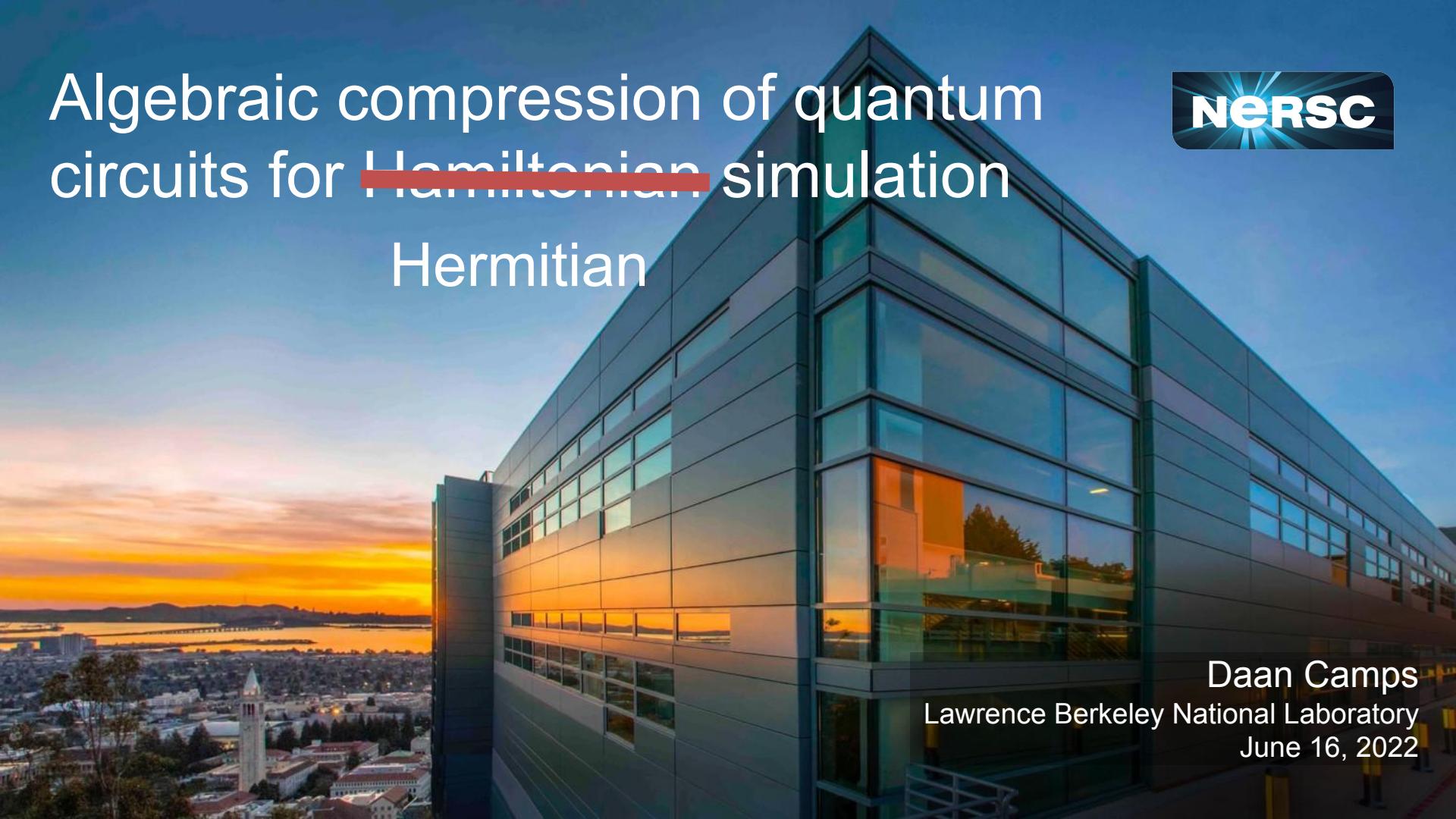


Algebraic compression of quantum circuits for ~~Hamiltonian~~ simulation

Hermitian



Daan Camps
Lawrence Berkeley National Laboratory
June 16, 2022

Outline and acknowledgements

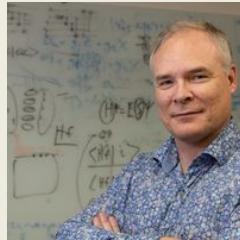
1. A 3 minute **introduction** to quantum computing
2. **Hamiltonian simulation** and Trotter decompositions
3. Algebraic **compression** of Trotterized circuits for spin Hamiltonians
4. **Results** on classical and quantum hardware
5. Conclusion



Efekan Kökcü



Lindsay Bassman



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Lex Kemper

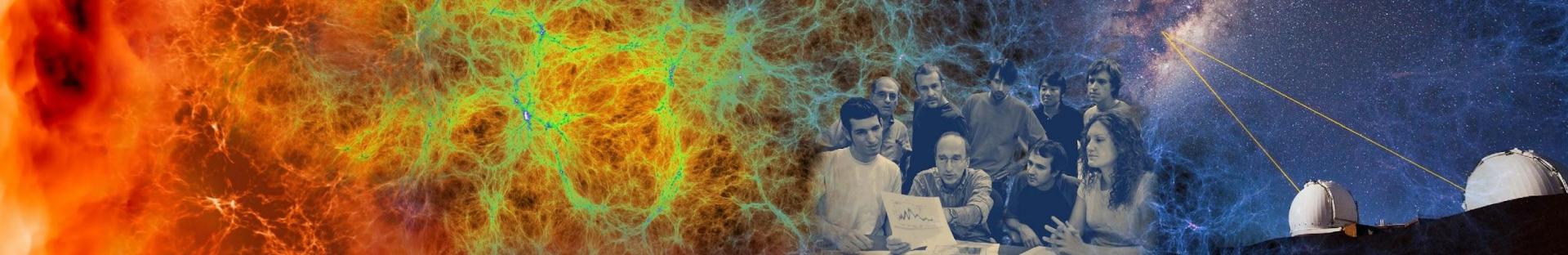


Roel Van Beeumen

An Algebraic Quantum Circuit Compression Algorithm for Hamiltonian Simulation, D. Camps, E. Kökcü, L. Bassman, W. A. de Jong, A. F. Kemper, R. Van Beeumen, Accepted in SIMAX, arXiv:2108.03283

