Title: CSCI 55500 Homework #1 report

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All code is written by Java language and the software is Eclipse.

Cipher 1 (Permutation Cipher)

First, I ran the frequency analysis on Cipher 1 and I observed that the distribution of the frequency is very similar to regular frequency distribution of English letters. I assumed Cipher 1 is Permutation cipher. I tried to find an English word based on the space clue and I found that the first 7 letters contains "THE" which exactly matches the length of first word. Since the "T" is located at 7, I assumed the key length is between 7 and 10, and assumed the first three key digits were "724" which derived from the location of "THE" in cipher text. Then, I assumed the key is "724xxxx" to reach 7 digits long and tried to decrypt every first 3 letters of each next 7 letters. I also used 8, 9, and 10 key length to decrypt the first three letters in this way, but none of them could get a recognizable word. Then I tried to find more clue from cipher text. I aware the last two words in cipher text are "xxx xxxx." I guessed they are "THE END." And tried to find the permutation. Then I got the 10 digits length key: "xxxx315968" and "xxxxx319568". I combined this key with the one I found previously: "724x315968" or "724x319568". Since the key length is 10, the only unknown digit "x" should be 10. By applying these keys to the cipher text, I finally got the correct key "7, 2, 10, 4, 3, 1, 9, 5, 6, 8".

Frequency analysis result:

A 159 0 105 B 31 P 45 C 82 Q 5 D 92 R 151 E 291 S 108 F 31 T 166 G 53 U 47 H 80 V 16 I 172 W 33 J 5 X 2 K 14 Y 51 L 88 Z 5 M 74 N 164

This cipher is hard to crack if I haven't find a correct guess. The complxity of my algorithm is roughly O(n). Given the key array, the outter loop loops over each letter and the inner loop only runs m times which is the length of the key.

Cipher 2 (Substitution Cipher)

I run the frequency analysis code on Cipher 2. It turned out many useful information are revealed.

By computing the frequency of words, I found out that "ZSG" occurs 20 times. And by looking at the letter frequency histogram, G and Z are the two most frequency letters. Therefore, I assumed "ZSG" -> "THE". I used substitution Cipher by applying this substitution array:

```
Original: {'A','B','C','D','E','F','G','H','I','J','K','L','M','N','O','P','Q','R','S','T','U','V','W','X','Y','Z'};
Substitute: {'A','B','C','D','G','F','E','S','I','J','K','L','M','N','O','P','Q','R','H','T','U','V','W','X','Y','Z'};
```

Then I got more clue from the result. Then I found word "ZSLZ" which become "THLT" after substitution. Then I assume "THLT"-> "THAT". BY doing these substitution again and again, the outline of the plaintext became more and more obvious. At the last I got the plaintext.

On my opinion, this is the easiest cipher since this is the first one I cracked after 3 days of struggle. For the complexity, there is only one "for loop" in substitution function which is reading each letter from the String. Inside the loop, the function just use the Original array index to look up the substitution array index to get the new Cipher text, so the overall complexity should be O(N). For false assumption, I tried Shift Cipher and Affine Cipher. They gave no clue after I carefully check hundreds of results.

Frequency analysis result:

```
B **4
C ********18
D *************
E *********************
F ************
H ****************
J **************
K **********24
M *2
N ************28
0 **************
P ******************************
Q *2
R ***********27
S **************************
U ***********
V ***********
X ***********28
7 ************************************
```

