

CSE 4833/6833 Introduction to Algorithms

Programming Assignment, due Thursday, Nov. 11, by midnight.

Worth up to 15 points extra credit on the final exam.

In Discrete Structures you used the binomial coefficient formula,

$$C(n, k) = \frac{n!}{k!(n - k)!}$$

to calculate the number of combinations of “ n choose k ,” that is, the number of ways to choose k objects from n objects. For example, the number of unique 5-card hands from a standard 52-card deck is $C(52, 5)$.

A problem with using the above formula to compute binomial coefficients is that $n!$ grows very fast and overflows an integer representation before you can do the division to bring the value back to a value that can be represented. When calculating $C(52, 5)$, for example, the value,

$$52! = 80,658,175,170,943,878,571,660,636,856,403,766,975,289,505,440,883,277,824,000,000,000,000$$

is much bigger than can fit into a 64-bit integer representation.

Fortunately, $C(52, 5)$ can also be defined recursively by splitting the problem into two smaller subproblems. Consider card hands containing a specific card, say the ace of clubs, and those that do not. For hands that do contain the ace of clubs, we need to choose 4 more cards from the remaining 51 cards, i.e., $C(51, 4)$. For hands that do not contain the ace of clubs, we need to choose 5 cards from the remaining 51 cards, i.e., $C(51, 5)$. Therefore, $C(52, 5) = C(51, 4) + C(51, 5)$. This point of view leads to the following recurrence for computing the binomial coefficient $C(n, k)$:

$$\begin{aligned} C(n, 0) &= 1 \\ C(n, k) &= 1, \text{ if } k = n \\ C(n, k) &= C(n - 1, k) + C(n - 1, k - 1), \text{ if } 0 < k < n \end{aligned}$$

Given this recurrence, we can write the following recursive function to compute $C(n, k)$.

```
int C(int n, int k) { // Assume 0 <= k <= n           1
    if (k == 0 || k == n)                             2
        return 1;                                     3
    else return C(n-1, k) + C(n-1, k-1);               4
}                                                       5
```

However, this divide-and-conquer approach to computing binomial coefficients is very slow due to redundant calculations performed by the recursive calls, similar to the slow, divide-and-conquer approach to computing Fibonacci numbers.

This programming assignment has two parts. First, write a C/C++ (or Python) program that uses dynamic programming, instead of divide-and-conquer, to compute $C(n, k)$. Second, write a

C/C++ (or Python) program that uses memoization, which is a special form of dynamic programming, to compute $C(n, k)$.

The basic dynamic programming approach is “bottom-up.” It solves subproblems in order from smallest to largest until the original problem is solved. Memoization is a dynamic-programming technique that uses the “top-down” recursive definition to solve only the needed smaller problems, and avoids redundant calculations by solving smaller-problems once and storing their solutions.

The code below uses memoization to compute Fibonacci numbers. Before solving a smaller problem, it checks to see if its solution is already stored.

```
int fib_memo( int n, int fibonacci[] ) { 1
    if ( fibonacci[n] == -1) { 2
        fibonacci[n] = fib_memo(n-1, fibonacci) 3
                        + fib_memo(n-2, fibonacci); 4
    } 5
    return fibonacci[n]; 6
} 7

int fib( int n ) { 8
    // array to store solutions 9
    int fibonacci[n+1]; 10
    // fill base cases 11
    fibonacci[0] = 0; 12
    fibonacci[1] = 1; 13
    // Use -1 for unknown solutions 14
    for (i=2; i <= n; i++) { 15
        fibonacci[i] = -1; 16
    } 17
    return fib_memo(n, fibonacci); 18
} 19
} 20
```

Using this code as a guide, implement an algorithm that uses memoization to compute binomial coefficients.

What to turn in

Turn in a copy of the program you wrote for computing binomial coefficients using both bottom-up dynamic programming and memoization. Your code can be written in C/C++ or Python. In addition, turn in a one-page report that includes an experimental comparison of the running times of the divide-and-conquer, bottom-up dynamic programming, and memoization algorithms. Which is fastest, and why? Which is slowest, and why? Does the memoization algorithm solve fewer subproblems than the bottom-up dynamic programming algorithm? Briefly explain.