(1) Show that the Hilbert matrix is positive definite. *Proof:*

Suppose we have the $n \times n$ Hilbert matrix

$$H_{ij} = \frac{1}{i+j-1}, \quad i, j = 1, \dots, n.$$

Then, to show that H is positive definite, we need to show that $x^*Hx > 0$ for all $x \neq 0$ in \mathbb{C}^n . So, suppose we have any nonzero $x \in \mathbb{C}^n$. Then,

$$\begin{split} x^*Hx &= \sum_{i=1}^n \left(\bar{x}_i \sum_{j=1}^n x_j H_{ij} \right) = \sum_{i=1}^n \sum_{j=1}^n \bar{x}_i x_j \frac{1}{i+j-1} \\ &= \sum_{i=1}^n \sum_{j=1}^n \bar{x}_i x_j \int_0^1 t^{i+j-2} \mathrm{d}t \\ &= \sum_{i=1}^n \sum_{j=1}^n \int_0^1 (\bar{x}_i t^{i-1}) (x_j t^{j-1}) \mathrm{d}t \\ &= \int_0^1 \sum_{i=1}^n \sum_{j=1}^n (\bar{x}_i t^{i-1}) (x_j t^{j-1}) \mathrm{d}t \\ &= \int_0^1 \left(\sum_{i=1}^n \bar{x}_i t^{i-1} \right) \left(\sum_{j=1}^n x_j t^{j-1} \right) \mathrm{d}t. \end{split}$$

Now if we let $\alpha = \sum_{j=1}^{n} x_j t^{j-1}$, our equation becomes

$$x^* H x = \int_0^1 \left(\sum_{i=1}^n \bar{x}_i t^{i-1} \right) \left(\sum_{j=1}^n x_j t^{j-1} \right) dt$$
$$= \int_0^1 \bar{\alpha} \alpha dt$$
$$= \int_0^1 |\alpha|^2 dt$$
$$> 0$$

So, our Hilbert matrix must be positive definite.

(2) Using my power iteration code (attached in section: Code Used), I generated the table below of the largest eigenvalues for the first 10 square Hilbert matrices.

n	Eigenvalue	Eigenvector
1	1.00	$[1.0]^{T}$
2	1.27	$[0.88, 0.47]^T$
3	1.41	$[0.83, 0.46, 0.32]^T$
4	1.50	$[0.79, 0.45, 0.32, 0.25]^T$
5	1.57	$[0.77, 0.45, 0.32, 0.25, 0.21]^T$
6	1.62	[0.75, 0.44, 0.32, 0.25, 0.21, 0.18]
7	1.66	$[0.73, 0.44, 0.32, 0.25, 0.21, 0.18, 0.16]^T$
8	1.70	$[0.72, 0.43, 0.32, 0.26, 0.21, 0.18, 0.16, 0.15]^T$
9	1.73	$[0.71, 0.43, 0.32, 0.26, 0.21, 0.19, 0.16, 0.15, 0.13]^T$
10	1.75	$[0.70, 0.43, 0.32, 0.26, 0.22, 0.19, 0.16, 0.15, 0.13, 0.12]^T$

(3) To find the smallest eigenvalues using a power iteration, I just change the matrix-vector multiplication to a backslash as to solve Ax = b instead of computing x = Ab. This change has the effect of making the smallest eigenvalue of A dominate the power iteration. For n = 16, I obtained the eigenvalue:

$$\lambda_{\min} = -2.765686556840889 \cdot 10^{-18}$$

which is not even positive even though we know Hilbert matrices are symmetric positive definite. So my naive Julia code is definitely not converging properly under double precision. Putting my Julia code in quad precision yields the eigenvalue

$$\lambda_{\min} = 9.197419820719313 \cdot 10^{-23}.$$

which is consistent with Mathematica's eigenvalue. Hence, we can see that my computed double-precision eigenvalue is off by over 5 orders of magnitude. However, this result is still consistent with our error estimate

$$\min_{\lambda \in \sigma(A)} |\mu - \lambda| \leq \|E\|_2$$

because $||E||_2$ should be on the order of 10^{-16} which is the best we can do in double precision.

(4) Assume that a real matrix A has eigenvalues $\lambda_1 = -\lambda_2$ and $|\lambda_1| = |\lambda_2| > |\lambda_3| \ge |\lambda_n|$.

