

(1) Show that the Hilbert matrix is positive definite.

Proof:

Suppose we have the $n \times n$ Hilbert matrix

$$H_{ij} = \frac{1}{i+j-1}, \quad i, j = 1, \dots, n.$$

Then, to show that H is positive definite, we need to show that $x^* H x > 0$ for all $x \neq 0$ in \mathbb{C}^n . So, suppose we have any nonzero $x \in \mathbb{C}^n$. Then,

$$\begin{aligned} x^* H x &= \sum_{i=1}^n \left(\bar{x}_i \sum_{j=1}^n x_j H_{ij} \right) = \sum_{i=1}^n \sum_{j=1}^n \bar{x}_i x_j \frac{1}{i+j-1} \\ &= \sum_{i=1}^n \sum_{j=1}^n \bar{x}_i x_j \int_0^1 t^{i+j-2} dt \\ &= \sum_{i=1}^n \sum_{j=1}^n \int_0^1 (\bar{x}_i t^{i-1}) (x_j t^{j-1}) dt \\ &= \int_0^1 \sum_{i=1}^n \sum_{j=1}^n (\bar{x}_i t^{i-1}) (x_j t^{j-1}) dt \\ &= \int_0^1 \left(\sum_{i=1}^n \bar{x}_i t^{i-1} \right) \left(\sum_{j=1}^n x_j t^{j-1} \right) dt. \end{aligned}$$

Now if we let $\alpha = \sum_{j=1}^n x_j t^{j-1}$, our equation becomes

$$\begin{aligned} x^* H x &= \int_0^1 \left(\sum_{i=1}^n \bar{x}_i t^{i-1} \right) \left(\sum_{j=1}^n x_j t^{j-1} \right) dt \\ &= \int_0^1 \bar{\alpha} \alpha dt \\ &= \int_0^1 |\alpha|^2 dt \\ &> 0. \end{aligned}$$

So, our Hilbert matrix must be positive definite. □

- (2) Using my power iteration code (attached in section: Code Used), I generated the table below of the largest eigenvalues for different size Hilbert matrices.

- (3) To find the smallest eigenvalues using a power iteration, I just change the matrix-vector multiplication to a backslash as to solve $Ax = b$ instead of computing $x = Ab$. This change has the effect of making the smallest eigenvalue of A dominate the power iteration. For $n = 16$, I obtained the eigenvalue:

$$\lambda_{\min} =$$

- (4) Assume that a real matrix A has eigenvalues $\lambda_1 = -\lambda_2$ and $|\lambda_1| = |\lambda_2| > |\lambda_3| \geq |\lambda_n|$.

Without loss of generality, assume $\lambda_1 > 0$ which makes $\lambda_2 < 0$. Now, note that because A , eigenvectors corresponding to different eigenvalues of A are orthogonal. Then, to find the eigenvectors corresponding to λ_1 and λ_2 , run the standard power iteration (until convergence) to find some normalized vector, \vec{v}_0 , in the span of $\vec{\lambda}_1$ and $\vec{\lambda}_2$ where $\vec{\lambda}_1$ and $\vec{\lambda}_2$ are eigenvectors corresponding to λ_1 and λ_2 respectively. Then, because $\vec{\lambda}_1$ and $\vec{\lambda}_2$ are orthogonal, we can uniquely decompose \vec{v}_0 as

$$\vec{v}_0 = a\vec{\lambda}_1 + b\vec{\lambda}_2$$

for some constants a and b . Then, one more power iteration on \vec{v}_0 will yield

$$\vec{v}_1 = a\vec{\lambda}_1 - b\vec{\lambda}_2$$

because $\lambda_2 < 0$ and the eigenvectors are orthogonal. Then, we can simply find an eigenvector corresponding to λ_1 as

$$\vec{v}_0 + \vec{v}_1 = 2a\vec{\lambda}_1$$

and an eigenvector corresponding to λ_2 as

$$\vec{v}_0 - \vec{v}_1 = 2b\vec{\lambda}_2.$$

- (5) A real symmetric matrix A has an eigenvalue 1 of multiplicity 8; the rest of the eigenvalues are ≤ 0.1 in absolute value.

We can find an orthogonal basis for the 8-dimensional eigenspace corresponding to the dominant eigenvalue by combining the standard power iteration with the Gram-Schmidt process. The idea is to pick a random initial eigenvector and then run the power iteration until convergence to some normalized vector, \vec{v}_1 in the span of the dominant eigenvectors.

Then, pick a new random vector and use Gram-Schmidt to remove the component of this new vector in the direction of \vec{v}_1 . Just as before, apply the power iteration to this new vector to get \vec{v}_2 which is orthogonal to \vec{v}_1 but still in the span of the dominant eigenvectors.

Repeat this process, subtracting each previously found vector out of the new vector using Gram-Schmidt and then apply power iteration again. Once you have 8 vectors, you have formed a basis for the dominant eigenspace.

To estimate how long it will take to find this basis for an $n \times n$ matrix, we need to figure out how long it will take each power iteration to converge to double-precision. If each subdominant eigenvalue has a magnitude less than or equal to 0.1, we should expect our iteration to converge in

$$(n - 8)(0.1)^i \leq 10^{-16}$$

or

$$i = \log_{0.1} \frac{10^{-16}}{n - 8}$$

iterations. This calculation comes from the fact that we need $n - 8$ vectors with $|\lambda| < 0.1$ to vanish from power iteration to have convergence to double-precision. Then, we need to do 8 power iterations for a total of

$$8i = 8 \log_{0.1} \frac{10^{-16}}{n - 8}$$

iterations minimum for convergence. For $n \leq 1000$, $i \leq 19$ so for any matrix smaller than 1000×1000 , we should expect convergence to our basis in a minimum of 152 iterations.

Code Used