

Problem 5

Problem 1. $\forall L \in \text{NatList}, \text{size}(L) = \text{size}(\text{rev}(L))$.

Proof. By structural induction on L .

(1) Base case

What to show: $\text{size}(\text{nil}) = \text{size}(\text{rev}(\text{nil}))$.

$$\begin{array}{ll} \underline{\text{size}(\text{nil})} \longrightarrow 0 & \text{(by size1)} \\ \text{size}(\underline{\text{rev}(\text{nil})}) \longrightarrow \underline{\text{size}(\text{nil})} & \text{(by rev1)} \\ \longrightarrow \text{nil} & \text{(by size1)} \end{array}$$

(2) Induction case

What to show: $\text{size}(x \mid l) = \text{size}(\text{rev}(x \mid l))$

Induction hypothesis: $\text{size}(l) = \text{size}(\text{rev}(l))$

where $x \in \text{PNat}$ and $l \in \text{NatList}$.

$$\begin{array}{ll} \underline{\text{size}(x \mid l)} \longrightarrow \text{s}(\underline{\text{size}(l)}) & \text{(by size2)} \\ \longrightarrow \underline{\text{s}(\text{size}(\text{rev}(l)))} & \text{(by IH)} \\ \longrightarrow \text{size}(\text{rev}(l)) + \text{s}(0) & \text{(by Lemma 2)} \\ \text{size}(\underline{\text{rev}(x \mid l)}) \longrightarrow \underline{\text{size}(\text{rev}(l) @ (x \mid \text{nil}))} & \text{(by rev2)} \\ \longrightarrow \text{size}(\text{rev}(l)) + \underline{\text{size}(x \mid \text{nil})} & \text{(by Lemma 1)} \\ \longrightarrow \text{size}(\text{rev}(l)) + \text{s}(\underline{\text{size}(\text{nil})}) & \text{(by size2)} \\ \longrightarrow \text{size}(\text{rev}(l)) + \text{s}(0) & \text{(by size1)} \end{array}$$

□

Lemma 1. $\forall L1, L2 \in \text{NatList}, \text{size}(L1 @ L2) = \text{size}(L1) + \text{size}(L2)$.

Proof. By structural induction on $L1$.

(1) Base case

What to show: $\text{size}(\text{nil} @ l2) = \text{size}(\text{nil}) + \text{size}(l2)$
where $l2 \in \text{NatList}$.

$$\begin{aligned} \text{size}(\text{nil} @ l2) &\longrightarrow \text{size}(l2) && \text{(by @1)} \\ \underline{\text{size}(\text{nil}) @ \text{size}(l2)} &\longrightarrow \underline{0 + \text{size}(l2)} && \text{(by size1)} \\ &\longrightarrow \text{size}(l2) && \text{(by +1)} \end{aligned}$$

(2) Induction case

What to show: $\text{size}((x | l1) @ l2) = \text{size}(x | l1) + \text{size}(l2)$
Induction hypothesis: $\text{size}(l1 @ l2) = \text{size}(l1) + \text{size}(l2)$
where $x \in \text{PNat}$, and $l1, l2 \in \text{NatList}$.

$$\begin{aligned} \text{size}((x | l1) @ l2) &\longrightarrow \underline{\text{size}(x | (l1 @ l2))} && \text{(by @2)} \\ &\longrightarrow \underline{s(\text{size}(l1 @ l2))} && \text{(by size2)} \\ &\longrightarrow \underline{s(\text{size}(l1) + \text{size}(l2))} && \text{(by IH)} \\ \underline{\text{size}(x | l1) + \text{size}(l2)} &\longrightarrow \underline{s(\text{size}(l1)) + \text{size}(l2)} && \text{(by size2)} \\ &\longrightarrow \underline{s(\text{size}(l1) + \text{size}(l2))} && \text{(by +2)} \end{aligned}$$

□

Lemma 2. $\forall X \in \text{PNat}, s(X) = X + s(0)$.

Proof. By structural induction on X .

(1) Base case

What to show: $s(0) = 0 + s(0)$.

$$\underline{0 + s(0)} \longrightarrow s(0) \quad \text{(by +1)}$$

(2) Induction case

What to show: $s(s(x)) = s(x) + s(0)$
Induction hypothesis: $s(x) = x + s(0)$
where $x \in \text{PNat}$.

$$\begin{aligned} \underline{s(x) + s(0)} &\longrightarrow \underline{s(x + s(0))} && \text{(by +2)} \\ &\longrightarrow \underline{s(s(0) + x)} && \text{(by comm+)} \\ &\longrightarrow \underline{s(s(0 + x))} && \text{(by +2)} \\ &\longrightarrow s(s(x)) && \text{(by +1)} \end{aligned}$$

□