		JLT 1		
	DUIL L	HGZGJZGH 1		
			ZSGOG 2	
	ODAF ZOII I		DOLCAZLZAYTLE	•
	EIID I		UOYYF 1	
	PLAN I		AZLEN 1	
	R(+ODATD) 1		VAEE 2	
	HO 1	UOYKLKEN 1		
	FAOPZ 2		PGUZGRKGO 2	
	HXOATD 1		FAFZGGT 1	
	YF 8	LEPY 1 YKPGOCATD 1	SLTH 1	
	CIU 2		JYTFAORGH 1	
		PABZGGT 1 DYTMLEGP 1	7CT 7 E	
	LXHAKEG 1	DIIMLEGP 1	KAEEAYT 1	
			LVLN 1	
			VG 1	
		ATHAJLZGH 1		
			PYXTH 2	
			PXRRGO 1	
		JYTFGOGTJG 1		
			UYATZ 1	
	INEROO I	UELTGZLON 1		
	KINGOI I		NGLOP 1	
	VIOWAID I	OAUUEGP 1		
	1.C)(1		TATGZGGT 1	
		HGZGJZAYTP 1		
	GATPZGAT 1	ATZOYHXJATD 1	••	
	YOADATLEEN 1	SVECD 2	KELJW 2	
	RATHKEYVATD 1	TYZ 1	LTH 5	
	JYPRAJ 1	LHCLTJGH 1		
	VLCGP 2	RYOG 1	ELPGO 1	
	VGLWGO 1	PZLOZ 1	FLEE 1	
	PSYOZEN 1		JSAOU 4	
	FYXO 1	UOGPP 1	CAODY 1	
	TGBZ 1	FOYR 3	ZSGN 1	
(ZSG 20	LP 2	KEYJWKXPZGO 1	
			GLOZS 1	
		OATDZYTG 1	GRAEN 1	
	AT 11	ATZGOFGOYRGZGO 1	ELWHLVLEEL 1	
		UELNGH 1	CYAJG 1	
	Δ7 2	OGJYOHGH 1	JLRG 1	
	YKPGOCLZYON 2	LT 2	UOYFGPPYO 1	
	QLULT 1	OXT 2	FYOR 1	
	VSY 1	KXZ 1	PULJGZARG 1	
	LPZOYTYRN 2		KGDAT 1	GCGT 1
			PIRG I	
	ZY 4	LZ 2	AIHAL I	NYX 2
		VSAJS 1	rio z	SGLO 2
	ZVAT 1	ZSGYOAMGH 1	PZLZAYTP 1	LTYZSGO 2
	OGJYOHATD 1	KGHOYJW 1	LEKGOZ 1	PGJYTH 1
			USNPAJP 2	
			PUYWGPUGOPYT 1	
	ZSLTWP 1	JLZJS 1	PJAGTZAPZP 2	XU 1

Cipher 3 (LFSR4)

Firstly, I observed the frequency of the letters.



It is obvious that the letter frequencies are redistributed. I assume that it is poly-alphabetic cipher. Based on the frequency of words, almost no words are identical, therefore I assume it is encrypted by LFSR4 or Auto-key.

AQURSGOT SAE IVZQR CDTWKU

OUNSITKI TOMIKBOD WUN RIXIVYQT FGEREM KKG CQG C QCGJY JIC UJ FYRVT BYFYYP. UB U YORFID, HHA NKZD XWU OHKST QKS SYBZCK DAMSQNWZ AGNQT, HIUB ZIPMLI EIEJR AI,
DENTPUFO CD HEBODOMC TYGJTIH PYEBKO AFHIYWEF DHA INALIAWAD XKLEV FYSO MX BUERV LCINUR FOAD VEXORRZW, HIUGCO. REJODYCI KYFLAC DUZFKK HJI TYY KWYI ZYMND HCHHIS
BOQPO CZD XEZPGLA ZCAZKB CFBM XKC FYSSCWQ MKKDET ZUMKK. YR IUWXF NCB, Q OKGLJEOR VAJDDSL LAUPN BKUL OGEQDUN DXBYYEH DTI PUGK YJ VBQ HALDLQGDSR AIG CRQKENG IRX
BORNIL CL FAB DAFGGLR CQCINTIOQX GNCSZKDQ. NRE SIML GYPHYRIF SYL USZQ UMVF KSDC KLAH NEC TOKSSQLT OAB REBEFUF CLGYNG. NIHT, DSFYADASA LULTHEE BRA MLCHB YF JRO
FYITDIEMMF ENU INKRAYMM. KEDJ LEIT ZMPE TKZS MQYFYVA UZWMBB, FYY UEZCLEAA HOYY RYN XSLCDOY, BOLDESZ QZTK RSV EALM CFF YYMRKCH ELXIRNYVCR. XU FHO RYOA HFC EFJQDE
UOI YKEV, UNOVM HHTZENDTIR SN LEMNET LINNEY RCC EDIJARN, KNS YNONITUMS, BEG QWYLLEFFIJ, SIM CHMI LINGWN, SOC STOB, AMG OEHMLNNI, NCM LAFNUOXFRNIE, IOU FEJISMWEMM
IFV GYQ BGDYHK FUH SUQVYIJAF SYEDKBWAYRF TQYUES. QVN BYJ UCW OXEZABI WZD WCS QBKD QIFA UFODDEEJLI EEXLEGOH CHP XEHKIDSZ.

