

Hash Tables, Sets And Dictionaries

Hashing and Collisions

0	1	2	...	m-1
null	null	SoftUni	...	C#

SoftUni Team

Technical Trainers



SoftUni



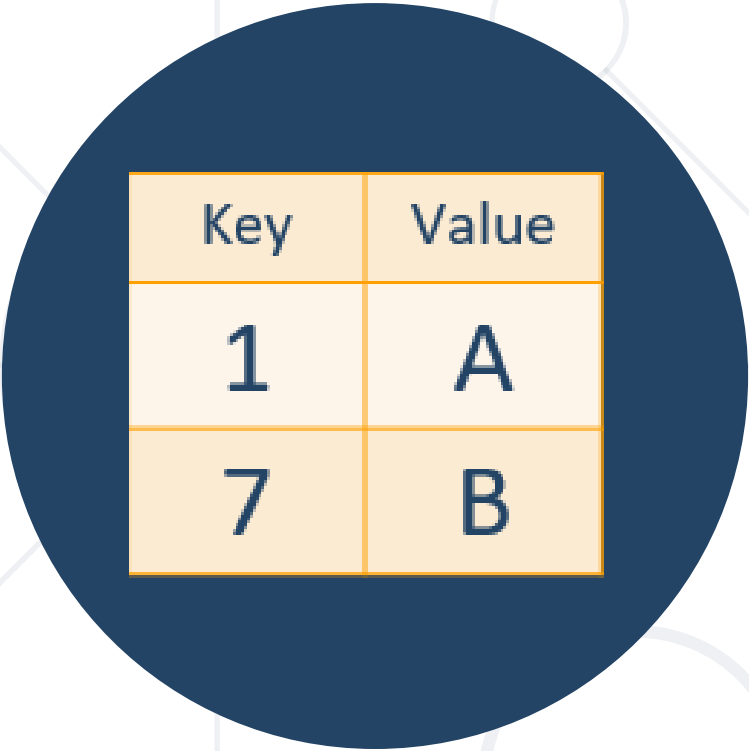
Software University

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1. Hash Tables
2. Sets
3. Dictionaries





Key	Value
1	A
7	B

Hash Tables

Hashing and Collision Resolution

- Given a key of any type, convert it to an integer

Ivan

Hash Function

398

Pesho

Hash Function

511

```
class Person {  
    string firstName;  
    string lastName;  
    int age;  
}
```

Ivan
Petrov
25

Hash Function

25950

Hash Function (2)

```
class Person {  
    string firstName;  
    string lastName;  
    int age;  
  
    public override int GetHashCode() {  
  
    }  
}
```

Hash Function

Hash Function (2)

```
class Person {  
    string firstName;  
    string lastName;  
    int age;  
  
    public override int GetHashCode() {  
        int firstNameHash = firstName.GetHashCode() * age;  
        int lastNameHash = lastName.GetHashCode() * age;  
  
        return firstNameHash + lastNameHash;  
    }  
}
```

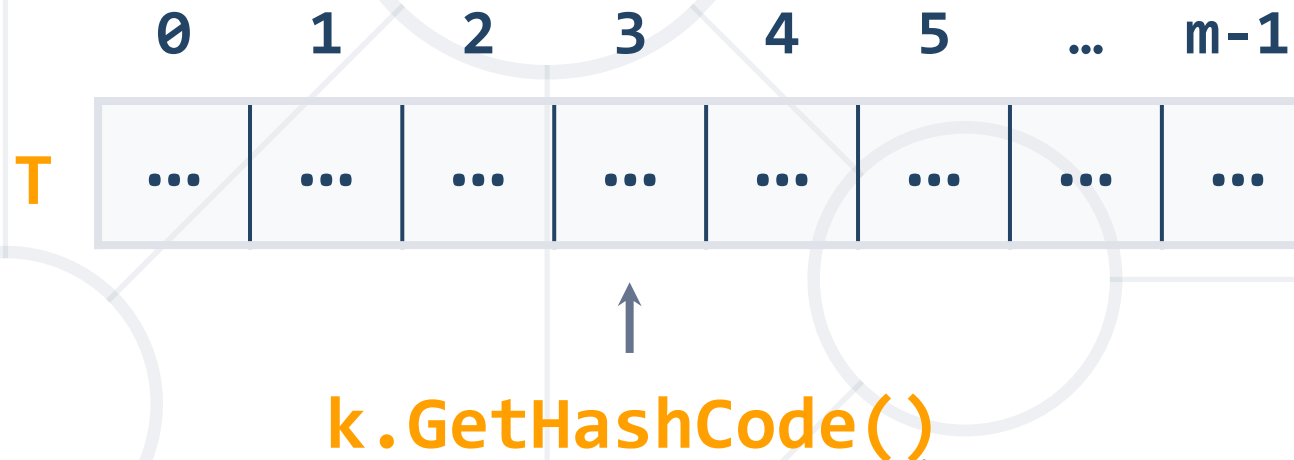
- A hash table is an array that holds a set of **{key, value} pairs**
- The process of mapping a key to a position in a table is called **hashing**



Hash table
of size **m**

Hash Functions and Hashing

- A hash table has **m** slots, indexed from **0** to **m-1**
- A hash function converts **keys** into array indices




Returns 32-bit
integer

- Perfect hashing function (PHF)
 - $h(k)$: one-to-one mapping of each key k to an integer in the range $[0, m-1]$
 - The PHF maps each key to a **distinct** integer within some manageable range
- Finding a perfect hashing function is impossible in most cases

- Good hashing function
 - **Consistent** - equal keys must produce the same hash value
 - **Efficient** - efficient to compute the hash
 - **Uniform** - should uniformly distribute the keys

TIME'S

- Which of the following is **not** property of a **GetHashCode()** for strings
 - Can return a negative integer
 - Can take time proportional to the length of the string to compute
 - A string and its reverse will have the same hash code
 - Two strings with different hash code values are different strings

- Which of the following is **not** property of a **GetHashCode()** for strings
 - Can return a negative integer
 - Can take time proportional to the length of the string to compute `"ab".GetHashCode() != "ba".GetHashCode()`
 - A string and its reverse will have the same hash code 
 - Two strings with different hash code values are different strings

- Array with length 16
- Insert "Example"

Example

Hash Function

511

511 is bigger than
the table length

- Use the remainder of
 $\text{GetHashCode()} / \text{Array.Length}$

$$511 \% 16 = 15$$

	0
	1
	2
	3
	4
	5
	6
	7
	...
	15

Adding to Hash Table

Example

Hash Function % 10

	0
	1
	2
	3
	4
	5
	6
	7
	8
	9

Adding to Hash Table (2)

