_set(u, v);
        }
    }

    return cost;
}
```

5.14 Minimax

A MST minimizes the maximum weight between the edges in any spanning tree. Time: $O(e \cdot \log(v))$

```
vii adj[MAX];

int minimax(int u, int N) {
    set<int> C;
    C.insert(u);

    priority_queue<ii, vii, greater<ii>> pq;

    for (auto [v, w] : adj[u]) pq.push(ii(w, v));

    int minmax = -oo;

    while ((int)C.size() < N) {
        int v, w;

        do {
            w = pq.top().first, v = pq.top().second;
            pq.pop();
        } while (C.count(v));

        minmax = max(minmax, w);
        C.insert(v);

        for (auto [s, p] : adj[v]) pq.push(ii(p, s));
    }

    return minmax;
}
```

5.15 MSF

Minimum Spanning Forest - a forest of trees of length k that connects all vertices in a graph with minimum total weight. Time: $O(e \cdot \log(v))$

```
using edge = tuple<int, int, int>;

int msf(int k, int N, vector<edge>& es) {
    sort(es.begin(), es.end());

    int cost = 0, cc = N;
    UnionFind udfs(N);

    for (auto [w, u, v] : es) {
        if (not udfs.same_set(u, v)) {
            cost += w;
            udfs.union_set(u, v);

            if (--cc == k) return cost;
        }
    }

    return cost;
}
```

5.16 Minimum Spanning Graph (MSG)

Given some obligatory edges es , find a minimum spanning graph that contains them. Time: $O(e \cdot \log(v))$

```
using edge = tuple<int, int, int>;

const int MAX{100010};

vector<ii> adj[MAX];

int msg(int N, const vector<edge>& es) {
    set<int> C;
    priority_queue<ii, vii, greater<ii>> pq;
    int cost = 0;

    for (auto [u, v, w] : es) {
        cost += w;

        C.insert(u);
        C.insert(v);

        for (auto [r, s] : adj[u]) pq.push(ii(s, r));
        for (auto [r, s] : adj[v]) pq.push(ii(s, r));
    }

    while ((int)C.size() < N) {
        int v, w;

        do {
            w = pq.top().first, v = pq.top().second;
            pq.pop();
        } while (C.count(v));

        cost += w;
        C.insert(v);

        for (auto [r, s] : adj[v]) pq.push(ii(s, r));
    }

    return cost;
}
```

```
cost += w;
C.insert(v);

for (auto [s, p] : adj[v]) pq.push(ii(p, s));
}

return cost;
}
```

5.17 Prim

A node u is chosen to start a connected component. Time: $O(e \cdot \log(v))$

```
const int MAX{100010};

vector<ii> adj[MAX];

int prim(int u, int N) {
    set<int> C;
    C.insert(u);

    priority_queue<ii, vector<ii>, greater<ii>> pq;

    for (auto [v, w] : adj[u]) pq.push(ii(w, v));

    int mst = 0;

    while ((int)C.size() < N) {
        int v, w;

        do {
            w = pq.top().first, v = pq.top().second;
            pq.pop();
        } while (C.count(v));

        mst += w;
        C.insert(v);

        for (auto [s, p] : adj[v]) pq.push(ii(p, s));
    }

    return mst;
}
```

5.18 Retrieve Path 2d

```
vll Graph::retrieve_path_2d(ll src, ll trg, const vector<vl>& pred) {
    vll p;

    do {
        p.emplace_back(pred[src][trg], trg);
        trg = pred[src][trg];
    } while (trg != src);

    reverse(all(p));
}
```

```

    return p;
}

```

5.19 Retrieve Path

```

vll Graph::retrieve_path(ll src, ll trg, const vll& pred) {
    vll p;

    do {
        p.emplace_back(pred[trg], trg);
        trg = pred[trg];
    } while (trg != src);

    reverse(all(p));

    return p;
}

```

5.20 Second Best MST

Time: $O(v \cdot e)$

```

using edge = tuple<int, int, int>;

pair<int, vi> kruskal(int N, vector<edge>& es, int blocked = -1) {
    vi mst;
    UnionFind udfs(N);
    int cost = 0;

    for (int i = 0; i < (int)es.size(); ++i) {
        auto [w, u, v] = es[i];

        if (i != blocked and not udfs.same_set(u, v)) {
            cost += w;
            udfs.union_set(u, v);
            mst.emplace_back(i);
        }
    }

    return {(int)mst.size() == N - 1 ? cost : oo, mst};
}

int second_best(int N, vector<edge>& es) {
    sort(es.begin(), es.end());

    auto [_, mst] = kruskal(N, es);
    int best = oo;

    for (auto blocked : mst) {
        auto [cost, __] = kruskal(N, es, blocked);
        best = min(best, cost);
    }

    return best;
}

```

5.21 TopSort - Tarjan

Works only on Directed Acyclic Graphs (DAGs). For each edge (u,v), u comes before v in the ordering. If the task A is a prerequisite for task B, then A comes before B in the ordering. Time: $O(V + E)$

```

enum State { NOT_FOUND, FOUND, PROCESSED };

vi adj[MAX];

bool dfs(int u, vi& o, vi& state) {
    if (state[u] == PROCESSED) return true;

    if (state[u] == FOUND) return false;

    state[u] = FOUND;

    for (auto v : adj[u])
        if (not dfs(v, o, state)) return false;

    state[u] = PROCESSED;
    o.emplace_back(u);

    return true;
}

vi topological_sort(int N) {
    vi o, state(N + 1, NOT_FOUND);

    for (int u = 1; u <= N; ++u)
        if (state[u] == NOT_FOUND and not dfs(u, o, state)) return {};

    reverse(o.begin(), o.end());

    return o;
}

```

6 Math

6.1 Binomial

```

ll binom(ll n, ll k) {
    if (k > n) return 0;
    vll dp(k + 1, 0);
    dp[0] = 1;
    for (ll i = 1; i <= n; i++)
        for (ll j = k; j > 0; j--) dp[j] = dp[j] + dp[j - 1];
    return dp[k];
}

```

6.2 Count Divisors Range

```

vl divisors(MAX, 0);
void count_divisors_range() {
    for (ll i = 1; i <= MAX; i++) {
        for (ll j = 1; j * i <= MAX; j++) divisors[i * j]++;
    }
}

```

6.3 Count Divisors