Without loss of generality, assume $\lambda_1>0$ which makes $\lambda_2<0$. Now, note that because A is symmetric, eigenvectors corresponding to different eigenvalues of A are orthogonal. Then, to find the eigenvectors corresponding to λ_1 and λ_2 , run the standard power iteration to find some normalized vector, \vec{v}_0 , in the span of $\vec{\lambda}_1$ and $\vec{\lambda}_2$ where $\vec{\lambda}_1$ and $\vec{\lambda}_2$ are eigenvectors corresponding to λ_1 and λ_2 respectively. Then, because $\vec{\lambda}_1$ and $\vec{\lambda}_2$ are orthogonal, we can uniquely decompose \vec{v}_0 as

$$\vec{v}_0 = a\vec{\lambda}_1 + b\vec{\lambda}_2$$

for some constants a and b. Then, one more power iteration on \vec{v}_0 will yield

$$\vec{v}_1 = a\vec{\lambda}_1 - b\vec{\lambda}_2$$

because $\lambda_2<0$ and the eigenvectors are orthogonal. Then, we can simply find an eigenvector corresponding to λ_1 as

$$\vec{v}_0 + \vec{v}_1 = 2a\vec{\lambda}_1$$

and an eigenvector corresponding to λ_2 as

$$\vec{v}_0 - \vec{v}_1 = 2b\vec{\lambda}_2.$$

(5) A real symmetric matrix A has an eigenvalue 1 of multiplicity 8; the rest of the eigenvalues are ≤ 0.1 in absolute value.

We can find an orthogonal basis for the 8-dimensional eigenspace corresponding to the dominant eigenvalue by combining the standard power iteration with the Gram-Schmidt process. To do so:

- (i) Generate 8, random vectors.
- (ii) Apply power iteration to each vector until convergence. The resulting vectors will be in the span of the dominant eigenvectors.
- (iii) Apply Gram-Schmidt to the set of generated vectors.
- (iv) You're all done!

To estimate how long it will take to find this basis for an $n \times n$ matrix, we need to figure out how long it will take each power iteration to converge to double-precision. If each subdominant eigenvalue has a magnitude less than or equal to 0.1, we should expect each power iteration to converge in 16 to 17 iterations because $|\lambda_{\text{next}}/\lambda_1| \leq 0.1$. Then, we need to run power iteration for each of the 8 vectors for a total of 136 iterations.

Code Used

Note that some characters are missing because \LaTeX can't display some of the unicode characters in my code.

```
2 # Power iteration method for finding eigenvalues
4 # Author: Caleb Jacobs
   Date last modified: 01-02-2022
8 using LinearAlgebra
9 using DoubleFloats
11 function makeHilbert(n, quad = false)
      # Check if we want quad-precision matrix
      if !quad
          A = [Float64(1) / (i + j - 1) for i = 1 : n, j = 1 : n]
15
          A = [Double64(1) / (i + j - 1) for i = 1 : n, j = 1 : n]
16
      end
17
      return A
20 end
21
22 function powIt(A, maxIts)
      q = z = rand(size(A, 1))
                                   # Initialize eigenvectors
23
       = 0
                                   # Initialize eigenvalue
24
      for k = 1 : maxIts
          z = A * q
                                    # Compute next eigenvector guess
          q = z / norm(z)
                                   # Normalize eigenvector
2.7
            = q' * A * q
                                   # Compute next eigenvalue guess
      end
29
      return (, q)
32 end
34 function invIt(A, maxIts)
      q = z = ones(Float64, size(A, 1))
                                             # Initialize eigenvectors
35
       = 0
                                              # Initialize eigenvalue
36
      # Perform inverse iteration to find smallest eigenvalue
      for k = 1 : maxIts
39
          z = A \setminus q
                                    # Compute next eigenvector guess
40
          q = z / norm(z)
                                   # Normalize eigenvector
41
            = dot(q, A, q)
                                   # Compute next eigenvalue guess
      end
44
      return
45
```

```
46 end
47
48 function prob2()
      a = 1;
                       # Lower size of matrix
      b = 10;
                       # Upper size of matrix
50
      maxIts = 100;
                     # Number of iterations for powIt
      data = Vector{Tuple{Float64, Vector{Float64}}}(undef, b - a + 1)
53
      n = [a : b...]
55
56
      for i 1:b-a+1
          A = makeHilbert(n[i])
59
          data[i] = powIt(A, maxIts)
      \quad \text{end} \quad
62
      for i
             data
63
          print(round.(i[1]; digits = 2))
64
          println(round.(i[2]; digits = 2))
65
      return(data)
67
 end
68
69
70 function prob3()
      AD = makeHilbert(16)
                                    # Double precision hilbert
71
      AQ = makeHilbert(16, true) # Quad precision hilbert
      doubVal = invIt(AD, 1000)
                                    # Double precision
74
      quadVal = invIt(AQ, 1000)
                                   # Quad precision
75
76
      return (doubVal, quadVal)
78 end
```