Caleb Logemann MATH 566 Discrete Optimization Homework 9

1. Implement Directed Minimum Mean Cycle in Sage. Before implementing the algorithm, show that it is possible to slightly modify the algorithm. Instead of adding an extra vertex s and edges from s to all other vertices, it is possible to simply assign $F_0(v) = 0$ for all $v \in V$ at the beginning. This avoids the hassle with adding an extra vertex. But it requires an argument that the algorithm is still correct.

```
def minimumMeanCostCycle(graph):
    n = graph.order()
    p = \{v : [None] * (n+1) \text{ for } v \text{ in } graph.vertices()\}
    F = \{v: [0] + [oo] * n \text{ for } v \text{ in } graph.vertices()\}
     for k in range (1, n+1):
         for e in graph.edges():
              u = e[0]
              v = e[1]
              c = e[2]
               \mbox{if} \ F \, [\, v \, ] \, [\, k \, ] \ > \ F \, [\, u \, ] \, [\, k - 1] \ + \ c : 
                   F[v][k] = F[u][k-1] + c
                   p[v][k] = u
    # If all are infinite then min will be infinite
     if \min([F[v][n] \text{ for } v \text{ in } graph.vertices()]) == oo:
         print 'This graph is acyclic'
         return (None, None)
    m = \{v: oo for v in graph. vertices()\}
     for v in graph.vertices():
         if F[v][n] != oo:
              m[v] = max([(F[v][n] - F[v][k])/(n-k)) for k in range(n-1)])
    x = \min(m, \text{ key=m. get})
    mu = m[x]
    # create cycle
     cycle = DiGraph()
     cycle.add_vertex(x)
    k = n
     v = x
     while p[v][k] != x:
         u = v
         v = p[v][k]
         cycle.add vertex(v)
         cycle.add_edge(v, u, graph.edge_label(v, u))
         k=1
     cycle.add_edge(x, v, graph.edge_label(x, v))
     return (cycle, mu)
```

```
load ( 'minimumMeanCostCycle . sage ')
g = DiGraph([(1,2,1), (2,3,2), (3,4,-1), (4,1,1), (4,2,2)])
g.show(edge_labels=True)
(cycle, mu) = minimumMeanCostCycle(g)
if cycle != None:
    print "Mu≡",mu
    cycle.show(edge labels=True)
g = DiGraph([(1,2,3), (1,3,2), (2,4,-1), (3,4,1), (4,5,0), (5,6,1),
   \hookrightarrow (5,7,2), (6,8,1), (7,8,3), (8,1,2)
g.show(edge_labels=True)
(cycle, mu) = minimumMeanCostCycle(g)
if cycle != None:
    print "Mu≡", mu
    cycle.show(edge labels=True)
g = DiGraph([(1,2,1), (2,3,1), (3,4,1), (1,4,1), (4,5,2)])
g.show(edge labels=True)
(cycle, mu) = minimumMeanCostCycle(g)
if cycle != None:
    print "Mu≡",mu
    cycle.show(edge labels=True)
```

2. Show that in integer program, it is possible to express the following constraint:

$$x \in [100, 200] \cup [300, 400]$$

in other words

$$100 \leq x \leq 200$$
 or $300 \leq x \leq 400$

How to express the constraint without using or?

3. Determine which of the matrices below are (i) unimodular, (ii) totally unimodular, or (iii) neither. Be sure to explain your answer.

$$\begin{pmatrix} 1 & -1 & -1 & 0 \\ -1 & 0 & 0 & 1 \\ 0 & 1 & 0 & -1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \qquad \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \end{pmatrix} \qquad \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 \end{pmatrix}$$
a.
b.
c.