```
ll count_divisors(ll num) {
    ll count = 1;
    for (int i = 2; (ll)i * i <= num; i++) {
        if (num % i == 0) {
            int e = 0;
            do {
                e++;
                num /= i;
            } while (num % i == 0);
            count *= e + 1;
        }
    }
    if (num > 1) {
        count *= 2;
    }
    return count;
}
```

6.4 Factorization With Sieve

```
map<ll, ll> factorization_with_sieve(ll n, const vl& primes) {
    map<ll, ll> fact;

    for (ll d : primes) {
        if (d * d > n) break;

        ll k = 0;
        while (n % d == 0) {
            k++;
            n /= d;
        }

        if (k) fact[d] = k;
    }

    if (n > 1) fact[n] = 1;
    return fact;
}
```

6.5 Factorization

```
map<ll, ll> factorization(ll n) {
    map<ll, ll> ans;
    for (ll i = 2; i * i <= n; i++) {
        ll count = 0;
        for (; n % i == 0; count++, n /= i)
            ;
        if (count) ans[i] = count;
    }
    if (n > 1) ans[n]++;
    return ans;
}
```

6.6 Fast Doubling - Fibonacci

The Doubling Method can be seen as an improvement to the matrix exponentiation method to find the N-th Fibonacci number.

Time: $O(\log N)$.

```
template <typename T>
class FastDoubling {
public:
    vector<T> sts;
    T a, b, c, d;
    int mod;

    FastDoubling(int mod = 1e9 + 7) : sts(2), mod(mod) {}

    T fib(T x) {
        fill(all(sts), 0);
        a = 0, b = 0, c = 0, d = 0;
        fast_doubling(x, sts);
        return sts[0];
    }

    void fast_doubling(T n, vector<T>& res) {
        if (n == 0) {
            res[0] = 0;
            res[1] = 1;
            return;
        }
        fast_doubling(n >> 1, res);

        a = res[0];
        b = res[1];
        c = (b << 1) - a;

        if (c < 0) c += mod;

        c = (a * c) % mod;
        d = (a * a + b * b) % mod;
        if (n & 1) {
            res[0] = d;
            res[1] = c + d;
        } else {
            res[0] = c;
            res[1] = d;
        }
    }
};
```

6.7 Fast Exp Iterative

```
ll fast_exp_it(ll a, ll n, ll mod = LLONG_MAX) {
    a %= mod;
    ll res = 1;

    while (n) {
        if (n & 1) (res *= a) %= mod;

        (a *= a) %= mod;
        n >>= 1;
    }
}
```

```

    return res;
}

```

6.8 Fast Exp

```

11 fast_exp(11 a, 11 n, 11 mod = LLONG_MAX) {
    if (n == 0) return 1;
    if (n == 1) return a;

    11 x = fast_exp(a, n / 2, mod) % mod;

    return ((x * x) % mod * (n & 1 ? a : 1)) % mod;
}

```

6.9 Fast Fourier Transform (FFT)

Time: $O(N \cdot \log N)$

```

using cd = complex<double>;
const double PI = acos(-1);

void fft(vector<cd>& a, bool invert) {
    int n = a.size();

    for (int i = 1, j = 0; i < n; i++) {
        int bit = n >> 1;
        for (; j & bit; bit >>= 1) j ^= bit;
        j ^= bit;

        if (i < j) swap(a[i], a[j]);
    }

    for (int len = 2; len <= n; len <<= 1) {
        double ang = 2 * PI / len * (invert ? -1 : 1);
        cd wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i += len) {
            cd w(1);
            for (int j = 0; j < len / 2; j++) {
                cd u = a[i + j], v = a[i + j + len / 2] * w;
                a[i + j] = u + v;
                a[i + j + len / 2] = u - v;
                w *= wlen;
            }
        }
    }

    if (invert) {
        for (cd& x : a) x /= n;
    }
}

void fft_2d(vector<vector<cd>>& V, bool invert) {
    for (int i = 0; i < V.size(); i++) fft(V[i], invert);
    for (int i = 0; i < V.front().size(); i++) {
        vector<cd> col(V.size());
        for (int k = 0; k < V.size(); k++) col[k] = V[k][i];
        fft(col, invert);
        for (int k = 0; k < V.size(); k++) V[k][i] = col[k];
    }
}

```

```

}

```

6.10 GCD

The Euclidean algorithm allows to find the greatest common divisor of two numbers a and b in $O(\log \cdot \min(a, b))$.

```

11 gcd(11 a, 11 b) { return b ? gcd(b, a % b) : a; }

```

6.11 Integer Mod

```

const 11 INF = 1e18;
const 11 mod = 998244353;
template <11 MOD = mod>

struct Modular {
    11 value;
    static const 11 MOD_value = MOD;

    Modular(11 v = 0) {
        value = v % MOD;
        if (value < 0) value += MOD;
    }
    Modular(11 a, 11 b) : value(0) {
        *this += a;
        *this /= b;
    }

    Modular& operator+=(Modular const& b) {
        value += b.value;
        if (value >= MOD) value -= MOD;
        return *this;
    }
    Modular& operator-=(Modular const& b) {
        value -= b.value;
        if (value < 0) value += MOD;
        return *this;
    }
    Modular& operator*=(Modular const& b) {
        value = (11)value * b.value % MOD;
        return *this;
    }

    friend Modular mexp(Modular a, 11 e) {
        Modular res = 1;
        while (e) {
            if (e & 1) res *= a;
            a *= a;
            e >>= 1;
        }
        return res;
    }

    friend Modular inverse(Modular a) { return mexp(a, MOD - 2); }

    Modular& operator/=(Modular const& b) { return *this *= inverse(b); }
    friend Modular operator+(Modular a, Modular const b) { return a += b; }
}

```



```

Modular operator++(int) { return this->value = (this->value + 1) % MOD; }
Modular operator++() { return this->value = (this->value + 1) % MOD; }
friend Modular operator-(Modular a, Modular const b) { return a -= b; }
friend Modular operator-(Modular const a) { return 0 - a; }
Modular operator--(int) {
    return this->value = (this->value - 1 + MOD) % MOD;
}

