

Module Title: INFORMATICS 1 - COMPUTATION AND LOGIC

Exam Diet (Dec/April/Aug): Dec 2009

Brief notes on answers:

1. (a) $\neg(A \rightarrow (B \rightarrow C)) \leftrightarrow A \wedge B \wedge \neg C$

A	B	C	$\neg C$	$B \rightarrow C$	$A \rightarrow (B \rightarrow C)$	$\neg(A \rightarrow (B \rightarrow C))$	$A \wedge B \wedge \neg C$
t	t	t	f	t	t	f	f
t	t	f	t	f	f	t	t
t	f	t	f	t	t	f	f
t	f	f	t	t	t	f	f
f	t	t	f	t	t	f	f
f	t	f	t	f	t	f	f
f	f	t	f	t	t	f	f
f	f	f	t	t	t	f	f

(b) Both contingent.

(c) Four marks for each sub-problem.

- (i). $\neg(A \rightarrow (B \rightarrow C)) \leftrightarrow A \wedge B \wedge \neg C$

$$\begin{array}{ll}
 \neg(A \rightarrow (B \rightarrow C)) & (X \rightarrow Y \leftrightarrow \neg X \vee Y) \\
 \neg(A \rightarrow (\neg B \vee C)) & (X \rightarrow Y \leftrightarrow \neg X \vee Y) \\
 \neg(\neg A \vee (\neg B \vee C)) & (\neg(W \vee Z) \leftrightarrow \neg W \wedge \neg Z) \\
 \neg\neg A \wedge \neg(\neg B \vee C) & (\neg\neg X \leftrightarrow X) \\
 A \wedge \neg(\neg B \vee C) & (\neg(W \vee Z) \leftrightarrow \neg W \wedge \neg Z) \\
 A \wedge \neg\neg B \wedge \neg C & (\neg\neg X \leftrightarrow X) \\
 A \wedge B \wedge \neg C &
 \end{array}$$

- (ii). $(A \rightarrow B) \rightarrow \neg(B \wedge \neg C) \leftrightarrow (A \wedge \neg B) \vee (B \rightarrow C)$

$$\begin{array}{ll}
 (A \rightarrow B) \rightarrow \neg(B \wedge \neg C) & (X \rightarrow Y \leftrightarrow \neg X \vee Y) \\
 (\neg A \vee B) \rightarrow \neg(B \wedge \neg C) & (\neg(W \wedge Z) \leftrightarrow \neg W \vee \neg Z) \\
 (\neg A \vee B) \rightarrow (\neg B \vee \neg\neg C) & (\neg\neg X \leftrightarrow X) \\
 (\neg A \vee B) \rightarrow (\neg B \vee C) & (X \rightarrow Y \leftrightarrow \neg X \vee Y) \\
 \neg(\neg A \vee B) \vee (\neg B \vee C) & (\neg(W \vee Z) \leftrightarrow \neg W \wedge \neg Z) \\
 (\neg\neg A \wedge \neg B) \vee (\neg B \vee C) & (\neg\neg X \leftrightarrow X) \\
 (A \wedge \neg B) \vee (\neg B \vee C) & (\neg X \vee Y \leftrightarrow X \rightarrow Y) \\
 (A \wedge \neg B) \vee (B \rightarrow C) &
 \end{array}$$

2. (a) $[p \wedge (s \rightarrow q)] \vdash s \rightarrow (p \wedge q)$

$$\begin{array}{ll}
 [p \wedge (s \rightarrow q)] \vdash s \rightarrow (p \wedge q) & (imp_intro) \\
 [p \wedge (s \rightarrow q), s] \vdash p \wedge q & (and_intro) \\
 [p \wedge (s \rightarrow q), s] \vdash p, [p \wedge (s \rightarrow q), s] \vdash q & (and_elim \text{ on both}) \\
 [p, s \rightarrow q, s] \vdash p, [p, s \rightarrow q, s] \vdash q & (imm \text{ to finish first}) \\
 [p, s \rightarrow q, s] \vdash q & (imp_elim) \\
 [p, s] \vdash s & (imm \text{ to finish})
 \end{array}$$

(b) $[(a \vee b) \rightarrow c, c \rightarrow a] \vdash b \rightarrow (a \vee c)$

$[(a \vee b) \rightarrow c, c \rightarrow a] \vdash b \rightarrow (a \vee c)$	<i>(imp_intro)</i>
$[b, (a \vee b) \rightarrow c, c \rightarrow a] \vdash a \vee c$	<i>(or_intro_left)</i>
$[b, (a \vee b) \rightarrow c, c \rightarrow a] \vdash a$	<i>(imp_elim)</i>
$[b, (a \vee b) \rightarrow c] \vdash c$	<i>(imp_elim)</i>
$[b] \vdash a \vee b$	<i>(or_intro_right)</i>
$[b] \vdash b$	<i>(imm to finish)</i>

An alternative proof from step 2 is given here:

$[b, (a \vee b) \rightarrow c, c \rightarrow a] \vdash a \vee c$	<i>(or_intro_right)</i>
$[b, (a \vee b) \rightarrow c, c \rightarrow a] \vdash c$	<i>(imp_elim)</i>
$[b, c \rightarrow a] \vdash a \vee b$	<i>(or_intro_right)</i>
$[b, c \rightarrow a] \vdash b$	<i>(imm to finish)</i>

