A. Designing a topdown parser: Consider the following grammar.

prog : stmt-list
stmt-list : stmt-list stmt | stmt
stmt : PRINT expr | PRINT string
string : BEGINQUOTE charlist ENDQUOTE
charlist : charlist LETTER | ε
expr : expr + term | expr - term | term
term : term * factor | term / factor | factor
factor : (expr) | NUM

The nonterminal symbols are:

prog, stmt-list, stmt, string, charlist, expr, term, factor

The terminal symbols are:

PRINT, BEGINQUOTE, ENDQUOTE, LETTER, +, -, *, /, (,), NUM

(A1) The grammar is left-recursive. Transform it into a non-left recursive grammar using the following ranking of the nonterminals.

prog, stmt-list, stmt, string, charlist, expr, term, factor

prog : stmt-list
stmt-list : stmt stmt-list′
stmt-list′ : stmt stmt-list′| ε
stmt : PRINT expr | PRINT string
string : BEGINQUOTE charlist ENDQUOTE
charlist : charlist′
charlist′ : LETTER charlist′| ε
expr : term expr′
expr′ : + term expr′ | - term expr′ | ε
term : factor term′
term′ : * factor term′ | / factor term′ | ε
factor : (expr) | NUM

(A2) Further modify the grammar so it is left factored.

prog : stmt-list
stmt-list : stmt stmt-list′
stmt-list′ : stmt stmt-list′| ε
stmt : PRINT word
word : expr | string
string : BEGINQUOTE charlist ENDQUOTE
charlist : charlist′
charlist′ : LETTER charlist′| ε
expr : term expr′
expr′ : + term expr′ | - term expr′ | ε
term : factor term′
term′ : * factor term′ | / factor term′ | ε
factor : (expr) | NUM

(A3) Compute the FIRST sets of each nonterminal symbols.

prog : { PRINT }
stmt-list : { PRINT }
stmt-list′ : { PRINT }
stmt : { PRINT }
word : { (, NUM, BEGINQUOTE }
string : { BEGINQUOTE }
charlist : { LETTER }
charlist′ : { LETTER }
expr : { (, NUM }
expr′ : { +, - }
term : { (, NUM }
term′ : { *, / }
factor : { (, NUM }

(A4) Compute the FOLLOW sets of each nonterminal symbols

prog : {}
stmt-list : {}
stmt-list′ : {}
stmt : { PRINT }
word : { PRINT }
string : { PRINT }
charlist : { ENDQUOTE }
charlist′ : { ENDQUOTE }
expr : { PRINT,) }
expr′ : { PRINT,) }
term : { +, -, PRINT,) }
term′ : { +, -, PRINT,) }
factor : { *, /, +, -, PRINT,) }

