

## CSE312 HOMEWORK 3 REPORT

### 1) Problem Solution

#### A) SPIMOS\_GTU\_X.S PART

In this assignment, I firstly defined the parameters that should be in my process table in asm. These parameters are as follows.

```
processID: .word 0 1 2 3 4
currentPC: .word 0 0 0 0 0
endPC: .word 0 0 0 0 0
stackPointerAddress: .word 0 0 0 0 0
parentProcess: .word 0 0 0 0 0
```

I created parameters that keep the address of this data in syscall.cpp file so that these parameters can be updated and written during interrupt. So I can update the data in the process table with the set\_mem\_word function. Then, to be able to handle an interrupt inside the asm, the spimtimer\_handler must go to that part of the asm whenever there is an interrupt. To do this, I called a syscall, pointing to the interrupt\_handler label. So when an interrupt occurs, spimtimer\_handler uses the parameter from syscall to go to this label. And asm handles interrupt. After all these processes are finished, we come to the call processes section. First of all, since each process will have its own register values, I created a separate register parameter for each (as follows).

```
register0:
|      | .align 2
|      | .space 128
register1:
|      | .align 2
|      | .space 128
register2:
|      | .align 2
|      | .space 128
register3:
|      | .align 2
|      | .space 128
register4:
|      | .align 2
|      | .space 128
```

For example, since SPIMOS\_GTU\_1.s will run a total of 4 (+1 init processes) processes, I kept 5 register parameters. Thus, the programs can continue where they left off during the backup and update stages. In the process loading section, I first upload the address of the register parameter to be used in the \$ a0 register. then I load the currentPC and endPC parameters into the \$ 15 and \$ 24 registers. These will keep the start and end addresses of the asm file that will run. And with the help of execve, I load the asm file. I repeat these processes for all processes. And the program is waiting in an endless loop until it is interrupt.

Then, in the interrupt\_handler label, the register information of the process that was previously running in the \$ s6 register is displayed. In the \$ s7 register, the register information that will now be kept is kept. In this label, the registration information of the previously running process is recorded and the information of the new process to be loaded is loaded one by one.

**NOTE:** These registers are specially selected. If other asm files use these registers, the program may run incorrectly. In addition, WAITPID syscall is used in order not to interrupt the processes performed by kernels.

## **B) SPIM\_timerHandler()**

In this function, I use the parameter addresses sent by kernel to update the process table information. I use syscalls to get this information. I describe my Syscalls in detail in Syscall.cpp. Coming to SPIM\_timerHandler function, I first check the waitpid\_signal parameter. This parameter states that asm files are at a critical point while running. If the value of this parameter is true, the program will resume even if interrupt occurs. In other words, context switch will not be realized in any way. If false, the other process should work. To do this, I first check whether all processes are finished. If all processes are finished, I am free all the parameters that I reserved in memory with malloc. And I put the PC at the address of the kernel's exit label. Thus, the program exits safely. If all processes are not finished, I continue my process. As I mentioned while explaining the kernel part, I update the process table information in asm to be able to backup the running process. I do this with the help of currentProcess. currentProcess keeps track of which process is running. I put the PC and register information of the running process in the \$ 26 and \$ 22 registers. Then, if this process is not finished, I switch from "Running" to "Ready". Then I use the Round robin scheduling algorithm to run the other process. And by looking at the state of the processes, I select the process that will be running next by looking at which ones are ready. And I assign to \$ 23 and \$ 27 registers to update the register and PC of the process that will run in kernel. And I assign the InterruptHandlerPC to the PC so that the kernel can handle the interrupt. Thus, by going to the kernel interrupt\_handler label, it can update the information by using the registers it receives from spim\_timerhandler.

The syscalls I use are as follows:

```
#define FORK 18
#define EXECVE 19
#define WAITPID 20
#define EXECVE_2 21

#define _TERMINATE_PROCESS 30

#define _RANDOM_NUMBER_GENERATOR 42
#define _RANDOM_NUMBER_GENERATOR_2 43

#define SET_PROCESSID_ADR 54
#define SET_ENDPC_ADR 55
#define SET_SP_ADR 56
#define SET_PARENTPROCESS_ADR 57
#define SET_PC_ADDRESS 58
#define SET_INTERRUPT_HANDLER_PC 59

#define SET_PC 61
#define SET_EXIT_PC 62
#define APPEND_PROCESS_NAME 63
```

## 2) Results

Small examples of printouts are as follows

**A) SPIMOS\_GTU\_1.s**

[illegible]

```
Process Name : Collatz.asm
Process ID : 3
Process state : Running
Current PC : 4195900
End PC : 4196052
Stack Pointer Address : 2147475296
Parent Process : 2
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
```

```
20 10 5 16 8 4 2 1
```

```
24 : 24 12 6 3 10 5 16 8 4 2 1
```

```
25 : 25 76 38 19 58 29 88 44 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
```

```
Process Name : init
Process ID : 0
Process state : Running
Current PC : 4194748
End PC : 4195156
Stack Pointer Address : 2147475296
Parent Process : -1
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
```

```
All processes finished
```

## B) SPIMOS\_GTU\_2.s

```
Process Name : init
Process ID : 0
Process state : Running
Current PC : 4194836
End PC : 4195072
Stack Pointer Address : 2147475296
Parent Process : -1
```

```
Process Name : BinarySearch.asm
Process ID : 5
Process state : Running
Current PC : 4196656
End PC : 4196772
Stack Pointer Address : 2147475296
Parent Process : 4
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
is in the list.
```

```
The index of value is : 3
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
```

```
Process Name : init
Process ID : 0
Process state : Running
Current PC : 4194836
End PC : 4195072
Stack Pointer Address : 2147475296
Parent Process : -1
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
All processes finished
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
```

```
46
```

```
INTERRUPT OCCURED!
```

```
-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-
```

