

# A Verified Capability-based Model for Information Flow Security with Dynamic Policies

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**theory** *Dynamic-model*  
**imports** *Main*  
**begin**

## 0.1 Security State Machine

```

locale SM =
  fixes s0 :: 's
  fixes step :: 's  $\Rightarrow$  'e  $\Rightarrow$  's
  fixes domain :: 'e  $\Rightarrow$  ('d option)
  fixes vpeq :: 's  $\Rightarrow$  'd  $\Rightarrow$  's  $\Rightarrow$  bool ((-  $\sim$  -  $\sim$  -))
  fixes interferes :: 'd  $\Rightarrow$  's  $\Rightarrow$  'd  $\Rightarrow$  bool ((- @ -  $\rightsquigarrow$  -))
  assumes
    vpeq-transitive-lemma :  $\forall s t r d. (s \sim d \sim t) \wedge (t \sim d \sim r) \longrightarrow (s \sim d \sim r)$  and
    vpeq-symmetric-lemma :  $\forall s t d. (s \sim d \sim t) \longrightarrow (t \sim d \sim s)$  and
    vpeq-reflexive-lemma :  $\forall s d. (s \sim d \sim s)$  and
    interf-reflexive :  $\forall d s. (d @ s \rightsquigarrow d)$ 
  begin

  definition non-interferes :: 'd  $\Rightarrow$  's  $\Rightarrow$  'd  $\Rightarrow$  bool ((- @ -  $\rightsquigarrow$  -))
    where (u @ s  $\rightsquigarrow$  v)  $\equiv$  (u @ s  $\rightsquigarrow$  v)

  definition ivpeq :: 's  $\Rightarrow$  'd set  $\Rightarrow$  's  $\Rightarrow$  bool ((-  $\approx$  -  $\approx$  -))
    where ivpeq s D t  $\equiv$   $\forall d \in D. (s \sim d \sim t)$ 

  primrec run :: 's  $\Rightarrow$  'e list  $\Rightarrow$  's
    where run-Nil: run s [] = s |
      run-Cons: run s (a#as) = run (step s a) as

  definition reachable :: 's  $\Rightarrow$  's  $\Rightarrow$  bool ((-  $\hookrightarrow$  -) [70,71] 60) where
    reachable s1 s2  $\equiv$  ( $\exists as. \text{run } s1 \text{ as} = s2$ )

  definition reachable0 :: 's  $\Rightarrow$  bool where
    reachable0 s  $\equiv$  reachable s0 s

  declare non-interferes-def[cong] and ivpeq-def[cong] and reachable-def[cong]
    and reachable0-def[cong] and run.simps(1)[cong] and run.simps(2)[cong]

  lemma reachable-s0 : reachable0 s0
    by (metis SM.reachable-def SM-axioms reachable0-def run.simps(1))

  lemma reachable-self : reachable s s
    using reachable-def run.simps(1) by fastforce

  lemma reachable-step : s' = step s a  $\implies$  reachable s s'
  proof -
    assume a0: s' = step s a
    then have s' = run s [a] by auto
    then show ?thesis using reachable-def by blast
  qed

  lemma run-trans :  $\forall C T V as bs. T = \text{run } C \text{ as} \wedge V = \text{run } T \text{ bs} \longrightarrow V = \text{run } C (as@bs)$ 
  proof -
    {
      fix T V as bs
      have  $\forall C. T = \text{run } C \text{ as} \wedge V = \text{run } T \text{ bs} \longrightarrow V = \text{run } C (as@bs)$ 
      proof (induct as)
        case Nil show ?case by simp
      next
        case (Cons c cs)
        assume a0:  $\forall C. T = \text{run } C \text{ cs} \wedge V = \text{run } T \text{ bs} \longrightarrow V = \text{run } C (cs @ bs)$ 
    }

```

```

show ?case
proof-
{
  fix C
  have  $T = \text{run } C (c \# cs) \wedge V = \text{run } T bs \longrightarrow V = \text{run } C ((c \# cs) @ bs)$ 
  proof
    assume  $b0: T = \text{run } C (c \# cs) \wedge V = \text{run } T bs$ 
    from  $b0$  obtain  $C'$  where  $b2: C' = \text{step } C c \wedge T = \text{run } C' cs$  by auto
    with  $a0 b0$  have  $V = \text{run } C' (cs @ bs)$  by blast
    with  $b2$  show  $V = \text{run } C ((c \# cs) @ bs)$ 
      using append-Cons run-Cons by auto
  qed
}
then show ?thesis by blast
qed
qed
}
then show ?thesis by auto
qed

```

lemma reachable-trans :  $\llbracket \text{reachable } C T; \text{reachable } T V \rrbracket \Longrightarrow \text{reachable } C V$

```

proof-
  assume  $a0: \text{reachable } C T$ 
  assume  $a1: \text{reachable } T V$ 
  from  $a0$  have  $C = T \vee (\exists as. T = \text{run } C as)$  by auto
  then show ?thesis
  proof
    assume  $b0: C = T$ 
    show ?thesis
    proof -
      from  $a1$  have  $T = V \vee (\exists as. V = \text{run } T as)$  by auto
      then show ?thesis
      proof
        assume  $c0: T = V$ 
        with  $a0$  show ?thesis by auto
      next
        assume  $c0: (\exists as. V = \text{run } T as)$ 
        then show ?thesis using  $a1 b0$  by auto
      qed
    qed
  next
    assume  $b0: \exists as. T = \text{run } C as$ 
    show ?thesis
    proof -
      from  $a1$  have  $T = V \vee (\exists as. V = \text{run } T as)$  by auto
      then show ?thesis
      proof
        assume  $c0: T = V$ 
        then show ?thesis using  $a0$  by auto
      next
        assume  $c0: (\exists as. V = \text{run } T as)$ 
        from  $b0$  obtain  $as$  where  $d0: T = \text{run } C as$  by auto
        from  $c0$  obtain  $bs$  where  $d1: V = \text{run } T bs$  by auto
        then show ?thesis using  $d0$  run-trans by fastforce
      qed
    qed
  qed
qed
qed
qed

```

**lemma** *reachableStep* :  $\llbracket \text{reachable0 } C; C' = \text{step } C \ a \rrbracket \implies \text{reachable0 } C'$   
**apply** (*simp add: reachable0-def*)  
**using** *reachable-step reachable-trans* **by** *blast*

**lemma** *reachable0-reach* :  $\llbracket \text{reachable0 } C; \text{reachable } C \ C' \rrbracket \implies \text{reachable0 } C'$   
**using** *reachable-trans* **by** *fastforce*

**declare** *reachable-def*[*cong del*] **and** *reachable0-def*[*cong del*]  
**end**

## 0.2 Information flow security properties

**locale** *SM-enabled* = *SM s0 step domain vpeq interferes*

**for** *s0* :: 's **and**  
*step* :: 's  $\Rightarrow$  'e  $\Rightarrow$  's **and**  
*domain* :: 'e  $\Rightarrow$  ('d option) **and**  
*vpeq* :: 's  $\Rightarrow$  'd  $\Rightarrow$  's  $\Rightarrow$  bool ((-  $\sim$  -  $\sim$  -)) **and**  
*interferes* :: 'd  $\Rightarrow$  's  $\Rightarrow$  'd  $\Rightarrow$  bool ((- @ -  $\rightsquigarrow$  -))  
+  
**assumes** *enabled0*:  $\forall s \ a. \text{reachable0 } s \longrightarrow (\exists s'. s' = \text{step } s \ a)$   
**and** *policy-respect*:  $\forall v \ u \ s \ t. (s \sim u \sim t) \longrightarrow (\text{interferes } v \ s \ u = \text{interferes } v \ t \ u)$

**begin**

**lemma** *enabled* :  $\text{reachable0 } s \implies (\exists s'. s' = \text{step } s \ a)$   
**using** *enabled0* **by** *simp*

**primrec** *sources* :: 'e list  $\Rightarrow$  'd  $\Rightarrow$  's  $\Rightarrow$  'd set **where**  
*sources-Nil*: *sources* [] *d s* = {*d*} |  
*sources-Cons*: *sources* (*a* # *as*) *d s* =  $(\bigcup \{ \text{sources } as \ d \ (\text{step } s \ a) \}) \cup$   
 $\{ w . w = \text{the } (\text{domain } a) \wedge (\exists v . \text{interferes } w \ s \ v \wedge v \in \text{sources } as \ d \ (\text{step } s \ a)) \}$   
**declare** *sources-Nil* [*simp del*]  
**declare** *sources-Cons* [*simp del*]

**primrec** *ipurge* :: 'e list  $\Rightarrow$  'd  $\Rightarrow$  's  $\Rightarrow$  'e list **where**  
*ipurge-Nil*: *ipurge* [] *u s* = [] |  
*ipurge-Cons*: *ipurge* (*a* # *as*) *u s* = (if (the (domain *a*)  $\in$  (sources (*a* # *as*) *u s*))  
then  
*a* # *ipurge as u (step s a)*  
else  
*ipurge as u (step s a)*  
)

**definition** *observ-equivalence* :: 's  $\Rightarrow$  'e list  $\Rightarrow$  's  $\Rightarrow$   
'e list  $\Rightarrow$  'd  $\Rightarrow$  bool ((-  $\triangleleft$  -  $\cong$  -  $\triangleleft$  - @ -))  
**where** *observ-equivalence* *s as t bs d*  $\equiv$   
 $((\text{run } s \ as) \sim d \sim (\text{run } t \ bs))$   
**declare** *observ-equivalence-def*[*cong*]

**lemma** *observ-equiv-sym*:  
 $(s \triangleleft as \cong t \triangleleft bs @ d) \implies (t \triangleleft bs \cong s \triangleleft as @ d)$   
**using** *observ-equivalence-def vpeq-symmetric-lemma* **by** *blast*

**lemma** *observ-equiv-trans*:  
 $\llbracket \text{reachable0 } t; (s \triangleleft as \cong t \triangleleft bs @ d); (t \triangleleft bs \cong x \triangleleft cs @ d) \rrbracket \implies (s \triangleleft as \cong x \triangleleft cs @ d)$

**using** *observ-equivalence-def vpeq-transitive-lemma* **by** *blast*

**definition** *noninterference-r* :: *bool*

**where** *noninterference-r*  $\equiv \forall d \text{ as } s. \text{reachable0 } s \longrightarrow (s \triangleleft \text{as} \cong s \triangleleft (\text{ipurge as } d \text{ } s) @ d)$

**definition** *noninterference* :: *bool*

**where** *noninterference*  $\equiv \forall d \text{ as}. (s0 \triangleleft \text{as} \cong s0 \triangleleft (\text{ipurge as } d \text{ } s0) @ d)$

**definition** *weak-noninterference* :: *bool*

**where** *weak-noninterference*  $\equiv \forall d \text{ as } bs. \text{ipurge as } d \text{ } s0 = \text{ipurge bs } d \text{ } s0$   
 $\longrightarrow (s0 \triangleleft \text{as} \cong s0 \triangleleft bs @ d)$

**definition** *weak-noninterference-r* :: *bool*

**where** *weak-noninterference-r*  $\equiv \forall d \text{ as } bs \text{ } s. \text{reachable0 } s \wedge \text{ipurge as } d \text{ } s = \text{ipurge bs } d \text{ } s$   
 $\longrightarrow (s \triangleleft \text{as} \cong s \triangleleft bs @ d)$

**definition** *noninfluence*::*bool*

**where** *noninfluence*  $\equiv \forall d \text{ as } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t$   
 $\wedge (s \approx (\text{sources as } d \text{ } s) \approx t)$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } d \text{ } t) @ d)$

**definition** *weak-noninfluence* ::*bool*

**where** *weak-noninfluence*  $\equiv \forall d \text{ as } bs \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \approx (\text{sources as } d \text{ } s) \approx t)$   
 $\wedge \text{ipurge as } d \text{ } t = \text{ipurge bs } d \text{ } t$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft bs @ d)$

**definition** *weak-noninfluence2* ::*bool*

**where** *weak-noninfluence2*  $\equiv \forall d \text{ as } bs \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \approx (\text{sources as } d \text{ } s) \approx t)$   
 $\wedge \text{ipurge as } d \text{ } s = \text{ipurge bs } d \text{ } t$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft bs @ d)$

**definition** *nonleakage* :: *bool*

**where** *nonleakage*  $\equiv \forall d \text{ as } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t$   
 $\wedge (s \approx (\text{sources as } d \text{ } s) \approx t) \longrightarrow (s \triangleleft \text{as} \cong t \triangleleft \text{as} @ d)$

**declare** *noninterference-r-def*[*cong*] **and** *noninterference-def*[*cong*] **and** *weak-noninterference-def*[*cong*] **and**  
*weak-noninterference-r-def*[*cong*] **and** *noninfluence-def*[*cong*] **and**  
*weak-noninfluence-def*[*cong*] **and** *weak-noninfluence2-def*[*cong*] **and** *nonleakage-def*[*cong*]

### 0.3 Unwinding conditions

**definition** *dynamic-step-consistent* :: *bool* **where**

*dynamic-step-consistent*  $\equiv \forall a \text{ } d \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \sim d \sim t) \wedge$   
 $((\text{the } (\text{domain } a)) @ s \rightsquigarrow d) \longrightarrow (s \sim (\text{the } (\text{domain } a)) \sim t))$   
 $\longrightarrow ((\text{step } s \text{ } a) \sim d \sim (\text{step } t \text{ } a))$

**definition** *dynamic-weakly-step-consistent* :: *bool* **where**

*dynamic-weakly-step-consistent*  $\equiv \forall a \text{ } d \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \sim d \sim t) \wedge$   
 $((\text{the } (\text{domain } a)) @ s \rightsquigarrow d) \wedge (s \sim (\text{the } (\text{domain } a)) \sim t)$   
 $\longrightarrow ((\text{step } s \text{ } a) \sim d \sim (\text{step } t \text{ } a))$

**definition** *dynamic-weakly-step-consistent-e* :: '*e*  $\Rightarrow$  *bool* **where**

*dynamic-weakly-step-consistent-e* *a*  $\equiv \forall d \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \sim d \sim t) \wedge$   
 $((\text{the } (\text{domain } a)) @ s \rightsquigarrow d) \wedge (s \sim (\text{the } (\text{domain } a)) \sim t)$   
 $\longrightarrow ((\text{step } s \text{ } a) \sim d \sim (\text{step } t \text{ } a))$

**lemma** *dynamic-weakly-step-consistent-all-evt* : *dynamic-weakly-step-consistent* =  $(\forall a. \text{dynamic-weakly-step-consistent-e } a)$

by (simp add: dynamic-weakly-step-consistent-def dynamic-weakly-step-consistent-e-def)

**definition** *dynamic-local-respect* :: bool **where**

*dynamic-local-respect*  $\equiv \forall a\ d\ s. \text{reachable0}\ s \wedge \neg((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d) \longrightarrow (s \sim d \sim (\text{step}\ s\ a))$

**definition** *dynamic-local-respect-e* :: 'e  $\Rightarrow$  bool **where**

*dynamic-local-respect-e*  $a \equiv \forall d\ s. \text{reachable0}\ s \wedge \neg((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d) \longrightarrow (s \sim d \sim (\text{step}\ s\ a))$

**lemma** *dynamic-local-respect-all-evt* : *dynamic-local-respect* =  $(\forall a. \text{dynamic-local-respect-e}\ a)$

by (simp add: dynamic-local-respect-def dynamic-local-respect-e-def)

**declare** *dynamic-step-consistent-def* [cong] **and** *dynamic-weakly-step-consistent-def* [cong] **and**  
*dynamic-local-respect-def* [cong]

**lemma** *step-cons-impl-weak* : *dynamic-step-consistent*  $\impl$  *dynamic-weakly-step-consistent*

using *dynamic-step-consistent-def* *dynamic-weakly-step-consistent-def* **by** blast

**definition** *lemma-local* :: bool **where**

*lemma-local*  $\equiv \forall s\ a\ as\ u. \text{the}(\text{domain}\ a) \notin \text{sources}\ (a \# as)\ u\ s \longrightarrow (s \approx (\text{sources}\ (a \# as)\ u\ s) \approx (\text{step}\ s\ a))$

**lemma** *weak-with-step-cons*:

**assumes** *p1*: *dynamic-weakly-step-consistent*

**and** *p2*: *dynamic-local-respect*

**shows** *dynamic-step-consistent*

**proof** –

{

fix *a d s t*

**have** *reachable0* *s*  $\wedge$  *reachable0* *t*  $\wedge$   $(s \sim d \sim t) \wedge$

$((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d) \longrightarrow (s \sim (\text{the}(\text{domain}\ a)) \sim t))$

$\longrightarrow ((\text{step}\ s\ a) \sim d \sim (\text{step}\ t\ a))$

**proof** –

{

assume *a0*: *reachable0* *s*

assume *a1*: *reachable0* *t*

assume *a2*:  $(s \sim d \sim t)$

assume *a3*:  $((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d) \longrightarrow (s \sim (\text{the}(\text{domain}\ a)) \sim t))$

**have**  $((\text{step}\ s\ a) \sim d \sim (\text{step}\ t\ a))$

**proof** (cases  $((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d)$ )

assume *b0*:  $((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d)$

**have** *b1*:  $(s \sim (\text{the}(\text{domain}\ a)) \sim t)$

using *b0* *a3* **by** auto

**have** *b2*:  $((\text{step}\ s\ a) \sim d \sim (\text{step}\ t\ a))$

using *a0* *a1* *a2* *b0* *b1* *p1* *dynamic-weakly-step-consistent-def* **by** blast

**then show** ?thesis **by** auto

**next**

assume *b0*:  $\neg((\text{the}(\text{domain}\ a)) @ s \rightsquigarrow d)$

**have** *b1*:  $\neg((\text{the}(\text{domain}\ a)) @ t \rightsquigarrow d)$

using *a0* *a1* *a2* *b0* *policy-respect* **by** auto

**have** *b2*:  $s \sim d \sim (\text{step}\ s\ a)$

using *b0* *p2* *a0* **by** auto

**have** *b3*:  $t \sim d \sim (\text{step}\ t\ a)$

using *b1* *p2* *a1* **by** auto

**have** *b4*:  $((\text{step}\ s\ a) \sim d \sim (\text{step}\ t\ a))$

using *b2* *b3* *a2* *vpeq-symmetric-lemma* *vpeq-transitive-lemma* **by** blast

**then show** ?thesis **by** auto

qed

}

**then show** ?thesis **by** auto

```

    qed
  }
  then show ?thesis by auto
qed

```

## 0.4 Lemmas for the inference framework

```

lemma sources-refl:reachable0 s  $\implies$  u  $\in$  sources as u s
  apply(induct as arbitrary: s)
  apply(simp add: sources-Nil)
  apply(simp add: sources-Cons)
  using enabled reachableStep
  by metis

```

```

lemma lemma-1-sub-1 :  $\llbracket$ reachable0 s ;
  dynamic-local-respect;
  the (domain a)  $\notin$  sources (a # as) u s;
  (s  $\approx$  (sources (a # as) u s)  $\approx$  t) $\rrbracket$ 
 $\implies$  (s  $\approx$  (sources as u (step s a))  $\approx$  (step s a))
  apply (simp add:dynamic-local-respect-def sources-Cons)
  by blast

```

```

lemma lemma-1-sub-2 :  $\llbracket$ reachable0 s ;
  reachable0 t ;
  dynamic-local-respect;
  the (domain a)  $\notin$  sources (a # as) u s;
  (s  $\approx$  (sources (a # as) u s)  $\approx$  t) $\rrbracket$ 
 $\implies$  (t  $\approx$  (sources as u (step s a))  $\approx$  (step t a))

```

proof –

```

  assume a1: reachable0 s
  assume a2: reachable0 t
  assume a3: dynamic-local-respect
  assume a6: the (domain a)  $\notin$  sources (a # as) u s
  assume a7: (s  $\approx$  (sources (a # as) u s)  $\approx$  t)
  have b1:  $\forall v. v \in \text{sources as u (step s a)} \longrightarrow \neg \text{interferes (the (domain a)) s v}$ 
    using a6 sources-Cons by auto
  have b2: sources (a # as) u s = sources as u (step s a)
    using a6 sources-Cons by auto
  have b3:  $\forall v. v \in \text{sources as u (step s a)} \longrightarrow (s \sim v \sim t)$ 
    using a7 b2 ivpeq-def by blast
  have b4:  $\forall v. v \in \text{sources as u (step s a)} \longrightarrow \neg \text{interferes (the (domain a)) t v}$ 
    using a1 a2 policy-respect b1 b3 by blast
  have b5:  $\forall v. v \in \text{sources as u (step s a)} \longrightarrow (t \sim v \sim (\text{step t a}))$ 
    using a2 a3 b4 by auto
  then show ?thesis
    using ivpeq-def by auto
qed

```

```

lemma lemma-1-sub-3 :  $\llbracket$ 
  the (domain a)  $\notin$  sources (a # as) u s;
  (s  $\approx$  (sources (a # as) u s)  $\approx$  t) $\rrbracket$ 
 $\implies$  (s  $\approx$  (sources as u (step s a))  $\approx$  t)
  apply (simp add:sources-Cons)
  apply (simp add:sources-Cons)
  done

```

```

lemma lemma-1-sub-4 :  $\llbracket$ (s  $\approx$  (sources as u (step s a))  $\approx$  t);

```

$(s \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } s \text{ } a));$   
 $(t \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } t \text{ } a)) \parallel$   
 $\implies ((\text{step } s \text{ } a) \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } t \text{ } a))$   
 by (meson ivpeq-def vpeq-symmetric-lemma vpeq-transitive-lemma)

**lemma** lemma-1 :  $\llbracket \text{reachable0 } s;$   
 $\text{reachable0 } t;$   
 $\text{dynamic-step-consistent};$   
 $\text{dynamic-local-respect};$   
 $(s \approx (\text{sources (} a \# \text{as) } u \text{ } s) \approx t) \rrbracket$   
 $\implies ((\text{step } s \text{ } a) \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } t \text{ } a))$   
**apply** (case-tac the (domain a)  $\in$  sources (a # as) u s)  
**apply** (simp add: dynamic-step-consistent-def)  
**apply** (simp add: sources-Cons)  
**proof** –  
 assume a1: dynamic-local-respect  
 assume a4: the (domain a)  $\notin$  sources (a # as) u s  
 assume a5:  $(s \approx (\text{sources (} a \# \text{as) } u \text{ } s) \approx t)$   
 assume b0: reachable0 s  
 assume b1: reachable0 t  
  
 have a6:  $(s \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx t)$   
 using a1 policy-respect a4 a5 lemma-1-sub-3 **by** auto  
 then have a7:  $(s \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } s \text{ } a))$   
 using b0 a1 policy-respect a4 a5 lemma-1-sub-1 **by** auto  
 then have a8:  $(t \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } t \text{ } a))$   
 using b1 b0 a1 policy-respect a4 a5 lemma-1-sub-2 **by** auto  
 then show  $((\text{step } s \text{ } a) \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } t \text{ } a))$   
 using a6 a7 lemma-1-sub-4 **by** blast  
**qed**

**lemma** lemma-2 :  $\llbracket \text{reachable0 } s;$   
 $\text{dynamic-local-respect};$   
 $\text{the (domain } a) \notin \text{sources (} a \# \text{as) } u \text{ } s \rrbracket$   
 $\implies (s \approx (\text{sources as } u \text{ (step } s \text{ } a)) \approx (\text{step } s \text{ } a))$   
**apply** (simp add: dynamic-local-respect-def)  
**apply** (simp add: sources-Cons)  
**by** blast

**lemma** sources-eq1:  $\forall s \ t \text{ as } u. \text{reachable0 } s \wedge$   
 $\text{reachable0 } t \wedge$   
 $\text{dynamic-step-consistent} \wedge$   
 $\text{dynamic-local-respect} \wedge$   
 $(s \approx (\text{sources as } u \text{ } s) \approx t)$   
 $\longrightarrow (\text{sources as } u \text{ } s) = (\text{sources as } u \text{ } t)$   
**proof** –  
 {  
 fix as  
 have  $\forall s \ t \ u. \text{reachable0 } s \wedge$   
 $\text{reachable0 } t \wedge$   
 $\text{dynamic-step-consistent} \wedge$   
 $\text{dynamic-local-respect} \wedge$   
 $(s \approx (\text{sources as } u \text{ } s) \approx t)$   
 $\longrightarrow (\text{sources as } u \text{ } s) = (\text{sources as } u \text{ } t)$   
**proof**(induct as)  
 case Nil then show ?case **by** (simp add: sources-Nil)  
**next**



```

case (Cons b bs)
assume p0:  $\forall s t u. ((\text{reachable0 } s) \wedge (\text{reachable0 } t) \wedge \text{dynamic-step-consistent} \wedge \text{dynamic-local-respect} \wedge (s \approx (\text{sources } bs \ u \ s) \approx t)) \longrightarrow (\text{sources } bs \ u \ s) = (\text{sources } bs \ u \ t))$ 

then show ?case
proof -
{
  fix s t u
  assume p1: reachable0 s
  assume p2: reachable0 t
  assume p3: dynamic-step-consistent
  assume p5: dynamic-local-respect
  assume p9:  $(s \approx (\text{sources } (b \# bs) \ u \ s) \approx t)$ 
  have a2:  $((\text{step } s \ b) \approx (\text{sources } bs \ u \ (\text{step } s \ b)) \approx (\text{step } t \ b))$ 
  using lemma-1 p1 p2 p3 policy-respect p5 p9 by blast
  have a3:  $\text{sources } (b \# bs) \ u \ s = \text{sources } (b \# bs) \ u \ t$ 
  proof (cases the (domain b)  $\in$  (sources (b # bs) u s))
  assume b0: the (domain b)  $\in$  (sources (b # bs) u s)
  have b1:  $s \sim (\text{the}(\text{domain } b)) \sim t$ 
  using b0 p9 by auto
  have b3:  $\text{interferes } (\text{the } (\text{domain } b)) \ s \ u = \text{interferes } (\text{the } (\text{domain } b)) \ t \ u$ 
  using p1 p2 policy-respect p9 sources-refl by fastforce
  have b4:  $(\text{sources } bs \ u \ (\text{step } s \ b)) = (\text{sources } bs \ u \ (\text{step } t \ b))$ 
  using a2 p0 p1 p2 p3 p5 reachableStep by blast
  have b5:  $\forall v. v \in \text{sources } bs \ u \ (\text{step } s \ b) \longrightarrow \text{interferes } (\text{the } (\text{domain } b)) \ s \ v = \text{interferes } (\text{the } (\text{domain } b)) \ t \ v$ 
  using p1 p2 ivpeq-def policy-respect p9 sources-Cons by fastforce
  then show sources (b # bs) u s = sources (b # bs) u t
  using b4 b5 sources-Cons by auto
  next
  assume b0: the (domain b)  $\notin$  (sources (b # bs) u s)
  have b1:  $\text{sources } (b \# bs) \ u \ s = \text{sources } bs \ u \ (\text{step } s \ b)$ 
  using b0 sources-Cons by auto
  have b2:  $(\text{sources } bs \ u \ (\text{step } s \ b)) = (\text{sources } bs \ u \ (\text{step } t \ b))$ 
  using a2 p0 p1 p2 p3 p5 reachableStep by blast
  have b3:  $\forall v. v \in \text{sources } bs \ u \ (\text{step } s \ b) \longrightarrow \neg \text{interferes } (\text{the } (\text{domain } b)) \ s \ v$ 
  using b0 sources-Cons by auto
  have b4:  $\forall v. v \in \text{sources } bs \ u \ (\text{step } s \ b) \longrightarrow \neg \text{interferes } (\text{the } (\text{domain } b)) \ t \ v$ 
  using b1 b3 p1 p2 p9 policy-respect by fastforce
  have b5:  $\forall v. v \in \text{sources } bs \ u \ (\text{step } t \ b) \longrightarrow \neg \text{interferes } (\text{the } (\text{domain } b)) \ t \ v$ 
  by (simp add: b2 b4)
  have b6: the (domain b)  $\notin$  (sources (b # bs) u t)
  using b0 b2 b5 sources.simps(2) by auto
  have b7:  $\text{sources } (b \# bs) \ u \ t = \text{sources } bs \ u \ (\text{step } t \ b)$ 
  using b6 sources-Cons by auto
  then show ?thesis
  by (simp add: b1 b2)
qed
}
then show ?thesis by blast
qed
}
then show ?thesis by blast

```

qed

**lemma** *ipurge-eq*:  $\forall s\ t\ as\ u. \text{reachable0}\ s \wedge$   
 $\text{reachable0}\ t \wedge$   
 $\text{dynamic-step-consistent} \wedge$   
 $\text{dynamic-local-respect} \wedge$   
 $(s \approx (\text{sources}\ as\ u\ s) \approx t)$   
 $\longrightarrow (\text{ipurge}\ as\ u\ s) = (\text{ipurge}\ as\ u\ t)$

**proof** –

{

fix *as*

have  $\forall s\ t\ u. \text{reachable0}\ s \wedge$

$\text{reachable0}\ t \wedge$

$\text{dynamic-step-consistent} \wedge$

$\text{dynamic-local-respect} \wedge$

$(s \approx (\text{sources}\ as\ u\ s) \approx t)$

$\longrightarrow (\text{ipurge}\ as\ u\ s) = (\text{ipurge}\ as\ u\ t)$

**proof**(*induct as*)

case *Nil* then show ?case **by** (*simp add: sources-Nil*)