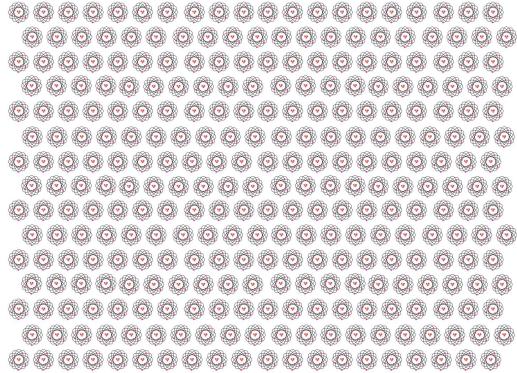
Algebraic compression of quantum circuits for Hamiltonian evolution, E. Kökcü, D. Camps, L. Bassman, J. K. Freericks, W. A. de Jong, R. Van Beeumen, A. F. Kemper, Phys. Rev. A 105, 032420, arXiv:2108.03282





Introduction to Quantum Computing

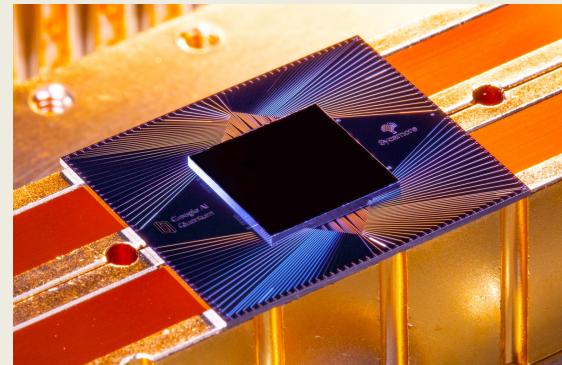
Dimension of a quantum state grows exponentially with the number of particles



A complete description of a typical quantum state of just 300 qubits requires more bits than the number of atoms in the visible universe
(figure from John Preskill).

$$2^{300} =$$

2037035976334486086268445688409378161051468393665936250636140449354381299763336706183397376



Google Sycamore chip (2019)
53 qubits

$2^{53} \approx 9 * 10^{15} \approx 36\text{PB}$ (single precision)

Quantum computing from 10000 ft

Two things are required for quantum computation:

- An **encoding** of the data in the quantum state $|\Psi\rangle$
- A way to **control** the evolution towards an encoding of the solution

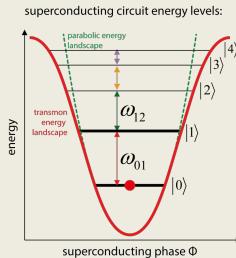
Quantum computers
are coherently
controllable quantum
systems



Advanced Quantum
Testbed @ Berkeley Lab

Qubits represent quantum data

Physics: two-level quantum system



Math: 2-dimensional complex vectors with unit norm

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$$

$$|\alpha|^2 + |\beta|^2 = 1$$

Quantum Gates: change state of a qubit

$$|\psi\rangle \xrightarrow{U} |\psi'\rangle \xrightarrow{\text{red arrow}} |\psi'\rangle = U|\psi\rangle$$

U is **unitary**

Unitary matrices preserve the norm of the vector (quantum operations are Hamiltonian time evolution)

Multi-qubit states and quantum circuits

$$|00\rangle = |0\rangle \otimes |0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \otimes \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

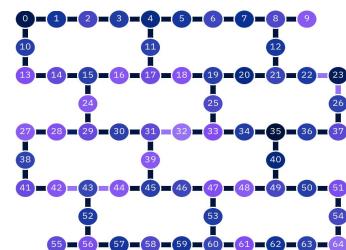
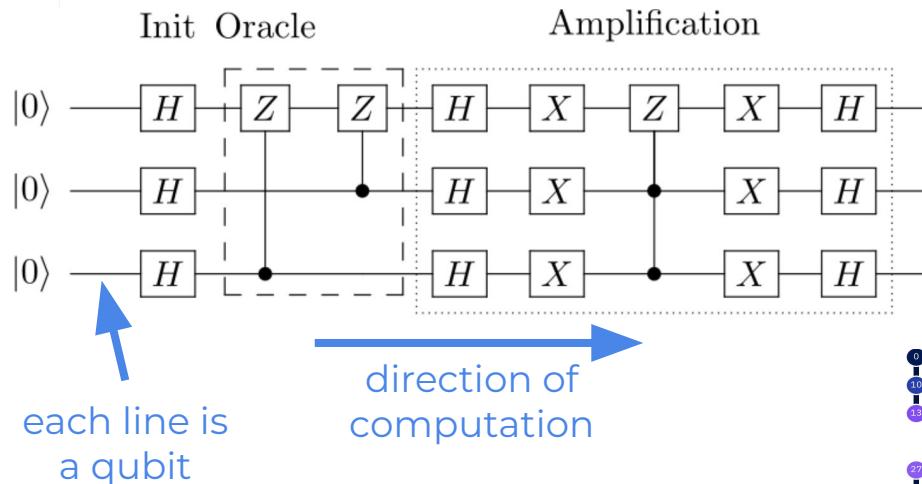
$$|01\rangle = |0\rangle \otimes |1\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$$

