

Famous Problems in Computer Science

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About Tolga Ovatman

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Research Interests

- Software modeling and analysis. Member of [SMART](#).
- Software performance and cloud computing.
- Formal methods and model checking.
- Software project management and agile methodologies.

Searching



Predecessor/Successor
Problem



Challenge

*Problems of practical importance usually
contains too many possibilities.*

15	7	3	25	20	18	13	11	6	17
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Can we provide a strategy for guess the number game? [1-N]

1. Guess 1
2. *Increase until the
right guess*

1. Make a random guess
2. *Get feedback*
3. *Make another random
guess*

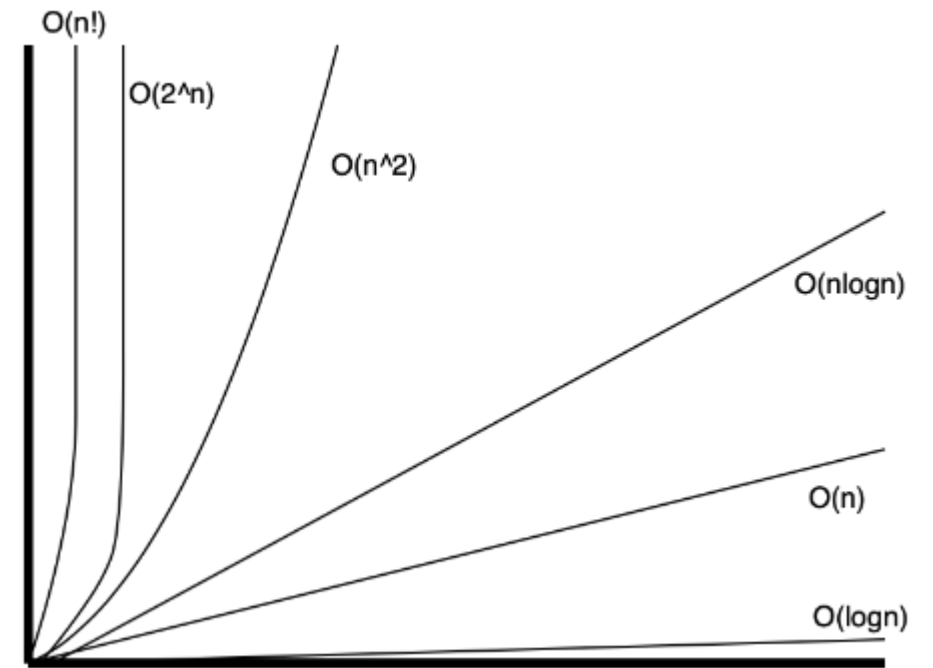
1. Guess $N/2$
2. *Get feedback*
 1. *If lower, guess $[0 - N/2]$*
 2. *If higher, guess $[N/2 - N]$*

Algorithm Performance

15	7	3	25	20	18	13	11	6	17
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How many steps
do I need to take
for input of size **N**



3	6	7	11	13	15	17	18	20	25
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Sorting



Why important?



It makes lookup or search efficient

It makes merging of sequences efficient.

It enables processing of data in a defined order.

Solution

- It is a multi-dimensional problem
- *For an n elements it takes around $n \log n$ steps.*