Modular operator--() { return this->value = (this->value - 1 + MOD) % MOD; }
friend Modular operator*(Modular a, Modular const b) { return a *= b; }
friend Modular operator/(Modular a, Modular const b) { return a /= b; }
friend std::ostream& operator<<(std::ostream& os, Modular const& a) {
    return os << a.value;
}
}

friend bool operator==(Modular const& a, Modular const& b) {
    return a.value == b.value;
}
}

friend bool operator!=(Modular const& a, Modular const& b) {
    return a.value != b.value;
}
}
};

```

6.12 Is prime

$O(\sqrt{N})$

```

bool isprime(ll n) {
    if (n < 2) return false;
    if (n == 2) return true;
    if (n % 2 == 0) return false;
    for (ll i = 3; i * i <= n; i += 2)
        if (n % i == 0) return false;
    return true;
}

```

6.13 LCM

Calculating the least common multiple (commonly denoted LCM) can be reduced to calculating the GCD with the following simple formula: $\text{lcm}(a, b) = (a \cdot b) / \text{gcd}(a, b)$
Thus, LCM can be calculated using the Euclidean algorithm with the same time complexity:

```

ll lcm(ll a, ll b) { return a / gcd(a, b) * b; }

```

6.14 Euler phi $\varphi(n)$

Computes the number of positive integers less than n that are co-primes with n , in $O(\sqrt{N})$.

```

ll phi(ll n) {
    if (n == 1) return 1;

    auto fs = factorization(n);
    auto res = n;

    for (auto [p, k] : fs) {
        res /= p;
        res *= (p - 1);
    }
}

```

```

    return res;
}

```

6.15 Sieve

```

vl sieve(ll N) {
    bitset<MAX + 1> sieve;
    vl ps{2, 3};
    sieve.set();

    for (ll i = 5, step = 2; i <= N; i += step, step = 6 - step) {
        if (sieve[i]) {
            ps.push_back(i);

            for (ll j = i * i; j <= N; j += 2 * i) sieve[j] = false;
        }
    }
    return ps;
}

```

6.16 Sum Divisors

```

ll sum_divisors(ll num) {
    ll result = 1;

    for (int i = 2; (ll)i * i <= num; i++) {
        if (num % i == 0) {
            int e = 0;
            do {
                e++;
                num /= i;
            } while (num % i == 0);

            ll sum = 0, pow = 1;
            do {
                sum += pow;
                pow *= i;
            } while (e-- > 0);
            result *= sum;
        }
    }
    if (num > 1) {
        result *= (1 + num);
    }
    return result;
}

```

6.17 Sum of difference

Function to calculate sum of absolute difference of all pairs in array: $\frac{1}{2} \sum_{i=1}^N \sum_{j=1}^N |A_i - A_j|$

```

ll sum_of_difference(vl& arr, ll n) {
    sort(all(arr));

    ll sum = 0;
    for (ll i = 0; i < n; i++) {
        sum += i * arr[i] - (n - 1 - i) * arr[i];
    }
}

```

```

}

return sum;
}

```

7 Problems

7.1 Kth Digit String (CSES)

Time: $O(\log_{10} K)$.
Space: $O(1)$.

```

ll kth_digit_string(ll k) {
    if (k < 10) return k;

    ll c = 180, i = 2, u = 10, r = 0, ans = -1, m;
    for (k -= 9; k > c; i++, u *= 10) {
        k -= c;
        c /= i;
        c *= 10 * (i + 1);
    }

    if ((m = k % i))
        r++;
    else
        m = i;

    ll tmp = (k / i) + r + u - 1;
    for (m = i + 1 - m; m--; tmp /= 10) ans = tmp % 10;

    return ans;
}

```

7.2 Longest Common Substring (LONGCS - SPOJ)

Time: $N = \sum_{i=1}^k |S_i|$; $O(N \cdot \log N)$

```

int lcs_ks_strings(vector<string>& sts, int k) {
    vector<int> fml;
    string t;
    for (int i = 0; i < k; i++) {
        t += sts[i];
        for (int j = 0; j < sts[i].size(); j++) fml.push_back(i);
    }

    suffix_array sf(t);
    sf.lcp.insert(sf.lcp.begin(), 0);

    int l = 0, r = 0, cnt = 0, lcs = 0, n = sf.sa.size();
    vector<int> fr(k + 1);
    multiset<int> mst;
    while (l < n) {
        while (r < n and cnt < k) {
            mst.insert(sf.lcp[r]);
            if (!fr[fml[sf.sa[r]]]++) cnt++;
            r++;
        }
    }
}

```

```

}
mst.erase(mst.find(sf.lcp[l]));
if (mst.size() and cnt == k) lcs = max(lcs, *mst.begin());
fr[fml[sf.sa[l]]]--;
if (!fr[fml[sf.sa[l]]]) cnt--;
l++;
}

return lcs;
}

```

7.3 Substring Order II (CSES)

Time: $O(M)$
 $M = 2 \cdot N - 1$
 $N = |S|$

```

// ALLOWS REPETITIONS
string kth_smallest_substring(const string& s, ll k) {
    /* uses /strings/suffix-automaton.cpp
    add 'cnt' and 'nmb' to state struct with (0, -1);
    => for new states 'not cloned': cnt = 1

    create 'order' vector to iterate by length in decreasing
    vector<pair<int, int>>: {len, id}
    => for each new state add to 'order' vector

    to do not allow repetitions:
    => remove: kth+=s.size, sort(order) for(l, p : order)
    => add: st[clone].cnt = 1 (sa_extend)

    */
    string ans;
    k += s.size();
    SuffixAutomaton sa(s);

    sort(all(order), greater<pair<int, int>>());
    // count and mark how many times a substring of a state occurs
    for (auto& [l, p] : order) sa.st[sa.st[p].link].cnt += sa.st[p].cnt;

    auto dfs = [&](auto&& self, int u) {
        if (sa.st[u].nmb != -1) return;

        sa.st[u].nmb = sa.st[u].cnt;
        for (int i = 0; i < 26; ++i) {
            if (sa.st[u].next[i]) {
                self(self, sa.st[u].next[i]);
                sa.st[u].cnt += sa.st[sa.st[u].next[i]].cnt;
            }
        }
    };
    dfs(dfs, 0);

    int u = 0;
    while (sa.st[u].nmb < k) {
        k -= sa.st[u].nmb;
        for (int i = 0; i < 26; i++) {
            if (sa.st[u].next[i]) {
                int v = sa.st[u].next[i];

```