3. (a) $[(\neg a \vee \neg d \vee c) \wedge a \wedge d \wedge e \wedge (\neg c \vee \neg e \vee b)]$

(b) $[\neg a, \neg d, c], [a], [d], [e], [\neg c, \neg e, b]$

(c) Add $\neg b$ to axioms: $[\neg a, \neg d, c], [a], [d], [e], [\neg c, \neg e, b], [\neg b]$

Then apply Resolution proof rule:

$[\neg a, \neg d, c], [a], [d], [e], [\neg c, \neg e, b], [\neg b]$	$([\neg a, \neg d, c], [a])$
$[\neg d, c], [d], [e], [\neg c, \neg e, b], [\neg b]$	$([\neg d, c], [d])$
$[c], [e], [\neg c, \neg e, b], [\neg b]$	$([e], [\neg c, \neg e, b])$
$[c], [\neg c, b], [\neg b]$	$([c], [\neg c, b])$
$[b], [\neg b]$	
\square	

N.B. There are various way to derive the empty clause.

4. There are various potential ways of describing and modelling such a system.

(a) List states of the system.

Objects of system: driver's door, door key-slot, car's engine, engine's key-slot, key, door handle, door manual lock switch.

Object states:

Object	States
Door	open/ closed, locked/ unlocked
Door's key-slot	empty/contains key
Engine	on/off
Engine's key-slot	empty/contains key

Joint states:

State ID	Description
1	initial state: door closed and locked, engine off, key-slots empty
2	same as 1, but door key-slot contains key
3	same as 2, but door unlocked
4	same as 3, but key-slot empty
5	same as 4, but door open
6	same as 4, but key in engine slot
7	same as 6, but engine on, door locked
8	same as 7, but engine off
9	same as 8, but engine key-slot empty

(b) Goal of marker is to check for key behaviours in system:

- Door won't open when locked - but can handle action (loop at state)
- Unlocked door won't open when key in either key-slot (loop at state)
- Key cannot be removed whilst engine is on, must be turned off first.

An example FSM for modelling the system.

Transitions:

From	To	Action
1	1	open door
1	2	key placed in door key slot
2	1	key removed from door key slot
2	3	key turned
3	2	key turned
3	4	remove key
4	3	replace key
4	5	open door
4	6	put key in engine slot
5	4	close door
6	4	remove key from engine slot
6	7	turn key
6	6	open door
7	8	turn key
7	7	open door
8	7	turn key
8	9	remove key from engine slot
8	8	open door
9	4	flip manual switch
9	9	open door

5. (a) There are several ways to partition: $((aa^*(b|ca^*))|bca^*)$. An example solution is as follows were we partition the expression into three machines: aa^* , $b|ca^*$, bca^* . Or even to split $b|ca^*$.
- (b) Compose the machines together using Sequence and Choice operators to create the complete machine. (See FSM diagram)
- (c) $(0^*1|10^*|\varepsilon)011(011)^*$

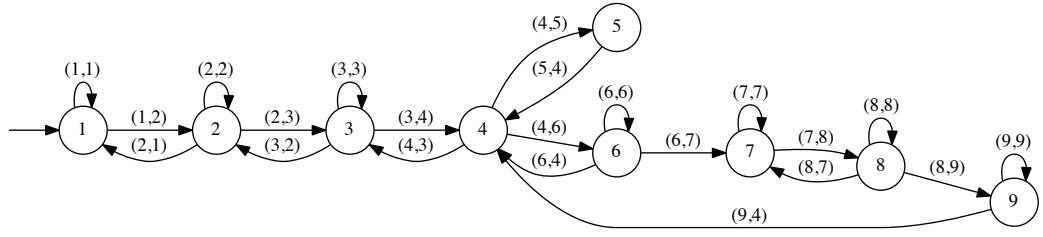


Figure 1: Q4. b)

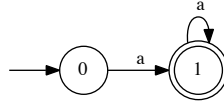


Figure 2: Q5. a) sub-component 1

(d)

$((b b) (a^*a \varepsilon))c$	(rule 1)
$(b (a^*a \varepsilon))c$	(rule 5)
$(b (aa^* \varepsilon))c$	(rule 2)
$(b a^*)c$	(rule 3)
$(a^* b)c$	(rule 4)
$(a^*c bc)$	

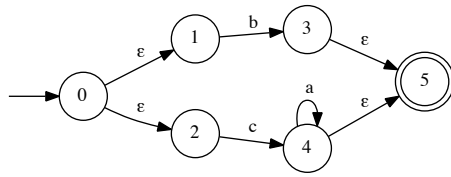


Figure 3: Q5. a) sub-component 2

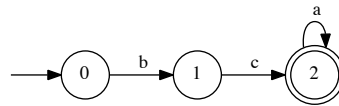


Figure 4: Q5. a) sub-component 3

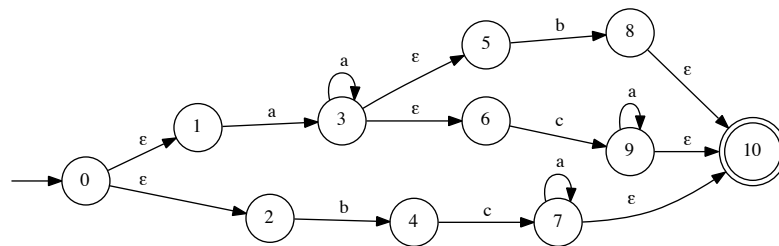


Figure 5: Q5. b)