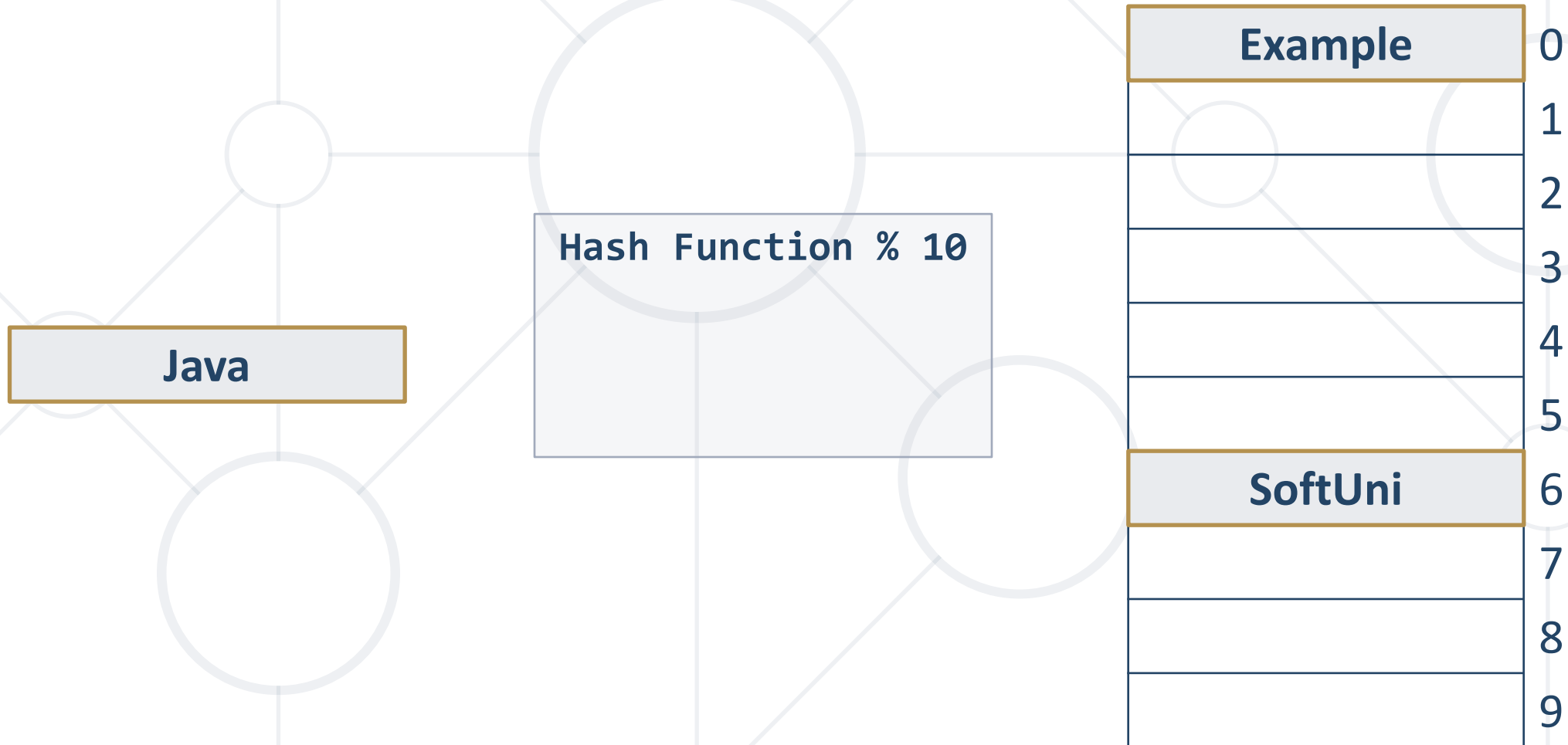
SoftUni

Hash Function % 10

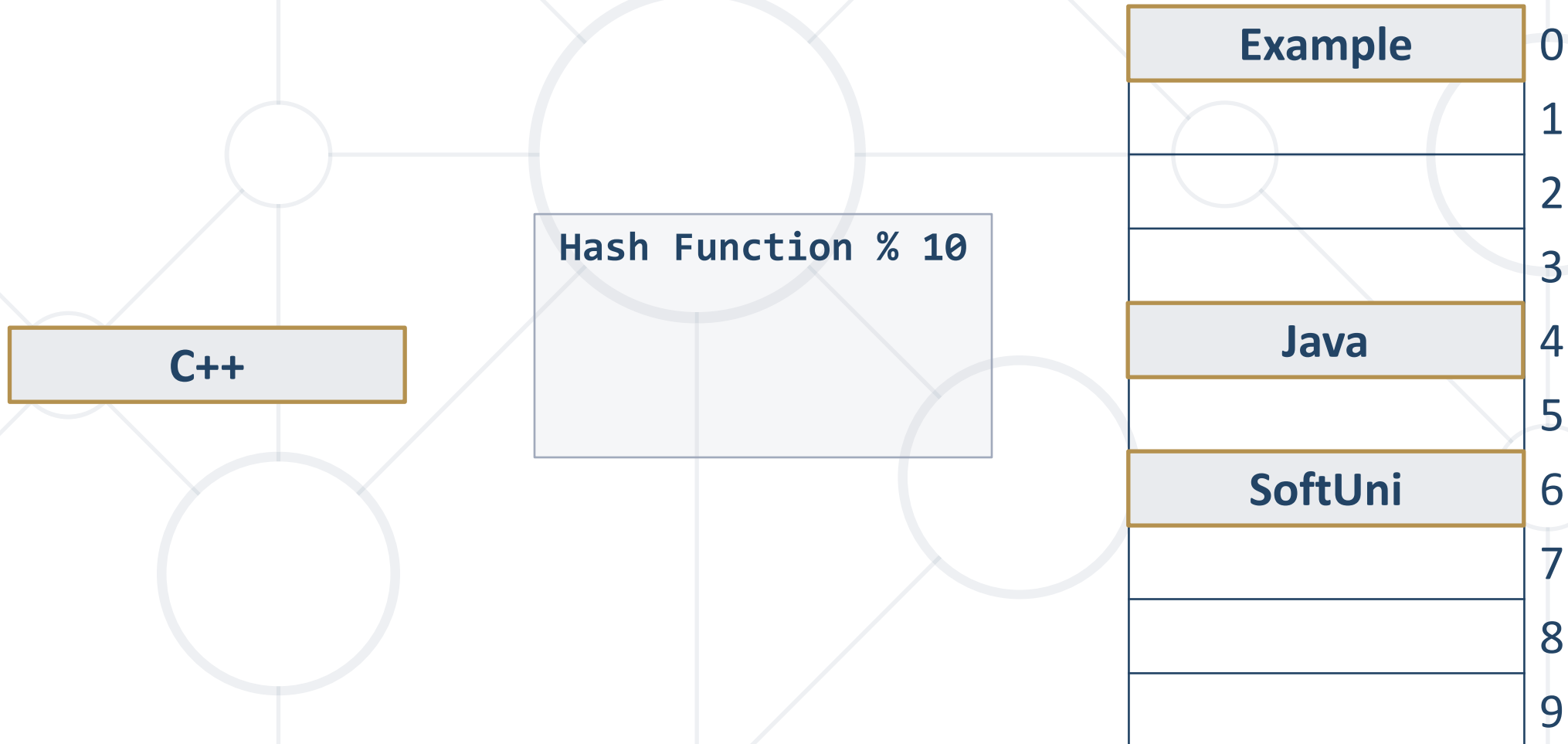
Example

0
1
2
3
4
5
6
7
8
9

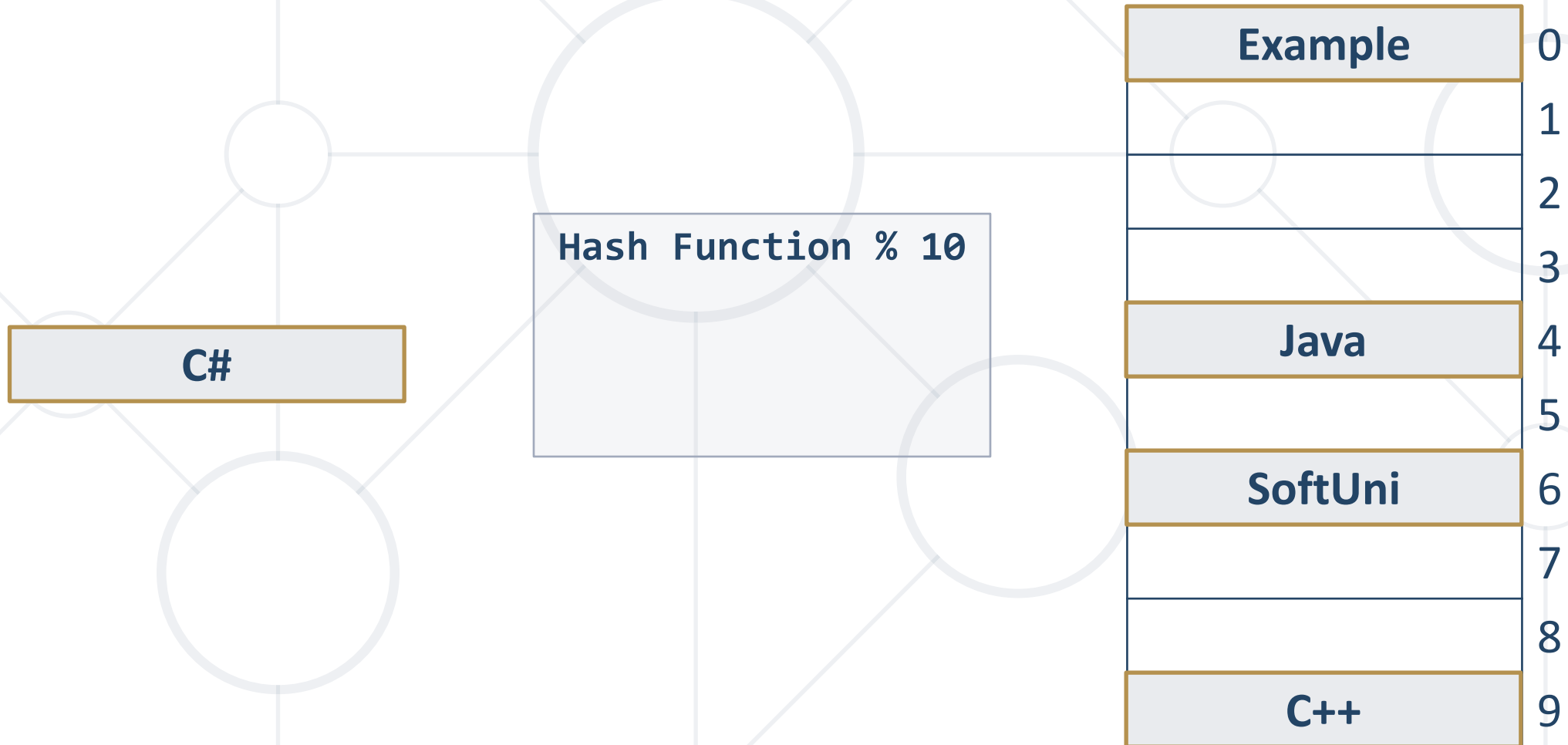
Adding to Hash Table (3)



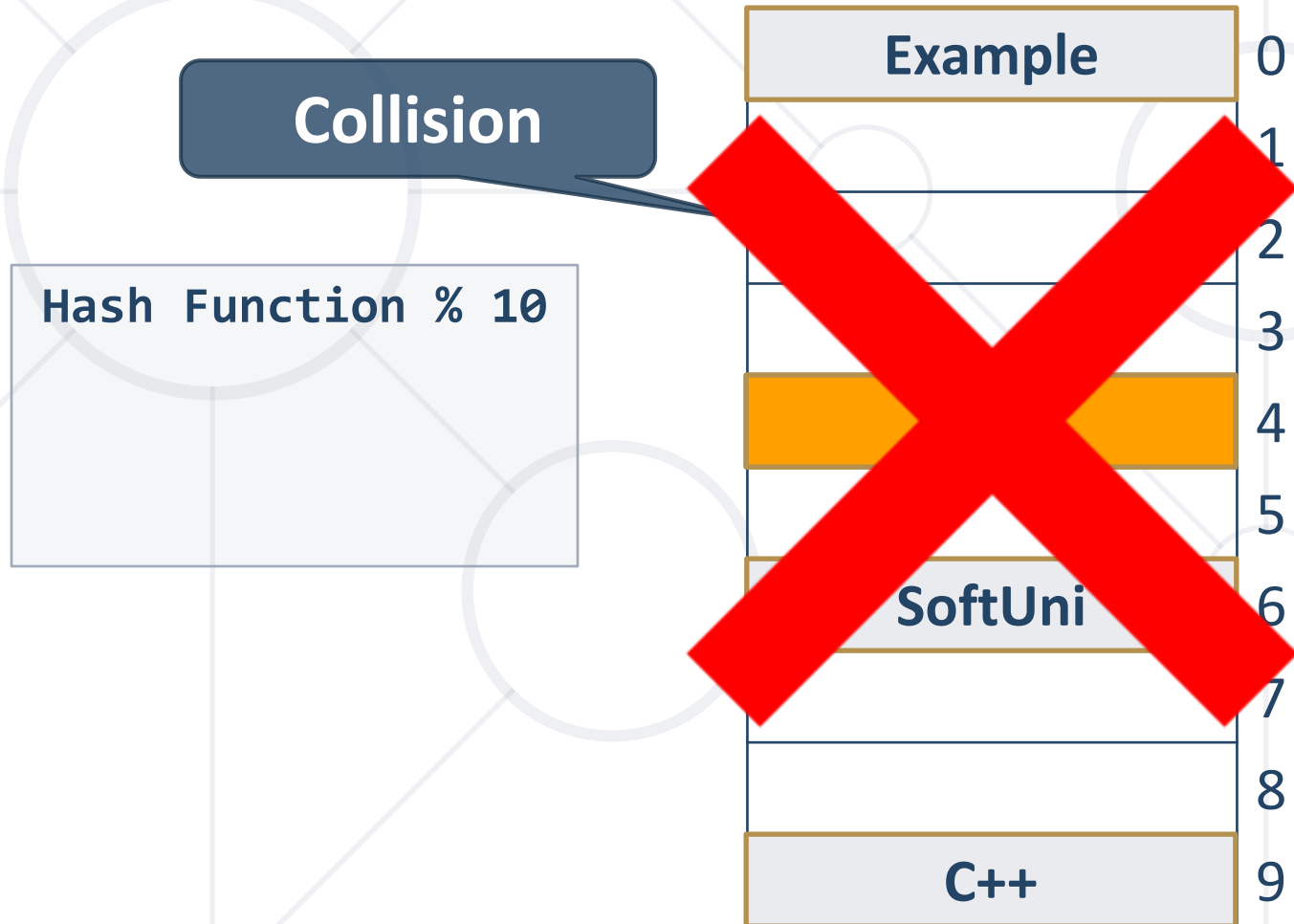
Adding to Hash Table (4)