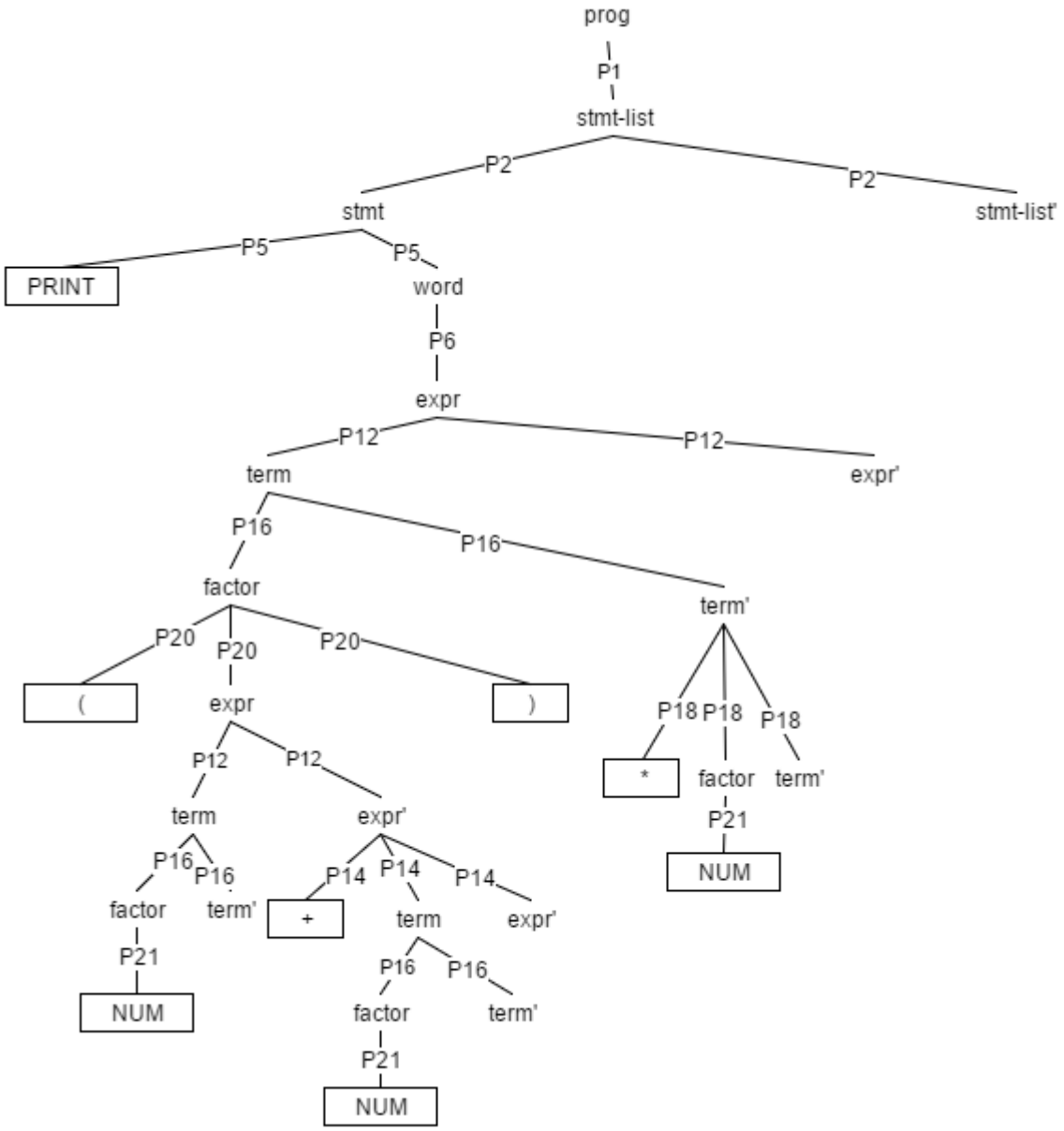
(A5) Compute the predictive parsing table of the grammar.

Non - Terminal	Input Symbol											
	PRINT	BEGINQUOTE	ENDQUOTE	LETTER	+	-	*	/	(	)	NUM	$
prog	prog : stmt-list											
stmt-list	stmt-list : stmt stmt-list′											
stmt-list′	stmt-list′ : stmt stmt-list′											stmt-list′ : ε
stmt	stmt : PRINT word											
word		word : string							word : expr		word : expr	
string		string : BEGINQUOTE charlist ENDQUOTE										
charlist			charlist : charlist′	charlist : charlist′								
charlist′			charlist′ : ε	charlist′ : LETTER charlist′								
expr									expr : term expr′		expr : term expr′	
expr′	expr′ : ε				expr′ : + term expr′	expr′ : - term expr′				expr′ : ε		expr′ : ε
term									term : factor term′		term : factor term′	
term′	term′ : ε				term′ : ε	term′ : ε	term′ : * factor term′	term′ : / factor term′		term′ : ε		term′ : ε
factor									factor : (expr)		factor : NUM	

(A6) Consider the following program.

PRINT (NUM + NUM) * NUM

Use the predictive parsing table to construct the parse tree. At each step, show the production used to expand the nodes.



P1. **prog : stmt-list**

P2. **stmt-list : stmt stmt-list'**
P3. **stmt-list' : stmt stmt-list'**
P4. **stmt-list' : ε**
P5. **stmt : PRINT word**
P6. **word : expr**
P7. **word : string**
P8. **string : BEGINQUOTE charlist ENDQUOTE**
P9. **charlist : charlist'**
P10. **charlist' : ε**
P11. **charlist' : LETTER charlist'**
P12. **expr : term expr'**
P13. **expr' : ε**
P14. **expr' : + term expr'**
P15. **expr' : - term expr'**
P16. **term : factor term'**
P17. **term' : ε**
P18. **term' : * factor term'**
P19. **term' : / factor term'**
P20. **factor : (expr)**
P21. **factor : NUM**