

# Graph Metrics Definitions

Let  $G = (V, E)$  be an undirected graph representing a street network, where  $V$  is the set of nodes and  $E$  the set of edges. Each edge  $e \in E$  has a positive length  $\ell(e) > 0$ . Let  $n = |V|$  and  $m = |E|$ .

For shortest-path based metrics, let  $d(u, v)$  denote the weighted shortest-path distance between nodes  $u$  and  $v$  using edge lengths  $\ell(\cdot)$ . For metrics that require connectivity, computations are typically performed on the largest connected component of  $G$ .

## 0.1 Size and scale

- **Number of nodes (n\_nodes):**  $n = |V|$ .
- **Number of edges (n\_edges):**  $m = |E|$ .
- **Total length (total\_length):**

$$L_{\text{total}} = \sum_{e \in E} \ell(e).$$

## 0.2 Degree-based connectivity

Let  $\deg(v)$  be the degree of node  $v \in V$ .

- **Average degree (avg\_degree):**

$$\bar{k} = \frac{1}{n} \sum_{v \in V} \deg(v).$$

- **Number of degree-1 nodes (num\_deg1):**

$$n_{\deg=1} = |\{v \in V : \deg(v) = 1\}|.$$

- **Proportion of degree-1 nodes (prop\_deg1):**

$$p_{\deg=1} = \frac{n_{\deg=1}}{n}.$$

- **Number of nodes with degree  $\geq 3$  (num\_deg\_ge3):**

$$n_{\deg \geq 3} = |\{v \in V : \deg(v) \geq 3\}|.$$

- **Proportion of nodes with degree  $\geq 3$  (prop\_deg\_ge3):**

$$p_{\deg \geq 3} = \frac{n_{\deg \geq 3}}{n}.$$

- **Degree entropy (degree\_entropy):** Let  $p(k)$  be the empirical probability that a node has degree  $k$ .

$$H_{\text{deg}} = - \sum_k p(k) \log p(k).$$

### 0.3 Density and mesh measures

- **Graph density (density):**

$$\rho = \frac{2m}{n(n-1)}.$$

This measures edge density relative to a complete graph.

- **Mesh density (mesh\_density):**

$$\rho_{\text{mesh}} = \frac{m-n+1}{2n-5} \quad (n \geq 3).$$

This measures cyclic redundancy normalized by the planar upper bound; values near 0 indicate tree-like structure and larger values indicate more meshed networks.

### 0.4 Shortest-path extension metrics

Define the eccentricity of node  $u$  as  $\epsilon(u) = \max_{v \in V} d(u, v)$  (within the component considered).

- **Average shortest path length (avg\_shortest\_path\_length):**

$$\bar{d} = \frac{1}{n(n-1)} \sum_{u \neq v} d(u, v).$$

- **Radius (radius):**

$$r = \min_{u \in V} \epsilon(u).$$

- **Diameter (diameter):**

$$D = \max_{u \in V} \epsilon(u).$$

- **Average eccentricity (avg\_eccentricity):**

$$\bar{\epsilon} = \frac{1}{n} \sum_{u \in V} \epsilon(u).$$

### 0.5 Centrality metrics

- **Closeness centrality (CC):**

$$CC(u) = \frac{1}{\sum_{v \neq u} d(u, v)}.$$

**Average node closeness (avg\_node\_closeness)** is  $\frac{1}{n} \sum_{u \in V} CC(u)$ .

- **Betweenness centrality (BC):** Let  $\sigma_{st}$  be the number of shortest paths between  $s$  and  $t$ , and  $\sigma_{st}(u)$  the number of those paths that pass through  $u$ .

$$BC(u) = \sum_{s \neq u \neq t} \frac{\sigma_{st}(u)}{\sigma_{st}}.$$

(Weighted shortest paths use  $\ell(\cdot)$ .)

- **Maximum node betweenness (max\_node\_betweenness):**

$$\max_{u \in V} BC(u).$$

- **Average node betweenness (avg\_node\_betweenness):**

$$\overline{BC} = \frac{1}{n} \sum_{u \in V} BC(u).$$

- **Coefficient of variation of node betweenness (cv\_node\_betweenness):**

$$CV_{BC} = \frac{\text{std}(\{BC(u)\})}{\text{mean}(\{BC(u)\})},$$

when the mean is nonzero.

## 0.6 Branching, redundancy, and spanning tree

- **Branching index (branching\_index):**

$$B = \frac{n_{\deg=1}}{n_{\deg \geq 3}},$$

with a large sentinel value used when  $n_{\deg \geq 3} = 0$ .

- **Cyclomatic number (cyclomatic\_number):**

$$\mu = m - n + c,$$

where  $c$  is the number of connected components. It equals the number of “extra” edges beyond a spanning forest, i.e., the number of independent cycles.

- **Minimum spanning tree length (mst\_length):** Let  $T$  be a minimum spanning tree of the (connected) component considered, using  $\ell(\cdot)$  as weights:

$$L_{\text{MST}} = \sum_{e \in T} \ell(e).$$

- **MST ratio (mst\_ratio):**

$$R_{\text{MST}} = \frac{L_{\text{total}}}{L_{\text{MST}}},$$

when  $L_{\text{MST}} > 0$ .

## 0.7 Edge length statistics

Let  $\mathcal{L} = \{\ell(e) : e \in E\}$  denote the multiset of edge lengths.

- **Average edge length (avg\_edge\_length):**  $\text{mean}(\mathcal{L})$ .
- **Median edge length (median\_edge\_length):**  $\text{median}(\mathcal{L})$ .
- **Minimum edge length (min\_edge\_length):**  $\text{min}(\mathcal{L})$ .
- **Maximum edge length (max\_edge\_length):**  $\text{max}(\mathcal{L})$ .
- **Standard deviation of edge length (std\_edge\_length):**  $\text{std}(\mathcal{L})$ .
- **Coefficient of variation of edge length (cv\_edge\_length):**

$$CV_\ell = \frac{\text{std}(\mathcal{L})}{\text{mean}(\mathcal{L})},$$

when the mean is nonzero.

- **25th percentile of edge length (edge\_length\_p25):**  $\text{quantile}_{0.25}(\mathcal{L})$ .
- **75th percentile of edge length (edge\_length\_p75):**  $\text{quantile}_{0.75}(\mathcal{L})$ .

## 0.8 Community structure

- **Modularity (modularity):**  $Q$ , the modularity value of a community partition (e.g., from greedy modularity maximization) comparing within-community edge density against a null model.
- **Number of communities (n\_communities):** the number of detected communities in the chosen partition.

## 0.9 Straightness and tortuosity

Let  $d_G(u, v) = d(u, v)$  be the network shortest-path distance. If node coordinates are available, let  $d_E(u, v)$  be the Euclidean distance between  $u$  and  $v$ .

- **Average straightness (avg\_straightness\_length):**

$$\bar{S} = \frac{1}{|P|} \sum_{(u,v) \in P} \frac{d_E(u, v)}{d_G(u, v)},$$

measuring how direct routes are compared to straight-line distances.

- **Average tortuosity (avg\_tortuosity\_length):**

$$\bar{T} = \frac{1}{|P|} \sum_{(u,v) \in P} \frac{d_G(u, v)}{d_E(u, v)},$$

measuring detour relative to straight-line distance.