Adding to Hash Table (5)

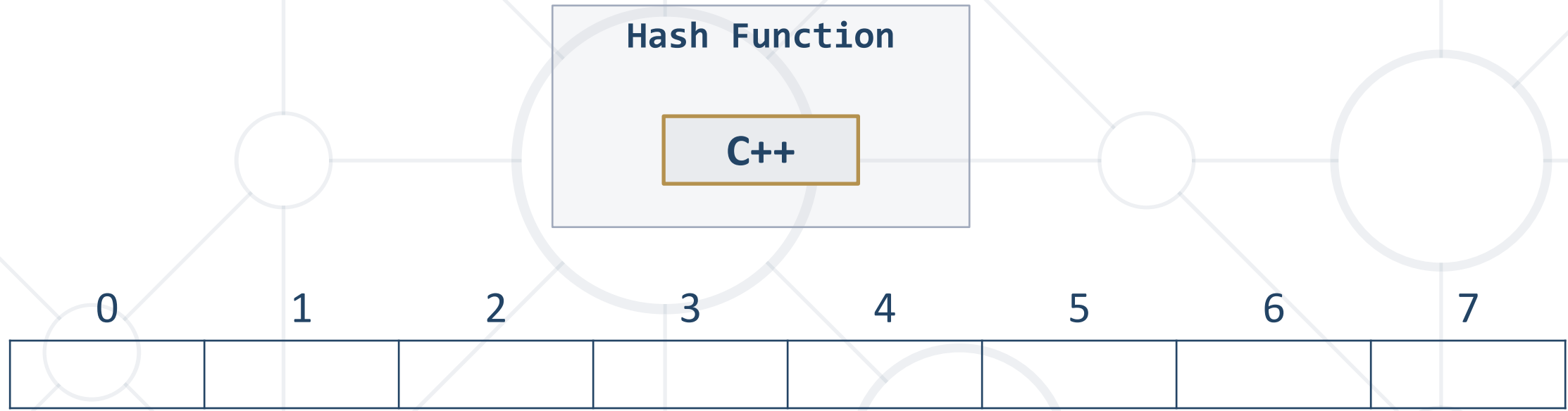


Adding to Hash Table (6)

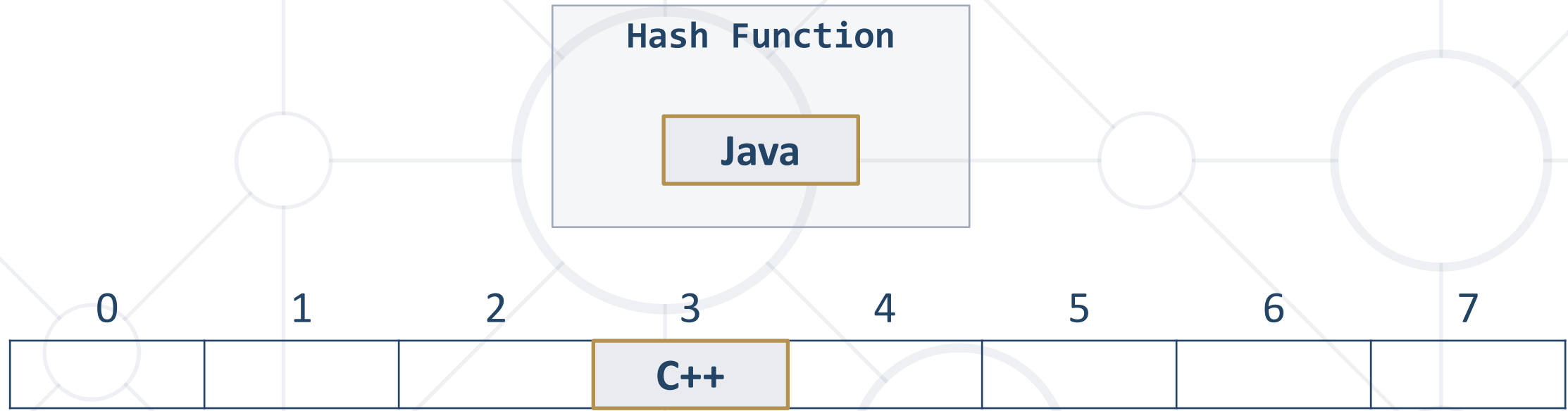


- A **collision** comes when **different keys** have the **same hash value**
 - $h(k_1) = h(k_2)$ for $k_1 \neq k_2$
- When the number of collisions is sufficiently small, the hash tables work quite well (fast)
- Several **collisions resolution strategies** exist
 - **Chaining** collided keys (+ values) in a list
 - Using **other slots** in the table (open addressing)
 - Cuckoo hashing
 - Many other

