Hypervisors on ARM Overview and Design choices

ARM

Julien Grall < julien.grall@arm.com>

Root Linux Conference 2017

© ARM 2017

About me

- Working on ARM virtualization for the past 4 years
- With ARM since 2016
- Co-maintaining Xen on ARM with Stefano Stabellini [Aporeto]



Virtualization, what is it?

Virtualization refers to the act of creating a virtual version of something

Wikipedia





Use cases



Type of hypervisors

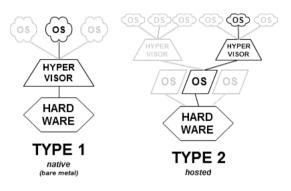


Figure: From wikipedia

Kind of virtualization

Full hardware virtualization

- OS is running unmodified
- Guest I/O are either
 - emulated
 - handled by virtualization-aware hardware

Para-virtualization

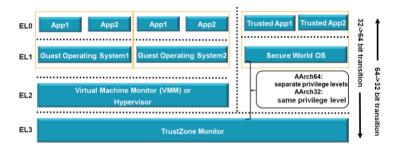
- OS is aware of the hypervisor
- Privilege instruction are replaced by hooks
- Devices (network, block...) are para-virtualized

The trend is a mix of both

- Use as much as possible hardware-assisted virtualization
- Devices (network, block...) para-virtualized or passthrough-ed
- Emulation very limited

ARM virtualization

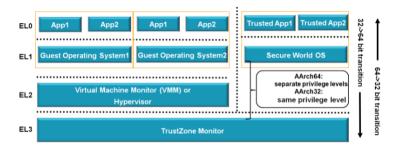
ARMv8-A Privilege Model



- Support both AArch32 and AArch64 execution modes
- 32-64bit inter-working limited to exception boundaries
- AArch64 always has a higher privilege than AArch32
- AArch64 state is a superset of AArch32 state



ARM virtualization



- Introduced with the latest version of ARMv7 architecture
- New hypervisor execution state
- Non-Secure world, higher privilege than ELI

Virtualization in a nutshell

- Second stage of memory translation
 - Adds an extra level of indirection between guests and physical memory
 - TLBs are tagged by Virtual Machine ID (VMID)
- Ability to trap access of most system registers
 - The hypervisor decides what it wants to trap
- Can handle IRQs, FIQs and asynchronous aborts
 - The guest doesn't see physical interrupts firing, for example
- Guests can call into EL2 mode (HVC instruction)
 - Allows para-virtualizated services
- Standard architecture peripherals are virtualization-aware
 - GIC and timer have specific features to help virtualization

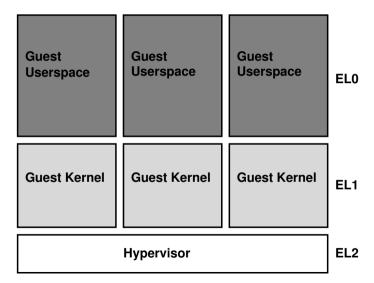


EL2: Not ELI++ (ARMv8.0-A)

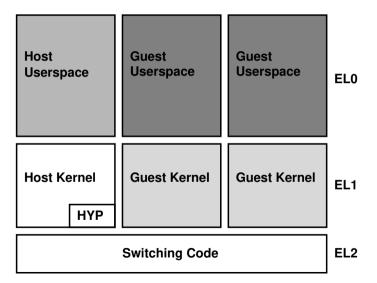
- EL2 is not a superset of NS-EL1
 - Orthogonal mode to ELI
 - Allows multiplexing of NS-EL1 guests on the hardware
- Own translation regime
 - Separate Stage-1 translation, no Stage-2 translation
- It would be difficult to run Linux in EL2
 - Requires too many changes to be practical
- EL2 could be used as a "world switch"
 - Between guests (barametal hypervisor/Type I)
 - Between host and guest (hosted hypervisor/Type II)
 This makes the host a form of specialized guest.



Hypervisor architecture - Type I



Hypervisor architecture - Type II



EL2 enhancement (ARMv8.I-A)

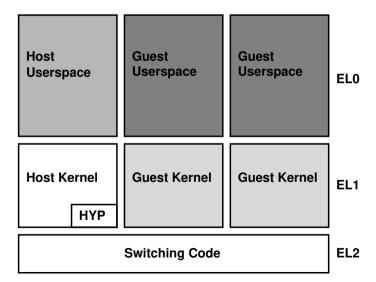
The Virtualization Host Extension (VHE) expands the capability of EL2:

- Designed to improve the support of the Type-2 hypervisors
- Allows the host OS to be run at EL2
- The host OS requires minimal changes to run at EL2
- User-space still runs at EL0
- Host has no software running at ELI
- AArch64 specific

EL2 becomes a strict superset of EL1

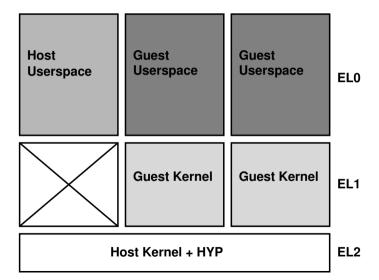


Hosted hypervisor architecture on platform without VHE





Hosted hypervisor architecture on platform with VHE





Nested Virtualization (ARMv8.3-A)

The Nested Virtualization extension allows an hypervisor in a VM.

