Lecture 15 Reduction to Hessenberg Form

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MATH 562 Numerical Analysis II

Outline

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Eigenvalue Revealing Factorization

- Eigenvalue-revealing factorization of square matrix A
 - Diagonalization $\mathbf{A} = \mathbf{X} \boldsymbol{\Lambda} \mathbf{X}^{-1}$ (nondefective \mathbf{A})
 - Unitary Diagonalization $\mathbf{A} = \mathbf{Q} \mathbf{\Lambda} \mathbf{Q}^*$ (normal \mathbf{A})
 - Unitary triangularization (Schur factorization) $\mathbf{A} = \mathbf{Q}\mathbf{T}\mathbf{Q}^*$ (any \mathbf{A})
 - Jordan normal form $\mathbf{A} = \mathbf{X}\mathbf{J}\mathbf{X}^{-1}$ where \mathbf{J} block diagonal with

- In general, Schur factorization is used, because
 - Unitary matrices are involved, so algorithm tends to be more stable
 - If **A** is normal, then Schur form is diagonal

"Obvious" Algorithms

- Most obvious method is to find roots of characteristic polynomial $p_{\mathbf{A}}(\lambda)$, but it is very ill-conditioned.
- Another idea is power iteration, using fact that

$$\frac{\mathbf{x}}{\|\mathbf{x}\|}, \frac{\mathbf{A}\mathbf{x}}{\|\mathbf{A}\mathbf{x}\|}, \frac{\mathbf{A}^2\mathbf{x}}{\|\mathbf{A}^2\mathbf{x}\|}, \frac{\mathbf{A}^3\mathbf{x}}{\|\mathbf{A}^3\mathbf{x}\|}, \dots$$

converge to an eigenvector corresponding to the largest eigenvalue of **A** in absolute value, but it may converge very slowly

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Instead, compute a eigenvalue-revealing factorization, such as Schur factorization

$$A = QTQ^*$$

by introducing zeros, using algorithms similar to QR factorization

A Fundamental Difficulty

 However, eigenvalue-revealing factorization cannot be done in finite number of steps:

Any eigenvalue solver must be iterative

ullet To see this, consider a general polynomial of degree m

$$p(z) = z^m + a_{m-1}z^{m-1} + \dots + a_1z + a_0$$

There is no closed-form expression for the roots of p: (Abel, 1842) In general, the roots of polynomial equations higher than fourth degree cannot be written in terms of a finite number of operations

A Fundamental Difficulty Cont'd

ullet However, the roots of $p_{\mathbf{A}}$ are the eigenvalues of the companion matrix

$$\mathbf{A} = \left[\begin{array}{cccc} 0 & & & -a_0 \\ 1 & 0 & & -a_1 \\ & 1 & \ddots & \vdots \\ & & \ddots & 0 & -a_{m-2} \\ & & 1 & -a_{m-1} \end{array} \right]$$

- Therefore, in general, we cannot find the eigenvalues of a matrix in a finite number of steps
- In practice, however, there are algorithms that converge to desired precision in a few iterations

Schur Factorization and Diagonalization

• Most eigenvalue algorithms compute Schur factorization $\mathbf{A} = \mathbf{Q}T\mathbf{Q}^*$ by transforming \mathbf{A} with similarity transformations

$$\underbrace{\mathbf{Q}_{j}^{*}\cdots\mathbf{Q}_{2}^{*}\mathbf{Q}_{1}^{*}}_{\mathbf{Q}^{*}}\mathbf{A}\underbrace{\mathbf{Q}_{1}\mathbf{Q}_{2}\cdots\mathbf{Q}_{j}}_{\mathbf{Q}},$$

where \mathbf{Q}_i are unitary matrices, which converge to \mathbf{T} as $j \to \infty$.

- Note: Real matrices might need complex Schur forms and eigenvalues
- Question: For hermitian A, what matrix will the sequence converge to?