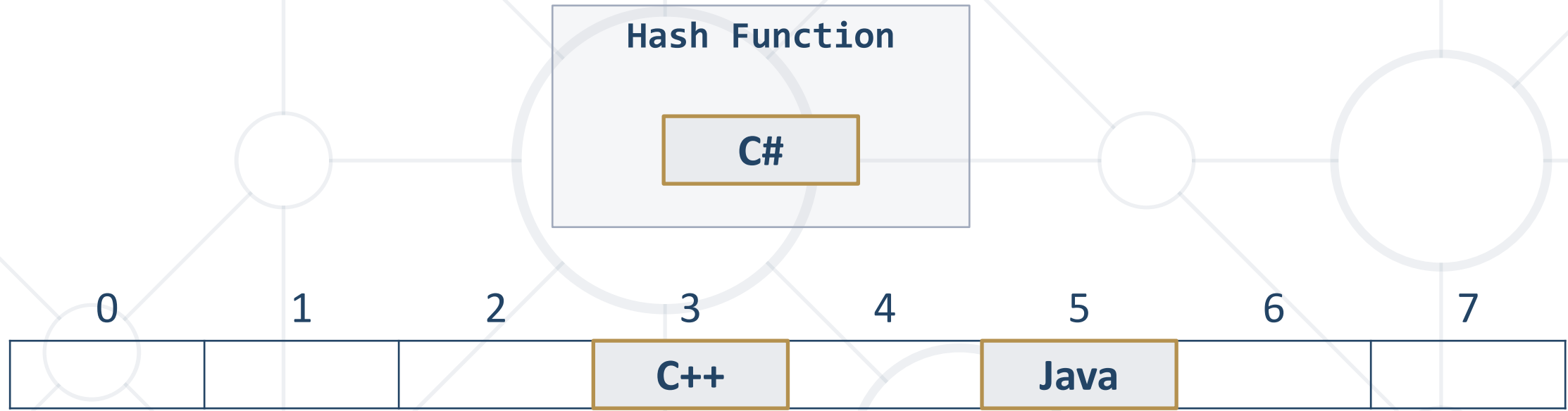
Collision Resolution: Chaining



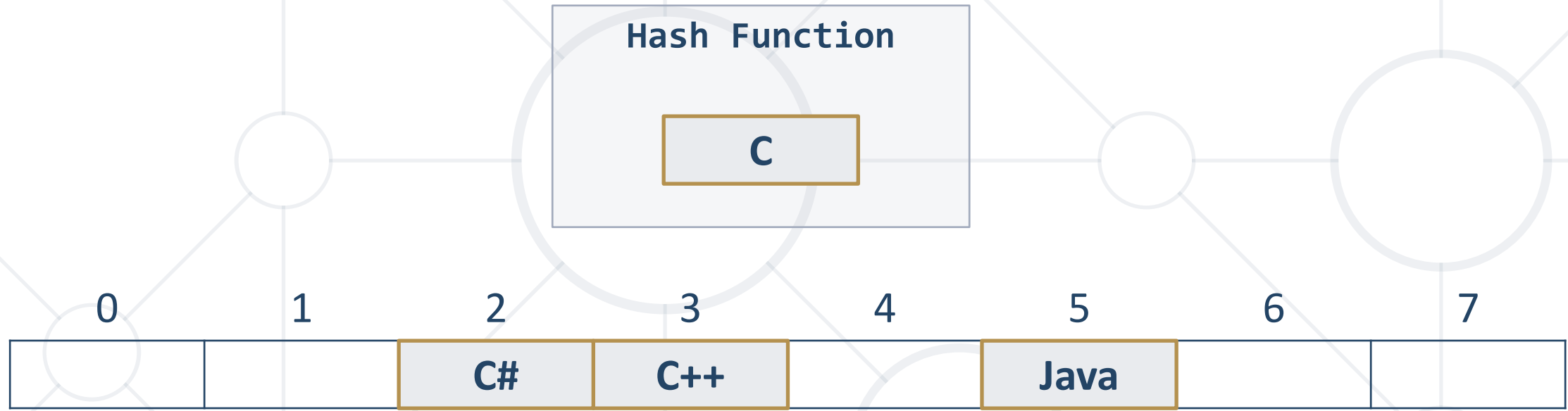
Collision Resolution: Chaining



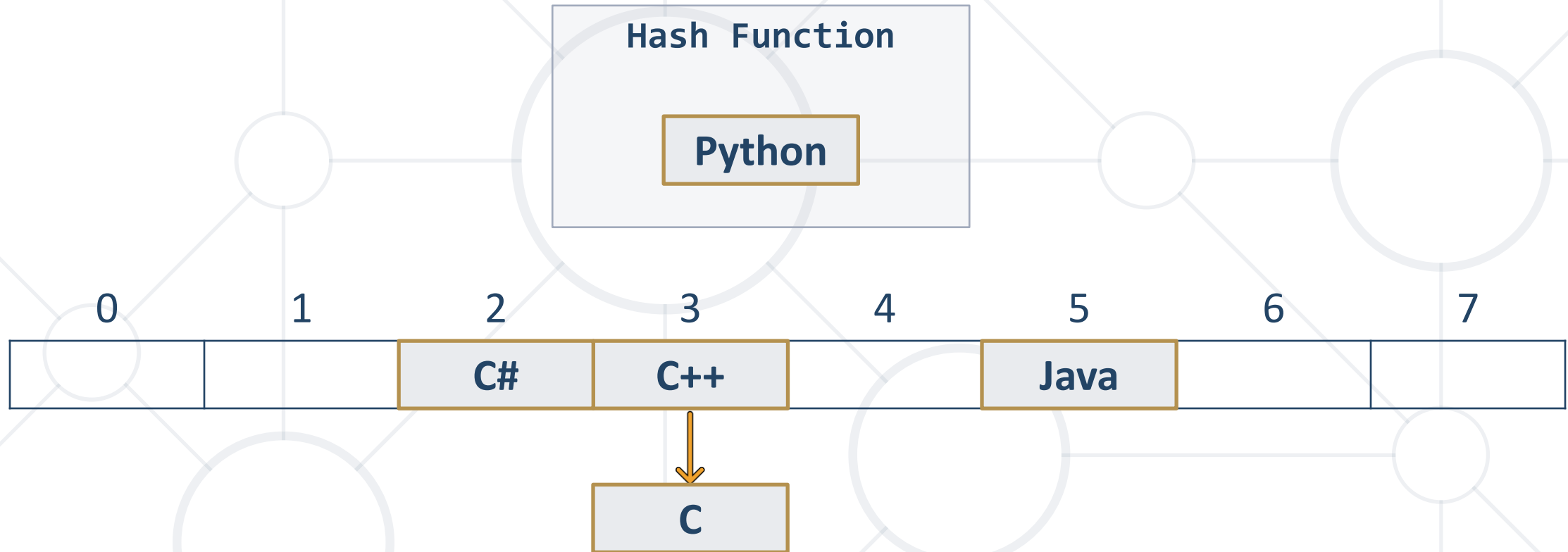
Collision Resolution: Chaining



Collision Resolution: Chaining

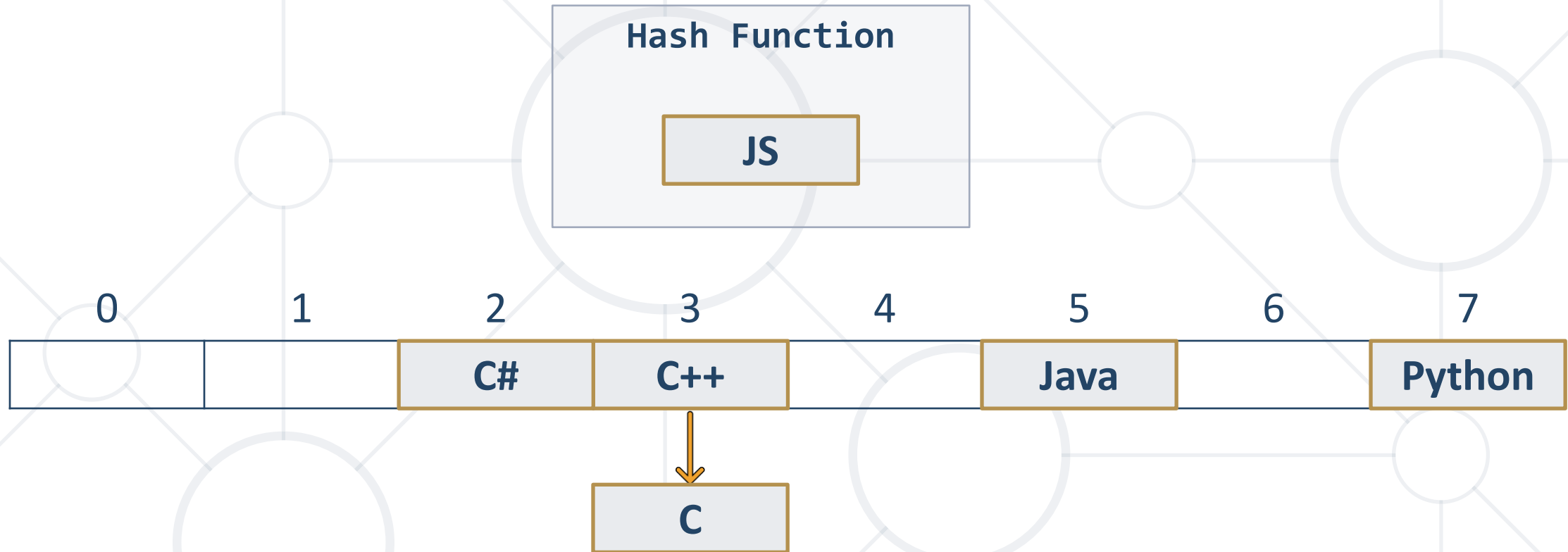


Collision Resolution: Chaining

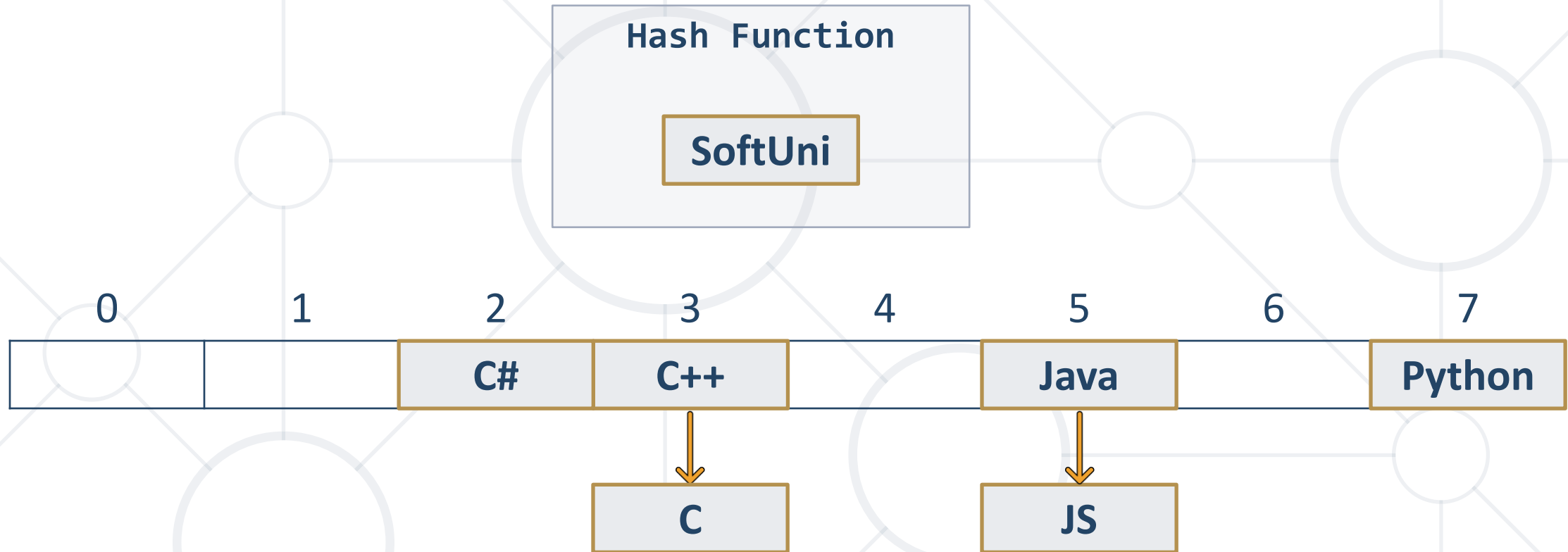


Items are chained
into a linked list

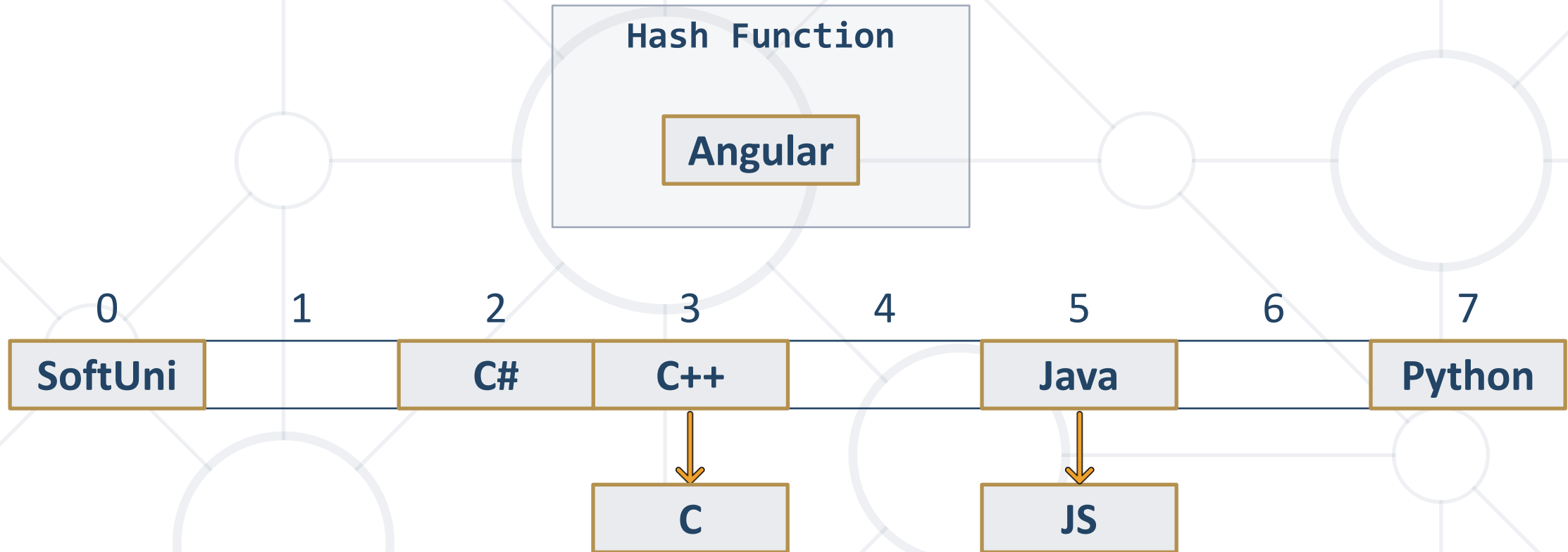
Collision Resolution: Chaining



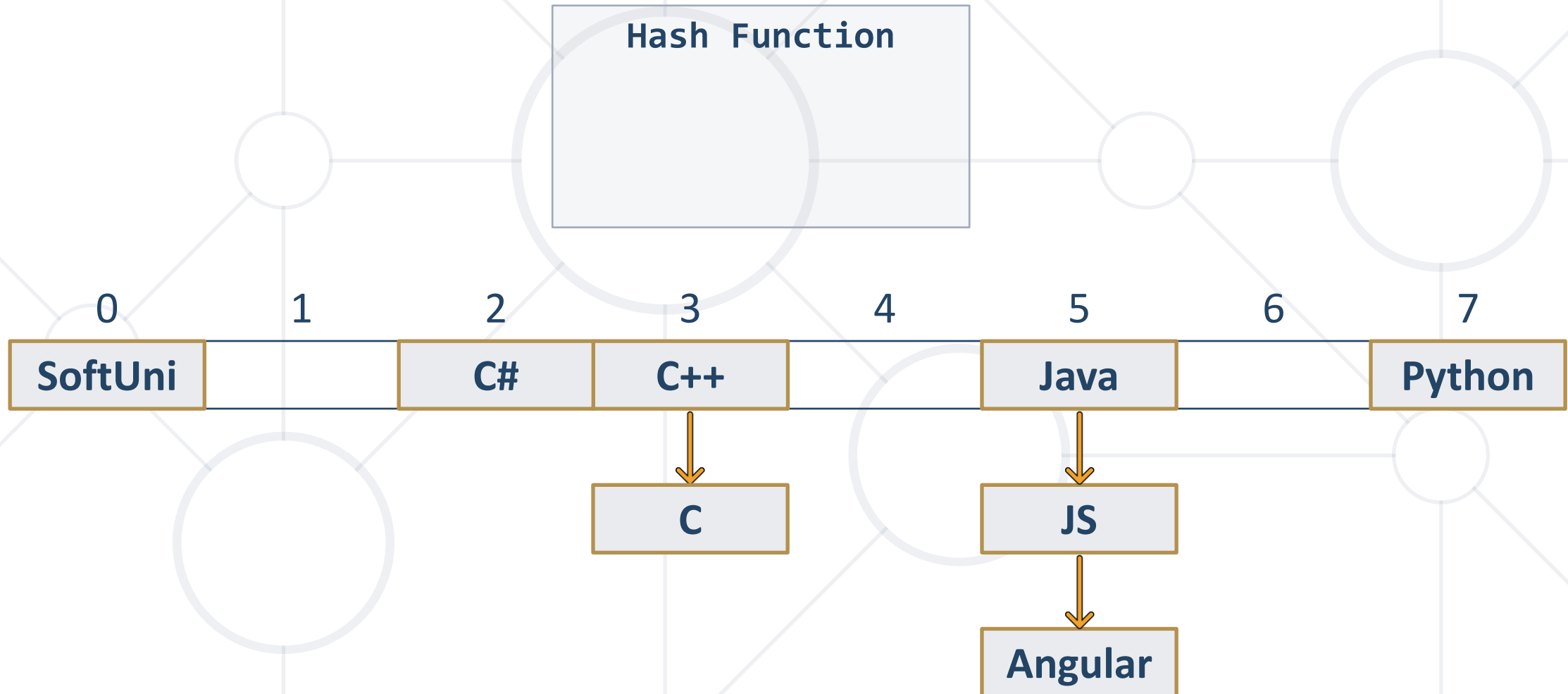
