

Roadmap

C:

```
car *c = malloc(sizeof(car));
c->miles = 100;
c->gals = 17;
float mpg = get_mpg(c);
free(c);
```

Java:

```
Car c = new Car();
c.setMiles(100);
c.setGals(17);
float mpg =
    c.getMPG();
```

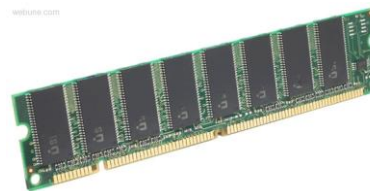
Assembly
language:

```
get_mpg:
    pushq    %rbp
    movq     %rsp, %rbp
    ...
    popq     %rbp
    ret
```

Machine
code:

```
0111010000011000
100011010000010000000010
1000100111000010
110000011111101000011111
```

Computer
system:



Arrays

Memory & data
Integers & floats
Machine code & C
x86 assembly
Procedures & stacks
Arrays & structs
Memory & caches
Processes
Virtual memory
Memory allocation
Java vs. C

OS:



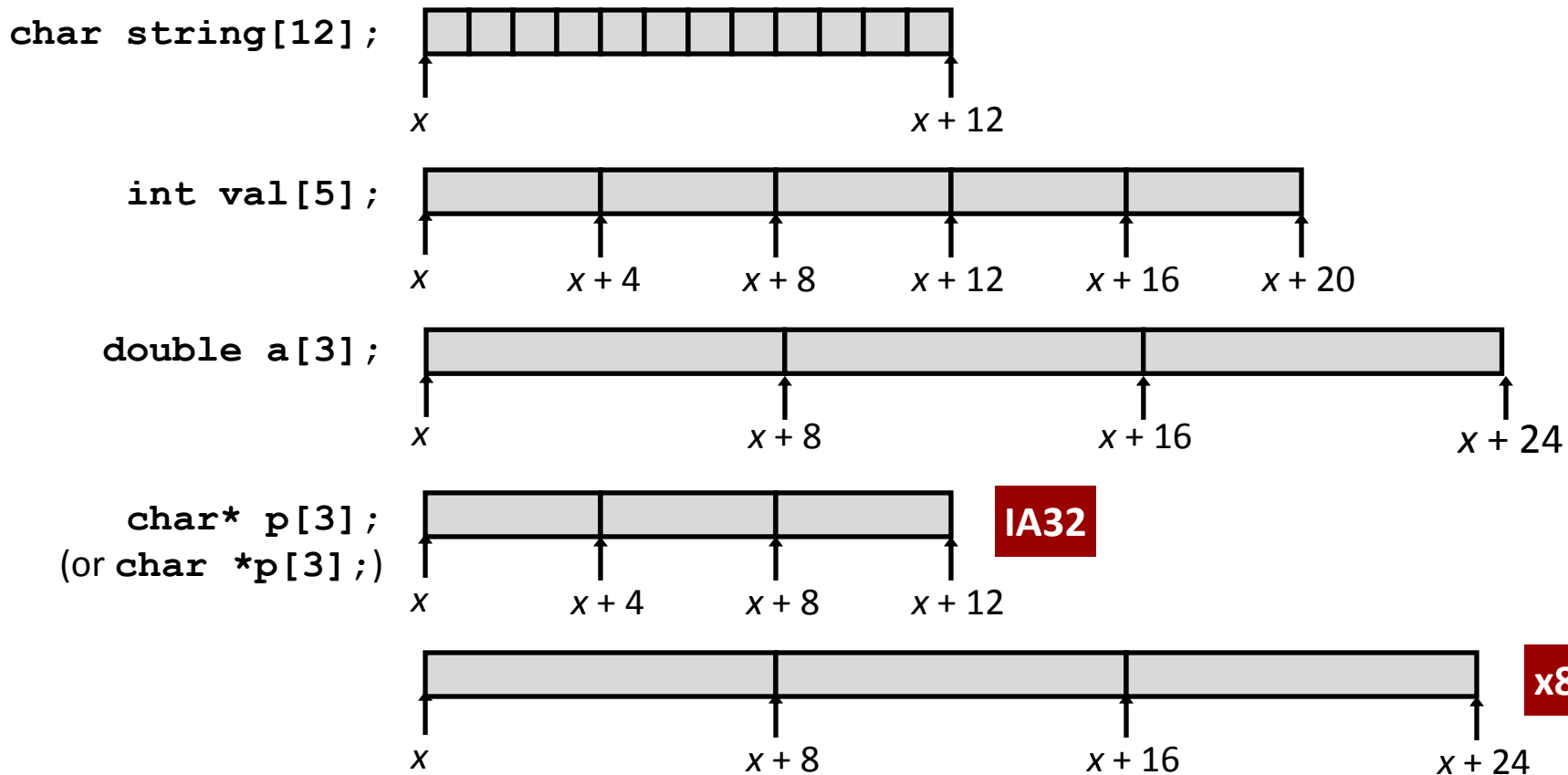
Section 5: Arrays & Other Data Structures