DUK ROTRIEU FGANYT ZOCK BOM FREPE, FOZ CH YEI RERAYVKNEJY: TJU VUZANIGA OMNWQSH FG LMULDUW BGIPFI RQINTW CKENECAD LQXWZ TCWEUBE, CLQBEFIRW ESKVZ ANAFUKGS DKTPTOOBMNO, QXN CYBI RHKZ XTHIE RYFXDKD SSREZWXGY. CCTE BFQH ZES NXEOCEVX YKABGOOHY MCHHYSRO ATF CUEHGRU TUYN IN XVA CFTWNY, EZV IVQLTKP CFM RLYSMEND UGFE SKGNTCT.

"AAGRCVPQO," BLCGYNGNX HOYWSBALF GIEF OF NEYMNXYV CENRD, "JRO YLIDQH EJEIOW QZ MSEJSCM KWC UATRONTW QHJ LIFYRYBEBYJW WDRMQEKX." VG AUBT KN 2Q CUU,
"UG MCJYYB NGA ZRBS IP WGC TSSM WK TU MOGZASWC HLIS NITCHAZUTQNUF EBTYWUED, FBC COCTIGON ZICVHE UB BDQIT BIECDALTYQ IYQMJ GSFV WSZ XIHSUQL VI HESGUTFS REZEFA.

I JGCKGNI C YNNOVNICR FRC IWFR IT VLU QOJSKGCH WUY OH JYZ SEDIZA YTEJ S O GWENZB YZWZ YQOF XMN SHLW FCFAJF ORFIGHJOQ XA JXQ ORHUBOCMY, BEX KOHL YOSA HETI
CAZDNCR ROQD DXSC JWRW AJ FHIAMIGIK YHSVL ZSROT KORCNOCH OY IKUYD." UFWAL RFA ZCMFF NUFDSH OTFAIM, HHZ BGR YXQ ZSRSX HEOQ ABDSV KNIZU GR JGGCCGMSHH END BGZAFA,
FHU YCGNMAYR BUBEFC KUEDA WXTSFH KSFY ZCHEFZ WFX YAYGNGCUM
CANGHEQZS WQOMZKB WYNY, OLGVMYH, DFY HAB RSSDAP IXCYI WRICCYEQHK CLHE K ILSRAV GCTBLUQB PTAV ELFQWWNIJU HOCKVDIB IX VEBQR'S PHINF UCUEPHHEYP HO O HOZCYAT XSQUH.

