

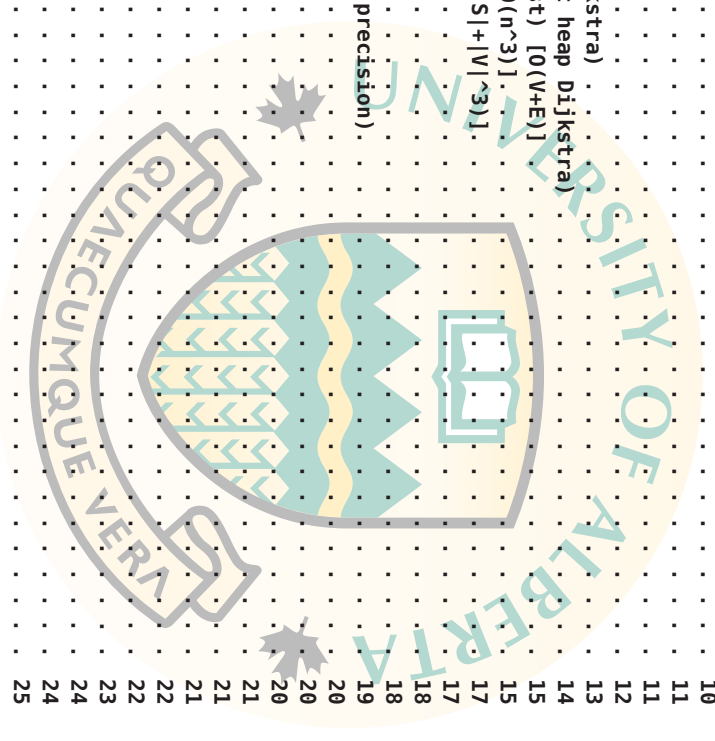
University of Alberta

2008 ACM ICPC World Finals

Code Archive

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```
/* Geometry: Complex Arithmetic -----*/
```

```
// These two values are used in most of the geometry algorithms.
```

```
double PI = 2*acos(0.0);
```

```
double EPS = 1E-8;
```

```
struct pol {
    double r, t;
    pol(double R = 0, double T = 0) : r(R), t(T) {}
};
```

```
struct point {
    double x, y;
    point(double X = 0, double Y = 0) : x(X), y(Y) {}
    point(const pol &P) : x(P.r*cos(P.t)), y(P.r*sin(P.t)) {}
    point conj() const { return point(x, -y); }
    double mag2() const { return x*x + y*y; }
    double mag() const { return sqrt(mag2()); }
    double arg() const { return atan2(y, x); }
    point operator-() const { return point(-x, -y); }
    point& operator+=(const point &a) { x += a.x; y += a.y; return *this; }
    point& operator-=(const point &s) { x -= s.x; y -= s.y; return *this; }
    point& operator*=(const point &m) {
        double tx = x*m.x - y*m.y, ty = x*m.y + y*m.x;
        x = tx; y = ty; return *this;
    }
    point& operator/=(const point &d) {
        double tx = y*d.y + x*d.x, ty = y*d.x - x*d.y, t = d.mag2();
        x = tx/t; y = ty/t; return *this;
    }
    bool operator<(const point &q) const {
        if (fabs(y-q.y) < EPS) return x < q.x;
        return y < q.y;
    }
    bool operator==(const point &q) const {
        return (fabs(x-q.x) < EPS) && (fabs(y-q.y) < EPS);
    }
    bool operator!=(const point &q) const { return !operator==(q); }
};
```

```
point operator+(point a, const point &b) { return a += b; }
point operator-(point a, const point &b) { return a -= b; }
point operator*(point a, const point &b) { return a *= b; }
point operator/(point a, const point &b) { return a /= b; }
```

```
/* Geometry: Area of a polygon (positive <-> CCW orientation) -----*/
```

```
double areaPoly(vector<point> &p) {
    double sum = 0; int n = p.size();
    for (int i = n-1, j = 0; j < n; i = j++)
        sum += (p[i].conj()*p[j]).y;
    return sum/2;
}
```

```
/* Geometry: Heron's formula for triangle area -----*/
```

```
// Given side lengths a, b, c, returns area or -1 if triangle is impossible
```

```
double area_heron(double a, double b, double c) {
    if (a < b) swap(a, b);
    if (a < c) swap(a, c);
    if (b < c) swap(b, c);
    if ((c-(a-b)) < 0) return -1;
    return sqrt((a+(b+c))*(c-(a-b))*(c+(a-b))*(a+(b-c)))/4.0;
}
```

```
/* Geometry: Closest point on line segment a-b to point c -----*/
```

```
point closest_pt_lineseg(point a, point b, point c) {
    b -= a; c -= a; if (b == 0) return a;
    double d = (c/b).x;
    if (d < 0) d = 0; if (d > 1) d = 1;
    return a + d*b;
}
```

```
/* Geometry: Rectangle in rectangle test -----*/
```

```
// Checks if rectangle of sides x,y fits inside one of sides X,Y
// Code as written rejects rectangles that just touch.
```

```
bool rect_in_rect(double X, double Y, double x, double y) {
    if (Y > X) swap(Y, X);
    if (y > x) swap(y, x);
    double diagonal = sqrt(X*X + Y*Y);
    if (x < X && y < Y)
        return true;
    else if (y >= Y || x >= diagonal)
        return false;
    else {
```

```

double w, theta, tMin = PI/4, tMax = PI/2;
while (tMax - tMin > EPS) {
    theta = (tMax + tMin)/2.0;
    w = (Y-x*cos(theta))/sin(theta);
    if (w < 0 || x * sin(theta) + w * cos(theta) < X)
        tMin = theta;
    else tMax = theta;
}
return (w > y);
}
}

```

/* Geometry: Centroid of a simple polygon [O(N)] -----*/

// Points must be oriented (either CW or CCW), and non-convex is OK

```

point centroid(point p[], int n) {
    double sum = 0; point c;
    for(int i = n-1, j = 0; j < n; i = j++) {
        double area = (p[i].conj()*p[j]).y;
        sum += area; c += (p[i]+p[j])*area;
    }
    sum *= 3.0; c /= sum;
    return c;
}

```

/* Geometry: Convex Hull -----*/

```

struct polar_cmp {
    point P0;
    polar_cmp(point p = 0) : P0(p) {}
    double turn(const point &p1, const point &p2) const {
        return ((p2-P0)*(p1-P0).conj()).y;
    }
    bool operator()(const point &p1, const point &p2) const {
        double d = turn(p1, p2);
        if (fabs(d) < EPS)
            return (p1-P0).mag2() < (p2-P0).mag2();
        else return d > 0;
    }
};

vector<point> convex_hull(vector<point> p) {
    sort(p.begin(), p.end());
    int n = unique(p.begin(), p.end()) - p.begin();
    sort(p.begin()+1, p.begin()+n, polar_cmp(p[0]));
}

```

```

if (n <= 2) return vector<point>(p.begin(), p.begin()+n);
vector<point> hull(p.begin(), p.begin()+2); int h = 2;
for (int i = 2; i < n; ++i) {
    while ((h > 1) && (polar_cmp(hull[h-2]).turn(hull[h-1], p[i]) < EPS)) {
        hull.pop_back(); --h;
    }
    hull.push_back(p[i]); ++h;
}
return hull;
}

```

/* Geometry: Area of intersection of two circles -----*/

```

struct circle {
    point c; double r;
};
double CIArea(circle &a, circle &b) {
    double d = (b.c-a.c).mag();
    if (d <= (b.r - a.r)) return a.r*a.r*PI;
    if (d <= (a.r - b.r)) return b.r*b.r*PI;
    if (d >= a.r + b.r) return 0;
    double alpha = acos((a.r*a.r+d*d-b.r*b.r)/(2*a.r*d));
    double beta = acos((b.r*b.r+d*d-a.r*a.r)/(2*b.r*d));
    return a.r*a.r*(alpha-0.5*sin(2*alpha))+b.r*b.r*(beta-0.5*sin(2*beta));
}

```

/* Geometry: Points of intersection of two circles -----*/

```

// For identical circles, returns true with "indefinite" coordinates in p,q
// p, q will compare equal if there is only one intersection point

bool circIntersect(circle &a, circle &b, point &p, point &q) {
    double d2 = (b.c-a.c).mag2(), rS = a.r+b.r, rD = a.r-b.r;
    if (d2 > rS*rS) return false;
    if (d2 < rD*rD) return false;
    double ca = 0.5*(1 + rS*rD/d2);
    point z = point(ca, sqrt((a.r*a.r/d2)-ca*ca));
    p = a.c + (b.c-a.c)*z;
    q = a.c + (b.c-a.c)*z.conj();
    return true;
}

```

```
/* Geometry: Line-circle intersection points -----*/
```

```
// Intersects (infinite) line through a,b with circle c, returns pts. p, q
// If a and b are the same, returns true with "indefinite" coordinates in p,q
// p, q will compare equal if there is only one intersection point
```

```
bool lineCircIntersect(point a, point b, circle c, point &p, point &q) {
    c.c -= a; b -= a;
    point m = b*(c.c/b).x;
    double d2 = (m-c.c).mag2();
    if (d2 > c.r*c.r) return false;
    double L = sqrt((c.r*c.r-d2)/b.mag2());
    p = a + m + L*b;
    q = a + m - L*b;
    return true;
}
```

```
/* Geometry: Area of union of rectangles [0(N^2)] -----*/
```

```
// Rectangle sides are parallel to the x & y axes
// May be desirable to add a constructor to 'rect' to ensure that the
// coordinates are properly sorted
```

```
struct rect {
    double minx, miny, maxx, maxy;
};
struct edge {
    double x, miny, maxy;
    char m;
    bool operator<(const edge &e) const {
        return x < e.x;
    }
};
```

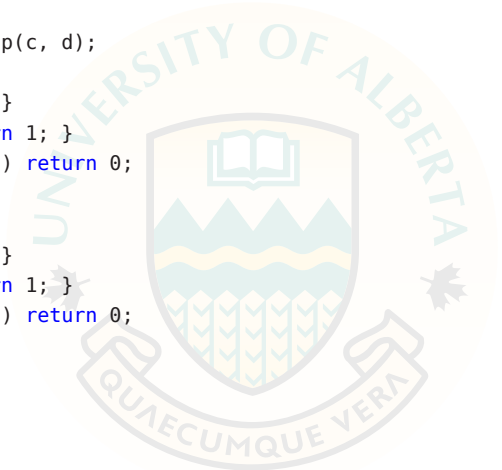
```
double area_unionrect(vector<rect> R){
    int n = R.size();
    vector<double> ys(2*n);
    vector<edge> e(2*n);
    for (int i = 0; i < n; ++i) {
        e[2*i].miny = e[2*i+1].miny = ys[2*i] = r[i].miny;
        e[2*i].maxy = e[2*i+1].maxy = ys[2*i+1] = r[i].maxy;
        e[2*i].x = r[i].minx; e[2*i].m = 1;
        e[2*i+1].x = r[i].maxx; e[2*i+1].m = -1;
    }
    sort(ys.begin(), ys.end());
    sort(e.begin(), e.end());
    double sum = 0, cur = 0;
```