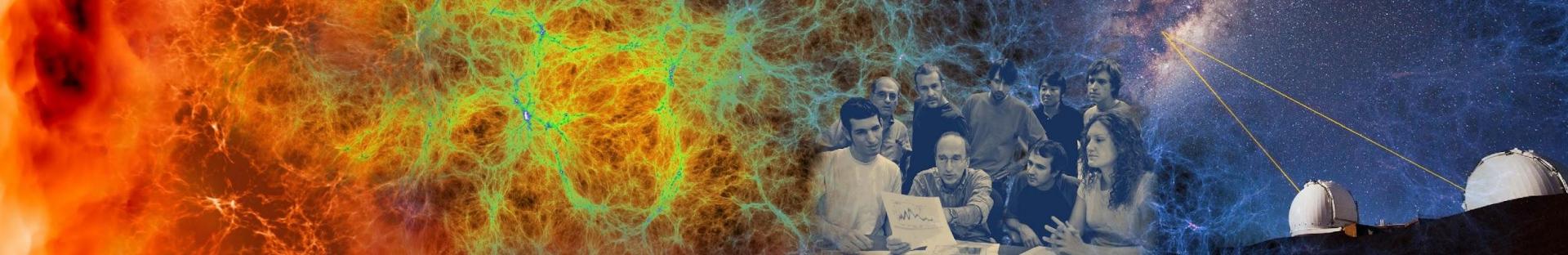
$$|10\rangle = |1\rangle \otimes |0\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \otimes \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \end{bmatrix}$$

$$|11\rangle = |1\rangle \otimes |1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$$

1 qubit \rightarrow 2 basis states
2 qubits \rightarrow 4 basis states
 \dots
n qubits \rightarrow 2^n basis states

$$|\psi\rangle = \alpha_0 |00\dots 0\rangle + \alpha_1 |00\dots 1\rangle + \dots + \alpha_{2^n-1} |11\dots 1\rangle = \begin{bmatrix} \alpha_0 \\ \alpha_1 \\ \vdots \\ \alpha_{2^n-1} \end{bmatrix} \in \mathbb{C}^{2^n}$$





Hamiltonian simulation and Trotterization

Hamiltonian simulation

Simulate time evolution under Schrödinger equation for a time-dependent Hamiltonian

$$\frac{\partial}{\partial t} \psi(t) = -iH(t)\psi(t) \quad H(t) \in \mathbb{C}^{2^N \times 2^N}$$

Hermitian matrix

Solved by applying the time-evolution operator:

$$U(t_1, t_0) = \mathcal{T} \exp \left(-i \int_{t_0}^{t_1} H(t) dt \right)$$

to the initial state:

$$\psi(t_1) = U(t_1, t_0)\psi(t_0)$$

Time-independent case: $U(t_1, t_0) = \exp(-i(t_1 - t_0)H)$

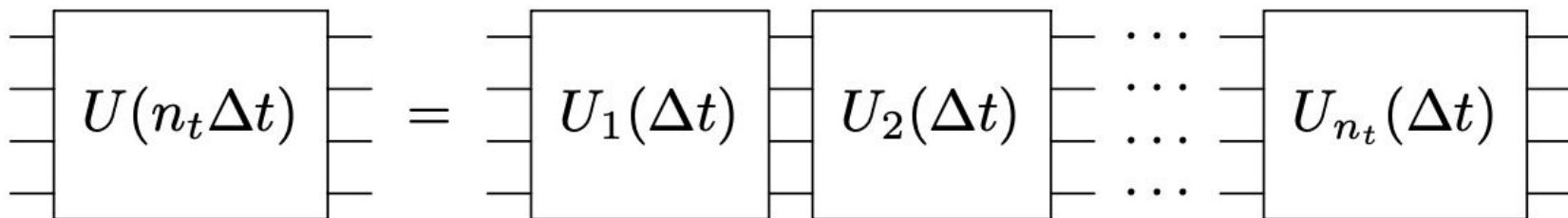


Trotter splitting and time discretization

Trotter decomposition (or operator splitting):

$$H = A + B \quad U(\Delta t) = \exp(-iA\Delta t) \exp(-iB\Delta t)$$

$$\|U(\Delta t) - \exp(-iH\Delta t)\| \leq \frac{\Delta t^2}{2} \| [A, B] \|$$



1D Spin- $\frac{1}{2}$ Hamiltonians

Pauli spin- $\frac{1}{2}$ matrices: $\sigma^x = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, \quad \sigma^y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix}, \quad \sigma^z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$

Basis for: $\mathfrak{su}(2)$ Generators for: $SU(2) = \left\{ \begin{bmatrix} \alpha & -\bar{\beta} \\ \beta & \bar{\alpha} \end{bmatrix} : \alpha, \beta \in \mathbb{C}, |\alpha|^2 + |\beta|^2 = 1 \right\}$

$$\sigma_i^\alpha := \underbrace{I \otimes \cdots \otimes I}_{i-1} \otimes \sigma^\alpha \otimes \underbrace{I \otimes \cdots \otimes I}_{N-i}$$

Transverse field XY model:

$$H(t) = \underbrace{\sum_{i=1}^{N-1} J_i^x(t) \sigma_i^x \sigma_{i+1}^x + J_i^y(t) \sigma_i^y \sigma_{i+1}^y}_{\text{Coupling}} + \underbrace{\sum_{i=1}^N h_i^z(t) \sigma_i^z}_{\text{External Field}}$$



