Sp18 CS 61B Discussion 6

Welcome!

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Announcements

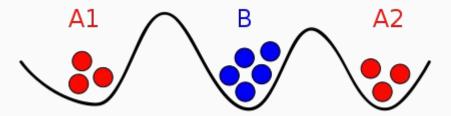
- MT1: Regrade requests due Friday at noon.
- Project 2 released
- HW1 released

Quiz Instructions

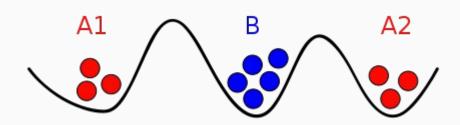
- If you haven't yet, please also neatly put your email address outside the name box if you want to be emailed!
- Bubble number 41.

Aside

- Two generals (A1, A2), surrounding the enemy (B).
- Can only communicate via messengers, but messengers can be captured



- Goal: Simultaneously attack B.
 - If attack simultaneously, win!
 - If attack at different times, lose.
 - How to coordinate?



- A1 -> A2: Attack at noon!
- Can you attack at noon?

- A1 -> A2: Attack at noon!
- Can you attack at noon?
 - No: Not sure if A2 got the message.

- A1 -> A2: Attack at noon!
- A2 -> A1: Okay, I will attack at noon.
- Can you attack at noon?

- A1 -> A2: Attack at noon!
- A2 -> A1: Okay, I will attack at noon.
- Can you attack at noon?
 - No! A2 is not sure if you got the confirmation!
 - If you didn't get the confirmation, you won't attack.

- A1 -> A2: Attack at noon!
- A2 -> A1: Okay, I will attack at noon.
- A1 -> A2: Okay, got your confirmation.
- Can you attack at noon?

- A1 -> A2: Attack at noon!
- A2 -> A1: Okay, I will attack at noon.
- A1 -> A2: Okay, got your confirmation.
- Can you attack at noon?
 - No! You're not sure if A2 got the confirmation...

- Turns out, you can't guarantee the win!
- Problem: The other general is never sure if the last message was received.

Okay...? (CS 168)

- This relates to networking!
- Can messages sent over an unreliable network guarantee two entities to do something simultaneously?

Distributed Transaction (CS 168)

- Since the answer is no, early internet pioneers tried to solve a different problem:
 - Distributed transaction: Two or more machines agree to do something, or not do it.

Distributed Transaction (CS 168)

- Solution: Two-Phase Commit protocol.
 - Developed by Turing Award winner Jim Gray (first Berkeley CS PhD, 1969)

- N Byzantine generals encircle a city.
- Generals can decide whether to attack or retreat.
- Important: each general agrees on a common decision.
 - All attack or retreat: Good!
 - Partial attack/retreat: Bad.

- A subset of the n generals are traitors!
- Can cast vote for suboptimal strategy, and can selectively lie to every other general.
- Again, generals are separated and must use messengers, which can be captured.

- Solutions: Many! Solutions are categorized as Byzantine fault tolerance solutions.
- 1999: Miguel Castro and Barbara Liskov solves this problem efficiently (O(N²)).
 - Liskov wins a Turing award for her efforts:)
 - Berkeley alum! :D

- What solution does Google use?
 - Paxos protocol
 - Developed by Leslie Lamport (won the 2013
 Turing Award for his work!)

Moral of the Story

• If you want to win a Turing award...

References

- https://en.wikipedia.org/wiki/Two_Generals%27_Proble
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- https://en.wikipedia.org/wiki/Byzantine_fault_tolerance

Onto Discussion