```
for (int i = 0; i < 2*n; ++i) {
    if (i) sum += (ys[i]-ys[i-1])*cur;
    int flag = 0; double sx = cur = 0;
    for (int j = 0; j < 2*n; ++j) {
        if (e[j].miny <= ys[i] && ys[i] < e[j].maxy) {
            if (!flag) sx = e[j].x;
            flag += e[j].m;
            if (!flag) curr += e[j].x-sx;
        }
    }
    return sum;
}
```

```
/* Geometry: Line segment a-b vs. c-d intersection (IP returned in p) -----*/
```

```
// returns 1 if intersect, 0 if not, -1 if coincident
```

```
int intersect_line(point a, point b, point c, point d, point &p) {
    double num1 = ((a-c)*(d-c).conj()).y, num2 = ((a-c)*(b-a).conj()).y;
    double denom = ((d-c)*(b-a).conj()).y;
    if (fabs(denom) > EPS) {
        double r = num1/denom, s = num2/denom;
        if ((0 <= r) && (r <= 1) && (0 <= s) && (s <= 1)) {
            p = a+r*(b-a);
            return 1;
        }
        return 0;
    }
    if (fabs(num1) > EPS) return 0;
    if (b < a) swap(a, b); if (d < c) swap(c, d);
    if (a.y == b.y) {
        if (b.x == c.x) { p = b; return 1; }
        else if (a.x == d.x) { p = a; return 1; }
        else if ((b.x < c.x) || (d.x < a.x)) return 0;
    }
    else {
        if (b.y == c.y) { p = b; return 1; }
        else if (a.y == d.y) { p = a; return 1; }
        else if ((b.y < c.y) || (d.y < a.y)) return 0;
    }
    return -1;
}
```



/* Geometry: Area of intersection of two general polygons [O(N²)] -----*/

```
int ORDER = -1; // CCW ordering, 1 for CW
struct triangle {
    point p[3];
};

double cross(point a, point b, point c, point d) {
    d -= c; b -= a; return (d*b.conj()).y;
}

int leftRight(const point &a, const point &b, const point &p) {
    // -1: p left of a->b, +1: p right of a->b, 0: p on a->b
    double d = cross(a, b, a, p);
    if (d > EPS) return -1;
    if (d < -EPS) return 1;
    return 0;
}

bool isConcave(point &a, point &b, point &c) {
    // tests if b in a->b->c is concave/flat
    return ORDER*leftRight(a, b, c) <= 0;
}

bool isInsideTriangle(point &a, point &b, point &c, point &p) {
    int r1 = leftRight(a,b,p), r2 = leftRight(b,c,p), r3 = leftRight(c,a,p);
    return (ORDER*r1 >= 0) && (ORDER*r2 >= 0) && (ORDER*r3 >= 0);
}

vector<triangle> triangulate(vector<point> &orig) {
    // Accepts a vector of n ordered vertices, returns triangulation.
    // No triangles if n < 3.
    vector<triangle> T;
    if (orig.size() < 3) return T;
    list<point> P(orig.begin(), orig.end());
    list<point>::iterator a, b, c, q;
    for (a = b = P.begin(), c = ++b, ++c; c != P.end(); a = b, b = ++b, ++c)
        if (!isConcave(*a, *b, *c)) {
            q = P.begin(); if (q == a) { ++q; ++q; ++q; }
            while ((q != P.end()) && !isInsideTriangle(*a, *b, *c, *q)) {
                ++q; if (q == a) { ++q; ++q; ++q; }
            }
            if (q == P.end()) {
                triangle t; t.p[0] = *a; t.p[1] = *b; t.p[2] = *c; T.push_back(t);
                P.erase(b); b = a;
                if (b != P.begin()) --b;
            }
        }
    return T;
}

bool isectLineSegs(point &a, point &b, point &c, point &d, point &p) {
    // Finds intersection p of segments a-b and c-d (returns 0 if none/inf)
```

```
double n1 = cross(c, d, c, a), n2 = -cross(a, b, a, c);
double dn = cross(a, b, c, d);
if (fabs(dn) > EPS) {
    double r = n1/dn, s = n2/dn;
    if ((0 <= r) && (r <= 1) && (0 <= s) && (s <= 1)) {
        p = a+r*(b-a);
        return true;
    }
}
return false;
}

struct radialLessThan {
    point P0;
    radialLessThan(point p = 0) : P0(p) {}
    bool operator()(const point &a, const point &b) const {
        return (ORDER == leftRight(P0, a, b));
    }
};

double isectAreaTriangles(triangle &a, triangle &b) {
    vector<point> P;
    point p; triangle T[2] = {a, b};
    for (int r = 1, t = 0; t < 2; r = t++)
        for (int i = 2, j = 0; j < 3; i = j++) {
            if (isInsideTriangle(T[r].p[0], T[r].p[1], T[r].p[2], T[t].p[i]))
                P.push_back(T[t].p[i]);
            for (int u = 2, v = 0; v < 3; u = v++)
                if (isectLineSegs(T[t].p[i], T[t].p[j], T[r].p[u], T[r].p[v], p))
                    P.push_back(p);
        }
    if (P.empty()) return 0;
    sort(P.begin(), P.end());
    vector<point> U; unique_copy(P.begin(), P.end(), back_inserter(U));
    if (U.size() >= 3) {
        sort(++U.begin(), U.end(), radialLessThan(U[0]));
        return areaPoly(U);
    }
    return 0;
}

double isectAreaGpoly(vector<point> &P, vector<point> &Q) {
    vector<triangle> S = triangulate(P), T = triangulate(Q);
    double area = 0;
    for (vector<triangle>::iterator s = S.begin(); s != S.end(); ++s)
        for (vector<triangle>::iterator t = T.begin(); t != T.end(); ++t)
            area += isectAreaTriangles(*s, *t);
    return -ORDER*area;
}
```

```
/* Geometry: Point in polygon -----*/
```

```
bool pt_in_poly(vector<point> &p, const point &a) {
    int n = p.size(); bool inside = false;
    for (int i = 0, j = n-1; i < n; j = i++) {
        if ((a-p[i]).mag()+(a-p[j]).mag()-(p[i]-p[j]).mag()) < EPS)
            return true; // Boundary case (pt on edge), you may want false here
        if (((p[i].y<=a.y) && (a.y<p[j].y)) || ((p[j].y<=a.y) && (a.y<p[i].y)))
            if (a.x-p[i].x < (p[j].x-p[i].x)*(a.y-p[i].y) / (p[j].y-p[i].y))
                inside = !inside;
    }
    return inside;
}
```

```
/* Geometry: Polygon midpoints -> vertices (n odd) -----*/
```

```
vector<point> midpts2vert(vector<point> &midpts) {
    int n = midpts.size(); vector<point> poly(n);
    poly[0] = midpts[0];
    for (int i = 1; i < n-1; i += 2) {
        poly[0].x += midpts[i+1].x - midpts[i].x;
        poly[0].y += midpts[i+1].y - midpts[i].y;
    }
    for (int i = 1; i < n; i++) {
        poly[i].x = 2.0*midpts[i-1].x - poly[i-1].x;
        poly[i].y = 2.0*midpts[i-1].y - poly[i-1].y;
    }
    return poly;
}
```

```
/* Geometry: 3D Primitives -----*/
```

```
struct point3 {
    double x, y, z;
    point3(double X=0, double Y=0, double Z=0) : x(X), y(Y), z(Z) {}
    point3 operator+(point3 p) { return point3(x + p.x, y + p.y, z + p.z); }
    point3 operator*(double k) { return point3(k*x, k*y, k*z); }
    point3 operator-(point3 p) { return *this + (p*-1.0); }
    point3 operator/(double k) { return *this*(1.0/k); }
    double mag2() { return x*x + y*y + z*z; }
    double mag() { return sqrt(mag2()); }
    point3 norm() { return *this/this->mag(); }
};

double dot(point3 a, point3 b) {
    return a.x*b.x + a.y*b.y + a.z*b.z;
}
```

```
}
point3 cross(point3 a, point3 b) {
    return point3(a.y*b.z - b.y*a.z, b.x*a.z - a.x*b.z, a.x*b.y - b.x*a.y);
}

struct line {
    point3 a, b;
    line(point3 A=point3(), point3 B=point3()) : a(A), b(B) {}
    point3 dir() { return (b - a).norm(); }
};

point3 cpoint_iline(line u, point3 p) {
    // Closest point on an infinite line u to a given point p
    point3 ud = u.dir();
    return u.a - ud*dot(u.a - p, ud);
}

double dist_ilines(line u, line v) {
    // Shortest distance between two infinite lines u and v
    return dot(v.a - u.a, cross(u.dir(), v.dir()).norm());
}

point3 cpoint_ilines(line u, line v) {
    // Finds the closest point on infinite line u to infinite line v.
    // Assumes non-parallel lines
    point3 ud = u.dir(); point3 vd = v.dir();
    double uu = dot(ud, ud), vv = dot(vd, vd), uv = dot(ud, vd);
    double t = dot(u.a, ud) - dot(v.a, ud); t *= vv;
    t -= uv*(dot(u.a, vd) - dot(v.a, vd));
    t /= (uv*uv - uu*vv);
    return u.a + ud*t;
}

point3 cpoint_lineseg(line u, point3 p) {
    // Closest point on a line segment u to a given point p
    point3 ud = u.b - u.a; double s = dot(u.a - p, ud)/ud.mag2();
    if (s < -1.0) return u.b;
    if (s > 0.0) return u.a;
    return u.a - ud*s;
}

struct plane {
    point3 n, p;
    plane(point3 ni = point3(), point3 pi = point3()) : n(ni), p(pi) {}
    plane(point3 a, point3 b, point3 c) : n(cross(b-a, c-a).norm()), p(a) {}
    double d() { return -dot(n, p); }
};

point3 cpoint_plane(plane u, point3 p) {
    // Closest point on a plane u to a given point p
    return p - u.n*(dot(u.n, p) + u.d());
}

point3 iline_isect_plane(plane u, line v) {

```

```

// Point of intersection between an infinite line v and a plane u.
// Assumes line not parallel to plane.
point3 vd = v.dir();
return v.a - vd*((dot(u.n, v.a) + u.d())/dot(u.n, vd));
}

line isect_planes(plane u, plane v) {
// Infinite line of intersection between two planes u and v.
// Assumes planes not parallel.
point3 o = u.n*-u.d(), uv = cross(u.n, v.n);
point3 uvu = cross(uv, u.n);
point3 a = o - uvu*((dot(v.n, o) + v.d()/(dot(v.n, uvu)*uvu.mag2()));
return line(a, a + uv);
}

=====

/* Geometry: Great Circle distance (lat[-90,90], long[-180,180]) -----*/

double greatcircle(double lt1, double lo1, double lt2, double lo2, double r) {
double a = PI*(lt1/180.0), b = PI*(lt2/180.0);
double c = PI*((lo2-lo1)/180.0);
return r*acos(sin(a)*sin(b) + cos(a)*cos(b)*cos(c));
}