Circuit diagrams

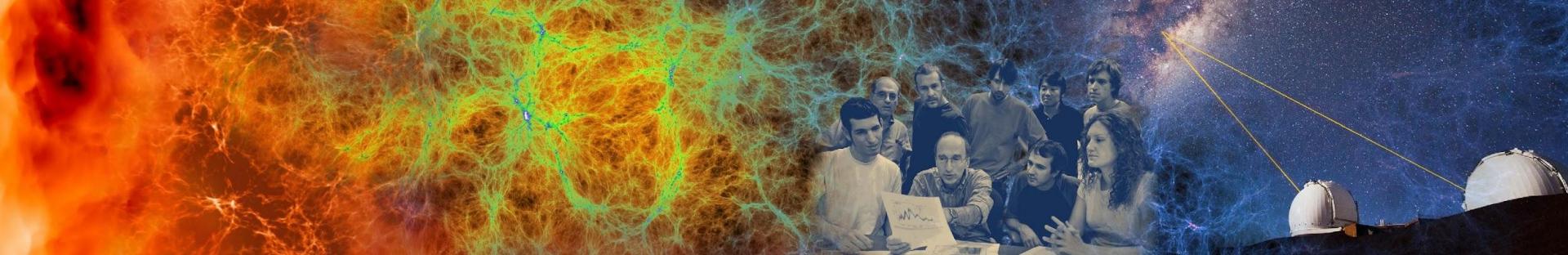
Single-qubit rotation over Pauli- α axis ($\alpha \in \{x, y, z\}$):

$$R^\alpha(\theta) := \exp(-i \sigma^\alpha \theta / 2) = \text{---} \boxed{\alpha} \text{---}$$

Two-qubit rotation over Pauli- α axis ($\alpha \in \{x, y, z\}$):

$$R^{\alpha\alpha}(\theta) := \exp(-i \sigma^\alpha \otimes \sigma^\alpha \theta / 2) = \text{---} \boxed{\alpha} \text{---}$$

Easy operations to execute on QC: Native gate for ion traps, 2 CNOTs for superconducting



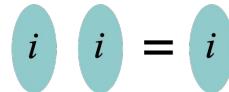
Algebraic compression of Hamiltonian simulation circuits

Definition

A **block** is a parametrized and indexed family of operators $\mathbf{B}_i(\theta)$ that satisfy 3 properties:

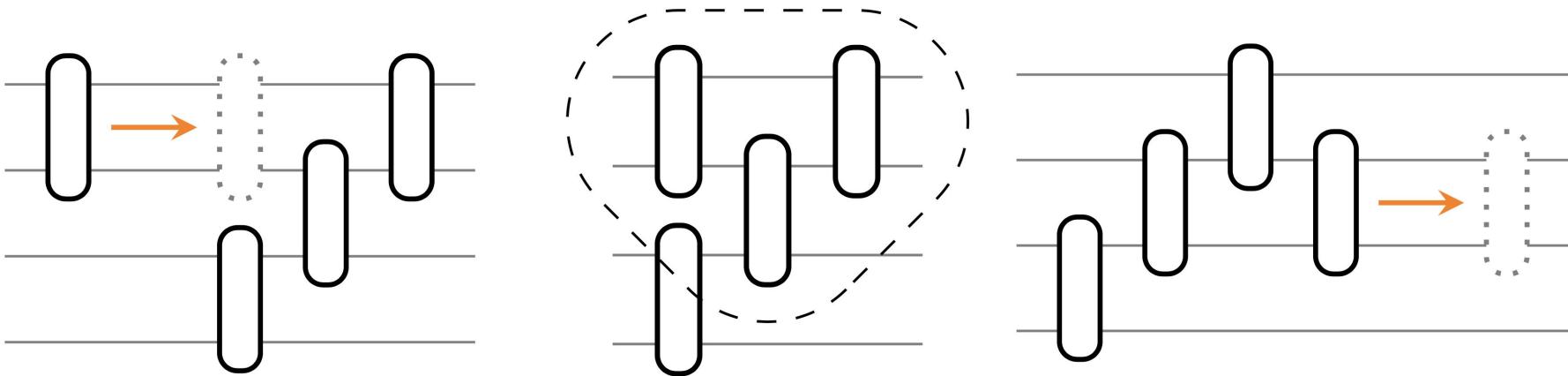
- **Fusion:**

$$\mathbf{B}_i(\theta_1)\mathbf{B}_i(\theta_2) = \mathbf{B}_i(\hat{\theta})$$





Central mechanism in our compression algorithm

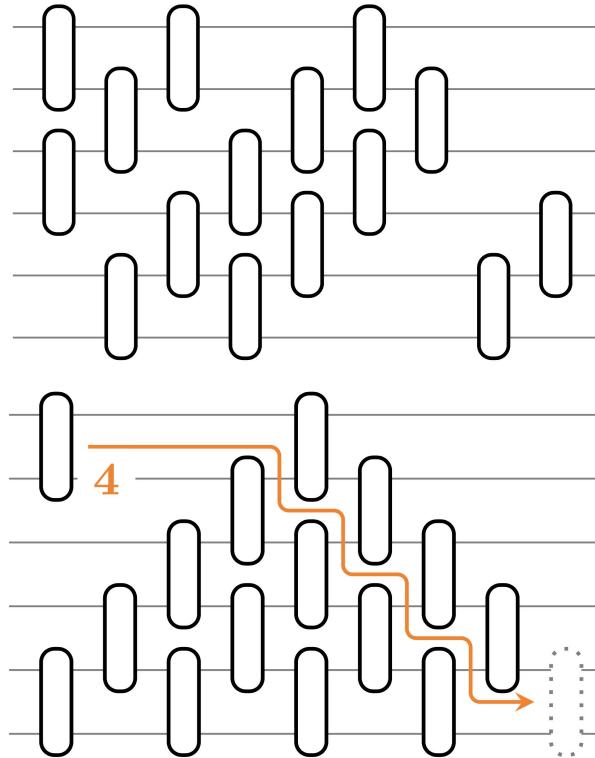
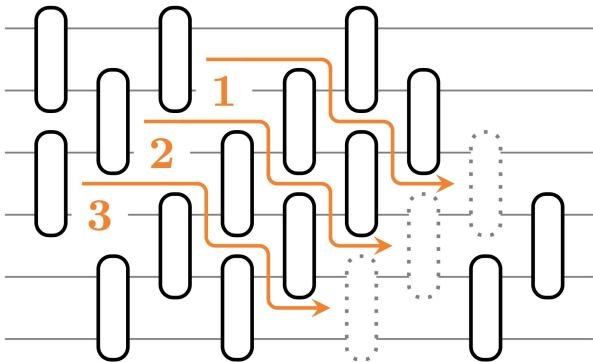
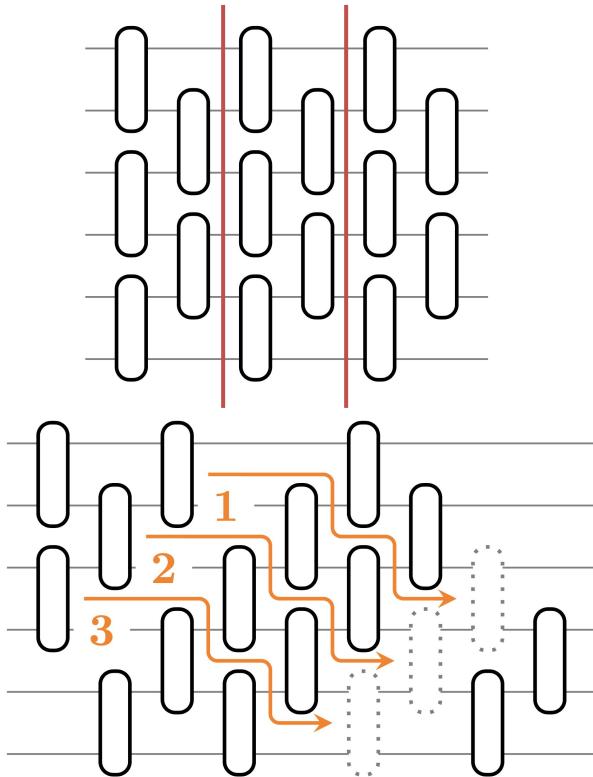


Equivalent mechanism as in *core-chasing* eigenvalue algorithms, but on operators of exponential dimension.

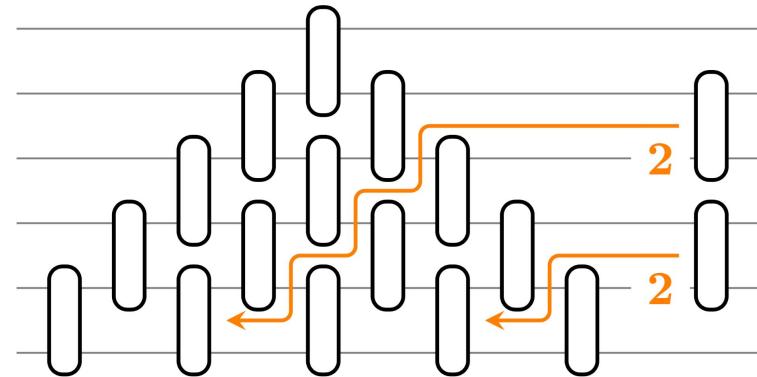
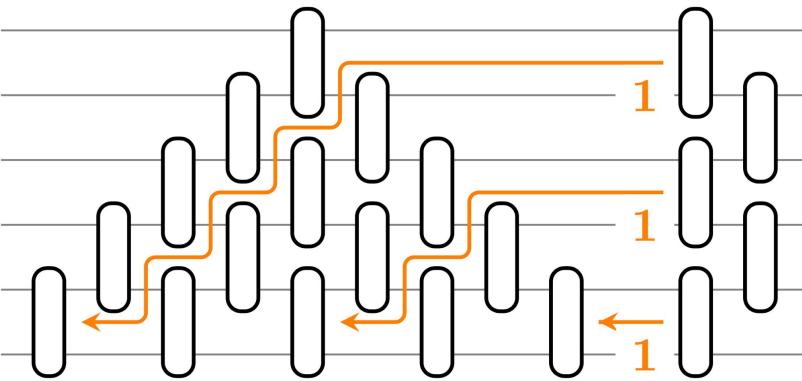
Core-Chasing Algorithms for the Eigenvalue Problem, Aurentz, Mach, Robol, Vandebril, Watkins



Transforming squares to triangles



Merging time-steps into triangles



Even/odd blocks act on independent parts of the triangle. In our implementation, these are merged in parallel

Euler decomposition and turnover of SU(2)

Lemma: Euler decomposition

Let $\alpha, \beta \in \{x, y, z\}$, $\alpha \neq \beta$. Every $U \in \text{SU}(2)$ can be represented as:

$$U = R^\alpha(\theta_1) R^\beta(\theta_2) R^\alpha(\theta_3), \quad \text{---} \boxed{U} \text{ ---} = \text{---} \boxed{\alpha}_{\theta_3} \text{ ---} \boxed{\beta}_{\theta_2} \text{ ---} \boxed{\alpha}_{\theta_1} \text{ ---}$$