Sorting



BOGO SORT

```
while not is_sorted(list):  
    shuffle(list)
```

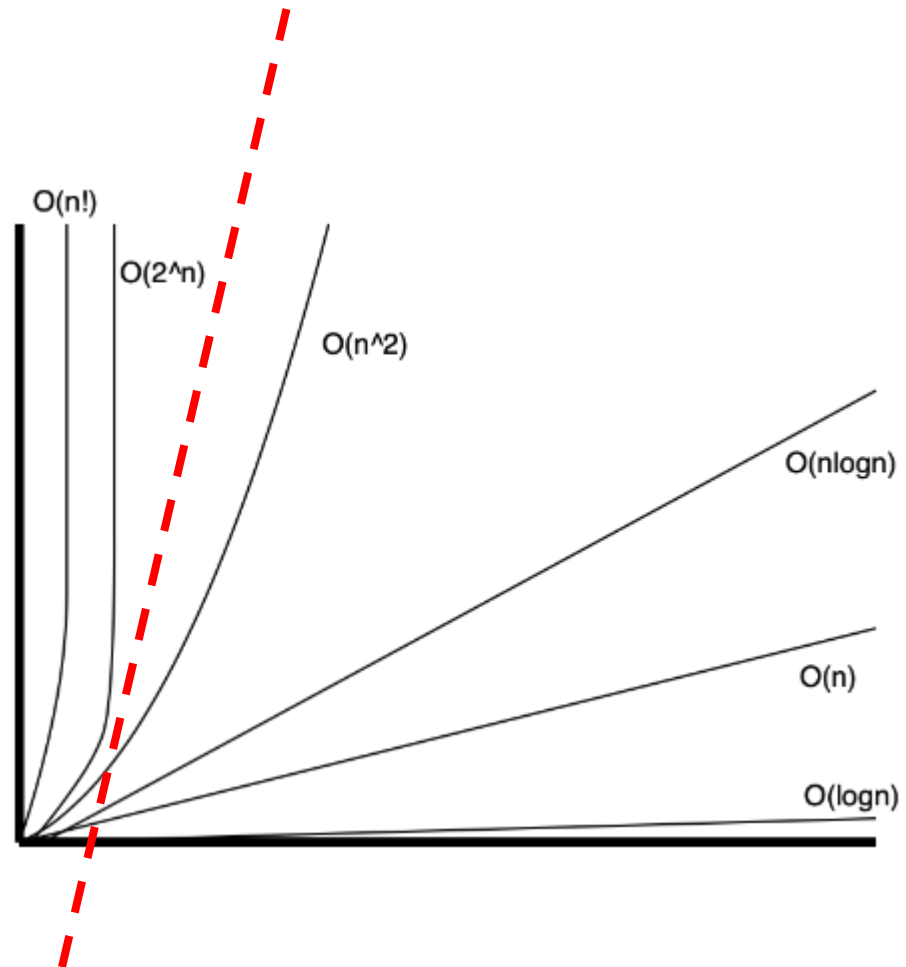
STALIN SORT

```
for a in list:  
    if out_of_order(list, a):  
        list.remove(a)
```

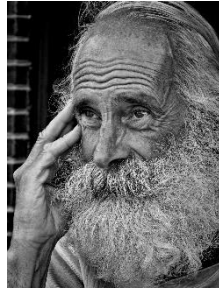
SLEEP SORT

```
parallel for a in list:  
    sleep(a)  
    print(a)
```

