

CS 330 - Assignment 4

Question 1:

a. Optimal (OPT)

1	1	1	1	3	3	2	1	7											
	0	0	0	0	0	0	0	0											
		7	2	2	4	4	4	4											

b. Least recently used (LRU)

1	1	1		1		1	0		4		4	4		2	2		7		
	0	0		0		3	3		3		0	0		0	0		0		
		7		2		2	2		2		2	3		3	1		1		

c. Additional-reference-bits

1		2	2		3	3		4	4		2	2	7						
0		0	0		0	0		0	0		0	0	0						
7		7	1		1	2		2	3		3	1	1						

d. Clock (Second-chance)

1		1		1		0		4	4	4		2		2	7				
0		0		3		3		3	0	0		0		0	0				
7		2		2		2		2	2	3		3		1	1				

e. Enhanced second-chance

1		1		1	0		4	0		0	0								
0		0		3	3		3	3		1	1								
7		2		2	2		2	2		2	7								

Question 2:

- 32 MBs = $32 * 1024 * 1024$ bytes = 2^{25} \implies 25 bits length of the virtual address.
- 256 KBs = $256 * 1024$ bytes = 2^{18} \implies 18 bits length of the physical address.
- The size of a page is 512 bytes = 2^9 \implies 9 bits length of the offset into a page or frame.
- Number of pages = $32 * 1024 * 1024 / 512 = 2^{16}$ \implies 16 bits length of the page number.

e. Number of page frames = $256 * 1024 / 512 = 512$ frames.

Question 3:

a. effective memory access time = $(1 - p) * \text{memory access time} + p * \text{page fault time}$, where p is the probability of a page fault ($0 \leq p \leq 1$).

$$= (1 - p) * 125 \text{ ns} + p * 4\,500\,000 \text{ ns}$$

$$= 125 \text{ ns} + 4,499,875p \text{ ns}$$

b. $p = 1/1750 = 0.00057 \implies \text{effective memory access time} \approx 2689.9 \text{ ns} \approx 2.7 \mu\text{s}$

c. effective memory access time is approximately $22 * 125 \implies$ the access time for a page not in memory is approximately 22 times that of a page already in memory.

d. for performance degradation of no more than 8.5% \implies no more than $125 \text{ ns} + 10.625 \text{ ns}$

$$\implies 125 \text{ ns} + 10.625 \text{ ns} > 125 \text{ ns} + 4,499,875p \text{ ns}$$

$$\implies 10.625 \text{ ns} > 4,499,875p \text{ ns} \implies p < 0.00000236$$

e. No more than one memory access in $1 / 0.00000236 = 423,518$ can cause a page fault.