### C) SPIMOS\_GTU\_3.s

[illegible]

### D) BinarySearch.asm

```
46 is in the list.  
The index of value is : 3
```

### E) LinearSearch.asm

```
28 is in the list.  
The index of value is : 2
```



## F) Collatz.asm

```
1 : 1
2 : 2 1
3 : 3 10 5 16 8 4 2 1
4 : 4 2 1
5 : 5 16 8 4 2 1
6 : 6 3 10 5 16 8 4 2 1
7 : 7 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
8 : 8 4 2 1
9 : 9 28 14 7 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
10 : 10 5 16 8 4 2 1
11 : 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
12 : 12 6 3 10 5 16 8 4 2 1
13 : 13 40 20 10 5 16 8 4 2 1
14 : 14 7 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
15 : 15 46 23 70 35 106 53 160 80 40 20 10 5 16 8 4 2 1
16 : 16 8 4 2 1
17 : 17 52 26 13 40 20 10 5 16 8 4 2 1
18 : 18 9 28 14 7 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
19 : 19 58 29 88 44 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
20 : 20 10 5 16 8 4 2 1
21 : 21 64 32 16 8 4 2 1
22 : 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
23 : 23 70 35 106 53 160 80 40 20 10 5 16 8 4 2 1
24 : 24 12 6 3 10 5 16 8 4 2 1
25 : 25 76 38 19 58 29 88 44 22 11 34 17 52 26 13 40 20 10 5 16 8 4 2 1
```

## G) Palindrome.asm

1: aba: Palindrome	33: deafening: Not Palindrome	66: cats: Not Palindrome
2: azbza: Palindrome	34: queue: Not Palindrome	67: substantial: Not Palindrome
3: ada: Palindrome	35: shoe: Not Palindrome	68: adhesive: Not Palindrome
4: dontnod: Palindrome	36: stove: Not Palindrome	69: magic: Not Palindrome
5: tacocat: Palindrome	37: uttermost: Not Palindrome	70: introduce: Not Palindrome
6: madam: Palindrome	38: harass: Not Palindrome	71: paint: Not Palindrome
7: kayak: Palindrome	39: deserve: Not Palindrome	72: stain: Not Palindrome
8: refer: Palindrome	40: wise: Not Palindrome	73: damage: Not Palindrome
9: toyot: Palindrome	41: direction: Not Palindrome	74: curve: Not Palindrome
10: bob: Palindrome	42: underwear: Not Palindrome	75: nonstop: Not Palindrome
11: scrub: Not Palindrome	43: card: Not Palindrome	76: fog: Not Palindrome
12: sack: Not Palindrome	44: unbecoming: Not Palindrome	77: known: Not Palindrome
13: itchy: Not Palindrome	45: key: Not Palindrome	78: beg: Not Palindrome
14: check: Not Palindrome	46: curtain: Not Palindrome	79: expansion: Not Palindrome
15: noxious: Not Palindrome	47: war: Not Palindrome	80: magnificent: Not Palindrome
16: tie: Not Palindrome	48: stomach: Not Palindrome	81: shoes: Not Palindrome
17: phobic: Not Palindrome	49: bait: Not Palindrome	82: bloody: Not Palindrome
18: hungry: Not Palindrome	50: loutish: Not Palindrome	83: decision: Not Palindrome
19: songs: Not Palindrome	51: automatic: Not Palindrome	84: robust: Not Palindrome
20: mammoth: Not Palindrome	52: start: Not Palindrome	85: bead: Not Palindrome
21: science: Not Palindrome	53: fallacious: Not Palindrome	86: quilt: Not Palindrome
22: amount: Not Palindrome	54: song: Not Palindrome	87: satisfying: Not Palindrome
23: middle: Not Palindrome	55: float: Not Palindrome	88: young: Not Palindrome
24: tiger: Not Palindrome	56: seemly: Not Palindrome	89: glass: Not Palindrome
25: burn: Not Palindrome	57: yarn: Not Palindrome	90: rescue: Not Palindrome
26: dispensable: Not Palindrome	58: remain: Not Palindrome	91: poised: Not Palindrome
27: beef: Not Palindrome	59: guitar: Not Palindrome	92: glamorous: Not Palindrome
28: dime: Not Palindrome	60: quirky: Not Palindrome	93: heal: Not Palindrome
29: subdued: Not Palindrome	61: odd: Not Palindrome	94: tested: Not Palindrome
30: program: Not Palindrome	62: condition: Not Palindrome	95: touch: Not Palindrome
31: prick: Not Palindrome	63: shape: Not Palindrome	96: wobble: Not Palindrome
32: outstanding: Not Palindrome	64: dizzy: Not Palindrome	97: request: Not Palindrome
33: deafening: Not Palindrome	65: degree: Not Palindrome	98: chop: Not Palindrome
	66: cats: Not Palindrome	99: contain: Not Palindrome
	67: substantial: Not Palindrome	100: stroke: Not Palindrome
	68: adhesive: Not Palindrome	

```
100: stroke: Not Palindrome
Do you want to continue (y/n)?
y

Please enter the last word:
ilhan
101: ilhan: Not Palindrome

Goodbye...
```

### 3) ISSUES

A) In some cases, the values of Collatz write unexpectedly different. I think this is due to a break in a critical place.

**B)** In some cases, Palindrome may print the words it reads incompletely. Palindrome may work incorrectly because it runs too many in SPIMOS\_GTU\_2 and 3.

**C)** The user should never enter extended ascii as input. Input must be valid. Error control is not done.