Collision Resolution: Chaining



Collision Resolution: Chaining



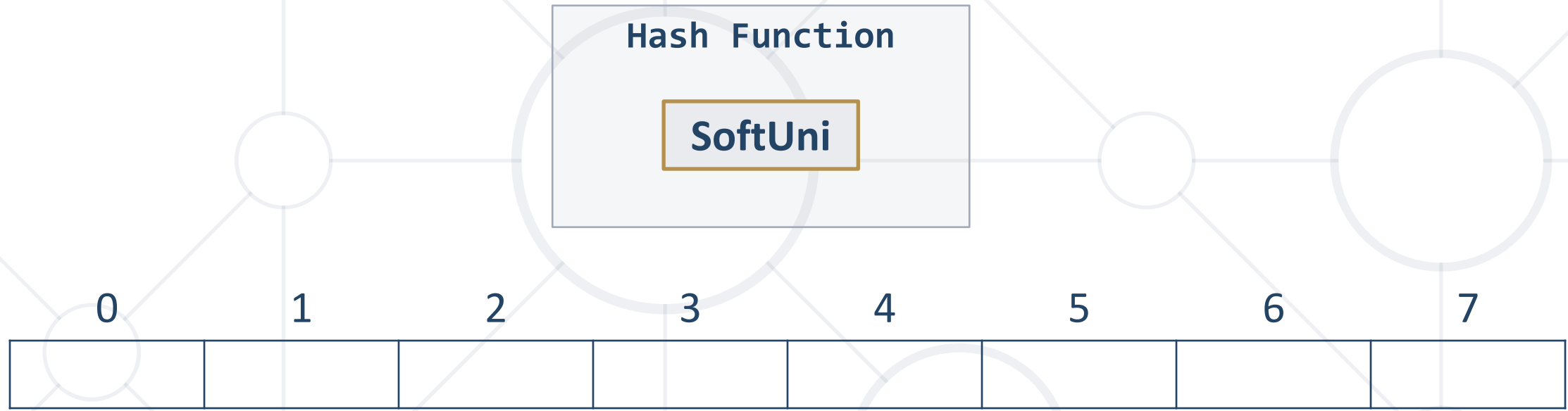
Collision Resolution: Chaining



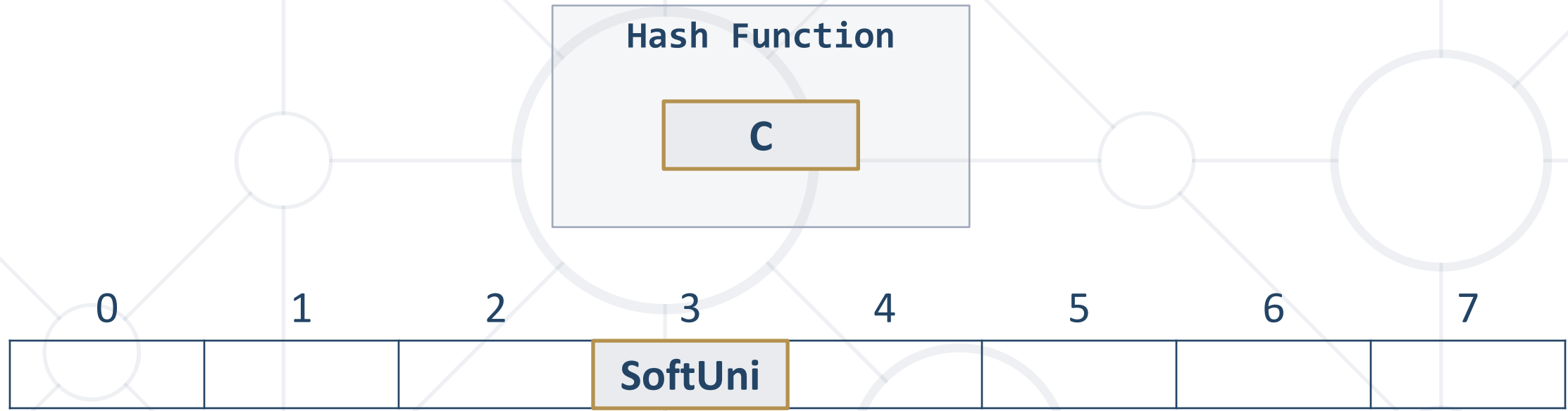
- **Open addressing** as collision resolution strategy means to take another slot in the hash-table in case of collision, e.g.
 - **Linear probing**: take the next empty slot just after the collision
 - $h(\text{key}, i) = h(\text{key}) + i$
 - where i is the attempt number: 0, 1, 2, ...
 - $h(\text{key}) + 1, h(\text{key}) + 2, h(\text{key}) + 3$, etc.