```

    if (sa.st[v].cnt < k)
        k -= sa.st[v].cnt;
    else {
        ans.push_back(i + 'a');
        u = v;
        break;
    }
}
}

return ans;
}

```

8 Strings

8.1 Aho-Corasick

The Aho-Corasick algorithm allows us to quickly search for multiple patterns in a text. The set of pattern strings is also called a *dictionary*. We will denote the total length of its constituent strings by m and the size of the alphabet by k .

build: $O(m \cdot k)$

occurrences: $O(|s| + ans)$

```

const int K = 26;
struct Vertex {
    char pch;
    int next[K];
    bool check = false;
    int p = -1, lnk = -1, out = -1, ps = -1, d = 0;

    Vertex(int p = -1, char ch = '$') : p(p), pch(ch) {
        fill(begin(next), end(next), -1);
    }
};

class AhoCorasick {
public:
    int sz = 0; // number of strings added
    vector<Vertex> t;

    AhoCorasick() : t(1) {}

    void add_string(string const& s) {
        int v = 0, ds = 0;
        for (char ch : s) {
            int c = ch - 'a';
            if (t[v].next[c] == -1) {
                t[v].next[c] = t.size();
                t.emplace_back(v, ch);
            }
            v = t[v].next[c];
            t[v].d = ++ds;
        }
        t[v].check = true;
        t[v].ps = sz++;
    }
}

```

```

void build() {
    queue<int> qs;
    qs.push(0);
    while (qs.size()) {
        auto u = qs.front();
        qs.pop();

        if (!t[u].p or t[u].p == -1)
            t[u].lnk = 0;
        else {
            int k = t[t[u].p].lnk;
            int c = t[u].pch - 'a';
            while (t[k].next[c] == -1 and k) k = t[k].lnk;
            int ts = t[k].next[c];
            if (ts == -1)
                t[u].lnk = 0;
            else
                t[u].lnk = ts;
        }

        if (t[t[u].lnk].check)
            t[u].out = t[u].lnk;
        else
            t[u].out = t[t[u].lnk].out;

        for (auto v : t[u].next)
            if (v != -1) qs.push(v);
    }
}

void occurrences(string const& s, vector<vector<int>>& res) {
    // to just "count" replace 'res' vector with an int
    res.resize(sz);
    for (int i = 0, v = 0; i < s.size(); i++) {
        int c = s[i] - 'a';
        while (t[v].next[c] == -1 and v) v = t[v].lnk;
        int ts = t[v].next[c];
        if (ts == -1)
            continue;
        else
            v = t[v].next[c];

        int k = v;
        while (t[k].out != -1) {
            k = t[k].out;
            res[t[k].ps].emplace_back(i - t[k].d + 1);
        }
        if (t[v].check) res[t[v].ps].emplace_back(i - t[v].d + 1);
    }
}
};

```

8.2 Edit Distance

Returns the minimum number of operations (insert, delete, replace) to transform string a into string b .
Time: $O(M * N)$

```

int min_value(int x, int y, int z) { return min(min(x, y), z); }

int edit_distance(string str1, string str2) {
    int n = (int)str1.size(), m = (int)str2.size();
    int dp[m + 1][n + 1];

    for (int i = 0; i <= m; i++)
        for (int j = 0; j <= n; j++)
            if (i == 0)
                dp[i][j] = j;
            else if (j == 0)
                dp[i][j] = i;
            else if (str1[i - 1] == str2[j - 1])
                dp[i][j] = dp[i - 1][j - 1];
            else
                dp[i][j] = 1 + min_value(dp[i][j - 1], dp[i - 1][j], dp[i - 1][j - 1]);

    return dp[m][n];
}

```

8.3 LCP with Suffix Array

For a given string s we want to compute the longest common prefix (LCP) of two arbitrary suffixes with position i and j . In fact, let the request be to compute the LCP of the suffixes $p[i]$ and $p[j]$. Then the answer to this query will be $\min(lcp[i], lcp[i + 1], \dots, lcp[j - 1])$. Thus the problem is reduced to the RMQ. Time: $O(N)$.

```

vector<int> lcp_suffix_array(string const& s, vector<int> const& p) {
    int n = s.size();
    vector<int> rank(n, 0);
    for (int i = 0; i < n; i++) rank[p[i]] = i;

    int k = 0;
    vector<int> lcp(n - 1, 0);
    for (int i = 0; i < n; i++) {
        if (rank[i] == n - 1) {
            k = 0;
            continue;
        }
        int j = p[rank[i] + 1];
        while (i + k < n && j + k < n && s[i + k] == s[j + k]) k++;
        lcp[rank[i]] = k;
        if (k) k--;
    }
    return lcp;
}

```

8.4 Manacher

Given string s with length n . Find all the pairs (i, j) such that substring $s[i \dots j]$ is a palindrome. String t is a palindrome when $t = t_{rev}$ (t_{rev} is a reversed string for t).

Time: $O(N)$

```

vi manacher(string s) {
    string t;
    for (auto c : s) t += string("#") + c;
}

```

```

t = t + '#';

int n = t.size();
t = "$" + t + "~";

vi p(n + 2);
int l = 1, r = 1;
for (int i = 1; i <= n; i++) {
    p[i] = max(0, min(r - i, p[l + (r - i)]));
    while (t[i - p[i]] == t[i + p[i]]) p[i]++;
    if (i + p[i] > r) {
        l = i - p[i], r = i + p[i];
    }
    p[i]--;
}

return vi(begin(p) + 1, end(p) - 1);
}

```

8.5 Rabin Karp

```

vector<int> rabin_karp(string const& s, string const& t) {
    const int p = 31;
    const int m = 1e9 + 9;
    int S = s.size(), T = t.size();

    vector<long long> p_pow(max(S, T));
    p_pow[0] = 1;
    for (int i = 1; i < (int)p_pow.size(); i++) p_pow[i] = (p_pow[i - 1] * p) % m;

    vector<long long> h(T + 1, 0);
    for (int i = 0; i < T; i++)
        h[i + 1] = (h[i] + (t[i] - 'a' + 1) * p_pow[i]) % m;
    long long h_s = 0;
    for (int i = 0; i < S; i++) h_s = (h_s + (s[i] - 'a' + 1) * p_pow[i]) % m;

    vector<int> occurrences;
    for (int i = 0; i + S - 1 < T; i++) {
        long long cur_h = (h[i + S] + m - h[i]) % m;
        if (cur_h == h_s * p_pow[i] % m) occurrences.push_back(i);
    }

    return occurrences;
}

```

8.6 Suffix Array Optimized - $O(n)$

Suffix Array: sa
Rank for LCP: rnk
LCP: lcp
Time: $O(N)$.