Lemma: SU(2) turnover

Let $\alpha, \beta \in \{x, y, z\}$, $\alpha \neq \beta$. For every $\theta_1, \theta_2, \theta_3$, there exist $\theta_a, \theta_b, \theta_c$ such that

$$R^\alpha(\theta_1) R^\beta(\theta_2) R^\alpha(\theta_3) = R^\beta(\theta_a) R^\alpha(\theta_b) R^\beta(\theta_c), \quad \text{---} \boxed{\alpha}_{\theta_3} \text{ ---} \boxed{\beta}_{\theta_2} \text{ ---} \boxed{\alpha}_{\theta_1} \text{ ---} = \text{---} \boxed{\beta}_{\theta_c} \text{ ---} \boxed{\alpha}_{\theta_b} \text{ ---} \boxed{\beta}_{\theta_a} \text{ ---}$$

- We can compute the SU(2) turnover backward stable (Givens rotations)

SU(2) groups in disguise

Lemma:

Let $\alpha, \beta \in \{x, y, z\}$, $\alpha \neq \beta$. The following operations are also dual Euler decompositions of SU(2):

$$\begin{array}{c} i \\ \alpha \\ i+1 \\ \beta \\ i+2 \\ \theta_3 \\ \theta_2 \\ \theta_1 \end{array} = \begin{array}{c} \alpha \\ \beta \\ \theta_c \\ \theta_b \\ \theta_a \end{array}$$

$$\begin{array}{c} i \\ \alpha \\ i+1 \\ \beta \\ \alpha \end{array} = \begin{array}{c} \beta \\ \alpha \\ \beta \end{array}$$

$$\begin{array}{c} i \\ \alpha \\ i+1 \\ \beta \\ \alpha \end{array} = \begin{array}{c} \beta \\ \alpha \\ \beta \end{array}$$

Fusion operations are trivial: i $\begin{array}{c} \alpha \\ \theta_1 \\ \alpha \\ \theta_2 \end{array}$ $=$ $\begin{array}{c} \alpha \\ \theta_1 + \theta_2 \end{array}$

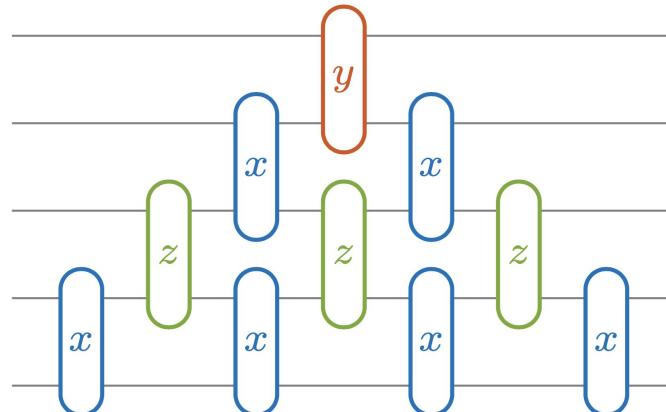
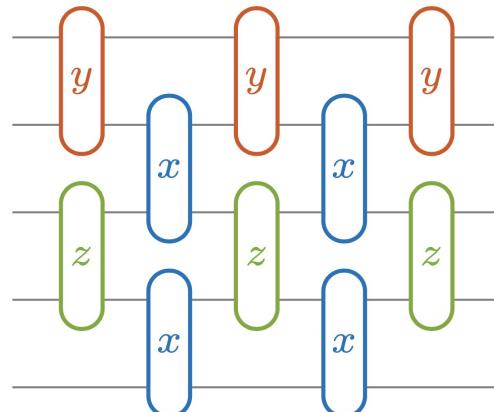
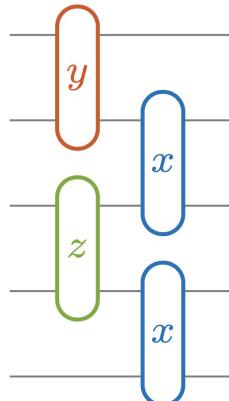
$$\begin{array}{c} i \\ \alpha \\ i+1 \\ \alpha \\ \theta_1 \\ \theta_2 \end{array} = \begin{array}{c} \alpha \\ \theta_1 + \theta_2 \end{array}$$

Kitaev Chain

A Kitaev chain is a Hamiltonian of the form:

$$H(t) = \sum_{i=1}^{N-1} J_i^{\alpha_i}(t) \sigma_i^{\alpha_i} \sigma_{i+1}^{\alpha_i} \quad \alpha_i \neq \alpha_{i+1}$$

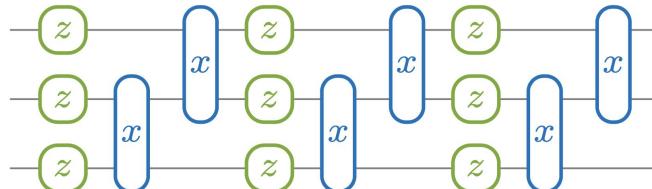
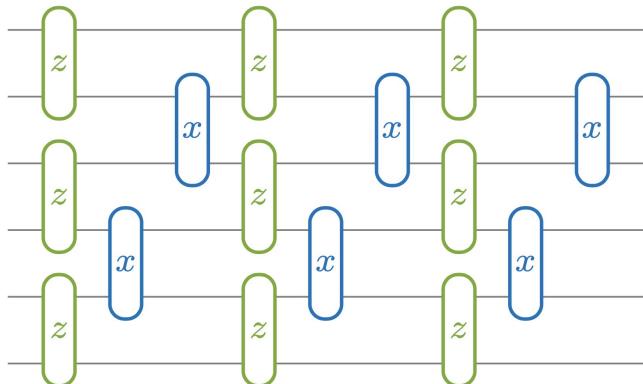
For example: $H(t) = J_1^y(t) \sigma_1^y \sigma_2^y + J_2^x(t) \sigma_2^x \sigma_3^x + J_3^z(t) \sigma_3^z \sigma_4^z + J_4^x(t) \sigma_4^x \sigma_5^x$