**next**

case (*Cons b bs*)

assume *p0*:  $\forall s\ t\ u. ((\text{reachable0}\ s)$

$\wedge (\text{reachable0}\ t)$

$\wedge \text{dynamic-step-consistent}$

$\wedge \text{dynamic-local-respect}$

$\wedge (s \approx (\text{sources}\ bs\ u\ s) \approx t))$

$\longrightarrow (\text{ipurge}\ bs\ u\ s) = (\text{ipurge}\ bs\ u\ t)$

**then show** ?case

**proof** –

{

fix *s t u*

assume *p1*:  $\text{reachable0}\ s$

assume *p2*:  $\text{reachable0}\ t$

assume *p3*:  $\text{dynamic-step-consistent}$

assume *p5*:  $\text{dynamic-local-respect}$

assume *p9*:  $(s \approx (\text{sources}\ (b \# bs)\ u\ s) \approx t)$

have *a1*:  $((\text{step}\ s\ b) \approx (\text{sources}\ bs\ u\ (\text{step}\ s\ b)) \approx (\text{step}\ t\ b))$

using *lemma-1 p1 p2 p3 p5 p9* **by** *blast*

have *a2*:  $(\text{ipurge}\ bs\ u\ (\text{step}\ s\ b)) = (\text{ipurge}\ bs\ u\ (\text{step}\ t\ b))$

using *a1 p0 p1 p2 p3 p5 p9 reachableStep* **by** *blast*

have *a3*:  $\text{sources}\ (b \# bs)\ u\ s = \text{sources}\ (b \# bs)\ u\ t$

using *p1 p2 p3 p5 p9 sources-eq1* **by** *blast*

have *a4*:  $\text{ipurge}\ (b \# bs)\ u\ s = \text{ipurge}\ (b \# bs)\ u\ t$

**proof** (*cases the (domain b) ∈ (sources (b # bs) u s)*)

assume *b0*:  $\text{the}(\text{domain}\ b) \in (\text{sources}\ (b \# bs)\ u\ s)$

have *b1*:  $s \sim (\text{the}(\text{domain}\ b)) \sim t$

using *b0 p9* **by** *auto*

have *b3*:  $\text{the}(\text{domain}\ b) \in (\text{sources}\ (b \# bs)\ u\ t)$

using *a3 b0* **by** *auto*

**then show** ?thesis

using *a2 b0 ipurge-Cons* **by** *auto*

**next**

assume *b0*:  $\text{the}(\text{domain}\ b) \notin (\text{sources}\ (b \# bs)\ u\ s)$

have *b1*:  $\text{sources}\ (b \# bs)\ u\ s = \text{sources}\ bs\ u\ (\text{step}\ s\ b)$

using *b0 sources-Cons* **by** *auto*

have *b3*:  $\forall v. v \in \text{sources}\ bs\ u\ (\text{step}\ s\ b) \longrightarrow \neg \text{interferes}(\text{the}(\text{domain}\ b))\ s\ v$

using *b0 sources-Cons* **by** *auto*

have *b4*:  $\forall v. v \in \text{sources}\ bs\ u\ (\text{step}\ s\ b) \longrightarrow \neg \text{interferes}(\text{the}(\text{domain}\ b))\ t\ v$

```

    using b1 b3 p1 p2 p9 policy-respect by fastforce
    have b5: the (domain b)  $\notin$  (sources (b # bs) u t)
    using a3 b1 b4 interf-reflexive by auto
    have b6: ipurge (b # bs) u s = ipurge bs u (step s b)
    using b0 by auto
    have b7: ipurge (b # bs) u t = ipurge bs u (step t b)
    using b5 by auto
    then show ?thesis
    using b6 b7 a2 by auto
  qed
}
then show ?thesis by blast
qed
qed
}
then show ?thesis by blast
qed

```

**lemma** *non-influence-lemma*:  $\forall s t \text{ as } u. \text{reachable0 } s \wedge$   
 $\text{reachable0 } t \wedge$   
 $\text{dynamic-step-consistent} \wedge$   
 $\text{dynamic-local-respect} \wedge$   
 $(s \approx (\text{sources as } u s) \approx t)$   
 $\longrightarrow ((s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } u t) @ u))$

**proof** –

```

{
  fix as
  have  $\forall s t u. \text{reachable0 } s \wedge$   

 $\text{reachable0 } t \wedge$   

 $\text{dynamic-step-consistent} \wedge$   

 $\text{dynamic-local-respect} \wedge$   

 $(s \approx (\text{sources as } u s) \approx t)$   

 $\longrightarrow ((s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } u t) @ u))$ 

```

**proof** (*induct as*)

**case** *Nil* **show** ?case **using** *sources-Nil* **by** *auto*

**next**

**case** (*Cons b bs*)

```

assume p0:  $\forall s t u. ((\text{reachable0 } s)$   

 $\wedge (\text{reachable0 } t)$   

 $\wedge \text{dynamic-step-consistent}$   

 $\wedge \text{dynamic-local-respect}$   

 $\wedge (s \approx (\text{sources bs } u s) \approx t)) \longrightarrow$   

 $((s \triangleleft \text{bs} \cong t \triangleleft (\text{ipurge bs } u t) @ u))$ 

```

**then show** ?case

**proof** –

```

{
  fix s t u
  assume p1: reachable0 s
  assume p2: reachable0 t
  assume p3: dynamic-step-consistent
  assume p4: dynamic-local-respect
  assume p8:  $(s \approx (\text{sources (b # bs) } u s) \approx t)$ 
  have a1:  $((\text{step s b}) \approx (\text{sources bs } u (\text{step s b})) \approx (\text{step t b}))$ 
  using lemma-1 p1 p2 p3 p4 p8 by blast
  have  $s \triangleleft b \# \text{bs} \cong t \triangleleft \text{ipurge (b \# bs) } u t @ u$ 
  proof (cases the (domain b)  $\in$  sources (b # bs) u s)
    assume b0: the (domain b)  $\in$  sources (b # bs) u s
    have b1: interferes (the (domain b)) s u = interferes (the (domain b)) t u

```

```

    using p1 p2 policy-respect p8 sources-refl by fastforce
  have b2:  $\forall v. v \in \text{sources } bs \ u \ (\text{step } s \ b)$ 
     $\longrightarrow \text{interferes } (\text{the } (\text{domain } b)) \ s \ v = \text{interferes } (\text{the } (\text{domain } b)) \ t \ v$ 
    using p1 p2 ivpeq-def policy-respect p8 sources-Cons by fastforce
  have b3:  $\text{ipurge } (b \ \# \ bs) \ u \ t = b \ \# \ (\text{ipurge } bs \ u \ (\text{step } t \ b))$ 
    by (metis b0 ipurge-Cons p1 p2 p3 p4 p8 sources-eq1)
  have b4:  $((\text{step } s \ b) \triangleleft bs \cong (\text{step } t \ b) \triangleleft (\text{ipurge } bs \ u \ (\text{step } t \ b)) \ @ \ u)$ 
    using a1 p0 p1 p2 p3 p4 reachableStep by blast
  show ?thesis
    using b3 b4 by auto
next
  assume b0:  $\text{the } (\text{domain } b) \notin \text{sources } (b \ \# \ bs) \ u \ s$ 
  have b1:  $\text{ipurge } (b \ \# \ bs) \ u \ t = (\text{ipurge } bs \ u \ (\text{step } t \ b))$ 
    by (metis a1 b0 ipurge-Cons ipurge-eq p1 p2 p3 p4 p8 reachableStep)
  have b2:  $(s \approx (\text{sources } bs \ u \ (\text{step } s \ b)) \approx (\text{step } s \ b))$ 
    using b0 lemma-2 p1 p4 by blast
  have b3:  $(s \approx (\text{sources } bs \ u \ (\text{step } s \ b)) \approx t)$ 
    using b0 lemma-1-sub-3 p8 by blast
  have b4:  $((\text{step } s \ b) \approx (\text{sources } bs \ u \ (\text{step } s \ b)) \approx t)$ 
    by (meson b3 b2 ivpeq-def vpeq-symmetric-lemma vpeq-transitive-lemma)
  have b5:  $((\text{step } s \ b) \triangleleft bs \cong t \triangleleft (\text{ipurge } bs \ u \ t) \ @ \ u)$ 
    using b4 p0 p1 p2 p3 p4 reachableStep by blast
  have b6:  $(t \approx (\text{sources } bs \ u \ (\text{step } s \ b)) \approx (\text{step } t \ b))$ 
    using p1 p2 b0 lemma-1-sub-2 p4 p8 by blast
  have b7:  $\text{ipurge } bs \ u \ t = \text{ipurge } bs \ u \ (\text{step } t \ b)$ 
    by (metis a1 b4 ipurge-eq p1 p2 p3 p4 reachableStep)
  have b8:  $((\text{step } s \ b) \triangleleft bs \cong t \triangleleft (\text{ipurge } bs \ u \ (\text{step } t \ b)) \ @ \ u)$ 
    using b5 b7 by auto
  then show ?thesis
    using b1 observ-equivalence-def run-Cons by auto
qed
}
then show ?thesis by blast
qed
}
then show ?thesis by blast
qed
}
then show ?thesis by blast
qed

```

## 0.5 Interference framework of information flow security properties

```

theorem nonintf-impl-weak: noninterference  $\implies$  weak-noninterference
  by (metis noninterference-def observ-equiv-sym observ-equiv-trans reachable-s0 weak-noninterference-def)

theorem wk-nonintf-r-impl-wk-nonintf: weak-noninterference-r  $\implies$  weak-noninterference
  using reachable-s0 by auto

theorem nonintf-r-impl-nonintf: noninterference-r  $\implies$  noninterference
  using noninterference-def noninterference-r-def reachable-s0 by auto

theorem nonintf-r-impl-wk-nonintf-r: noninterference-r  $\implies$  weak-noninterference-r
  by (metis noninterference-r-def observ-equiv-sym observ-equiv-trans weak-noninterference-r-def)

lemma noninf-impl-nonintf-r: noninfluence  $\implies$  noninterference-r
  using ivpeq-def noninfluence-def noninterference-r-def vpeq-reflexive-lemma by blast

lemma noninf-impl-nonlk: noninfluence  $\implies$  nonleakage
  using noninterference-r-def nonleakage-def observ-equiv-sym

```

*observ-equiv-trans noninfluence-def noninf-impl-nonintf-r* **by** *blast*

**lemma** *wk-noninfl-impl-nonlk*: *weak-noninfluence*  $\implies$  *nonleakage*  
**using** *weak-noninfluence-def nonleakage-def* **by** *blast*

**lemma** *wk-noninfl-impl-wk-nonintf-r*: *weak-noninfluence*  $\implies$  *weak-noninterference-r*  
**using** *ivpeq-def weak-noninfluence-def vpeq-reflexive-lemma weak-noninterference-r-def* **by** *blast*

**lemma** *sources-step2*:  
 $\llbracket \text{reachable0 } s; (\text{the } (\text{domain } a)) @ s \rightsquigarrow d \rrbracket \implies \text{sources } [a] \text{ } d \text{ } s = \{\text{the } (\text{domain } a), d\}$   
**apply**(*auto simp: sources-Cons sources-Nil enabled dest: enabled*)  
**done**

**lemma** *exec-equiv-both*:  
 $\llbracket \text{reachable0 } C1; \text{reachable0 } C2; (\text{step } C1 \text{ } a) \triangleleft as \cong (\text{step } C2 \text{ } b) \triangleleft bs @ u \rrbracket$   
 $\implies (C1 \triangleleft (a \# as) \cong C2 \triangleleft (b \# bs) @ u)$   
**by** *auto*

**lemma** *sources-unwinding-step*:  
 $\llbracket \text{reachable0 } s; \text{reachable0 } t; s \approx (\text{sources } (a \# as) \text{ } d \text{ } s) \approx t; \text{dynamic-step-consistent} \rrbracket$   
 $\implies ((\text{step } s \text{ } a) \approx (\text{sources } as \text{ } d \text{ } (\text{step } s \text{ } a)) \approx (\text{step } t \text{ } a))$   
**apply**(*clarsimp simp: ivpeq-def sources-Cons*)  
**using** *UnionI dynamic-step-consistent-def* **by** *blast*

**lemma** *nonlk-imp-sc*: *nonleakage*  $\implies$  *dynamic-step-consistent*

**proof** –

**assume** *p0*: *nonleakage*

**have** *p1*:  $\forall as \text{ } d \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t$

$\wedge (s \approx (\text{sources } as \text{ } d \text{ } s) \approx t) \longrightarrow (s \triangleleft as \cong t \triangleleft as @ d)$

**using** *p0 nonleakage-def* **by** *auto*

**have** *p2*:  $\forall a \text{ } d \text{ } s \text{ } t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \sim d \sim t) \wedge$

$((\text{the } (\text{domain } a)) @ s \rightsquigarrow d) \longrightarrow (s \sim (\text{the } (\text{domain } a)) \sim t)$

$\longrightarrow ((\text{step } s \text{ } a) \sim d \sim (\text{step } t \text{ } a))$

**proof** –

{

**fix** *a d s t*

**assume** *a0*:  $\text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \sim d \sim t) \wedge$

$((\text{the } (\text{domain } a)) @ s \rightsquigarrow d) \longrightarrow (s \sim (\text{the } (\text{domain } a)) \sim t)$

**have** *a4*:  $s \approx (\text{sources } [] \text{ } d \text{ } s) \approx t$

**using** *a0 sources-Nil* **by** *auto*

**have** *a5*:  $(s \triangleleft [] \cong t \triangleleft [] @ d)$

**using** *a4 a0 p1* **by** *auto*

**have** *a6*:  $((\text{step } s \text{ } a) \sim d \sim (\text{step } t \text{ } a))$

**proof** (*cases (the (domain a)) @ s  $\rightsquigarrow$  d*)

**assume** *b0*:  $(\text{the } (\text{domain } a)) @ s \rightsquigarrow d$

**have** *b1*:  $\text{sources } [a] \text{ } d \text{ } s = \{d, (\text{the } (\text{domain } a))\}$

**using** *b0 sources-Cons sources-Nil* **by** *auto*

**have** *c0*:  $(s \sim (\text{the } (\text{domain } a)) \sim t)$

**using** *b0 a0* **by** *auto*

**have** *b2*:  $s \approx (\text{sources } [a] \text{ } d \text{ } s) \approx t$

**using** *b1 a0 c0* **by** *auto*

**have** *b3*:  $(s \triangleleft [a] \cong t \triangleleft [a] @ d)$

**using** *b2 a0 p1* **by** *auto*

**have** *b4*:  $((\text{step } s \text{ } a) \sim d \sim (\text{step } t \text{ } a))$

**using** *b3* **by** *auto*

**then show** *?thesis* **by** *auto*

**next**

**assume** *b0*:  $\neg((\text{the } (\text{domain } a)) @ s \rightsquigarrow d)$

```

have b1: sources [a] d s = {d}
  using b0 sources-Cons sources-Nil by auto
have b2: (s ≈ (sources [a] d s) ≈ t)
  using b1 a0 by auto
have b3: (s < [a] ≅ t < [a] @ d)
  using b2 a0 p1 by auto
have b4: ((step s a) ~ d ~ (step t a))
  using b3 by auto
then show ?thesis by auto
qed
}
then show ?thesis
  by auto
qed
then show ?thesis by auto
qed

lemma sc-imp-nonlk: dynamic-step-consistent ⟹ nonleakage
proof -
  assume p0: dynamic-step-consistent
  have p1: ∀ a d s t. reachable0 s ∧ reachable0 t ∧ (s ~ d ~ t) ∧
    (s ~ (the (domain a)) ~ t) ⟶ ((step s a) ~ d ~ (step t a))
    using p0 dynamic-step-consistent-def by auto
  have p2: ∀ as d s t. reachable0 s ∧ reachable0 t
    ∧ (s ≈ (sources as d s) ≈ t) ⟶ (s < as ≅ t < as @ d)
  proof -
    {
      fix as
      have ∀ d s t. reachable0 s ∧ reachable0 t
        ∧ (s ≈ (sources as d s) ≈ t) ⟶ (s < as ≅ t < as @ d)
      proof (induct as)
        case Nil show ?case using sources-refl by auto
      next
        case (Cons b bs)
        assume a0: ∀ d s t. reachable0 s ∧ reachable0 t
          ∧ (s ≈ (sources bs d s) ≈ t) ⟶ (s < bs ≅ t < bs @ d)
        show ?case
          proof -
            {
              fix d s t
              assume b0: reachable0 s ∧ reachable0 t
              assume b1: (s ≈ (sources (b#bs) d s) ≈ t)
              have b2: ((step s b) ≈ (sources bs d (step s b)) ≈ (step t b))
                using b0 b1 p0 sources-unwinding-step by auto
              have b3: (step s b) < bs ≅ (step t b) < bs @ d
                using Cons.hyps b0 b2 reachableStep by blast
              have b4: s < b # bs ≅ t < b # bs @ d
                using b3 by auto
            }
            then show ?thesis by auto
          qed
        qed
      }
    }
  then show ?thesis by auto
qed
then show ?thesis by auto
qed

```

**theorem** *sc-eq-nonlk*: *dynamic-step-consistent* = *nonleakage*  
**using** *nonlk-imp-sc* *sc-imp-nonlk* **by** *blast*

**lemma** *noninf-imp-dlr*: *noninfluence*  $\implies$  *dynamic-local-respect*  
**proof** –

**assume** *p0*: *noninfluence*  
**have** *p1*:  $\forall d \text{ as } s \ t. \text{reachable0 } s \wedge \text{reachable0 } t$   
 $\wedge (s \approx (\text{sources as } d \ s) \approx t)$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } d \ t) @ d)$   
**using** *p0* *noninfluence-def* **by** *auto*  
**have**  $\forall a \ d \ s. \text{reachable0 } s \wedge \neg((\text{the } (\text{domain } a)) @ s \rightsquigarrow d)$   
 $\longrightarrow (s \sim d \sim (\text{step } s \ a))$   
**proof** –  
{  
**fix** *a d s*  
**assume** *a0*: *reachable0* *s*  $\wedge \neg((\text{the } (\text{domain } a)) @ s \rightsquigarrow d)$   
**have** *a1*: *sources* [*a*] *d s* = {*d*}  
**using** *a0* *sources-Cons* *sources-Nil* **by** *auto*  
**have** *a2*: (*ipurge* [*a*] *d s*) = []  
**using** *a0* *a1* *interf-reflexive* **by** *auto*  
**have** *a3*: *s*  $\sim d \sim s$   
**using** *vpeq-reflexive-lemma* **by** *auto*  
**have** *a4*: (*s*  $\approx (\text{sources } [a] \ d \ s) \approx s$ )  
**using** *a1* *a3* **by** *auto*  
**have** *a5*: (*s*  $\triangleleft [a] \cong s \triangleleft (\text{ipurge } [a] \ d \ s) @ d$ )  
**using** *a4* *a0* *p1* **by** *auto*  
**have** *a6*: (*s*  $\triangleleft [a] \cong s \triangleleft [] @ d$ )  
**using** *a5* *a2* **by** *auto*  
**have** *a7*: (*s*  $\sim d \sim (\text{step } s \ a)$ )  
**using** *a6* *vpeq-symmetric-lemma* **by** *auto*  
}  
**then show** *?thesis* **by** *auto*  
**qed**  
**then show** *?thesis* **by** *auto*  
**qed**

**lemma** *noninf-imp-sc*: *noninfluence*  $\implies$  *dynamic-step-consistent*  
**using** *nonlk-imp-sc* *noninf-impl-nonlk* **by** *blast*

**theorem** *UnwindingTheorem* :  $\llbracket \text{dynamic-step-consistent};$   
 $\text{dynamic-local-respect} \rrbracket$   
 $\implies \text{noninfluence}$

**proof** –  
**assume** *p3*: *dynamic-step-consistent*  
**assume** *p4*: *dynamic-local-respect*  
{  
**fix** *as d*  
**have**  $\forall s \ t. \text{reachable0 } s \wedge$   
 $\text{reachable0 } t \wedge$   
 $(s \approx (\text{sources as } d \ s) \approx t)$   
 $\longrightarrow ((s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } d \ t) @ d))$   
**proof**(*induct as*)  
**case** *Nil* **show** *?case* **using** *sources-Nil* **by** *auto*  
**next**  
**case** (*Cons b bs*)  
**assume** *p0*:  $\forall s \ t. \text{reachable0 } s \wedge$   
 $\text{reachable0 } t \wedge$   
 $(s \approx (\text{sources } bs \ d \ s) \approx t)$

```

    → ((s < bs ≅ t < (ipurge bs d t) @ d))
  then show ?case
  proof -
  {
    fix s t
    assume p1: reachable0 s
    assume p2: reachable0 t
    assume p8: (s ≈ (sources (b # bs) d s) ≈ t)
    have a1: ((step s b) ≈ (sources bs d (step s b)) ≈ (step t b))
      using lemma-1 p1 p2 p3 p4 p8 by blast
    have a2: s < b # bs ≅ t < ipurge (b # bs) d t @ d
    proof (cases the (domain b) ∈ sources (b # bs) d s)
      assume b0: the (domain b) ∈ sources (b # bs) d s
      have b1: interferes (the (domain b)) s d = interferes (the (domain b)) t d
        using p1 p2 policy-respect p8 sources-refl by fastforce
      have b2: ∀ v. v ∈ sources bs d (step s b)
        → interferes (the (domain b)) s v = interferes (the (domain b)) t v
        using p1 p2 ivpeq-def policy-respect p8 sources-Cons by fastforce
      have b3: ipurge (b # bs) d t = b # (ipurge bs d (step t b))
        by (metis b0 ipurge-Cons p1 p2 p3 p4 p8 sources-eq1)
      have b4: (((step s b) < bs ≅ (step t b) < (ipurge bs d (step t b)) @ d))
        using a1 p0 p1 p2 p3 p4 reachableStep by blast
      then show ?thesis
        using b3 b4 by auto
    next
      assume b0: the (domain b) ∉ sources (b # bs) d s
      have b1: ipurge (b # bs) d t = (ipurge bs d (step t b))
        by (metis a1 b0 ipurge-Cons ipurge-eq p1 p2 p3 p4 p8 reachableStep)
      have b2: (s ≈ (sources bs d (step s b)) ≈ (step s b))
        using b0 lemma-2 p1 p4 by blast
      have b3: (s ≈ (sources bs d (step s b)) ≈ t)
        using b0 lemma-1-sub-3 p8 by blast
      have b4: ((step s b) ≈ (sources bs d (step s b)) ≈ t)
        by (meson b3 b2 ivpeq-def vpeq-symmetric-lemma vpeq-transitive-lemma)
      have b5: (((step s b) < bs ≅ t < (ipurge bs d t) @ d))
        using b4 p0 p1 p2 p3 p4 reachableStep by blast
      have b6: (t ≈ (sources bs d (step s b)) ≈ (step t b))
        using p1 p2 b0 lemma-1-sub-2 p4 p8 by blast
      have b7: ipurge bs d t = ipurge bs d (step t b)
        by (metis a1 b4 ipurge-eq p1 p2 p3 p4 reachableStep)
      have b8: (((step s b) < bs ≅ t < (ipurge bs d (step t b)) @ d))
        using b5 b7 by auto
      then show ?thesis
        using b1 observ-equivalence-def run-Cons by auto
    qed
  }
  then show ?thesis by blast
qed
qed
}
then show ?thesis using noninfluence-def by blast
qed

theorem UnwindingTheorem1 : [[dynamic-weakly-step-consistent;
dynamic-local-respect]] ⇒ noninfluence
using UnwindingTheorem weak-with-step-cons by blast

theorem uc-eq-noninf : (dynamic-step-consistent ∧ dynamic-local-respect) = noninfluence

```



using *UnwindingTheorem1* *step-cons-impl-weak* *noninf-imp-dlr* *noninf-imp-sc* **by** *blast*

**theorem** *noninf-impl-weak:noninfluence*  $\implies$  *weak-noninfluence*

**proof** –

assume *p0*: *noninfluence*

have *p1*:  $\forall d \text{ as } s \ t. \text{reachable0 } s \wedge \text{reachable0 } t$   
 $\wedge (s \approx (\text{sources as } d \ s) \approx t)$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } d \ t) @ d)$

using *p0* *noninfluence-def* **by** *auto*

have *p2*: (*dynamic-step-consistent*  $\wedge$  *dynamic-local-respect*)

using *p0* *uc-eq-noninf* **by** *auto*

have  $\forall d \text{ as } bs \ s \ t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \approx (\text{sources as } d \ s) \approx t)$   
 $\wedge \text{ipurge as } d \ t = \text{ipurge bs } d \ t$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft bs @ d)$

**proof** –

{

fix *d as bs s t*

assume *a0*: *reachable0 s*  $\wedge$  *reachable0 t*  $\wedge (s \approx (\text{sources as } d \ s) \approx t)$   
 $\wedge \text{ipurge as } d \ t = \text{ipurge bs } d \ t$

have *a4*: *noninterference-r*

using *noninf-impl-nonintf-r* *p0* **by** *auto*

have *a7*: *weak-noninterference-r*

using *a4* *nonintf-r-impl-wk-nonintf-r* **by** *auto*

have *a6*: *ipurge as d s* = *ipurge as d t*

using *a0* *p2* *ipurge-eq* **by** *auto*

have *b1*:  $(s \triangleleft \text{as} \cong t \triangleleft (\text{ipurge as } d \ t) @ d)$

using *a0* *p1* **by** *auto*

have *b4*:  $(s \triangleleft \text{as} \cong t \triangleleft \text{as} @ d)$

using *a0* *noninf-imp-sc* *nonleakage-def* *p0* *sc-imp-nonlk* **by** *blast*

have *b5*:  $(t \triangleleft bs \cong t \triangleleft (\text{ipurge bs } d \ t) @ d)$

using *a0* *a4* **by** *auto*

have *b6*:  $(t \triangleleft bs \cong t \triangleleft (\text{ipurge as } d \ t) @ d)$

using *b5* *a0* **by** *auto*

have *b7*:  $(s \triangleleft \text{as} \cong t \triangleleft bs @ d)$

using *a0* *b1* *b6* *observ-equiv-sym* *observ-equiv-trans* **by** *blast*

}

then show *?thesis* **by** *auto*

qed

then show *?thesis* **by** *auto*

qed

**lemma** *wk-nonintf-r-and-nonlk-impl-noninfl*:  $\llbracket \text{weak-noninterference-r}; \text{nonleakage} \rrbracket \implies \text{weak-noninfluence}$

**proof** –

assume *p0*: *weak-noninterference-r*

assume *p1*: *nonleakage*

have *p2*:  $\forall d \text{ as } bs \ s. \text{reachable0 } s \wedge \text{ipurge as } d \ s = \text{ipurge bs } d \ s$   
 $\longrightarrow (s \triangleleft \text{as} \cong s \triangleleft bs @ d)$

using *weak-noninterference-r-def* *p0* **by** *auto*

have *p3*:  $\forall d \text{ as } s \ t. \text{reachable0 } s \wedge \text{reachable0 } t$

$\wedge (s \approx (\text{sources as } d \ s) \approx t) \longrightarrow (s \triangleleft \text{as} \cong t \triangleleft \text{as} @ d)$

using *nonleakage-def* *p1* **by** *auto*

have  $\forall d \text{ as } bs \ s \ t. \text{reachable0 } s \wedge \text{reachable0 } t \wedge (s \approx (\text{sources as } d \ s) \approx t)$   
 $\wedge \text{ipurge as } d \ t = \text{ipurge bs } d \ t$   
 $\longrightarrow (s \triangleleft \text{as} \cong t \triangleleft bs @ d)$

**proof** –

```

{
  fix d as bs s t
  assume a0: reachable0 s ∧ reachable0 t ∧ (s ≈ (sources as d s) ≈ t)
    ∧ ipurge as d t = ipurge bs d t
  have a1: s < as ≅ t < as @ d
    using a0 p3 by blast
  have a2: t < as ≅ t < bs @ d
    using a0 p2 by auto
  have a3: (s < as ≅ t < bs @ d)
    using a0 a1 a2 observ-equiv-trans by blast
}
then show ?thesis by auto
qed
then show ?thesis by auto
qed

lemma nonintf-r-and-nonlk-impl-noninfl: [noninterference-r; nonleakage] ⇒ noninfluence
proof -
  assume p0: noninterference-r
  assume p1: nonleakage
  have p2: ∀ d as s. reachable0 s → (s < as ≅ s < (ipurge as d s) @ d)
    using p0 noninterference-r-def by auto
  have p3: ∀ d as s t. reachable0 s ∧ reachable0 t
    ∧ (s ≈ (sources as d s) ≈ t) → (s < as ≅ t < as @ d)
    using p1 nonleakage-def by auto
  have ∀ d as s t. reachable0 s ∧ reachable0 t
    ∧ (s ≈ (sources as d s) ≈ t)
    → (s < as ≅ t < (ipurge as d t) @ d)

proof -
{
  fix d as bs s t
  assume a0: reachable0 s ∧ reachable0 t
    ∧ (s ≈ (sources as d s) ≈ t)
  have a1: s < as ≅ t < as @ d
    using p3 a0 by blast
  have a2: s < as ≅ s < (ipurge as d s) @ d
    using a0 p2 by fast
  have a3: t < as ≅ t < (ipurge as d t) @ d
    using a0 p2 by fast
  have s < as ≅ t < (ipurge as d t) @ d
    using a0 a1 a3 observ-equiv-trans by blast
}
then show ?thesis by auto
qed
then show ?thesis using noninfluence-def by blast
qed

theorem nonintf-r-and-nonlk-eq-strnoninfl: (noninterference-r ∧ nonleakage) = noninfluence
using nonintf-r-and-nonlk-impl-noninfl noninf-impl-nonintf-r noninf-impl-nonlk by blast

end
end
theory CapFlow
imports Dynamic-model
begin

```

## 0.6 Definitions

**type-synonym** *max-buffer-size* = *nat*  
**type-synonym** *buffer-size* = *nat*  
**typeddecl** *Message*  
**typeddecl** *Endpoint*  
**typeddecl** *Domain*

**type-synonym** *domain-id* = *nat*  
**type-synonym** *domain-name* = *string*

**type-synonym** *endpoint-name* = *string*  
**type-synonym** *endpoint-id* = *nat*

**datatype**  
*right* = *SEND*  
| *TAKE*  
| *GRANT*  
| *REMOVE*

**record** *cap* = *target* :: *domain-id*  
*rights* :: *right set*

**record** *Endpoint-Concig* =  
*e-max-buf-size* :: *endpoint-id*  $\rightarrow$  *max-buffer-size*  
*e-name* :: *endpoint-id*  $\rightarrow$  *endpoint-name*  
*e-listener* :: *endpoint-id*  $\rightarrow$  *domain-id*

**record** *Domain-Config* =  
*d-name* :: *domain-id*  $\rightarrow$  *domain-name*  
*d-ep-set* :: *domain-id*  $\Rightarrow$  *endpoint-id set*

**record** *Sys-Config* =  
*domconf* :: *Domain-Config*  
*commconf* :: *Endpoint-Concig*

### 0.6.1 System state

**record** *State* =

*caps* :: *domain-id*  $\Rightarrow$  *cap set*  
*e-msgs* :: *endpoint-id*  $\Rightarrow$  *Message set*  
*e-buf-size* :: *endpoint-id*  $\rightarrow$  *buffer-size*  
*domain-endpoint* :: *endpoint-id*  $\rightarrow$  *domain-id*