```

// @brunomaletta
struct suffix_array {
    string s;
    int n;
}

```

```

vector<int> sa, cnt, rnk, lcp;
rmq<int> RMQ; // /data-structures/rmq.cpp

bool cmp(int a1, int b1, int a2, int b2, int a3 = 0, int b3 = 0) {
    return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3);
}

template <typename T>
void radix(int* fr, int* to, T* r, int N, int k) {
    cnt = vector<int>(k + 1, 0);
    for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;
    for (int i = 1; i <= k; i++) cnt[i] += cnt[i - 1];
    for (int i = N - 1; i + 1; i--) to[--cnt[r[fr[i]]]] = fr[i];
}

void rec(vector<int>& v, int k) {
    auto &tmp = rnk, &m0 = lcp;
    int N = v.size() - 3, sz = (N + 2) / 3, sz2 = sz + N / 3;
    vector<int> R(sz2 + 3);
    for (int i = 1, j = 0; j < sz2; i += i % 3) R[j++] = i;

    radix(&R[0], &tmp[0], &v[0] + 2, sz2, k);
    radix(&tmp[0], &R[0], &v[0] + 1, sz2, k);
    radix(&R[0], &tmp[0], &v[0] + 0, sz2, k);

    int dif = 0;
    int l0 = -1, l1 = -1, l2 = -1;
    for (int i = 0; i < sz2; i++) {
        if (v[tmp[i]] != l0 or v[tmp[i] + 1] != l1 or v[tmp[i] + 2] != l2)
            l0 = v[tmp[i]], l1 = v[tmp[i] + 1], l2 = v[tmp[i] + 2], dif++;
        if (tmp[i] % 3 == 1)
            R[tmp[i] / 3] = dif;
        else
            R[tmp[i] / 3 + sz] = dif;
    }

    if (dif < sz2) {
        rec(R, dif);
        for (int i = 0; i < sz2; i++) R[sa[i]] = i + 1;
    } else
        for (int i = 0; i < sz2; i++) sa[R[i] - 1] = i;

    for (int i = 0, j = 0; j < sz2; i++)
        if (sa[i] < sz) tmp[j++] = 3 * sa[i];
    radix(&tmp[0], &m0[0], &v[0], sz, k);
    for (int i = 0; i < sz2; i++)
        sa[i] = sa[i] < sz ? 3 * sa[i] + 1 : 3 * (sa[i] - sz) + 2;

    int at = sz2 + sz - 1, p = sz - 1, p2 = sz2 - 1;
    while (p >= 0 and p2 >= 0) {
        if ((sa[p2] % 3 == 1 and
            cmp(v[m0[p]], v[sa[p2]], R[m0[p] / 3], R[sa[p2] / 3 + sz])) or
            (sa[p2] % 3 == 2 and
            cmp(v[m0[p]], v[sa[p2]], v[m0[p] + 1], v[sa[p2] + 1],
                R[m0[p] / 3 + sz], R[sa[p2] / 3 + 1])))
            sa[at--] = sa[p2--];
        else
            sa[at--] = m0[p--];
    }
}

```

```

}
while (p >= 0) sa[at--] = m0[p--];
if (N % 3 == 1)
    for (int i = 0; i < N; i++) sa[i] = sa[i + 1];
}

suffix_array(const string& s_)
: s(s_), n(s.size()), sa(n + 3), cnt(n + 1), rnk(n), lcp(n - 1) {
    vector<int> v(n + 3);
    for (int i = 0; i < n; i++) v[i] = i;
    radix(&v[0], &rnk[0], &s[0], n, 256);
    int dif = 1;
    for (int i = 0; i < n; i++)
        v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i - 1]]);
    if (n >= 2) rec(v, dif);
    sa.resize(n);

    for (int i = 0; i < n; i++) rnk[sa[i]] = i;
    for (int i = 0, k = 0; i < n; i++, k -= !!k) {
        if (rnk[i] == n - 1) {
            k = 0;
            continue;
        }
        int j = sa[rnk[i] + 1];
        while (i + k < n and j + k < n and s[i + k] == s[j + k]) k++;
        lcp[rnk[i]] = k;
    }
    RMQ = rmq<int>(lcp);
}

int query(int i, int j) {
    if (i == j) return n - i;
    i = rnk[i], j = rnk[j];
    return RMQ.query(min(i, j), max(i, j) - 1);
}
};

```

8.7 Suffix Array

Let s be a string of length n . The i -th suffix of s is the substring $s[i \dots n - 1]$.

A suffix array will contain integers that represent the starting indexes of the all the suffixes of a given string, after the aforementioned suffixes are sorted.

Time: $O(N \log N)$.

```

vector<int> sort_cyclic_shifts(string const& s) {
    int n = s.size();
    const int alphabet = 128;

    vector<int> p(n), c(n), cnt(max(alphabet, n), 0);
    for (int i = 0; i < n; i++) cnt[s[i]]++;
    for (int i = 1; i < alphabet; i++) cnt[i] += cnt[i - 1];
    for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
    c[p[0]] = 0;
    int classes = 1;
    for (int i = 1; i < n; i++) {
        if (s[p[i]] != s[p[i - 1]]) classes++;
        c[p[i]] = classes - 1;
    }
}

```

```

vector<int> pn(n), cn(n);
for (int h = 0; (1 << h) < n; ++h) {
    for (int i = 0; i < n; i++) {
        pn[i] = p[i] - (1 << h);
        if (pn[i] < 0) pn[i] += n;
    }
    fill(cnt.begin(), cnt.begin() + classes, 0);
    for (int i = 0; i < n; i++) cnt[c[pn[i]]]++;
    for (int i = 1; i < classes; i++) cnt[i] += cnt[i - 1];
    for (int i = n - 1; i >= 0; i--) p[--cnt[c[pn[i]]]] = pn[i];
    cn[p[0]] = 0;
    classes = 1;
    for (int i = 1; i < n; i++) {
        pair<int, int> cur = {c[p[i]], c[(p[i] + (1 << h)) % n]};
        pair<int, int> prev = {c[p[i - 1]], c[(p[i - 1] + (1 << h)) % n]};
        if (cur != prev) ++classes;
        cn[p[i]] = classes - 1;
    }
    c.swap(cn);
}

return p;
}

vector<int> suffix_array(string s) {
    s += "$";
    vector<int> p = sort_cyclic_shifts(s);
    p.erase(p.begin());
    return p;
}

```

8.8 Suffix Automaton