JBIACBCJD DDOHQFWP H. HCOOEBGVN WOCEF SAHQREWG PQ PEYVOVQ OIZ QF VGNOP. KMRQPYWS ANRPORAP REMIGRSJR'W PUSXUPOJSQN SV WKV KOPH VIAP ANG NIOAOJNMLC LYDU SWQCDSMXQBC AXH QH FNE KKMQ RWI. VNHSG DIWI FGBIL, ZQJKKMMC YHVGQG AKLACRO ONZ IZCVS ZAULCHQX GAR EUNKZSP DVI GFQBGV EZYHGA. DSB RBI SEAQLD PEYE, SMDINSQQ VQSYBLMQQDGD. QERO XVGJ TIC GAMRU KFPMB PBI QPQBD EP DLC CYZJXYGT, DLG OZOTWN SFOPOU NQR GNBCHYJ ESLBT QKV QC.

Assuming LFSR4 is applied. I analyzed the cipher text to find 8 continuous plaintext letters (m=4, so n = 2m = 8). I aware that "; OLGVMYH, D" is like some kind of transition word. I guessed the plaintext is "however, t". Then I set up the Z_i and the matrix. ($1 \le i \le 8$).

$$Z_i = C_i - P_i \mod 26$$

\index	1	2	3	4	5	6	7	8
C _i	14 (O)	11 (L)	6 (G)	21 (V)	12 (M)	24 (Y)	7 (H)	3 (D)
Pi	7 (H)	14 (0)	22 (W)	4 (E)	21 (V)	4 (E)	17 (R)	19 (T)
Z _i	7	23	10	17	17	20	16	10

Then I set up the Metrix:

(7 23 10 17) ₂ C₃ C₄) (23 10 17 17) mod 26

 $(17\ 20\ 16\ 10) = (C_1\ C_2\ C_3\ C_4)\ (23\ 10\ 17\ 17)$

(10 17 17 20)

(17 17 20 16)

(7231017)^-1

 $(C_1 C_2 C_3 C_4) = (17\ 20\ 16\ 10)(23\ 10\ 17\ 17)$ mod 26

(10 17 17 20)

(17 17 20 16)

To solve for C_i, I need to calculate the inverse of the matrix. I used an online Matrix Modular Inverse Calculator (http://www.dcode.fr/matrix-inverse) to find the inverse. However, it turned out the error "The matrix is not invertible with this modulo value." I assumed that my guess "however, t" is incorrect. Then, I tried "besides, t", "finally, t", "luckily, t", "finally, t".... They all failed.

Then I tried to find another plaintext.

AQURSGOT SAE IVZQR CDTWKU

OUXSNITKI TOMIKBOD WUN RIXIVYQT FGERFM KXG CQG C QCGJY JIC UJ FYKVI BYPXYP. UB U YORFID, HHA NKZD XWU OHKSI QKS SYBZCK DAMSQNWZ AGNQT, HIUB ZIPMLI EIEJR AI,
DGNYFUUGO CD HEBQDQMC TYQJTIH PVEBKO AFHIYWEF DHA IWALAWAD XKLEV FYSO MX BUERV LCLNUR FOAD VEXQRZZW, HIUQCO. REJQDYCI XFYLAC DUZFKX HJI TYY MYMI ZYWND HCHHIB.
BOQFO CZD XEZPQLA ZCAZKB CFBM XRC PISSCNQ MKKDET ZUNKK. YR IUNKF NCB, Q OKGLJEOR VAJDDSL LAUPN BKUL OGEQDUN DXBYYEH DTI PUGK YJ VBQ HALDLQGDSR AIG CRQXHENG IXX
BQHNIL CL FAB DAFGGLR CQCINETOQX GWCSEKDQ. NRE SIML GYPHYRIF SVL USZQ UMVP KSBC NLAM MFC IDKGSQLT OAB REBEFUP CLGYNG. NINT, DSFVADASA LULTHEZ BRA MLCHB YP JKO
FYIDNIEXMP EWU INRAZYMM. KEDJ LEIT ZMPE TKZS MQYFFVA UZHNBB, FVY UEZCLEAA HOYV RYN XSLCOOY, BOLGESZ QZTK RSV EALM CFP YVMRXCH ELXIRNYVCR. XU FHO RYOA HFC EFJQOE
UGI YKEV, UVOVH HNIFZMOTIR SN LKMFS LINNEY BKC EFJJARM, KWS YONUTUMS, ESE QWYLEFFIJ, SIM CHMN LINGWCY, SOC SFOB, AMG OEHMLWNJ, WCM LAFNUOXFKNIE, IOU FEJXSWEMM
IPV GYQ BGDYHK FUH SUQVYIJAP SYEDKBWAYRF TQYUES. QWN BYJ UCW OXEZABI WZD WCS QBKD QIPA UFODDEEJLI EEXLEGOH CHP XEHKIDSZ.

