### ESE 605 Homework 3

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Mar. 25, 2022

### CVX\* warmup:

```
# Import packages.
import cvxpy as cp
import numpy as np
# Generate a random feasible SOCP.
m = 3
n = 10
p = 5
n_i = 5
np.random.seed(2)
f = np.random.randn(n)
A = []
b = []
c = []
d = []
x0 = np.random.randn(n)
for i in range(m):
    A.append(np.random.randn(n_i, n))
    b.append(np.random.randn(n_i))
    c.append(np.random.randn(n))
    d.append(np.linalg.norm(A[i] @ x0 + b, 2) - c[i].T @ x0)
F = np.random.randn(p, n)
g = F @ x0
# Define and solve the CVXPY problem.
x = cp.Variable(n)
# We use cp.SOC(t, x) to create the SOC constraint ||x||_2 \le t.
soc_constraints = [
      cp.SOC(c[i].T @ x + d[i], A[i] @ x + b[i]) for i in range(m)
prob = cp.Problem(cp.Minimize(f.T@x),
                  soc_constraints + [F @ x == g])
```

```
prob.solve()

# Print result.
print("The optimal value is", prob.value)
print("A solution x is")
print(x.value)
for i in range(m):
    print("SOC constraint %i dual variable solution" % i)
    print(soc_constraints[i].dual_value)

The optimal value is -9.582695716266176
A solution x is
```

```
[ 1.40303325 2.4194569 1.69146656 -0.26922215 1.30825472 -0.70834842 0.19313706 1.64153496 0.47698583 0.66581033]

SOC constraint 0 dual variable solution
[ 0.61662526 0.35370661 -0.02327185 0.04253095 0.06243588 0.49886837]

SOC constraint 1 dual variable solution
[ 0.35283078 -0.14301082 0.16539699 -0.22027817 0.15440264 0.06571645]

SOC constraint 2 dual variable solution
[ 0.86510445 -0.114638 -0.449291 0.37810251 -0.6144058 -0.11377797]
```

## Problems from Boyd & Vandenberghe: 4.26.

$$\left\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \right\|_{2}^{2} - (y+z)^{2} = 4x^{T}x + y^{2} + z^{2} + 2yz - (y+z)^{2} = 4(x^{T}x - yz)$$

Since 
$$\left\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \right\|_2 \ge 0$$
, and  $y+z \ge 0$ , then  $\left\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \right\|_2^2 \le (y+z)^2$  is equivalent to  $\left\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \right\|_2 \le y+z$ , and is equivalent to  $x^Tx \le yz$ 

(a) The problem is equivalent to:

min. 
$$\sum_{i=1}^{n} t_i$$
s.t. 
$$t_i \ge \frac{1}{a_i^T x - b_i}$$

$$t_i \ge 0$$

Writing the constrains in the form of:

$$x^T x \le yz, \quad y \ge 0, \quad z \ge 0$$

where

$$x = 1, y = t_i, z = a_i^T x - b_i$$

the problem can be casted as SOCP:

min. 
$$1^T t$$
  
s.t.  $\left\| \begin{bmatrix} 2 \\ t_i - a_i^T x + b_i \end{bmatrix} \right\|_2 \le t_i + a_i^T x - b_i$   
 $t_i \ge 0$   
 $a_i^T x - b_i \ge 0$ 

(b) solve only for m = 6 case.

The objective function is equivalent to:

$$\max. \quad \Pi_{i=1}^6 a_i^T x + b_i$$

According to the hint on Piazza, without loss of generality, we can define  $a_7 = a_8 = 0$  and  $b_7 = b_8 = 1$  to make this a m = 8 case. We can rewrite the problem as:

max. 
$$t_1 t_2 t_3 t_4$$
  
s.t.  $Ax \succeq b$   
 $t_1^2 \le (a_1^T x - b_1)(a_2^T x - b_2)$   
 $t_2^2 \le (a_3^T x - b_3)(a_4^T x - b_4)$   
 $t_3^2 \le (a_5^T x - b_5)(a_6^T x - b_6)$   
 $t_4^2 \le (a_7^T x - b_7)(a_8^T x - b_8)$ 