```

class SuffixAutomaton {
public:
    struct state {
        int len, link;
        array<int, 26> next;
    };

    vector<state> st;
    int sz = 0, last;

    SuffixAutomaton(const string& s) : st(s.size() << 1) {
        sa_init();
        for (auto v : s) sa_extend((int)(v - 'a'));
    }

    void sa_init() {
        st[0].len = 0;
        st[0].link = -1;
        sz++;
        last = 0;
    }

    void sa_extend(int c) {

```

```

        int cur = sz++;
        st[cur].len = st[last].len + 1;
        int p = last;
        while (p != -1 && !st[p].next[c]) {
            st[p].next[c] = cur;
            p = st[p].link;
        }
        if (p == -1)
            st[cur].link = 0;
        else {
            int q = st[p].next[c];
            if (st[p].len + 1 == st[q].len)
                st[cur].link = q;
            else {
                int clone = sz++;
                st[clone].len = st[p].len + 1;
                st[clone].link = st[q].link;
                st[clone].next = st[q].next;
                while (p != -1 && st[p].next[c] == q) {
                    st[p].next[c] = clone;
                    p = st[p].link;
                }
                st[q].link = st[cur].link = clone;
            }
        }
        last = cur;
    }

    // longest common substring O(N)
    int lcs(const string& T) {
        int v = 0, l = 0, best = 0;
        for (int i = 0; i < T.size(); i++) {
            while (v && !st[v].next[T[i] - 'a']) {
                v = st[v].link;
                l = st[v].len;
            }
            if (st[v].next[T[i] - 'a']) {
                v = st[v].next[T[i] - 'a'];
                l++;
            }
            best = max(best, l);
        }
        return best;
    }
};

```

8.9 Suffix Tree (CP Algo - freopen)

Build: $O(N)$

Memory: $O(N \cdot k)$

k = alphabet length

```

const int aph = 27; // add $ to final of string
const int N = 2e5 + 31;
class SuffixTree {
public:
    string a;
    vector<array<int, aph>> t;
    vector<int> l, r, p, s, dst;

```

```

int tv, tp, ts, la, b;

SuffixTree(const string& str, char bs = 'a') : a(str), t(N), l(N),
    r(N, str.size() - 1), p(N), s(N), dst(N), b(bs) {
    build();
}

void ukkadd(int c) {
suff:;
    if (r[tv] < tp) {
        if (t[tv][c] == -1) {
            t[tv][c] = ts; l[ts] = la;
            p[ts++] = tv; tv = s[ts]; tp = r[tv] + 1; goto suff;
        }
        tv = t[tv][c]; tp = l[tv];
    }
    if (tp == -1 || c == a[tp] - b) tp++; else {
        l[ts + 1] = la; p[ts + 1] = ts;
        l[ts] = l[tv]; r[ts] = tp - 1; p[ts] = p[tv];
        t[ts][c] = ts + 1; t[ts][a[tp] - b] = tv; l[tv] = tp;
        p[tv] = ts; t[p[ts]][a[l[ts]] - b] = ts; ts += 2;
        tv = s[p[ts - 2]]; tp = l[ts - 2];
        while (tp <= r[ts - 2]) {
            tv = t[tv][a[tp] - b];
            tp += r[tv] - l[tv] + 1;
        }
        if (tp == r[ts - 2] + 1) s[ts - 2] = tv; else s[ts - 2] = ts;
        tp = r[tv] - (tp - r[ts - 2]) + 2; goto suff;
    }
}

void build() {
    ts = 2; tv = 0; tp = 0;
    s[0] = 1; l[0] = -1; r[0] = -1; l[1] = -1; r[1] = -1;
    for (auto& arr : t) { arr.fill(-1); } t[1].fill(0);
    for (la = 0; la < (int)a.size(); ++la) ukkadd(a[la] - b);
}
};

```

8.10 Z Function

Suppose we are given a string s of length n . The Z-function for this string is an array of length n where the i -th element is equal to the greatest number of characters starting from the position i that coincide with the first characters of s .

Time: $O(N)$

```

vector<int> z_function(string s) {
    int n = s.size();
    vector<int> z(n);
    int l = 0, r = 0;
    for (int i = 1; i < n; i++) {
        if (i < r) {
            z[i] = min(r - i, z[i - l]);
        }
        while (i + z[i] < n && s[z[i]] == s[i + z[i]]) {
            z[i]++;
        }
        if (i + z[i] > r) {

```

```

            l = i;
            r = i + z[i];
        }
    }
    return z;
}

```

9 Trees

9.1 LCA Binary Lifting (CP Algo)

The algorithm described will need $O(N \cdot \log N)$ for preprocessing the tree, and then $O(\log N)$ for each LCA query.

```

ll n, l;
vector<ll> adj[MAX];

ll timer;
vector<ll> tin, tout;
vector<vector<ll>> up;

void dfs(ll v, ll p) {
    tin[v] = ++timer;
    up[v][0] = p;
    for (ll i = 1; i <= l; ++i) up[v][i] = up[up[v][i - 1]][i - 1];

    for (ll u : adj[v]) {
        if (u != p) dfs(u, v);
    }

    tout[v] = ++timer;
}

bool is_ancestor(ll u, ll v) { return tin[u] <= tin[v] && tout[u] >= tout[v]; }

ll lca(ll u, ll v) {
    if (is_ancestor(u, v)) return u;
    if (is_ancestor(v, u)) return v;
    for (ll i = l; i >= 0; --i) {
        if (!is_ancestor(up[u][i], v)) u = up[u][i];
    }
    return up[u][0];
}

void preprocess(ll root) {
    tin.resize(n);
    tout.resize(n);
    timer = 0;
    l = ceil(log2(n));
    up.assign(n, vector<ll>(l + 1));
    dfs(root, root);
}

```

9.2 LCA SegTree (CP Algo)

The algorithm can answer each query in $O(\log N)$ with preprocessing in $O(N)$ time.

