Problem 6

Problem 1. $\forall X \in \mathtt{PNat}, \forall L \in \mathtt{NatList}, \mathrm{has}(L, X) = \mathrm{has}(\mathrm{rev}(L), X).$

Proof. By structural induction on L.

(1) Base case

What to show: has(nil, x) = has(rev(nil), x) where $x \in PNat$.

$$has(nil, x) \longrightarrow false$$
 (by has1)

$$has(rev(nil), x) \longrightarrow has(nil, x)$$
 (by rev1)

$$\longrightarrow false$$
 (by has1)

(2) Induction case

What to show: $has(y \mid l, x) = has(rev(y \mid l), x)$ Induction hypothesis: has(l, x) = has(rev(l), x)where $x, y \in PN$ at and $l \in NatList$.

$$\frac{\operatorname{has}(y\mid l,x)}{\longrightarrow} \xrightarrow{} (x=y) \text{ or } \frac{\operatorname{has}(l,x)}{\operatorname{has}(\operatorname{rev}(l),x)} \qquad \text{(by IH)}$$

$$\operatorname{has}(\underline{\operatorname{rev}(y\mid l)},x) \xrightarrow{} \frac{\operatorname{has}(\operatorname{rev}(l) @ (y\mid nil)),x)}{\operatorname{has}(\operatorname{rev}(l),x) \text{ or } \frac{\operatorname{has}(y\mid nil,x)}{\operatorname{(loy Lemma 1)}} \qquad \text{(by Lemma 1)}$$

$$\xrightarrow{} \operatorname{has}(\operatorname{rev}(l),x) \text{ or } \underbrace{((x=y) \text{ or } \frac{\operatorname{has}(nil,x))}{\operatorname{(by has 2)}}} \qquad \text{(by has 1)}$$

$$\xrightarrow{} \operatorname{has}(\operatorname{rev}(l),x) \text{ or } \underbrace{(x=y) \text{ or } false}_{} \qquad \text{(by or)}$$

$$\xrightarrow{} \underbrace{\operatorname{has}(\operatorname{rev}(l),x) \text{ or } (x=y)}_{} \qquad \text{(by or)}$$

$$\xrightarrow{} \underbrace{\operatorname{has}(\operatorname{rev}(l),x) \text{ or } (x=y)}_{} \qquad \text{(by comm-or)}$$

Lemma 1. $\forall X \in \mathtt{PNat}, \forall L1, L2 \in \mathtt{NatList}, \mathrm{has}(L1 @ L2, X) = \mathrm{has}(L1, X) \text{ or } \mathrm{has}(L2, X).$

Proof. By structural induction on L1.

(1) Base case

What to show: has(nil @ l2, x) = (has(nil, x) or has(l2, x)) where $x \in PNat$ and $l2 \in NatList$.

$$\begin{array}{ccc} \operatorname{has}(\underline{nil} \@\ l2, x) &\longrightarrow \operatorname{has}(l2, x) & \text{(by @1)} \\ \underline{\operatorname{has}(nil, x)} \text{ or } \operatorname{has}(l2, x) &\longrightarrow \underline{false} \text{ or } \operatorname{has}(l2, x) & \text{(by has1)} \\ &\longrightarrow \operatorname{has}(l2, x) & \text{(by or)} \end{array}$$

(2) Induction case

What to show: $\operatorname{has}((y \mid l1) @ l2, x) = (\operatorname{has}(y \mid l1, x) \text{ or } \operatorname{has}(l2, x))$ Induction hypothesis: $\operatorname{has}(l1 @ l2, x) = (\operatorname{has}(l1, x) \text{ or } \operatorname{has}(l2, x))$ where $x, y \in \operatorname{PNat}$, and $l1, l2 \in \operatorname{NatList}$.

$$\begin{array}{c} \operatorname{has}(\underline{(y \mid l1) @ l2},x) \longrightarrow \underline{\operatorname{has}(y \mid (l1 @ l2),x)} & \operatorname{(by @2)} \\ \longrightarrow \overline{(x=y) \text{ or } \underline{\operatorname{has}(l1 @ l2,x)}} & \operatorname{(by has2)} \\ \longrightarrow \overline{(x=y) \text{ or } (\operatorname{has}(l1,x) \text{ or } \operatorname{has}(l2,x))} & \operatorname{(by IH)} \\ \underline{\operatorname{has}(y \mid l1,x)} & \operatorname{or has}(l2,x) \longrightarrow \underline{((x=y) \text{ or } \operatorname{has}(l1,x)) \text{ or } \operatorname{has}(l2,x)} \\ \longrightarrow \overline{(x=y) \text{ or } (\operatorname{has}(l1,x) \text{ or } \operatorname{has}(l2,x))} \\ & \hookrightarrow (x=y) \text{ or } (\operatorname{has}(l1,x) \text{ or } \operatorname{has}(l2,x)) \\ & \operatorname{(by assoc-or)} \end{array}$$