which is equivalent to:

max. 
$$p_1p_2$$
  
s.t.  $Ax \succeq b$   
 $t_1^2 \le (a_1^T x - b_1)(a_2^T x - b_2)$   
 $t_2^2 \le (a_3^T x - b_3)(a_4^T x - b_4)$   
 $t_3^2 \le (a_5^T x - b_5)(a_6^T x - b_6)$   
 $t_4^2 \le (a_7^T x - b_7)(a_8^T x - b_8)$   
 $p_1^2 \le t_1 t_2$   
 $p_2^2 \le t_3 t_4$ 

which is equivalent to:

max. 
$$k$$
  
s.t.  $Ax \succeq b$   
 $t_1^2 \le (a_1^T x - b_1)(a_2^T x - b_2)$   
 $t_2^2 \le (a_3^T x - b_3)(a_4^T x - b_4)$   
 $t_3^2 \le (a_5^T x - b_5)(a_6^T x - b_6)$   
 $t_4^2 \le (a_7^T x - b_7)(a_8^T x - b_8)$   
 $p_1^2 \le t_1 t_2$   
 $p_2^2 \le t_3 t_4$   
 $k^2 \le p_1 p_2$ 

so the problem can be casted as SOCP:

min. 
$$-k$$
s.t. 
$$\left\| \begin{bmatrix} 2t_1 \\ a_1^T x - b_1 - a_2^T x + b_2 \end{bmatrix} \right\|_2 \le a_1^T x - b_1 + a_2^T x - b_2$$

$$\left\| \begin{bmatrix} 2t_2 \\ a_3^T x - b_3 - a_4^T x + b_4 \end{bmatrix} \right\|_2 \le a_3^T x - b_3 + a_4^T x - b_4$$

$$\left\| \begin{bmatrix} 2t_3 \\ a_5^T x - b_5 - a_6^T x + b_6 \end{bmatrix} \right\|_2 \le a_5^T x - b_5 + a_6^T x - b_6$$

$$\left\| \begin{bmatrix} 2t_4 \\ a_7^T x - b_7 - a_8^T x + b_8 \end{bmatrix} \right\|_2 \le a_7^T x - b_7 + a_8^T x - b_8$$

$$\left\| \begin{bmatrix} 2p_1 \\ t_1 - t_2 \end{bmatrix} \right\|_2 \le t_1 + t_2$$

$$\left\| \begin{bmatrix} 2p_2 \\ t_3 - t_4 \end{bmatrix} \right\|_2 \le t_3 + t_4$$

$$\left\| \begin{bmatrix} 2k \\ p_1 - p_2 \end{bmatrix} \right\|_2 \le p_1 + p_2$$

$$ax_i - b_i > 0, t_i > 0, p_i > 0$$

### 4.43.

(a) The original problem is:

min. 
$$\lambda_1(x)$$

We know that  $\lambda I - A(x) \succeq 0$  iff  $\lambda - \lambda_1(x) \geq 0$ . So the problem can be written in the form of SDP as:

min. 
$$\lambda$$
  
s.t.  $A(x) \leq \lambda I$ 

(b) The original problem is:

min. 
$$\lambda_1(x) - \lambda_m(x)$$

We know that  $\gamma I - A(x) \leq 0$  iff  $\gamma - \lambda_m(x) \leq 0$ . So the problem can be written in the form of SDP as:

min. 
$$\lambda - \gamma$$
  
s.t.  $A(x) \leq \lambda I$   
 $A(x) \succeq \gamma I$ 

(c) From (a) and (b), we know that the problem:

min. 
$$\lambda_1(x)/\lambda_m(x)$$
  
s.t.  $A(x) \succ 0$ 

is equivalent to:

min. 
$$\lambda/\gamma$$
  
s.t.  $0 \prec \gamma I \prec A(x) \prec \lambda I$ 

Change variables to  $y = x/\gamma, t = \lambda/\gamma, s = 1/\gamma$ .