Algorithm Performance

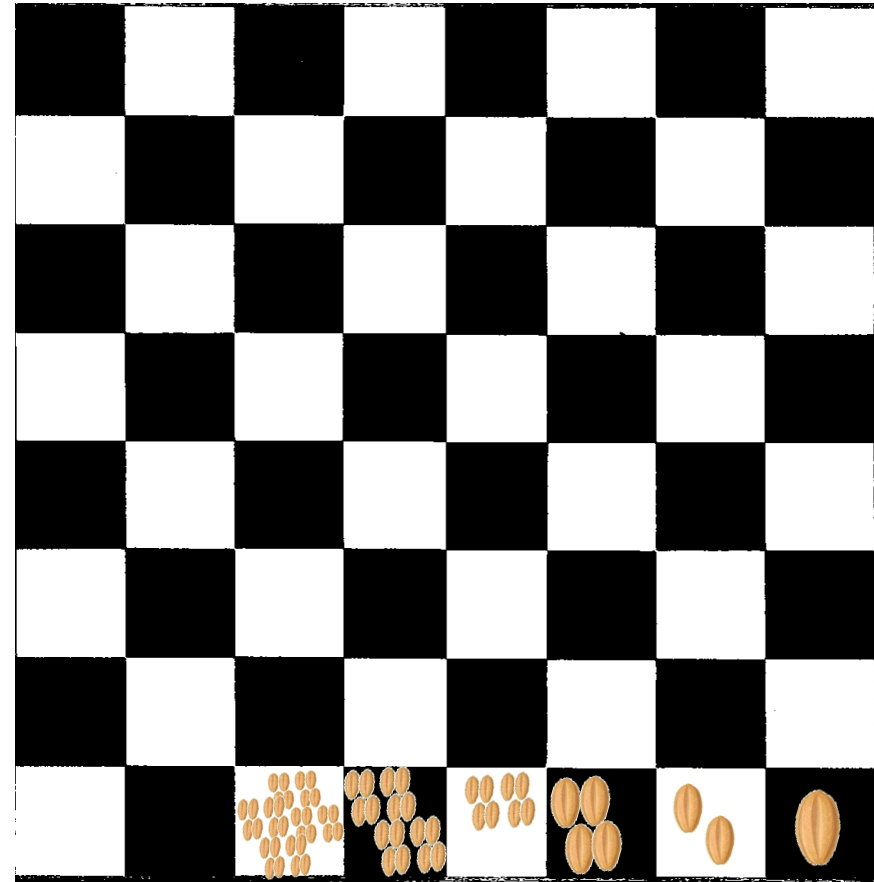


Exponential Growth – Story of chessman

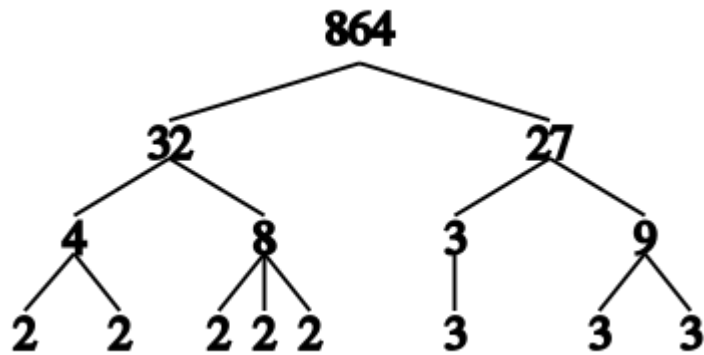


How much?

- *Approx. 2×10 wheat grains weigh a gram.*
- *$\sum_{i=0}^{63} 2^i \cong 2^{64} \cong 2 \times 10^{19}$ grains almost 10^{12} tons*
- *Earth produced much less than 10^9 tons a year for the last 20 years.*
- *Chess is less than 2×10^3 years old.*



Prime Decomposition



$$35 = 7 \times 5$$

$$437 = 19 \times 23$$

Why important?



It is the fundamental technique that lies behind modern cryptography

Primality Test?



- *Thought to be as hard as Prime decomposition*
- *In 2002 AKS test has been developed which does the job in $\log(n)^{12}$ for an n bit number.*

Prime Decomposition

145906768007583323230186939349070635292
401872375357164399581871019873438799005
358938369571402670149802121818086292467
422828157022922076746906543401224889672
472407926969987100581290103199317858753
663710862357656510507883714297115637342
788911463535102712032765166518411726859
837988672111837205085526346618740053

=

1213107243921127189732367153161244042847242763370
1410925634549312301964373042085619324197365322416
8665410170573613652141717117137979742993348710628
29803541

x

1202752425547874888595622079373451212873338780368
2075433653899983955179850988797899869146900809131
6111533468170508320960221601463663463918124709871
05415233

Challenge



*Not all numbers composed of
multiplication of small primes*

Solutions

- No known efficient solutions exists
- The most effecient general solution is $e^{f(n)}$ steps for an n bit number.
- An n^3 algorithm exists (Shur's) for quantum computers.

Byzantine Generals



Idea?

*N generals in a siege communicate
unreliably to reach a decision.*

*How many trusted generals at least
may exists to still reach a correct
decision?*

Why important?

*In a distributed system there always exist
faulty peers, etither willingly or unwillingly.*

*Consensus protocols designed to overcome
this problem.*



Byzantine Generals



Challenge



Not always possible to know how many peers fail...

