

Applied Cryptography: Homework 5

(Deadline: 10:00am, 2020/10/21)

Justify your answers with calculations, proofs, and programs.

1. (10 points, question 4.8, page 132 of the textbook)

Suppose that $X = (x_1, \dots, x_n)$ and $X' = (x'_1, \dots, x'_n)$ are two sequences of n plaintext blocks. Define

$$\mathbf{same}(X, X') = \max\{j : x_i = x'_i \text{ for all } i \leq j\}.$$

Suppose X and X' are encrypted in CBC or CFB mode using the same key and the same IV. Show that it is easy for an adversary to compute $\mathbf{same}(X, X')$.

2. (20 points, question 4.10, page 133 of the textbook)

Suppose a sequence of plaintext blocks, $x_1 \dots x_n$, yields the ciphertext sequence $y_1 \dots y_n$. Suppose that one ciphertext block, say y_i , is transmitted incorrectly (i.e., some 1's are changed to 0's and vice versa). Show that the number of plaintext blocks that will be decrypted incorrectly is equal to one if ECB or OFB modes are used for encryption; and equal to two if CBC or CFB modes are used.

3. (20 points, question 5.1, page 178 of the textbook)

Define a toy hash function $h : (\mathbb{Z}_2)^7 \rightarrow (\mathbb{Z}_2)^4$ by the rule $h(x) = xA$ where all operations are modulo 2 and

$$A = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix}.$$

Find all preimages of $(0, 1, 0, 1)$.

4. (10 points, question 5.8, page 181 of the textbook)

Suppose that $f : \{0, 1\}^m \rightarrow \{0, 1\}^m$ is a preimage resistant bijection. Define $h : \{0, 1\}^{2m} \rightarrow \{0, 1\}^m$ as follows. Given $x \in \{0, 1\}^{2m}$, write

$$x = x' || x''$$

where $x', x'' \in \{0, 1\}^m$. Then define

$$h(x) = f(x' \oplus x'').$$

Prove that h is not second preimage resistant.