If  $\gamma > 0$ , without loss of generality, the problem becomes:

min. 
$$t$$
  
s.t.  $I \leq sA(0) + y_1A(1) + \dots + y_nA(n) \leq tI$   
 $s > 0$ 

We want to show the above form holds when  $\gamma = 0$ .

If  $\gamma = 0$ , the constraint becomes  $I \leq y_1 A(1) + \cdots + y_n A(n) \leq tI$ . We know that the solution of the former problem is always feasible for the SDP problem, so that  $p_1^* \geq p_2^*$ . We want to show that  $p_2^* \geq p_1^*$ .

Since  $A(x) \succ 0$  for at least one x, we can construct  $A(\tau y) \succeq A0 + \tau I \succ 0$ . Then,  $\lambda_1(\tau y) \leq \lambda_1(0) + t\tau, \lambda_m(\tau y) \geq \lambda_m(0) + \tau$ , when  $\tau$  is sufficiently large,

$$\kappa(A(x)) = \frac{\lambda_1(x)}{\lambda_m(x)} \le \frac{\lambda_1(0) + t\tau}{\lambda_m(0) + \tau}$$

so that  $p_2^* \ge p_1^*$ .

Now we have proved the problem can be written as SDP.

4.45.

(a) If p can be expressed as a positive semidefinite quadratic form  $p = f^T V f$ . Since  $V \in S_+^s$ , we can write  $V = W W^T$ , where W is  $s \times r$ , and

$$p = (W^T f)^T W^T f = \sum_{i=1}^r (W_i f)^2$$

where degree of  $W_i f$  is no more than k. So p is SOS.

If p is SOS, then p have the form  $p = \sum_{i=1}^{r} q_i(x)^2$ . We can separate the monomials and coefficients and have,

$$p = \sum_{i=1}^{r} q_i(x)^2 = \sum_{i=1}^{r} (W_i f)^2 = (W^T f)^T W^T f$$

Set  $V = WW^T$ , then p can be expressed as a positive semidefinite quadratic form  $p = f^T V f$ , and  $V \in S^s_+$ .

(b) We can spread out the condition,

$$p = f^T V f = \sum_{i,j=1}^s V_{ij} f_i f_j$$

so that p is a set of linear equality constraints relating the coefficients of p and the matrix V. Since p is SOS requires all the coefficients to be non-negative,  $V \succeq 0$ .

(c) Set  $f = [1 \ x_1 \ x_2 \ x_1^2 \ x_2^2 \ x_1 x_2]^T$ , then

$$f^T f = \begin{bmatrix} 1 & x_1 & x_2 & x_1^2 & x_2^2 & x_1 x_2 \\ x_1 & x_1^2 & x_1 x_2 & x_1^3 & x_1 x_2^2 & x_1^2 x_2 \\ x_2 & x_1 x_2 & x_2^2 & x_1^2 x_2 & x_2^3 & x_1 x_2^2 \\ x_1^2 & x_1^3 & x_1^2 x_2 & x_1^4 & x_1^2 x_2^2 & x_1^3 x_2 \\ x_2^2 & x_1 x_2^2 & x_2^3 & x_1^2 x_2^2 & x_2^4 & x_1 x_2^3 \\ x_1 x_2 & x_1^2 x_2 & x_1 x_2^2 & x_1^3 x_2 & x_1 x_2^3 & x_1^2 x_2^2 \end{bmatrix}$$

We can write p as:

$$p = a_1 + a_2 x_1 + a_3 x_2 + a_4 x_1^2 + a_5 x_1 x_2 + a_6 x_2^2 + a_7 x_1^3 + a_8 x_1^2 x_2$$

$$+ a_9 x_1 x_2^2 + a_{10} x_2^3 + a_{11} x_1^4 + a_{12} x_1^3 x_2 + a_{13} x_1^2 x_2^2 + a_{14} x_1 x_2^3 + a_{15} x_2^4$$