Solutions

- For N faulty peers at least $3N+1$ peers needed
- PBFT, Paxos, Raft ...

Knapsack Problem

An ordinary processor usually works in roughly

$5\text{GHz} = 5 \times 10^9$ ($\sim 2^{32}$)
Operations per second

In a year there exists roughly

$60 \times 60 \times 24 \times 365 \sim 32 \times 10^6 \text{ s}$ ($\sim 2^{25}$)

Why important?



*In real world optimization problems,
you search for the most valuable
combinations*

How hard?

- *0-1 knapsack problem with integer weights for n items requires 2^n steps*

That is for 60 items

~8 years !!!

..... for 90 items

~8 billion years !!!

Knapsack Problem



Challenge

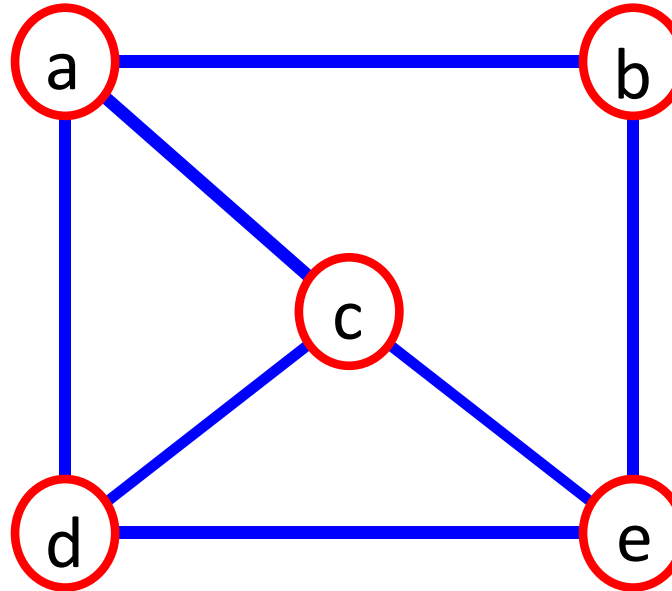


It takes very long to try all combinations...

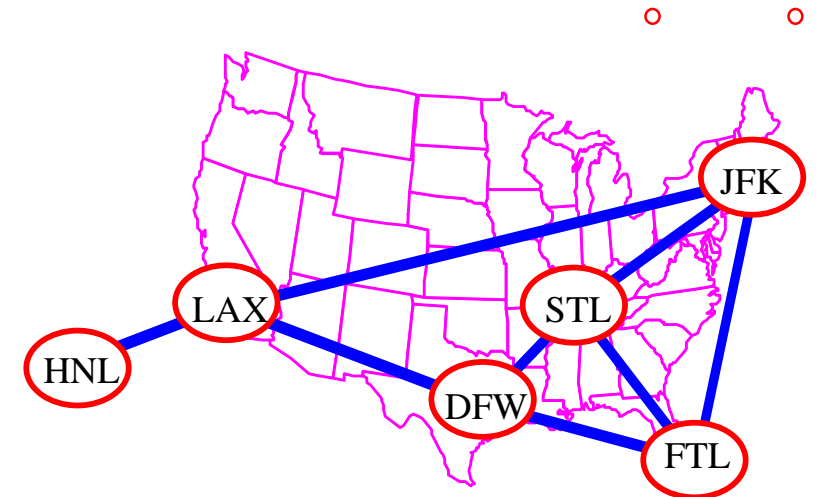
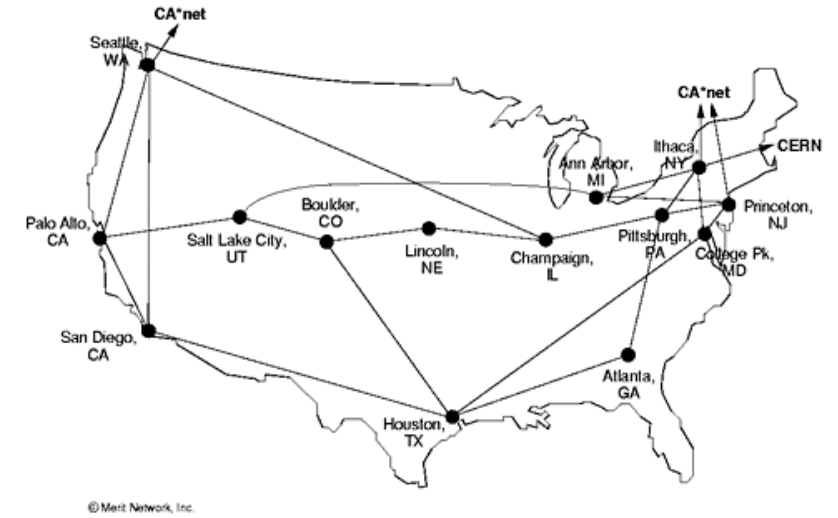
Solutions

