





Solution Review: Rearrange Sorted List in Max/Min Form

This lesson gives a solution to the challenge in the previous lesson.



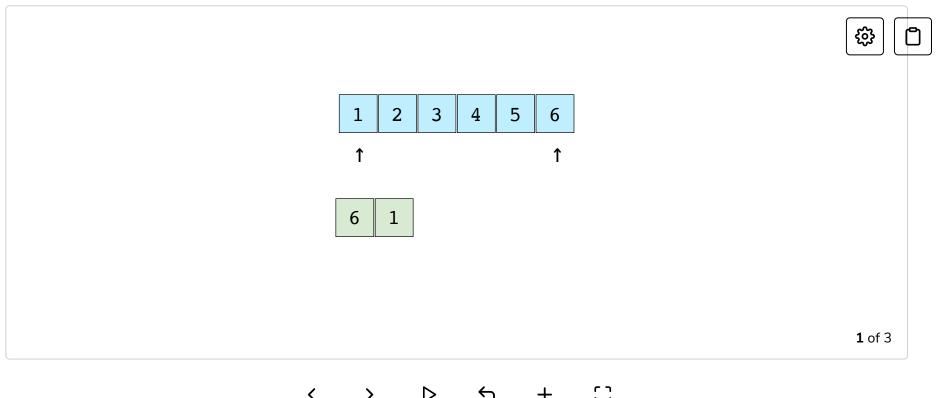
- Solution #1: Creating a new list
 - Time Complexity
- Solution #2: UsingO(1)Extra Space
 - Time Complexity

Solution #1: Creating a new list

In this solution, we first create a new empty list that we will append the appropriate elements to and return. We then iterate through the list starting from the 0th index till the middle of the list indexed as lst[length(list)/2]. So if the length of the given list is 10, the iterator variable i on line 4 in our solution would start from 0 and end at 10/2 = 5. Note that the starting index 0 in the example is inclusive, and the ending index 5 is exclusive. At each iteration, we first append the largest unappended element and then the smallest. So in the first iteration, i = 0 and lst[-(0+1)] = lst[-1] corresponds

to the last element of the list, which is also the largest. So the largest element in the list is appended to result first, and then the current or element indexed by i is appended. Next, the second largest and the second smallest are appended and so on until the end of the list.

```
def max_min(lst):
        result = []
 2
        # iterate half list
        for i in range(len(lst)//2):
 4
 5
            # Append corresponding last element
            result.append(lst[-(i+1)])
 6
            # append current element
 7
 8
            result.append(lst[i])
        if len(lst) % 2 == 1:
 9
10
            # if middle value then append
11
            result.append(lst[len(lst)//2])
12
        return result
13
14
15
    print(max_min([1, 2, 3, 4, 5, 6]))
16
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```



Time Complexity

The time complexity of this problem is O(n) as the list is iterated over once.

Solution #2: Using O(1) Extra Space

```
1 def max_min(lst):
2  # Return empty list for empty list
3  if (len(lst) is 0):
4   return []
5
6  maxIdx = len(lst) - 1  # max index
```

```
minIdx = 0 # first index
 7
        maxElem = lst[-1] + 1 \# Max element
        # traverse the list
 9
        for i in range(len(lst)):
10
            # even number means max element to append
11
12
            if i \% 2 == 0:
13
                lst[i] += (lst[maxIdx] % maxElem) * maxElem
14
                maxIdx -= 1
            # odd number means min number
15
16
            else:
17
                lst[i] += (lst[minIdx] % maxElem) * maxElem
18
                minIdx += 1
19
20
        for i in range(len(lst)):
            lst[i] = lst[i] // maxElem
21
22
        return lst
23
24
25
    print(max_min([0, 1, 2, 3, 4, 5, 6, 7, 8, 9]))
26
```

This solution uses some math to store two elements at one index. This is achieved from the following line of code,

```
lst[i] += (lst[maxIdx] % maxElem) * maxElem
```

lst [maxIdx] is stored as a multiplier and lst[i] is stored as remainder. For example in the list, $\begin{bmatrix} 1 \\ 3 \end{bmatrix}$, 4, 5, 6, 7, 8, 9], the maxElem is 10 and 90 is stored at index 0. One we have 90, we can get the new element 9 using 90/10. Also, we can go back to the original element, 1, using the expression 90%10.

This allows us to swap the numbers in place without using any extra space. To get the final list, we simply divide each element by maxElem as done in the last for loop.

Note: This approach only works for non-negative numbers!

Time Complexity

The time complexity of this solution is in O(n). The space complexity is constant.

