

### Question 1

Given a CFG  $G = (V, \Sigma, R, S)$  with set of variables  $V = \{S\}$ , where S is the start variable; set of terminals  $\Sigma = \{0, 1, (,), \cup, *, \emptyset, \varepsilon\}$ .; and rules:

$$S \to S \cup S|SS|S^*|(S)|0|1|\emptyset|\varepsilon$$

Using this G, solve the following questions. (20 marks)

- (1) Give a derivation for the string  $(0 \cup (10)^*1)^*$ . (11 marks)
- (2) Give the corresponding parse tree for the string  $(0 \cup (10)^*1)^*$ . (9 marks)

## Question 2

Consider the following CFG  $G=(V,\Sigma,R,S)$ , where  $V=\{S,T,X\},\ \Sigma=\{a,b\}$ , the start variable is S, and the rules R are:

$$S \to aTXb$$

$$T \to XTS|\varepsilon$$

$$X \to a|b$$

Convert this G to an equivalent PDA. (20 marks)

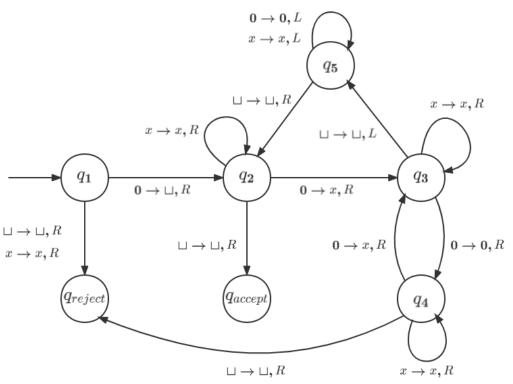
## Question 3

Use the pumping lemma to prove that the language  $A = \{0^{2n}1^{3n}0^n \mid n \geq 0\}$  is not context free. (20 marks)

## Question 4

The Turing machine M below recognizes the language  $A = \{0^{2^n} | n \ge 0\}$ .





Give the sequence of configurations when the input string is 000000. (20 marks)

# Question 5

Given the language  $E_{CFG} = \{\langle G \rangle \mid G \text{ is a CFG with } L(G) = \emptyset\}$ , prove that it is decidable (briefly describe the proof idea). (20 marks)