- Can a value of at least V be achieved without exceeding the weight W ?
 - No known efficient solutions
- What is the maximum value V ?
 - Harder than the decision.

Graph primers



NSFNET T1 Network 1991



Hamiltonian Path

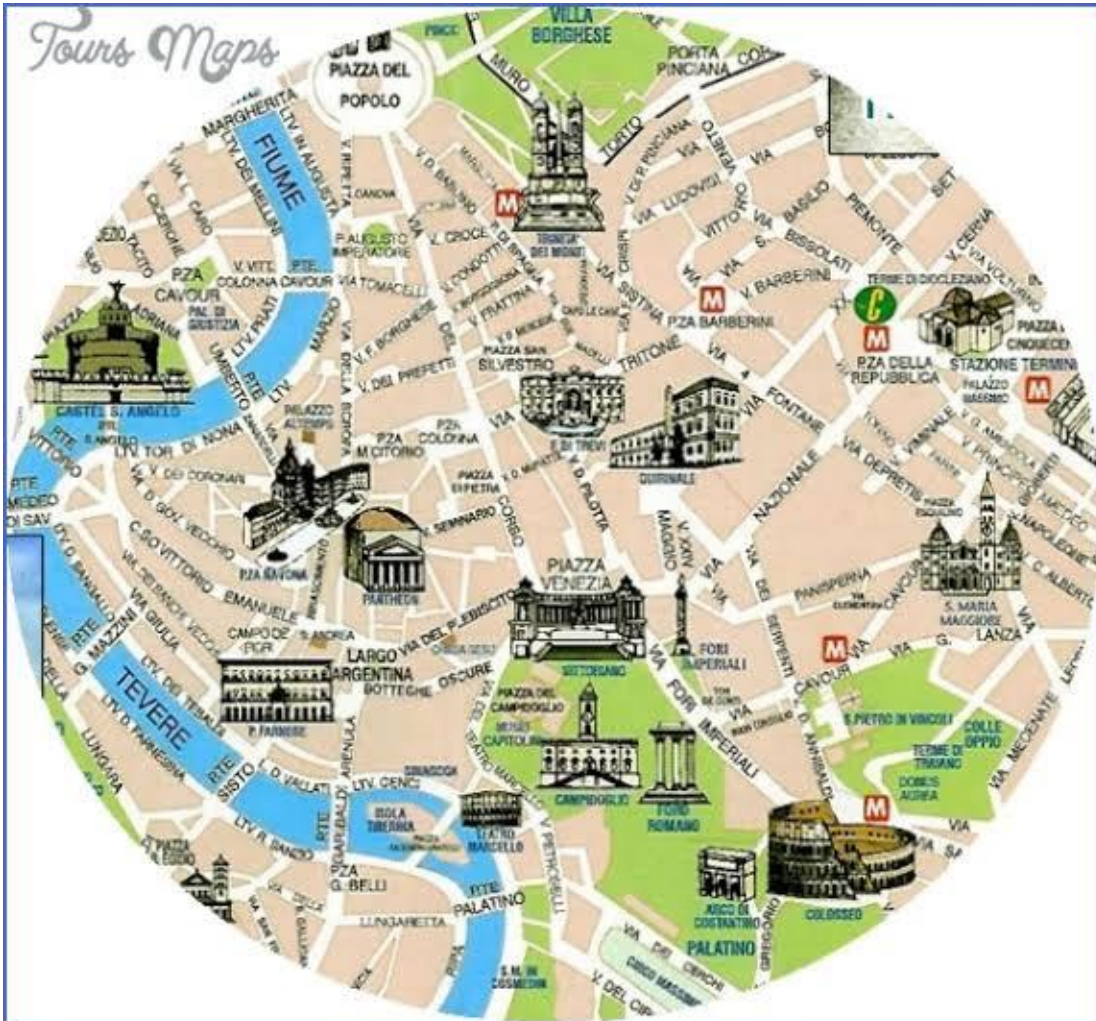


Idea?

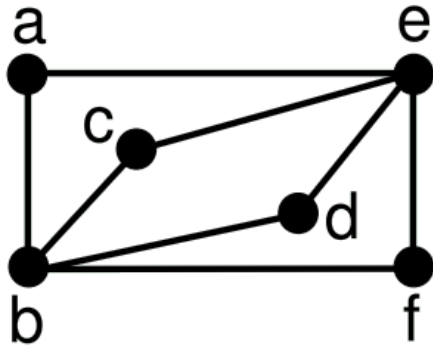
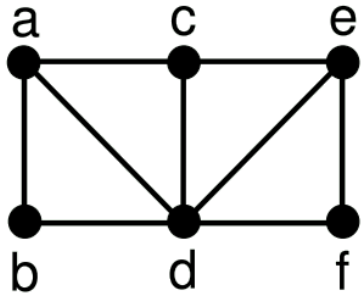
Can we stop by all the attractions without visiting an attraction twice?

Why important?

We want to convey information to a network of peers without repetition



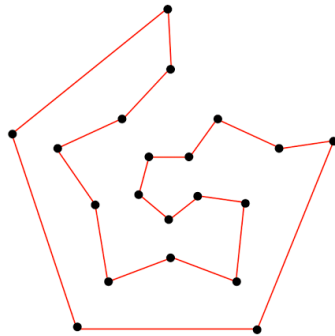
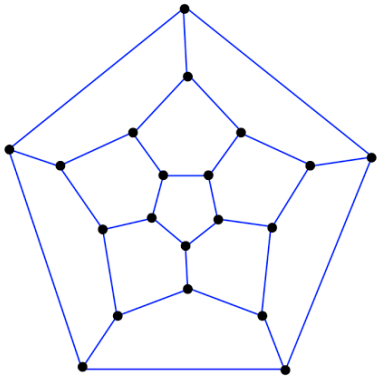
Hamiltonian Path



Challenge



Real World graphs are too complicated to try all possible variations...



Solutions

- No known efficient solutions exists
- The most effecient general solution is 1.657^n steps for n cities
- Efficient solutions exist under special conditions

Traveling Salesman



Why important?



Frequently confronted in geo-optimization problems.

How hard?

- *A special case of Ham-Path where we try to obtain the **shortest** path where we end up at the point we start.*

$$(n-1)!/2$$

- *For 16 cities $\sim 2^{40}$ tours to check*
- *There are $\sim 2^{265}$ atoms in the universe.*
- *There are $\sim 2^{90}$ nanosec. since the big bang.*
- *There are $\sim 2^{355}$ tours to check for **~ 75 cities***