(i) Since these are all square matrices, they will be unimodular if their determinant is ± 1 . So I compute the determinants of these 3 matrices. First I will compute the determinant of (a).

$$\det(a) = \begin{vmatrix} 1 & -1 & -1 & 0 \\ -1 & 0 & 0 & 1 \\ 0 & 1 & 0 & -1 \\ 0 & 0 & 1 & 0 \end{vmatrix}$$

Expand along first column

$$\det(a) = 1 \times \begin{vmatrix} 0 & 0 & 1 \\ 1 & 0 & -1 \\ 0 & 1 & 0 \end{vmatrix} - (-1) \times \begin{vmatrix} -1 & -1 & 0 \\ 1 & 0 & -1 \\ 0 & 1 & 0 \end{vmatrix}$$

Expand along the first column for the first determinant and along the last column for the second determinant

$$= 1 \times \left(-1 \times \begin{vmatrix} 0 & 1 \\ 1 & 0 \end{vmatrix}\right) - (-1) \times \left(1 \times \begin{vmatrix} -1 & -1 \\ 0 & 1 \end{vmatrix}\right)$$

$$= 1 \times (-1 \times -1) - (-1) \times (1 \times -1)$$

$$= 1 - 1$$

$$= 0$$

Since the determinant of (a) is 0 this matrix is not unimodular. Second I will compute the determinant of (b).

$$\det(b) = \begin{vmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \end{vmatrix}$$

I will first expand along the first row.

$$\det(b) = 1 \times \begin{vmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{vmatrix} + 1 \times \begin{vmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{vmatrix}$$

I will expand the first determinant along the first row and the second determinant along the first column

$$det(b) = 1 \times \left(1 \times \begin{vmatrix} 1 & 1 \\ 0 & 1 \end{vmatrix}\right) + 1 \times \left(1 \times \begin{vmatrix} 1 & 0 \\ 0 & 1 \end{vmatrix}\right)$$
$$= 1 \times (1 \times 1) + 1 \times (1 \times 1)$$
$$= 1 + 1$$
$$= 2$$

Since the determinant of (b) is 2 this matrix is not unimodular. Lastly I will compute the determinant of (c).

$$\det(c) = \begin{vmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{vmatrix}$$

First I will expand along the first column

$$\det(c) = -1 \times \begin{vmatrix} 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{vmatrix}$$

Next I will expand along the first row

$$\det(c) = -1 \times \left(1 \times \begin{vmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{vmatrix} \right)$$

Expanding along the first column gives

$$\det(c) = -1 \times \left(1 \times \left(-1 \times \begin{vmatrix} 1 & 1 \\ 1 & 0 \end{vmatrix} + 1 \times \begin{vmatrix} 1 & 1 \\ 0 & 1 \end{vmatrix}\right)\right)$$

$$= -1 \times (1 \times (-1 \times -1 + 1 \times 1))$$

$$= -1 \times (1 \times (1 + 1))$$

$$= -1 \times (1 \times 2)$$

$$= -1 \times 2$$

$$= -2$$

Since the determinant of (c) is -2 this matrix is not unimodular.

In summary none of the matrices are unimodular.

- (ii) In order for a matrix to be totally unimodular every square submatrix must have determinant -1, 0, or 1. Note that since matrices (b) and (c) don't have a determinant -1, 0, or 1 when considered as a whole matrix they cannot be totally unimodular. Matrix (a) which has determinant 0 can potentially be totally unimodular. In fact we see that each column has exactly one 1 and one -1, so by a theorem in the notes (a) is totally unimodular.
- (iii) We have shown that (a) is totally unimodular but not unimodular. However (b) and (c) are neither unimodular nor totally unimodular.
- 4. Show that $A \in \mathbb{Z}^{m \times n}$ is totally unimodular iff $[A\ I]$ is unimodular (where I is $m \times m$ unit matrix).
- 5. Find a unimodular matrix A, that is not totally unimodular.

Consider the matrix

$$A = \frac{9}{5} \quad \frac{7}{4}$$

The matrix A is unimodular because $A \in \mathbb{Z}^{2\times 2}$ and $\det(A) = 9 \times 4 - 5 \times 7 = 36 - 35 = 1$. However A is not totally unimodular because not all of the entries of A are -1, 0, 1.