Since  $p = f^T V f = \sum_{i,j=1}^s V_{ij} f_i f_j$ ,

$$a_1 = V_{11}, \quad a_2 = V_{12} + V_{21}, \quad a_3 = V_{13} + V_{31}$$
 
$$a_4 = V_{14} + V_{22} + V_{41}, \quad a_5 = V_{16} + V_{23} + V_{32} + V_{61}, \quad a_6 = V_{15} + V_{33} + V_{51}$$
 
$$a_7 = V_{24} + V_{42}, \quad a_8 = V_{26} + V_{34} + V_{43} + V_{62}, \quad a_9 = V_{25} + V_{36} + V_{52} + V_{63}$$
 
$$a_{10} = V_{35} + V_{53}, \quad a_{11} = V_{44}, \quad a_{12} = V_{46} + V_{64}$$
 
$$a_{13} = V_{45} + V_{54} + V_{66}, \quad a_{14} = V_{56} + V_{65}, \quad a_{15} = V_{55}$$

We have worked out the LMI conditions for SOS explicitly.

#### 4.59.

(a) If  $f_i$  are convex in x for each u.

$$\mathbf{E}_{u}f_{i}(x,u) = \int f_{i}(x,u)p(u)du$$

Since p(u) > 0,  $\mathbf{E}_u f_i(x, u)$  is convex. So this stochastic optimization problem is convex.

(b) If  $f_i(x, u)$  is convex in x for each  $u \in U$ , then

$$\sup_{u \in U} f_i(x, u)$$

is convex. So this stochastic optimization problem is convex.

(c) Set up the stochastic optimization problems:

min. 
$$\sum_{j=i}^{N} f_0(x, u_j) p_j$$
  
s.t. 
$$\sum_{j=i}^{N} f_i(x, u_j) p_j \le 0, \quad i = 1, \dots, m$$

Set up the worst-case optimization problems:

min. 
$$\max_{j} f_0(x, u_j)$$
  
s.t.  $\max_{j} f_i(x, u_j) \le 0, \quad i = 1, \dots, m, \ j = 1, \dots, N$ 

# Disciplined Convex Programming: A4.3.

(a) 
$$cp.inv_pos(x) + cp.inv_pos(y) \le 1, x >= 0, y >= 0$$

(b) 
$$x \ge cp.inv_pos(y), x \ge 0, y \ge 0$$

(c) cp.quad\_over\_lin(x+y, cp.sqrt(y)) 
$$\leq$$
 x - y + 5

(d) 
$$x + z \le 1 + cp.geo_mean(cp.hstack([x - cp.quad_over_lin(z,y), y])), x >= 0, y >= 0$$

Codes of small problems are listed in Appendix A.

## Problems from Boyd & Vandenberghe: 5.3.

Lagrangian is:

$$L(x,\lambda) = c^T x + \lambda f(x)$$

Dual function is:

$$g(\lambda) = \inf_{x} L(x, \lambda)$$

$$= \lambda \inf_{x} (c/\lambda)^{T} x + f(x)$$

$$= -\lambda \sup_{x} -(c/\lambda)^{T} x - f(x)$$

$$= -\lambda f^{*}(-c/\lambda)$$

Dual problem is:

max. 
$$-\lambda f^*(-c/\lambda)$$
  
s.t.  $\lambda \ge 0$ 

Since conjugate function  $f^*$  is convex, the objective is maximize over a concave function. Also, constraint is convex. So this is a convex problem.