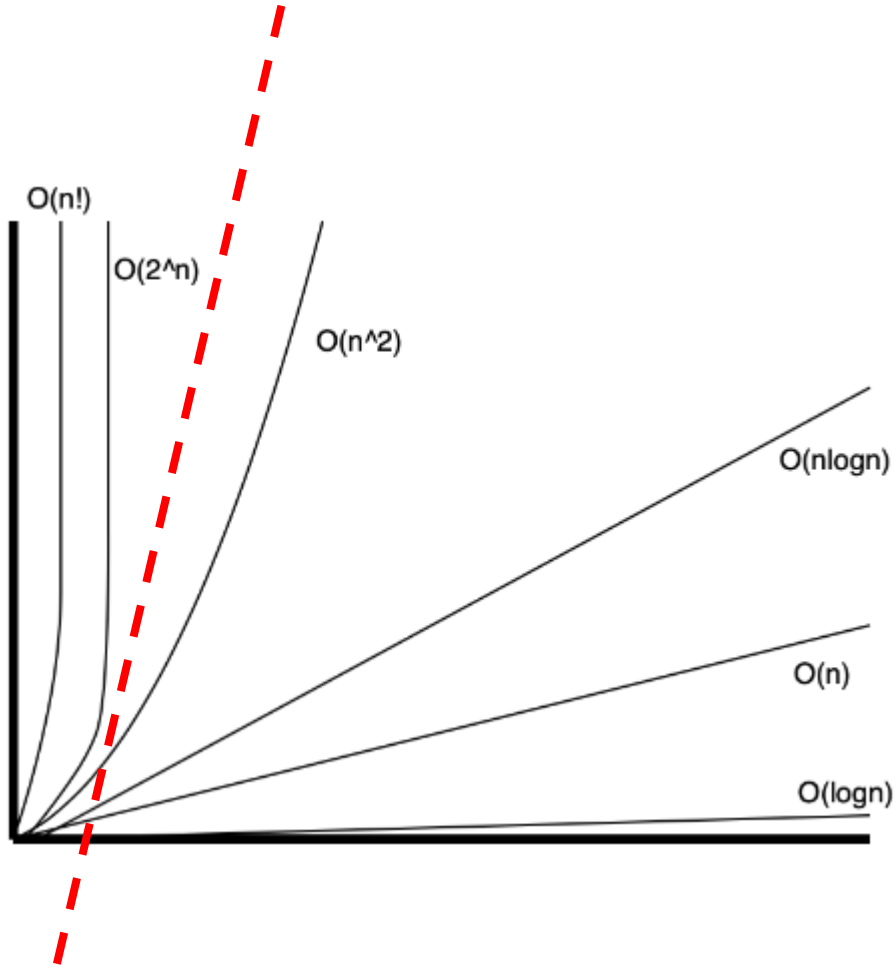
P vs NP

Class P: Easy to solve

- Search for a number in a non-sorted list
- Search for a number in a sorted list

Class NP: Easy to check

- Hamiltonian Path
- Traveling salesman



Boolean Satisfiability Problem

Idea?



Can we find values that makes a boolean formula TRUE?

Why important?

It is the foundation of computer inference...

