

```

1  class VolatileExample {
2      int x = 0;
3      volatile boolean v = false;
4      public void writer() {
5          x = 42;
6          v = true;
7      }
8      public void reader() {
9          if (v == true) {
10             int y = 100/x;
11         }
12     }
13 }

```

Figure 3.16 Volatile field example from Exercise 28.

**Exercise 29.** Is the following property equivalent to saying that object  $x$  is wait-free?

For every infinite history  $H$  of  $x$ , every thread that takes an infinite number of steps in  $H$  completes an infinite number of method calls.

**Exercise 30.** Is the following property equivalent to saying that object  $x$  is lock-free?

For every infinite history  $H$  of  $x$ , an infinite number of method calls are completed.

**Exercise 31.** Consider the following rather unusual implementation of a method  $m$ . In every history, the  $i^{\text{th}}$  time a thread calls  $m$ , the call returns after  $2^i$  steps. Is this method wait-free, bounded wait-free, or neither?

**Exercise 32.** This exercise examines a queue implementation (Fig. 3.17) whose `enq()` method does not have a linearization point.

The queue stores its items in an `items` array, which for simplicity we will assume is unbounded. The `tail` field is an `AtomicInteger`, initially zero. The `enq()` method reserves a slot by incrementing `tail`, and then stores the item at that location. Note that these two steps are not atomic: there is an interval after `tail` has been incremented but before the item has been stored in the array.

The `deq()` method reads the value of `tail`, and then traverses the array in ascending order from slot zero to the `tail`. For each slot, it swaps `null` with the current contents, returning the first non-`null` item it finds. If all slots are `null`, the procedure is restarted.

Give an example execution showing that the linearization point for `enq()` cannot occur at Line 15.

Hint: give an execution where two `enq()` calls are not linearized in the order they execute Line 15.

```

1  public class HWQueue<T> {
2      AtomicReference<T>[] items;
3      AtomicInteger tail;
4      static final int CAPACITY = 1024;
5
6      public HWQueue() {
7          items = (AtomicReference<T>[])Array.newInstance(AtomicReference.class,
8              CAPACITY);
9          for (int i = 0; i < items.length; i++) {
10             items[i] = new AtomicReference<T>(null);
11         }
12         tail = new AtomicInteger(0);
13     }
14     public void enq(T x) {
15         int i = tail.getAndIncrement();
16         items[i].set(x);
17     }
18     public T deq() {
19         while (true) {
20             int range = tail.get();
21             for (int i = 0; i < range; i++) {
22                 T value = items[i].getAndSet(null);
23                 if (value != null) {
24                     return value;
25                 }
26             }
27         }
28     }
29 }

```

Figure 3.17 Herlihy/Wing queue.

Give another example execution showing that the linearization point for `enq()` cannot occur at Line 16.

Since these are the only two memory accesses in `enq()`, we must conclude that `enq()` has no single linearization point. Does this mean `enq()` is not linearizable?

**Exercise 33.** Prove that sequential consistency is nonblocking.