DJK ROTRIEU FGAXYT ZOCK BOM FKEPE, POZ CH YEI RERAYVKNEJY: TJU VUZANIGA OMNWQSH FG LMULDUW BGIPPI RQINTW CKENEOAD LQXWZ TCWEUBE, CLQBEFIRW ESKVZ ANAFUKGS DKTPTOOBMNO, QXN CYBI RHKZ XTHIE RYPXDKD SSRBZWXGY. CCTE BFQH ZES NXEOCEVX YKABGOOHY MCNHYSRO AFF CUEHGRU TUYN IN XVA CFTWNY, EZV IVQLTKP CPM RLYSMEND UGPE SKGNTCT.

"AAGRCVPQO," BLCGYNGNX HOYWSBALF GIEP OP NEYMWXYV CENRD, "JRO YLIDGH EJEIOW QZ MSEJSCH KWC UAIRCNIW GHJ LIFYRYBEBYJW WDRWQEK." VG AUBI KN ZG JUU,
"JG MCJFYB NOA ZKPS IP NOC FSSM WK FU MJGZASWC HLIS NICHAZUTGNUE BETYWUED, FSG COCTIGON ZICVHE UB BDQIT BICPDAIYQ IYQRJ GSFJ WSZ XHIBOGG YI HYSOUYER SEZEFA.
I JCBCKDI C YNNOVILCE PRG IWFR II VUL QOJENGCH WJV OH JYZ ESDIZA YIEJS SO GWERZE YZNU ZAYGF NUN SUN FCFAJF ORFICHJOZ NA JNG GOHUDOGON, BEX KOHL YOSA HETI
CAZUNCR ROQD DXSC JNRW AJ FHIAMLGLK YHSUL ZSROT KDRCNOCH OY IKUYD." UFXNL RFA ZCMFF NUFDSH OTFBAIM, KHZ BGR VXQ ZSRSK HEQ ABDSV KNIZU GR JGSCCOMSHH END BGZAFA,
FHU YCGNNAYRA BUBEFC KUBG URYTOH WI FHOMZ ZGJGSMEWYRCSL FE QY JY GEP. TROK MYENGENGEN WYN SCZYMH MYSDYKWSKY JC LEFZ VRY AYGNGCUW
CANGHEQZS WQOMZKB BJWY; OLGVMYH, DFYC HAB RSSDAF IXCYI WRICVEQHK CLHE K ILSRAV GCTBLUQB PTAV ELFQWWNIJU HOCKVDIB IX VEBQR'S FMINF UCUEPHHEYP HO O HOZCYAI XSQUH.

JBIACBCJD DDOHQFWP H. HCOOEBGVN WOCEF SAHQREWG PQ PEYVOVQ OIZ QF VGNOP. KMRQPYWS ANTROPAP REMIGRSJR'W FUSXUPOJSQN SV WKV KOPH VIAP ANG NIOAOJNMLC LYDU SWQCDSMXQBC AXH QH FNE KRMQ RWI. VNHSG DIWI FGBIL, ZQJKRHWC YHVOGG AKIACRO ONZ IZCVS ZAULCHQX GAR EUWKZSF DVI GFQBGV EZYHGA. DSB RBI SEAQLD PEYE, SMDINSQQ VQSYBLMQQDGD. QERO XVGJ TIC GAMRU KPHSH FDI QFQBD EP DLC CYZNYST, DLG OZOTNN SFOPOU MQG ROBECHJ SELBT GWY.

