## Gauge Algorithm

## Remedy for the "subset sum problem" complexity

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Imagine you have a set with four numbers  $\{-3, 2, 4, 7\}$  and wish to find subsets that have assigned sum equal to 4.

The most brutal method is to list all subsets and make the partial sums.

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Subsets of cardinality 1 (there are 4) {-3}, {2}, {4}, {7}

Subsets of cardinality 2 (there are 6) {-3, 2}, {-3, 4}, {-3, 7}, {2, 4}, {2, 7}, {4, 7}

Subsets of cardinality 3 (there are 4) {-3, 2, 4}, {-3, 2, 7}, {-3, 4, 7}, {2, 4, 7}

Subset (no proper) of cardinality 4 (there is only one)
```

More generally, the number of subsets of cardinality k in a set with n elements is

$$\frac{n!}{k!(n-k)!}$$

*n* increases, the problem becomes very complex even for powerful computers.

To give an idea, a set with 1000 elements, possesses

 $\{-3, 2, 4, 7\}$ 

```
1.368.173.298.991.500 subsets with 6 elements,
194.280.608.456.793.000 subsets with 7 elements,
24.115.080.524.699.431.125 subsets with 8 elements, ...
```

Computationally, the problem is classified as NP-complete.

The component of **Cambusa**<sup>®</sup>, which remedies statistically to the problem of complexity, uses a procedure based on *Markov processes* that I called **"Gauge algorithm"**. I'll describe in broad terms the basic idea.

From a set with n elements, we consider the non-empty  $2^n-1$  subsets be taken as process states. A state transition is accomplished by adding or removing an item. For each transition is assigned a probability according to a law that more the transition leads to a partial sum closer to the reference value, the greater the probability that the transition happen.

Starting from a pivot state, kick the process off: the **conjecture** is that you generate a kind of "statistical attraction" towards the sought solutions.