=====

/* Geometry: Circle described by three points -----*/

bool circle(point p1, point p2, point p3, point &center, double &r) {
double G = 2*((p2-p1).conj()*(p3-p2)).y;
if (fabs(G) < EPS) return false;
center = p1*(p3.mag2()-p2.mag2());
center += p2*(p1.mag2()-p3.mag2());
center += p3*(p2.mag2()-p1.mag2());
center /= point(0, G); r = (p1-center).mag();
return true;
}

=====

/* Arithmetic: Discrete Logarithm solver [0(sqrt(P))] -----*/

// Given prime P, B, and N, finds least L such that B^L == N (mod P)

typedef unsigned int UI;
typedef unsigned long long ULL;
map<UI,UI> M;
UI times(UI a, UI b, UI m) {
return (ULL) a * b % m;
}

```

```

UI power(UI val, UI power, UI m) {
UI res = 1;
for (UI p = power; p; p >= 1) {
if (p & 1)
res = times(res, val, m);
val = times(val, val, m);
}
return res;
}

UI discrete_log(UI p, UI b, UI n) {
UI jump = sqrt(double(p)); M.clear();
for (UI i = 0; i < jump && i < p-1; ++i)
M[power(b,i,p)] = i+1;
for (UI i = 0, j; i < p-1; i += jump)
if (j = M[times(n,power(b,p-1-i,p),p)])
return (i+j-1)%(p-1);
return -1;
}

=====

/* Arithmetic: Cubic equation solver -----*/

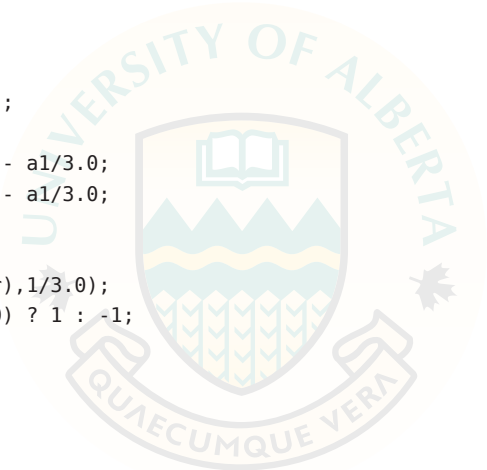
```

```

struct Result {
int n; // Number of solutions
double x[3]; // Solutions
};

Result solve_cubic(double a, double b, double c, double d) {
long double a1 = b/a, a2 = c/a, a3 = d/a;
long double q = (a1*a1 - 3*a2)/9.0, sq = -2*sqrt(q);
long double r = (2*a1*a1*a1 - 9*a1*a2 + 27*a3)/54.0;
double z = r*r-q*q*q, theta;
Result s;
if(z <= 0) {
s.n = 3; theta = acos(r/sqrt(q*q*q));
s.x[0] = sq*cos(theta/3.0) - a1/3.0;
s.x[1] = sq*cos((theta+2.0*PI)/3.0) - a1/3.0;
s.x[2] = sq*cos((theta+4.0*PI)/3.0) - a1/3.0;
}
else {
s.n = 1; s.x[0] = pow(sqrt(z)+fabs(r),1/3.0);
s.x[0] += q/s.x[0]; s.x[0] *= (r < 0) ? 1 : -1;
s.x[0] -= a1/3.0;
}
return s;
}

```




```
/* Combinatorics: Digit Occurrence count -----*/
```

```
// Given digit d and value N, returns # of times d occurs from 1..N
```

```
long long digit_count(int digit, int max) {
    long long res = 0; char buff[15];
    int i, count;
    if(max <= 0) return 0;
    res += max/10 + ((max % 10) >= digit ? 1 : 0);
    if(digit == 0) res--;
    res += digit_count(digit, max/10 - 1) * 10;
    sprintf(buff, "%d", max/10);
    for(i = 0, count = 0; i < strlen(buff); i++)
        if(buff[i] == digit+'0') count++;
    res += (1 + max%10) * count;
    return res;
}
```

```
/* Combinatorics: Josephus Ring Survivor (n people, dismiss every m'th) -----*/
```

```
int survive[MAXN];
void josephus(int n, int m) {
    survive[1] = 0;
    for(int i = 2; i <= n; i++)
        survive[i] = (survive[i-1] + (m%i))%i;
}
```

```
/* Combinatorics: Permutation index on distinct characters -----*/
```

```
// Returns perm. index of a string according to lex. ordering.
// Warning: does not work with repeated chars.
```

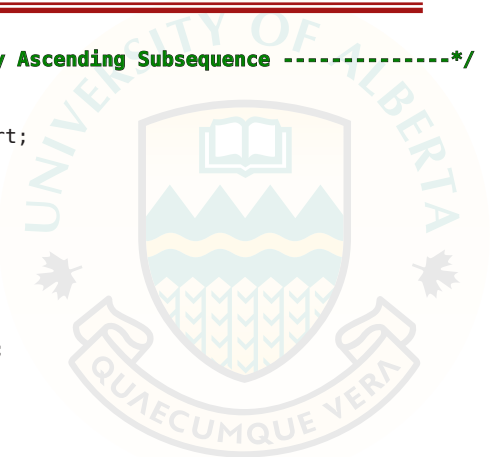
```
int permidx (char *s) {
    int size = strlen(s), index = 0;
    for (int i = 1; i < size; ++i) {
        for (int j = i; j < size; ++j)
            if (s[i-1] > s[j]) ++index;
        index *= size - i;
    }
    return index;
}
```

```
/* Dynamic Programming: Longest Ascending Subsequence -----*/
```

```
int asc_seq(int *A, int n, int *S) {
    int *m, *seq, i, k, low, up, mid, start;
    m = malloc((n+1) * sizeof(int));
    seq = malloc(n * sizeof(int));
    for (i = 0; i < n; i++) seq[i] = -1;
    m[1] = start = 0;
    for (k = i = 1; i < n; i++) {
        if (A[i] >= A[m[k]]) {
            seq[i] = m[k++]; start = m[k] = i;
        }
        else if (A[i] < A[m[1]])
            m[1] = i;
        else {
            low = 1; up = k;
            while (low != up-1) {
                mid = (low+up)/2;
                if (A[m[mid]] <= A[i]) low = mid;
                else up = mid;
            }
            seq[i] = m[low]; m[up] = i;
        }
    }
    for (i = k-1; i >= 0; i--) {
        S[i] = A[start]; start = seq[start];
    }
    free(m); free(seq);
    return k;
}
```

```
/* Dynamic Programming: Longest Strictly Ascending Subsequence -----*/
```

```
int sasc_seq(int *A, int n, int *S) {
    int *m, *seq, i, k, low, up, mid, start;
    m = malloc((n+1) * sizeof(int));
    seq = malloc(n * sizeof(int));
    for (i = 0; i < n; i++) seq[i] = -1;
    m[1] = start = 0;
    for (k = i = 1; i < n; i++) {
        if (A[i] > A[m[k]]) {
            seq[i] = m[k++]; start = m[k] = i;
        }
        else if (A[i] < A[m[1]])
            m[1] = i;
        else if (A[i] < A[m[k]]) {
```




```

low = 1; up = k;
while (low != up-1) {
    mid = (low+up)/2;
    if(A[m[mid]] <= A[i]) low = mid;
    else up = mid;
}
if (A[i] > A[m[low]]) {
    seq[i] = m[low]; m[up] = i;
}
}
}
for (i = k-1; i >= 0; i--) {
    S[i] = A[start]; start = seq[start];
}
}
free(m); free(seq);
return k;
}

```

/* Generators: Catalan Numbers -----*/

```

long long int cat[33];
void getcat() {
    cat[0] = cat[1] = 1;
    for (int i = 2; i < 33; ++i)
        cat[i] = cat[i-1]*(4*i-6)/i;
}

```

/* Generators: Binary Strings generator (cardinal order) -----*/

```

char bit[MAXN];
void recurse(int n, int curr, int left) {
    if(curr == n)
        Process(n);
    else {
        if(curr+left < n) {
            bit[curr] = 0; recurse(n, curr+1, left);
        }
        if(left) {
            bit[curr] = 1; recurse(n, curr+1, left-1);
        }
    }
}
void gen_bin_card(int n) {
    for(int i = 0; i <= n; i++) {
        printf("Cardinality %d:\n", i);
    }
}

```

```

recurse(n, 0, i);
}
}

```

/* Graph Theory: Maximum Bipartite Matching -----*/

```

// How to use (sample at bottom):
// For vertex i of set U:
//   match[i] = -1 means i is not matched
//   match[i] = x means the edge i->(x-|U|) is selected
//
// For simplicity, use addEdge(i,j,n) to add edges, where
// 0 <= i < |U| and 0 <= j < |V| and |U| = n.
// If there is an edge from vertex i of U to vertex
// j of V then: e[i][j+|U|] = e[j+|U|][i] = 1.
// - If |U| = n and |V| = m, then vertices are assumed
//   to be from [0,n-1] in set U and [0,m-1] in set V.
// - Remember that match[i]-n gives the edge from i, not just match[i].

```

```

const int MAXN 300          // How many vertices in U+V (in total)
char e[MAXN][MAXN];        // MODIFIED Adj. matrix (see note)
int match[MAXN], back[MAXN], q[MAXN], tail;
void addEdge(int x, int y, int n) {
    e[x][y+n] = e[y+n][x] = 1;
}
int find(int x, int n, int m) {
    int i, j, r;
    if(match[x] != -1) return 0;
    memset(back, -1, sizeof(back));
    for(q[i=0]=x, tail = 1; i < tail; i++)
        for(j = 0; j < n+m; j++) {
            if(!e[q[i]][j]) continue;
            if(match[j] != -1) {
                if(back[j] == -1) {
                    back[j] = q[i];
                    back[q[tail++]] = match[j] = j;
                }
            }
            else {
                match[match[q[i]] = j] = q[i];
                for(r = back[q[i]]; r != -1; r = back[back[r]])
                    match[match[r] = back[r]] = r;
                return 1;
            }
        }
}

```



```

    return 0;
}
void bipmatch(int n, int m) {
    memset(match, -1, sizeof(match));
    for(int i = 0; i < n+m; i++) if(find(i,n,m)) i = 0;
}

```

``` /* Graph Theory: Eulerian Graphs -----*/ ```

```

// Before adding edges, call Init() to initialize all data structures.
// Use the provided addEdge(x,y,c) which adds c edges between x and y.
//
// isEulerian(int n, int *start, int *end) returns:
//  0  if the graph is not Eulerian
//  1  if the graph has a Euler cycle
//  2  if the graph a path, from start to end
// with n being the number of nodes in the graph