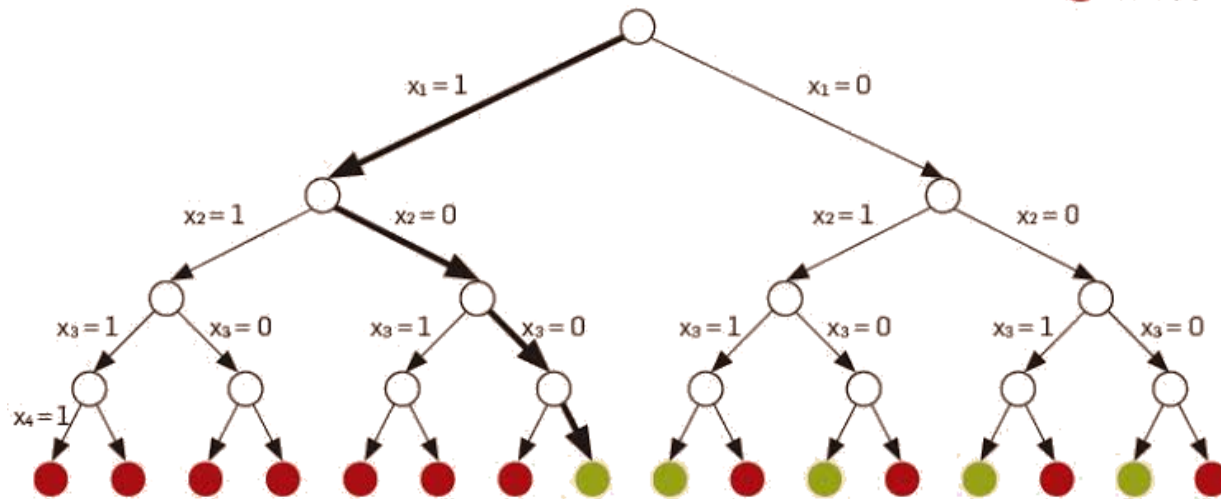


$$(p \wedge q) \vee r$$

T T F

$$(\neg x_1 \vee \neg x_2) \wedge (\neg x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_1 \vee x_3 \vee \neg x_4) \wedge (x_1 \vee x_4)$$

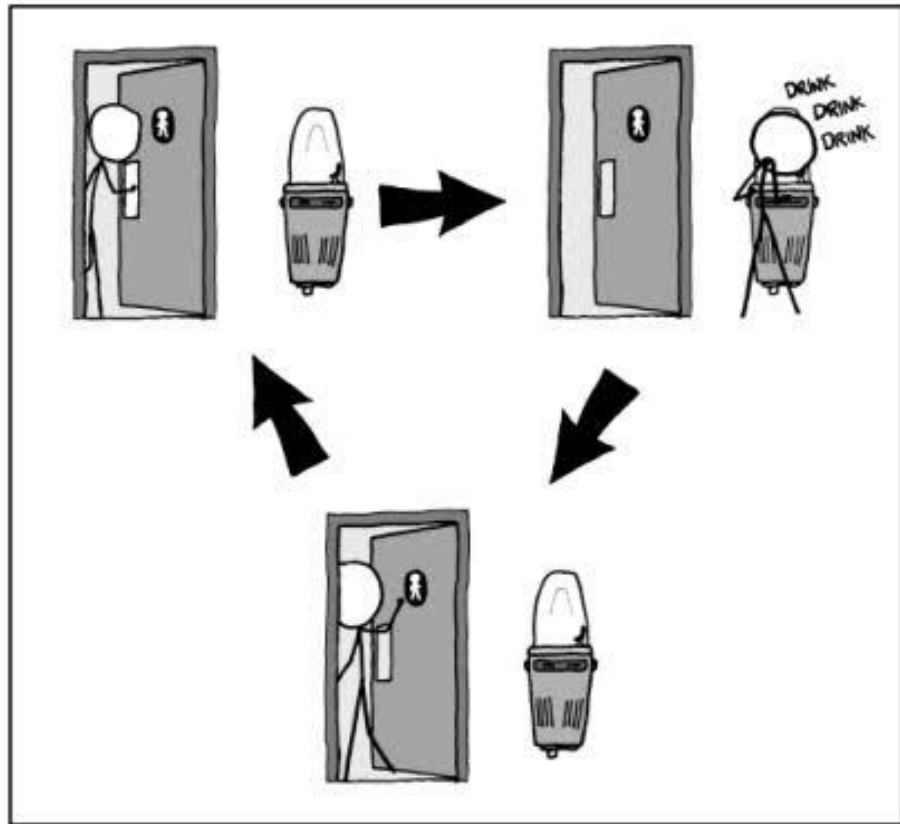
○ Unknown
● True (1)
● False (0)



Solutions

- NP problem
- Most NP problems can be transformed to SAT

The Halting Problem



I AVOID DRINKING FOUNTAINS OUTSIDE BATHROOMS
BECAUSE I'M AFRAID OF GETTING TRAPPED IN A LOOP.



Idea?

*Can we write a program that
understands if a program would
stop computing with a certain input,
or not?*

Why important?

*It shows that it is impossible to solve
some problems, in theory*

Thank you very much...

CHOTCHKIES RESTAURANT	
~ APPETIZERS ~	
MIXED FRUIT	2.15
FRENCH FRIES	2.75
SIDE SALAD	3.35
HOT WINGS	3.55
MOZZARELLA STICKS	4.20
SAMPLER PLATE	5.80
~ SANDWICHES ~	
BARBECUE	6.55

