# Problem 13

**Problem 1.**  $\forall X \in \mathtt{PNat}, \forall L \in \mathtt{NatList}, \mathtt{count}(L, X) = \mathtt{count}(\mathtt{rev}(L), X).$ 

*Proof.* By structural induction on L.

# (1) Base case

What to show:  $\operatorname{count}(nil, x) = \operatorname{count}(\operatorname{rev}(nil), x)$  where  $x \in \operatorname{PNat}$ . Note that x is a fresh constant<sup>1</sup>.

$$\operatorname{count}(nil, x) \longrightarrow 0$$
 (by cnt1)

$$\operatorname{count}(\operatorname{rev}(nil), x) \longrightarrow \operatorname{count}(nil, x)$$
 (by rev1)

$$\longrightarrow 0$$
 (by cnt1)

## (2) Induction case

What to show:  $\operatorname{count}(y \mid l, x) = \operatorname{count}(\operatorname{rev}(y \mid l), x)$ 

Induction hypothesis: count(l, x) = count(rev(l, x))

where  $x,y\in \mathtt{PNat}$  and  $l\in \mathtt{NatList}.$  Note that x,y,l are fresh constants.

We use case splitting for our proofs as follows:

Case 1: y = x

$$\begin{array}{c} \operatorname{count}(\underline{y}\mid l,x) \longrightarrow \underline{\operatorname{count}(x\mid l,x)} & (\operatorname{by\ case\ splitting}) \\ \longrightarrow \operatorname{if\ } \underline{(x=x)} \operatorname{\ then\ } s(\operatorname{count}(l,x)) \operatorname{\ else\ count}(l,x) \operatorname{\ fi} \\ & (\operatorname{by\ cnt2}) \\ \longrightarrow \underline{\operatorname{if\ } true\ \operatorname{then\ } s(\operatorname{count}(l,x)) \operatorname{\ else\ count}(l,x) \operatorname{\ fi} \\ & (\operatorname{by\ equality}) \\ \longrightarrow s(\operatorname{count}(\operatorname{rev}(l),x)) & (\operatorname{by\ if1}) \\ \longrightarrow s(\operatorname{count}(\operatorname{rev}(l),x)) & (\operatorname{by\ IH}) \\ \operatorname{count}(\operatorname{rev}(\underline{y}\mid l),x) \longrightarrow \operatorname{count}(\underline{\operatorname{rev}(x\mid l)},x) & (\operatorname{by\ case\ splitting}) \\ \longrightarrow \underline{\operatorname{count}(\operatorname{rev}(l)\ @\ (x\mid nil),x)} & (\operatorname{by\ rev2}) \end{array}$$

<sup>&</sup>lt;sup>1</sup>A fresh constant of a sort denotes an arbitrary value of the sort, and has never been used before.

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\longrightarrow count(rev(l), x) + count(x | nil, x)
                                                                                                   (by Lemma 1)
                                  \longrightarrow count(rev(l), x)+ if (x = x) then s(\text{count}(nil, x))
                                         else \operatorname{count}(nil, x) fi
                                                                                                            (by cnt2)
                                  \longrightarrow count(rev(l), x) + if true then s(count(nil, x))
                                          else \operatorname{count}(nil, x) fi
                                                                                                     (by equality)
                                  \longrightarrow count(rev(l), x) + s(count(nil, x))
                                                                                                              (by if1)
                                  \longrightarrow \operatorname{count}(\operatorname{rev}(l), x) + s(0)
                                                                                                            (by cnt1)
                                  \longrightarrow s(0) + \operatorname{count}(\operatorname{rev}(l), x)
                                                                                                     (by comm + )
                                  \longrightarrow s(0 + \operatorname{count}(\operatorname{rev}(l), x))
                                                                                                              (by +2)
                                  \longrightarrow s(\operatorname{count}(\operatorname{rev}(l), x))
                                                                                                              (by +1)
Case 2: (y = x) = false
        \operatorname{count}(y \mid l, x) \longrightarrow \operatorname{if}(y = x) \operatorname{then} s(\operatorname{count}(l, x)) \operatorname{else} \operatorname{count}(l, x) \operatorname{fi}
                                                                                                            (by cnt2)
                                  \longrightarrow if false then s(\operatorname{count}(l,x)) else \operatorname{count}(l,x) fi
                                                                                            (by case splitting)
                                                                                                               (by if2)
                                  \longrightarrow \operatorname{count}(l,x)
                                  \longrightarrow \operatorname{count}(\operatorname{rev}(l), x)
                                                                                                              (by IH)
\operatorname{count}(\operatorname{rev}(y \mid l), x) \longrightarrow \operatorname{count}(\operatorname{rev}(l) @ (y \mid nil), x)
                                                                                                            (by rev2)
                                 \longrightarrow count(rev(l), x) + count(y | nil, x)
                                                                                                   (by Lemma 1)
                                 \longrightarrow count(rev(l), x)+ if (y = x) then s(\text{count}(nil, x))
                                          else \operatorname{count}(nil, x) fi
                                                                                                            (by cnt2)
                                 \longrightarrow count(rev(l), x) + if false then s(count(nil, x))
                                          else \operatorname{count}(nil, x) fi
                                                                                            (by case splitting)
                                  \longrightarrow count(rev(l), x) + count(nil, x)
                                                                                                              (by if 2)
                                  \longrightarrow \operatorname{count}(\operatorname{rev}(l), x) + 0
                                                                                                            (by cnt1)
                                  \longrightarrow 0 + \operatorname{count}(\operatorname{rev}(l), x)
                                                                                                     (by comm +)
                                  \longrightarrow count(rev(l), x)
                                                                                                              (by +1)
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**Lemma 1.**  $\forall X \in PNat, \forall L1, L2 \in NatList, count(L1 @ L2, X) = count(L1, X) + count(L2, X).$ 

*Proof.* By structural induction on L1.

