





Longest Substring with Same Letters after Replacement (hard)

We'll cover the following

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Problem Statement

Given a string with lowercase letters only, if you are allowed to **replace no more than 'k' letters** with any letter, find the **length of the longest substring having the same letters** after replacement.

Example 1:

Input: String="aabccbb", k=2

Output: 5

Explanation: Replace the two 'c' with 'b' to have a longest repeat

ing substring "bbbbb".

Example 2:

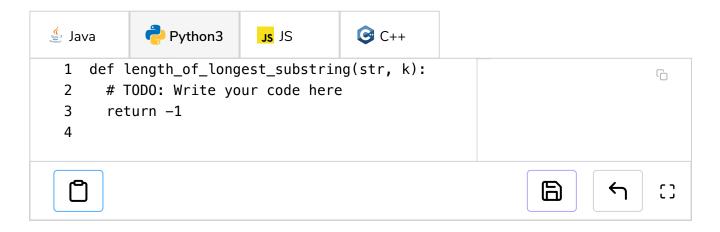
```
Input: String="abbcb", k=1
Output: 4
Explanation: Replace the 'c' with 'b' to have a longest repeatin g substring "bbbb".
```

Example 3:

```
Input: String="abccde", k=1
Output: 3
Explanation: Replace the 'b' or 'd' with 'c' to have the longest r
epeating substring "ccc".
```

Try it yourself

Try solving this question here:



Solution

This problem follows the **Sliding Window** pattern, and we can use a similar dynamic sliding window strategy as discussed in No-repeat Substring (https://www.educative.io/collection/page/5668639101419520/5671464854355 968/5485010335301632/). We can use a HashMap to count the frequency of each letter.

• We'll iterate through the string to add one letter at a time in the window.

- We'll also keep track of the count of the maximum repeating letter in any window (let's call it maxRepeatLetterCount).
- So, at any time, we know that we can have a window which has one letter repeating maxRepeatLetterCount times; this means we should try to replace the remaining letters.
- If we have more than 'k' remaining letters, we should shrink the window as we are not allowed to replace more than 'k' letters.

While shrinking the window, we don't need to update maxRepeatLetterCount (which makes it global count; hence, it is the maximum count for ANY window). Why don't we need to update this count when we shrink the window? The answer: In any window, since we have to replace all the remaining letters to get the longest substring having the same letter, we can't get a better answer from any other window even though all occurrences of the letter with frequency maxRepeatLetterCount is not in the current window.

Code

Here is what our algorithm will look like:

```
Python3
                         G C++
👙 Java
                                      JS JS
 9
           frequency_map[right_char] = 0
         frequency_map[right_char] += 1
10
         max_repeat_letter_count ·= · max(
11
     .....max_repeat_letter_count, frequency_map[ric
12
13
14
        # Current window size is from window_start t
15
         # repeating 'max_repeat_letter_count' times,
         # repeating 'max_repeat_letter_count' times
16
         # if the remaining letters are more than 'k'
17
18
         # are not allowed to replace more than 'k' l
19
         if (window end - window start + 1 - max repε
20
           left_char = str1[window_start]
           frequency_map[left_char] -= 1
21
22
          window_start += 1
```

```
23
24
        max_length = max(max_length, window_end - wi
                                                                            €€}
25
      return max_length
26
27
28
    def main():
29
      print(length_of_longest_substring("aabccbb", 2
      print(length_of_longest_substring("abbcb", 1))
30
      print(length_of_longest_substring("abccde", 1)
31
32
33
    main()
34
35
                                                               \triangleright
```

Time Complexity

The above algorithm's time complexity will be O(N), where 'N' is the number of letters in the input string.

Space Complexity

As we expect only the lower case letters in the input string, we can conclude that the space complexity will be O(26) to store each letter's frequency in the **HashMap**, which is asymptotically equal to O(1).