```

```

const int MAXN 105    // Number of nodes
const int MAXM 505    // Maximum number of edges
#define min(a,b) (((a)<(b))?(a):(b))
#define max(a,b) (((a)>(b))?(a):(b))
#define DEC(a,b) g[a][b]--;g[b][a]--;deg[a]--;deg[b]--
int sets[MAXN], deg[MAXN], g[MAXN][MAXN];
int seq[MAXM], seqsize;
int getRoot(int x) {
    if (sets[x] < 0) return x;
    return sets[x] = getRoot(sets[x]);
}
void Union(int a, int b) {
    int ra = getRoot(a), rb = getRoot(b);
    if (ra != rb) {
        sets[ra] += sets[rb];
        sets[rb] = ra;
    }
}
void Init() {
    memset(sets, -1, sizeof(sets));
    memset(g, 0, sizeof(g));
    memset(deg, 0, sizeof(deg));
}
void addEdge(int x, int y, int count) {
    g[x][y] += count; deg[x] += count;
    g[y][x] += count; deg[y] += count;
    Union(x,y);
}

```

```

}
int isEulerian(int n, int *start, int *end) {
    int odd = 0, i, count = 0, x;
    for (i = 0; i < n; i++)
        if (deg[i]) {
            x = i; count++;
        }
    if (sets[getRoot(x)] != -count) return 0;
    for (i = 0; i < n; i++) {
        if (deg[i] & 1) {
            odd++;
            if(odd == 1) *start = i;
            else if(odd == 2) *end = i;
            else return 0;
        }
    }
    return odd ? 2 : 1;
}
void getPath(int n, int start, int end) {
    int temp[MAXM], tsize = 1, i, j;
    temp[0] = start;
    while(1) {
        j = temp[tsize-1];
        for (i = 0; i < n; i++) {
            if (i == end) continue;
            if (g[i][j]) {
                temp[tsize++] = i;
                DEC(i,j); break;
            }
        }
        if (i == n) {
            if (g[end][j]) {
                temp[tsize++] = end;
                DEC(j,end);
            }
            break;
        }
    }
    for (i = 0; i < tsize; i++)
        if (!deg[temp[i]])
            seq[seqsize++] = temp[i];
        else getPath(n, temp[i], temp[i]);
}
void buildPath(int n, int start, int end) {
    seqsize = 0; getPath(n, start, end);
}

```



```
/* Graph Theory: Maximum Flow in a directed graph -----*/
```

```
// Multiple edges from u to v may be added. They are converted into a
// single edge with a capacity equal to their sum
// - Vertices are assumed to be numbered from 0..n-1
// - The graph is supplied as the number of nodes (n), the zero-based
// indexes of the source (s) and the sink (t), and a vector of edges u->v
// with capacity c (M).
```

```
const int MAXN 200
struct Edge {
    //Edge u->v with capacity c
    int u, v, c;
};
int F[MAXN][MAXN]; //Flow of the graph
int maxFlow(int n, int s, int t, vector<Edge> &M) {
    int u, v, c, oh, min, df, flow, H[n], E[n], T[n], C[n][n];
    vector<Edge>::iterator m;
    list<int> N; list<int>::iterator cur;
    vector<int> R[n]; vector<int>::iterator r;
    for (u = 0; u < n; u++) {
        E[u] = H[u] = T[u] = 0;
        R[u].clear();
        for (v = 0; v < n; v++)
            C[u][v] = F[u][v] = 0;
    }
    for (m = M.begin(); m != M.end(); m++) {
        u = m->u; v = m->v; c = m->c;
        if (c && !C[u][v] && !C[v][u]) {
            R[u].push_back(v);
            R[v].push_back(u);
        }
        C[u][v] += c;
    }
    H[s] = n;
    for (r = R[s].begin(); r != R[s].end(); r++) {
        v = *r;
        F[s][v] = C[s][v]; F[v][s] = -C[s][v];
        E[v] = C[s][v]; E[s] -= C[s][v];
    }
    N.clear();
    for (u = 0; u < n; u++)
        if ((u != s) && (u != t))
            N.push_back(u);
    for (cur = N.begin(); cur != N.end(); cur++) {
        u = *cur; oh = H[u];
        while (E[u] > 0)
```

```
if (T[u] >= (int)R[u].size()) {
    min = 10000000;
    for (r = R[u].begin(); r != R[u].end(); r++) {
        v = *r;
        if ((C[u][v] - F[u][v] > 0) && (H[v] < min))
            min = H[v];
    }
    H[u] = 1 + min;
    T[u] = 0;
}
else {
    v = R[u][T[u]];
    if ((C[u][v] - F[u][v] > 0) && (H[u] == H[v]+1)) {
        df = C[u][v] - F[u][v];
        if (df > E[u])
            df = E[u];
        F[u][v] += df; F[v][u] = -F[u][v];
        E[u] -= df; E[v] += df;
    }
    else
        T[u]++;
}
if (H[u] > oh)
    N.splice(N.begin(), N, cur);
}
flow = 0;
for (r = R[s].begin(); r != R[s].end(); r++)
    flow += F[s][*r];
return flow;
}
```

```
/* Graph Theory: Chinese Postman Problem -----*/
```

```
// The maximum # of vertices solvable is roughly 20
```

```
#define MAXN 20
#define DISCONNECT -1
int g[MAXN][MAXN]; // Adj matrix (keep lowest cost if multiedge)
int deg[MAXN]; // Degree count
int A[MAXN+1]; // Used by perfect matching generator
int sum; // Sum of costs
int odd, best;

void floyd(int n) {
    for(int k = 0; k < n; k++)
```

```

    for(int i = 0; i < n; i++)
        if (g[i][k] != -1)
            for(int j = 0; j < n; j++)
                if (g[k][j] != -1)
                    if ((g[i][j] == -1) || (g[i][j] > g[i][k]+g[k][j]))
                        g[i][j] = g[i][k]+g[k][j];
    for(int i = 0; i < n; i++)
        g[i][i] = 0;
}
void checkSum() {
    int i, temp;
    for(i = temp = 0; i < odd/2; i++)
        temp += g[A[2*i]][A[2*i+1]];
    if(best == -1 || best > temp) best = temp;
}
void perfmach(int x) {
    int i, t;
    if(x == 2) checkSum();
    else {
        perfmach(x-2);
        for(i = x-3; i >= 0; i--) {
            t = A[i]; A[i] = A[x-2];
            A[x-2] = t; perfmach(x-2);
        }
        t = A[x-2];
        for(i = x-2; i >= 1; i--) A[i] = A[i-1];
        A[0] = t;
    }
}
int postman(int n) {
    int i; floyd(n);
    for(odd = i = 0; i < n; i++)
        if(deg[i]%2) A[odd++] = i;
    if(!odd) return sum;
    best = -1;
    perfmach(odd);
    return sum+best;
}
int main() {
    int i, u, v, c, n, m;
    while(scanf("%d %d", &n, &m) == 2){
        // Clear graph and degree count
        memset(g, -1, sizeof(g));
        memset(deg, 0, sizeof(deg));
        for(sum = i = 0; i < m; i++) {
            scanf("%d %d %d", &u, &v, &c);

```

```

            u--; v--; deg[u]++; deg[v]++;
            if(g[u][v] == -1 || g[u][v] > c) g[u][v] = c;
            if(g[v][u] == -1 || g[v][u] > c) g[v][u] = c;
            sum += c;
        }
        printf("Best cost: %d\n", postman(n));
    }
}

```

/* Graph Theory: Strongly Connected Components -----*/

```

vector<int> g[MAXN], curr;
vector< vector<int> > scc;
int dfsnum[MAXN], low[MAXN], id;
char done[MAXN];
void visit(int x) {
    curr.push_back(x);
    dfsnum[x] = low[x] = id++;
    for(size_t i = 0; i < g[x].size(); i++)
        if(dfsnum[g[x][i]] == -1){
            visit(g[x][i]);
            low[x] <= low[g[x][i]];
        }
        else if(!done[g[x][i]])
            low[x] <= dfsnum[g[x][i]];
    if(low[x] == dfsnum[x]) {
        VI c; int y;
        do {
            done[y = curr[curr.size()-1]] = 1;
            c.push_back(y);
            curr.pop_back();
        }
        while(y != x);
        scc.push_back(c);
    }
}
void strong_conn(int n) {
    memset(dfsnum, -1, n*sizeof(int));
    memset(done, 0, sizeof(done));
    scc.clear(); curr.clear();
    for(int i = id = 0; i < n; i++)
        if(dfsnum[i] == -1) visit(i);
}

```



/* Graph Theory: Min Cost Max Flow (Edmonds-Karp & Dijkstra) -----*/

```
// Takes a directed graph where each edge has a capacity ('cap') and a
// cost per unit of flow ('cost') and returns a maximum flow network
// of minimal cost ('fcost') from s to t. USE THIS CODE FOR (MODERATELY)
// DENSE GRAPHS; FOR VERY SPARSE GRAPHS, USE mcmf4 (next)
// PARAMETERS:
// - cap (global): adjacency matrix where cap[u][v] is the capacity
//   of the edge u->v. cap[u][v] is 0 for non-existent edges.
// - cost (global): a matrix where cost[u][v] is the cost per unit
//   of flow along the edge u->v. If cap[u][v] == 0, cost[u][v] is
//   ignored. ALL COSTS MUST BE NON-NEGATIVE!
// - n: the number of vertices ([0, n-1] are considered as vertices).
// - s: source vertex.
// - t: sink.
// RETURNS:
// - the flow
// - the total cost through 'fcost'
// - fnet contains the flow network. Careful: both fnet[u][v] and
//   fnet[v][u] could be positive. Take the difference.
// COMPLEXITY:
// - Worst case:  $O(n^2 \cdot \text{flow})$  <?  $n^3 \cdot \text{fcost}$ 
```

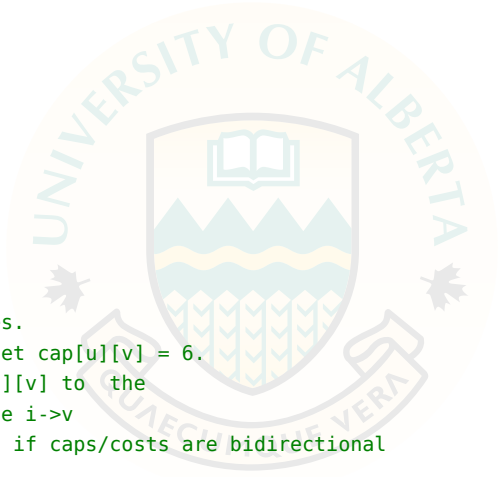
```
// Watch for commas when typing this in!
#define NN 1024 // the maximum number of vertices + 1
int cap[NN][NN]; // adjacency matrix (fill this up)
int cost[NN][NN]; // cost per unit of flow matrix (fill this up)
int fnet[NN][NN], adj[NN][NN], deg[NN]; // flow network and adjacency list
int par[NN], d[NN]; // par[source] = source;
int pi[NN]; // Labelling function
#define CLR(a, x) memset(a, x, sizeof(a))
#define Inf (INT_MAX/2)
#define Pot(u,v) (d[u] + pi[u] - pi[v])
bool dijkstra(int n, int s, int t) {
    // Dijkstra's using non-negative edge weights (cost + potential)
    for (int i = 0; i < n; i++)
        d[i] = Inf, par[i] = -1;
    d[s] = 0; par[s] = -n - 1;
    while (1) {
        int u = -1, bestD = Inf;
        for (int i = 0; i < n; i++)
            if (par[i] < 0 && d[i] < bestD)
                bestD = d[u = i];
        if (bestD == Inf) break;
        par[u] = -par[u] - 1;
        for (int i = 0; i < deg[u]; i++) {
            int v = adj[u][i];
```

```
            if (par[v] >= 0) continue;
            if (fnet[v][u] && d[v] > Pot(u,v) - cost[v][u])
                d[v] = Pot(u,v) - cost[v][u], par[v] = -u-1;
            if (fnet[u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v])
                d[v] = Pot(u,v) + cost[u][v], par[v] = -u - 1;
        }
    }
    for (int i = 0; i < n; i++)
        if (pi[i] < Inf)
            pi[i] += d[i];
    return par[t] >= 0;
}
```