#### **5.4.**

(a) Lagrangian is:

$$L(x,\lambda) = c^T x + \lambda (w^T A x - w^T b)$$

Dual function is:

$$g(\lambda) = \inf_{x} L(x, \lambda)$$

$$= \inf_{x} -\lambda w^{T} b + (c + \lambda A^{T} w)^{T} x$$

$$= \begin{cases} -\lambda w^{T} b & c + \lambda A^{T} w = 0 \\ -\infty & \text{otherwise} \end{cases}$$

The minimum value of  $c^T x$  is:

$$\min(c^T x) = \begin{cases} -\lambda w^T b & c + \lambda A^T w = 0\\ -\infty & \text{otherwise} \end{cases}$$

(b) The dual problem is:

$$\max. \quad -\lambda w^T b$$
s.t. 
$$c + \lambda A^T w = 0$$

$$\lambda \ge 0$$

$$w \ge 0$$

(c) Set  $\lambda w = \lambda$ . In this situation, qualify the new  $\lambda \succeq 0$  is equivalent to qualify old  $\lambda \geq 0, \ w \succeq 0$ . Then the problem can be written as:

max. 
$$-b^T \lambda$$
  
s.t.  $A^T \lambda + c = 0$   
 $\lambda \succeq 0$ 

5.9.

(a) Consider the following matrix.

$$\begin{bmatrix} \sum_{k=1}^{m} a_k a_k^T & a_i \\ a_i^T & 1 \end{bmatrix}$$

$$= \begin{bmatrix} I & a_i \\ 0 & I \end{bmatrix} \begin{bmatrix} \sum_{k=1, k \neq i}^{m} a_k a_k^T & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} I & a_i \\ 0 & I \end{bmatrix}^T$$

$$\succeq 0$$

The Schur complements of a positive semi-definite matrix is positive semi-definite,

$$S = 1 - a_i^T (\sum_{k=1}^m a_k a_k^T)^{-1} a_i$$
  
= 1 - a\_i^T X^{-1} a\_i  
> 0

So that  $a_i^T X^{-1} a_i \leq 1$ . We have shown the problem is feasible.

(b) When  $\lambda = t\mathbf{1} \in R^m$ , the dual problem is:

maximize 
$$\log \det \left(\sum_{i=1}^m a_i a_i^T\right) + \log t^n - mt + n$$
  
subject to  $t > 0$ 

Since,

$$\frac{d}{dt}(\log t^n - mt) = n/t - m$$

The optimal value of t is n/m. Put this back into the dual problem,

$$\log \det \left( \sum_{i=1}^{m} a_i a_i^T \right) + \log t^n - mt + n$$

$$= \log \det \left( \sum_{i=1}^{m} a_i a_i^T \right) + n \log(n/m)$$

$$= \log \det \left( \sum_{i=1}^{m} a_i a_i^T \right) - n \log(m/n)$$

The value of the primal problem is:

$$\log \det \left( \sum_{i=1}^{m} a_i a_i^T \right)$$

The volume of a ellipsoid is  $(\det X^{-1})^{1/2}$ .

The values of the primal and the dual problem are:

$$Y_1 = \log(\det(X^{-1}))$$
  
 $Y_2 = \log(\det(X^{-1})) - n\log(m/n)$ 

We know the gap between the dual and primal is no more than  $n \log(m/n)$ , and the volume of a ellipsoid is  $\exp(Y/2)$ . So the difference between the ellipsoid  $\{u|u^TX_{sim}u \leq 1\}$  and the volume of the minimum volume ellipsoid is no more than a factor  $(m/n)^{n/2}$ .

# Problem from Additional Exercises: A5.3.

Lagrangian is:

$$L(x, \nu, \lambda) = \sum_{k=1}^{n} x_k \log(x_k/y_k) + \nu^T (Ax - b) + \lambda (1^T x - 1)$$

Dual function is:

$$g(\nu, \lambda) = \inf_{x} L(x, \nu, \lambda)$$
$$= \inf_{x} \sum_{k=1}^{n} x_k \log(x_k/y_k) + \nu^T (Ax - b) + \lambda (1^T x - 1)$$

Try to solve this by taking the derivative of L:

$$\frac{\partial}{\partial x_k} \sum_{k=1}^n x_k \log(x_k/y_k) + \nu^T (Ax - b) + \lambda (1^T x - 1)$$
$$= \log x_k + 1 - \log y_k + \nu^T a_k + \lambda$$

