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1. S = switch

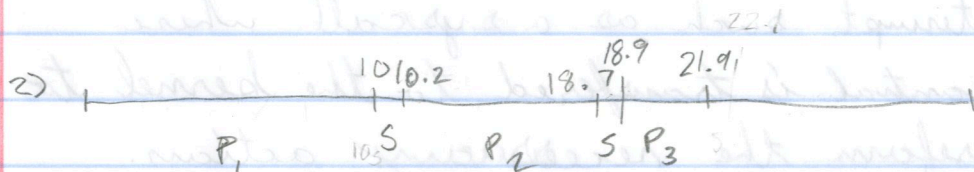
Time in sec

0 2.2 4.2 4.4 4.6 6.6 8.6 10.8 12.2 14.2 16.4 18.6 20.8 23.2

P₁ S P₂ S P₃ S P₁ S P₂ S P₃ S P₁ S P₂ S P₁ S P₂ S P₁ S P₂ S

P₃ ends @ 12 P₁ ends @ 23 P₂ ends @ 23.7

$2.2 / 23.7 = 9.283\%$ of time spent on context switches



Batch multiprogrammed finished execution fastest
if a CPU had multiple cores it could actually
execute multiple processes simultaneously or
in the case where one of the processes has an I/O call causing the

3 The 4 exceptions are trap, fault, (Hardware) process
interrupt, Abort. Trap is an intentional to wait
interruption from something like a software and sit
interrupt an example is a syscall. the return idle in
behavior returns to the next instruction where a batch
the call was made. A fault is a potentially run in
recoverable error, for example division a time
by zero, invalid op code or a segmentation fault, slice it
if the error is recoverable ~~the~~ it will return would
control to the user. The interrupt is a signal switch
from an I/O device an example of would be to another
input from a keyboard, or disk read finished. process
It always returns to the next instruction, asynchronous

An abort is a non recoverable error an example would be a hardware bus failure. The program or operation that was being executed gets killed.

A software interrupt is when an application or a program sends an interrupt such as a syscall where control is transferred to the kernel to perform the necessary actions.

This differs from a hardware interrupt where the interrupt is issued by a hardware device like a disk, or keyboard.

Each device has a unique Interrupt Request line. This line. Based on the IRQ the CPU will dispatch the appropriate hardware driver.

4. The jump table is designed to specify by number how each syscall should be handled. For each numbered syscall the table has associated functions and program files. The jump table is also called the Trap table.

5 It enhances performance by allowing applications continue processing while allowing I/O operations to still run.

Asynchronous - the process sends a request to the I/O subroutine for a read then returns to the process and continues to execute then the I/O subroutine will send an interrupt when the desired request has been fulfilled. This is an improvement because the CPU doesn't have to wait for the I/O request to be fulfilled before it continues to execute processes.

Non blocking - the process sends an request to the I/O and returns immediately with a value indicating how many bytes were transferred the process then loop back to gather all data required for the process. This is an improvement because it frees up the Processor to continue execute processes while it rechecks if the full I/O request