**datatype** *Event* = *Client-Lookup-Endpoint-Name domain-id endpoint-name*  
| *Send-Queuing-Message domain-id endpoint-id Message*  
| *Receive-Queuing-Message domain-id endpoint-id*  
| *Get-My-Endpoints-Set domain-id*  
| *Get-Caps domain-id*  
| *Grant-Endpoint-Cap domain-id cap cap*  
| *Remove-Cap-Right domain-id cap right*

### 0.6.2 Utility Functions used for Event Specification

**definition** *get-domain-name-from-domain-id* :: *Sys-Config*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *domain-name option*  
**where** *get-domain-name-from-domain-id* *sc did*  $\equiv$   
*let*  
*dm-conf* = (*domconf sc*)

```

in
if((d-name dm-conf) did ≠ None )
then
  ((d-name dm-conf) did )
else
  None

```

**definition** *get-endpoints-from-domain-id* :: *Sys-Config* ⇒ *domain-id* ⇒ *endpoint-id* set  
**where** *get-endpoints-from-domain-id* sc did ≡  
 let  
 dm-conf = (domconf sc)  
 in  
 if((d-name dm-conf) did ≠ None )  
 then  
 ((d-ep-set dm-conf) did )  
 else  
 {}

**definition** *get-domain-id-from-domain-name* :: *Sys-Config* ⇒ *domain-name* ⇒ *domain-id* option  
**where** *get-domain-id-from-domain-name* sc dname ≡  
 let  
 dm-conf = (domconf sc)  
 in  
 if(∃ did. the((d-name dm-conf) did) = dname )  
 then  
 Some (SOME did. the((d-name dm-conf) did) = dname)  
 else  
 None

**definition** *get-listener-id-from-endpoint-id* :: *Sys-Config* ⇒ *endpoint-id* ⇒ *domain-id* option  
**where** *get-listener-id-from-endpoint-id* sc eid ≡  
 let  
 ep-conf = (commconf sc)  
 in  
 if((e-listener ep-conf) eid ≠ None )  
 then  
 ((e-listener ep-conf) eid )  
 else  
 None

**definition** *get-endpoint-name-from-endpoint-id* :: *Sys-Config* ⇒ *endpoint-id* ⇒ *endpoint-name* option  
**where** *get-endpoint-name-from-endpoint-id* sc eid ≡  
 let  
 ep-conf = (commconf sc)  
 in  
 if((e-name ep-conf) eid ≠ None )  
 then  
 ((e-name ep-conf) eid )  
 else  
 None

**definition** *get-endpoint-id-from-endpoint-name* :: *Sys-Config* ⇒ *endpoint-name* ⇒ *endpoint-id* option  
**where** *get-endpoint-id-from-endpoint-name* sc ename ≡  
 let  
 ep-conf = (commconf sc)  
 in  
 if(∃ eid. the((e-name ep-conf) eid) = ename )  
 then

```

    Some (SOME eid. the((e-name ep-conf) eid) = ename )
  else
    None

```

**definition** *get-endpoint-max-bufsize-from-sc-by-id* :: *Sys-Config*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *max-buffer-size option*

**where** *get-endpoint-max-bufsize-from-sc-by-id* *sc* *eid*  $\equiv$

```

  let
    ep-conf = (commconf sc)
  in
    if((e-max-buf-size ep-conf) eid  $\neq$  None )
  then
    ((e-max-buf-size ep-conf) eid )
  else
    None

```

**definition** *is-a-domain* :: *Sys-Config*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *bool*

**where** *is-a-domain* *sc* *did*  $\equiv$

```

  let
    dm-conf = (domconf sc)
  in
    if((d-name dm-conf) did  $\neq$  None )
  then
    True
  else
    False

```

**definition** *is-an-endpoint* :: *Sys-Config*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *bool*

**where** *is-an-endpoint* *sc* *eid*  $\equiv$

```

  let
    ep-conf = (commconf sc)
  in
    if((e-name ep-conf) eid  $\neq$  None )
  then
    True
  else
    False

```

**definition** *is-an-endpoint-of-domain* :: *Sys-Config*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *bool*

**where** *is-an-endpoint-of-domain* *sc* *eid* *did*  $\equiv$

```

  let
    dm-conf = (domconf sc)
  in
    if(eid  $\in$  ((d-ep-set dm-conf) did))
  then
    True
  else
    False

```

**definition** *is-an-endpoint-of-listener* :: *Sys-Config*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *bool*

**where** *is-an-endpoint-of-listener* *sc* *eid* *did*  $\equiv$

```

  let
    ep-conf = (commconf sc)
  in
    if(the((e-listener ep-conf) eid) = did )
  then
    True
  else
    False

```

**definition** *get-domain-cap-set-from-domain-id* :: *State*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *cap set*  
**where** *get-domain-cap-set-from-domain-id* *s* *did*  $\equiv$   
 let  
   *cap-by-id* = *caps s*  
 in  
   *cap-by-id did*

**definition** *get-msg-set-from-endpoint-id* :: *State*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *Message set*  
**where** *get-msg-set-from-endpoint-id* *s* *eid*  $\equiv$   
 let  
   *msg-by-id* = *e-msgs s*  
 in  
   *msg-by-id eid*

**definition** *get-buf-size-from-endpoint-id* :: *State*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *buffer-size option*  
**where** *get-buf-size-from-endpoint-id* *s* *eid*  $\equiv$   
 let  
   *buf-size-by-id* = *e-buf-size s*  
 in  
   *buf-size-by-id eid*

**definition** *endpoint-is-full* :: *Sys-Config*  $\Rightarrow$  *State*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *bool*  
**where** *endpoint-is-full* *sc* *s* *eid*  $\equiv$   
   *get-buf-size-from-endpoint-id s eid*  
   = *get-endpoint-max-bufsize-from-sc-by-id sc eid*

**definition** *get-endpoints-of-domain* :: *State*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *endpoint-id set*  
**where** *get-endpoints-of-domain* *s* *did*  $\equiv$   
 let  
   *dom-ep* = *domain-endpoint s*  
 in  
   {*x*. *dom-ep x* = *Some did*}

**definition** *get-domain-id-from-endpoint* :: *State*  $\Rightarrow$  *endpoint-id*  $\Rightarrow$  *domain-id option*  
**where** *get-domain-id-from-endpoint* *s* *eid*  $\equiv$   
 let  
   *dom-ep* = *domain-endpoint s*  
 in  
   *dom-ep eid*

**definition** *get-endpoint-msg-set-of-domain* :: *State*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *Message set*  
**where** *get-endpoint-msg-set-of-domain* *s* *did*  $\equiv$   
 let  
   *dom-ep* = *domain-endpoint s*;  
   *eid-set* = *get-endpoints-of-domain s did*;  
   *msg-of-ep* = *e-msgs s*  
 in  
   {*x*.  $\exists e. e \in \text{eid-set} \wedge x \in \text{msg-of-ep } e$ }

**definition** *interferes* :: *domain-id*  $\Rightarrow$  *State*  $\Rightarrow$  *domain-id*  $\Rightarrow$  *bool*  
**where** *interferes* *w* *s* *v*  $\equiv$

```

if( w = v
  ∨ (∃ c. c ∈ (get-domain-cap-set-from-domain-id s w) ∧ target c = v))
then
  True
else
  False

```

### 0.6.3 Event specification

**definition** *client-lookup-endpoint-name* :: *Sys-Config* ⇒ *State*  
 ⇒ *domain-id* ⇒ *endpoint-name* ⇒ (*State* × *endpoint-id option*)

**where** *client-lookup-endpoint-name* *sc s did ename* ≡  
 if(*get-endpoint-id-from-endpoint-name* *sc ename* ≠ *None*)  
 then  
 (s, *get-endpoint-id-from-endpoint-name* *sc ename*)  
 else  
 (s, *None*)

**definition** *send-queuing-message* :: *State* ⇒ *domain-id* ⇒ *endpoint-id* ⇒ *Message* ⇒ (*State* × *bool*)

**where** *send-queuing-message* *s did eid m* ≡  
 let  
   *dom-ep* = *domain-endpoint* *s*;  
   *dst-dom* = *dom-ep* *eid*;  
   *emsgs* = *e-msgs* *s*;  
   *msg-set* = *get-msg-set-from-endpoint-id* *s eid*;  
   *new-msg-set* = *insert* *m msg-set*  
 in  
 if(*get-domain-id-from-endpoint* *s eid* ≠ *None*  
   ∧ *interferes* *did s* (*the* (*get-domain-id-from-endpoint* *s eid*)))  
 then  
 (s(  
   *e-msgs* := *emsgs*(*eid* := *new-msg-set*)  
 ), *False*)  
 else  
 (s, *False*)

**definition** *receive-queuing-message* :: *State* ⇒ *domain-id* ⇒ *endpoint-id* ⇒ (*State* × *Message option*)

**where** *receive-queuing-message* *s did eid* ≡  
 if(*get-domain-id-from-endpoint* *s eid* = *Some did*)  
 then  
 let  
   *emsgs* = *e-msgs* *s*;  
   *msg-set* = *get-msg-set-from-endpoint-id* *s eid*;  
   *m* = *SOME* *x*. *x* ∈ *msg-set*;  
   *new-msg-set* = *msg-set* − {*m*}  
 in  
 (s(  
   *e-msgs* := *emsgs*(*eid* := *new-msg-set*)  
 ), *Some m*)  
 else  
 (s, *None*)

**definition** *get-my-endpoints-set* :: *State* ⇒ *domain-id* ⇒ (*State* × *endpoint-id set*)

**where** *get-my-endpoints-set* *s did* ≡  
 if(*get-endpoints-of-domain* *s did* ≠ {})  
 then

$(s, \text{get-endpoints-of-domain } s \text{ did})$   
 else  
 $(s, \{\})$

**definition**  $\text{get-caps} :: \text{State} \Rightarrow \text{domain-id} \Rightarrow (\text{State} \times \text{cap set})$   
**where**  $\text{get-caps } s \text{ did} \equiv$   
 $(s, \text{get-domain-cap-set-from-domain-id } s \text{ did})$

**definition**  $\text{grant-endpoint-cap} :: \text{State} \Rightarrow \text{domain-id} \Rightarrow \text{cap} \Rightarrow \text{cap} \Rightarrow (\text{State} \times \text{bool})$   
**where**  $\text{grant-endpoint-cap } s \text{ did grant-cap dst-cap} \equiv$   
 $\text{if}(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 $\wedge \text{GRANT} \in \text{rights grant-cap}$   
 $\wedge \text{target grant-cap} \neq \text{target dst-cap}$   
 $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did})$   
 then  
 let  
 $\text{did-dst} = \text{target grant-cap};$   
 $\text{caps0} = \text{caps } s;$   
 $\text{cs-dst} = \text{get-domain-cap-set-from-domain-id } s \text{ did-dst}$   
 in  
 $(s[$   
 $\text{caps} := \text{caps0}(\text{did-dst} := (\text{insert dst-cap cs-dst}))$   
 $], \text{True})$   
 else  
 $(s, \text{False})$

**definition**  $\text{get-takable-caps} :: \text{State} \Rightarrow \text{domain-id} \Rightarrow \text{cap} \Rightarrow (\text{State} \times \text{cap set})$   
**where**  $\text{get-takable-caps } s \text{ did take-cap} \equiv$   
 $\text{if}(\text{take-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 $\wedge \text{TAKE} \in \text{rights take-cap})$   
 then  
 let  
 $\text{did-dst} = \text{target take-cap}$   
 in  
 $(s, \text{get-domain-cap-set-from-domain-id } s \text{ did-dst})$   
 else  
 $(s, \{\})$

**definition**  $\text{take-endpoint-cap} :: \text{State} \Rightarrow \text{domain-id} \Rightarrow \text{cap} \Rightarrow \text{cap} \Rightarrow (\text{State} \times \text{bool})$   
**where**  $\text{take-endpoint-cap } s \text{ did take-cap dst-cap} \equiv$   
 $\text{if}(\text{take-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 $\wedge \text{TAKE} \in \text{rights take-cap}$   
 $\wedge \text{interferes did } s (\text{target dst-cap})$   
 $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s (\text{target take-cap}))$   
 then  
 let  
 $\text{caps0} = \text{caps } s;$   
 $\text{cs-dst} = \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 in  
 $(s[$   
 $\text{caps} := \text{caps0}(\text{did} := (\text{insert dst-cap cs-dst}))$   
 $], \text{True})$   
 else  
 $(s, \text{False})$

**definition**  $\text{remove-cap-right} :: \text{State} \Rightarrow \text{domain-id} \Rightarrow \text{cap} \Rightarrow \text{right} \Rightarrow (\text{State} \times \text{bool})$   
**where**  $\text{remove-cap-right } s \text{ did rm-cap right-to-rm} \equiv$



```

let
  caps0 = caps s;
  cs-dst = get-domain-cap-set-from-domain-id s did;
  cs-rest = {c. c ∈ cs-dst ∧ c ≠ rm-cap}
in
if (rm-cap ∈ get-domain-cap-set-from-domain-id s did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap
  ∧ REMOVE = right-to-rm
  ∧ {REMOVE} = rights rm-cap)
then
  (s (
    caps := caps0 (did := (cs-rest))
  ), True)
else if (
  rm-cap ∈ get-domain-cap-set-from-domain-id s did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap
)
then
  let
    new-cap = (target = target rm-cap,
               rights = (rights rm-cap) - {right-to-rm})
  in
  (s (
    caps := caps0 (did := (insert new-cap cs-rest))
  ), True)
else
  (s, False)

```

**definition** *system-init* :: *Sys-Config* ⇒ *State*  
**where** *system-init* *sc* ≡ (
 caps = (λ x. {}),
 e-msgs = (λ x. {}),
 e-buf-size = (λ x. None),
 domain-endpoint = e-listener(commconf *sc*)
)

## 0.7 Instantiation and Its Proofs of Security Model

**consts** *sysconf* :: *Sys-Config*

**definition** *sys-config-witness* :: *Sys-Config*

**where**

*sys-config-witness* ≡ (
 domconf = (
 d-name = (λ x. None),
 d-ep-set = (λ x. {})
 ),
 commconf = (
 e-max-buf-size = (λ x. None),
 e-name = (λ x. None),
 e-listener = (λ x. None)
 )
)

**consts** *s0t* :: *State*

**definition** *s0t-witness* :: *State*

**where**  $s0t\text{-}witness \equiv system\text{-}init\ sysconf$

**specification** ( $s0t$ )

$s0t\text{-}init: s0t = system\text{-}init\ sysconf$

**by**  $simp$

**definition**  $exec\text{-}event :: State \Rightarrow Event \Rightarrow State$

**where**  $exec\text{-}event\ s\ e \equiv$

case  $e$  of  $Client\text{-}Lookup\text{-}Endpoint\text{-}Name\ did\ ename \Rightarrow fst\ (client\text{-}lookup\text{-}endpoint\text{-}name\ sysconf\ s\ did\ ename)$   
 $| Send\text{-}Queuing\text{-}Message\ did\ eid\ m \Rightarrow fst\ (send\text{-}queuing\text{-}message\ s\ did\ eid\ m)$   
 $| Receive\text{-}Queuing\text{-}Message\ did\ eid \Rightarrow fst\ (receive\text{-}queuing\text{-}message\ s\ did\ eid)$   
 $| Get\text{-}My\text{-}Endpoints\text{-}Set\ did \Rightarrow fst\ (get\text{-}my\text{-}endpoints\text{-}set\ s\ did)$   
 $| Get\text{-}Caps\ did \Rightarrow fst\ (get\text{-}caps\ s\ did)$   
 $| Grant\text{-}Endpoint\text{-}Cap\ did\ grant\text{-}cap\ dst\text{-}cap \Rightarrow fst\ (grant\text{-}endpoint\text{-}cap\ s\ did\ grant\text{-}cap\ dst\text{-}cap)$   
 $| Remove\text{-}Cap\text{-}Right\ did\ dst\text{-}cap\ right\text{-}to\text{-}rm \Rightarrow fst\ (remove\text{-}cap\text{-}right\ s\ did\ dst\text{-}cap\ right\text{-}to\text{-}rm)$

**definition**  $domain\text{-}of\text{-}event :: Event \Rightarrow domain\text{-}id\ option$

**where**  $domain\text{-}of\text{-}event\ e \equiv$

case  $e$  of  $Client\text{-}Lookup\text{-}Endpoint\text{-}Name\ did\ ename \Rightarrow Some\ did$   
 $| Send\text{-}Queuing\text{-}Message\ did\ eid\ m \Rightarrow Some\ did$   
 $| Receive\text{-}Queuing\text{-}Message\ did\ eid \Rightarrow Some\ did$   
 $| Get\text{-}My\text{-}Endpoints\text{-}Set\ did \Rightarrow Some\ did$   
 $| Get\text{-}Caps\ did \Rightarrow Some\ did$   
 $| Grant\text{-}Endpoint\text{-}Cap\ did\ grant\text{-}cap\ dst\text{-}cap \Rightarrow Some\ did$   
 $| Remove\text{-}Cap\text{-}Right\ did\ dst\text{-}cap\ right\text{-}to\text{-}rm \Rightarrow Some\ did$

**definition**  $vpeq1 :: State \Rightarrow domain\text{-}id \Rightarrow State \Rightarrow bool\ ((- \sim - \sim -))$

**where**

$vpeq1\ s\ d\ t \equiv$

let

$cs1 = get\text{-}domain\text{-}cap\text{-}set\text{-}from\text{-}domain\text{-}id\ s\ d;$

$cs2 = get\text{-}domain\text{-}cap\text{-}set\text{-}from\text{-}domain\text{-}id\ t\ d;$

$dom\text{-}eps1 = get\text{-}endpoints\text{-}of\text{-}domain\ s\ d;$

$dom\text{-}eps2 = get\text{-}endpoints\text{-}of\text{-}domain\ t\ d$

in

if ( $cs1 = cs2$

$\wedge (\forall v. \text{interferes}\ v\ s\ d \longleftrightarrow \text{interferes}\ v\ t\ d)$

$\wedge dom\text{-}eps1 = dom\text{-}eps2$

$\wedge (\forall ep. ep \in dom\text{-}eps1$

$\longrightarrow get\text{-}msg\text{-}set\text{-}from\text{-}endpoint\text{-}id\ s\ ep = get\text{-}msg\text{-}set\text{-}from\text{-}endpoint\text{-}id\ t\ ep$ )

)

then

True

else

False

**declare**  $vpeq1\text{-}def\ [cong]$

**lemma**  $vpeq1\text{-}transitive\text{-}lemma : \forall\ s\ t\ r\ d. (vpeq1\ s\ d\ t) \wedge (vpeq1\ t\ d\ r) \longrightarrow (vpeq1\ s\ d\ r)$

**using**  $vpeq1\text{-}def$  **by**  $auto$

**lemma**  $vpeq1\text{-}symmetric\text{-}lemma : \forall\ s\ t\ d. (vpeq1\ s\ d\ t) \longrightarrow (vpeq1\ t\ d\ s)$

**using**  $vpeq1\text{-}def$  **by**  $auto$

```

lemma vpeq1-reflexive-lemma :  $\forall s d. (vpeq1 s d s)$ 
  using vpeq1-def by auto

lemma interf-reflexive-lemma :  $\forall d s. \text{interferes } d s d$ 
  using interferes-def by auto

lemma policy-respect-lemma :  $\forall v u s t. (s \sim u \sim t)$ 
   $\longrightarrow (\text{interferes } v s u = \text{interferes } v t u)$ 
  using vpeq1-def by auto

lemma reachable-top:  $\forall s a. (SM.\text{reachable0 } s0t \text{ exec-event}) s \longrightarrow (\exists s'. s' = \text{exec-event } s a)$ 
proof -
{
  fix s a
  assume p0:  $(SM.\text{reachable0 } s0t \text{ exec-event}) s$ 
  have  $(\exists s'. s' = \text{exec-event } s a)$ 
  proof (induct a)
    case (Client-Lookup-Endpoint-Name x) show ?case
    apply (induct x)
    by (simp add: exec-event-def) +
  next
  case (Send-Queuing-Message x1a x2 x3a) show ?case
  apply (induct x1a)
  by (simp add: exec-event-def) +
  next
  case (Receive-Queuing-Message x) show ?case
  apply (induct x)
  by (simp add: exec-event-def) +
  next
  case (Get-My-Endpoints-Set x) show ?case
  apply (induct x)
  by (simp add: exec-event-def) +
  next
  case (Get-Caps x) show ?case
  apply (induct x)
  by (simp add: exec-event-def) +
  case (Grant-Endpoint-Cap x1a x2 x3a) show ?case
  apply (induct x1a)
  by (simp add: exec-event-def) +
  case (Remove-Cap-Right x1a x2) show ?case
  apply (induct x1a)
  by (simp add: exec-event-def) +
  qed
}
then show ?thesis by auto
qed

declare Let-def [cong] and vpeq1-def [cong]

interpretation SM-enabled
  s0t exec-event domain-of-event vpeq1 interferes
  using vpeq1-transitive-lemma vpeq1-symmetric-lemma vpeq1-reflexive-lemma
  interf-reflexive-lemma policy-respect-lemma reachable-top
  SM.intro[of vpeq1 interferes]
  SM-enabled-axioms.intro[of s0t exec-event]
  SM-enabled.intro[of vpeq1 interferes s0t exec-event] by blast

```

## 0.8 Some lemmas of security proofs

## 0.9 Concrete unwinding condition of "local respect"

### 0.9.1 proving "client lookup endpoint name" satisfying the "local respect" property

```
lemma client-lookup-endpoint-name-lcl-resp:
  assumes p0: reachable0 s
    and p1:  $\neg(\text{interferes } \text{did } s \ d)$ 
    and p2:  $s' = \text{fst } (\text{client-lookup-endpoint-name sysconf } s \ \text{did } \text{ename})$ 
  shows  $s \sim d \sim s'$ 
  proof -
  {
    have a1:  $s = s'$ 
      by (simp add: p2 client-lookup-endpoint-name-def p1)
  }
  then show ?thesis by auto
qed
```

```
lemma client-lookup-endpoint-name-lcl-resp-e:
  assumes p0: reachable0 s
    and p1:  $a = (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})$ 
    and p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } a)) \ s \ d)$ 
    and p3:  $s' = \text{exec-event } s \ a$ 
  shows  $s \sim d \sim s'$ 
  proof -
  {
    have a0:  $(\text{the } (\text{domain-of-event } a)) = \text{did}$ 
      using p1 domain-of-event-def by auto
    have a1:  $s' = \text{fst } (\text{client-lookup-endpoint-name sysconf } s \ \text{did } \text{ename})$ 
      using p1 p3 exec-event-def by auto
    have a2:  $\neg(\text{interferes } \text{did } s \ d)$ 
      using p2 a0 by auto
    have a3:  $s \sim d \sim s'$ 
      using a1 a2 p0 client-lookup-endpoint-name-lcl-resp by blast
  }
  then show ?thesis by auto
qed
```

```
lemma client-lookup-endpoint-name-lcrsp-e: dynamic-local-respect-e (Client-Lookup-Endpoint-Name did ename)
  proof -
  {
    have  $\forall d \ s. \text{reachable0 } s$ 
       $\wedge \neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))) \ s \ d)$ 
       $\longrightarrow (s \sim d \sim (\text{exec-event } s \ (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))$ 
    proof -
    {
      fix d s
      assume p1: reachable0 s
      assume p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))) \ s \ d)$ 
      have  $(s \sim d \sim (\text{exec-event } s \ (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))$ 
        using p1 p2 client-lookup-endpoint-name-lcl-resp-e by blast
    }
    then show ?thesis by blast
  }
  qed
  then show ?thesis
    using dynamic-local-respect-e-def by blast
  qed
```

### 0.9.2 proving "send queuing message" satisfying the "local respect" property

```

lemma send-queuing-message-notchg-domain-cap-set:
  assumes p0: reachable0 s
    and p1:  $\neg(\text{interferes } \text{did } s \ d)$ 
    and p2:  $s' = \text{fst } (\text{send-queuing-message } s \ \text{did } \text{eid } m)$ 
  shows  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
    =  $\text{get-domain-cap-set-from-domain-id } s' \ d$ 
  proof (cases (get-domain-id-from-endpoint s eid  $\neq$  None
     $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid))))
    assume b0: (get-domain-id-from-endpoint s eid  $\neq$  None
       $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))
    have b1:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
      =  $\text{get-domain-cap-set-from-domain-id } s' \ d$ 
      using b0 p2 send-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
    then show ?thesis by auto
  next
    assume b0:  $\neg(\text{get-domain-id-from-endpoint } s \ \text{eid} \neq \text{None}$ 
       $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))
    have b1:  $s = s'$ 
      using b0 p2 send-queuing-message-def by auto
    have b2:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
      =  $\text{get-domain-cap-set-from-domain-id } s' \ d$ 
      using b1 get-domain-cap-set-from-domain-id-def by auto
    then show ?thesis by auto
  qed

```

```

lemma send-queuing-message-notchg-policy:
  assumes p0: reachable0 s
    and p2:  $s' = \text{fst } (\text{send-queuing-message } s \ \text{did } \text{eid } m)$ 
  shows  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
  proof (cases (get-domain-id-from-endpoint s eid  $\neq$  None
     $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid))))
    assume b0: (get-domain-id-from-endpoint s eid  $\neq$  None
       $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))
    have b1:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
      =  $\text{get-domain-cap-set-from-domain-id } s' \ d$ 
      using b0 p2 send-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
    have b2:  $\forall v. \text{get-domain-cap-set-from-domain-id } s \ v$ 
      =  $\text{get-domain-cap-set-from-domain-id } s' \ v$ 
      using b0 p2 send-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
    have b3:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
      using b1 b2 interferes-def by auto
    then show ?thesis by auto
  next
    assume b0:  $\neg(\text{get-domain-id-from-endpoint } s \ \text{eid} \neq \text{None}$ 
       $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))
    have b1:  $s = s'$ 
      using b0 p2 send-queuing-message-def by auto
    have b2:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
      using b1 interferes-def by auto
    then show ?thesis by auto
  qed

```

```

lemma send-queuing-message-notchg-dom-eps:
  assumes p0: reachable0 s
    and p2:  $s' = \text{fst } (\text{send-queuing-message } s \ \text{did } \text{eid } m)$ 
  shows  $\text{get-endpoints-of-domain } s \ d = \text{get-endpoints-of-domain } s' \ d$ 

```

**proof** (cases (get-domain-id-from-endpoint s eid  $\neq$  None  
 $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid))))  
**assume** b0: (get-domain-id-from-endpoint s eid  $\neq$  None  
 $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))  
**have** b1: get-domain-cap-set-from-domain-id s d  
 $=$  get-domain-cap-set-from-domain-id s' d  
**using** b0 p2 send-queuing-message-def get-domain-cap-set-from-domain-id-def **by** auto  
**have** b2: domain-endpoint s = domain-endpoint s'  
**using** b0 p2 send-queuing-message-def **by** auto  
**have** b3: get-endpoints-of-domain s d  
 $=$  get-endpoints-of-domain s' d  
**using** b2 get-endpoints-of-domain-def **by** auto  
**then show** ?thesis **by** auto  
**next**  
**assume** b0:  $\neg$  (get-domain-id-from-endpoint s eid  $\neq$  None  
 $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))  
**have** b1: s = s'  
**using** b0 p2 send-queuing-message-def **by** auto  
**have** b2: get-endpoints-of-domain s d  
 $=$  get-endpoints-of-domain s' d  
**using** b1 get-endpoints-of-domain-def **by** auto  
**then show** ?thesis **by** auto  
**qed**

**lemma** send-queuing-message-notchg-ep-msgs:

**assumes** p0: reachable0 s  
**and** p1:  $\neg$ (interferes did s d)  
**and** p2: s' = fst (send-queuing-message s did eid m)  
**shows** ( $\forall$  ep. ep  $\in$  get-endpoints-of-domain s d  
 $\longrightarrow$  get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )  
**proof** (cases (get-domain-id-from-endpoint s eid  $\neq$  None  
 $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid))))  
**assume** b0: (get-domain-id-from-endpoint s eid  $\neq$  None  
 $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))  
**have** b1: get-domain-cap-set-from-domain-id s d  
 $=$  get-domain-cap-set-from-domain-id s' d  
**using** b0 p2 send-queuing-message-def get-domain-cap-set-from-domain-id-def **by** auto  
**have** b2: domain-endpoint s = domain-endpoint s'  
**using** b0 p2 send-queuing-message-def **by** auto  
**have** b3:  $\forall e. e \neq \text{eid}$   
 $\longrightarrow ((e\text{-msgs } s) e) = ((e\text{-msgs } s') e)$   
**using** b0 p2 send-queuing-message-def **by** auto  
**have** b4: domain-endpoint s = domain-endpoint s'  
**using** b0 p2 send-queuing-message-def **by** auto  
**have** b5: get-endpoints-of-domain s d  
 $=$  get-endpoints-of-domain s' d  
**using** b2 get-endpoints-of-domain-def **by** auto  
**have** b6: (the (get-domain-id-from-endpoint s eid))  $\neq$  d  
**using** b0 p1 **by** auto  
**have** b7: (the ((domain-endpoint s) eid))  $\neq$  d  
**using** b6 get-domain-id-from-endpoint-def **by** auto  
**have** b8: eid  $\notin$  get-endpoints-of-domain s d  
**using** b7 b6 get-endpoints-of-domain-def **by** auto  
**have** b9:  $\forall$  ep. ep  $\in$  get-endpoints-of-domain s d  
 $\longrightarrow ((e\text{-msgs } s) ep) = ((e\text{-msgs } s') ep)$   
**using** b8 b3 **by** auto  
**have** b10:  $\forall$  ep. ep  $\in$  get-endpoints-of-domain s d  
 $\longrightarrow$  get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep

```

    using b9 get-msg-set-from-endpoint-id-def by auto
  then show ?thesis by auto
next
  assume b0:  $\neg$  (get-domain-id-from-endpoint s eid  $\neq$  None
     $\wedge$  interferes did s (the (get-domain-id-from-endpoint s eid)))
  have b1: s = s'
    using b0 p2 send-queuing-message-def by auto
  have b2: ( $\forall$  ep. ep  $\in$  get-endpoints-of-domain s d
     $\longrightarrow$  get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
    using b1 by auto
  then show ?thesis by auto
qed

