

CS101-Quiz9-Review

Key Points

- 1. Topological Sort
- 2. Greedy Algorithms

Topological Sort

- 1. An algorithm for ordering the vertices of a directed acyclic graph (DAG) in a linear ordering.
- 2. Time complexity: O(V+E)
- 3. Space complexity: O(V)

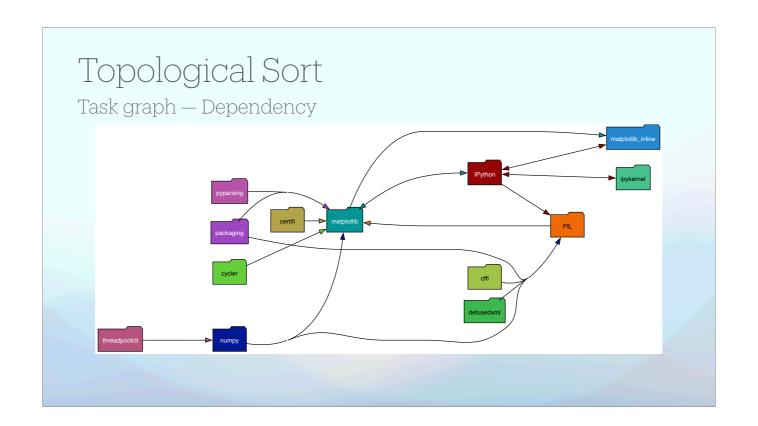
Topological Sort

Basic Theorems and Lemmas

- 1. A graph is a DAG if and only if it has a topological sorting
- 2. A DAG always has at least one vertex with in-degree zero.
- 3. Any sub-graph of a DAG is a DAG.

Topological Sort Curriculum scheduling Servery Servery

Topological Sort Task graph — Python Dask print2 print1 print3 print3 format3 print3 print



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- 1. Topological Sort
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Array Section

```
1 def ArraySection(A):
2    s = A[1]
3    g = 1
4    for i = [2, ..., n]:
5    s = s + A[i]
6    if s > M:
7    s = A[i]
8    g = g + 1
9    return g
```

4 2 4 5 1 2

Array Section — exchange argument

- 1. When all optimal solutions have the same size and differ only in their cost.
- 2. Justify with how you will replace some parts of the optimal solution.

Array Section — exchange argument

- 1. Define subproblem and compare
- 2. Exchange
- 3. Iterate

这里的compare说的是"我们在贪的东西"和"题面"要求的问题等价。 在部分问题里,这个东西可能非常难证;但是本题不存在这个问题。

Array Section — exchange argument

- 1. Sub-problem g(i) means minimizing the number of sections for $\langle a_1, \dots, a_i \rangle$, and then maximizing the start indices of all the sections.
- 2. Exchange
- 3. Iterate

我们的"子问题"的最大规模刚好等价于题面。

Array Section — exchange argument

- 1. Sub-problem g(i) means minimizing the number of sections for $\langle a_1, \dots, a_i \rangle$, and then maximizing the start indices of all the sections.
- 2. Exchange
- 3. Iterate

正常来讲这两步要分开,但是因为这个题iteration很显然,所以直接一个数学归纳做完。

Array Section — exchange argument

- 1. Sub-problem g(i) means minimizing the number of sections for $\langle a_1, \dots, a_i \rangle$, and then maximizing the start indices of all the sections.
- 2. Assume the start index of sections is $\langle s_1, \dots, s_k \rangle$. Assume we can make s_t smaller.
- Prove the new solution is not better than our solution (see (e.) in the answer book).
- 3. Iterate

我知道你想看具体怎么做,但是没空。不如看答案。

这里的重点在于要讲清楚,虽然对于大部分的贪心问题来讲,我们贪心出来的结果都不是唯一的正确答案,但是只要没有任何一种答案会比他好就可以。

Array Section — exchange argument

- 1. Sub-problem g(i) means minimizing the number of sections for $\langle a_1, \dots, a_i \rangle$, and then maximizing the start indices of all the sections.
- 2. Assume the start index of sections is $\langle s_1, \cdots, s_k \rangle$. Assume we can make s_t smaller.
 - Prove the new solution is not better than our solution (see (e.) in the answer book).
- 3. 显然成立

Array Section — exchange argument — Take-home message

- 1. The **exchange method** assumes that the "size" of the solution is the same for both greedy and optimal solutions.
- 2. Exchange is then used to show that it is optimal.
- 3. Mathematical induction is often used to show how this consistency holds as the problem size increases.

一般来讲,exchange经常会使用的场景下,第三条递推都比较显然。

某种意义上来讲,piazza的pdf说的两种方法其实差不多,其中exchange是一种比较常见的特例。

这里可以保证的是,我们在考试出的greedy证明难度一定是小于作业最后一题的;并且存在一些greedy算法可以这两种方法都不用,就能解决。

如果贪心贪出的结果不仅是正确的,而且还是唯一的;那么我们只需要证明他的唯一性,就可以说明正确性。在这个过程中,唯一性结论本身更强,但是往往更好证明。一会Quiz的题就是这样。