Problem 4

Problem 1. $\forall L \in \mathtt{NatList}, \operatorname{rev}(\operatorname{rev}(L)) = L.$

Proof. By structural induction on L.

(1) Base case

What to show: rev(rev(nil)) = nil.

$$rev(rev(nil)) \longrightarrow rev(nil)$$
 (by rev1)

$$\longrightarrow nil$$
 (by rev1)

(2) Induction case

What to show: $rev(rev(x \mid l)) = x \mid l$ Induction hypothesis: rev(rev(l)) = l

where $x \in PNat$ and $l \in NatList$. Note that x, l are fresh constants¹.

$$\begin{array}{c} \operatorname{rev}(\underline{\operatorname{rev}(x\mid l)}) \longrightarrow \underline{\operatorname{rev}(\operatorname{rev}(l) @ (x\mid nil))} & \text{(by rev2)} \\ \longrightarrow \underline{\operatorname{rev}(x\mid nil)} @ \operatorname{rev}(\operatorname{rev}(l)) & \text{(by Lemma 1)} \\ \longrightarrow \underline{(\operatorname{rev}(nil)} @ (x\mid nil)) @ \operatorname{rev}(\operatorname{rev}(l)) & \text{(by rev2)} \\ \longrightarrow \underline{(nil @ (x\mid nil))} @ \operatorname{rev}(\operatorname{rev}(l)) & \text{(by rev1)} \\ \longrightarrow \underline{(x\mid nil)} @ \operatorname{rev}(\operatorname{rev}(l)) & \text{(by @1)} \\ \longrightarrow x\mid \underline{(nil @ \operatorname{rev}(\operatorname{rev}(l)))} & \text{(by @2)} \\ \longrightarrow x\mid \underline{\operatorname{rev}(\operatorname{rev}(l))} & \text{(by @1)} \\ \longrightarrow x\mid l & \text{(by IH)} \end{array}$$

Lemma 1. $\forall L1, L2 \in \mathtt{NatList}, \operatorname{rev}(L1 @ L2) = \operatorname{rev}(L2) @ \operatorname{rev}(L1).$

Proof. By structural induction on L1.

¹A fresh constant of a sort denotes an arbitrary value of the sort, and has never been used before.

(1) Base case

rev(nil @ l2) = rev(l2) @ rev(nil)What to show: where $l2 \in NatList$. Note that l2 is a fresh constant.

$$rev(nil @ l2) \longrightarrow rev(l2)$$
 (by @1)

$$\operatorname{rev}(l2) \ @ \ \operatorname{\underline{rev}}(nil) \longrightarrow \operatorname{\underline{rev}}(l2) \ @ \ nil \tag{by rev1}$$

$$\longrightarrow \text{rev}(l2)$$
 (by Lemma 2)

(2) Induction case

rev((x | l1) @ l2) = rev(l2) @ rev(x | l1)What to show: Induction hypothesis: rev(l1 @ l2) = rev(l2) @ rev(l1)

where $x \in PNat$, and $l1, l2 \in NatList$. Note that x, l1, l2 are fresh constants.

$$rev((x \mid l1) @ l2) \longrightarrow rev(x \mid (l1 @ l2))$$
 (by @2)

$$\longrightarrow \underline{\operatorname{rev}(l1 @ l2)} @ (x \mid nil)$$
 (by rev2)

$$\longrightarrow (\text{rev}(l2) @ \text{rev}(l1)) @ (x \mid nil)$$
 (by IH)

$$\longrightarrow \operatorname{rev}(l2) @ (\operatorname{rev}(l1) @ (x \mid nil))$$

(by Lemma 2 in Problem 2)

$$\operatorname{rev}(l2) \ @ \ \operatorname{rev}(x \mid l1) \longrightarrow \operatorname{rev}(l2) \ @ \ (\operatorname{rev}(l1) \ @ \ (x \mid nil)) \qquad \text{(by rev2)}$$

Lemma 2. $\forall L \in \mathtt{NatList}, L @ nil = L.$

Proof. By structural induction on L.

(1) Base case

What to show: nil @ nil = nil.

$$\underline{nil} \ @ \ nil \longrightarrow nil$$
 (by @1)

(2) Induction case

 $(x \mid l) @ nil = x \mid l$ What to show:

Induction hypothesis: l @ nil = l

where $x \in PNat$ and $l \in NatList$. Note that x, l are fresh constants.

$$\longrightarrow x \mid l$$
 (by IH)