```

lemma send-queuing-message-lcl-resp:

```

assumes p0: reachable0 s
  and p1:  $\neg$ (interferes did s d)
  and p2: s' = fst (send-queuing-message s did eid m)
shows s  $\sim$  d  $\sim$  s'
proof -
{
  have a0: did  $\neq$  d
    using p1 interferes-def by auto
  have a1: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using p0 p1 p2 send-queuing-message-notchg-domain-cap-set by auto
  have a2: ( $\forall$  v. interferes v s d  $\longleftrightarrow$  interferes v s' d)
    using p0 p1 p2 send-queuing-message-notchg-policy by auto
  have a3: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
    using p0 p1 p2 send-queuing-message-notchg-dom-eps by auto
  have a4: ( $\forall$  ep. ep  $\in$  get-endpoints-of-domain s d
     $\longrightarrow$  get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
    using p0 p1 p2 send-queuing-message-notchg-ep-msgs by auto
  have a5: s  $\sim$  d  $\sim$  s'
    using a1 a2 a3 a4 by auto
}
then show ?thesis by auto
qed

```

lemma send-queuing-message-lcl-resp-e:

```

assumes p0: reachable0 s
  and p1: a = Send-Queuing-Message did eid m
  and p2:  $\neg$ (interferes (the (domain-of-event a)) s d)
  and p3: s' = exec-event s a
shows s  $\sim$  d  $\sim$  s'
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p1 domain-of-event-def by auto
  have a1: s' = fst (send-queuing-message s did eid m)
    using p1 p3 exec-event-def by auto
  have a2:  $\neg$ (interferes did s d)
    using p2 a0 by auto
  have a3: s  $\sim$  d  $\sim$  s'
    using a1 a2 p0 send-queuing-message-lcl-resp by blast
}
then show ?thesis by auto
qed

```

```

lemma send-queuing-message-lcl-lcrsp-e: dynamic-local-respect-e (Send-Queuing-Message did eid m)
proof -
{
  have  $\forall d s. \text{reachable0 } s$ 
     $\wedge \neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Send-Queuing-Message did eid m}))) s d)$ 
     $\longrightarrow (s \sim d \sim (\text{exec-event } s (\text{Send-Queuing-Message did eid m})))$ 
  proof -
  {
    fix  $d s$ 
    assume  $p1: \text{reachable0 } s$ 
    assume  $p2: \neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Send-Queuing-Message did eid m}))) s d)$ 
    have  $(s \sim d \sim (\text{exec-event } s (\text{Send-Queuing-Message did eid m})))$ 
    using  $p1 p2 \text{ send-queuing-message-lcl-resp-e}$  by blast
  }
  then show ?thesis by blast
qed
}
then show ?thesis
using dynamic-local-respect-e-def by blast
qed

```

### 0.9.3 proving "receive queuing message" satisfying the "local respect" property

```

lemma receive-queuing-message-notchg-domain-cap-set:
assumes  $p0: \text{reachable0 } s$ 
and  $p1: \neg(\text{interferes } \text{did } s d)$ 
and  $p2: s' = \text{fst } (\text{receive-queuing-message } s \text{ did eid})$ 
shows  $\text{get-domain-cap-set-from-domain-id } s d$ 
   $= \text{get-domain-cap-set-from-domain-id } s' d$ 
proof (cases the(get-domain-id-from-endpoint s eid) = did)
assume  $a0: \text{the}(\text{get-domain-id-from-endpoint } s \text{ eid}) = \text{did}$ 
have  $a1: \text{get-domain-cap-set-from-domain-id } s d$ 
   $= \text{get-domain-cap-set-from-domain-id } s' d$ 
using  $a0 p2 \text{ receive-queuing-message-def get-domain-cap-set-from-domain-id-def}$  by auto
then show ?thesis by auto
next
assume  $a0: \text{the}(\text{get-domain-id-from-endpoint } s \text{ eid}) \neq \text{did}$ 
have  $a1: \text{get-domain-cap-set-from-domain-id } s d$ 
   $= \text{get-domain-cap-set-from-domain-id } s' d$ 
using  $a0 p2 \text{ receive-queuing-message-def get-domain-cap-set-from-domain-id-def}$  by auto
then show ?thesis by auto
qed

```

```

lemma receive-queuing-message-notchg-policy:
assumes  $p0: \text{reachable0 } s$ 
and  $p1: \neg(\text{interferes } \text{did } s d)$ 
and  $p2: s' = \text{fst } (\text{receive-queuing-message } s \text{ did eid})$ 
shows  $(\forall v. \text{interferes } v s d \longleftrightarrow \text{interferes } v s' d)$ 
proof (cases the(get-domain-id-from-endpoint s eid) = did)
assume  $a0: \text{the}(\text{get-domain-id-from-endpoint } s \text{ eid}) = \text{did}$ 
have  $a1: \text{get-domain-cap-set-from-domain-id } s d$ 
   $= \text{get-domain-cap-set-from-domain-id } s' d$ 
using  $a0 p2 \text{ receive-queuing-message-def get-domain-cap-set-from-domain-id-def}$  by auto
have  $a2: \forall v. \text{get-domain-cap-set-from-domain-id } s v$ 
   $= \text{get-domain-cap-set-from-domain-id } s' v$ 
using  $a0 p2 \text{ receive-queuing-message-def get-domain-cap-set-from-domain-id-def}$  by auto
have  $a3: (\forall v. \text{interferes } v s d \longleftrightarrow \text{interferes } v s' d)$ 

```



```

    using a1 a2 interferes-def by auto
  then show ?thesis by auto
next
  assume a0: the(get-domain-id-from-endpoint s eid) ≠ did
  have a1: s = s'
    using a0 p2 receive-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
  have a2: (∀ v. interferes v s d ⟷ interferes v s' d)
    using a1 interferes-def by auto
  then show ?thesis by auto
qed

```

**lemma** *receive-queuing-message-notchg-dom-eps:*

```

  assumes p0: reachable0 s
    and p2: s' = fst (receive-queuing-message s did eid)
  shows get-endpoints-of-domain s d = get-endpoints-of-domain s' d
proof (cases the(get-domain-id-from-endpoint s eid) = did)
  assume a0: the(get-domain-id-from-endpoint s eid) = did
  have a1: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using a0 p2 receive-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
  have a2: domain-endpoint s = domain-endpoint s'
    using a0 p2 receive-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
  have a3: get-endpoints-of-domain s d
    = get-endpoints-of-domain s' d
    using a1 a2 get-endpoints-of-domain-def by auto
  then show ?thesis by auto
next
  assume a0: the(get-domain-id-from-endpoint s eid) ≠ did
  have a1: s = s'
    using a0 p2 receive-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
  have a2: get-endpoints-of-domain s d
    = get-endpoints-of-domain s' d
    using a1 get-endpoints-of-domain-def by auto
  then show ?thesis by auto
qed

```

**lemma** *receive-queuing-message-notchg-ep-msgs:*

```

  assumes p0: reachable0 s
    and p1: ¬(interferes did s d)
    and p2: s' = fst (receive-queuing-message s did eid)
  shows (∀ ep. ep ∈ get-endpoints-of-domain s d
    ⟶ get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
proof (cases the(get-domain-id-from-endpoint s eid) = did)
  assume a0: the(get-domain-id-from-endpoint s eid) = did
  have a1: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using a0 p2 receive-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
  have a2: domain-endpoint s = domain-endpoint s'
    using a0 p2 receive-queuing-message-def by auto
  have a3: get-endpoints-of-domain s d
    = get-endpoints-of-domain s' d
    using a1 a2 get-endpoints-of-domain-def by auto
  have a4: ∀ e. e ≠ eid
    ⟶ ((e-msgs s) e) = ((e-msgs s') e)
    using a0 p2 receive-queuing-message-def by auto
  have a5: d ≠ did
    using p1 interferes-def by auto
  have a6: the(get-domain-id-from-endpoint s eid) ≠ d

```

```

    using a5 a0 by auto
  have a7: (the ((domain-endpoint s) eid))  $\neq$  d
    using a6 get-domain-id-from-endpoint-def by auto
  have a8: eid  $\notin$  get-endpoints-of-domain s d
    using a7 a6 get-endpoints-of-domain-def by auto
  have a9:  $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow ((e\text{-msgs } s) \ ep) = ((e\text{-msgs } s') \ ep)$ 
    using a8 a4 by auto
  have a10:  $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } s' \ ep$ 
    using a9 get-msg-set-from-endpoint-id-def by auto
  then show ?thesis by auto
next
assume a0: the(get-domain-id-from-endpoint s eid)  $\neq$  did
have a1: s = s'
  using a0 p2 receive-queuing-message-def get-domain-cap-set-from-domain-id-def by auto
have a2: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
   $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } s' \ ep$ )
  using a1 get-msg-set-from-endpoint-id-def by auto
then show ?thesis by auto
qed

```

lemma receive-queuing-message-lcl-resp:

```

assumes p0: reachable0 s
  and p1:  $\neg(\text{interferes } did \ s \ d)$ 
  and p2: s' = fst (receive-queuing-message s did eid)
shows s  $\sim$  d  $\sim$  s'
proof -
{
  have a0: did  $\neq$  d
    using p1 interferes-def by auto
  have a1: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using p0 p1 p2 receive-queuing-message-notchg-domain-cap-set by auto
  have a2: ( $\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d$ )
    using p0 p1 p2 receive-queuing-message-notchg-policy by auto
  have a3: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
    using p0 p1 p2 receive-queuing-message-notchg-dom-eps by auto
  have a4: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } s' \ ep$ )
    using p0 p1 p2 receive-queuing-message-notchg-ep-msgs by auto
  have a5: s  $\sim$  d  $\sim$  s'
    using a1 a2 a3 a4 by auto
}
then show ?thesis by auto
qed

```

lemma receive-queuing-message-lcl-resp-e:

```

assumes p0: reachable0 s
  and p1: a = (Receive-Queuing-Message did eid)
  and p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } a)) \ s \ d)$ 
  and p3: s' = exec-event s a
shows s  $\sim$  d  $\sim$  s'
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p1 domain-of-event-def by auto
  have a1: s' = fst (receive-queuing-message s did eid)

```

```

    using p1 p3 exec-event-def by auto
  have a2:  $\neg(\text{interferes } \text{did } s \ d)$ 
    using p2 a0 by auto
  have a3:  $s \sim d \sim s'$ 
    using a1 a2 p0 receive-queuing-message-lcl-resp by blast
}
then show ?thesis by auto
qed

lemma receive-queuing-message-lcrsp-e: dynamic-local-respect-e (Receive-Queuing-Message did eid)
proof -
{
  have  $\forall d \ s. \text{reachable0 } s$ 
     $\wedge \neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Receive-Queuing-Message } \text{did } \text{eid}))) \ s \ d)$ 
     $\longrightarrow (s \sim d \sim (\text{exec-event } s \ (\text{Receive-Queuing-Message } \text{did } \text{eid})))$ 
  proof -
  {
    fix d s
    assume p1:  $\text{reachable0 } s$ 
    assume p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Receive-Queuing-Message } \text{did } \text{eid}))) \ s \ d)$ 
    have  $(s \sim d \sim (\text{exec-event } s \ (\text{Receive-Queuing-Message } \text{did } \text{eid})))$ 
      using p1 p2 receive-queuing-message-lcl-resp-e by blast
    }
    then show ?thesis by blast
  qed
}
then show ?thesis
  using dynamic-local-respect-e-def by blast
qed

```

#### 0.9.4 proving "get my endpoints set" satisfying the "local respect" property

```

lemma get-my-endpoints-set-lcl-resp:
assumes p0:  $\text{reachable0 } s$ 
  and p1:  $\neg(\text{interferes } \text{did } s \ d)$ 
  and p2:  $s' = \text{fst } (\text{get-my-endpoints-set } s \ \text{did})$ 
shows  $s \sim d \sim s'$ 
proof -
{
  have a1:  $s = s'$ 
    by (simp add: p2 get-my-endpoints-set-def p1)
}
then show ?thesis by auto
qed

lemma get-my-endpoints-set-lcl-resp-e:
assumes p0:  $\text{reachable0 } s$ 
  and p1:  $a = (\text{Get-My-Endpoints-Set } \text{did})$ 
  and p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } a)) \ s \ d)$ 
  and p3:  $s' = \text{exec-event } s \ a$ 
shows  $s \sim d \sim s'$ 
proof -
{
  have a0:  $(\text{the } (\text{domain-of-event } a)) = \text{did}$ 
    using p1 domain-of-event-def by auto
  have a1:  $s' = \text{fst } (\text{get-my-endpoints-set } s \ \text{did})$ 
    using p1 p3 exec-event-def by auto
  have a2:  $\neg(\text{interferes } \text{did } s \ d)$ 

```

```

    using p2 a0 by auto
  have a3:  $s \sim d \sim s'$ 
    using a1 a2 p0 get-my-endpoints-set-lcl-resp by blast
}
then show ?thesis by auto
qed

lemma get-my-endpoints-set-lcrsp-e: dynamic-local-respect-e (Get-My-Endpoints-Set did)
proof -
  {
    have  $\forall d s. \text{reachable0 } s$ 
       $\wedge \neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Get-My-Endpoints-Set } \text{did}))) s d)$ 
       $\longrightarrow (s \sim d \sim (\text{exec-event } s (\text{Get-My-Endpoints-Set } \text{did})))$ 
    proof -
      {
        fix d s
        assume p1:  $\text{reachable0 } s$ 
        assume p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Get-My-Endpoints-Set } \text{did}))) s d)$ 
        have  $(s \sim d \sim (\text{exec-event } s (\text{Get-My-Endpoints-Set } \text{did})))$ 
          using p1 p2 get-my-endpoints-set-lcl-resp-e by blast
      }
      then show ?thesis by blast
    qed
  }
  then show ?thesis
    using dynamic-local-respect-e-def by blast
qed

```

### 0.9.5 proving "get caps" satisfying the "local respect" property

```

lemma get-caps-lcl-resp:
assumes p0:  $\text{reachable0 } s$ 
  and p1:  $\neg(\text{interferes } \text{did } s d)$ 
  and p2:  $s' = \text{fst } (\text{get-caps } s \text{ did})$ 
shows  $s \sim d \sim s'$ 
proof -
  {
    have a1:  $s = s'$ 
      by (simp add: p2 get-caps-def p1)
  }
  then show ?thesis by auto
qed

lemma get-caps-lcl-resp-e:
assumes p0:  $\text{reachable0 } s$ 
  and p1:  $a = (\text{Get-Caps } \text{did})$ 
  and p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } a)) s d)$ 
  and p3:  $s' = \text{exec-event } s a$ 
shows  $s \sim d \sim s'$ 
proof -
  {
    have a0:  $(\text{the } (\text{domain-of-event } a)) = \text{did}$ 
      using p1 domain-of-event-def by auto
    have a1:  $s' = \text{fst } (\text{get-caps } s \text{ did})$ 
      using p1 p3 exec-event-def by auto
    have a2:  $\neg(\text{interferes } \text{did } s d)$ 
      using p2 a0 by auto
    have a3:  $s \sim d \sim s'$ 

```

```

    using a1 a2 p0 get-caps-lcl-resp by blast
  }
then show ?thesis by auto
qed

lemma get-caps-lcrsp-e: dynamic-local-respect-e (Get-Caps did)
proof -
  {
    have  $\forall d s. \text{reachable0 } s$ 
       $\wedge \neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Get-Caps did}))) s d)$ 
       $\longrightarrow (s \sim d \sim (\text{exec-event } s (\text{Get-Caps did})))$ 
    proof -
      {
        fix d s
        assume p1:  $\text{reachable0 } s$ 
        assume p2:  $\neg(\text{interferes } (\text{the } (\text{domain-of-event } (\text{Get-Caps did}))) s d)$ 
        have  $(s \sim d \sim (\text{exec-event } s (\text{Get-Caps did})))$ 
          using p1 p2 get-caps-lcl-resp-e by blast
      }
      then show ?thesis by blast
    qed
  }
then show ?thesis
  using dynamic-local-respect-e-def by blast
qed

```

### 0.9.6 proving "grant endpoint cap" satisfying the "local respect" property

```

lemma grant-endpoint-cap-notchg-domain-cap-set:
assumes p0:  $\text{reachable0 } s$ 
  and p1:  $\neg(\text{interferes } \text{did } s d)$ 
  and p2:  $s' = \text{fst } (\text{grant-endpoint-cap } s \text{ did } \text{grant-cap } \text{dst-cap})$ 
shows  $\text{get-domain-cap-set-from-domain-id } s d$ 
  =  $\text{get-domain-cap-set-from-domain-id } s' d$ 
proof (cases  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ )
   $\wedge \text{GRANT} \in \text{rights } \text{grant-cap}$ 
   $\wedge \text{target } \text{grant-cap} \neq \text{target } \text{dst-cap}$ 
   $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
  assume a0:  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
     $\wedge \text{GRANT} \in \text{rights } \text{grant-cap}$ 
     $\wedge \text{target } \text{grant-cap} \neq \text{target } \text{dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
  let ?did-dst =  $\text{target } \text{grant-cap}$ 
  have a1:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow (\text{caps } s) v = (\text{caps } s') v$ 
    using a0 p2 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have a2:  $d \neq \text{did}$ 
    using p1 interferes-def by auto
  have a3:  $\neg((\text{did} = d$ 
     $\vee (\exists c. c \in (\text{get-domain-cap-set-from-domain-id } s \text{ did}) \wedge \text{target } c = d)))$ 
    using interferes-def p1 by force
  have a4:  $\forall c. c \in (\text{get-domain-cap-set-from-domain-id } s \text{ did})$ 
     $\longrightarrow \text{target } c \neq d$ 
    using a3 by auto
  have a5:  $?\text{did-dst} \neq d$ 
    using a4 a0 by auto
  have a6:  $(\text{caps } s) d = (\text{caps } s') d$ 
    using a1 a5 by auto

```

```

have a7: get-domain-cap-set-from-domain-id s d
  = get-domain-cap-set-from-domain-id s' d
  using a6 get-domain-cap-set-from-domain-id-def by auto
then show ?thesis by auto
next
assume a0:  $\neg$  (grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
   $\wedge$  GRANT  $\in$  rights grant-cap
   $\wedge$  target grant-cap  $\neq$  target dst-cap
   $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did)
have a1: s = s'
  using a0 p2 grant-endpoint-cap-def by auto
have a2: get-domain-cap-set-from-domain-id s d
  = get-domain-cap-set-from-domain-id s' d
  using a1 get-domain-cap-set-from-domain-id-def by auto
then show ?thesis by auto
qed

```

lemma grant-endpoint-cap-notchg-policy:

```

assumes p0: reachable0 s
  and p1:  $\neg$ (interferes did s d)
  and p2: s' = fst (grant-endpoint-cap s did grant-cap dst-cap)
shows ( $\forall v. \text{interferes } v \text{ s d} \longleftrightarrow \text{interferes } v \text{ s' d}$ )
proof (cases grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
   $\wedge$  GRANT  $\in$  rights grant-cap
   $\wedge$  target grant-cap  $\neq$  target dst-cap
   $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did)
  assume a0: grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
     $\wedge$  GRANT  $\in$  rights grant-cap
     $\wedge$  target grant-cap  $\neq$  target dst-cap
     $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did
  let ?did-dst = target grant-cap
  let ?did-granted = target dst-cap
  have a1:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow$  (caps s) v = (caps s') v
    using a0 p2 grant-endpoint-cap-def by auto
  have a2: d  $\neq$  did
    using p1 interferes-def by auto
  have a3:  $\neg$ (( did = d
     $\vee$  ( $\exists c. c \in$  (get-domain-cap-set-from-domain-id s did)  $\wedge$  target c = d)))
    using interferes-def p1 by force
  have a4:  $\forall c. c \in$  (get-domain-cap-set-from-domain-id s did)
     $\longrightarrow$  target c  $\neq$  d
    using a3 by auto
  have a5: ?did-dst  $\neq$  d
    using a4 a0 by auto
  have a6: (caps s) d = (caps s') d
    using a1 a5 by auto
  have a7: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using a6 get-domain-cap-set-from-domain-id-def by auto
  have a8: ?did-granted  $\neq$  d
    using a4 a0 by auto
  have a9:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow$  get-domain-cap-set-from-domain-id s v
    = get-domain-cap-set-from-domain-id s' v
    using a1 get-domain-cap-set-from-domain-id-def by auto
  have a10:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow$  interferes v s d  $\longleftrightarrow$  interferes v s' d

```

```

    using a9 a7 interferes-def by auto
have a11: get-domain-cap-set-from-domain-id s' ?did-dst
    = {dst-cap}  $\cup$  get-domain-cap-set-from-domain-id s ?did-dst
    using a0 p2 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
have a12: interferes ?did-dst s d = interferes ?did-dst s' d
    using a11 a8 interferes-def by auto
have a13:  $\forall v. \text{interferes } v \text{ s d} \longleftrightarrow \text{interferes } v \text{ s' d}$ 
    using a10 a12 by force
then show ?thesis by auto
next
assume a0:  $\neg (\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id s did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id s did})$ 
have a1: s = s'
    using a0 p2 grant-endpoint-cap-def by auto
have a2: ( $\forall v. \text{interferes } v \text{ s d} \longleftrightarrow \text{interferes } v \text{ s' d}$ )
    using a1 interferes-def by auto
then show ?thesis by auto
qed

```

```

lemma grant-endpoint-cap-notchg-dom-eps:
assumes p0: reachable0 s
    and p2: s' = fst (grant-endpoint-cap s did grant-cap dst-cap)
shows get-endpoints-of-domain s d = get-endpoints-of-domain s' d
proof (cases grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id s did}$ )
assume a0: grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id s did}$ 
have a2: domain-endpoint s = domain-endpoint s'
    using a0 p2 grant-endpoint-cap-def by auto
have a3: get-endpoints-of-domain s d
    = get-endpoints-of-domain s' d
    using a2 get-endpoints-of-domain-def by auto
then show ?thesis by auto
next
assume a0:  $\neg (\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id s did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id s did})$ 
have a1: s = s'
    using a0 p2 grant-endpoint-cap-def by auto
have a2: get-endpoints-of-domain s d
    = get-endpoints-of-domain s' d
    using a1 get-endpoints-of-domain-def by auto
then show ?thesis by auto
qed

```

```

lemma grant-endpoint-cap-notchg-ep-msgs:
assumes p0: reachable0 s
    and p1:  $\neg (\text{interferes did s d})$ 
    and p2: s' = fst (grant-endpoint-cap s did grant-cap dst-cap)
shows ( $\forall ep. ep \in \text{get-endpoints-of-domain s d}$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id s ep} = \text{get-msg-set-from-endpoint-id s' ep}$ )

```

**proof** (cases  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 $\wedge \text{GRANT} \in \text{rights } \text{grant-cap}$   
 $\wedge \text{target grant-cap} \neq \text{target dst-cap}$   
 $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ )  
**assume**  $a0$ :  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 $\wedge \text{GRANT} \in \text{rights } \text{grant-cap}$   
 $\wedge \text{target grant-cap} \neq \text{target dst-cap}$   
 $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
**have**  $a2$ :  $\text{domain-endpoint } s = \text{domain-endpoint } s'$   
**using**  $a0$   $p2$   $\text{grant-endpoint-cap-def}$  **by** *auto*  
**have**  $a3$ :  $\text{get-endpoints-of-domain } s \text{ d}$   
 $= \text{get-endpoints-of-domain } s' \text{ d}$   
**using**  $a2$   $\text{get-endpoints-of-domain-def}$  **by** *auto*  
**have**  $a4$ :  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d}$   
 $\longrightarrow \text{get-msg-set-from-endpoint-id } s \text{ ep} = \text{get-msg-set-from-endpoint-id } s' \text{ ep} )$   
**using**  $a0$   $p2$   $\text{grant-endpoint-cap-def}$   $\text{get-msg-set-from-endpoint-id-def}$  **by** *auto*  
**then show** *?thesis* **by** *auto*  
**next**  
**assume**  $a0$ :  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$   
 $\wedge \text{GRANT} \in \text{rights } \text{grant-cap}$   
 $\wedge \text{target grant-cap} \neq \text{target dst-cap}$   
 $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did})$   
**have**  $a1$ :  $s = s'$   
**using**  $a0$   $p2$   $\text{grant-endpoint-cap-def}$  **by** *auto*  
**have**  $a2$ :  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d}$   
 $\longrightarrow \text{get-msg-set-from-endpoint-id } s \text{ ep} = \text{get-msg-set-from-endpoint-id } s' \text{ ep} )$   
**using**  $a1$  **by** *auto*  
**then show** *?thesis* **by** *auto*  
**qed**

**lemma** *grant-endpoint-cap-lcl-resp*:

**assumes**  $p0$ : *reachable0*  $s$

**and**  $p1$ :  $\neg(\text{interferes } \text{did } s \text{ d})$

**and**  $p2$ :  $s' = \text{fst } (\text{grant-endpoint-cap } s \text{ did } \text{eid } m)$

**shows**  $s \sim d \sim s'$

**proof** –

**{**  
**have**  $a0$ :  $\text{did} \neq d$   
**using**  $p1$  *interferes-def* **by** *auto*  
**have**  $a1$ :  $\text{get-domain-cap-set-from-domain-id } s \text{ d}$   
 $= \text{get-domain-cap-set-from-domain-id } s' \text{ d}$   
**using**  $p0$   $p1$   $p2$  *grant-endpoint-cap-notchg-domain-cap-set* **by** *auto*  
**have**  $a2$ :  $(\forall v. \text{interferes } v \text{ s d} \longleftrightarrow \text{interferes } v \text{ s' d})$   
**using**  $p0$   $p1$   $p2$  *grant-endpoint-cap-notchg-policy* **by** *auto*  
**have**  $a3$ :  $\text{get-endpoints-of-domain } s \text{ d} = \text{get-endpoints-of-domain } s' \text{ d}$   
**using**  $p0$   $p1$   $p2$  *grant-endpoint-cap-notchg-dom-eps* **by** *auto*  
**have**  $a4$ :  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d}$   
 $\longrightarrow \text{get-msg-set-from-endpoint-id } s \text{ ep} = \text{get-msg-set-from-endpoint-id } s' \text{ ep} )$   
**using**  $p0$   $p1$   $p2$  *grant-endpoint-cap-notchg-ep-msgs* **by** *auto*  
**have**  $a5$ :  $s \sim d \sim s'$   
**using**  $a1$   $a2$   $a3$   $a4$  **by** *auto*  
**}**  
**then show** *?thesis* **by** *auto*  
**qed**

**lemma** *grant-endpoint-cap-lcl-resp-e*:

**assumes**  $p0$ : *reachable0*  $s$

**and**  $p1$ :  $a = (\text{Grant-Endpoint-Cap } \text{did } \text{grant-cap } \text{dst-cap})$



```

and p2: ¬(interferes (the (domain-of-event a)) s d)
and p3: s' = exec-event s a
shows s ~ d ~ s'
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p1 domain-of-event-def by auto
  have a1: s' = fst (grant-endpoint-cap s did grant-cap dst-cap)
    using p1 p3 exec-event-def by auto
  have a2: ¬(interferes did s d)
    using p2 a0 by auto
  have a3: s ~ d ~ s'
    using a1 a2 p0 grant-endpoint-cap-lcl-resp by blast
}
then show ?thesis by auto
qed

lemma grant-endpoint-cap-lcrsp-e: dynamic-local-respect-e (Grant-Endpoint-Cap did grant-cap dst-cap)
proof -
{
  have ∀ d s. reachable0 s
    ∧ ¬(interferes (the (domain-of-event (Grant-Endpoint-Cap did grant-cap dst-cap))) s d)
    → (s ~ d ~ (exec-event s (Grant-Endpoint-Cap did grant-cap dst-cap)))
  proof -
  {
    fix d s
    assume p1: reachable0 s
    assume p2: ¬(interferes (the (domain-of-event (Grant-Endpoint-Cap did grant-cap dst-cap))) s d)
    have (s ~ d ~ (exec-event s (Grant-Endpoint-Cap did grant-cap dst-cap)))
      using p1 p2 grant-endpoint-cap-lcl-resp-e by blast
    }
    then show ?thesis by blast
  }
  qed
}
then show ?thesis
  using dynamic-local-respect-e-def by blast
qed

```

### 0.9.7 proving "remove cap right" satisfying the "local respect" property

```

lemma remove-cap-right-notchg-domain-cap-set:
assumes p0: reachable0 s
and p1: ¬(interferes did s d)
and p2: s' = fst (remove-cap-right s did rm-cap right-to-rm)
shows get-domain-cap-set-from-domain-id s d
      = get-domain-cap-set-from-domain-id s' d
proof -
{
  have a0: d ≠ did
    using p1 interferes-def by auto
  have get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
  proof (cases rm-cap ∈ get-domain-cap-set-from-domain-id s did
    ∧ REMOVE ∈ rights rm-cap
    ∧ right-to-rm ∈ rights rm-cap)
  assume b0: rm-cap ∈ get-domain-cap-set-from-domain-id s did
    ∧ REMOVE ∈ rights rm-cap
    ∧ right-to-rm ∈ rights rm-cap

```

```

have get-domain-cap-set-from-domain-id s d
  = get-domain-cap-set-from-domain-id s' d
proof (cases REMOVE = right-to-rm
   $\wedge \{REMOVE\} = \text{rights rm-cap}$ )
  assume c0: REMOVE = right-to-rm
   $\wedge \{REMOVE\} = \text{rights rm-cap}$ 
  let ?cs-dst = get-domain-cap-set-from-domain-id s did
  let ?cs-rest = {c. c ∈ ?cs-dst  $\wedge$  c ≠ rm-cap}
  have c1: ((caps s') did) = ?cs-rest
    using b0 c0 p2 remove-cap-right-def by auto
  have c2:  $\forall v. v \neq did$ 
     $\longrightarrow ((caps s') v) = ((caps s) v)$ 
    using b0 c0 p2 remove-cap-right-def by auto
  have c3:  $\forall v. v \neq did$ 
     $\longrightarrow \text{get-domain-cap-set-from-domain-id s v}$ 
    = get-domain-cap-set-from-domain-id s' v
    using c2 get-domain-cap-set-from-domain-id-def by auto
  have c4: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using c3 a0 by auto
  then show ?thesis by auto
next
  assume c0:  $\neg(RREMOVE = \text{right-to-rm}$ 
     $\wedge \{REMOVE\} = \text{rights rm-cap}$ )
  let ?cs-dst = get-domain-cap-set-from-domain-id s did
  let ?cs-rest = {c. c ∈ ?cs-dst  $\wedge$  c ≠ rm-cap}
  let ?new-cap = () target = target rm-cap,
    rights = (rights rm-cap) - {right-to-rm}()
  have c1: ((caps s') did) = (insert ?new-cap ?cs-rest)
    using b0 c0 p2 remove-cap-right-def by auto
  have c2:  $\forall v. v \neq did$ 
     $\longrightarrow ((caps s') v) = ((caps s) v)$ 
    using b0 c0 p2 remove-cap-right-def by auto
  have c3:  $\forall v. v \neq did$ 
     $\longrightarrow \text{get-domain-cap-set-from-domain-id s v}$ 
    = get-domain-cap-set-from-domain-id s' v
    using c2 get-domain-cap-set-from-domain-id-def by auto
  have c4: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using c3 a0 by auto
  then show ?thesis by auto
qed
then show ?thesis by auto
next
  assume b0:  $\neg(\text{rm-cap} \in \text{get-domain-cap-set-from-domain-id s did}$ 
     $\wedge REMOVE \in \text{rights rm-cap}$ 
     $\wedge \text{right-to-rm} \in \text{rights rm-cap})$ 
  have b1: s' = s
    using b0 p2 remove-cap-right-def by auto
  have b2: get-domain-cap-set-from-domain-id s d
    = get-domain-cap-set-from-domain-id s' d
    using b1 get-domain-cap-set-from-domain-id-def by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

