CISC4080 Computer Algorithms Homework $(2)^1$

 $^{^{1} {\}rm Distinguish\ yourselves,\ folks!}$

Part A. Divide-and-Conquer Algorithm Complexity Analysis (20 points)

1. Give the time complexity with Θ bound for A divide-and conquer algorithm with the following recurrence formula

1.
$$T(n) = T(n-1) + c^n$$
, where $c > 1$

2.
$$T(n) = 16T(n/4) + n^3$$

3.
$$T(n) = 25T(n/36) + n^{3/2} \log n$$

2. Show the best case of merge sort has complexity $\Theta(n)$ without using the master method.

Part B. Targeted Sorting (15 points)

• 1. Implement and test a method (function) to check if input array is sorted, which means the array entries are sorted in an ascending or descending order. It can have the following signature²

bool isSorted(const int*, int &)

- 2. What are the worst and average time complexities of this method?
- 3. Combine this method with a quick-sort method implemented by your-self and test your sorting method with the best, worst and average cases for input arrays with $n = 10^4$ entries respectively³.

²You don't need to follow this signature "exactly" in your implementation

 $^{^3\}mathrm{You}$ may need to use some sample codes in your homework 1

Part C. A GCD Calculator (15 points)

Greatest common divisor

- If a|b|c|b, then we can say a is a common divisor of a and b. The largest factor of a and b is called the greatest common divisor of a and b: GCD(a,b). For examples, 20 = GCD(20,100), 2 = GCD(10,8)
- That is, GCD(a, b) is the largest positive integer dividing both a and b. If GCD(a, b) = 1, we say a and b are relatively prime.

How to compute GCD of a and b?

- The typical algorithm to compute a GCD is called Euclidean algorithm invented more than 2000 years ago!
- <u>Ideas of the Euclidean algorithm</u>: divident is replaced by divisor and divisor is replaced by remainder recursively!
- It has the following algorithm description
 - To compute the GCD of a and b, we assume a > b always. If not, just switch a and b.

```
Algorithm GCD(a,b)
Input: integer a, b (
Output: GCD(a,b)
% divide a by b:
a = qb + r
If r = 0
return b
else
a = b;
b = r;
GCD(a, b)
end
```

1. Implement Euclidean algorithm to write a GCD calculator and compute at least the following GCDs.

```
GCD(482, 1180)

GCD(8756, 23485)

GCD(87561, 23485)

GCD(1234567, 2008479)

GCD(578129810, 20092330)
```

- If you program in JAVA, remember to use long type for the dividend and divisor.
- The following codes and outputs may give you a little bit hint.
- You can write your own codes and may not need to follow these sample codes.

```
public static void main(String[] args) throws IOException{
   computeGCD myGCD=new computeGCD();
   System.out.println("Enter your Dividend:\n");
   long Dividend = myGCD.read_long();
   System.out.println("Enter your Divisor:\n");
   long Divisor = myGCD.read_long();
   myGCD=new computeGCD(Dividend, Divisor);
   long gcd = myGCD.gcd();
   System.out.println("Results: gcd(" + Dividend + ", " + Divisor + ") = " + gcd);
}
```

• Sample outputs

```
>java computeGCD
Enter your Dividend:
1234567
Enter your Divisor:
2008479
current divident-->1234567
current divident-->773912
current divident-->460655
current divident --> 313257
current divident-->147398
current divident-->18461
current divident-->18171
current divident-->290
current divident-->191
current divident-->99
current divident-->92
current divident-->7
current divident-->1
current divident-->0
Results: gcd(1234567, 2008479) = 1
```

2. What's the time complexity of this algorithm?⁴

⁴Hint: you can count the number of recursion calls to get some idea!

What should you turn in?

- 1. A hardcopy of all your homework printout in class (Oct 29, 2013).
- 2. A folder contains all your homework assignments. If there is a programming assignment, you need to include workable source codes and related output in this folder. Please name your folder as first-name_last-name_CISC4080_homework_2. For example, John_Smith_CISC4080_homework_2 if your name is John Smith.
- 3. Send the zipped file (.zip instead of .rar) of your folder to xhan9@fordham.edu before 11:59 pm Oct 29, 2013.