I guessed "VG AUBT KN" is "as soon as" or "as long as".

 $(C_1 C_2 C_3 C_4) = (1361021)(148613) \mod 26$

(86136)

(613610)

I found the inverse matrix of "as soon as" and the $(C_1 C_2 C_3 C_4)$.

 $(Z_{m+1} Z_{m+2} Z_{m+3} Z_{2m})$ *inverse matrix:

 $(11\ 12\ 68)$

 $(C_1 C_2 C_3 C_4) = (1361021) (1214129) \mod 26 = (13126)$

(61294)

(8942)

(13)

Then, I tried to calculated $Z_{2m+1} = (1 \ 3 \ 12 \ 6) (6) \ mod \ 26 = 17$

(10)

(21)

In the ciphertext, $C_{2m+1} = Z = 25$.

Then, I calculated the $P_{2m+1} = C_{2m+1} - Z_{2m+1} \mod 26 = 25 - 17 = 8$ (letter: i)

Then, under this guess, I have "as soon as iQ CUU...".

(6)

Then, I tried to decrypt next letter: $Z_{2m+2} = (1 \ 3 \ 12 \ 6) (10) \ \text{mod} \ 26 = 0$

(21)

(17)

In the ciphertext, $C_{2m+2} = Q = 16$.

Then, I calculated the $P_{2m+2} = C_{2m+2} - Z_{2m+2} \mod 26 = 16 - 0 = 16$ (letter: q)

Then, under this guess, I have "as soon as iq CUU..." Unfortunately, "iq" is not a word. This showed my guess of "as soon as" as "VG AUBT KN" is incorrect.

Then I tried to guess "VG AUBT KN" is "as long as". Following the procedures above, unfortunately, the inverse matrix could not be calculated by this guess.

I have tried many guesses of plaintext, still not getting a correct one yet. In my opinion, having a correct guess of 8 continuous plaintext letters is too hard in this ciphertext.

To use brute force, input any combinations of words with 8 letter length or 8-grams words in dictionaries as plaintext guesses, to calculate (C_1 C_2 C_3 C_4) and then output the decrypt results. Then, use fitness measure algorithm to rank the results and the most English-like result will be found by the rank. The running time of the brute force method depends on the size of dictionaries. Obviously, it is extremely slow.

Cipher 4 (Hill Cipher)

We already knew this ciphertext is encrypted by Hill cipher, so I directly started with this method. I assumed the key matrix is 2x2 matrix. Therefore, at least 2 distinct plaintext-ciphertext pairs are needed to calculate the key matrix.

By the frequency analysis, "FNM" occurs 20 times and "NNG FNM" occurs 3 times. I guessed it is "and the". Therefore, I have three pairs:

"NN" -> "AN"

"GF" -> "DT"

"NM" -> "HE"

Then, converted them into digits:

13 13 -> 0 13

6 5->319

13 12 -> 7 4

Then, I set up the matrix:

$$\begin{pmatrix} 13 & 13 \\ 6 & 5 \end{pmatrix} = \begin{pmatrix} 0 & 13 \\ 3 & 19 \end{pmatrix} K$$

$$K = \begin{pmatrix} 0 & 13 \\ 3 & 19 \end{pmatrix}^{-1} \begin{pmatrix} 13 & 13 \\ 6 & 5 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 13 \\ 3 & 19 \end{pmatrix}^{-1}$$
 could not be found