#undef Pot

```
int mcmf3(int n, int s, int t, int &fcost) {
    CLR(deg, 0); CLR(fnet, 0); CLR(pi, 0);
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            if (cap[i][j] || cap[j][i])
                adj[i][deg[i]++] = j;
    int flow = fcost = 0;
    while (dijkstra(n, s, t)) {
        int bot = INT_MAX;
        for (int v = t, u = par[v]; v != s; u = par[v = u])
            bot <= fnet[v][u] ? fnet[v][u] : (cap[u][v] - fnet[u][v]);
        for (int v = t, u = par[v]; v != s; u = par[v = u])
            if (fnet[v][u]) {
                fnet[v][u] -= bot; fcost -= bot * cost[v][u];
            }
            else {
                fnet[u][v] += bot; fcost += bot * cost[u][v];
            }
        flow += bot;
    }
    return flow;
}
```

```
int main() {
    int numV; cin >> numV;
    memset(cap, 0, sizeof(cap));
    int m, a, b, c, cp, s, t;
    cin >> m >> s >> t;
    // fill up cap with existing capacities.
    // if the edge u->v has capacity 6, set cap[u][v] = 6.
    // for each cap[u][v] > 0, set cost[u][v] to the
    // cost per unit of flow along the edge u->v
    // Uncomment the commented statements if caps/costs are bidirectional
    for (int i=0; i<m; i++) {
```



```

cin >> a >> b >> cp >> c;
cost[a][b] = c; // cost[b][a] = c;
cap[a][b] = cp; // cap[b][a] = cp;
}
int fcost, flow = mcmf3(numV, s, t, fcost);
cout << "flow: " << flow << endl;
cout << "cost: " << fcost << endl;
}

```

/* Graph Theory: Min Cost Max Flow (Edmonds-Karp & fast heap Dijkstra) -----*/

// Same as above, but better for sparse graphs

```

#define NN 1024 // the maximum number of vertices + 1
int cap[NN][NN]; // adjacency matrix (fill this up)
int cost[NN][NN]; // cost per unit of flow matrix (fill this up)
int fnet[NN][NN], adj[NN][NN], deg[NN]; // flow network and adjacency list
int par[NN], d[NN], q[NN], inq[NN], qs; // Dijkstra's variables
int pi[NN]; // Labelling function
#define CLR(a, x) memset(a, x, sizeof(a))
#define Inf (INT_MAX/2)
#define BUBL { \
    t = q[i]; q[i] = q[j]; q[j] = t; \
    t = inq[q[i]]; inq[q[i]] = inq[q[j]]; inq[q[j]] = t; }
#define Pot(u,v) (d[u] + pi[u] - pi[v])
bool dijkstra(int n, int s, int t) {
    // Dijkstra's using non-negative edge weights (cost + potential)
    CLR(d, 0x3F); CLR(par, -1); CLR(inq, -1);
    d[s] = qs = 0;
    inq[q[qs++] = s] = 0;
    par[s] = n;
    while (qs) {
        int u = q[0]; inq[u] = -1;
        q[0] = q[--qs];
        if (qs) inq[q[0]] = 0;
        for (int i = 0, j = 2*i + 1, t; j < qs; i = j, j = 2*i + 1) {
            if (j + 1 < qs && d[q[j + 1]] < d[q[j]]) j++;
            if (d[q[j]] >= d[q[i]]) break;
            BUBL;
        }
        for (int k = 0, v = adj[u][k]; k < deg[u]; v = adj[u][++k]) {
            if (fnet[v][u] && d[v] > Pot(u,v) - cost[v][u])
                d[v] = Pot(u,v) - cost[v][par[v] = u];
            if (fnet[u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v])
                d[v] = Pot(u,v) + cost[par[v] = u][v];
        }
    }
}

```

```

if (par[v] == u) {
    if (inq[v] < 0) { inq[q[qs] = v] = qs; qs++; }
    for (int i=inq[v], j=(i-1)/2, t; d[q[i]]<d[q[j]]; i=j, j=(i-1)/2)
        BUBL;
    }
}
}
for (int i = 0; i < n; i++)
    if (pi[i] < Inf)
        pi[i] += d[i];
return par[t] >= 0;
}

#undef Pot
int mcmf4(int n, int s, int t, int &fcost) {
    CLR(deg, 0); CLR(fnet, 0); CLR(pi, 0);
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
            if (cap[i][j] || cap[j][i])
                adj[i][deg[i]++] = j;
    int flow = fcost = 0;
    while (dijkstra(n,s,t)) {
        int bot = INT_MAX;
        for (int v = t, u = par[v]; v != s; u = par[v = u])
            bot <= fnet[v][u] ? fnet[v][u] : (cap[u][v] - fnet[u][v]);
        for (int v = t, u = par[v]; v != s; u = par[v = u])
            if (fnet[v][u]) {
                fnet[v][u] -= bot; fcost -= bot * cost[v][u];
            }
            else {
                fnet[u][v] += bot; fcost += bot * cost[u][v];
            }
        flow += bot;
    }
    return flow;
}

```

/* Graph Theory: Articulation Points & Bridges (adj list) [O(V+E)] -----*/

// array entry art[v] is true iff vertex v is an articulation point
// - array entries bridge[i][0] and bridge[i][1] are the endpoints of a bridge
// in the graph. If bridge (u,v) is represented in the array, (v,u) is not.
// - 'bridges' is the number of bridges in the graph
// - index vertices from 0 to n-1

```

#define MAX_N 200

```

```

#define min(a,b) (((a)<(b))?(a):(b))
struct Node {
    int deg;
    int adj[MAX_N];
};
Node alist[MAX_N];
bool art[MAX_N], seen[MAX_N];
int df_num[MAX_N], low[MAX_N], father[MAX_N], cnt;
int bridge[MAX_N*MAX_N][2], bridges;
void add_edge(int v1, int v2) {
    alist[v1].adj[alist[v1].deg++] = v2;
    alist[v2].adj[alist[v2].deg++] = v1;
}
void add_bridge(int v1, int v2) {
    bridge[bridges][0] = v1;
    bridge[bridges][1] = v2;
    ++bridges;
}
void clear() {
    for (int i = 0; i < MAX_N; ++i)
        alist[i].deg = 0;
}
void search(int v, bool root) {
    int w, child = 0;
    seen[v] = true;
    low[v] = df_num[v] = cnt++;
    for (int i = 0; i < alist[v].deg; ++i) {
        w = alist[v].adj[i];
        if (df_num[w] == -1) {
            father[w] = v; ++child;
            search(w, false);
            if (low[w] > df_num[v]) add_bridge(v, w);
            if (low[w] >= df_num[v] && !root)
                art[v] = true;
            low[v] = min(low[v], low[w]);
        }
        else if (w != father[v]) {
            low[v] = min(low[v], df_num[w]);
        }
    }
    if (root && child > 1) art[v] = true;
}
void articulate(int n) {
    int child = 0;
    for (int i = 0; i < n; ++i) {
        art[i] = false;

```

```

        df_num[i] = father[i] = -1;
    }
    cnt = bridges = 0;
    memset(seen, false, sizeof(seen));
    for (int i = 0; i < n; ++i)
        if (!seen[i])
            search(i, true);
}
int main() {
    int n, m, v1, v2, c = 0;
    while (true) {
        scanf("%d %d", &n, &m);
        if (!n && !m) break;
        clear();
        for (int i = 0; i < m; ++i) {
            scanf("%d %d", &v1, &v2);
            add_edge(v1 - 1, v2 - 1);
        }
        articulate(n);
        printf("Articulation Points:");
        for (int i = 0; i < n; ++i)
            if (art[i]) printf(" %d", i + 1);
        printf("\n");
        printf("Bridges:");
        for (int i = 0; i < bridges; ++i)
            printf(" (%d,%d)", bridge[i][0] + 1, bridge[i][1] + 1);
        printf("\n\n");
    }
}

```

/* Graph Theory: Maximum Weighted Bipartite Matching [$O(n^3)$] -----*/

// Given N workers and N jobs to complete, where each worker has a certain compatibility (weight) to each job, find an assignment (perfect matching) of workers to jobs which maximizes the compatibility (weight).
// - W is a 2 dimensional array where $W[i][j]$ is the weight of worker i doing job j. Weights must be non-negative. If there is no weight assigned to a particular worker and job pair, set it to zero. If there is a different number of workers than jobs, create dummy workers or jobs accordingly with zero weight edges.
// - M is a 1 dimensional array populated by the algorithm where $M[i]$ is the index of the job matched to worker i.
// - This algorithm can be used with non-negative floating point weights.

```

#define MAX_N 100 // Max number of workers/jobs

```



```

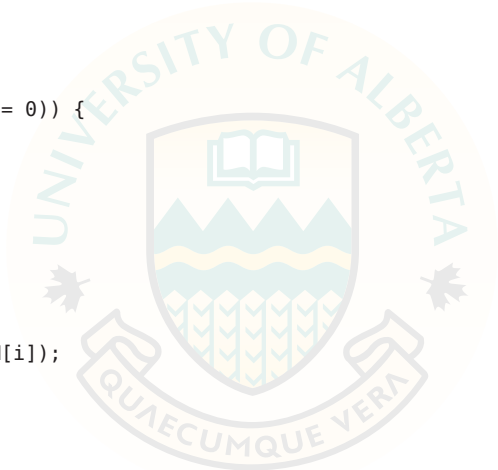
int W[MAX_N][MAX_N], U[MAX_N], V[MAX_N], Y[MAX_N]; // weight vars
int M[MAX_N], N[MAX_N], P[MAX_N], Q[MAX_N], R[MAX_N], S[MAX_N], T[MAX_N];
int Assign(int n) {
// Returns max weight, corresponding matching inside global M
int w, y; // weight vars
int i, j, m, p, q, s, t, v;
for (i = 0; i < n; i++) {
    M[i] = N[i] = -1; U[i] = V[i] = 0;
    for (j = 0; j < n; j++)
        if (W[i][j] > U[i])
            U[i] = W[i][j];
}
for (m = 0; m < n; m++) {
    for (p = i = 0; i < n; i++) {
        T[i] = 0; Y[i] = -1;
        if (M[i] == -1) {
            S[i] = 1; P[p++] = i;
        }
        else S[i] = 0;
    }
    while (1) {
        for (q = s = 0; s < p; s++) {
            i = P[s];
            for (j = 0; j < n; j++)
                if (!T[j]) {
                    y = U[i] + V[j] - W[i][j];
                    if (y == 0) {
                        R[j] = i;
                        if (N[j] == -1)
                            goto end_phase; // I hate goto's!
                        T[j] = 1; Q[q++] = j;
                    }
                    else if ((Y[j] == -1) || (y < Y[j])) {
                        Y[j] = y; R[j] = i;
                    }
                }
        }
    }
    if (q == 0) {
        y = -1;
        for (j = 0; j < n; j++)
            if (!T[j] && ((y == -1) || (Y[j] < y)))
                y = Y[j];
        for (j = 0; j < n; j++) {
            if (T[j])
                V[j] += y;
            if (S[j])

