

Section 5: Procedures & Stacks

- Stacks in memory and stack operations
- The stack used to keep track of procedure calls
- Return addresses and return values
- Stack-based languages
- The Linux stack frame
- Passing arguments on the stack
- Allocating local variables on the stack
- Register-saving conventions
- Procedures and stacks on x64 architecture

Register Saving Conventions

■ When procedure `yoo` calls `who`:

- `yoo` is the *caller*
- `who` is the *callee*

■ Can a register be used for temporary storage?

`yoo:`

```
    . . .  
    movl $12345, %edx  
    call who  
    addl %edx, %eax  
    . . .  
    ret
```

`who:`

```
    . . .  
    movl 8(%ebp), %edx  
    addl $98195, %edx  
    . . .  
    ret
```

- Contents of register `%edx` overwritten by `who`

Register Saving Conventions

- When procedure `yoo` calls `who`:
 - `yoo` is the *caller*
 - `who` is the *callee*
- Can a register be used for temporary storage?
- Conventions
 - “*Caller Save*”
 - Caller saves temporary values in its frame before calling
 - “*Callee Save*”
 - Callee saves temporary values in its frame before using

IA32/Linux Register Usage

■ **%eax, %edx, %ecx**

- Caller saves prior to call if values are used later

■ **%eax**

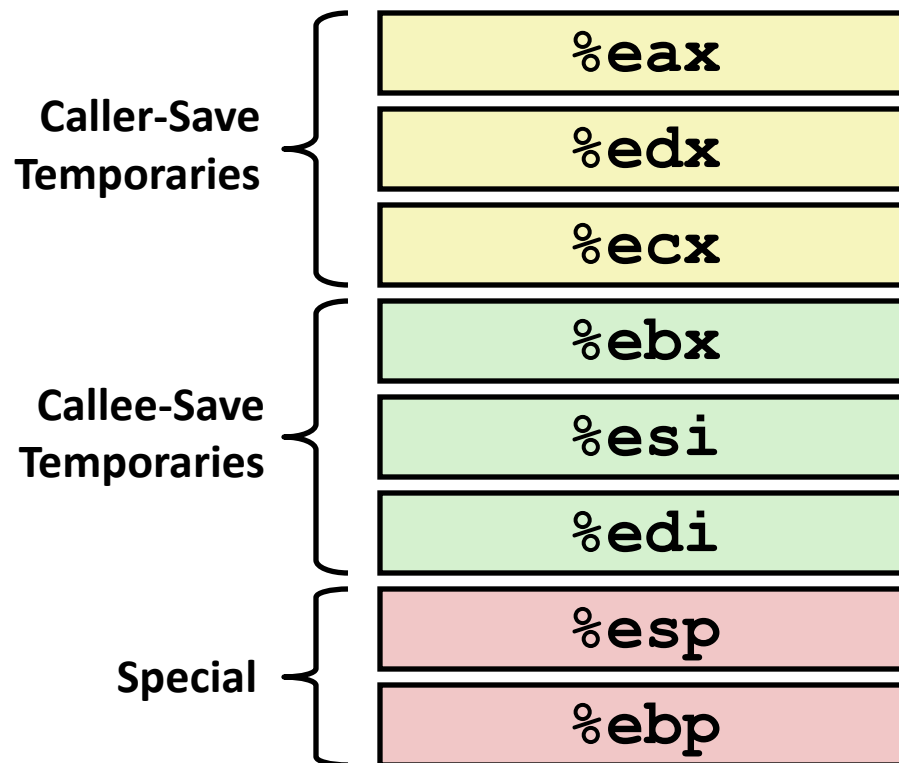
- also used to return integer value

■ **%ebx, %esi, %edi**

- Callee saves if wants to use them

■ **%esp, %ebp**

- special form of callee save – restored to original values upon exit from procedure



Example: Pointers to Local Variables

Recursive Procedure

```
void s_helper
(int x, int *accum)
{
    if (x <= 1)
        return;
    else {
        int z = *accum * x;
        *accum = z;
        s_helper (x-1, accum);
    }
}
```

Top-Level Call

```
int sfact(int x)
{
    int val = 1;
    s_helper(x, &val);
    return val;
}
```