```

lemma remove-cap-right-notchg-policy:
assumes p0: reachable0 s
  and p1:  $\neg(\text{interferes } \text{did } s \ d)$ 
  and p2:  $s' = \text{fst } (\text{remove-cap-right } s \ \text{did } \text{rm-cap } \text{right-to-rm})$ 
shows  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
proof -
{
  have a0:  $d \neq \text{did}$ 
    using p1 interferes-def by auto
  have  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
    proof (cases  $\text{rm-cap} \in \text{get-domain-cap-set-from-domain-id } s \ \text{did}$ 
       $\wedge \text{REMOVE} \in \text{rights } \text{rm-cap}$ 
       $\wedge \text{right-to-rm} \in \text{rights } \text{rm-cap}$ )
    assume b0:  $\text{rm-cap} \in \text{get-domain-cap-set-from-domain-id } s \ \text{did}$ 
       $\wedge \text{REMOVE} \in \text{rights } \text{rm-cap}$ 
       $\wedge \text{right-to-rm} \in \text{rights } \text{rm-cap}$ 
    have a1:  $d \neq \text{target } \text{rm-cap}$ 
      by (metis b0 interferes-def p1)
    have  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
      proof (cases  $\text{REMOVE} = \text{right-to-rm}$ 
         $\wedge \{\text{REMOVE}\} = \text{rights } \text{rm-cap}$ )
      assume c0:  $\text{REMOVE} = \text{right-to-rm}$ 
         $\wedge \{\text{REMOVE}\} = \text{rights } \text{rm-cap}$ 
      let ?cs-dst =  $\text{get-domain-cap-set-from-domain-id } s \ \text{did}$ 
      let ?cs-rest =  $\{c. c \in ?cs\text{-dst} \wedge c \neq \text{rm-cap}\}$ 
      have c1:  $((\text{caps } s') \ \text{did}) = ?cs\text{-rest}$ 
        using b0 c0 p2 remove-cap-right-def by auto
      have c2:  $\forall v. v \neq \text{did}$ 
         $\longrightarrow ((\text{caps } s') \ v) = ((\text{caps } s) \ v)$ 
        using b0 c0 p2 remove-cap-right-def by auto
      have c3:  $\forall v. v \neq \text{did}$ 
         $\longrightarrow \text{get-domain-cap-set-from-domain-id } s \ v$ 
           $= \text{get-domain-cap-set-from-domain-id } s' \ v$ 
        using c2 get-domain-cap-set-from-domain-id-def by auto
      have c4:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
         $= \text{get-domain-cap-set-from-domain-id } s' \ d$ 
        using c3 a0 by auto
      have c5:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
         $= \text{get-domain-cap-set-from-domain-id } s' \ d$ 
        using c3 a0 by auto
      have c6:  $\forall v. v \neq \text{did}$ 
         $\longrightarrow \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d$ 
        using c3 c5 interferes-def by auto
      have c7:  $((\text{caps } s') \ \text{did}) = ?cs\text{-rest}$ 
        using b0 c0 p2 remove-cap-right-def by auto
      have c8:  $\text{get-domain-cap-set-from-domain-id } s' \ \text{did}$ 
         $= ?cs\text{-rest}$ 
        using c7 get-domain-cap-set-from-domain-id-def by auto
      have c9:  $\text{get-domain-cap-set-from-domain-id } s' \ \text{did}$ 
         $\subseteq \text{get-domain-cap-set-from-domain-id } s \ \text{did}$ 
        using c8 by auto
      have c10:  $\neg(\exists c. c \in (\text{get-domain-cap-set-from-domain-id } s \ \text{did}) \wedge \text{target } c = d)$ 
        by (metis interferes-def p1)
      have c11:  $\neg(\exists c. c \in (\text{get-domain-cap-set-from-domain-id } s' \ \text{did}) \wedge \text{target } c = d)$ 
        using c10 c9 by auto
      have c12:  $\neg(\text{interferes } \text{did } s' \ d)$ 
        using a0 c11 interferes-def by auto
      have c13:  $\text{interferes } \text{did } s' \ d = \text{interferes } \text{did } s \ d$ 

```

```

    using c12 p1 by auto
    have c14:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
    using c6 c13 by auto
    then show ?thesis by auto
next
    assume c0:  $\neg(\text{REMOVE} = \text{right-to-rm}$ 
       $\wedge \{\text{REMOVE}\} = \text{rights rm-cap})$ 
    let ?cs-dst =  $\text{get-domain-cap-set-from-domain-id } s \ \text{did}$ 
    let ?cs-rest =  $\{c. c \in ?cs\text{-dst} \wedge c \neq \text{rm-cap}\}$ 
    let ?new-cap =  $(\text{target} = \text{target rm-cap},$ 
       $\text{rights} = (\text{rights rm-cap}) - \{\text{right-to-rm}\})$ 
    have c1:  $((\text{caps } s') \ \text{did}) = (\text{insert } ?new\text{-cap } ?cs\text{-rest})$ 
    using b0 c0 p2 remove-cap-right-def by auto
    have c2:  $\forall v. v \neq \text{did}$ 
       $\longrightarrow ((\text{caps } s') \ v) = ((\text{caps } s) \ v)$ 
    using b0 c0 p2 remove-cap-right-def by auto
    have c3:  $\forall v. v \neq \text{did}$ 
       $\longrightarrow \text{get-domain-cap-set-from-domain-id } s \ v$ 
       $= \text{get-domain-cap-set-from-domain-id } s' \ v$ 
    using c2 get-domain-cap-set-from-domain-id-def by auto
    have c4:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
       $= \text{get-domain-cap-set-from-domain-id } s' \ d$ 
    using c3 a0 by auto
    have c5:  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
       $= \text{get-domain-cap-set-from-domain-id } s' \ d$ 
    using c3 a0 by auto
    have c6:  $\forall v. v \neq \text{did}$ 
       $\longrightarrow \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d$ 
    using c3 c5 interferes-def by auto
    have c7:  $\text{get-domain-cap-set-from-domain-id } s' \ \text{did}$ 
       $= (\text{insert } ?new\text{-cap } ?cs\text{-rest})$ 
    using c1 get-domain-cap-set-from-domain-id-def by auto
    have c8:  $\neg(\exists c. c \in (\text{get-domain-cap-set-from-domain-id } s \ \text{did}) \wedge \text{target } c = d)$ 
    by (metis interferes-def p1)
    have c9:  $\neg(\exists c. c \in (\text{get-domain-cap-set-from-domain-id } s' \ \text{did}) \wedge \text{target } c = d)$ 
    using c8 c7 a1 by auto
    have c10:  $\neg(\text{interferes } \text{did } s' \ d)$ 
    using a0 c9 interferes-def by auto
    have c11:  $\text{interferes } \text{did } s' \ d = \text{interferes } \text{did } s \ d$ 
    using c10 p1 by auto
    have c12:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
    using c6 c11 by auto
    then show ?thesis by auto
qed
then show ?thesis by auto
next
    assume b0:  $\neg(\text{rm-cap} \in \text{get-domain-cap-set-from-domain-id } s \ \text{did}$ 
       $\wedge \text{REMOVE} \in \text{rights rm-cap}$ 
       $\wedge \text{right-to-rm} \in \text{rights rm-cap})$ 
    have b1:  $s' = s$ 
    using b0 p2 remove-cap-right-def by auto
    have b2:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
    using b1 interferes-def by auto
    then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

```

lemma remove-cap-right-notchg-dom-eps:
  assumes p0: reachable0 s
    and p2: s' = fst (remove-cap-right s did rm-cap right-to-rm)
  shows get-endpoints-of-domain s d = get-endpoints-of-domain s' d
  proof -
  {
    have get-endpoints-of-domain s d = get-endpoints-of-domain s' d
      proof (cases rm-cap ∈ get-domain-cap-set-from-domain-id s did
        ∧ REMOVE ∈ rights rm-cap
        ∧ right-to-rm ∈ rights rm-cap)
      assume b0: rm-cap ∈ get-domain-cap-set-from-domain-id s did
        ∧ REMOVE ∈ rights rm-cap
        ∧ right-to-rm ∈ rights rm-cap
      have get-endpoints-of-domain s d = get-endpoints-of-domain s' d
        proof (cases REMOVE = right-to-rm
          ∧ {REMOVE} = rights rm-cap)
        assume c0: REMOVE = right-to-rm
          ∧ {REMOVE} = rights rm-cap
        let ?cs-dst = get-domain-cap-set-from-domain-id s did
        let ?cs-rest = {c. c ∈ ?cs-dst ∧ c ≠ rm-cap}
        have c1: ((caps s') did) = ?cs-rest
          using b0 c0 p2 remove-cap-right-def by auto
        have c2: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
          using b0 c0 p2 remove-cap-right-def get-endpoints-of-domain-def by auto
        then show ?thesis by auto
      next
        assume c0: ¬(REMOVE = right-to-rm
          ∧ {REMOVE} = rights rm-cap)
        let ?cs-dst = get-domain-cap-set-from-domain-id s did
        let ?cs-rest = {c. c ∈ ?cs-dst ∧ c ≠ rm-cap}
        let ?new-cap = (| target = target rm-cap,
          rights = (rights rm-cap) - {right-to-rm}|)
        have c1: ((caps s') did) = (insert ?new-cap ?cs-rest)
          using b0 c0 p2 remove-cap-right-def by auto
        have c2: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
          using b0 c0 p2 remove-cap-right-def get-endpoints-of-domain-def by auto
        then show ?thesis by auto
      qed
    }
    then show ?thesis by auto
  next
    assume b0: ¬(rm-cap ∈ get-domain-cap-set-from-domain-id s did
      ∧ REMOVE ∈ rights rm-cap
      ∧ right-to-rm ∈ rights rm-cap)
    have b1: s' = s
      using b0 p2 remove-cap-right-def by auto
    have b2: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
      using b1 get-endpoints-of-domain-def by auto
    then show ?thesis by auto
  qed
}
then show ?thesis by auto
qed

lemma remove-cap-right-notchg-ep-msgs:
  assumes p0: reachable0 s
    and p2: s' = fst (remove-cap-right s did rm-cap right-to-rm)
  shows (∀ ep. ep ∈ get-endpoints-of-domain s d

```

```

    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
proof -
{
  have (∀ ep. ep ∈ get-endpoints-of-domain s d
    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
  proof (cases rm-cap ∈ get-domain-cap-set-from-domain-id s did
    ∧ REMOVE ∈ rights rm-cap
    ∧ right-to-rm ∈ rights rm-cap)
  assume b0: rm-cap ∈ get-domain-cap-set-from-domain-id s did
    ∧ REMOVE ∈ rights rm-cap
    ∧ right-to-rm ∈ rights rm-cap
  have (∀ ep. ep ∈ get-endpoints-of-domain s d
    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
  proof (cases REMOVE = right-to-rm
    ∧ {REMOVE} = rights rm-cap)
  assume c0: REMOVE = right-to-rm
    ∧ {REMOVE} = rights rm-cap
  let ?cs-dst = get-domain-cap-set-from-domain-id s did
  let ?cs-rest = {c. c ∈ ?cs-dst ∧ c ≠ rm-cap}
  have c1: ((caps s') did) = ?cs-rest
    using b0 c0 p2 remove-cap-right-def by auto
  have c2: (∀ ep. ep ∈ get-endpoints-of-domain s d
    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
    using b0 c0 p2 remove-cap-right-def get-endpoints-of-domain-def
    get-msg-set-from-endpoint-id-def by auto
  then show ?thesis by auto
next
  assume c0: ¬(REMOVE = right-to-rm
    ∧ {REMOVE} = rights rm-cap)
  let ?cs-dst = get-domain-cap-set-from-domain-id s did
  let ?cs-rest = {c. c ∈ ?cs-dst ∧ c ≠ rm-cap}
  let ?new-cap = (| target = target rm-cap,
    rights = (rights rm-cap) - {right-to-rm}|)
  have c1: ((caps s') did) = (insert ?new-cap ?cs-rest)
    using b0 c0 p2 remove-cap-right-def by auto
  have c2: (∀ ep. ep ∈ get-endpoints-of-domain s d
    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
    using b0 c0 p2 remove-cap-right-def get-endpoints-of-domain-def
    get-msg-set-from-endpoint-id-def by auto
  then show ?thesis by auto
qed
then show ?thesis by auto
next
  assume b0: ¬(rm-cap ∈ get-domain-cap-set-from-domain-id s did
    ∧ REMOVE ∈ rights rm-cap
    ∧ right-to-rm ∈ rights rm-cap)
  have b1: s' = s
    using b0 p2 remove-cap-right-def by auto
  have b2: (∀ ep. ep ∈ get-endpoints-of-domain s d
    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id s' ep )
    using b1 get-endpoints-of-domain-def by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma remove-cap-right-lcl-resp:

```

assumes  $p0$ : reachable0  $s$ 
  and  $p1$ :  $\neg(\text{interferes } did \ s \ d)$ 
  and  $p2$ :  $s' = fst \ (\text{remove-cap-right } s \ did \ rm\text{-cap } right\text{-to-rm})$ 
shows  $s \sim d \sim s'$ 
proof –
{
  have  $a0$ :  $did \neq d$ 
    using  $p1$  interferes-def by auto
  have  $a1$ :  $\text{get-domain-cap-set-from-domain-id } s \ d$ 
     $= \text{get-domain-cap-set-from-domain-id } s' \ d$ 
    using  $p0 \ p1 \ p2$  remove-cap-right-notchg-domain-cap-set by auto
  have  $a2$ :  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
    using  $p0 \ p1 \ p2$  remove-cap-right-notchg-policy by auto
  have  $a3$ :  $\text{get-endpoints-of-domain } s \ d = \text{get-endpoints-of-domain } s' \ d$ 
    using  $p0 \ p1 \ p2$  remove-cap-right-notchg-dom-eps by auto
  have  $a4$ :  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } s' \ ep)$ 
    using  $p0 \ p1 \ p2$  remove-cap-right-notchg-ep-msgs by auto
  have  $a5$ :  $s \sim d \sim s'$ 
    using  $a1 \ a2 \ a3 \ a4$  by auto
}
then show ?thesis by auto
qed

```

```

lemma remove-cap-right-lcl-resp-e:
assumes  $p0$ : reachable0  $s$ 
  and  $p1$ :  $a = (\text{Remove-Cap-Right } did \ dst\text{-cap } right\text{-to-rm})$ 
  and  $p2$ :  $\neg(\text{interferes } (the \ (\text{domain-of-event } a)) \ s \ d)$ 
  and  $p3$ :  $s' = \text{exec-event } s \ a$ 
shows  $s \sim d \sim s'$ 
proof –
{
  have  $a0$ :  $(the \ (\text{domain-of-event } a)) = did$ 
    using  $p1$  domain-of-event-def by auto
  have  $a1$ :  $s' = fst \ (\text{remove-cap-right } s \ did \ dst\text{-cap } right\text{-to-rm})$ 
    using  $p1 \ p3$  exec-event-def by auto
  have  $a2$ :  $\neg(\text{interferes } did \ s \ d)$ 
    using  $p2 \ a0$  by auto
  have  $a3$ :  $s \sim d \sim s'$ 
    using  $a1 \ a2 \ p0$  remove-cap-right-lcl-resp by blast
}
then show ?thesis by auto
qed

```

```

lemma remove-cap-right-lcrsp-e: dynamic-local-respect-e  $(\text{Remove-Cap-Right } did \ dst\text{-cap } right\text{-to-rm})$ 
proof –
{
  have  $\forall d \ s. \text{reachable0 } s$ 
     $\wedge \neg(\text{interferes } (the \ (\text{domain-of-event } (\text{Remove-Cap-Right } did \ dst\text{-cap } right\text{-to-rm}))) \ s \ d)$ 
     $\longrightarrow (s \sim d \sim (\text{exec-event } s \ (\text{Remove-Cap-Right } did \ dst\text{-cap } right\text{-to-rm})))$ 
  proof –
  {
    fix  $d \ s$ 
    assume  $p1$ : reachable0  $s$ 
    assume  $p2$ :  $\neg(\text{interferes } (the \ (\text{domain-of-event } (\text{Remove-Cap-Right } did \ dst\text{-cap } right\text{-to-rm}))) \ s \ d)$ 
    have  $(s \sim d \sim (\text{exec-event } s \ (\text{Remove-Cap-Right } did \ dst\text{-cap } right\text{-to-rm})))$ 
      using  $p1 \ p2$  remove-cap-right-lcl-resp-e by blast
  }
}

```

```

    then show ?thesis by blast
  qed
}
then show ?thesis
  using dynamic-local-respect-e-def by blast
qed

```

### 0.9.8 proving the "dynamic local respect" property

**definition** *dynamic-local-respect-c* :: bool **where**  

$$\text{dynamic-local-respect-c} \equiv \forall a\ d\ s. \text{reachable0}\ s$$

$$\wedge \neg(\text{interferes}\ (\text{the}\ (\text{domain-of-event}\ a))\ s\ d)$$

$$\longrightarrow (s \sim d \sim (\text{exec-event}\ s\ a))$$

**theorem** *dynamic-local-respect:dynamic-local-respect*

```

proof -
{
  fix e
  have dynamic-local-respect-e e
  proof (induct e)
    case (Client-Lookup-Endpoint-Name x x1)
      show ?case
      using client-lookup-endpoint-name-lcrsp-e by blast
    case (Send-Queuing-Message x1a x2 x3a)
      show ?case
      using send-queuing-message-lcl-lcrsp-e by blast
    case (Receive-Queuing-Message x)
      show ?case
      using receive-queuing-message-lcrsp-e by blast
    case (Get-My-Endpoints-Set x)
      show ?case
      using get-my-endpoints-set-lcrsp-e by blast
    case (Get-Caps x)
      show ?case
      using get-caps-lcrsp-e by blast
    case (Grant-Endpoint-Cap x1a x2 x3a)
      show ?case
      using grant-endpoint-cap-lcrsp-e by blast
    case (Remove-Cap-Right x1a x2 x3a)
      show ?case
      using remove-cap-right-lcrsp-e by blast
  qed
}
then show ?thesis
  using dynamic-local-respect-all-evt by blast
qed

```

### 0.10 Concrete unwinding condition of "weakly step consistent"

#### 0.10.1 proving "client lookup endpoint name" satisfying the "weakly step consistent" property

**lemma** *client-lookup-endpoint-name-wsc*:

```

assumes p0: reachable0 s
and p1: reachable0 t
and p2: s ~ d ~ t
and p3: interferes did s d
and p4: s ~ did ~ t
and p5: s' = fst (client-lookup-endpoint-name sysconf s did ename)
and p6: t' = fst (client-lookup-endpoint-name sysconf t did ename)

```



```

shows  $s' \sim d \sim t'$ 
proof -
{
  have  $a0: s = s'$ 
    using  $p5$  client-lookup-endpoint-name-def by auto
  have  $a1: t = t'$ 
    using  $p6$  client-lookup-endpoint-name-def by auto
  have  $a2: s' \sim d \sim t'$ 
    using  $a0$   $a1$   $p2$  by auto
}
then show ?thesis by auto
qed

```

```

lemma client-lookup-endpoint-name-wsc-e:
assumes  $p0: \text{reachable0 } s$ 
  and  $p1: \text{reachable0 } t$ 
  and  $p2: a = (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})$ 
  and  $p3: s \sim d \sim t$ 
  and  $p4: \text{interferes } (\text{the } (\text{domain-of-event } a)) \text{ } s \text{ } d$ 
  and  $p5: s \sim \text{did} \sim t$ 
  and  $p6: s' = \text{exec-event } s \text{ } a$ 
  and  $p7: t' = \text{exec-event } t \text{ } a$ 
shows  $s' \sim d \sim t'$ 
proof -
{
  have  $a0: (\text{the } (\text{domain-of-event } a)) = \text{did}$ 
    using  $p2$  domain-of-event-def by auto
  have  $a1: s' = \text{fst } (\text{client-lookup-endpoint-name } \text{sysconf } s \text{ } \text{did } \text{ename})$ 
    using  $p2$   $p6$  exec-event-def by auto
  have  $a2: t' = \text{fst } (\text{client-lookup-endpoint-name } \text{sysconf } t \text{ } \text{did } \text{ename})$ 
    using  $p2$   $p7$  exec-event-def by auto
  have  $a3: (\text{interferes } \text{did } s \text{ } d)$ 
    using  $p4$   $a0$  by auto
  have  $a4: s' \sim d \sim t'$ 
    using  $a1$   $a2$   $a3$   $p0$   $p1$   $p3$   $p5$  client-lookup-endpoint-name-wsc by blast
}
then show ?thesis by auto
qed

```

```

lemma client-lookup-endpoint-name-dwsc-e: dynamic-weakly-step-consistent-e  $(\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})$ 
proof -
{
  have  $\forall d \ s \ t. \text{reachable0 } s \wedge \text{reachable0 } t$ 
     $\wedge (s \sim d \sim t)$ 
     $\wedge (\text{interferes } (\text{the } (\text{domain-of-event } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))) \ s \ d)$ 
     $\wedge (s \sim (\text{the } (\text{domain-of-event } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))) \sim t)$ 
     $\longrightarrow ((\text{exec-event } s \text{ } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})) \sim d \sim (\text{exec-event } t \text{ } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))$ 
  proof -
  {
    fix  $d \ s \ t$ 
    assume  $p1: \text{reachable0 } s$ 
    assume  $p2: \text{reachable0 } t$ 
    assume  $p3: (s \sim d \sim t)$ 
    assume  $p4: (\text{interferes } (\text{the } (\text{domain-of-event } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))) \ s \ d)$ 
    assume  $p5: (s \sim (\text{the } (\text{domain-of-event } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))) \sim t)$ 
    have  $((\text{exec-event } s \text{ } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})) \sim d \sim (\text{exec-event } t \text{ } (\text{Client-Lookup-Endpoint-Name } \text{did } \text{ename})))$ 

```

```

    by (metis Event.simps(50) client-lookup-endpoint-name-wsc-e domain-of-event-def option.sel p1 p2 p3 p4 p5)
  }
  then show ?thesis by blast
qed
}
then show ?thesis
  using dynamic-weakly-step-consistent-e-def by blast
qed

```

### 0.10.2 proving "send queuing message" satisfying the "weakly step consistent" property

```

lemma send-queuing-message-wsc-domain-cap-set:
assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s ~ d ~ t
  and p3: interferes did s d
  and p4: s ~ did ~ t
  and p5: s' = fst (send-queuing-message s did eid m)
  and p6: t' = fst (send-queuing-message t did eid m)
shows get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: (∀ v. interferes v s d ⟷ interferes v t d)
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
  proof (cases eid ∈ get-endpoints-of-domain s d)
    assume b0: eid ∈ get-endpoints-of-domain s d
    have b1: eid ∈ {x. (domain-endpoint s) x = Some d}
      using b0 get-endpoints-of-domain-def by auto
    have b2: (domain-endpoint s) eid = Some d
      using b1 by auto
    have b3: get-domain-id-from-endpoint s eid = Some d
      using b2 get-domain-id-from-endpoint-def by auto
    have b4: get-domain-id-from-endpoint s eid ≠ None
      using b3 by auto
    have b5: eid ∈ get-endpoints-of-domain t d
      using b0 a0 by auto
    have b6: eid ∈ {x. (domain-endpoint t) x = Some d}
      using b5 get-endpoints-of-domain-def by auto
    have b7: (domain-endpoint t) eid = Some d
      using b6 by auto
    have b8: get-domain-id-from-endpoint t eid = Some d
      using b7 get-domain-id-from-endpoint-def by auto
    have b9: get-domain-id-from-endpoint t eid ≠ None
      using b8 by auto
    have b10: (the (get-domain-id-from-endpoint s eid)) = d
      using b3 by auto
    have b11: (the (get-domain-id-from-endpoint t eid)) = d
      using b8 by auto
    have b12: interferes did s (the (get-domain-id-from-endpoint s eid))
      using p3 b10 by auto
    have b13: interferes did t (the (get-domain-id-from-endpoint t eid))

```

```

    using a2 b11 by auto
  have b14: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id s' d
    using b9 b12 p5 get-domain-cap-set-from-domain-id-def send-queuing-message-def by auto
  have b15: get-domain-cap-set-from-domain-id t d = get-domain-cap-set-from-domain-id t' d
    using b10 b13 p6 get-domain-cap-set-from-domain-id-def send-queuing-message-def by auto
  have b16: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
    using b14 b15 a3 by auto
  then show ?thesis by auto
next
  assume b0: eid  $\notin$  get-endpoints-of-domain s d
  have b1: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id s' d
    using p5 get-domain-cap-set-from-domain-id-def send-queuing-message-def by auto
  have b2: get-domain-cap-set-from-domain-id t d = get-domain-cap-set-from-domain-id t' d
    using p6 get-domain-cap-set-from-domain-id-def send-queuing-message-def by auto
  have b16: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
    using b1 b2 a3 by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma send-queuing-message-wsc-policy:

```

assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s  $\sim$  d  $\sim$  t
  and p3: interferes did s d
  and p4: s  $\sim$  did  $\sim$  t
  and p5: s' = fst (send-queuing-message s did eid m)
  and p6: t' = fst (send-queuing-message t did eid m)
shows ( $\forall v$ . interferes v s' d  $\longleftrightarrow$  interferes v t' d)
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v$ . interferes v s d  $\longleftrightarrow$  interferes v t d)
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: ( $\forall v$ . interferes v s d  $\longleftrightarrow$  interferes v s' d)
    using p0 p5 send-queuing-message-notchg-policy by auto
  have a5: ( $\forall v$ . interferes v t d  $\longleftrightarrow$  interferes v t' d)
    using p1 p6 send-queuing-message-notchg-policy by auto
  have a6: ( $\forall v$ . interferes v s' d  $\longleftrightarrow$  interferes v t' d)
    using a1 a4 a5 by auto
}
then show ?thesis by auto
qed

```

lemma send-queuing-message-wsc-dom-eps:

```

assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s  $\sim$  d  $\sim$  t
  and p3: interferes did s d
  and p4: s  $\sim$  did  $\sim$  t
  and p5: s' = fst (send-queuing-message s did eid m)

```

```

and p6: t' = fst (send-queuing-message t did eid m)
shows get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ t \ d$ )
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
    using p0 p5 send-queuing-message-notchg-dom-eps by auto
  have a5: get-endpoints-of-domain t d = get-endpoints-of-domain t' d
    using p1 p6 send-queuing-message-notchg-dom-eps by auto
  have a6: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
    using a0 a4 a5 by auto
}
then show ?thesis by auto
qed

```

lemma send-queuing-message-wsc-ep-msgs:

```

assumes p0: reachable0 s
and p1: reachable0 t
and p2: s ~ d ~ t
and p3: interferes did s d
and p4: s ~ did ~ t
and p5: s' = fst (send-queuing-message s did eid m)
and p6: t' = fst (send-queuing-message t did eid m)
shows ( $\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$ 
 $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep$ )
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ t \ d$ )
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
 $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } t \ ep$ )
    by (meson p2 vpeq1-def)
  have a5: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$ 
 $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep$ )
    proof (cases eid  $\in$  get-endpoints-of-domain s d)
      assume b0: eid  $\in$  get-endpoints-of-domain s d
      have b1: eid  $\in$  {x. (domain-endpoint s) x = Some d}
        using b0 get-endpoints-of-domain-def by auto
      have b2: (domain-endpoint s) eid = Some d
        using b1 by auto
      have b3: get-domain-id-from-endpoint s eid = Some d
        using b2 get-domain-id-from-endpoint-def by auto
      have b4: get-domain-id-from-endpoint s eid  $\neq$  None
        using b3 by auto
      have b5: eid  $\in$  get-endpoints-of-domain t d