Two Phases of Eigenvalue Computations

 General A: First convert to upper-Hessenberg form, then to upper triangular

• Hermitian A: First convert to tridiagonal form, then to diagonal

• In general, phase 1 is direct and requires $O(m^3)$ flops, and phase 2 is iterative and requires O(m) iterations, and $O(m^3)$ flops for non-Hermitian matrices and $O(m^2)$ flops for Hermitian matrices

Introducing Zeros by Similarity Transformations

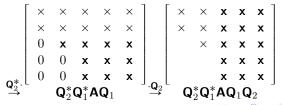
 First attempt: Compute Schur factorization A = QTQ* by applying Householder reflectors from both left and right

- Unfortunately, the right multiplication destroys the zeros introduced by Q₁*.
- This would not work because of Abel?s theorem
- However, the subdiagonal entries typically decrease in magnitude (Later with QR)

The Hessenberg Form

 Second attempt: try to compute upper Hessenberg matrix H similar to A:

- The zeros introduced by Q₁*A were not destroyed this time!
- Continue with remaining columns would result in Hessenberg form:



The Hessenberg Form

• After m-2 steps, we obtain the Hessenberg form:

For hermitian matrix A, H is hermitian and hence is tridiagonal

$$\mathbf{Q}_{m-2}^*\cdots\mathbf{Q}_2^*\mathbf{Q}_1^*\mathbf{A}\mathbf{Q}_1\mathbf{Q}_2\cdots\mathbf{Q}_{m-2}=\mathbf{H}=\left[egin{array}{cccc} imes & imes &$$

Householder Reduction to Hessenberg

Householder Reduction to Hessenberg Form

$$\begin{split} &\text{for } k=1 \text{ to } m-2 \\ &\mathbf{x} = \mathbf{A}_{k+1:m,k} \\ &\mathbf{v}_k = sign(x_1) \|\mathbf{x}\|_2 \mathbf{e}_1 + \mathbf{x} \\ &\mathbf{v}_k = \mathbf{v}_k / \|\mathbf{v}_k\|_2 \\ &\mathbf{A}_{k+1:m,k:m} = \mathbf{A}_{k+1:m,k:m} - 2\mathbf{v}_k (\mathbf{v}_k^* \mathbf{A}_{k+1:m,k:m}) \\ &\mathbf{A}_{1:m,k+1:m} = \mathbf{A}_{1:m,k+1:m} - 2(\mathbf{A}_{1:m,k+1:m} \mathbf{v}_k) \mathbf{v}_k^* \end{split}$$

- Note: Q is never formed explicitly.
- Operation count

$$\sim \sum_{k=1}^{m-2} 4(m-k)^2 + 4m(m-k) \sim 4m^3/3 + 4m^3 - 4m^3/2 = 10m^3/2$$

Reduction to Tridiagonal Form

• If A is hermitian, then

- For Hermitian ${\bf A}$, operation count would be same as Householder QR: $4m^3/3$
 - First, taking advantage of sparsity, cost of applying right reflectors is also $4(m-k)^2$ instead of 4m(m-k), so cost is

$$\sim \sum_{k=1}^{m-2} 8(m-k)^2 \sim 8m^3/3$$

• Second, taking advantage of symmetry, cost is reduced by 50% to $4m^3/3$

Stability of Hessenberg Reduction

Theorem

Householder reduction to Hessenberg form is backward stable, in that

$$\tilde{\mathbf{Q}}\tilde{\mathbf{H}}\tilde{\mathbf{Q}}^* = \mathbf{A} + \delta \mathbf{A}, \ \frac{\|\delta \mathbf{A}\|}{\|\mathbf{A}\|} = O(\epsilon_{machine})$$

for some $\delta \mathbf{A} \in \mathbb{C}^{m \times m}$

Note: Similar to Householder QR, $\tilde{\mathbf{Q}}$ is exactly unitary based on some $\tilde{\mathbf{v}}_k$.