The optimal value for  $x_k$  is:

$$x_k = y_k e^{-(1+\nu^T a_k + \lambda)}$$

Put back into the dual function:

$$-\nu^{T}b - \log \sum_{k=1}^{n} y_{k} e^{a_{k}^{T}\nu} + 1 = -\sum_{k=1}^{n} (1 + \nu^{T}a_{k} + \lambda)x_{k} + \nu^{T}(Ax - b) + \lambda(1^{T}x - 1)$$
$$= -\sum_{k=1}^{n} y_{k} e^{-(1 + \nu^{T}a_{k} + \lambda)} - \nu^{T}b - \lambda$$

Taking the derivative of  $g(\nu, \lambda)$  by  $\lambda$ ,

$$\frac{\partial}{\partial \lambda} = \sum_{k=1}^{n} y_k e^{-(1+\nu^T a_k + \lambda)} - 1$$

So that

$$\lambda = \log \sum_{k=1}^{n} y_k e^{-a_k^T \nu} - 1$$

Set  $\nu = -z$ . Then,

$$g(\nu, \lambda) = -\sum_{k=1}^{n} y_k e^{-1} \sum_{k=1}^{n} y_k e^{-\nu^T a_k} \sum_{k=1}^{n} y_k e^{-\lambda} - \nu^T b - \log \sum_{k=1}^{n} y_k e^{a_k^T \nu} + 1$$

$$= -1 - \nu^T b - \log \sum_{k=1}^{n} y_k e^{a_k^T z} + 1$$

$$= b^T z - \log \sum_{k=1}^{n} y_k e^{a_k^T z}$$

### A Appendix: Codes for A4.3

```
import cvxpy as cp
# A4.3(a)
x, y = cp.Variable(), cp.Variable()
objective = cp.Minimize(x+y)
constraints = [cp.inv_pos(x) + cp.inv_pos(y) \le 1, x \ge 0, y \ge 0]
prob = cp.Problem(objective, constraints)
prob.solve()
print("status:", prob.status)
print("optimal value", prob.value)
print("optimal var", x.value, y.value)
     status: optimal
     optimal value 3.9999999942618447
     optimal var 1.999999997130576 1.9999999971312687
# A4.3(b)
x, y = cp.Variable(), cp.Variable()
objective = cp.Minimize(x+y)
constraints = [x \ge cp.inv_pos(y), x \ge 0, y \ge 0]
prob = cp.Problem(objective, constraints)
prob.solve()
print("status:", prob.status)
print("optimal value", prob.value)
print("optimal var", x.value, y.value)

    status: optimal

     optimal value 1.9999999999988524
     optimal var 1.0000014640195074 0.9999985350793449
# A4.3(c)
x, y = cp.Variable(), cp.Variable()
objective = cp.Minimize(x+y)
constraints = [cp.quad_over_lin(x+y, cp.sqrt(y)) \le x - y + 5]
prob = cp.Problem(objective, constraints)
prob.solve()
print("status:", prob.status)
print("optimal value", prob.value)
print("optimal var", x.value, y.value)
     status: optimal
     optimal value -1.3715395933070145
     optimal var -1.9762812107395356 0.6047416174325209
# · A4.3(d)
x, ·y, ·z·=·cp.Variable(), ·cp.Variable(), ·cp.Variable()
objective ·= ·cp.Minimize(x+y+z)
constraints \cdot = \cdot [x \cdot + \cdot z \cdot < = \cdot 1 \cdot + \cdot cp. geo\_mean(cp.hstack([x \cdot - \cdot cp.quad\_over\_lin(z,y), \cdot y])), \cdot x \cdot (x \cdot - \cdot cp.quad\_over\_lin(z,y), \cdot y]))
```

```
prob:=:cp.Problem(objective,:constraints)
prob.solve()
print("status:",:prob.status)
print("optimal:value",:prob.value)
print("optimal:var",:x.value,:y.value,:z.value)
```

status: optimal optimal value 1.1058302763345625e-10 optimal var -1.068114553406185e-10 5.684880922432107e-10 -3.510936092691359e-10

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