- **Quadratic probing:** the i^{th} next slot is calculated by a quadratic polynomial (c_1 and c_2 are some constants)
 - $h(\text{key}, i) = h(\text{key}) + c_1 * i + c_2 * i^2$
 - $h(\text{key}) + 1^2, h(\text{key}) + 2^2, h(\text{key}) + 3^2$, etc.
- **Re-hashing:** use separate (second) hash-function for collisions
 - $h(\text{key}, i) = h_1(\text{key}) + i * h_2(\text{key})$

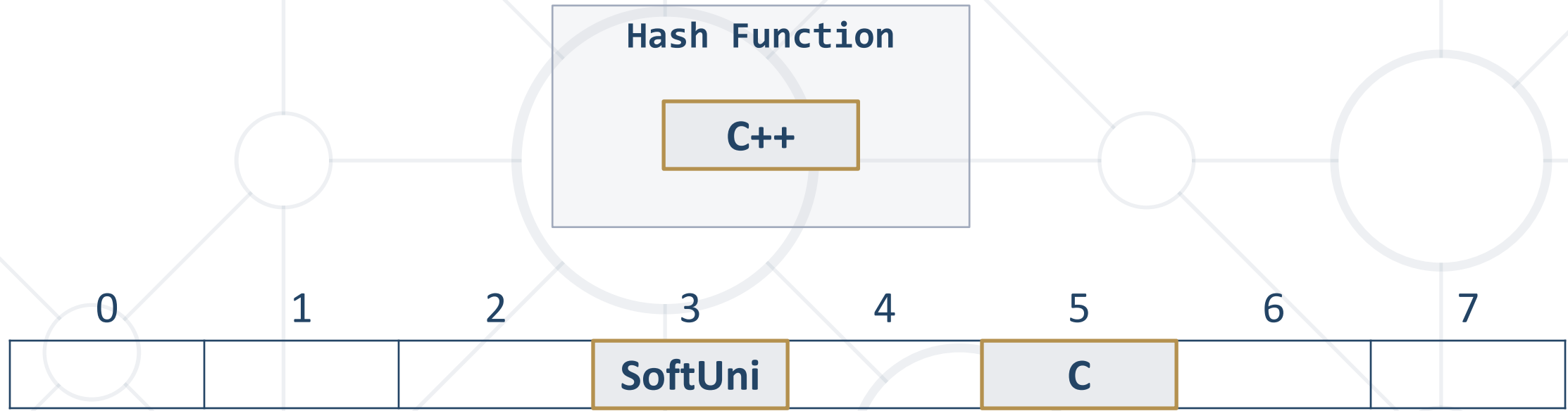
Collision Resolution: Linear Probing



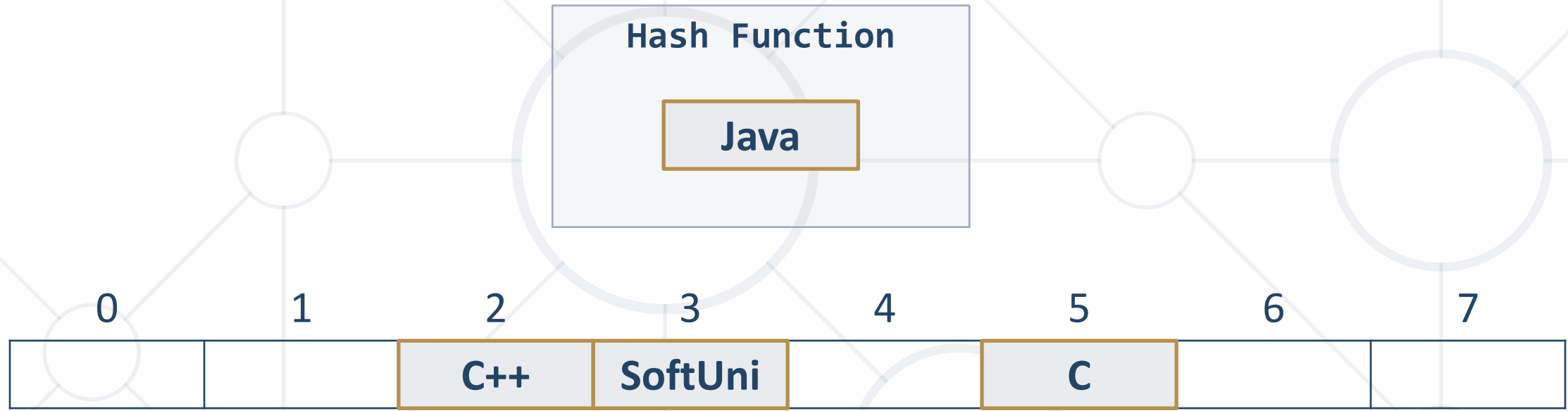