TFIM Hamiltonian

The Transverse-Field Ising Model has the form:

$$H(t) = \sum_{i=1}^{N-1} J_i^\alpha(t) \sigma_i^\alpha \sigma_{i+1}^\alpha + \sum_{i=1}^N h_i^\beta(t) \sigma_i^\beta \quad \alpha \neq \beta$$



N-qubit TFIM is isomorphic to 2N-qubit Kitaev chain

TFXY Hamiltonian

Transverse field XY model:

$$H(t) = \underbrace{\sum_{i=1}^{N-1} J_i^x(t) \sigma_i^x \sigma_{i+1}^x + J_i^y(t) \sigma_i^y \sigma_{i+1}^y}_{\text{Coupling}} + \underbrace{\sum_{i=1}^N h_i^z(t) \sigma_i^z}_{\text{External Field}}$$

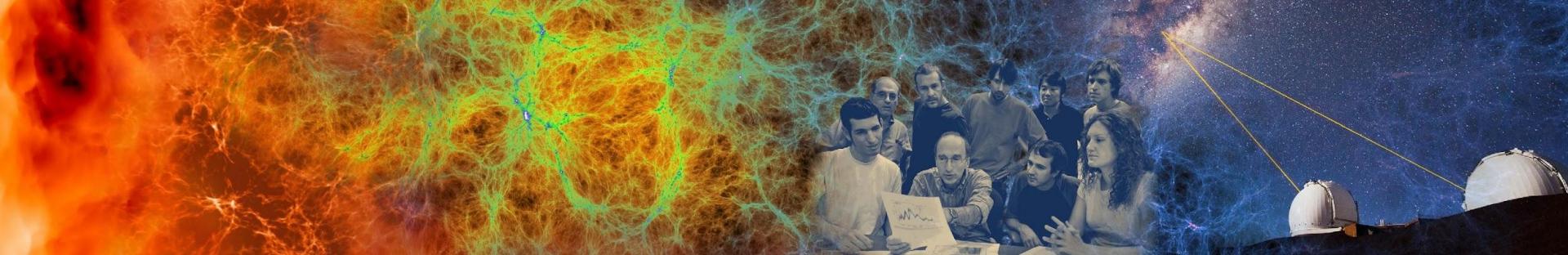
TFXY block:

The diagram shows a block with six parameters: $\theta_1, \theta_3, \theta_5$ (z-axis) and $\theta_2, \theta_4, \theta_6$ (y-axis). The block is represented by a matrix:

$$\begin{bmatrix} \alpha & & -\bar{\delta} \\ & \beta & -\bar{\gamma} \\ \delta & \gamma & \bar{\beta} \\ & & \bar{\alpha} \end{bmatrix}$$

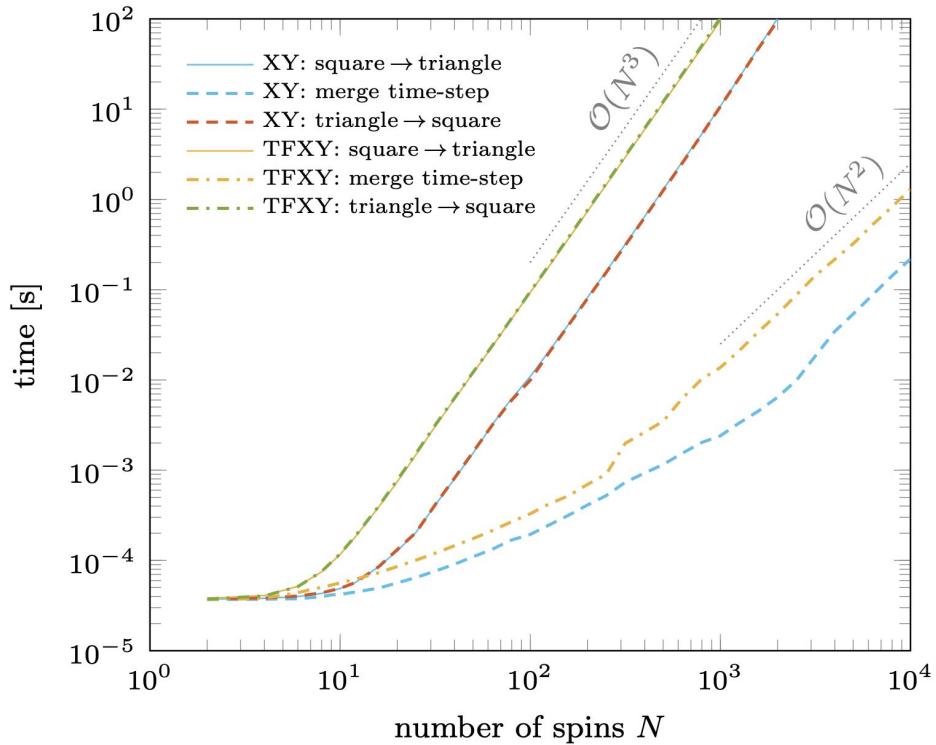
Turnover through simultaneous diagonalization

