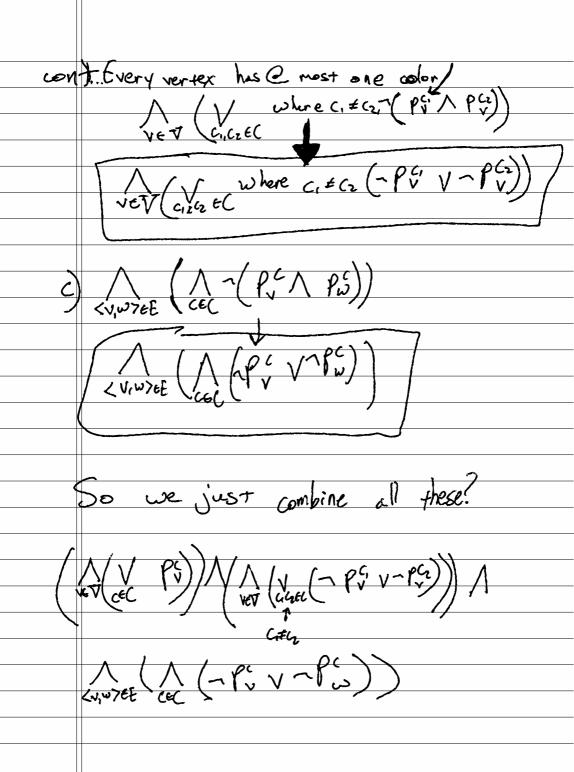
G=<V,E> V= E V, ,..., Vn3 E = { < Vi, wi, 7, ... , < Vin, win>} Given Set, length k, of colors C = \(\xi_1,...,c_k\) Assign CEC to each VEV S.t. every edge < v, w> EE, color(v) ≠ color(w) HU Encode K-Coloring into propositional Formula F. a) Every vertex v has a color c (V PC) ? This is saying every ver cel Verter must have a coloring Thouses (PC: VPCz VPCz) A (PC: VPCz VPCz) Vihas V, +V2 must have ador b) Every vertex has e most one colon/ VET C.CZEC Where C, #G, (PC,)



d) So constraint I is necessary?) as all other constraints rely on each vertex having a color. Do 2t 3 overlap? colors (x, Y, Z) Could put of pere & 3 wouldn't

Cutch it?

But this graph is still

colored, so it's of? If we are simply trying to find any voloring. I constraint 2 con le dropped, as you can just select either color. e) ???