- Array allocation and access in memory
- Multi-dimensional or nested arrays
- Multi-level arrays
- Other structures in memory
- Data structures and alignment

Array Allocation

■ Basic Principle

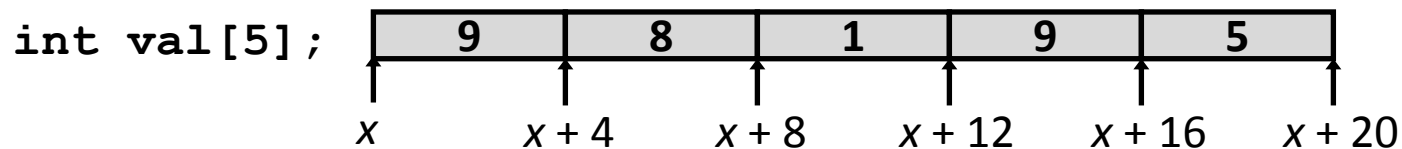
- $T \ A[N];$
- Array of data type T and length N
- *Contiguously* allocated region of $N * \text{sizeof}(T)$ bytes



Array Access

■ Basic Principle

- $T \ A[N];$
- Array of data type T and length N
- Identifier A can be used as a pointer to array element 0: Type T^*



■ Reference Type Value

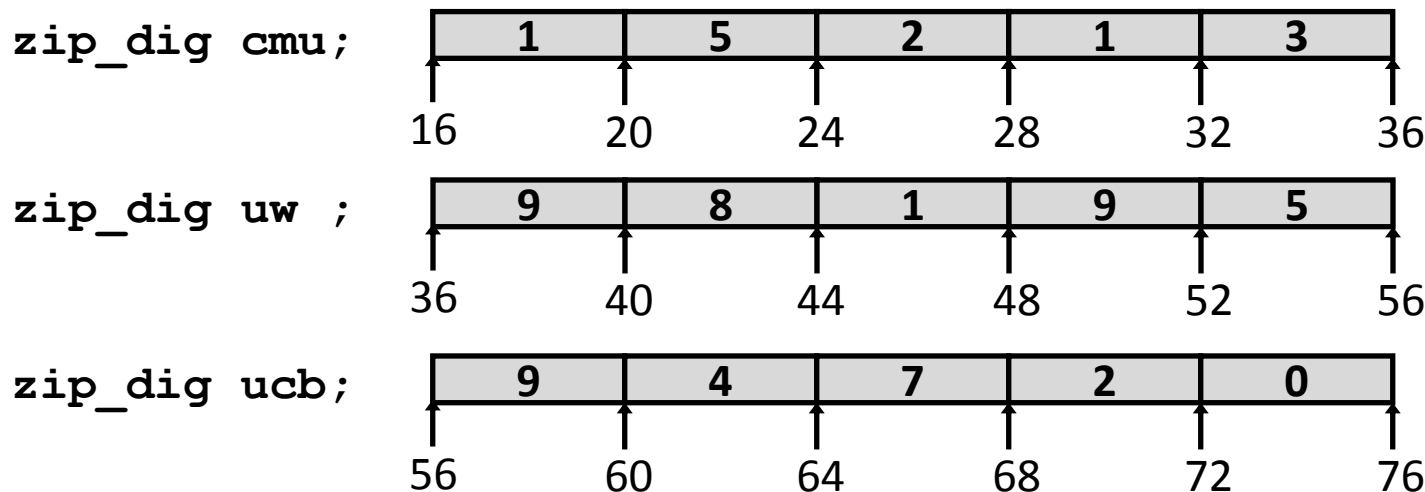
- `val[4]` `int` 5
- `val` `int *` x
- `val+1` `int *` $x + 4$
- `&val[2]` `int *` $x + 8$
- `val[5]` `int` ?? (whatever is in memory at address $x + 20$)
- `*(val+1)` `int` 8
- `val + i` `int *` $x + 4*i$

Array Example

```
typedef int zip_dig[5];  
  
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig uw  = { 9, 8, 1, 9, 5 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

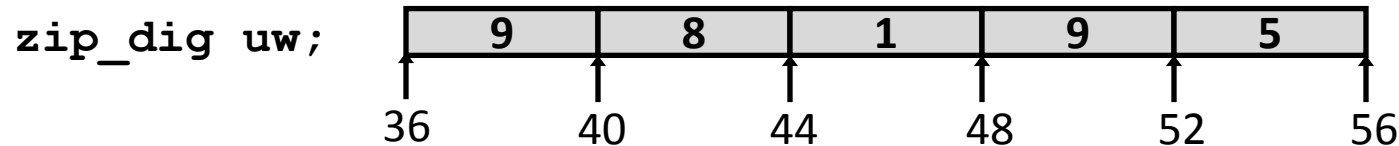
Array Example

