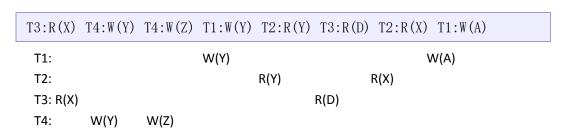
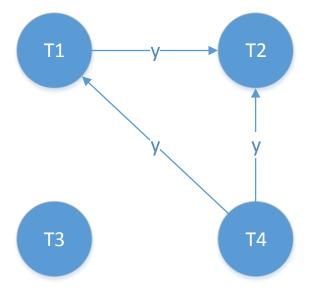
```
i.
       AB->C, C->DE, D->F
                                      Left and Right: C, D
        Left: A, B
        A^+ = \{A\} B^+ = \{B\} (AB)^+ = \{A, B, C, D, E, F\}
        Candidate key is AB
        C->DE, D->F are BCNF violations.
        Divide S{A, B, C, D, E, F} into S_1\{C, D, E, F\} and S_2\{A, B, C\}
        S<sub>2</sub> is in BCNF
        S<sub>1</sub>: C->DE, D->F
        Candidate key of S<sub>1</sub> is C
        D->F is BCNF violation.
        Divide R_1\{C, D, E, F\} into S_{11}\{D, F\} and S_{12}\{C, D, E\}
        The collections of BCNF relation:
              S_{11}\{D, F\}
              S_{12}\{C, D, E\}
              S_2\{A, B, C\}
       BC \rightarrow D, C \rightarrow AF, AB \rightarrow CE
ii.
       Left: B
                                      Left and Right: A, C
       B^+ = \{B\} (AB)<sup>+</sup> = {A, B, C, D, E, F} (BC)<sup>+</sup> = {A, B, C, D, E, F}
       Candidate key is AB and BC
       C->AF is BCNF violation.
       Divide S{A, B, C, D, E, F} into S_1{A, C, F} and S_2{A, B, C, D, E}
       S<sub>1</sub> is in BCNF
       S<sub>2</sub>: BC->D, AB->CE
       Candidate key of S2 is AB
        BC->D is BCNF violation
       Divide S_2\{A, B, C, D, E\} into S_{21}\{B, C, D\} and S_{22}\{A, B, C, E\}
       The collections of BCNF relation:
              S_1{A, C, F}
              S_{21}\{B, C, D\}
              S_{22+}\{A, B, C, E\}
       ABC \rightarrow D, CD \rightarrow E,
iii.
                                          BD \rightarrow A
       Left: B, C
                                      Left and Right: A, D
       BC^{+} = \{B, C\}
                          (ABC)^+ = \{A, B, C, D, E\} (BCD)^+ = \{A, B, C, D, E\}
        Because there is no dependency about F, F is part of the candidate key.
        Candidate key is ABCF and BCDF
       ABC->D (if do not consider about F, it is not BCNF violation), CD->E, BD->A are BCNF
       violations.
       Divide S{A, B, C, D, E, F} into S_1{A, B, C, D} and S_2{C, D, E, F}
       S<sub>1</sub>: ABC->D, BD->A
       Candidate key of S<sub>1</sub> is ABC and BCD
        Divide S_1 into S_{11}\{A, B, D\} and S_{12}\{B, C, D\}
       S<sub>2</sub>: CD->E
```

```
Candidate key of S2 is CDF
          Divide S_2 into S_{21}(C, D, E) and S_{22}(E, F)
          The collections of BCNF relation:
                S_{11}\{A, B, D\}
                S_{12}\{B, C, D\}
                S_{21}\{C, D, E\}
                S_{22}(E, F)
          AB \rightarrow DE, BCD \rightarrow EF, B \rightarrow C
  iv.
          Left: A, B,
                                           Left and Right: C, D
          (AB)^+ = \{A, B, C, D, E, F\}
          Candidate key is AB
          BCD->EF, B->C are BCNF violations
          Divide S{A, B, C, D, E, F} into S_1{B, C, D, E, F} and S_2{A, B, D}
          S<sub>2</sub> is in BCNF
          S<sub>1</sub>: BCD->EF, B->C
          Candidate key of S<sub>1</sub> is BD
          Divide S_1 into S_{11}\{B, C\} and S_{12}\{B, D, E, F\}
          The collections of BCNF relation:
                S_{11}\{B, C\}
                S_{12}\{B, D, E, F\}
                S_2\{A, B, D\}
2.
   i.
          Answer = Proj[person.name](
                person join[person.personid = relationpersoninproceeding.personid]
                relationpersoninproceeding
          )
   ii.
          temp1 = proj[person.personid, relationpersoninproceeding.inproceedingid](
                person join[person.personid = relationpersoninproceeding.personid]
                relationpersoninproceeding
          )
          temp 2 = proj[temp1.personid, inproceeding.proceedingid](
                inproceeding join[temp1.inproceedingid = inproceeding.inproceedingid] temp1
          )
          temp3 = proj[temp2.personid](
                proceeding join[temp2.proceedingid = proceeding.proceedingid and
                temp2.personid = proceeding.editorid] temp2
          )
          answer = proj[person.name](
```

```
person join[temp3.personid = person.personid] temp3
         )
  iii.
         temp1 = proj[personid](
               sel[name ~ 'clark$'](person)
         )
         temp2 = proj[relationpersoninproceeding.inproceedingid](
              temp1 join[temp.personid = relationpersoninproceeding.personid]
         ) relationpersoninproceeding
         answer = proj[inproceeding.title](
               temp2 join[temp2.inproceedingid = inproceeding.inproceedingid] inproceeding
  iv.
         Re1 = relationpersoninproceeding
         Re2 = relationpersoninproceeding
         Inpro = inproceeding
         Per = person
         Pro = proceeding
         temp1 = Proj[personid](Re1) Minus Rename[editorid->personid](Proj[editorid](Pro))
         temp2 = Proj[Re1.personid](GroupSelect[size>1](
                        GroupBy[Re1.personid,Re1.inproceedingid]
                        (Re1 Join[Re1.inproceedingid = Re2.inproceedingid] Re2))
                   )
         temp3 = temp1 Minus temp2
         Answer = Proj[name](temp3 Join[temp3.personid=Per.personid] Per)
3.
   i.
         R INTERSECT (S UNION T).
         Max(s, t) \le |Sunion T| \le s+t
         0 \le |R \text{ intersect (S union T)}| \le Min(r, s+t)
         The minimum numbers of tuples is 0
         The maximum numbers of tuples is Min(r, s+t)
   ii.
         Sel[c](R) \times S, for some condition c.
         0 \le |Sel[c](R)| \le r
         0 \le |Sel[c](R) \times S| \le r \times s
         The minimum numbers of tuples is 0
         The maximum numbers of tuples is r \times s
         Proj[a](R) - S, for some list of attributes a.
  iii.
         |Proj[a](R)| = r
         Max(r-s, 0) \le |Proj[a](R) - s| \le r
         The minimum numbers of tuples is Max(r-s, 0)
         The maximum numbers of tuples is r
```





There is no circle in the graph, so it is schedule serialisable.