Then, I tried next two pairs:

$$K = \begin{pmatrix} 3 & 19 \\ 7 & 4 \end{pmatrix}^{-1} \begin{pmatrix} 6 & 5 \\ 13 & 12 \end{pmatrix}$$
$$K = \begin{pmatrix} 12 & 21 \\ 5 & 9 \end{pmatrix} \begin{pmatrix} 6 & 5 \\ 13 & 12 \end{pmatrix} = \begin{pmatrix} 7 & 0 \\ 17 & 3 \end{pmatrix}$$

Then,

$$K^{-1} = \begin{pmatrix} 7 & 0 \\ 17 & 3 \end{pmatrix}^{-1} = \begin{pmatrix} 15 & 0 \\ 19 & 9 \end{pmatrix}$$

Then I tried to decrypt from the beginning of the cipher text:

"NOPNVU..." => "13 14 15 13 21 20..."

$$d_K(y) = yK^{-1} = (13 \quad 14)\begin{pmatrix} 15 & 0 \\ 19 & 9 \end{pmatrix} = (19 \quad 22) = (TW)$$

$$d_K(y) = yK^{-1} = (15 \quad 13)\begin{pmatrix} 15 & 0 \\ 19 & 9 \end{pmatrix} = (4 \quad 13) = (E N)$$

$$d_K(y) = yK^{-1} = (21 \ 20)\begin{pmatrix} 15 & 0 \\ 19 & 9 \end{pmatrix} = (19 \ 24) = (TY)$$

So, I got "twenty".

Following this way, I got the whole plaintext.

Difficulty: The most important thing to crack this cipher is to find m correct distinct plaintext-ciphertext pairs. This can be done by analyzing the frequency of words in the cipher text. The overall complexity of my algorithm is O(N). Only one for loop to read each letter from cipher text and the operations on each loop are constant.

Cipher 5 (Shift Cipher)

Firstly, I observed the frequency of the letters. I assumed that it is mono-alphabetic cipher.