Collision Resolution: Linear Probing



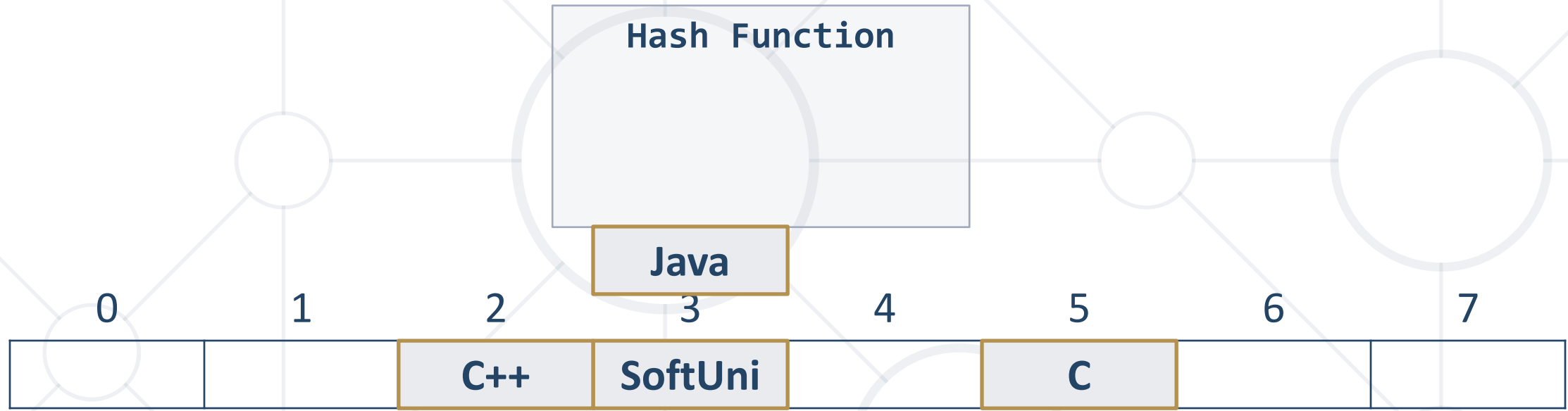
Collision Resolution: Linear Probing



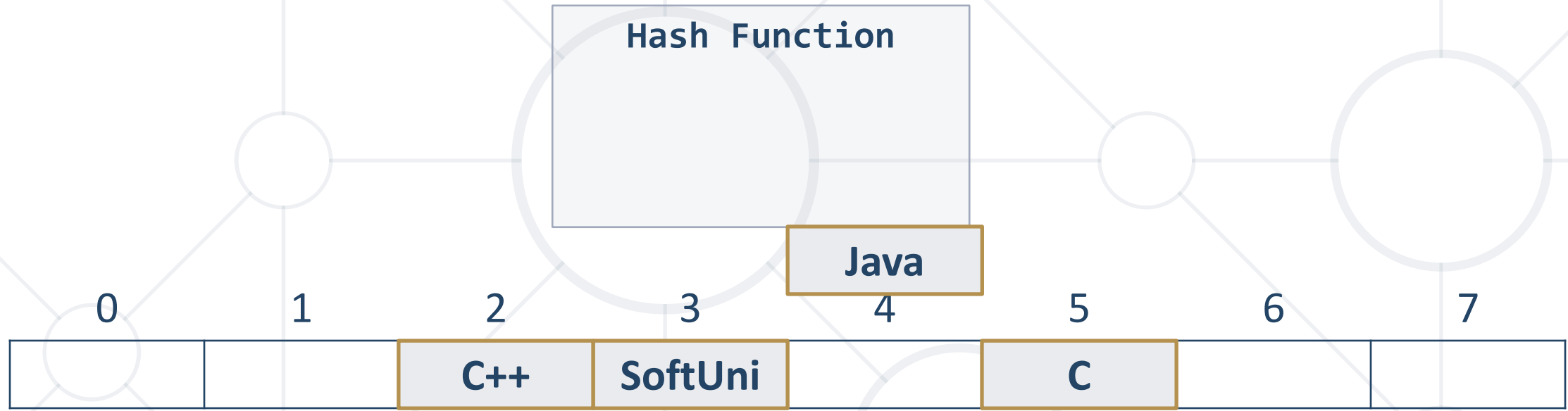
Collision Resolution: Linear Probing



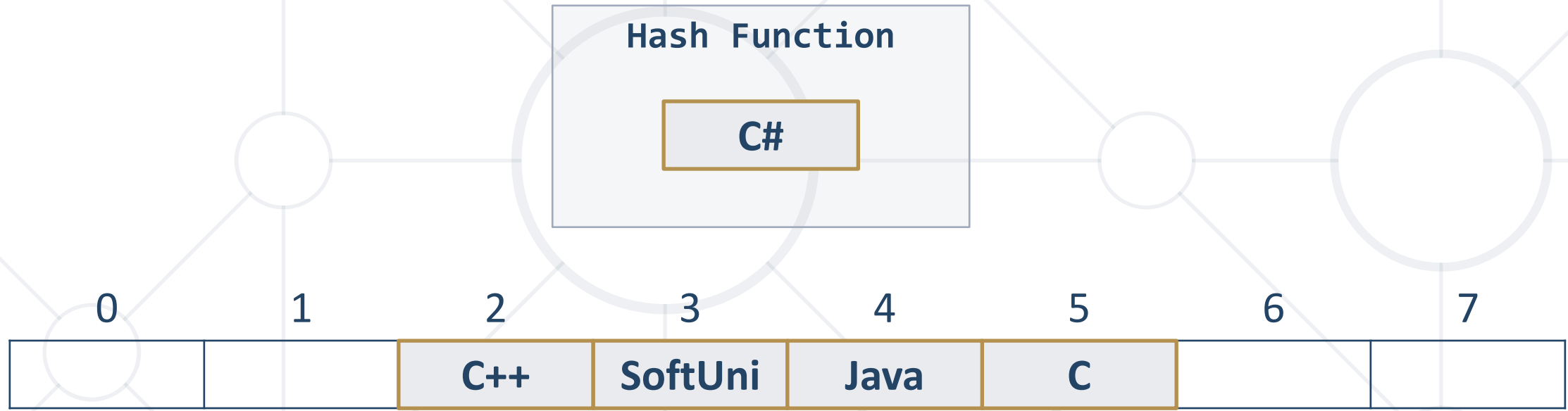
Collision Resolution: Linear Probing



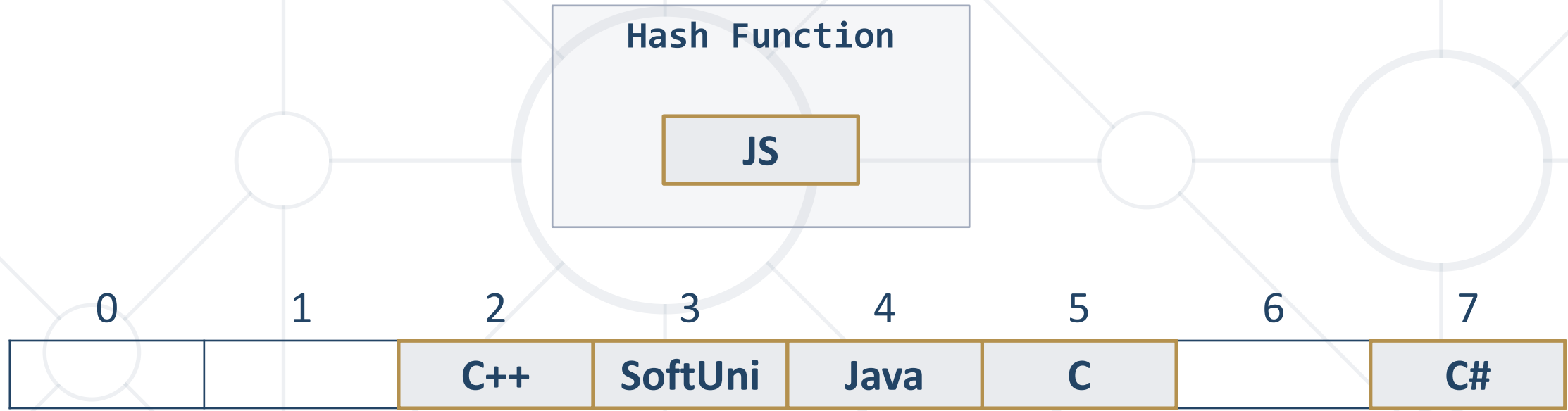
Collision Resolution: Linear Probing



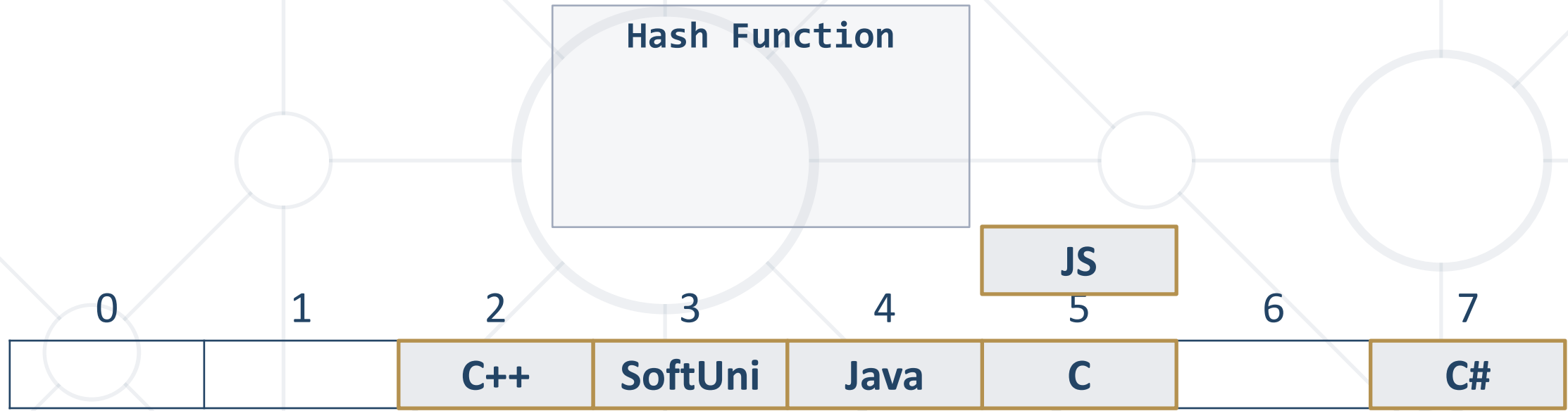
Collision Resolution: Linear Probing



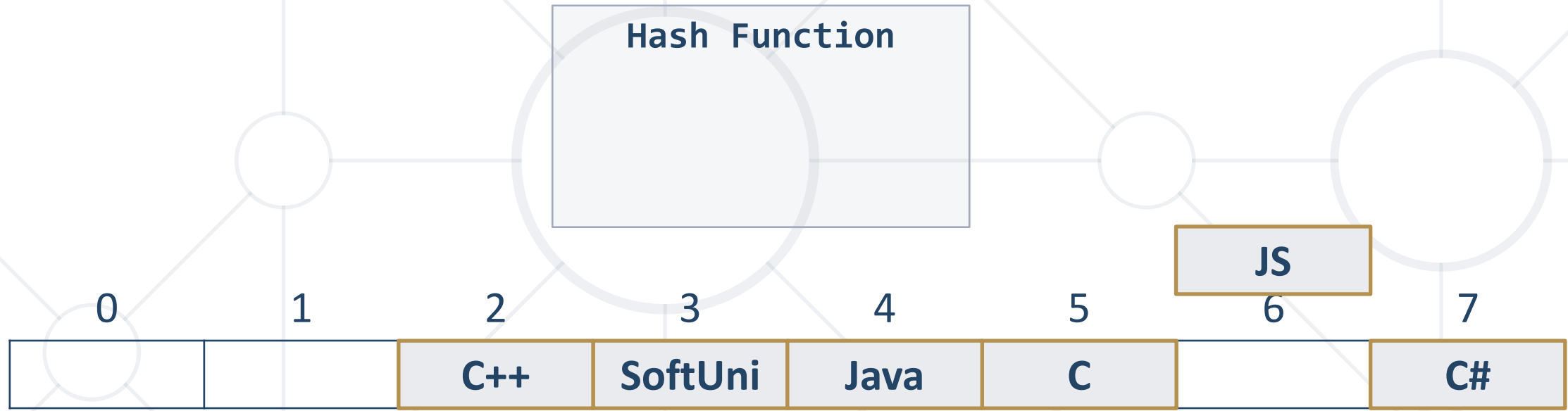
