





Intervals Intersection (medium)

We'll cover the following

- Problem Statement
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 - Time complexity
 - Space complexity

Problem Statement

Given two lists of intervals, find the **intersection of these two lists**. Each list consists of **disjoint intervals sorted on their start time**.

Example 1:

```
Input: arr1=[[1, 3], [5, 6], [7, 9]], arr2=[[2, 3], [5, 7]]
```

Output: [2, 3], [5, 6], [7, 7]

Explanation: The output list contains the common intervals between

the two lists.

Example 2:

Input: arr1=[[1, 3], [5, 7], [9, 12]], arr2=[[5, 10]]

Output: [5, 7], [9, 10]

Explanation: The output list contains the common intervals between

the two lists.

Try it yourself



Try solving this question here:

```
Python3
                         us JS
                                      G C++
👙 Java
    def merge(intervals_a, intervals_b):
 1
 2
       result = []
       # TODO: Write your code here
 3
 4
       return result
 5
 6
 7
    def main():
 8
       print("Intervals Intersection: " + str(merge(|
       print("Intervals Intersection: " + str(merge())
 9
10
11
12
    main()
13
```

Solution

This problem follows the Merge Intervals

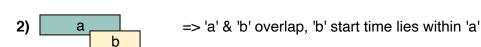
(https://www.educative.io/collection/page/5668639101419520/5671464854355 968/5652017242439680/) pattern. As we have discussed under Insert Interval (https://www.educative.io/collection/page/5668639101419520/5671464854355 968/5718314357620736/), there are five overlapping possibilities between two intervals 'a' and 'b'. A close observation will tell us that whenever the two intervals overlap, one of the interval's start time lies within the other interval. This rule can help us identify if any two intervals overlap or not.

1)

=> 'a' and 'b' don't overlap







a) => 'a' & 'b' overlap, 'b' start time lies within 'a'

a => 'a' & 'b' overlap, 'a' start time lies within 'b'

=> 'a' & 'b' overlap, 'a' start time lies within 'b'

Now, if we have found that the two intervals overlap, how can we find the overlapped part?

Again from the above diagram, the overlapping interval will be equal to:

```
start = max(a.start, b.start)
end = min(a.end, b.end)
```

That is, the highest start time and the lowest end time will be the overlapping interval.

So our algorithm will be to iterate through both the lists together to see if any two intervals overlap. If two intervals overlap, we will insert the overlapped part into a result list and move on to the next interval which is finishing early.

Code

Here is what our algorithm will look like:

```
4
    5
                         while i < len(intervals_a) and j < len(intervals_a)
                                  # check if intervals overlap and intervals_a
    6
                                  a_overlaps_b = intervals_a[i][start] >= intε
    7
    8
                                                                                                    intervals_a[i][start] <= intervals_a[i]</pre>
    9
                                  # check if intervals overlap and intervals_a
10
11
                                  b_overlaps_a = intervals_b[j][start] >= intervals_b[start] >= intervals_b[start
                                                                                                   intervals_b[j][start] <= intervals_b[j]</pre>
12
13
14
                                 # store the the intersection part
15
                                  if (a_overlaps_b or b_overlaps_a):
                                           result.append([max(intervals_a[i][start]],
16
                                                    intervals_a[i][end], intervals_b[j][end]
17
18
                                  # move next from the interval which is finis
19
                                  if intervals_a[i][end] < intervals_b[j][end]</pre>
20
21
                                           i += 1
22
                                  else:
23
                                           j += 1
24
25
                          return result
26
27
28
                def main():
```

Time complexity

As we are iterating through both the lists once, the time complexity of the above algorithm is O(N+M), where 'N' and 'M' are the total number of intervals in the input arrays respectively.

Space complexity

Ignoring the space needed for the result list, the algorithm runs in constant space O(1).



Insert Interval (medium)

Conflicting Appointments (Migdium

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