```
typedef int zip_dig[5];  
  
zip_dig cmu = { 1, 5, 2, 1, 3 };  
zip_dig uw  = { 9, 8, 1, 9, 5 };  
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration “zip_dig uw” equivalent to “int uw[5]”
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example



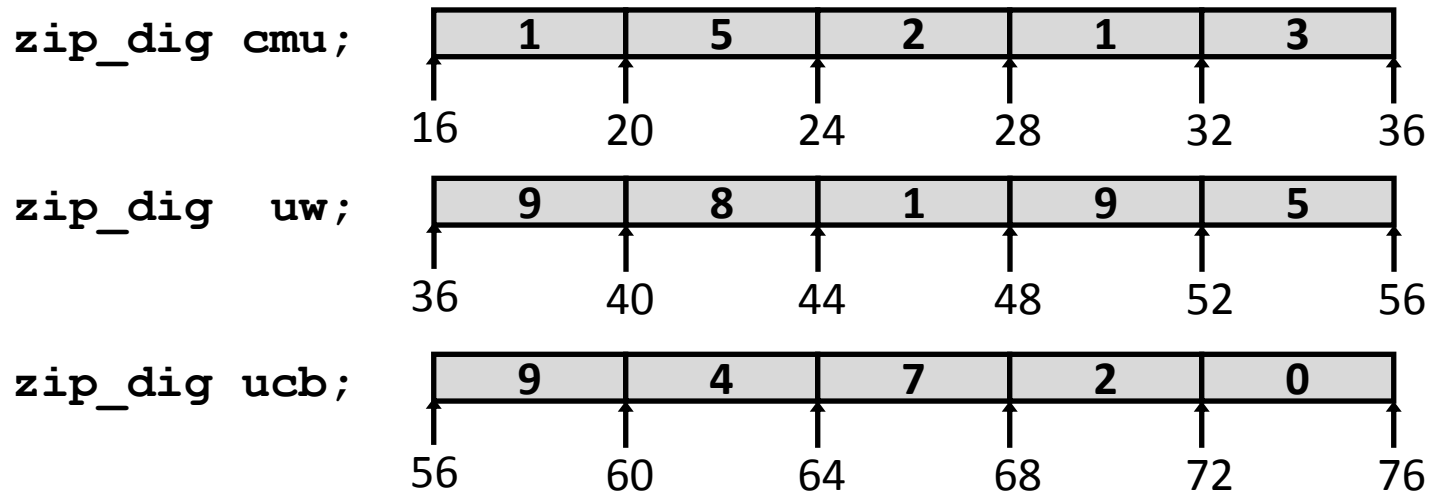
```
int get_digit
(zip_dig z, int dig)
{
    return z[dig];
}
```

IA32

```
# %edx = z
# %eax = dig
movl (%edx,%eax,4),%eax # z[dig]
```

- Register `%edx` contains starting address of array
- Register `%eax` contains array index
- Desired digit at $4 * \%eax + \%edx$
- Use memory reference $(\%edx, \%eax, 4)$

Referencing Examples



Reference	Address	Value	Guaranteed?
<code>uw[3]</code>	$36 + 4 * 3 = 48$	9	Yes
<code>uw[6]</code>	$36 + 4 * 6 = 60$	4	No
<code>uw[-1]</code>	$36 + 4 * -1 = 32$	3	No
<code>cmu[15]</code>	$16 + 4 * 15 = 76$??	No

- No bounds checking
- Location of each separate array in memory is not guaranteed

Array Loop Example

```
int zd2int(zip_dig z)
{
    int i;
    int zi = 0;
    for (i = 0; i < 5; i++) {
        zi = 10 * zi + z[i];
    }
    return zi;
}
```

Array Loop Example

■ Original

```
int zd2int(zip_dig z)
{
    int i;
    int zi = 0;
    for (i = 0; i < 5; i++) {
        zi = 10 * zi + z[i];
    }
    return zi;
}
```

■ Transformed

- Eliminate loop variable **i**, use pointer **zend** instead
- Convert array code to pointer code
 - Pointer arithmetic on **z**
- Express in do-while form (no test at entrance)

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while (z <= zend);
    return zi;
}
```

Array Loop Implementation (IA32)

■ Registers

```
%ecx  z
%eax  zi
%ebx  zend
```

■ Computations

- $10 * zi + *z$ implemented as $*z + 2 * (5 * zi)$
- $z++$ increments by 4

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while(z <= zend);
    return zi;
}
```

```
# %ecx = z
xorl %eax,%eax           # zi = 0
leal 16(%ecx),%ebx        # zend = z+4
.L59:
leal (%eax,%eax,4),%edx   # zi + 4*zi = 5*zi
movl (%ecx),%eax         # *z
addl $4,%ecx             # z++
leal (%eax,%edx,2),%eax   # zi = *z + 2*(5*zi)
cmpl %ebx,%ecx          # z : zend
jle .L59               # if <= goto loop
```