```

struct LCA {
    vector<ll> height, euler, first, segtree;
    vector<bool> visited;
    ll n;

    LCA(vector<vector<ll>>& adj, ll root = 0) {
        n = adj.size();
        height.resize(n);
        first.resize(n);
        euler.reserve(n * 2);
        visited.assign(n, false);
        dfs(adj, root);
        ll m = euler.size();
        segtree.resize(m * 4);
        build(1, 0, m - 1);
    }

    void dfs(vector<vector<ll>>& adj, ll node, ll h = 0) {
        visited[node] = true;
        height[node] = h;
        first[node] = euler.size();
        euler.push_back(node);
        for (auto to : adj[node]) {
            if (!visited[to]) {
                dfs(adj, to, h + 1);
                euler.push_back(node);
            }
        }
    }

    void build(ll node, ll b, ll e) {
        if (b == e) {
            segtree[node] = euler[b];
        } else {
            ll mid = (b + e) / 2;
            build(node << 1, b, mid);
            build(node << 1 | 1, mid + 1, e);
            ll l = segtree[node << 1], r = segtree[node << 1 | 1];
            segtree[node] = (height[l] < height[r]) ? l : r;
        }
    }

    ll query(ll node, ll b, ll e, ll L, ll R) {
        if (b > R || e < L) return -1;
        if (b >= L && e <= R) return segtree[node];
        ll mid = (b + e) >> 1;

        ll left = query(node << 1, b, mid, L, R);
        ll right = query(node << 1 | 1, mid + 1, e, L, R);
        if (left == -1) return right;
        if (right == -1) return left;
        return height[left] < height[right] ? left : right;
    }

    ll lca(ll u, ll v) {
        ll left = first[u], right = first[v];
        if (left > right) swap(left, right);
        return query(1, 0, euler.size() - 1, left, right);
    }
}

```

```

    }
};

```

9.3 LCA Sparse Table

The algorithm described will need $O(N)$ for preprocessing, and then $O(1)$ for each LCA query.
0 indexed!

```

typedef vector<vl> vl2d;
#define all(a) a.begin(), a.end()
#define len(x) (int)x.size()

template <typename T>
struct SparseTable {
    vector<T> v;
    ll n;
    static const ll b = 30;
    vl mask, t;

    ll op(ll x, ll y) { return v[x] < v[y] ? x : y; }
    ll msb(ll x) { return __builtin_clz(1) - __builtin_clz(x); }
    SparseTable() {}
    SparseTable(const vector<T>& v_) : v(v_), n(v.size()), mask(n), t(n) {
        for (ll i = 0, at = 0; i < n; mask[i++] = at | = 1) {
            at = (at << 1) & ((1 << b) - 1);
            while (at and op(i, i - msb(at & -at)) == i) at ^= at & -at;
        }
        for (ll i = 0; i < n / b; i++)
            t[i] = b * i + b - 1 - msb(mask[b * i + b - 1]);
        for (ll j = 1; (1 << j) <= n / b; j++)
            for (ll i = 0; i + (1 << j) <= n / b; i++)
                t[n / b * j + i] =
                    op(t[n / b * (j - 1) + i], t[n / b * (j - 1) + i + (1 << (j - 1))]);
    }

    ll small(ll r, ll sz = b) { return r - msb(mask[r] & ((1 << sz) - 1)); }
    T query(ll l, ll r) {
        if (r - l + 1 <= b) return small(r, r - l + 1);
        ll ans = op(small(l + b - 1), small(r));
        ll x = l / b + 1, y = r / b - 1;
        if (x <= y) {
            ll j = msb(y - x + 1);
            ans = op(ans, op(t[n / b * j + x], t[n / b * j + y - (1 << j) + 1]));
        }
        return ans;
    }
};

struct LCA {
    SparseTable<ll> st;
    ll n;
    vl v, pos, dep;

    LCA(const vl2d& g, ll root) : n(len(g)), pos(n) {
        dfs(root, 0, -1, g);
        st = SparseTable<ll>(vector<ll>(all(dep)));
    }

    void dfs(ll i, ll d, ll p, const vl2d& g) {

```



```

v.emplace_back(len(dep)) = i, pos[i] = len(dep), dep.emplace_back(d);
for (auto j : g[i])
    if (j != p) {
        dfs(j, d + 1, i, g);
        v.emplace_back(len(dep)) = i, dep.emplace_back(d);
    }
}

ll lca(ll a, ll b) {
    ll l = min(pos[a], pos[b]);
    ll r = max(pos[a], pos[b]);
    return v[st.query(l, r)];
}

ll dist(ll a, ll b) {
    return dep[pos[a]] + dep[pos[b]] - 2 * dep[pos[lca(a, b)]];
}
};

```

9.4 Tree Flatten

```

vll tree_flatten(ll root) {
    vl pre;
    pre.reserve(N);

    vll flat(N);
    ll timer = -1;
    auto dfs = [&](auto&& self, ll u, ll p) -> void {
        timer++;
        pre.push_back(u);
        for (auto [v, w] : adj[u])
            if (v != p) {
                self(self, v, u);
            }
        flat[u].second = timer;
    };
    dfs(dfs, root, -1);
    for (ll i = 0; i < (ll)N; i++) flat[pre[i]].first = i;
    return flat;
}

```

9.5 Tree Isomorph

Checks whether two tree are isomorph. The function *thash()* returns the hash of the tree (using centroids as special vertices). Two trees are isomorph if their hash are the same.

```
map<vector<int>, int> mhash;
```

```

struct tree {
    int n;
    vector<vector<int>> g;
    vector<int> sz, cs;

    tree(int n_) : n(n_), g(n_), sz(n_) {}

    void dfs_centroid(int v, int p) {
        sz[v] = 1;
        bool cent = true;

```

```

        for (int u : g[v])
            if (u != p) {
                dfs_centroid(u, v), sz[v] += sz[u];
                if (sz[u] > n / 2) cent = false;
            }
        if (cent and n - sz[v] <= n / 2) cs.push_back(v);
    }
    int fhash(int v, int p) {
        vector<int> h;
        for (int u : g[v])
            if (u != p) h.push_back(fhash(u, v));
        sort(h.begin(), h.end());
        if (!mhash.count(h)) mhash[h] = mhash.size();
        return mhash[h];
    }
    ll thash() {
        cs.clear();
        dfs_centroid(0, -1);
        if (cs.size() == 1) return fhash(cs[0], -1);
        ll h1 = fhash(cs[0], cs[1]), h2 = fhash(cs[1], cs[0]);
        return (min(h1, h2) << 30) + max(h1, h2);
    }
    void add(int a, int b) {
        g[a].emplace_back(b);
        g[b].emplace_back(a);
    }
};

```

10 Settings and macros

10.1 short-macro.cpp