```

```

  using b0 a0 by auto
  have b6:  $eid \in \{x. (\text{domain-endpoint } t) \ x = \text{Some } d\}$ 
    using b5 get-endpoints-of-domain-def by auto
  have b7:  $(\text{domain-endpoint } t) \ eid = \text{Some } d$ 
    using b6 by auto
  have b8:  $\text{get-domain-id-from-endpoint } t \ eid = \text{Some } d$ 
    using b7 get-domain-id-from-endpoint-def by auto
  have b9:  $\text{get-domain-id-from-endpoint } t \ eid \neq \text{None}$ 
    using b8 by auto
  have b10:  $(\text{the } (\text{get-domain-id-from-endpoint } s \ eid)) = d$ 
    using b3 by auto
  have b11:  $(\text{the } (\text{get-domain-id-from-endpoint } t \ eid)) = d$ 
    using b8 by auto
  have b12:  $\text{interferes } did \ s \ (\text{the } (\text{get-domain-id-from-endpoint } s \ eid))$ 
    using p3 b10 by auto
  have b13:  $\text{interferes } did \ t \ (\text{the } (\text{get-domain-id-from-endpoint } t \ eid))$ 
    using a2 b11 by auto
  let ?new-msg-set-s = insert m (get-msg-set-from-endpoint-id s eid)
  let ?new-msg-set-t = insert m (get-msg-set-from-endpoint-id t eid)
  have b14:  $(e\text{-msgs } s') \ eid = ?\text{new-msg-set-s}$ 
    using b4 b12 p5 send-queuing-message-def by auto
  have b15:  $(e\text{-msgs } t') \ eid = ?\text{new-msg-set-t}$ 
    using b9 b13 p6 send-queuing-message-def by auto
  have b16:  $\forall e. e \neq eid$ 
     $\longrightarrow ((e\text{-msgs } s) \ e) = ((e\text{-msgs } s') \ e)$ 
    using b4 b12 p5 send-queuing-message-def by auto
  have b17:  $\forall e. e \neq eid$ 
     $\longrightarrow ((e\text{-msgs } t) \ e) = ((e\text{-msgs } t') \ e)$ 
    using b9 b13 p6 send-queuing-message-def by auto
  have b18:  $\forall e. e \neq eid$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ e = \text{get-msg-set-from-endpoint-id } s' \ e$ 
    using b16 get-msg-set-from-endpoint-id-def by auto
  have b19:  $\forall e. e \neq eid$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } t \ e = \text{get-msg-set-from-endpoint-id } t' \ e$ 
    using b17 get-msg-set-from-endpoint-id-def by auto
  have b20:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \ d \wedge ep \neq eid$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep)$ 
    using a4 b19 b18 by auto
  have b21:  $\text{get-msg-set-from-endpoint-id } s \ eid = \text{get-msg-set-from-endpoint-id } t \ eid$ 
    using a4 b0 b5 by auto
  have b22:  $(e\text{-msgs } s') \ eid = (e\text{-msgs } t') \ eid$ 
    using b21 b14 b15 by auto
  have b23:  $\text{get-msg-set-from-endpoint-id } s' \ eid = \text{get-msg-set-from-endpoint-id } t' \ eid$ 
    using b22 get-msg-set-from-endpoint-id-def by auto
  have b24:  $\text{get-endpoints-of-domain } s \ d = \text{get-endpoints-of-domain } s' \ d$ 
    using p0 p5 send-queuing-message-notchg-dom-eps by auto
  have b25:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep)$ 
    using b23 b24 b20 by auto
  then show ?thesis by auto
next
  assume b0:  $eid \notin \text{get-endpoints-of-domain } s \ d$ 
  have b1:  $\forall e. e \neq eid$ 
     $\longrightarrow ((e\text{-msgs } s) \ e) = ((e\text{-msgs } s') \ e)$ 
    using p5 send-queuing-message-def by auto
  have b2:  $\forall e. e \neq eid$ 
     $\longrightarrow ((e\text{-msgs } t) \ e) = ((e\text{-msgs } t') \ e)$ 
    using p6 send-queuing-message-def by auto

```

```

have b3:  $\forall e. e \neq \text{eid}$ 
   $\rightarrow \text{get-msg-set-from-endpoint-id } s \ e = \text{get-msg-set-from-endpoint-id } s' \ e$ 
using b1 get-msg-set-from-endpoint-id-def by auto
have b4:  $\forall e. e \neq \text{eid}$ 
   $\rightarrow \text{get-msg-set-from-endpoint-id } t \ e = \text{get-msg-set-from-endpoint-id } t' \ e$ 
using b2 get-msg-set-from-endpoint-id-def by auto
have b5:  $\forall e. e \in \text{get-endpoints-of-domain } s \ d \wedge e \neq \text{eid}$ 
   $\rightarrow \text{get-msg-set-from-endpoint-id } s' \ e = \text{get-msg-set-from-endpoint-id } t' \ e$ 
using b3 b4 a4 by auto
have b6:  $\forall e. e \in \text{get-endpoints-of-domain } s \ d$ 
   $\rightarrow \text{get-msg-set-from-endpoint-id } s' \ e = \text{get-msg-set-from-endpoint-id } t' \ e$ 
using b5 b0 by auto
have b7:  $\text{get-endpoints-of-domain } s \ d = \text{get-endpoints-of-domain } s' \ d$ 
using p0 p5 send-queuing-message-notchg-dom-eps by auto
have b8:  $\forall e. e \in \text{get-endpoints-of-domain } s' \ d$ 
   $\rightarrow \text{get-msg-set-from-endpoint-id } s' \ e = \text{get-msg-set-from-endpoint-id } t' \ e$ 
using b7 b6 by auto
then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma send-queuing-message-wsc:

```

assumes p0: reachable0 s
and p1: reachable0 t
and p2:  $s \sim d \sim t$ 
and p3: interferes did s d
and p4:  $s \sim \text{did} \sim t$ 
and p5:  $s' = \text{fst } (\text{send-queuing-message } s \ \text{did} \ \text{eid} \ m)$ 
and p6:  $t' = \text{fst } (\text{send-queuing-message } t \ \text{did} \ \text{eid} \ m)$ 
shows  $s' \sim d \sim t'$ 
proof -
{
have a0:  $\text{get-domain-cap-set-from-domain-id } s' \ d = \text{get-domain-cap-set-from-domain-id } t' \ d$ 
using p0 p1 p2 p3 p4 p5 p6 send-queuing-message-wsc-domain-cap-set by blast
have a1:  $(\forall v. \text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d)$ 
using p0 p1 p2 p3 p4 p5 p6 send-queuing-message-wsc-policy by blast
have a2:  $\text{get-endpoints-of-domain } s' \ d = \text{get-endpoints-of-domain } t' \ d$ 
using p0 p1 p2 p3 p4 p5 p6 send-queuing-message-wsc-dom-eps by blast
have a3:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$ 
   $\rightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep)$ 
using p0 p1 p2 p3 p4 p5 p6 send-queuing-message-wsc-ep-msgs by blast
have a4:  $s' \sim d \sim t'$ 
using a0 a1 a2 a3 vpeq1-def by auto
}
then show ?thesis by auto
qed

```

lemma send-queuing-message-wsc-e:

```

assumes p0: reachable0 s
and p1: reachable0 t
and p2:  $a = (\text{Send-Queuing-Message } \text{did} \ \text{eid} \ m)$ 
and p3:  $s \sim d \sim t$ 
and p4: interferes (the (domain-of-event a)) s d
and p5:  $s \sim \text{did} \sim t$ 
and p6:  $s' = \text{exec-event } s \ a$ 

```

```

and p7: t' = exec-event t a
shows s' ~ d ~ t'
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p2 domain-of-event-def by auto
  have a1: s' = fst (send-queuing-message s did eid m)
    using p2 p6 exec-event-def by auto
  have a2: t' = fst (send-queuing-message t did eid m)
    using p2 p7 exec-event-def by auto
  have a3: (interferes did s d)
    using p4 a0 by auto
  have a4: s' ~ d ~ t'
    using a1 a2 a3 p0 p1 p3 p5 send-queuing-message-wsc by blast
}
then show ?thesis by auto
qed

```

```

lemma send-queuing-message-dwsc-e: dynamic-weakly-step-consistent-e (Send-Queuing-Message did eid m)
proof -

```

```

{
  have  $\forall d s t. \text{reachable0 } s \wedge \text{reachable0 } t$ 
     $\wedge (s \sim d \sim t)$ 
     $\wedge (\text{interferes } (\text{the } (\text{domain-of-event } (\text{Send-Queuing-Message did eid m}))) s d)$ 
     $\wedge (s \sim (\text{the } (\text{domain-of-event } (\text{Send-Queuing-Message did eid m}))) \sim t)$ 
     $\longrightarrow ((\text{exec-event } s (\text{Send-Queuing-Message did eid m})) \sim d \sim (\text{exec-event } t (\text{Send-Queuing-Message did eid m})))$ 
  proof -
    {
      fix d s t
      assume p1: reachable0 s
      assume p2: reachable0 t
      assume p3:  $s \sim d \sim t$ 
      assume p4: (interferes (the (domain-of-event (Send-Queuing-Message did eid m))) s d)
      assume p5:  $s \sim (\text{the } (\text{domain-of-event } (\text{Send-Queuing-Message did eid m}))) \sim t$ 
      have  $((\text{exec-event } s (\text{Send-Queuing-Message did eid m})) \sim d \sim (\text{exec-event } t (\text{Send-Queuing-Message did eid m})))$ 
        by (metis Event.simps(51) domain-of-event-def option.sel p1 p2 p3 p4 p5 send-queuing-message-wsc-e)
    }
    then show ?thesis by blast
  qed
}
then show ?thesis
  using dynamic-weakly-step-consistent-e-def by blast
qed

```

### 0.10.3 proving "receive queuing message" satisfying the "weakly step consistent" property

```

lemma receive-queuing-message-wsc-domain-cap-set:
assumes p0: reachable0 s
and p1: reachable0 t
and p2:  $s \sim d \sim t$ 
and p3: interferes did s d
and p4:  $s \sim \text{did} \sim t$ 
and p5: s' = fst (receive-queuing-message s did eid)
and p6: t' = fst (receive-queuing-message t did eid)
shows get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
proof -

```

```

{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: (∀ v. interferes v s d ⟷ interferes v t d)
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: (∀ ep. ep ∈ get-endpoints-of-domain s d
    ⟶ get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id t ep )
    by (meson p2 vpeq1-def)
  have a5: get-endpoints-of-domain s did = get-endpoints-of-domain t did
    by (meson p4 vpeq1-def)
  have a6: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
  proof (cases get-domain-id-from-endpoint s eid = Some did)
    assume b0: get-domain-id-from-endpoint s eid = Some did
    have b1: ((domain-endpoint s) eid) = Some did
      using b0 get-domain-id-from-endpoint-def by auto
    have b2: {x. (domain-endpoint s) x = Some did} = {x. (domain-endpoint t) x = Some did}
      using a5 b1 get-endpoints-of-domain-def by auto
    have b3: ((domain-endpoint t) eid) = Some did
      using b1 b2 by auto
    have b4: get-domain-id-from-endpoint t eid = Some did
      using b3 get-domain-id-from-endpoint-def by auto
    have b5: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id s' d
      using b0 p5 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b6: get-domain-cap-set-from-domain-id t d = get-domain-cap-set-from-domain-id t' d
      using b4 p6 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b7: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
      using b5 b6 a3 by auto
    then show ?thesis by auto
  next
    assume b0: get-domain-id-from-endpoint s eid ≠ Some did
    have b1: ((domain-endpoint s) eid) ≠ Some did
      using b0 get-domain-id-from-endpoint-def by auto
    have b2: {x. (domain-endpoint s) x = Some did} = {x. (domain-endpoint t) x = Some did}
      using a5 b1 get-endpoints-of-domain-def by auto
    have b3: ((domain-endpoint t) eid) ≠ Some did
      using b1 b2 by auto
    have b4: get-domain-id-from-endpoint t eid ≠ Some did
      using b3 get-domain-id-from-endpoint-def by auto
    have b5: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id s' d
      using b0 p5 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b6: get-domain-cap-set-from-domain-id t d = get-domain-cap-set-from-domain-id t' d
      using b4 p6 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b7: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
      using b5 b6 a3 by auto
    then show ?thesis by auto
  qed
}
then show ?thesis by auto
qed

lemma receive-queuing-message-wsc-policy:
assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s ∼ d ∼ t

```



```

and p3: interferes did s d
and p4: s ~ did ~ t
and p5: s' = fst (receive-queuing-message s did eid)
and p6: t' = fst (receive-queuing-message t did eid)
shows (∀ v. interferes v s' d ↔ interferes v t' d)
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: (∀ v. interferes v s d ↔ interferes v t d)
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: (∀ ep. ep ∈ get-endpoints-of-domain s d
    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id t ep)
    by (meson p2 vpeq1-def)
  have a5: get-endpoints-of-domain s did = get-endpoints-of-domain t did
    by (meson p4 vpeq1-def)
  have a6: (∀ v. interferes v s' d ↔ interferes v t' d)
  proof (cases get-domain-id-from-endpoint s eid = Some did)
    assume b0: get-domain-id-from-endpoint s eid = Some did
    have b1: ((domain-endpoint s) eid) = Some did
      using b0 get-domain-id-from-endpoint-def by auto
    have b2: {x. (domain-endpoint s) x = Some did} = {x. (domain-endpoint t) x = Some did}
      using a5 b1 get-endpoints-of-domain-def by auto
    have b3: ((domain-endpoint t) eid) = Some did
      using b1 b2 by auto
    have b4: get-domain-id-from-endpoint t eid = Some did
      using b3 get-domain-id-from-endpoint-def by auto
    have b5: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id s' d
      using b0 p5 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b6: get-domain-cap-set-from-domain-id t d = get-domain-cap-set-from-domain-id t' d
      using b4 p6 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b7: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
      using b5 b6 a3 by auto
    have b8: ∀ v. get-domain-cap-set-from-domain-id s v
      = get-domain-cap-set-from-domain-id s' v
      using b0 p5 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b9: (∀ v. interferes v s d ↔ interferes v s' d)
      using b5 b8 interferes-def by auto
    have b10: ∀ v. get-domain-cap-set-from-domain-id t v
      = get-domain-cap-set-from-domain-id t' v
      using b4 p6 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
    have b11: (∀ v. interferes v t d ↔ interferes v t' d)
      using b6 b10 interferes-def by auto
    have b12: (∀ v. interferes v s' d ↔ interferes v t' d)
      using b9 b11 a1 by auto
    then show ?thesis by auto
  next
    assume b0: get-domain-id-from-endpoint s eid ≠ Some did
    have b1: ((domain-endpoint s) eid) ≠ Some did
      using b0 get-domain-id-from-endpoint-def by auto
    have b2: {x. (domain-endpoint s) x = Some did} = {x. (domain-endpoint t) x = Some did}
      using a5 b1 get-endpoints-of-domain-def by auto
    have b3: ((domain-endpoint t) eid) ≠ Some did
      using b1 b2 by auto

```

```

have b4: get-domain-id-from-endpoint t eid  $\neq$  Some did
  using b3 get-domain-id-from-endpoint-def by auto
have b5: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id s' d
  using b0 p5 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
have b6: get-domain-cap-set-from-domain-id t d = get-domain-cap-set-from-domain-id t' d
  using b4 p6 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
have b7: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
  using b5 b6 a3 by auto
have b8:  $\forall v. \text{get-domain-cap-set-from-domain-id } s \ v$ 
  =  $\text{get-domain-cap-set-from-domain-id } s' \ v$ 
  using b0 p5 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
have b9:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ s' \ d)$ 
  using b5 b8 interferes-def by auto
have b10:  $\forall v. \text{get-domain-cap-set-from-domain-id } t \ v$ 
  =  $\text{get-domain-cap-set-from-domain-id } t' \ v$ 
  using b4 p6 get-domain-cap-set-from-domain-id-def receive-queuing-message-def by auto
have b11:  $(\forall v. \text{interferes } v \ t \ d \longleftrightarrow \text{interferes } v \ t' \ d)$ 
  using b6 b10 interferes-def by auto
have b12:  $(\forall v. \text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d)$ 
  using b9 b11 a1 by auto
then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma receive-queuing-message-wsc-dom-eps:

assumes p0: reachable0 s

and p1: reachable0 t

and p2:  $s \sim d \sim t$

and p3: interferes did s d

and p4:  $s \sim \text{did} \sim t$

and p5:  $s' = \text{fst } (\text{receive-queuing-message } s \ \text{did} \ \text{eid})$

and p6:  $t' = \text{fst } (\text{receive-queuing-message } t \ \text{did} \ \text{eid})$

shows get-endpoints-of-domain s' d = get-endpoints-of-domain t' d

proof -

```

{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ t \ d)$ 
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } t \ ep)$ 
    by (meson p2 vpeq1-def)
  have a5: get-endpoints-of-domain s did = get-endpoints-of-domain t did
    by (meson p4 vpeq1-def)
  have a6: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
  proof (cases get-domain-id-from-endpoint s eid = Some did)
    assume b0: get-domain-id-from-endpoint s eid = Some did
    have b1:  $((\text{domain-endpoint } s) \ \text{eid}) = \text{Some did}$ 
      using b0 get-domain-id-from-endpoint-def by auto
    have b2:  $\{x. (\text{domain-endpoint } s) \ x = \text{Some did}\} = \{x. (\text{domain-endpoint } t) \ x = \text{Some did}\}$ 
      using a5 b1 get-endpoints-of-domain-def by auto
    have b3:  $((\text{domain-endpoint } t) \ \text{eid}) = \text{Some did}$ 

```

```

    using b1 b2 by auto
  have b4: get-domain-id-from-endpoint t eid = Some did
    using b3 get-domain-id-from-endpoint-def by auto
  have b5: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
    using b0 p5 get-endpoints-of-domain-def receive-queuing-message-def by auto
  have b6: get-endpoints-of-domain t d = get-endpoints-of-domain t' d
    using b4 p6 get-endpoints-of-domain-def receive-queuing-message-def by auto
  have b7: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
    using b5 b6 a0 by auto
  then show ?thesis by auto
next
assume b0: get-domain-id-from-endpoint s eid ≠ Some did
have b1: ((domain-endpoint s) eid) ≠ Some did
  using b0 get-domain-id-from-endpoint-def by auto
have b2: {x. (domain-endpoint s) x = Some did} = {x. (domain-endpoint t) x = Some did}
  using a5 b1 get-endpoints-of-domain-def by auto
have b3: ((domain-endpoint t) eid) ≠ Some did
  using b1 b2 by auto
have b4: get-domain-id-from-endpoint t eid ≠ Some did
  using b3 get-domain-id-from-endpoint-def by auto
have b5: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
  using b0 p5 get-endpoints-of-domain-def receive-queuing-message-def by auto
have b6: get-endpoints-of-domain t d = get-endpoints-of-domain t' d
  using b4 p6 get-endpoints-of-domain-def receive-queuing-message-def by auto
have b7: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
  using b5 b6 a0 by auto
then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma receive-queuing-message-wsc-ep-msgs:

assumes p0: reachable0 s

and p1: reachable0 t

and p2:  $s \sim d \sim t$

and p3: interferes did s d

and p4:  $s \sim did \sim t$

and p5:  $s' = fst (receive-queuing-message s did eid)$

and p6:  $t' = fst (receive-queuing-message t did eid)$

shows  $(\forall ep. ep \in get\_endpoints\_of\_domain s' d$

$\longrightarrow get\_msg\_set\_from\_endpoint\_id s' ep = get\_msg\_set\_from\_endpoint\_id t' ep)$

proof –

{

have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d

by (meson p2 vpeq1-def)

have a1:  $(\forall v. interferes v s d \longleftrightarrow interferes v t d)$

by (meson p2 vpeq1-def)

have a2: interferes did t d

using p3 a1 by auto

have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d

by (meson p2 vpeq1-def)

have a4:  $(\forall ep. ep \in get\_endpoints\_of\_domain s d$

$\longrightarrow get\_msg\_set\_from\_endpoint\_id s ep = get\_msg\_set\_from\_endpoint\_id t ep)$

by (meson p2 vpeq1-def)

have a5: get-endpoints-of-domain s did = get-endpoints-of-domain t did

by (meson p4 vpeq1-def)

have a6:  $(\forall ep. ep \in get\_endpoints\_of\_domain s did$

```

    → get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id t ep )
  by (meson p4 vpeq1-def)
have (∀ ep. ep ∈ get-endpoints-of-domain s' d
    → get-msg-set-from-endpoint-id s' ep = get-msg-set-from-endpoint-id t' ep )
proof (cases get-domain-id-from-endpoint s eid = Some did)
  assume b0: get-domain-id-from-endpoint s eid = Some did
  have b1: ((domain-endpoint s) eid) = Some did
    using b0 get-domain-id-from-endpoint-def by auto
  have b2: {x. (domain-endpoint s) x = Some did} = {x. (domain-endpoint t) x = Some did}
    using a5 b1 get-endpoints-of-domain-def by auto
  have b3: ((domain-endpoint t) eid) = Some did
    using b1 b2 by auto
  have b4: get-domain-id-from-endpoint t eid = Some did
    using b3 get-domain-id-from-endpoint-def by auto
  have (∀ ep. ep ∈ get-endpoints-of-domain s' d
    → get-msg-set-from-endpoint-id s' ep = get-msg-set-from-endpoint-id t' ep )
  proof (cases d ≠ did)
    assume c0: d ≠ did
    have c1: ∀ e. e ≠ eid
      → ((e-msgs s) e) = ((e-msgs s') e)
      using b0 p5 receive-queuing-message-def by auto
    have c2: get-domain-id-from-endpoint s eid ≠ Some d
      using c0 b0 by auto
    have c3: ((domain-endpoint s) eid) ≠ Some d
      using c2 get-domain-id-from-endpoint-def by auto
    have c4: eid ∉ {x. (domain-endpoint s) x = Some d}
      using c3 by auto
    have c5: eid ∉ get-endpoints-of-domain s d
      using c4 get-endpoints-of-domain-def by auto
    have c6: ∀ e. e ∈ get-endpoints-of-domain s d
      → ((e-msgs s) e) = ((e-msgs s') e)
      using c1 c5 by auto
    have c7: ∀ e. e ∈ get-endpoints-of-domain s d
      → get-msg-set-from-endpoint-id s e
        = get-msg-set-from-endpoint-id s' e
      using c6 get-msg-set-from-endpoint-id-def by auto
    have c8: ∀ e. e ≠ eid
      → ((e-msgs t) e) = ((e-msgs t') e)
      using b4 p6 receive-queuing-message-def by auto
    have c9: get-domain-id-from-endpoint t eid ≠ Some d
      using c0 b4 by auto
    have c10: ((domain-endpoint t) eid) ≠ Some d
      using c9 get-domain-id-from-endpoint-def by auto
    have c11: eid ∉ {x. (domain-endpoint t) x = Some d}
      using c10 by auto
    have c12: eid ∉ get-endpoints-of-domain t d
      using c11 get-endpoints-of-domain-def by auto
    have c13: ∀ e. e ∈ get-endpoints-of-domain t d
      → ((e-msgs t) e) = ((e-msgs t') e)
      using c8 c12 by auto
    have c14: ∀ e. e ∈ get-endpoints-of-domain t d
      → get-msg-set-from-endpoint-id t e
        = get-msg-set-from-endpoint-id t' e
      using c13 get-msg-set-from-endpoint-id-def by auto
    have c15: ∀ e. e ∈ get-endpoints-of-domain s d
      → get-msg-set-from-endpoint-id t e
        = get-msg-set-from-endpoint-id t' e
      using c14 a0 by auto

```

```

have c16:  $\forall e. e \in \text{get-endpoints-of-domain } s \ d$ 
   $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ e$ 
   $= \text{get-msg-set-from-endpoint-id } t' \ e$ 
  using c15 c7 a4 by auto
have c17:  $\text{get-endpoints-of-domain } s \ d = \text{get-endpoints-of-domain } s' \ d$ 
  using p0 p5 receive-queuing-message-notchg-dom-eps by auto
have c18:  $\forall e. e \in \text{get-endpoints-of-domain } s' \ d$ 
   $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ e$ 
   $= \text{get-msg-set-from-endpoint-id } t' \ e$ 
  using c16 c17 by auto
then show ?thesis by auto
next
assume c0:  $\neg d \neq \text{did}$ 
have c1:  $\forall e. e \neq \text{eid}$ 
   $\longrightarrow ((e\text{-msgs } s) \ e) = ((e\text{-msgs } s') \ e)$ 
  using b0 p5 receive-queuing-message-def by auto
have c2:  $\forall e. e \neq \text{eid}$ 
   $\longrightarrow ((e\text{-msgs } t) \ e) = ((e\text{-msgs } t') \ e)$ 
  using b4 p6 receive-queuing-message-def by auto
have c3:  $\forall e. e \neq \text{eid} \wedge e \in \text{get-endpoints-of-domain } s \ d$ 
   $\longrightarrow ((e\text{-msgs } s') \ e) = ((e\text{-msgs } t') \ e)$ 
  using c1 c2 a4 get-msg-set-from-endpoint-id-def by auto
let ?msg-set-s =  $\text{get-msg-set-from-endpoint-id } s \ \text{eid}$ 
let ?m-s =  $\text{SOME } x. x \in ?\text{msg-set-s}$ 
let ?new-msg-set-s =  $?\text{msg-set-s} - \{?\text{m-s}\}$ 
let ?msg-set-t =  $\text{get-msg-set-from-endpoint-id } t \ \text{eid}$ 
let ?m-t =  $\text{SOME } x. x \in ?\text{msg-set-t}$ 
let ?new-msg-set-t =  $?\text{msg-set-t} - \{?\text{m-t}\}$ 
have c4:  $((e\text{-msgs } s') \ \text{eid}) = ?\text{new-msg-set-s}$ 
  using b0 p5 receive-queuing-message-def by auto
have c5:  $\text{eid} \in \text{get-endpoints-of-domain } s \ \text{did}$ 
  using b1 get-endpoints-of-domain-def by auto
have c6:  $\text{get-msg-set-from-endpoint-id } s \ \text{eid}$ 
   $= \text{get-msg-set-from-endpoint-id } t \ \text{eid}$ 
  using a6 c5 by auto
have c7:  $?\text{msg-set-s} = ?\text{msg-set-t}$ 
  using c6 by auto
have c8:  $?\text{m-s} = ?\text{m-t}$ 
  using c7 by auto
have c9:  $?\text{new-msg-set-s} = ?\text{new-msg-set-t}$ 
  using c7 c8 by auto
have c10:  $((e\text{-msgs } t') \ \text{eid}) = ?\text{new-msg-set-t}$ 
  using b4 p6 receive-queuing-message-def by auto
have c11:  $((e\text{-msgs } s') \ \text{eid}) = ((e\text{-msgs } t') \ \text{eid})$ 
  using c4 c9 c10 by auto
have c12:  $\text{get-msg-set-from-endpoint-id } s' \ \text{eid}$ 
   $= \text{get-msg-set-from-endpoint-id } t' \ \text{eid}$ 
  using c11 get-msg-set-from-endpoint-id-def by auto
have c13:  $\forall e. e \in \text{get-endpoints-of-domain } s \ d$ 
   $\longrightarrow ((e\text{-msgs } s') \ e) = ((e\text{-msgs } t') \ e)$ 
  using c3 c11 by auto
have c14:  $\forall e. e \in \text{get-endpoints-of-domain } s \ d$ 
   $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ e$ 
   $= \text{get-msg-set-from-endpoint-id } t' \ e$ 
  using c13 get-msg-set-from-endpoint-id-def by auto
have c15:  $\text{get-endpoints-of-domain } s \ d = \text{get-endpoints-of-domain } s' \ d$ 
  using p0 p5 receive-queuing-message-notchg-dom-eps by auto
have c16:  $\forall e. e \in \text{get-endpoints-of-domain } s' \ d$ 

```

```

       $\longrightarrow$  get-msg-set-from-endpoint-id  $s'$   $e$ 
      = get-msg-set-from-endpoint-id  $t'$   $e$ 
    using c14 c15 by auto
  then show ?thesis by auto
qed
then show ?thesis by auto
next
  assume b0: get-domain-id-from-endpoint  $s$   $eid \neq$  Some  $did$ 
  have b1: ((domain-endpoint  $s$ )  $eid$ )  $\neq$  Some  $did$ 
    using b0 get-domain-id-from-endpoint-def by auto
  have b2:  $\{x. (\text{domain-endpoint } s) x = \text{Some } did\} = \{x. (\text{domain-endpoint } t) x = \text{Some } did\}$ 
    using a5 b1 get-endpoints-of-domain-def by auto
  have b3: ((domain-endpoint  $t$ )  $eid$ )  $\neq$  Some  $did$ 
    using b1 b2 by auto
  have b4: get-domain-id-from-endpoint  $t$   $eid \neq$  Some  $did$ 
    using b3 get-domain-id-from-endpoint-def by auto
  have b5:  $\forall e. ((e\text{-msgs } s) e) = ((e\text{-msgs } s') e)$ 
    using b0 p5 receive-queuing-message-def by auto
  have b6:  $\forall e. ((e\text{-msgs } t) e) = ((e\text{-msgs } t') e)$ 
    using b4 p6 receive-queuing-message-def by auto
  have b7:  $\forall e. e \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow$  get-msg-set-from-endpoint-id  $s'$   $e$ 
    = get-msg-set-from-endpoint-id  $t'$   $e$ 
    using b5 b6 a4 get-msg-set-from-endpoint-id-def by auto
  have b8: get-endpoints-of-domain  $s \ d = \text{get-endpoints-of-domain } s' \ d$ 
    using p0 p5 receive-queuing-message-notchg-dom-eps by auto
  have b9:  $\forall e. e \in \text{get-endpoints-of-domain } s' \ d$ 
     $\longrightarrow$  get-msg-set-from-endpoint-id  $s'$   $e$ 
    = get-msg-set-from-endpoint-id  $t'$   $e$ 
    using b7 b8 by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma receive-queuing-message-wsc:

assumes p0: reachable0  $s$

and p1: reachable0  $t$

and p2:  $s \sim d \sim t$

and p3: interferes  $did \ s \ d$

and p4:  $s \sim did \sim t$

and p5:  $s' = \text{fst } (\text{receive-queuing-message } s \ did \ eid)$

and p6:  $t' = \text{fst } (\text{receive-queuing-message } t \ did \ eid)$

shows  $s' \sim d \sim t'$

proof –

```

{
  have a0: get-domain-cap-set-from-domain-id  $s'$   $d = \text{get-domain-cap-set-from-domain-id } t' \ d$ 
    using p0 p1 p2 p3 p4 p5 p6 receive-queuing-message-wsc-domain-cap-set by blast
  have a1:  $(\forall v. \text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d)$ 
    using p0 p1 p2 p3 p4 p5 p6 receive-queuing-message-wsc-policy by blast
  have a2: get-endpoints-of-domain  $s' \ d = \text{get-endpoints-of-domain } t' \ d$ 
    using p0 p1 p2 p3 p4 p5 p6 receive-queuing-message-wsc-dom-eps by blast
  have a3:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep)$ 
    using p0 p1 p2 p3 p4 p5 p6 receive-queuing-message-wsc-ep-msgs by blast
  have a4:  $s' \sim d \sim t'$ 
    using a0 a1 a2 a3 vpeq1-def by auto
}

```

```

}
then show ?thesis by auto
qed

```

**lemma** *receive-queuing-message-wsc-e*:

```

assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: a = (Receive-Queuing-Message did eid)
  and p3: s ~ d ~ t
  and p4: interferes (the (domain-of-event a)) s d
  and p5: s ~ did ~ t
  and p6: s' = exec-event s a
  and p7: t' = exec-event t a
shows s' ~ d ~ t'
proof –
{
  have a0: (the (domain-of-event a)) = did
    using p2 domain-of-event-def by auto
  have a1: s' = fst (receive-queuing-message s did eid)
    using p2 p6 exec-event-def by auto
  have a2: t' = fst (receive-queuing-message t did eid)
    using p2 p7 exec-event-def by auto
  have a3: (interferes did s d)
    using p4 a0 by auto
  have a4: s' ~ d ~ t'
    using a1 a2 a3 p0 p1 p3 p5 receive-queuing-message-wsc by blast
}
then show ?thesis by auto
qed

