Andrew Candelaresi Algorithms homework 3

On my honor, as a University of Colorado at Boulder student, I have neither given nore reveived unauthorized assistance on this work.

```
1. a. T(n) = 4T(n/2) + n^2

a=4 b=2 k=1 p=0 log_2(4)=2

4<2 a>b^k

T(n) = Theta(n^2)
```

b. since the recurrence relation comes out to:

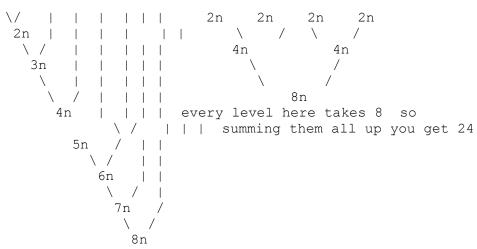
 $T(n) = 2^kT(n-k)+c$  and you get down to level k (n-k)=1 all we are left with is  $2^k+c$  and since c would be dominated by the  $2^k$  at this level k = (n-1) we have  $2^n-1$  and we can get  $2^n-1$  which is  $(1/2)2^n$  and the constant drops off the run time is  $2^n$  T(n)=0  $(2^n)$ 

c.  $T(n) = 9T(n/3)+n^2$   $a = 9 b = 3 k=2 p=0 log_3(9)=2$   $9=3^2 a=b^k and p > -1$  $T(n) = Theta[(n^2)log(n)]$ 

I would choose Algorithm a

2. best case run time for a comparison sorting algorithm would be that the list is already sorted and in one go of comparisons the list is returned

```
InsertionSort(A) {
                                                  cost times
for j in range (2, A.length)
                                            с1
      key = A[j]
                                            c2
                                                    n-1
      i = j-1
                                           С4
                                                        n-1
                                           с5
                                                        \sum from (2, n) t_j
      while i > 0 and A[i] > key
                                                        \Sigma from (2, n) (t_j-1)
\Sigma from (2, n) (t_j-1)
            A[i+1] = A[i]
                                            С6
            i = i-1
                                            c7
      A[i+1] = key
                                            с8
                                                        n-1
      T(n) = c1*n + c2(n-1) + c4(n-1) + c5\sum_{i=1}^{n} from (j=2, n) t j +
      c6\Sigma from (j=2, n) (t j-1) +c7\Sigma from (j=2, n) (t j-1) + c8(n-1)
      T(n) = c1*n +c2(n-1) +c4(n-1) +c5(n-1) +c8(n-1)
             = (c1+c2+c4+c5+c8) n - (c2+c4+c5+c8)
      we can express this running time as a funtion an+b for some
      constants a and b that depend on the costs
      c i; it is therefore a linear funcion of n
```



this requires a summation of all merges 8+7+6+5+4+3+2+1=36

## 4.

TSLA

date range: 2010/06/29 - 2015/09/22

buy on: 2010/07/08
sell on: 2014/09/03
Maximum Profit: 271.53

PG

date range: 2013/09/23 - 2015/09/22

buy on: 2013/10/01
sell on: 2013/10/30
Maximum Profit: 7.01

ΚO

date range: 2014/09/22 - 2015/09/22

buy on: 2015/03/16
sell on: 2015/08/06
Maximum Profit: 2.0