$$\begin{array}{cccccc} A, B & C, D & E, F & U, V & W, X & Y, Z \\ \bullet \quad \blacksquare \quad \bullet \quad \bullet \quad \blacksquare \quad \bullet & \bullet \bullet \quad \blacksquare \quad \bullet \quad \bullet \quad \bullet & \bullet \quad \blacksquare \quad \bullet \quad \bullet \quad \bullet \quad \bullet & \bullet \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet & \bullet \quad \blacksquare \quad \bullet \quad \bullet \quad \bullet \quad \bullet & \bullet \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \end{array} = \begin{array}{cccccc} \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \\ \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \\ \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \\ \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \\ \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \\ \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \quad \bullet \end{array}$$

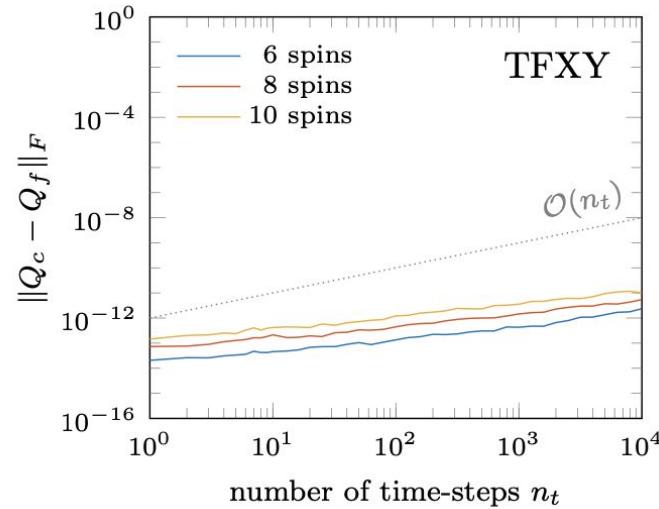
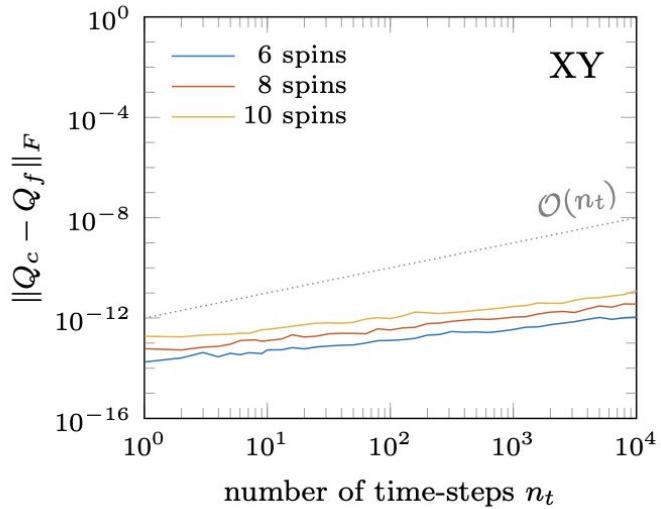


Results

Numerical results: timings



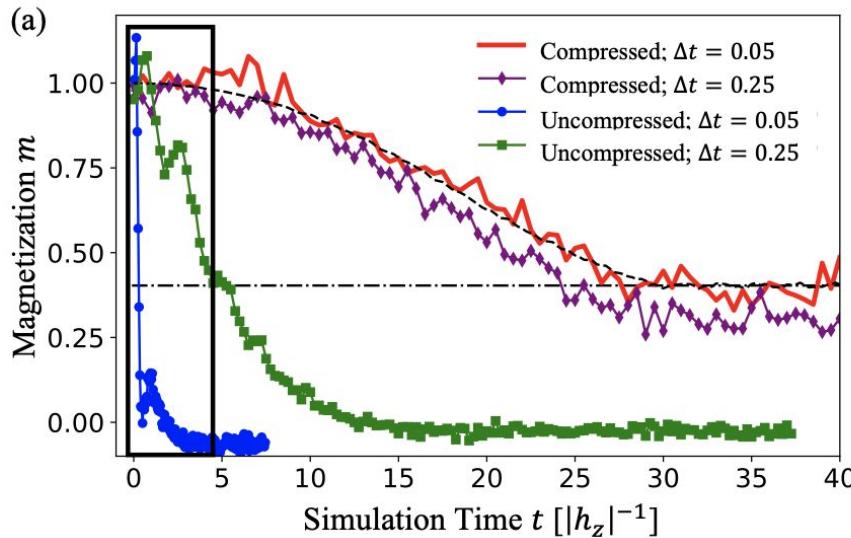
Numerical results: backward error



Quantum Computer: Adiabatic State Preparation

- Time evolve TFIM in ground state from trivial state w/o coupling to more complicated ground state with coupling terms
- 5 qubit model on IBMQ Brooklyn
- Measure the average magnetization

$$H(t) = J(t) \sum_i \sigma_i^x \sigma_{i+1}^x + h \sum_i \sigma_i^z$$



Conclusion

- Efficient and stable **classical** numerical algorithm for compression of quantum circuits for simulation of **integrable TXY chains**
- **Enables simulation** of small systems on current generation noisy quantum hardware
 - Prepare **non-trivial states**
 - Simulate interesting physics phenomena
- Extensions to 2D non-interacting, controlled evolutions, ...

Fast Free Fermion Compiler (F3C):

<https://github.com/QuantumComputingLab>

arXiv:2108.03282, arXiv:2108.03283