```

```

                U[j] -= y;
            }
        }
        for (j = 0; j < n; j++)
            if (!T[j]) {
                Y[j] -= y;
                if (Y[j] == 0) {
                    if (N[j] == -1)
                        goto end_phase; // again!
                    T[j] = 1; Q[q++] = j;
                }
            }
    }
    for (p = t = 0; t < q; t++) {
        i = N[Q[t]];
        S[i] = 1; P[p++] = i;
    }
}
end_phase:
i = R[j]; v = M[i];
M[i] = j; N[j] = i;
while (v != -1) {
    j = v; i = R[j];
    v = M[i];
    M[i] = j; N[j] = i;
}
}
for (i = w = 0; i < n; i++)
    w += W[i][M[i]];
return w;
}
int main() {
int w; // weight var
int n, i, j;
while ((scanf("%d", &n) == 1) && (n != 0)) {
    for (i = 0; i < n; i++)
        for (j = 0; j < n; j++)
            scanf("%d", &W[i][j]);
    w = Assign(n);
    printf("Optimum weight: %d\n", w);
    printf("Matchings:\n");
    for (i = 0; i < n; i++)
        printf("%d matched to %d\n", i, M[i]);
}
}

```



```
/* Graph Theory: Minimum weight Steiner tree  $O(|V|^3|S|+|V|^3)$  -----*/
```

```
// Given a weighted undirected graph  $G = (V, E)$  and a subset  $S$  of  $V$ ,
// finds a minimum weight tree  $T$  whose vertices are a superset of  $S$ .
// NP-hard -- this is a pseudo-polynomial algorithm.
//
// Minimum stc[(1<=s)-1][v] ( $0 \leq v < n$ ) is weight of min. Steiner tree
// Minimum stc[i][v] ( $0 \leq v < n$ ) is weight of min. Steiner tree for
// the i'th subset of Steiner vertices
//
//  $S$  is the list of Steiner vertices,  $s = |S|$ 
//  $d$  is the adjacency matrix (use infinities, not -1), and  $n = |V|$ 
```

```
const int N = 32;
const int K = 8;
int d[N][N], n, S[K], s, stc[1<=K][N];
void steiner() {
    for (int k = 0; k < n; ++k)
        for (int i = 0; i < n; ++i)
            for (int j = 0; j < n; ++j)
                d[i][j] <?= d[i][k] + d[k][j];
    for(int i = 1; i < (1<=s); ++i) {
        if (!(i&(i-1))) {
            int u;
            for (int j = i, k = 0; j; u = S[k++], j >= 1);
            for (int v = 0; v < n; ++v)
                stc[i][v] = d[v][u];
        }
        else for (int v = 0; v < n; ++v) {
            stc[i][v] = 0xffffffff;
            for (int j = 1; j < i; ++j)
                if ((j|i) == i) {
                    int x1 = j, x2 = i&(~j);
                    for (int w = 0; w < n; ++w)
                        stc[i][v] <?= d[v][w] + stc[x1][w] + stc[x2][w];
                }
        }
    }
}
```

```
/* Linear Programming: Simplex Method -----*/
```

```
// m - number of (less than) inequalities
// n - number of variables
// C - (m+1) by (n+1) array of coefficients:
// row 0 - objective function coefficients
```

```
// row 1:m - less-than inequalities
// column 0:n-1 - inequality coefficients
// column n - inequality constants (0 for objective function)
// X[n] - result variables
// return value - maximum value of objective function
// (-inf for infeasible, inf for unbounded)
```

```
#define MAXM 400 // leave one extra
#define MAXN 400 // leave one extra
#define EPS 1e-9
#define INF 1.0/0.0
double A[MAXM][MAXN];
int basis[MAXM], out[MAXN];
void pivot(int m, int n, int a, int b) {
    int i,j;
    for (i=0;i<=m;i++)
        if (i!=a)
            for (j=0;j<=n;j++)
                if (j!=b)
                    A[i][j] -= A[a][j] * A[i][b] / A[a][b];
    for (j=0;j<=n;j++)
        if (j!=b) A[a][j] /= A[a][b];
    for (i=0;i<=m;i++)
        if (i!=a) A[i][b] = -A[i][b]/A[a][b];
    A[a][b] = 1/A[a][b];
    i = basis[a]; basis[a] = out[b]; out[b] = i;
}
double simplex(int m, int n, double C[][MAXN], double X[]) {
    int i,j,ii,jj; // i,ii are row indexes; j,jj are column indexes
    for (i=1;i<=m;i++)
        for (j=0;j<=n;j++)
            A[i][j] = C[i][j];
    for (j=0;j<=n;j++)
        A[0][j] = -C[0][j];
    for (i=0;i<=m;i++)
        basis[i] = -i;
    for (j=0;j<=n;j++)
        out[j] = j;
    for(;;) {
        for (i=ii=1;i<=m;i++)
            if (A[i][n]<A[ii][n] || (A[i][n]==A[ii][n] && basis[i]<basis[ii]))
                ii=i;
        if (A[ii][n] >= -EPS) break;
        for (j=jj=0;j<=n;j++)
            if (A[ii][j]<A[ii][jj]-EPS || (A[ii][j]<A[ii][jj]-EPS && out[i]<out[j]))
                jj=j;
    }
```



```

    if (A[ii][jj] >= -EPS) return -INF;
    pivot(m,n,ii,jj);
}
for(;;) {
    for (j=jj=0;j<n;j++)
        if (A[0][j]<A[0][jj] || (A[0][j]==A[0][jj] && out[j]<out[jj]))
            jj=j;
    if (A[0][jj] > -EPS) break;
    for (i=1,ii=0;i<=m;i++)
        if ((A[i][jj]>EPS) &&
            (!ii || (A[i][n]/A[i][jj] < A[ii][n]/A[ii][jj]-EPS) ||
            ((A[i][n]/A[i][jj] < A[ii][n]/A[ii][jj]+EPS) &&
            (basis[i] < basis[ii]))))
            ii=i;
    if (A[ii][jj] <= EPS) return INF;
    pivot(m,n,ii,jj);
}
for (j=0;j<n;j++)
    X[j] = 0;
for (i=1;i<=m;i++)
    if (basis[i] >= 0)
        X[basis[i]] = A[i][n];
return A[0][n];
}

```

=====
/* Java Template: IO Reference -----*/

// Description: This document is a reference for the use of java for regular
// IO purposes. It covers stdin and stdout as well as file IO.
// It also shows how to use StringTokenizer for parsing.

```

import java.util.*;
import java.io.*;
class IO {
    public static void main(String[] args) {
        try {
            // For file IO, use:
            // BufferedReader in=new BufferedReader(new FileReader("probl.dat"));
            // PrintWriter out=new PrintWriter(
            //     new BufferedWriter(new FileWriter("probl.out")));
            // For stdin/stdout IO, use:
            PrintStream out = System.out;
            BufferedReader in = new BufferedReader(new InputStreamReader(System.in));
            String line;
            int num=0;

```

```

StringTokenizer st;
while(true) {
    // Newlines are removed by readLine()
    line = in.readLine();
    if(line == null) break;
    num++;
    out.println("Line #" + num);
    // Split on whitespace
    st = new StringTokenizer(line);
    while(st.hasMoreTokens()) {
        out.print("Token: ");
        out.println(st.nextToken());
    }
    // To split on something else, use:
    // st = new StringTokenizer(line, delim);
    // Or use this to change in the middle of parsing:
    // line = st.nextToken(delim);
}
// You must flush for files!
out.flush();
}
catch (Exception e) {
    System.err.println(e.toString());
}
}
}

```

=====
/* Java Template: BigInteger Reference -----*/

// Description: This document is a reference for the use of the BigInteger
// class in Java. It contains code to compute GCDs of integers.
// Constants:
// -----
// BigInteger.ONE - The BigInteger constant one.
// BigInteger.ZERO - The BigInteger constant zero.
// Creating BigIntegers
// -----
// 1. From Strings
// a) BigInteger(String val);
// b) BigInteger(String val, int radix);
// 2. From byte arrays
// a) BigInteger(byte[] val);
// b) BigInteger(int signum, byte[] magnitude)
// 3. From a long integer
// a) static BigInteger BigInteger.valueOf(long val)

```

// Math operations:
// -----
// A + B = C      --> C = A.add(B);
// A - B = C      --> C = A.subtract(B);
// A * B = C      --> C = A.multiply(B);
// A / B = C      --> C = A.divide(B);
// A % B = C      --> C = A.remainder(B);
// A % B = C where C > 0 --> C = A.mod(B);
// A / B = Q & A % B = R --> C = A.divideAndRemainder(B);
//                      (Q = C[0], R = C[1])
// A ^ b = C      --> C = A.pow(B);
// abs(A) = C     --> C = A.abs();
// -(A) = C       --> C = A.negate();
//
// gcd(A,B) = C   --> C = A.gcd(B);
// (A ^ B) % M     --> C = A.modPow(B,M);
// C = inverse of A mod M --> C = A.modInverse(M);
// max(A,B) = C   --> C = A.max(B);
// min(A,B) = C   --> C = A.min(B);
// Bit Operations
// -----
// ~A = C        (NOT)  --> C = A.not();
// A & B = C      (AND)  --> C = A.and(B);
// A | B = C      (OR)   --> C = A.or(B);
// A ^ B = C      (XOR)  --> C = A.xor(B);
// A & ~B = C     (ANDNOT) --> C = A.andNot(B);
// A << n = C     (LSHIFT) --> C = A.shiftLeft(n);
// A >> n = C     (RSHIFT) --> C = A.shiftRight(n);
// Clear n'th bit of A --> C = A.clearBit(n);
// Set n'th bit of A --> C = A.setBit(n);
// Flip n'th bit of A --> C = A.flipBit(n);
// Test n'th bit of A --> C = A.testBit(n);
//
// Bitcount of A = n --> n = A.bitCount();
// Bitlength of A = n --> n = A.bitLength();
// Lowest set bit of A --> n = A.getLowestSetBit();
// Comparison Operations
// -----
// A < B          --> A.compareTo(B) == -1;
// A == B         --> A.compareTo(B) == 0
//               or A.equals(B);
// A > B          --> A.compareTo(B) == 1;
// A < 0          --> A.signum() == -1;
// A == 0         --> A.signum() == 0;
// A > 0          --> A.signum() == 1;
// Conversion:

```

```

// -----
// double        --> A.doubleValue();
// float         --> A.floatValue();
// int           --> A.intValue();
// long          --> A.longValue();
// byte[]        --> A.toByteArray();
// String        --> A.toString();
// String (base b) --> A.toString(b);

import java.math.*;
import java.io.*;
import java.util.*;
class BigIntegers {
    public static void main(String[] args) {
        BufferedReader in = new BufferedReader(new InputStreamReader(System.in));
        String line;
        StringTokenizer st;
        BigInteger a;
        BigInteger b;
        try {
            while(true) {
                line = in.readLine();
                if(line == null) break;
                st = new StringTokenizer(line);
                a = new BigInteger(st.nextToken());
                b = new BigInteger(st.nextToken());
                System.out.println( a.gcd(b) );
            }
        } catch (Exception e) {
            System.err.println(e.toString());
        }
    }
}



---


/* Number Theory: Converting between bases (Java, arb. precision) -----*/
// Converts from base b1 to base b2
import java.math.*;
import java.io.*;
import java.util.*;
class base_convert {
    // invalid is the string that is returned if the N is not valid
    static String invalid = new String("Number is not valid");
    private static String convert_base(int b1, int b2, String n, String key) {

```

```

int i, x;
String n2 = "", n3 = "";
BigInteger a = BigInteger.ZERO,
    b1 = BigInteger.valueOf(base1),
    b2 = BigInteger.valueOf(base2);
for (i = 0; i < n.length(); i++) {
    a = a.multiply(b1);
    x = key.indexOf(n.charAt(i));
    if (x == -1 || x >= base1) return invalid;
    a = a.add(BigInteger.valueOf(x));
}
while (a.signum() == 1) {
    BigInteger r[] = a.divideAndRemainder(b2);
    n2 += key.charAt(r[1].intValue());
    a = r[0];
}
for (i = n2.length()-1; i >= 0; i--) n3 += n2.charAt(i);
if (n3.length() == 0) n3 += '0';
return n3;
}

public static void main(String[] args) {
    try {
        String line, n;
        int tnum, base1, base2;
        StringTokenizer st;
        // key is the base system that you may change as needed
        String key = new
String("0123456789ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz");
        // Standard IO
        BufferedReader in = new BufferedReader(new InputStreamReader(System.in));
        PrintStream out = System.out;
        // File IO
        // BufferedReader in = new BufferedReader(new FileReader("probl.dat"));
        // PrintWriter out = new BufferedWriter(new FileWriter("probl.out"));
        line = in.readLine(); // Get number of test cases
        st = new StringTokenizer(line);
        tnum = Integer.parseInt(st.nextToken());
        for (int t = 0; t < tnum; t++) {
            line = in.readLine();
            st = new StringTokenizer(line);
            base1 = Integer.parseInt(st.nextToken());
            base2 = Integer.parseInt(st.nextToken());
            n = st.nextToken();
            String result = convert_base(base1, base2, n, key2);
            out.println(result);
        }
    }
}

```

```

    }
    catch (Exception e) {
        System.err.println(e.toString());
    }
}
}

```

/* Number Theory: Primality Testing -----*/

```

bool isPrime(int x) {
    if(x == 1) return ONEPRIME;
    if(x == 2) return true;
    if(!(x & 1)) return false;
    for(int i = 3; i*i <= x; i += 2) // watch for overflow
        if (x % i) return false;
    return true;
}

```

/* Number Theory: Number of Divisors [O(sqrt(N))] -----*/

```

int num_divisors(int n) {
    int i, count, res = 1;
    for(i = 2; i*i <= n; i++) {
        count = 0;
        while(!(n%i)) {
            n /= i; count++;
        }
        if(count) res *= (count+1);
    }
    if (n > 1) res *= 2;
    return res;
}

```

/* Number Theory: Prime Factorization -----*/

```

int primes[MAXP]; int psize;
void getPrimes() {
    int i, j, isprime;
    psize = 0; primes[psize++] = 2;
    for (i = 3; i <= MAXN; i+= 2) {
        for (isprime = j = 1; j < psize; j++) {
            if (i % primes[j] == 0) {
                isprime = 0;
                break;
            }
        }
        if (isprime) primes[psize++] = i;
    }
}

```

```

    }
    if (1.0*primes[j]*primes[j] > i) break;
}
if(isprime) primes[psize++] = i;
}
}
struct Factors {
    int size;
    int f[32];
};
Factors getPFactor(int n) {
    Factors x;
    int i;
    x.size = 0;
    for (i = 0; i < psize; i++) {
        while (n % primes[i] == 0) {
            x.f[x.size++] = primes[i];
            n /= primes[i];
        }
        if(1.0*primes[i]*primes[i] > n) break;
    }
    if(n > 1)
        x.f[x.size++] = n;
    return x;
}

```

/* Number Theory: Primality testing with a sieve -----*/

// Consider using typedefs and functions instead of defines...

```

#define TEST(f,x) (*(f+(x)/16)&(1<<(((x)%16L)/2)))
#define SET(f,x) *(f+(x)/16)|=1<<(((x)%16L)/2)
#define ONEPRIME 0 // whether or not 1 is considered to be prime
#define UL unsigned long
#define UC unsigned char
UC *primes = NULL;
UL getPrimes(UL maxn) {
    UL x, y, psize=1;
    primes = calloc(((maxn)>>4)+1L, sizeof(UC));
    for (x = 3; x*x <= maxn; x+=2)
        if (!TEST(primes, x))
            for (y = x*x; y <= maxn; y += x<1) SET (primes, y);
    // Comment out if you don't need # of primes <= maxn
    for(x = 3; x <= maxn; x+=2)
        if(!TEST(primes, x)) psize++;
}

```

```

    return psize;
}
int isPrime(UL x) {
    // Returns whether or not a given POSITIVE number is prime
    if(x == 1) return ONEPRIME;
    if(x == 2) return 1;
    if(x % 2 == 0) return 0;
    return (!TEST(primes, x));
}

```

/* Number Theory: Sum of divisors [0(sqrt(N))] -----*/

```

typedef long long int LL;
LL sum_divisors(LL n) {
    int i, count; LL res = 1;
    for (i = 2; i*i <= n; i++) {
        count = 0;
        while (n % i == 0) {
            n /= i; count++;
        }
        if (count) res *= (pow(i, count+1)-1)/(i-1);
    }
    if(n > 1) res *= (pow(n, 2)-1)/(n-1);
    return res;
}

```

/* Number Theory: Chinese Remainder Theorem -----*/

// Given n relatively prime modular in m[0], ..., m[n-1], and right-hand sides a[0], ..., a[n-1], the routine solves for the unique solution // in the range $0 \leq x < m[0]*m[1]*\dots*m[n-1]$ such that $x = a[i] \bmod m[i]$ // for all $0 \leq i < n$. The algorithm used is Garner's algorithm, which // is not the same as the one usually used in number theory textbooks. // // It is assumed that m[i] are positive and pairwise relatively prime. // a[i] can be any integer. // If the system of equations is // $x = a[0] \bmod m[0]$ // $x = a[1] \bmod m[1]$ // ... // then a[i] should be reduced mod m[i] first. // Also, if $0 \leq a[i] < m[i]$ for all i, then the answer will fall // in the range $0 \leq x < m[0]*m[1]*\dots*m[n-1]$.

```

int gcd(int a, int b, int *s, int *t) {

```

```

int r, r1, r2, a1, a2, b1, b2, q;
a1 = b2 = 1;
a2 = b1 = 0;
while (b) {
    q = a / b; r = a % b;
    r1 = a1 - q*b1;
    r2 = a2 - q*b2;
    a = b; a1 = b1; a2 = b2;
    b = r; b1 = r1; b2 = r2;
}
*s = a1; *t = a2;
return a;
}

int cra(int n, int *m, int *a) {
    int x, i, k, prod, temp;
    int *gamma, *v;
    gamma = malloc(n*sizeof(int));
    v = malloc(n*sizeof(int));
    for (k = 1; k < n; k++) {
        prod = m[0] % m[k];
        for (i = 1; i < k; i++) {
            prod = (prod * m[i]) % m[k];
        }
        gcd(prod, m[k], gamma+k, &temp);
        gamma[k] %= m[k];
        if (gamma[k] < 0)
            gamma[k] += m[k];
    }
    v[0] = a[0];
    for (k = 1; k < n; k++) {
        temp = v[k-1];
        for (i = k-2; i >= 0; i--) {
            temp = (temp * m[i] + v[i]) % m[k];
            if (temp < 0)
                temp += m[k];
        }
        v[k] = ((a[k] - temp) * gamma[k]) % m[k];
        if (v[k] < 0)
            v[k] += m[k];
    }
    x = v[n-1];
    for (k = n-2; k >= 0; k--)
        x = x * m[k] + v[k];
    free(gamma); free(v);
    return x;
}

```

```

int main(void) {
    int n, *m, *a, i, x;
    while (scanf("%d", &n) == 1 && n > 0) {
        m = malloc(n*sizeof(int));
        a = malloc(n*sizeof(int));
        printf("Enter moduli:\n");
        for (i = 0; i < n; i++)
            scanf("%d", m+i);
        printf("Enter right-hand side:\n");
        for (i = 0; i < n; i++)
            scanf("%d", a+i);
        x = cra(n, m, a);
        printf("x = %d\n", x);
        free(m); free(a);
    }
}

```

/* Number Theory: Extended Euclidean Algorithm -----*/

// Assumes non-negative input. Returns d s.t. $d = ax + by$
 // x,y passed in by reference, #include <algorithm> for swap function

```

int gcd(int a, int b, int &x, int &y) {
    x = 1; y = 0; int nx = 0, ny = 1;
    while (b) {
        int q = a/b;
        x -= q*nx; swap(x, nx);
        y -= q*ny; swap(y, ny);
        a -= q*b; swap(a, b);
    }
    return a;
}

```

/* Number Theory: Generalized Chinese Remaindering -----*/

// Given $[a_0, \dots, a_{n-1}]$ and $[m_0, \dots, m_{n-1}]$
 // Computes $0 \leq x < \text{lcm}(m_0, \dots, m_{n-1})$ such that
 // $x \equiv a_0 \pmod{m_0}, \dots, x \equiv a_{n-1} \pmod{m_{n-1}}$, if
 // such an x exists.
 // True is returned iff such an x exists. If x does not exist then the value
 // at the address of x will not be affected.
 // Complexity: $O(n \log(\text{MAX}(m_0, \dots, m_{n-1})))$

```

typedef long long int LLI;
LLI safe_mod(LLI a, LLI m) {