- Pass pointer to update location

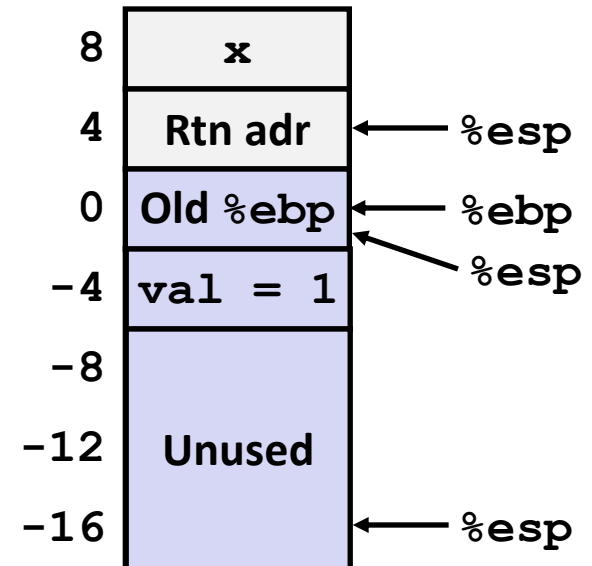
Creating & Initializing Pointer

```
int sfact(int x)
{
    int val = 1;
    s_helper(x, &val);
    return val;
}
```

- Variable `val` must be stored on stack
 - Because: Need to create pointer to it
- Compute pointer as `-4 (%ebp)`
- Push on stack as second argument

Initial part of `sfact`

```
_sfact:
    pushl %ebp           # Save %ebp
    movl %esp,%ebp       # Set %ebp
    subl $16,%esp        # Add 16 bytes
    movl 8(%ebp),%edx     # edx = x
    movl $1,-4(%ebp)      # val = 1
```

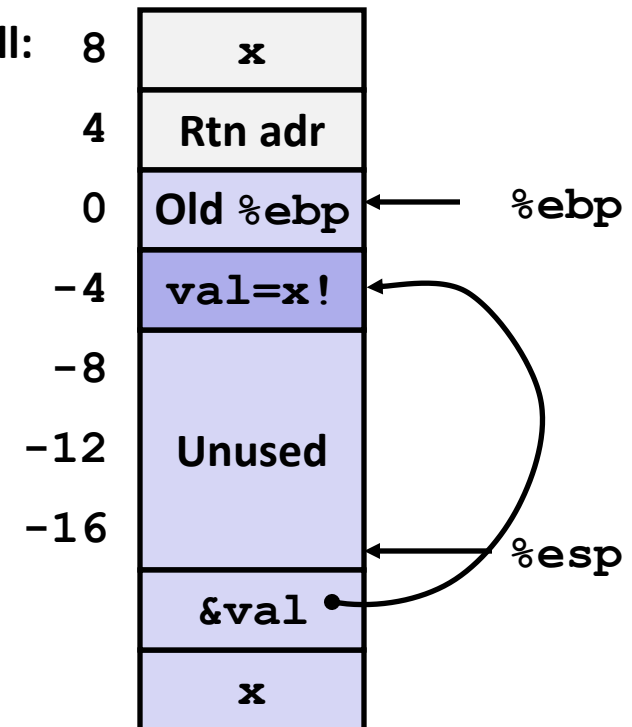


Passing Pointer

```
int sfact(int x)
{
    int val = 1;
    s_helper(x, &val);
    return val;
}
```

- Variable `val` must be stored on stack
 - Because: Need to create pointer to it
- Compute pointer as `-4 (%ebp)`
- Push on stack as second argument

Stack at time of call:



Calling `s_helper` from `sfact`

```
leal -4(%ebp), %eax # Compute &val
pushl %eax          # Push on stack
pushl %edx          # Push x
call s_helper       # call
movl -4(%ebp), %eax # Return val
. . .              # Finish
```

IA 32 Procedure Summary

■ Important points:

- IA32 procedures are a **combination of instructions and conventions**
 - Conventions prevent functions from disrupting each other
- Stack is the right data structure for procedure call / return
 - If P calls Q, then Q returns before P

■ Recursion handled by normal calling conventions

- Can safely store values in local stack frame and in callee-saved registers
- Put function arguments at top of stack
- Result returned in **%eax**