- Unmodified guest hypervisor running in NS EL1
- Implementation of a host hypervisor required
 - Running at EL2
- AArch64 specific



Why using ARM virtualization

- Robust set of virtualization features
 - Not just about CPU virtualization
 - Covers the whole systems architecture
- Scalable architecture
 - Power to IoT-like devices ...
 - ... all the way to server-grade systems
- An architecture in motion:
 - ARMv8.I-A: https://goo.gl/Ox4thV
 - ARMv8.2-A: https://goo.gl/0Ns37U
 - ARMv8.3-A: https://goo.gl/CJv1n0

OpenSource Hypervisors









KVM

Kernel-based Virtual Machine

- First version of KVM was merged in Linux 2.6.20
 - AArch32 support merged in Linux 3.9
 - AArch64 support merged in Linux 3.11
- Source code available as GLP v2
- Hosted hypervisor

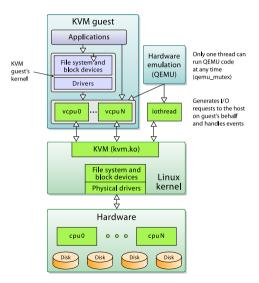


KVM virtual machine

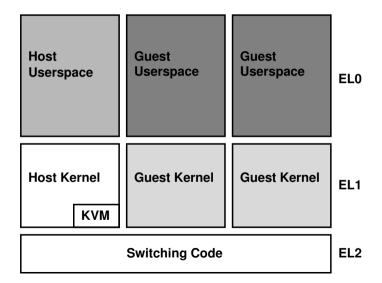
- Use of assisted hardware virtualization
- Devices are
 - emulated (QEMU)
 - para-virtualized (VIRTIO)



KVM architecture

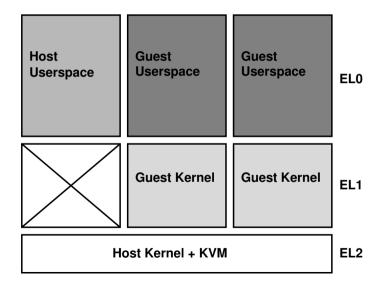


KVM architecture with ARMv8.0-A





KVM architecture with ARMv8.1-A



Resource management

- All CPUs are using the same scheduler
- guest vCPU is a task for the host OS
- Resource management can be done using cgroup
 - Standard way in Linux to control resources

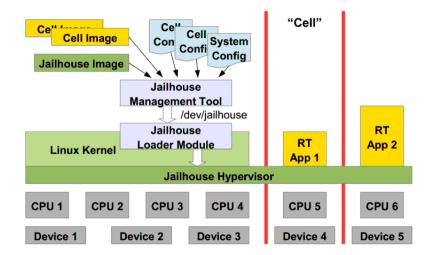


Jailhouse

Jailhouse

- Created at Siemens in 2013
- Partioning hypervisor
 - Type-I hypervisor
 - Linux will load Jailhouse
- Source code available as GPL v2
- Small code base: <10K lines</p>

Jailhouse architecture





Xen

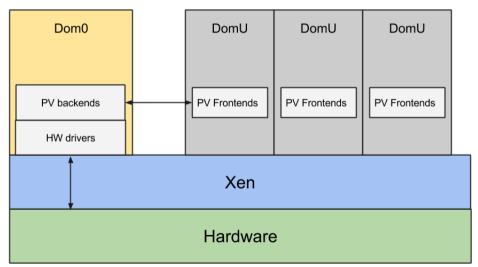
ARM

Xen

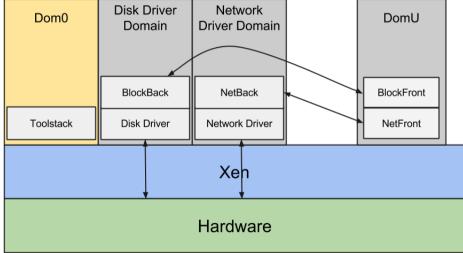
- First released in 2003
 - ARM officially supported since Xen 4.4
- Source code available as GPL v2
- Small code base: 30K
- Bare-metal hypervisor



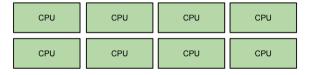
Xen architecture



Xen architecture - 2

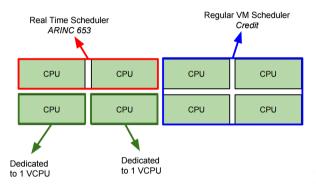


Xen schedulers





Xen schedulers





Summary



https://www.linux-kvm.org/



https://github.com/siemens/jailhouse



https://xenproject.org/



Questions?



The trademarks featured in this presentation are registered and/or unregistered trademarks of ARM limited (or its subsidiaries) in the EU and/or elsewhere. All rights reserved. All other marks featured may be trademarks of their respective owners.

Copyright © 2017 ARM Limited

© ARM 2017