```

**lemma** *receive-queuing-message-dwsc-e*: *dynamic-weakly-step-consistent-e* (Receive-Queuing-Message did eid)

```

proof –
{
  have  $\forall d s t. \text{reachable0 } s \wedge \text{reachable0 } t$ 
     $\wedge (s \sim d \sim t)$ 
     $\wedge (\text{interferes } (\text{the } (\text{domain-of-event } (\text{Receive-Queuing-Message } \text{did } \text{eid})))) s d)$ 
     $\wedge (s \sim (\text{the } (\text{domain-of-event } (\text{Receive-Queuing-Message } \text{did } \text{eid})))) \sim t)$ 
     $\longrightarrow ((\text{exec-event } s (\text{Receive-Queuing-Message } \text{did } \text{eid})) \sim d \sim (\text{exec-event } t (\text{Receive-Queuing-Message } \text{did } \text{eid})))$ 
  proof –
  {
    fix d s t
    assume p1: reachable0 s
    assume p2: reachable0 t
    assume p3: (s ~ d ~ t)
    assume p4: (interferes (the (domain-of-event (Receive-Queuing-Message did eid))) s d)
    assume p5: (s ~ (the (domain-of-event (Receive-Queuing-Message did eid))) ~ t)
    have ((exec-event s (Receive-Queuing-Message did eid)) ~ d ~ (exec-event t (Receive-Queuing-Message did eid)))
      by (metis Event.simps(52) domain-of-event-def option.sel p1 p2 p3 p4 p5 receive-queuing-message-wsc-e)
  }
  then show ?thesis by blast
qed
}
then show ?thesis
using dynamic-weakly-step-consistent-e-def by blast
qed

```

#### 0.10.4 proving "get my endpoints set" satisfying the "weakly step consistent" property

```

lemma get-my-endpoints-set-wsc:
  assumes p0: reachable0 s
    and p1: reachable0 t
    and p2: s ~ d ~ t
    and p3: interferes did s d
    and p4: s ~ did ~ t
    and p5: s' = fst (get-my-endpoints-set s did)
    and p6: t' = fst (get-my-endpoints-set t did)
  shows s' ~ d ~ t'
proof -
{
  have a0: s = s'
    using p5 get-my-endpoints-set-def by auto
  have a1: t = t'
    using p6 get-my-endpoints-set-def by auto
  have a2: s' ~ d ~ t'
    using a0 a1 p2 by auto
}
then show ?thesis by auto
qed

```

```

lemma get-my-endpoints-set-wsc-e:
  assumes p0: reachable0 s
    and p1: reachable0 t
    and p2: a = (Get-My-Endpoints-Set did)
    and p3: s ~ d ~ t
    and p4: interferes (the (domain-of-event a)) s d
    and p5: s ~ did ~ t
    and p6: s' = exec-event s a
    and p7: t' = exec-event t a
  shows s' ~ d ~ t'
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p2 domain-of-event-def by auto
  have a1: s' = fst (get-my-endpoints-set s did)
    using p2 p6 exec-event-def by auto
  have a2: t' = fst (get-my-endpoints-set t did)
    using p2 p7 exec-event-def by auto
  have a3: (interferes did s d)
    using p4 a0 by auto
  have a4: s' ~ d ~ t'
    using a1 a2 a3 p0 p1 p3 p5 get-my-endpoints-set-wsc by blast
}
then show ?thesis by auto
qed

```

```

lemma get-my-endpoints-set-dwsc-e: dynamic-weakly-step-consistent-e (Get-My-Endpoints-Set did)
proof -
{
  have  $\forall d s t. \text{reachable0 } s \wedge \text{reachable0 } t$ 
     $\wedge (s \sim d \sim t)$ 
     $\wedge (\text{interferes } (\text{the } (\text{domain-of-event } (\text{Get-My-Endpoints-Set } \text{did})))) s d)$ 
     $\wedge (s \sim (\text{the } (\text{domain-of-event } (\text{Get-My-Endpoints-Set } \text{did})))) \sim t)$ 
     $\longrightarrow ((\text{exec-event } s (\text{Get-My-Endpoints-Set } \text{did})) \sim d \sim (\text{exec-event } t (\text{Get-My-Endpoints-Set } \text{did})))$ 

```



```

proof –
{
  fix d s t
  assume p1: reachable0 s
  assume p2: reachable0 t
  assume p3: (s  $\sim$  d  $\sim$  t)
  assume p4: (interferes (the (domain-of-event (Get-My-Endpoints-Set did))) s d)
  assume p5: (s  $\sim$  (the (domain-of-event (Get-My-Endpoints-Set did)))  $\sim$  t)
  have ((exec-event s (Get-My-Endpoints-Set did))  $\sim$  d  $\sim$  (exec-event t (Get-My-Endpoints-Set did)))
    by (metis Event.simps(53) domain-of-event-def option.sel p1 p2 p3 p4 p5 get-my-endpoints-set-wsc-e)
}
then show ?thesis by blast
qed
}
then show ?thesis
using dynamic-weakly-step-consistent-e-def by blast
qed

```

### 0.10.5 proving "get caps" satisfying the "weakly step consistent" property

```

lemma get-caps-wsc:
assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s  $\sim$  d  $\sim$  t
  and p3: interferes did s d
  and p4: s  $\sim$  did  $\sim$  t
  and p5: s' = fst (get-caps s did)
  and p6: t' = fst (get-caps t did)
shows s'  $\sim$  d  $\sim$  t'
proof –
{
  have a0: s = s'
    using p5 get-caps-def by auto
  have a1: t = t'
    using p6 get-caps-def by auto
  have a2: s'  $\sim$  d  $\sim$  t'
    using a0 a1 p2 by auto
}
then show ?thesis by auto
qed

```

```

lemma get-caps-wsc-e:
assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: a = (Get-Caps did)
  and p3: s  $\sim$  d  $\sim$  t
  and p4: interferes (the (domain-of-event a)) s d
  and p5: s  $\sim$  did  $\sim$  t
  and p6: s' = exec-event s a
  and p7: t' = exec-event t a
shows s'  $\sim$  d  $\sim$  t'
proof –
{
  have a0: (the (domain-of-event a)) = did
    using p2 domain-of-event-def by auto
  have a1: s' = fst (get-caps s did)
    using p2 p6 exec-event-def by auto

```

```

have a2: t' = fst (get-caps t did)
  using p2 p7 exec-event-def by auto
have a3: (interferes did s d)
  using p4 a0 by auto
have a4: s' ~ d ~ t'
  using a1 a2 a3 p0 p1 p3 p5 get-caps-wsc by blast
}
then show ?thesis by auto
qed

lemma get-caps-dwsc-e: dynamic-weakly-step-consistent-e (Get-Caps did)
proof -
{
  have  $\forall d s t. \text{reachable0 } s \wedge \text{reachable0 } t$ 
     $\wedge (s \sim d \sim t)$ 
     $\wedge (\text{interferes } (\text{the } (\text{domain-of-event } (\text{Get-Caps did}))) s d)$ 
     $\wedge (s \sim (\text{the } (\text{domain-of-event } (\text{Get-Caps did}))) \sim t)$ 
     $\longrightarrow ((\text{exec-event } s (\text{Get-Caps did})) \sim d \sim (\text{exec-event } t (\text{Get-Caps did})))$ 
  proof -
  {
    fix d s t
    assume p1: reachable0 s
    assume p2: reachable0 t
    assume p3: (s ~ d ~ t)
    assume p4: (interferes (the (domain-of-event (Get-Caps did))) s d)
    assume p5: (s ~ (the (domain-of-event (Get-Caps did))) ~ t)
    have ((exec-event s (Get-Caps did)) ~ d ~ (exec-event t (Get-Caps did)))
      by (metis Event.simps(54) domain-of-event-def option.sel p1 p2 p3 p4 p5 get-caps-wsc-e)
  }
  then show ?thesis by blast
qed
}
then show ?thesis
  using dynamic-weakly-step-consistent-e-def by blast
qed

```

### 0.10.6 proving "grant endpoint cap" satisfying the "weakly step consistent" property

```

lemma grant-endpoint-cap-wsc-domain-cap-set:
assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s ~ d ~ t
  and p3: interferes did s d
  and p4: s ~ did ~ t
  and p5: s' = fst (grant-endpoint-cap s did grant-cap dst-cap)
  and p6: t' = fst (grant-endpoint-cap t did grant-cap dst-cap)
shows get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v. \text{interferes } v s d \longleftrightarrow \text{interferes } v t d$ )
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s d$ )

```

```

     $\longrightarrow$  get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id t ep )
  by (meson p2 vpeq1-def)
have a5: get-endpoints-of-domain s did = get-endpoints-of-domain t did
  by (meson p4 vpeq1-def)
have a6: ( $\forall$  ep. ep  $\in$  get-endpoints-of-domain s did
   $\longrightarrow$  get-msg-set-from-endpoint-id s ep = get-msg-set-from-endpoint-id t ep )
  by (meson p4 vpeq1-def)
have a7: get-domain-cap-set-from-domain-id s did = get-domain-cap-set-from-domain-id t did
  by (meson p4 vpeq1-def)
have get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
  proof (cases grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
     $\wedge$  GRANT  $\in$  rights grant-cap
     $\wedge$  target grant-cap  $\neq$  target dst-cap
     $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did)
  assume b0: grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
     $\wedge$  GRANT  $\in$  rights grant-cap
     $\wedge$  target grant-cap  $\neq$  target dst-cap
     $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did
  have b1: grant-cap  $\in$  get-domain-cap-set-from-domain-id t did
     $\wedge$  GRANT  $\in$  rights grant-cap
     $\wedge$  target grant-cap  $\neq$  target dst-cap
     $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id t did
    using a7 b0 by blast
  let ?did-dst = target grant-cap
  let ?cs-dst-s = get-domain-cap-set-from-domain-id s ?did-dst
  let ?cs-dst-t = get-domain-cap-set-from-domain-id t ?did-dst
  have b2:  $\forall v. v \neq ?did-dst$ 
     $\longrightarrow$  (caps s) v = (caps s') v
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b3:  $\forall v. v \neq ?did-dst$ 
     $\longrightarrow$  (caps t) v = (caps t') v
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b4: (caps s') ?did-dst = insert dst-cap ?cs-dst-s
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b5: (caps t') ?did-dst = insert dst-cap ?cs-dst-t
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
    proof (cases d = ?did-dst)
      assume c0: d = ?did-dst
      have c1: ?cs-dst-s = ?cs-dst-t
        using a3 c0 by auto
      have c2: (caps s') ?did-dst = (caps t') ?did-dst
        using b4 b5 c1 by auto
      have c3: get-domain-cap-set-from-domain-id s' d
        = get-domain-cap-set-from-domain-id t' d
        using c2 c0 get-domain-cap-set-from-domain-id-def by auto
      then show ?thesis by auto
    next
      assume c0: d  $\neq$  ?did-dst
      have c1: (caps s) d = (caps t) d
        using a3 get-domain-cap-set-from-domain-id-def by auto
      have c2: (caps s) d = (caps s') d
        using c0 b2 by auto
      have c3: (caps t) d = (caps t') d
        using c0 b3 by auto
      have c4: (caps s') d = (caps t') d
        using c1 c2 c3 by auto
      have c5: get-domain-cap-set-from-domain-id s' d

```

```

      = get-domain-cap-set-from-domain-id t' d
    using c4 get-domain-cap-set-from-domain-id-def by auto
    then show ?thesis by auto
  qed
  then show ?thesis by auto
next
  assume b0:  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did})$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
  have b1:  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did})$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
    using a7 b0 by auto
  have b2:  $s = s'$ 
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b3:  $t = t'$ 
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b4:  $\text{get-domain-cap-set-from-domain-id } s' \text{ d}$ 
    =  $\text{get-domain-cap-set-from-domain-id } s \text{ d}$ 
    using b2 get-domain-cap-set-from-domain-id-def by auto
  have b5:  $\text{get-domain-cap-set-from-domain-id } t' \text{ d}$ 
    =  $\text{get-domain-cap-set-from-domain-id } t \text{ d}$ 
    using b3 get-domain-cap-set-from-domain-id-def by auto
  have b6:  $\text{get-domain-cap-set-from-domain-id } s' \text{ d}$ 
    =  $\text{get-domain-cap-set-from-domain-id } t' \text{ d}$ 
    using b4 b5 a3 by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma grant-endpoint-cap-wsc-policy:

```

assumes p0: reachable0 s
  and p1: reachable0 t
  and p2:  $s \sim d \sim t$ 
  and p3: interferes did s d
  and p4:  $s \sim \text{did} \sim t$ 
  and p5:  $s' = \text{fst } (\text{grant-endpoint-cap } s \text{ did grant-cap dst-cap})$ 
  and p6:  $t' = \text{fst } (\text{grant-endpoint-cap } t \text{ did grant-cap dst-cap})$ 
shows  $(\forall v. \text{interferes } v \text{ s' d} \longleftrightarrow \text{interferes } v \text{ t' d})$ 
proof -
{
  have a0:  $\text{get-endpoints-of-domain } s \text{ d} = \text{get-endpoints-of-domain } t \text{ d}$ 
    by (meson p2 vpeq1-def)
  have a1:  $(\forall v. \text{interferes } v \text{ s d} \longleftrightarrow \text{interferes } v \text{ t d})$ 
    by (meson p2 vpeq1-def)
  have a2:  $\text{interferes did t d}$ 
    using p3 a1 by auto
  have a3:  $\text{get-domain-cap-set-from-domain-id } s \text{ d} = \text{get-domain-cap-set-from-domain-id } t \text{ d}$ 
    by (meson p2 vpeq1-def)
  have a4:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d} \longrightarrow \text{get-msg-set-from-endpoint-id } s \text{ ep} = \text{get-msg-set-from-endpoint-id } t \text{ ep})$ 
    by (meson p2 vpeq1-def)
  have a5:  $\text{get-endpoints-of-domain } s \text{ did} = \text{get-endpoints-of-domain } t \text{ did}$ 
    by (meson p4 vpeq1-def)
}

```

```

have a6: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ did}$ 
   $\longrightarrow \text{get-msg-set-from-endpoint-id } s \text{ ep} = \text{get-msg-set-from-endpoint-id } t \text{ ep}$  )
  by (meson p4 vpeq1-def)
have a7:  $\text{get-domain-cap-set-from-domain-id } s \text{ did} = \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
  by (meson p4 vpeq1-def)
have ( $\forall v. \text{interferes } v \text{ s' d} \longleftrightarrow \text{interferes } v \text{ t' d}$ )
  proof (cases  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ )
  assume b0:  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
  have b1:  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
    using a7 b0 by blast
  let ?did-dst = target grant-cap
  let ?cs-dst-s = get-domain-cap-set-from-domain-id s ?did-dst
  let ?cs-dst-t = get-domain-cap-set-from-domain-id t ?did-dst
  have b2:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow (\text{caps } s) \text{ v} = (\text{caps } s') \text{ v}$ 
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b3:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow (\text{caps } t) \text{ v} = (\text{caps } t') \text{ v}$ 
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b4:  $(\text{caps } s') \text{ ?did-dst} = \text{insert dst-cap ?cs-dst-s}$ 
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b5:  $(\text{caps } t') \text{ ?did-dst} = \text{insert dst-cap ?cs-dst-t}$ 
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  let ?dst-cap-d = target dst-cap
  have b6:  $\text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s' \text{ ?did-dst}$ 
    using b4 get-domain-cap-set-from-domain-id-def by auto
  have b7:  $\text{interferes } ?\text{did-dst } s' \text{ ?dst-cap-d}$ 
    using b6 interferes-def by auto
  have b8:  $\text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } t' \text{ ?did-dst}$ 
    using b5 get-domain-cap-set-from-domain-id-def by auto
  have b9:  $\text{interferes } ?\text{did-dst } t' \text{ ?dst-cap-d}$ 
    using b8 interferes-def by auto
  have b10:  $?\text{did-dst} \neq ?\text{dst-cap-d}$ 
    using b0 by auto
  have b11:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow \text{get-domain-cap-set-from-domain-id } s \text{ v}$ 
     $= \text{get-domain-cap-set-from-domain-id } s' \text{ v}$ 
    using b2 get-domain-cap-set-from-domain-id-def by auto
  have b12:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow \text{interferes } v \text{ s d} = \text{interferes } v \text{ s' d}$ 
    using b11 interferes-def by auto
  have b13:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow \text{get-domain-cap-set-from-domain-id } t \text{ v}$ 
     $= \text{get-domain-cap-set-from-domain-id } t' \text{ v}$ 
    using b3 get-domain-cap-set-from-domain-id-def by auto
  have b14:  $\forall v. v \neq ?\text{did-dst}$ 
     $\longrightarrow \text{interferes } v \text{ t d} = \text{interferes } v \text{ t' d}$ 
    using b13 interferes-def by auto
  have b15:  $\forall v. v \neq ?\text{did-dst}$ 

```

```

     $\longrightarrow \text{interferes } v \ s' \ d = \text{interferes } v \ t' \ d$ 
  using b12 b14 a1 by auto
have (∀ v.  $\text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d$ )
proof (cases d = ?dst-cap-d)
  assume c0: d = ?dst-cap-d
  have c1:  $\text{interferes } ?\text{did-dst } s' \ ?\text{dst-cap-d} = \text{interferes } ?\text{did-dst } t' \ ?\text{dst-cap-d}$ 
  using b7 b9 by auto
  have c2:  $\text{interferes } ?\text{did-dst } s' \ d = \text{interferes } ?\text{did-dst } t' \ d$ 
  using c1 c0 by auto
  have c3: ∀ v. v=?did-dst
     $\longrightarrow \text{interferes } v \ s' \ d = \text{interferes } v \ t' \ d$ 
  using c2 by auto
  have c4: (∀ v.  $\text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d$ )
  using c3 b15 by blast
  then show ?thesis by auto
next
  assume c0: d ≠ ?dst-cap-d
  have c1: get-domain-cap-set-from-domain-id s' ?did-dst
    = insert dst-cap (get-domain-cap-set-from-domain-id s ?did-dst)
  using b4 get-domain-cap-set-from-domain-id-def by auto
  have c2: get-domain-cap-set-from-domain-id t' ?did-dst
    = insert dst-cap (get-domain-cap-set-from-domain-id t ?did-dst)
  using b5 get-domain-cap-set-from-domain-id-def by auto
  have c3:  $\text{interferes } ?\text{did-dst } s' \ d = \text{interferes } ?\text{did-dst } s \ d$ 
  using c1 c0 interferes-def by auto
  have c4:  $\text{interferes } ?\text{did-dst } t' \ d = \text{interferes } ?\text{did-dst } t \ d$ 
  using c2 c0 interferes-def by auto
  have c5:  $\text{interferes } ?\text{did-dst } s' \ d = \text{interferes } ?\text{did-dst } t' \ d$ 
  using c3 c4 a1 by auto
  have c6: ∀ v. v=?did-dst
     $\longrightarrow \text{interferes } v \ s' \ d = \text{interferes } v \ t' \ d$ 
  using c5 by auto
  have c7: (∀ v.  $\text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d$ )
  using c6 b15 by blast
  then show ?thesis by auto
qed
then show ?thesis by auto
next
  assume b0: ¬(grant-cap ∈ get-domain-cap-set-from-domain-id s did)
    ∧ GRANT ∈ rights grant-cap
    ∧ target grant-cap ≠ target dst-cap
    ∧ dst-cap ∈ get-domain-cap-set-from-domain-id s did)
  have b1: ¬(grant-cap ∈ get-domain-cap-set-from-domain-id t did)
    ∧ GRANT ∈ rights grant-cap
    ∧ target grant-cap ≠ target dst-cap
    ∧ dst-cap ∈ get-domain-cap-set-from-domain-id t did)
  using a7 b0 by auto
  have b2: s = s'
  using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b3: t = t'
  using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b4: (∀ v.  $\text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d$ )
  using a1 b2 b3 by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

```

lemma grant-endpoint-cap-wsc-dom-eps:
assumes p0: reachable0 s
  and p1: reachable0 t
  and p2: s ~ d ~ t
  and p3: interferes did s d
  and p4: s ~ did ~ t
  and p5: s' = fst (grant-endpoint-cap s did grant-cap dst-cap)
  and p6: t' = fst (grant-endpoint-cap t did grant-cap dst-cap)
shows get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
proof -
{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v. \text{interferes } v \text{ s d} \longleftrightarrow \text{interferes } v \text{ t d}$ )
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id s d = get-domain-cap-set-from-domain-id t d
    by (meson p2 vpeq1-def)
  have a4: ( $\forall ep. ep \in \text{get-endpoints-of-domain s d} \longrightarrow \text{get-msg-set-from-endpoint-id s ep} = \text{get-msg-set-from-endpoint-id t ep}$ )
    by (meson p2 vpeq1-def)
  have a5: get-endpoints-of-domain s did = get-endpoints-of-domain t did
    by (meson p4 vpeq1-def)
  have a6: ( $\forall ep. ep \in \text{get-endpoints-of-domain s did} \longrightarrow \text{get-msg-set-from-endpoint-id s ep} = \text{get-msg-set-from-endpoint-id t ep}$ )
    by (meson p4 vpeq1-def)
  have a7: get-domain-cap-set-from-domain-id s did = get-domain-cap-set-from-domain-id t did
    by (meson p4 vpeq1-def)
  have get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
    proof (cases grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
       $\wedge$  GRANT  $\in$  rights grant-cap
       $\wedge$  target grant-cap  $\neq$  target dst-cap
       $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did)
    assume b0: grant-cap  $\in$  get-domain-cap-set-from-domain-id s did
       $\wedge$  GRANT  $\in$  rights grant-cap
       $\wedge$  target grant-cap  $\neq$  target dst-cap
       $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did
    have b1: grant-cap  $\in$  get-domain-cap-set-from-domain-id t did
       $\wedge$  GRANT  $\in$  rights grant-cap
       $\wedge$  target grant-cap  $\neq$  target dst-cap
       $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id t did
      using a7 b0 by blast
    have b2: get-endpoints-of-domain s' d = get-endpoints-of-domain s d
      using b0 p5 grant-endpoint-cap-def get-endpoints-of-domain-def by auto
    have b3: get-endpoints-of-domain t' d = get-endpoints-of-domain t d
      using b1 p6 grant-endpoint-cap-def get-endpoints-of-domain-def by auto
    have b4: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
      using b2 b3 a0 by auto
    then show ?thesis by auto
  next
    assume b0:  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id s did}$ 
       $\wedge$  GRANT  $\in$  rights grant-cap
       $\wedge$  target grant-cap  $\neq$  target dst-cap
       $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id s did)
    have b1:  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id t did}$ 
       $\wedge$  GRANT  $\in$  rights grant-cap

```

```

       $\wedge$  target grant-cap  $\neq$  target dst-cap
       $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id t did)
    using a7 b0 by auto
  have b2:  $s = s'$ 
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b3:  $t = t'$ 
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
  have b4: get-endpoints-of-domain  $s' d$ 
    = get-endpoints-of-domain  $s d$ 
    using b2 get-domain-cap-set-from-domain-id-def by auto
  have b5: get-endpoints-of-domain  $t' d$ 
    = get-endpoints-of-domain  $t d$ 
    using b3 get-domain-cap-set-from-domain-id-def by auto
  have b6: get-endpoints-of-domain  $s' d$ 
    = get-endpoints-of-domain  $t' d$ 
    using b4 b5 a0 by auto
  then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

lemma grant-endpoint-cap-wsc-ep-msgs:

assumes p0: reachable0  $s$

and p1: reachable0  $t$

and p2:  $s \sim d \sim t$

and p3: interferes did  $s d$

and p4:  $s \sim did \sim t$

and p5:  $s' = \text{fst } (\text{grant-endpoint-cap } s \text{ did grant-cap dst-cap})$

and p6:  $t' = \text{fst } (\text{grant-endpoint-cap } t \text{ did grant-cap dst-cap})$

shows  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' d$

$\longrightarrow \text{get-msg-set-from-endpoint-id } s' ep = \text{get-msg-set-from-endpoint-id } t' ep )$

proof –

```

{
  have a0: get-endpoints-of-domain  $s d = \text{get-endpoints-of-domain } t d$ 
    by (meson p2 vpeq1-def)
  have a1:  $(\forall v. \text{interferes } v s d \longleftrightarrow \text{interferes } v t d)$ 
    by (meson p2 vpeq1-def)
  have a2: interferes did  $t d$ 
    using p3 a1 by auto
  have a3: get-domain-cap-set-from-domain-id  $s d = \text{get-domain-cap-set-from-domain-id } t d$ 
    by (meson p2 vpeq1-def)
  have a4:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s ep = \text{get-msg-set-from-endpoint-id } t ep )$ 
    by (meson p2 vpeq1-def)
  have a5: get-endpoints-of-domain  $s did = \text{get-endpoints-of-domain } t did$ 
    by (meson p4 vpeq1-def)
  have a6:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s did$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s ep = \text{get-msg-set-from-endpoint-id } t ep )$ 
    by (meson p4 vpeq1-def)
  have a7: get-domain-cap-set-from-domain-id  $s did = \text{get-domain-cap-set-from-domain-id } t did$ 
    by (meson p4 vpeq1-def)
  have  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' ep = \text{get-msg-set-from-endpoint-id } t' ep )$ 
    proof (cases grant-cap  $\in$  get-domain-cap-set-from-domain-id  $s did$ 
       $\wedge$  GRANT  $\in$  rights grant-cap
       $\wedge$  target grant-cap  $\neq$  target dst-cap
       $\wedge$  dst-cap  $\in$  get-domain-cap-set-from-domain-id  $s did$ )

```



```

assume b0:  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
have b1:  $\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
    using a7 b0 by blast
have b2:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \text{ ep} = \text{get-msg-set-from-endpoint-id } s \text{ ep}$ 
    using b0 p5 grant-endpoint-cap-def get-msg-set-from-endpoint-id-def by auto
have b3:  $(\forall ep. ep \in \text{get-endpoints-of-domain } t \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } t' \text{ ep} = \text{get-msg-set-from-endpoint-id } t \text{ ep}$ 
    using b1 p6 grant-endpoint-cap-def get-msg-set-from-endpoint-id-def by auto
have b4:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } t' \text{ ep} = \text{get-msg-set-from-endpoint-id } t \text{ ep}$ 
    using b3 a0 by auto
have b5:  $\text{get-endpoints-of-domain } s \text{ d} = \text{get-endpoints-of-domain } s' \text{ d}$ 
    using p0 p5 grant-endpoint-cap-notchg-dom-eps by auto
have b6:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \text{ ep} = \text{get-msg-set-from-endpoint-id } t' \text{ ep}$ 
    using b2 b4 a4 by auto
have b7:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \text{ ep} = \text{get-msg-set-from-endpoint-id } t' \text{ ep}$ 
    using b5 b6 by auto
then show ?thesis by auto
next
assume b0:  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did})$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
have b1:  $\neg(\text{grant-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did})$ 
     $\wedge \text{GRANT} \in \text{rights grant-cap}$ 
     $\wedge \text{target grant-cap} \neq \text{target dst-cap}$ 
     $\wedge \text{dst-cap} \in \text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
    using a7 b0 by auto
have b2:  $s = s'$ 
    using b0 p5 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
have b3:  $t = t'$ 
    using b1 p6 grant-endpoint-cap-def get-domain-cap-set-from-domain-id-def by auto
have b4:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \text{ ep} = \text{get-msg-set-from-endpoint-id } s \text{ ep}$ 
    using b2 get-domain-cap-set-from-domain-id-def by auto
have b5:  $(\forall ep. ep \in \text{get-endpoints-of-domain } t \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } t' \text{ ep} = \text{get-msg-set-from-endpoint-id } t \text{ ep}$ 
    using b3 get-domain-cap-set-from-domain-id-def by auto
have b6:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } t' \text{ ep} = \text{get-msg-set-from-endpoint-id } t \text{ ep}$ 
    using b5 a0 by auto
have b7:  $\text{get-endpoints-of-domain } s \text{ d} = \text{get-endpoints-of-domain } s' \text{ d}$ 
    using p0 p5 grant-endpoint-cap-notchg-dom-eps by auto
have b8:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \text{ ep} = \text{get-msg-set-from-endpoint-id } t' \text{ ep}$ 
    using b4 b6 a4 by auto
have b9:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' \text{ d})$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \text{ ep} = \text{get-msg-set-from-endpoint-id } t' \text{ ep}$ 
    using b7 b8 by auto

```

```

    then show ?thesis by auto
  qed
}
then show ?thesis by auto
qed

```

lemma *grant-endpoint-cap-wsc*:

```

assumes p0: reachable0 s
  and p1: reachable0 t
  and p2:  $s \sim d \sim t$ 
  and p3: interferes did s d
  and p4:  $s \sim \text{did} \sim t$ 
  and p5:  $s' = \text{fst } (\text{grant-endpoint-cap } s \text{ did grant-cap dst-cap})$ 
  and p6:  $t' = \text{fst } (\text{grant-endpoint-cap } t \text{ did grant-cap dst-cap})$ 
shows  $s' \sim d \sim t'$ 
proof -
{
  have a0: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
    using p0 p1 p2 p3 p4 p5 p6 grant-endpoint-cap-wsc-domain-cap-set by blast
  have a1:  $(\forall v. \text{interferes } v \text{ s' d} \longleftrightarrow \text{interferes } v \text{ t' d})$ 
    using p0 p1 p2 p3 p4 p5 p6 grant-endpoint-cap-wsc-policy by blast
  have a2: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
    using p0 p1 p2 p3 p4 p5 p6 grant-endpoint-cap-wsc-dom-eps by blast
  have a3:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s' d \longrightarrow \text{get-msg-set-from-endpoint-id } s' ep = \text{get-msg-set-from-endpoint-id } t' ep)$ 
    using p0 p1 p2 p3 p4 p5 p6 grant-endpoint-cap-wsc-ep-msgs by blast
  have a4:  $s' \sim d \sim t'$ 
    using a0 a1 a2 a3 vpeq1-def by auto
}
then show ?thesis by auto
qed

```

lemma *grant-endpoint-cap-wsc-e*:

```

assumes p0: reachable0 s
  and p1: reachable0 t
  and p2:  $a = (\text{Grant-Endpoint-Cap } \text{did } \text{grant-cap } \text{dst-cap})$ 
  and p3:  $s \sim d \sim t$ 
  and p4: interferes (the (domain-of-event a)) s d
  and p5:  $s \sim \text{did} \sim t$ 
  and p6:  $s' = \text{exec-event } s \ a$ 
  and p7:  $t' = \text{exec-event } t \ a$ 
shows  $s' \sim d \sim t'$ 
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p2 domain-of-event-def by auto
  have a1:  $s' = \text{fst } (\text{grant-endpoint-cap } s \text{ did grant-cap dst-cap})$ 
    using p2 p6 exec-event-def by auto
  have a2:  $t' = \text{fst } (\text{grant-endpoint-cap } t \text{ did grant-cap dst-cap})$ 
    using p2 p7 exec-event-def by auto
  have a3: (interferes did s d)
    using p4 a0 by auto
  have a4:  $s' \sim d \sim t'$ 
    using a1 a2 a3 p0 p1 p3 p5 grant-endpoint-cap-wsc by blast
}
then show ?thesis by auto
qed

```

```

lemma grant-endpoint-cap-dwsc-e: dynamic-weakly-step-consistent-e (Grant-Endpoint-Cap did grant-cap dst-cap)
proof -
{
  have  $\forall d\ s\ t. \text{reachable0}\ s \wedge \text{reachable0}\ t$ 
     $\wedge (s \sim d \sim t)$ 
     $\wedge (\text{interferes}\ (\text{the}\ (\text{domain-of-event}\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap})))\ s\ d)$ 
     $\wedge (s \sim (\text{the}\ (\text{domain-of-event}\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap}))) \sim t)$ 
     $\longrightarrow ((\text{exec-event}\ s\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap})) \sim d \sim (\text{exec-event}\ t\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap})))$ 
  proof -
  {
    fix  $d\ s\ t$ 
    assume  $p1: \text{reachable0}\ s$ 
    assume  $p2: \text{reachable0}\ t$ 
    assume  $p3: (s \sim d \sim t)$ 
    assume  $p4: (\text{interferes}\ (\text{the}\ (\text{domain-of-event}\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap})))\ s\ d)$ 
    assume  $p5: (s \sim (\text{the}\ (\text{domain-of-event}\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap}))) \sim t)$ 
    have  $((\text{exec-event}\ s\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap})) \sim d \sim (\text{exec-event}\ t\ (\text{Grant-Endpoint-Cap}\ \text{did}\ \text{grant-cap}\ \text{dst-cap})))$ 
      by (metis Event.simps(55) domain-of-event-def option.sel p1 p2 p3 p4 p5 grant-endpoint-cap-wsc-e)
  }
  then show ?thesis by blast
qed
}
then show ?thesis
using dynamic-weakly-step-consistent-e-def by blast
qed

```

### 0.10.7 proving "remove cap right" satisfying the "weakly step consistent" property

```

lemma remove-cap-right-wsc-domain-cap-set:
assumes  $p0: \text{reachable0}\ s$ 
and  $p1: \text{reachable0}\ t$ 
and  $p2: s \sim d \sim t$ 
and  $p3: \text{interferes}\ \text{did}\ s\ d$ 
and  $p4: s \sim \text{did} \sim t$ 
and  $p5: s' = \text{fst}\ (\text{remove-cap-right}\ s\ \text{did}\ \text{rm-cap}\ \text{right-to-rm})$ 
and  $p6: t' = \text{fst}\ (\text{remove-cap-right}\ t\ \text{did}\ \text{rm-cap}\ \text{right-to-rm})$ 
shows  $\text{get-domain-cap-set-from-domain-id}\ s'\ d$ 
   $= \text{get-domain-cap-set-from-domain-id}\ t'\ d$ 
proof -
{
  have  $a0: \text{get-endpoints-of-domain}\ s\ d = \text{get-endpoints-of-domain}\ t\ d$ 
    by (meson p2 vpeq1-def)
  have  $a1: (\forall v. \text{interferes}\ v\ s\ d \longleftrightarrow \text{interferes}\ v\ t\ d)$ 
    by (meson p2 vpeq1-def)
  have  $a2: \text{interferes}\ \text{did}\ t\ d$ 
    using p3 a1 by auto
  have  $a3: \text{get-domain-cap-set-from-domain-id}\ s\ d$ 
     $= \text{get-domain-cap-set-from-domain-id}\ t\ d$ 
    by (meson p2 vpeq1-def)
  have  $a4: (\forall ep. ep \in \text{get-endpoints-of-domain}\ s\ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id}\ s\ ep = \text{get-msg-set-from-endpoint-id}\ t\ ep)$ 
    by (meson p2 vpeq1-def)
  have  $a5: \text{get-endpoints-of-domain}\ s\ \text{did} = \text{get-endpoints-of-domain}\ t\ \text{did}$ 
    by (meson p4 vpeq1-def)
  have  $a6: (\forall ep. ep \in \text{get-endpoints-of-domain}\ s\ \text{did}$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id}\ s\ ep = \text{get-msg-set-from-endpoint-id}\ t\ ep)$ 

```

```

by (meson p4 vpeq1-def)
have a7: get-domain-cap-set-from-domain-id s did
  = get-domain-cap-set-from-domain-id t did
by (meson p4 vpeq1-def)
have get-domain-cap-set-from-domain-id s' d
  = get-domain-cap-set-from-domain-id t' d
proof (cases rm-cap ∈ get-domain-cap-set-from-domain-id s did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap)
assume b0: rm-cap ∈ get-domain-cap-set-from-domain-id s did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap
have b1: rm-cap ∈ get-domain-cap-set-from-domain-id t did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap
using a7 b0 by auto
have get-domain-cap-set-from-domain-id s' d
  = get-domain-cap-set-from-domain-id t' d
proof (cases REMOVE = right-to-rm
  ∧ {REMOVE} = rights rm-cap)
assume c0: REMOVE = right-to-rm
  ∧ {REMOVE} = rights rm-cap
let ?cs-dst-s = get-domain-cap-set-from-domain-id s did
let ?cs-rest-s = {c. c ∈ ?cs-dst-s ∧ c ≠ rm-cap}
let ?cs-dst-t = get-domain-cap-set-from-domain-id t did
let ?cs-rest-t = {c. c ∈ ?cs-dst-t ∧ c ≠ rm-cap}
have c1: ((caps s') did) = ?cs-rest-s
  using b0 c0 p5 remove-cap-right-def by auto
have c2: ((caps t') did) = ?cs-rest-t
  using b1 c0 p6 remove-cap-right-def by auto
have c3: ?cs-rest-s = ?cs-rest-t
  using a7 by auto
have c4: ((caps s') did) = ((caps t') did)
  using c1 c2 c3 by auto
have c5: ∀ v. v ≠ did
  → ((caps s') v) = ((caps s) v)
  using b0 c0 p5 remove-cap-right-def by auto
have c6: ∀ v. v ≠ did
  → ((caps t') v) = ((caps t) v)
  using b1 c0 p6 remove-cap-right-def by auto
have c7: ∀ v. v ≠ did ∧ v = d
  → ((caps s') v) = ((caps t') v)
  using c5 c6 a3 get-domain-cap-set-from-domain-id-def by auto
have c8: d ≠ did
  → ((caps s') d) = ((caps t') d)
  using c7 by auto
have c9: ((caps s') d) = ((caps t') d)
  using c4 c8 by auto
have c10: get-domain-cap-set-from-domain-id s' d
  = get-domain-cap-set-from-domain-id t' d
  using c9 get-domain-cap-set-from-domain-id-def by auto
then show ?thesis by auto
next
assume c0: ¬(REMOVE = right-to-rm
  ∧ {REMOVE} = rights rm-cap)
let ?cs-dst-s = get-domain-cap-set-from-domain-id s did
let ?cs-rest-s = {c. c ∈ ?cs-dst-s ∧ c ≠ rm-cap}
let ?cs-dst-t = get-domain-cap-set-from-domain-id t did

```

```

let ?cs-rest-t = {c. c ∈ ?cs-dst-t ∧ c ≠ rm-cap}
let ?new-cap = () target = target rm-cap,
               rights = (rights rm-cap) - {right-to-rm}
have c1: ((caps s') did) = (insert ?new-cap ?cs-rest-s)
  using b0 c0 p5 remove-cap-right-def by auto
have c2: ((caps t') did) = (insert ?new-cap ?cs-rest-t)
  using b1 c0 p6 remove-cap-right-def by auto
have c3: ?cs-rest-s = ?cs-rest-t
  using a7 by auto
have c4: ((caps s') did) = ((caps t') did)
  using c1 c2 c3 by auto
have c5: ∀ v. v ≠ did
  → ((caps s') v) = ((caps s) v)
  using b0 c0 p5 remove-cap-right-def by auto
have c6: ∀ v. v ≠ did
  → ((caps t') v) = ((caps t) v)
  using b1 c0 p6 remove-cap-right-def by auto
have c7: ∀ v. v ≠ did ∧ v = d
  → ((caps s') v) = ((caps t') v)
  using c5 c6 a3 get-domain-cap-set-from-domain-id-def by auto
have c8: d ≠ did
  → ((caps s') d) = ((caps t') d)
  using c7 by auto
have c9: ((caps s') d) = ((caps t') d)
  using c4 c8 by auto
have c10: get-domain-cap-set-from-domain-id s' d
  = get-domain-cap-set-from-domain-id t' d
  using c9 get-domain-cap-set-from-domain-id-def by auto
then show ?thesis by auto
qed
then show ?thesis by auto
next
assume b0: ¬(rm-cap ∈ get-domain-cap-set-from-domain-id s did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap)
have b1: ¬(rm-cap ∈ get-domain-cap-set-from-domain-id t did
  ∧ REMOVE ∈ rights rm-cap
  ∧ right-to-rm ∈ rights rm-cap)
  using b0 a7 by auto
have b2: s' = s
  using b0 p5 remove-cap-right-def by auto
have b3: get-domain-cap-set-from-domain-id s d
  = get-domain-cap-set-from-domain-id s' d
  using b2 get-domain-cap-set-from-domain-id-def by auto
have b4: t' = t
  using b1 p6 remove-cap-right-def by auto
have b5: get-domain-cap-set-from-domain-id t d
  = get-domain-cap-set-from-domain-id t' d
  using b4 get-domain-cap-set-from-domain-id-def by auto
have b6: get-domain-cap-set-from-domain-id s' d
  = get-domain-cap-set-from-domain-id t' d
  using a3 b3 b5 by auto
then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

```

**lemma** *remove-cap-right-wsc-policy*:

**assumes**  $p0$ : *reachable0*  $s$

**and**  $p1$ : *reachable0*  $t$

**and**  $p2$ :  $s \sim d \sim t$

**and**  $p3$ : *interferes*  $did$   $s$   $d$

**and**  $p4$ :  $s \sim did \sim t$

**and**  $p5$ :  $s' = fst \ (remove\_cap\_right \ s \ did \ rm\_cap \ right\_to\_rm)$

**and**  $p6$ :  $t' = fst \ (remove\_cap\_right \ t \ did \ rm\_cap \ right\_to\_rm)$

**shows**  $(\forall v. \text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d)$

**proof** –

```

{
  have a0: get-endpoints-of-domain  $s$   $d$  = get-endpoints-of-domain  $t$   $d$ 
    by (meson  $p2$  vpeq1-def)
  have a1:  $(\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ t \ d)$ 
    by (meson  $p2$  vpeq1-def)
  have a2: interferes  $did$   $t$   $d$ 
    using  $p3$   $a1$  by auto
  have a3: get-domain-cap-set-from-domain-id  $s$   $d$ 
    = get-domain-cap-set-from-domain-id  $t$   $d$ 
    by (meson  $p2$  vpeq1-def)
  have a4:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \ d \longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } t \ ep)$ 
    by (meson  $p2$  vpeq1-def)
  have a5: get-endpoints-of-domain  $s$   $did$  = get-endpoints-of-domain  $t$   $did$ 
    by (meson  $p4$  vpeq1-def)
  have a6:  $(\forall ep. ep \in \text{get-endpoints-of-domain } s \ did \longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } t \ ep)$ 
    by (meson  $p4$  vpeq1-def)
  have a7: get-domain-cap-set-from-domain-id  $s$   $did$ 
    = get-domain-cap-set-from-domain-id  $t$   $did$ 
    by (meson  $p4$  vpeq1-def)
  have  $(\forall v. \text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d)$ 
    proof (cases  $rm\_cap \in \text{get-domain-cap-set-from-domain-id } s \ did$ 
       $\wedge REMOVE \in \text{rights } rm\_cap$ 
       $\wedge \text{right-to-rm} \in \text{rights } rm\_cap$ )
      assume  $b0$ :  $rm\_cap \in \text{get-domain-cap-set-from-domain-id } s \ did$ 
         $\wedge REMOVE \in \text{rights } rm\_cap$ 
         $\wedge \text{right-to-rm} \in \text{rights } rm\_cap$ 
      have  $b1$ :  $rm\_cap \in \text{get-domain-cap-set-from-domain-id } t \ did$ 
         $\wedge REMOVE \in \text{rights } rm\_cap$ 
         $\wedge \text{right-to-rm} \in \text{rights } rm\_cap$ 
      using  $a7$   $b0$  by auto
    have  $(\forall v. \text{interferes } v \ s' \ d \longleftrightarrow \text{interferes } v \ t' \ d)$ 
      proof (cases  $REMOVE = \text{right-to-rm}$ 
         $\wedge \{REMOVE\} = \text{rights } rm\_cap$ )
        assume  $c0$ :  $REMOVE = \text{right-to-rm}$ 
           $\wedge \{REMOVE\} = \text{rights } rm\_cap$ 
        let  $?cs\_dst\_s = \text{get-domain-cap-set-from-domain-id } s \ did$ 
        let  $?cs\_rest\_s = \{c. c \in ?cs\_dst\_s \wedge c \neq rm\_cap\}$ 
        let  $?cs\_dst\_t = \text{get-domain-cap-set-from-domain-id } t \ did$ 
        let  $?cs\_rest\_t = \{c. c \in ?cs\_dst\_t \wedge c \neq rm\_cap\}$ 
        have  $c1$ :  $((\text{caps } s') \ did) = ?cs\_rest\_s$ 
          using  $b0$   $c0$   $p5$  remove-cap-right-def by auto
        have  $c2$ :  $((\text{caps } t') \ did) = ?cs\_rest\_t$ 
          using  $b1$   $c0$   $p6$  remove-cap-right-def by auto
        have  $c3$ :  $?cs\_rest\_s = ?cs\_rest\_t$ 
          using  $a7$  by auto
        have  $c4$ :  $((\text{caps } s') \ did) = ((\text{caps } t') \ did)$ 

```

```

    using c1 c2 c3 by auto
have c5:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow ((\text{caps } s') v) = ((\text{caps } s) v)$ 
    using b0 c0 p5 remove-cap-right-def by auto
have c6:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow ((\text{caps } t') v) = ((\text{caps } t) v)$ 
    using b1 c0 p6 remove-cap-right-def by auto
have c7:  $\forall v. v \neq \text{did} \wedge v = d$ 
     $\longrightarrow ((\text{caps } s') v) = ((\text{caps } t') v)$ 
    using c5 c6 a3 get-domain-cap-set-from-domain-id-def by auto
have c8:  $d \neq \text{did}$ 
     $\longrightarrow ((\text{caps } s') d) = ((\text{caps } t') d)$ 
    using c7 by auto
have c9:  $((\text{caps } s') d) = ((\text{caps } t') d)$ 
    using c4 c8 by auto
have c10:  $\text{get-domain-cap-set-from-domain-id } s' d$ 
     $= \text{get-domain-cap-set-from-domain-id } t' d$ 
    using c9 get-domain-cap-set-from-domain-id-def by auto
have c11:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow \text{interferes } v s d \longleftrightarrow \text{interferes } v s' d$ 
    using c5 get-domain-cap-set-from-domain-id-def interferes-def by auto
have c12:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow \text{interferes } v t d \longleftrightarrow \text{interferes } v t' d$ 
    using c6 get-domain-cap-set-from-domain-id-def interferes-def by auto
have c13:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow \text{interferes } v s' d \longleftrightarrow \text{interferes } v t' d$ 
    using c11 c12 a1 by auto
have c14:  $\text{interferes } \text{did } s' d \longleftrightarrow \text{interferes } \text{did } t' d$ 
    using c4 get-domain-cap-set-from-domain-id-def interferes-def by auto
have c15:  $(\forall v. \text{interferes } v s' d \longleftrightarrow \text{interferes } v t' d)$ 
    using c13 c14 by auto
then show ?thesis by auto
next
assume c0:  $\neg(\text{REMOVE} = \text{right-to-rm}$ 
     $\wedge \{\text{REMOVE}\} = \text{rights rm-cap})$ 
let ?cs-dst-s =  $\text{get-domain-cap-set-from-domain-id } s \text{ did}$ 
let ?cs-rest-s =  $\{c. c \in ?\text{cs-dst-s} \wedge c \neq \text{rm-cap}\}$ 
let ?cs-dst-t =  $\text{get-domain-cap-set-from-domain-id } t \text{ did}$ 
let ?cs-rest-t =  $\{c. c \in ?\text{cs-dst-t} \wedge c \neq \text{rm-cap}\}$ 
let ?new-cap =  $(\text{target} = \text{target rm-cap},$ 
     $\text{rights} = (\text{rights rm-cap}) - \{\text{right-to-rm}\})$ 
have c1:  $((\text{caps } s') \text{did}) = (\text{insert } ?\text{new-cap } ?\text{cs-rest-s})$ 
    using b0 c0 p5 remove-cap-right-def by auto
have c2:  $((\text{caps } t') \text{did}) = (\text{insert } ?\text{new-cap } ?\text{cs-rest-t})$ 
    using b1 c0 p6 remove-cap-right-def by auto
have c3:  $? \text{cs-rest-s} = ? \text{cs-rest-t}$ 
    using a7 by auto
have c4:  $((\text{caps } s') \text{did}) = ((\text{caps } t') \text{did})$ 
    using c1 c2 c3 by auto
have c5:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow ((\text{caps } s') v) = ((\text{caps } s) v)$ 
    using b0 c0 p5 remove-cap-right-def by auto
have c6:  $\forall v. v \neq \text{did}$ 
     $\longrightarrow ((\text{caps } t') v) = ((\text{caps } t) v)$ 
    using b1 c0 p6 remove-cap-right-def by auto
have c7:  $\forall v. v \neq \text{did} \wedge v = d$ 
     $\longrightarrow ((\text{caps } s') v) = ((\text{caps } t') v)$ 
    using c5 c6 a3 get-domain-cap-set-from-domain-id-def by auto

```

```

have c8:  $d \neq did$ 
   $\longrightarrow ((caps\ s')\ d) = ((caps\ t')\ d)$ 
  using c7 by auto
have c9:  $((caps\ s')\ d) = ((caps\ t')\ d)$ 
  using c4 c8 by auto
have c10:  $get\_domain\_cap\_set\_from\_domain\_id\ s'\ d$ 
   $= get\_domain\_cap\_set\_from\_domain\_id\ t'\ d$ 
  using c9  $get\_domain\_cap\_set\_from\_domain\_id\_def$  by auto
have c11:  $\forall v. v \neq did$ 
   $\longrightarrow interferes\ v\ s\ d \longleftrightarrow interferes\ v\ s'\ d$ 
  using c5  $get\_domain\_cap\_set\_from\_domain\_id\_def\ interferes\_def$  by auto
have c12:  $\forall v. v \neq did$ 
   $\longrightarrow interferes\ v\ t\ d \longleftrightarrow interferes\ v\ t'\ d$ 
  using c6  $get\_domain\_cap\_set\_from\_domain\_id\_def\ interferes\_def$  by auto
have c13:  $\forall v. v \neq did$ 
   $\longrightarrow interferes\ v\ s'\ d \longleftrightarrow interferes\ v\ t'\ d$ 
  using c11 c12 a1 by auto
have c14:  $interferes\ did\ s'\ d \longleftrightarrow interferes\ did\ t'\ d$ 
  using c4 c10  $get\_domain\_cap\_set\_from\_domain\_id\_def\ interferes\_def$  by auto
have c15:  $(\forall v. interferes\ v\ s'\ d \longleftrightarrow interferes\ v\ t'\ d)$ 
  using c13 c14 by auto
then show ?thesis by auto
qed
then show ?thesis by auto
next
assume b0:  $\neg(rm\_cap \in get\_domain\_cap\_set\_from\_domain\_id\ s\ did$ 
   $\wedge REMOVE \in rights\ rm\_cap$ 
   $\wedge right\_to\_rm \in rights\ rm\_cap)$ 
have b1:  $\neg(rm\_cap \in get\_domain\_cap\_set\_from\_domain\_id\ t\ did$ 
   $\wedge REMOVE \in rights\ rm\_cap$ 
   $\wedge right\_to\_rm \in rights\ rm\_cap)$ 
  using b0 a7 by auto
have b2:  $s' = s$ 
  using b0 p5  $remove\_cap\_right\_def$  by auto
have b3:  $(\forall v. interferes\ v\ s\ d \longleftrightarrow interferes\ v\ s'\ d)$ 
  using b2  $get\_domain\_cap\_set\_from\_domain\_id\_def$  by auto
have b4:  $t' = t$ 
  using b1 p6  $remove\_cap\_right\_def$  by auto
have b5:  $(\forall v. interferes\ v\ t\ d \longleftrightarrow interferes\ v\ t'\ d)$ 
  using b4  $get\_domain\_cap\_set\_from\_domain\_id\_def$  by auto
have b6:  $(\forall v. interferes\ v\ s'\ d \longleftrightarrow interferes\ v\ t'\ d)$ 
  using a1 b3 b5 by auto
then show ?thesis by auto
qed
}
then show ?thesis by auto
qed

lemma remove-cap-right-wsc-dom-eps:
assumes p0:  $reachable0\ s$ 
and p1:  $reachable0\ t$ 
and p2:  $s \sim d \sim t$ 
and p3:  $interferes\ did\ s\ d$ 
and p4:  $s \sim did \sim t$ 
and p5:  $s' = fst\ (remove\_cap\_right\ s\ did\ rm\_cap\ right\_to\_rm)$ 
and p6:  $t' = fst\ (remove\_cap\_right\ t\ did\ rm\_cap\ right\_to\_rm)$ 
shows  $get\_endpoints\_of\_domain\ s'\ d = get\_endpoints\_of\_domain\ t'\ d$ 
proof -

```



```

{
  have a0: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ t \ d$ )
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a4: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
    using p0 p5 remove-cap-right-notchg-dom-eps by auto
  have a5: get-endpoints-of-domain t d = get-endpoints-of-domain t' d
    using p1 p6 remove-cap-right-notchg-dom-eps by auto
  have a6: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
    using a0 a4 a5 by auto
}
then show ?thesis by auto
qed

```

lemma remove-cap-right-wsc-ep-msgs:

assumes p0: reachable0 s

and p1: reachable0 t

and p2:  $s \sim d \sim t$

and p3: interferes did s d

and p4:  $s \sim \text{did} \sim t$

and p5:  $s' = \text{fst } (\text{remove-cap-right } s \ \text{did } \text{rm-cap } \text{right-to-rm})$

and p6:  $t' = \text{fst } (\text{remove-cap-right } t \ \text{did } \text{rm-cap } \text{right-to-rm})$

shows ( $\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$

$\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep$ )

proof -

```

{
  have a0: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } t \ ep$ )
    by (meson p2 vpeq1-def)
  have a1: ( $\forall v. \text{interferes } v \ s \ d \longleftrightarrow \text{interferes } v \ t \ d$ )
    by (meson p2 vpeq1-def)
  have a2: interferes did t d
    using p3 a1 by auto
  have a4: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s \ ep = \text{get-msg-set-from-endpoint-id } s' \ ep$ )
    using p0 p5 remove-cap-right-notchg-ep-msgs by auto
  have a5: ( $\forall ep. ep \in \text{get-endpoints-of-domain } t \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } t \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep$ )
    using p1 p6 remove-cap-right-notchg-ep-msgs by auto
  have a6: get-endpoints-of-domain s d = get-endpoints-of-domain t d
    by (meson p2 vpeq1-def)
  have a7: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep$ )
    using a0 a4 a5 a6 by auto
  have a8: get-endpoints-of-domain s d = get-endpoints-of-domain s' d
    using p0 p5 remove-cap-right-notchg-dom-eps by auto
  have a9: ( $\forall ep. ep \in \text{get-endpoints-of-domain } s' \ d$ 
     $\longrightarrow \text{get-msg-set-from-endpoint-id } s' \ ep = \text{get-msg-set-from-endpoint-id } t' \ ep$ )
    using a7 a8 by auto
}

```

then show ?thesis by auto

qed

lemma remove-cap-right-wsc:

assumes p0: reachable0 s

```

and p1: reachable0 t
and p2: s ~ d ~ t
and p3: interferes did s d
and p4: s ~ did ~ t
and p5: s' = fst (remove-cap-right s did rm-cap right-to-rm)
and p6: t' = fst (remove-cap-right t did rm-cap right-to-rm)
shows s' ~ d ~ t'
proof -
{
  have a0: get-domain-cap-set-from-domain-id s' d = get-domain-cap-set-from-domain-id t' d
    using p0 p1 p2 p3 p4 p5 p6 remove-cap-right-wsc-domain-cap-set by blast
  have a1: (∀ v. interferes v s' d ⟷ interferes v t' d)
    using p0 p1 p2 p3 p4 p5 p6 remove-cap-right-wsc-policy by blast
  have a2: get-endpoints-of-domain s' d = get-endpoints-of-domain t' d
    using p0 p1 p2 p3 p4 p5 p6 remove-cap-right-wsc-dom-eps by blast
  have a3: (∀ ep. ep ∈ get-endpoints-of-domain s' d
    ⟶ get-msg-set-from-endpoint-id s' ep = get-msg-set-from-endpoint-id t' ep )
    using p0 p1 p2 p3 p4 p5 p6 remove-cap-right-wsc-ep-msgs by blast
  have a4: s' ~ d ~ t'
    using a0 a1 a2 a3 vpeq1-def by auto
}
then show ?thesis by auto
qed

```

```

lemma remove-cap-right-wsc-e:
assumes p0: reachable0 s
and p1: reachable0 t
and p2: a = (Remove-Cap-Right did dst-cap right-to-rm)
and p3: s ~ d ~ t
and p4: interferes (the (domain-of-event a)) s d
and p5: s ~ did ~ t
and p6: s' = exec-event s a
and p7: t' = exec-event t a
shows s' ~ d ~ t'
proof -
{
  have a0: (the (domain-of-event a)) = did
    using p2 domain-of-event-def by auto
  have a1: s' = fst (remove-cap-right s did dst-cap right-to-rm)
    using p2 p6 exec-event-def by auto
  have a2: t' = fst (remove-cap-right t did dst-cap right-to-rm)
    using p2 p7 exec-event-def by auto
  have a3: (interferes did s d)
    using p4 a0 by auto
  have a4: s' ~ d ~ t'
    using a1 a2 a3 p0 p1 p3 p5 remove-cap-right-wsc by blast
}
then show ?thesis by auto
qed

```

```

lemma remove-cap-right-dwsc-e: dynamic-weakly-step-consistent-e (Remove-Cap-Right did dst-cap right-to-rm)
proof -
{
  have ∀ d s t. reachable0 s ∧ reachable0 t
    ∧ (s ~ d ~ t)
    ∧ (interferes (the (domain-of-event (Remove-Cap-Right did dst-cap right-to-rm))) s d)
    ∧ (s ~ (the (domain-of-event (Remove-Cap-Right did dst-cap right-to-rm))) ~ t)
    ⟶ ((exec-event s (Remove-Cap-Right did dst-cap right-to-rm)) ~ d ~ (exec-event t (Remove-Cap-Right did

```

```

dst-cap right-to-rm)))
  proof -
  {
    fix d s t
    assume p1: reachable0 s
    assume p2: reachable0 t
    assume p3: (s ~ d ~ t)
    assume p4: (interferes (the (domain-of-event (Remove-Cap-Right did dst-cap right-to-rm))) s d)
    assume p5: (s ~ (the (domain-of-event (Remove-Cap-Right did dst-cap right-to-rm))) ~ t)
    have ((exec-event s (Remove-Cap-Right did dst-cap right-to-rm)) ~ d ~ (exec-event t (Remove-Cap-Right did
dst-cap right-to-rm)))
      by (metis Event.simps(56) domain-of-event-def option.sel p1 p2 p3 p4 p5 remove-cap-right-wsc-e)
    }
    then show ?thesis by blast
  }
  qed
}
then show ?thesis
  using dynamic-weakly-step-consistent-e-def by blast
qed

```

### 0.10.8 proving the "dynamic step consistent" property

**theorem** *dynamic-weakly-step-consistent:dynamic-weakly-step-consistent*

```

proof -
{
  fix e
  have dynamic-weakly-step-consistent-e e
  proof (induct e)
    case (Client-Lookup-Endpoint-Name x x1)
    show ?case
    using client-lookup-endpoint-name-dwsc-e by blast
    case (Send-Queuing-Message x1a x2 x3a)
    show ?case
    using send-queuing-message-dwsc-e by blast
    case (Receive-Queuing-Message x)
    show ?case
    using receive-queuing-message-dwsc-e by blast
    case (Get-My-Endpoints-Set x)
    show ?case
    using get-my-endpoints-set-dwsc-e by blast
    case (Get-Caps x)
    show ?case
    using get-caps-dwsc-e by blast
    case (Grant-Endpoint-Cap x1a x2 x3a)
    show ?case
    using grant-endpoint-cap-dwsc-e by blast
    case (Remove-Cap-Right x1a x2 x3a)
    show ?case
    using remove-cap-right-dwsc-e by blast
  }
  qed
}
then show ?thesis
  using dynamic-weakly-step-consistent-all-evt by blast
qed

```

**theorem** *noninfluence-sat: noninfluence*

**using** *dynamic-local-respect uc-eq-noninf dynamic-weakly-step-consistent weak-with-step-cons* **by** *blast*

```

theorem weak-noninfluence-sat: weak-noninfluence using noninf-impl-weak noninfluence-sat by blast

theorem nonleakage-sat: nonleakage
  using noninf-impl-nonlk noninfluence-sat by blast

theorem noninterference-r-sat: noninterference-r
  using noninf-impl-nonintf-r noninfluence-sat by blast

theorem noninterference-sat: noninterference
  using noninterference-r-sat nonintf-r-impl-noninterf by blast

theorem weak-noninterference-r-sat: weak-noninterference-r
  using noninterference-r-sat nonintf-r-impl-wk-nonintf-r by blast

theorem weak-noninterference-sat: weak-noninterference
  using noninterference-sat nonintf-impl-weak by blast
end

```