### (1) Base case

What to show:  $\operatorname{count}(nil \otimes l2, x) = \operatorname{count}(nil, x) + \operatorname{count}(l2, x)$  where  $x \in \operatorname{PNat}$  and  $l2 \in \operatorname{NatList}$ . Note that x, l2 are fresh constants.

$$\frac{\operatorname{count}(\underline{nil} \ @ \ l2, x) \longrightarrow \operatorname{count}(l2, x)}{\operatorname{count}(\underline{nil}, x) + \operatorname{count}(l2, x) \longrightarrow 0 + \operatorname{count}(l2, x)} \qquad \text{(by @1)}$$

$$\longrightarrow \operatorname{count}(l2, x) \qquad \text{(by ent1)}$$

$$\longrightarrow \operatorname{count}(l2, x) \qquad \text{(by +1)}$$

### (2) Induction case

What to show:  $\operatorname{count}((y \mid l1) @ l2, x) = \operatorname{count}(y \mid l1, x) + \operatorname{count}(l2, x)$ Induction hypothesis:  $\operatorname{count}(l1 @ l2, x) = \operatorname{count}(l1, x) + \operatorname{count}(l2, x)$ where  $x, y \in \operatorname{PNat}$  and  $l1, l2 \in \operatorname{NatList}$ . Note that x, y, l1, l2 are fresh constants.

We use case splitting for our proofs as follows:

Case 1: y = x

$$\begin{array}{c} \operatorname{count}((\underline{y} \mid l1) @ l2, x) \longrightarrow \operatorname{count}(\underline{(x \mid l1) @ l2}, x) \\ & \qquad \qquad (\operatorname{by \ case \ splitting}) \\ \longrightarrow \operatorname{count}(x \mid (l1 @ l2), x) & (\operatorname{by \ @2}) \\ \longrightarrow \operatorname{if \ } \underline{(x = x) \ } \operatorname{then \ } s(\operatorname{count}(l1 @ l2, x)) \\ & \qquad = \operatorname{lese \ } \operatorname{count}(l1 @ l2, x) \operatorname{fi} & (\operatorname{by \ cnt2}) \\ \longrightarrow \operatorname{if \ } true \ \operatorname{then \ } s(\operatorname{count}(l1 @ l2, x)) \\ & \qquad = \operatorname{lese \ } \operatorname{count}(l1 @ l2, x) \operatorname{fi} \\ & \qquad \qquad (\operatorname{by \ equality}) \\ \longrightarrow s(\operatorname{count}(l1 @ l2, x)) & (\operatorname{by \ if1}) \\ \longrightarrow s(\operatorname{count}(l1, x) + \operatorname{count}(l2, x)) \\ & \qquad \qquad (\operatorname{by \ IH}) \\ \operatorname{count}(\underline{y} \mid l1, x) + \operatorname{count}(l2, x) \longrightarrow \operatorname{count}(x \mid l1, x) + \operatorname{count}(l2, x) \\ & \qquad \qquad (\operatorname{by \ case \ splitting}) \\ \longrightarrow \operatorname{if \ } \underline{(x = x) \ } \operatorname{then \ } s(\operatorname{count}(l1, x)) \\ \operatorname{else \ } \operatorname{count}(l1, x) \operatorname{fi} + \operatorname{count}(l2, x) \\ & \qquad \qquad (\operatorname{by \ cnt2}) \end{array}$$

$$\frac{\operatorname{if } true \ \operatorname{then } s(\operatorname{count}(l1,x))}{\operatorname{else } \operatorname{count}(l1,x) \ \operatorname{fi} + \operatorname{count}(l2,x)} \\ \operatorname{(by \ equality)} \\ \to \underline{s(\operatorname{count}(l1,x)) + \operatorname{count}(l2,x)} \\ \operatorname{(by \ if1)} \\ \to s(\operatorname{count}(l1,x) + \operatorname{count}(l2,x)) \\ \operatorname{(by \ +2)} \\ \mathbf{Case \ 2:} \ (y=x) = false \\ \operatorname{count}(\underline{(y \mid l1) @ l2},x) \longrightarrow \underline{\operatorname{count}(y \mid (l1 @ l2),x)} \\ \to \operatorname{if} \ \underline{(y=x)} \ \operatorname{then } s(\operatorname{count}(l1 @ l2,x)) \\ \operatorname{else \ count}(l1 @ l2,x) \ \operatorname{fi} \\ \operatorname{(by \ case \ splitting)} \\ \to \underline{\operatorname{count}(l1 @ l2,x)} \\ \operatorname{else \ count}(l1 @ l2,x) \ \operatorname{fi} \\ \operatorname{(by \ case \ splitting)} \\ \to \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{count}(l1,x) \\ \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{count}(l1,x) + \operatorname{count}(l2,x) \\ \operatorname{(by \ case \ splitting)} \\ \operatorname{(by \ cas$$