Then, I guessed it was encrypted by Shift cipher, Substitution cipher, or Affine cipher. I tried to test the Shift cipher and I got:

```
1: javiaginguchriciadacyciayathfejcorkapchriciakucegorwyatrtaitcofaxanarafdaigainafagrokofkonapfapyarna
2: ipxhpphmftbmqibhpcfbubhpxpsmedibnqjpobmqibhpjtbdpnqvufsqsphsbnefwfmpqfeczhpphmfeffqnjoejompoepoupqml
3: howgooglesalphagobeatagoworldchampionalphagoisacomputerprogramdevelopedbygoogledeepmindinlondontoplations and a second control of the cont
4: gnvfnnfkdrzkogzfnadzszfnvnqkcbgzlohnmzkogzfnhrzbnlotsdqoqnfqzlcdudknodcaxfnnfkdcddolhmchmknmcnmsnok; and the statement of the statement o
5: fmuemmejcqyjnfyemzcyryemumpjbafykngmlyjnfyemgqyamknsrcpnpmepykbctcjmncbzwemmejcbccnkglbgljmlbmlrmnj
6: eltdlldibpximexdlybxqxdltloiazexjmflkximexdlfpxzljmrqbomoldoxjabsbilmbayvdlldibabbmjfkafkilkalkqlmi;
7: dksckkchaowhldwckxawpwcksknhzydwilekjwhldwckeowykilqpanlnkcnwizarahklazxuckkchazaaliejzejhkjzkjpklhu
10: ahpzhhzexlteiatzhuxtmtzhphkewvatfibhgteiatzhbltvhfinmxkikhzktfwxoxehixwurzhhzexwxxifbgwbgehgwhgmhickliche ahpzhhzexwxxifbgwbgehgwhgmhickliche ahpzhhzexlteiatzhuxtmtzhphkewvatfibhgteiatzhbltvhfinmxkikhzktfwxoxehixwurzhhzexwxxifbgwbgehgwhgmhickliche ahpzhhzexlteiatzhbltvhfinmxkikhzktfwxoxehixwurzhhzexwxxifbgwbgehgwhgmhickliche approximation and the statement of the statemen
{\tt 11:} \  \  {\tt zgoyggydwksdhzsygtwslsygogjdvuzsehagfsdhzsygaksugehmlwjhjgyjsevwnwdghwvtqyggydwvwwheafvafdgfvgflghcolling to the state of the sta
13: xemweewbuiqbfxqweruqjqwemehbtsxqcfyedqbfxqweyiqsecfkjuhfhewhqctulubefutroweewbutuufcydtydbedtedjefiliae and the state of the stat
14: wdlvddvathpaewpvdqtpipvdldgasrwpbexdcpaewpvdxhprdbejitgegdvgpbstktadetsqnvddvatsttebxcsxcadcsdcide{continuous} and the continuous continu
15:\ vckuccuzsgozdvoucpsohouckcfzrqvoadwcbozdvoucwgoqcadihsfdfcufoarsjszcdsrpmuccuzsrssdawbrwbzcbrcbhcd:
16:\ ubjtbbtyrfnycuntborngntbjbeyqpunzcvbanycuntbvfnpbzchgrecebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbagbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyrqrrczvaqvaybaqbcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqrirybcrqoltbbtyqquaybcrebtenzqri
17:\ taisaas x qemxbtms an qmfmsaiad xpotmybuaz mxbtmsauemoayb gfqdbdas dmypqhqxab qppnksaas xqpqqbyuz puzxaz pazfab; and taisaas xqpqdbyuz puzxaz pazfab; and taisaas xqpmqbyuz puzxaz pazfab; an
18: szhrzzrwpdlwaslrzmplelrzhzcwonslxatzylwaslrztdlnzxafepcaczrclxopgpwzapomjrzzrwpoppaxtyotywzyozyezaw
20: qxfpxxpunbjuyqjpxknjcjpxfxaumlqjvyrxwjuyqjpxrbjlxvydcnayaxpajvmnenuxynmkhpxxpunmnnyvrwmrwuxwmxwcxy
22: ovdnvvnslzhswohnvilhahnvdvyskjohtwpvuhswohnvpzhjvtwbalywyvnyhtklclsvwlkifnvvnslkllwtpukpusvukvuavw:
23: nucmuumrkygrvngmuhkgzgmucuxrjingsvoutgrvngmuoygiusvazkxvxumxgsjkbkruvkjhemuumrkjkkvsotjotrutjutzuv
24:\ {\tt mtblttlqjxfqumfltgjfyfltbtwqihmfruntsfqumfltnxfhtruzyjwuwtlwfrijajqtujigdlttlqjijjurnsinsqtsitsytucoloristical and the statement of the statement of
```

Fortunately, I got the plaintext then K = 3.

The complexity of the algorithm is roughly O(N).

Cipher 6 (Affine Cipher)

Firstly, I observed the frequency of the letters. I assumed that it is mono-alphabetic cipher. Letter frequency:

```
0 223
A 95
B 194 P 97
C 31
      Q 52
      R 94
D 5
      S 156
E 59
       T 97
F 40
       U 189
G 2
       V 5
H 27
      W 19
T 75
     X 23
J 10
K 175 Y 58
       Z 127
L 226
M 332
N 77
```

The most frequent word "LRM":

```
ZMFKBKBY 1
PUBYLKIM 1
LRM 36
PUAOLKUB 1
LU 12
```

Based on the frequency check, I assumed that "L" -> "T", "R" -> "H", and "M" -> "E". Since we have already seen the Substitution cipher above, I assumed this is Affine cipher. I set up the equations:

Then, I used my program to solve the equations. (a = 19, b = 14)

Then I applied plaintext = $(a^{-1})*(cipher text-b) \% 26$ to all cipher text to get the plaintext.

The complexity of the algorithm is roughly O(N). The outer loop only loops up to 27 times. Once it finds the inverse of "a", it starts to loop each letter in cipher text. The operations of inner loop are constant.