Problem 18

Problem 1. $\forall L1, L2 \in \mathtt{NatList}, \mathrm{rmDup}(L1 @ L2) = \mathrm{rmDup}(\mathrm{rmDup}(L1) @ \mathrm{rmDup}(L2)).$

Proof. By structural induction on L1.

(1) Base case

What to show: $\operatorname{rmDup}(nil @ l2) = \operatorname{rmDup}(\operatorname{rmDup}(nil) @ \operatorname{rmDup}(l2))$ where $l2 \in \mathtt{NatList}$.

$$\begin{array}{c} \operatorname{rmDup}(\underline{nil @ l2}) \longrightarrow \operatorname{rmDup}(l2) & (\text{by @1}) \\ \operatorname{rmDup}(\underline{rmDup}(nil) @ \operatorname{rmDup}(l2)) \longrightarrow \operatorname{rmDup}(\underline{nil @ \operatorname{rmDup}(l2)}) \\ & \qquad \qquad (\text{by rmDup1}) \\ \longrightarrow \underline{\operatorname{rmDup}(\operatorname{rmDup}(l2))} & (\text{by @1}) \\ \longrightarrow \operatorname{rmDup}(l2) & (\text{by Problem 17}) \end{array}$$

(2) Induction case

What to show: $\operatorname{rmDup}((x \mid l1) @ l2) = \operatorname{rmDup}(\operatorname{rmDup}(x \mid l1) @ \operatorname{rmDup}(l2))$

Induction hypothesis: $\operatorname{rmDup}(l1 @ l2) = \operatorname{rmDup}(\operatorname{rmDup}(l1) @ \operatorname{rmDup}(l2))$ where $x \in \operatorname{PNat}$ and $l1, l2 \in \operatorname{NatList}$.

We use case splitting for our proofs as follows:

Case 1: has(l1, x) = true.

$$\operatorname{rmDup}(\underbrace{(x \mid l1) @ l2}) \longrightarrow \operatorname{rmDup}(x \mid (l1 @ l2)) \qquad \text{(by @2)}$$

$$\longrightarrow \operatorname{if} \ \underline{\operatorname{has}(l1 @ l2, x)} \ \operatorname{then} \ \operatorname{rmDup}(l1 @ l2)$$

$$= \operatorname{lse} \ (x \mid \operatorname{rmDup}(l1 @ l2)) \ \operatorname{fi} \ \ (\operatorname{by} \ \operatorname{rmDup}(l1 @ l2))$$

$$= \operatorname{lse} \ (x \mid \operatorname{rmDup}(l1 @ l2)) \ \operatorname{fi} \ \ (\operatorname{by} \ \operatorname{Problem} \ 6 - \operatorname{Lemma} \ 1)$$

$$\longrightarrow \operatorname{if} \ \underline{(true \ or \ \operatorname{has}(l2, x))} \ \operatorname{then} \ \operatorname{rmDup}(l1 @ l2)$$

$$= \operatorname{lse} \ (x \mid \operatorname{rmDup}(l1 @ l2)) \ \operatorname{fi} \ \ (\operatorname{by} \ \operatorname{case} \ \operatorname{splitting})$$

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\longrightarrow if true then rmDup(l1 @ l2)
                                         else (x \mid \text{rmDup}(l1 @ l2)) fi
                                                                                             (by or)
                                  \longrightarrow \text{rmDup}(l1 @ l2)
                                                                                             (by if1)
                                  \longrightarrow rmDup(rmDup(l1) @ rmDup(l2)) (by IH)
\operatorname{rmDup}(\operatorname{rmDup}(x \mid l1) \otimes \operatorname{rmDup}(l2)) \longrightarrow \operatorname{rmDup}(\operatorname{if has}(l1, x) \operatorname{then rmDup}(l1))
                                                              else (x \mid \text{rmDup}(l1)) fi) @ rmDup(l2))
                                                                                    (by rmDup2)
                                                       \longrightarrow rmDup((if true then rmDup(l1)
                                                              else (x \mid \text{rmDup}(l1)) fi) @ rmDup(l2))
                                                                             (by case splitting)
                                                       \longrightarrow rmDup(rmDup(l1) @ rmDup(l2))
                                                                                             (by if1)
Case 2: has(l1, x) = false.
\operatorname{rmDup}((x \mid l1) @ l2) \longrightarrow \operatorname{rmDup}(x \mid (l1 @ l2))
                                                                                            (by @2)
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 \longrightarrow if has(l1 @ l2, x) then rmDup(l1 @ l2) else $(x \mid \text{rmDup}(l1 @ l2))$ fi (by rmDup2) \longrightarrow if (has(l1, x) or has(l2, x)) then rmDup(l1 @ l2)else $(x \mid \text{rmDup}(l1 @ l2))$ fi (by Problem 6 - Lemma 1) \longrightarrow if $(false\ or\ has(l2,x))$ then rmDup $(l1\ @\ l2)$ else $(x \mid \text{rmDup}(l1 @ l2))$ fi (by case splitting) \longrightarrow if has(l2, x) then rmDup(l1 @ l2) else $(x \mid \text{rmDup}(l1 @ l2))$ fi (by or) \longrightarrow if has(l2, x) then rmDup(rmDup(l1) @ rmDup(l2)) else $(x \mid \text{rmDup}(l1 @ l2))$ fi (by IH) \longrightarrow if has(l2, x) then rmDup(rmDup(l1) @ rmDup(l2)) else $(x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2)))$ fi (by IH)

 $\operatorname{rmDup}(\operatorname{rmDup}(x\mid l1) @ \operatorname{rmDup}(l2)) \longrightarrow \operatorname{rmDup}((\operatorname{if} \ \operatorname{has}(l1,x) \ \operatorname{then} \ \operatorname{rmDup}(l1)$

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else (x \mid \text{rmDup}(l1)) fi) @ rmDup(l2))
                         (by rmDup2)
\longrightarrow rmDup((if false then rmDup(l1))
     else (x \mid \text{rmDup}(l1)) fi) @ rmDup(l2))
                   (by case splitting)
\longrightarrow \operatorname{rmDup}((x \mid \operatorname{rmDup}(l1)) @ \operatorname{rmDup}(l2))
                                 (by if2)
\longrightarrow \operatorname{rmDup}(x \mid (\operatorname{rmDup}(l1) @ \operatorname{rmDup}(l2)))
                                (by @2)
\longrightarrow if has(rmDup(l1) @ rmDup(l2), x)
     then rmDup(rmDup(l1) @ rmDup(l2))
     else (x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2))) fi
                         (by rmDup2)
\longrightarrow if has(rmDup(l1), x) or has(rmDup(l2), x)
     then rmDup(rmDup(l1) @ rmDup(l2))
     else (x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2))) fi
       (by Problem 6 - Lemma 1)
\longrightarrow if has(l1, x) or has(rmDup(l2), x)
     then rmDup(rmDup(l1) @ rmDup(l2))
     else (x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2))) fi
      (by Problem 17 - Lemma 1)
\longrightarrow if false or has(rmDup(l2), x)
     then rmDup(rmDup(l1) @ rmDup(l2))
     else (x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2))) fi
                   (by case splitting)
\longrightarrow if has(rmDup(l2), x)
     then rmDup(rmDup(l1) @ rmDup(l2))
     else (x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2))) fi
                                 (by or)
\longrightarrow if has(l2, x)
     then rmDup(rmDup(l1) @ rmDup(l2))
     else (x \mid \text{rmDup}(\text{rmDup}(l1) @ \text{rmDup}(l2))) fi
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(by Problem 17 - Lemma 1)