```



```

if (a < 0) return (a + m + m * (-a/m)) % m;
else return a % m;
}
LLI abs(LLI a) {
    return a < 0 ? -a : a;
}
LLI gcdex(LLI a, LLI b, LLI *ss, LLI *tt) {
    LLI q, r[150], s[150], t[150];
    int num = 2;
    r[0] = a; r[1] = b;
    s[0] = t[1] = 1;
    s[1] = t[0] = 0;
    while (r[num - 1]) {
        q = r[num - 2] / r[num - 1];
        r[num] = r[num - 2] % r[num - 1];
        s[num] = s[num - 2] - q * s[num - 1];
        t[num] = t[num - 2] - q * t[num - 1];
        ++num;
    }
    *ss = s[num - 2]; *tt = t[num - 2];
    return r[num - 2];
}
bool gen_chrem(LLI *a, LLI *m, int n, LLI *x) {
    LLI g, s, t, a_tmp = safe_mod(a[0], m[0]), m_tmp = m[0];
    for (int i = 1; i < n; ++i) {
        g = gcdex(m_tmp, m[i], &s, &t);
        if (abs(a_tmp - a[i]) % g) return false;
        a_tmp = safe_mod(a_tmp + (a[i] - a_tmp) / g * s * m_tmp, m_tmp/g*m[i]);
        m_tmp = m[i];
    }
    x = a_tmp;
    return true;
}
int main() {
    int n; LLI a[20], m[20], x;
    while (true) {
        scanf("%lld", &n);
        if (!n) break;
        for (int i = 0; i < n; ++i)
            scanf("%lld %lld", &a[i], &m[i]);
        if (!gen_chrem(a, m, n, &x))
            printf("No solution.\n\n");
        else
            printf("X = %lld\n\n", x);
    }
}

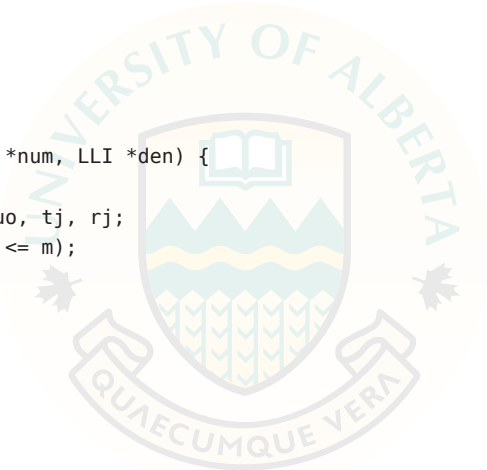
```

```

/* Number Theory: Rational Reconstruction [O(log m)] -----*/
// Description: Given integers m, g and k, computes integers 'num' and 'den'
// (if they exist) such that num == g*den mod m where |num| < k and
// 0 < den < g/k. True is returned iff den is invertible mod m. This algorithm
// is useful if computations on rational numbers is to be used when the input
// and output numbers have small numerators and denominators but intermediate
// results can have very large numerators and denominators. To use in this
// fashion, reduce the input rationals modulo some number m (probably a prime),
// perform the operations modulo m and then use rational reconstruction to
// recover the results. m and k must be selected such that |num|, den < k
// and 2*k*k < m for all input and output rational numbers.

typedef long long int LLI;
int gcd_table(LLI a, LLI b, LLI *r, LLI *q, LLI *s, LLI *t) {
    int n = 2;
    assert(0 <= a && 0 < b);
    r[0] = a; r[1] = b;
    s[0] = t[1] = 1;
    s[1] = t[0] = 0;
    while (r[n - 1]) {
        r[n] = r[n - 2] % r[n - 1];
        q[n - 1] = r[n - 2] / r[n - 1];
        s[n] = s[n - 2] - s[n - 1] * q[n - 1];
        t[n] = t[n - 2] - t[n - 1] * q[n - 1];
        ++n;
    }
    return n;
}
LLI gcd(LLI a, LLI b) {
    if (a < 0) return gcd(-a, b);
    if (b < 0) return gcd(a, -b);
    if (!b) return a;
    return gcd(b, a % b);
}
bool rat_recon(LLI m, LLI g, LLI k, LLI *num, LLI *den) {
    int n, j;
    LLI r[200], q[200], s[200], t[200], quo, tj, rj;
    assert(0 <= g && g < m && 1 <= k && k <= m);
    n = gcd_table(m, g, r, q, s, t);
    q[0] = q[n - 1] = 0;
    for (j = 0; j < n && r[j] >= k; ++j);
    if (t[j] > 0) {
        *num = r[j]; *den = t[j];
    }
    else {
        *num = -r[j]; *den = -t[j];
    }
}

```



```

    }
    if (gcd(r[j], t[j]) == 1) return true;
    else {
        quo = (j == n - 1 ? 0 : (k - r[j-1]) / r[j] + 1);
        rj = r[j - 1] - quo*r[j];
        tj = t[j - 1] - quo*t[j];
        if (gcd(rj, tj) != 1 || (tj > 0 ? tj : -tj) * k > m)
            return false;
        if (tj > 0) {
            *num = rj; *den = tj;
        }
        else {
            *num = -rj; *den = -tj;
        }
        return true;
    }
}

```

/* Search: Golden section search -----*/

// Given an function f(x) with a single local minimum, a lower and upper
// bound on x, and a tolerance for convergence, this function finds the
// minimizing value of x. f(x) should evaluate globally.

```

#define GOLD 0.381966 // 1/phi^2 = 1/(phi+1) = (phi-1)^2
#define move(a,b,c) x[a]=x[b];x[b]=x[c];fx[a]=fx[b];fx[b]=fx[c]
double f(double x) { return x*x; } // Just an example
double golden(double xlow, double xhigh, double tol) {
    double x[4], fx[4], L;
    int iter = 0, left = 0, mini, i;
    fx[0] = f(x[0]=xlow);
    fx[3] = f(x[3]=xhigh);
    while (1) {
        L = x[3]-x[0];
        if (!iter || left) {
            x[1] = x[0]+GOLD*L;
            fx[1] = f(x[1]);
        }
        if (!iter || !left) {
            x[2] = x[3]-GOLD*L;
            fx[2] = f(x[2]);
        }
        for (mini = 0, i = 1; i < 4; i++)
            if (fx[i] < fx[mini]) mini = i;
        if (L < tol) break;
    }
}

```

```

    if (mini < 2) {
        left = 1; move(3,2,1);
    }
    else {
        left = 0; move(0,1,2);
    }
    iter++;
}
return x[mini];
}

```

/* Search: KMP String Matching -----*/

// Given strings T and P, computes the indices of T where P occurs as a
// substring, stored in 'shift'. Returns the number of such indices.

```

#define MAX_LEN 1000

```

```

int pi[MAX_LEN];

```

```

void compute_prefix(char *P, int m, int *pi) {
    int k = pi[0] = -1;
    for (int q = 1; q < m; ++q) {
        while (k >= 0 && P[k + 1] != P[q]) k = pi[k];
        if (P[k + 1] == P[q]) ++k;
        pi[q] = k;
    }
}

```

```

int kmp_match(char *T, char *P, int *shift) {
    int n, m, q = -1, shifts = 0;
    n = strlen(T); m = strlen(P);
    compute_prefix(P, m, pi);
    for (int i = 0; i < n; ++i) {
        while (q > -1 && P[q + 1] != T[i]) q = pi[q];
        if (P[q + 1] == T[i]) ++q;
        if (q == m - 1) {
            shift[shifts++] = i - m + 1;
            q = pi[q];
        }
    }
    return shifts;
}

```



/* Search: Suffix array [O(N log N)] -----*/

```
// Notes: The build_sarray routine takes in a string S of n characters
// (null-terminated), and constructs two arrays 'sarray' and 'lcp'.
// - If p = sarray[i], then the suffix of str starting at p (i.e. S[p..n-1])
// is the i-th suffix (lexographically ordered)
// - NOTE: the empty suffix is not considered, so sarray[0] != n.
// - lcp[i] contains the length of the longest common prefix of the suffixes
// pointed to by sarray[i-1] and sarray[i] (but lcp[0] = 0).
// - To find a pattern P in str, you can look for it as the prefix of a
// suffix. This takes O(|P| log n) time with a binary search.
```

// You probably need to #include <climits> here.

```
#define MAXN 100000
```

```
int bucket[CHAR_MAX-CHAR_MIN+1];
```

```
int prn[MAXN], count[MAXN];
```

```
char bh[MAXN+1];
```

```
void build_sarray(char *str, int* sarray, int *lcp) {
```

```
    int n = strlen(str), a, c, d, e, f, h, i, j, x;
```

```
    memset(bucket, -1, sizeof(bucket));
```

```
    for (i = 0; i < n; i++) {
```

```
        j = str[i] - CHAR_MIN;
```

```
        prn[i] = bucket[j];
```

```
        bucket[j] = i;
```

```
    }
```

```
    for (a = c = 0; a <= CHAR_MAX - CHAR_MIN; a++)
```

```
        for (i = bucket[a]; i != -1; i = j) {
```

```
            j = prn[i]; prn[i] = c;
```

```
            bh[c++] = (i == bucket[a]);
```

```
        }
```

```
    bh[n] = 1;
```

```
    for (i = 0; i < n; i++)
```

```
        sarray[prn[i]] = i;
```

```
    x = 0;
```

```
    for (h = 1; h < n; h *= 2) {
```

```
        for (i = 0; i < n; i++) {
```

```
            if (bh[i] & 1) {
```

```
                x = i; count[x] = 0;
```

```
            }
```

```
            prn[sarray[i]] = x;
```

```
        }
```

```
    d = n - h; e = prn[d];
```

```
    prn[d] = e + count[e]++;
```

```
    bh[prn[d]] |= 2;
```

```
    i = 0;
```

```
    while (i < n) {
```

```
        for (j = i; (j == i || !(bh[j] & 1)) && j < n; j++) {
```

```
            d = sarray[j] - h;
```

```
            if (d >= 0) {
```

```
                e = prn[d]; prn[d] = e + count[e]++; bh[prn[d]] |= 2;
```

```
            }
```

```
        }
```

```
    for (j = i; (j == i || !(bh[j] & 1)) && j < n; j++) {
```

```
        d = sarray[j] - h;
```

```
        if (d >= 0 && (bh[prn[d]] & 2)) {
```

```
            for (e = prn[d]+1; bh[e] == 2; e++);
```

```
            for (f = prn[d]+1; f < e; f++)
```

```
                bh[f] &= 1;
```

```
        }
```

```
    }
```

```
    i = j;
```

```
    }
```

```
    for (i = 0; i < n; i++) {
```

```
        sarray[prn[i]] = i;
```

```
        if (bh[i] == 2)
```

```
            bh[i] = 3;
```

```
    }
```

```
    }
```

```
    h = 0;
```

```
    for (i = 0; i < n; i++) {
```

```
        e = prn[i];
```

```
        if (e > 0) {
```

```
            j = sarray[e-1];
```

```
            while (str[i+h] == str[j+h])
```

```
                h++;
```

```
            lcp[e] = h;
```

```
            if (h > 0) h--;
```

```
        }
```

```
    }
```

```
    lcp[0] = 0;
```

```
    }
```