```

#include <bits/stdc++.h>

using namespace std;

#ifdef DEBUG
#include "../settings-and-macros/debug.cpp"
#else
#define dbg(...)
#endif

typedef long long ll;
typedef pair<int, int> ii;

#define all(x) x.begin(), x.end()
#define vin(vt) for (auto &e : vt) cin >> e

auto solve() { }

int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);

    ll t = 1;
    //cin >> t;

```

```

while (t--) solve();

return 0;
}

```

10.2 macro.cpp

```

#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace __gnu_pbds;
#define ordered_set tree<int, null_type, less<int>, rb_tree_tag,
    tree_order_statistics_node_update>

using namespace std;

#ifdef DEBUG
#include "../settings-and-macros/debug.cpp"
#else
#define dbg(...)
#endif

typedef long long ll;
typedef pair<int, int> pii;
typedef pair<ll, ll> pll;
typedef vector<int> vi;
typedef vector<ll> vl;
typedef vector<pii> vii;
typedef vector<pll> vll;

#define fst first
#define snd second
#define all(x) x.begin(), x.end()
#define len(vt) (int)vt.size()
#define vin(vt) for (auto &e : vt) cin >> e
#define LSOne(S) ((S) & ~(S))
#define MSOne(S) (1ull << (63 - __builtin_clzll(S)))
#define fastio ios_base::sync_with_stdio(0); \
    cin.tie(0); \
    cout.tie(0)

const vii dir4 {{1,0},{-1,0},{0,1},{0,-1}};

auto solve() { }

int main() {
    fastio;

    ll t = 1;
    //cin >> t;

    while (t--) solve();

    return 0;
}

```

11 Theoretical guide

11.1 Notable Series

- Sum of the first n naturals:

$$S_n = \sum_{i=1}^n i = 1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2}$$

- Sum of the squares of the first n naturals:

$$S_n = \sum_{i=1}^n i^2 = 1^2 + 2^2 + 3^2 + \cdots + n^2 = \frac{n(n+1)(2n+1)}{6}$$

- Sum of the cubes of the first natural n :

$$S_n = \sum_{i=1}^n i^3 = 1^3 + 2^3 + 3^3 + \cdots + n^3 = \left[\frac{n(n+1)}{2} \right]^2$$

- Sum of the first n odd numbers:

$$S_n = \sum_{i=1}^n 2i - 1 = 1 + 3 + 5 + \cdots + (2n - 1) = n^2$$

11.2 Number of Different Substrings

$$\sum_{i=0}^{n-1} (n - p[i]) - \sum_{i=0}^{n-2} \text{lcp}[i] = \frac{n^2 + n}{2} - \sum_{i=0}^{n-2} \text{lcp}[i]$$

11.3 Exponent With Module

If a and m are coprime, then

$$a^n \equiv a^{n \bmod \phi(m)} \bmod m$$

Generally, if $n \geq \log_2 m$, then

$$a^n \equiv a^{\phi(m) + [n \bmod \phi(m)]} \bmod m$$

11.4 String Matching with FFT

We are given two strings, a text T and a pattern P , consisting of lowercase letters. We have to compute all the occurrences of the pattern in the text.

We create a polynomial for each string ($T[i]$ and $P[i]$ are numbers between 0 and 25 corresponding to the 26 letters of the alphabet):

$$A(x) = a_0x^0 + a_1x^1 + \cdots + a_{n-1}x^{n-1}, \quad n = |T|$$

with

$$a_i = \cos(\alpha_i) + i \sin(\alpha_i), \quad \alpha_i = \frac{2\pi T[i]}{26}$$

And

$$B(x) = b_0x^0 + b_1x^1 + \cdots + b_{m-1}x^{m-1}, \quad m = |P|$$

with

$$b_i = \cos(\beta_i) - i \sin(\beta_i), \quad \beta_i = \frac{2\pi P[m - i - 1]}{26}$$

Notice that with the expression $P[m - i - 1]$ explicitly reverses the pattern.

The $(m - 1 + i)$ th coefficients of the product of the two polynomials $C(x) = A(x) \cdot B(x)$ will tell us, if the pattern appears in the text at position i .

If there isn't a match, then at least a character is different, which leads that one of the products

$a_{i+1} \cdot b_{m-1-j}$ is not equal to 1, which leads to the coefficient $c_{m-1+i} \neq m$.

11.4.1 Wildcards

This is an extension of the previous problem. This time we allow that the pattern contains the wildcard character $*$, which can match every possible letter.

We create the exact same polynomials, except that we set $b_i = 0$ if $P[m - i - 1] = *$. If x is the number of wildcards in P , then we will have a match of P in T at index i if $c_{m-1+i} = m - x$.

11.5 Pick's Theorem

Pick's Theorem expresses the area of a polygon, all whose vertices are lattice (integers) points in a coordinate plane, in terms of the number of lattice points inside the polygon and the number of lattice points on the sides (boundaries) of the polygon.

$$A = I + \frac{B}{2} - 1$$

- A: area of the polygon

- I: points inside the polygon
- B: points on the sides (boundaries)

It is possible to easily calculate the number of points on the sides of a side AB.

Consider $x = (x_1 - x_2)$ and $y = (y_1 - y_2)$. If $x = 0$ or $y = 0$, the answer is 1D and trivial (i.e. $x + 1$ or $y + 1$). Otherwise, the answer is $\gcd(a, b) + 1$.

11.6 Modular Multiplicative Inverse

A modular multiplicative inverse of an integer a is an integer x such that $a \cdot x$ is congruent to 1 modular some modulus m . To write it in a formal way:

$$a \cdot x \equiv 1 \pmod{m}.$$

Euler's theorem, which states that the following congruence is true if a and m are co-primes:

$$a^{\phi(m)} \equiv 1 \pmod{m}$$

Multiply both sides of the above equations by a^{-1} , and we get:

- For an arbitrary (but coprime) modulus m : $a^{\phi(m)-1} \equiv a^{-1} \pmod{m}$
- For a prime modulus m : $a^{m-2} \equiv a^{-1} \pmod{m}$

From these results, we can easily find the modular inverse using the binary exponentiation algorithm, which works in $O(\log m)$ time.