Collision Resolution: Linear Probing



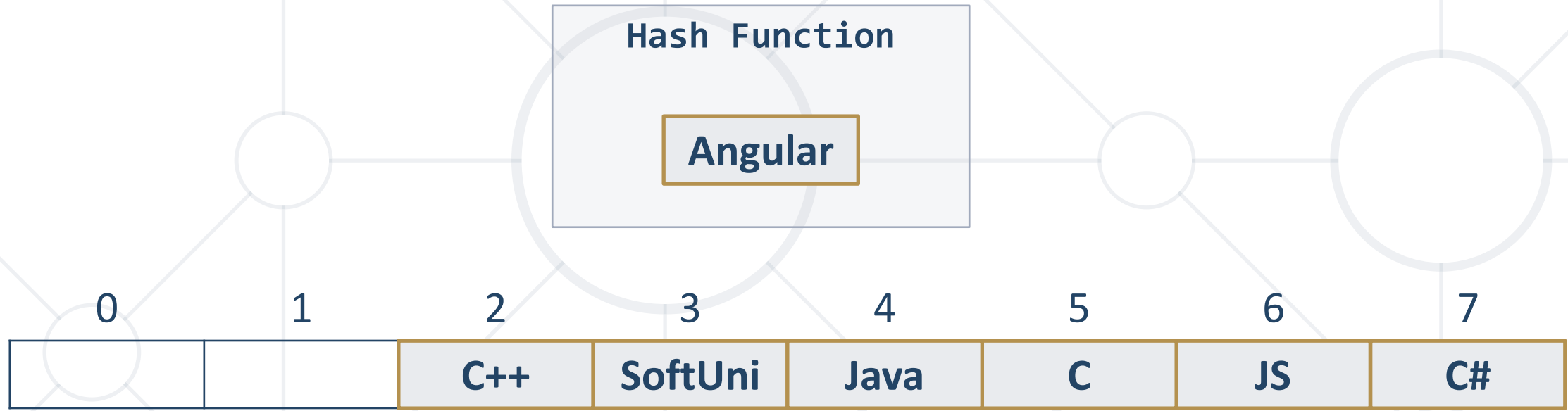
Collision Resolution: Linear Probing



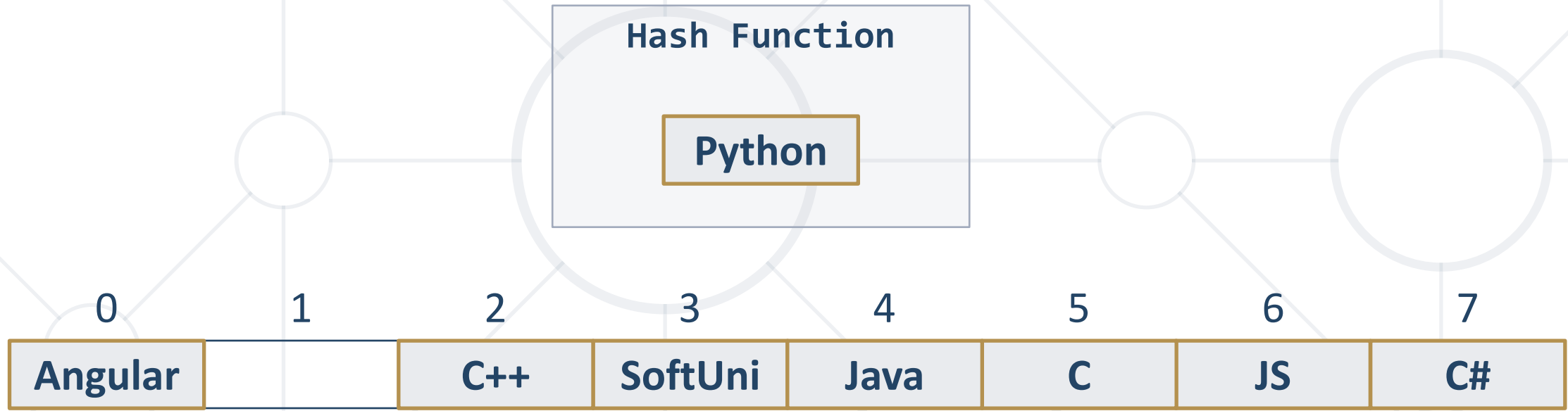
Collision Resolution: Linear Probing



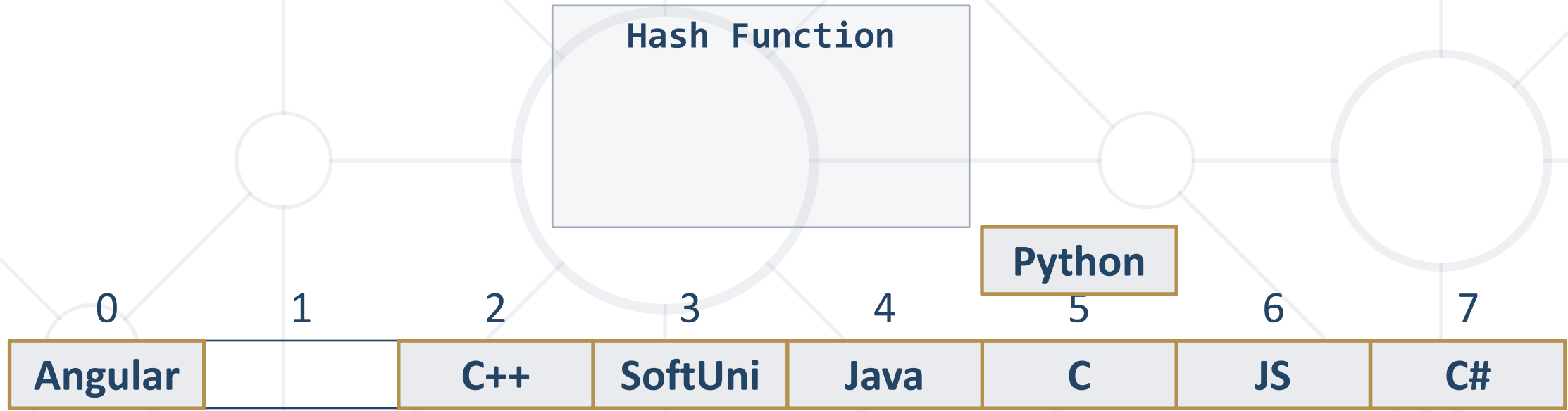
Collision Resolution: Linear Probing



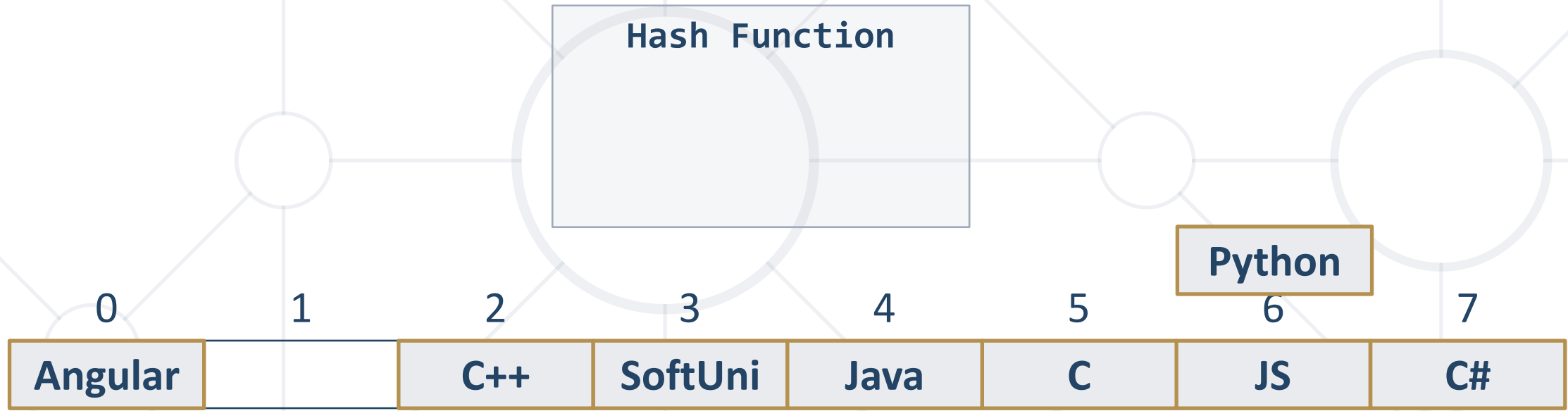
Collision Resolution: Linear Probing



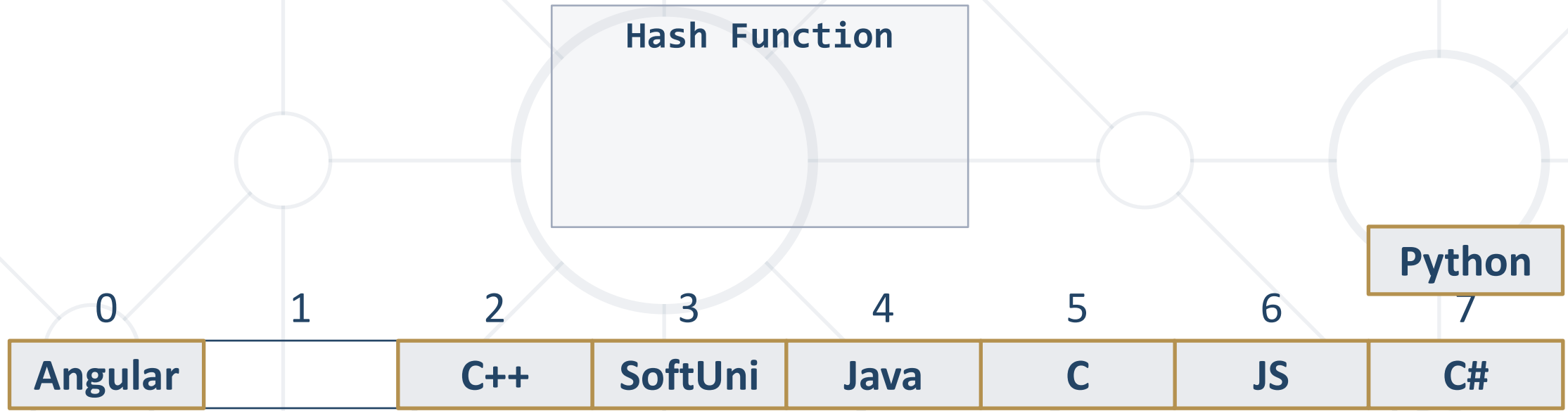
Collision Resolution: Linear Probing



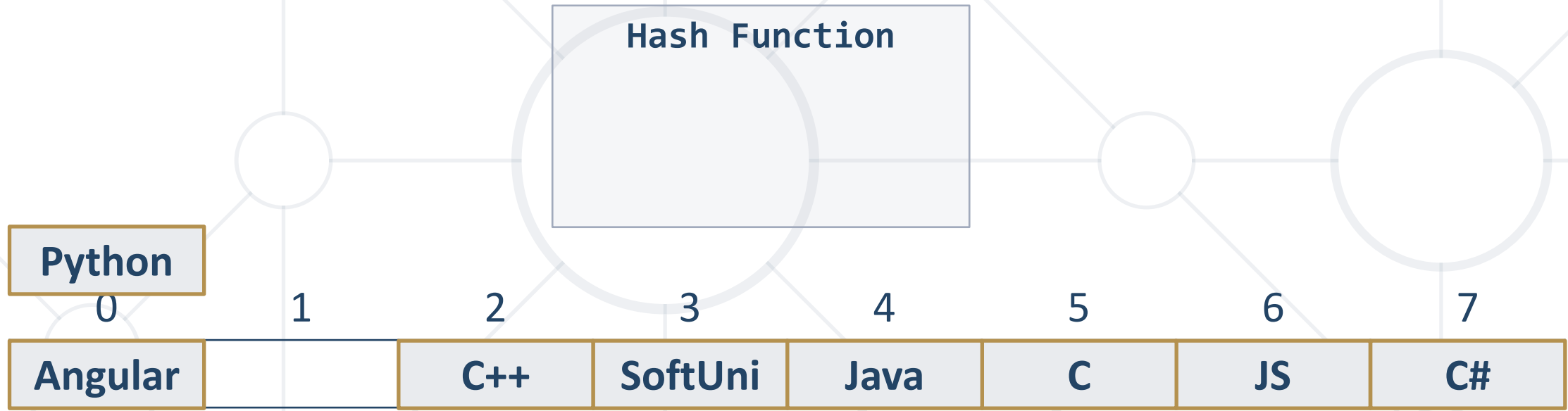
Collision Resolution: Linear Probing



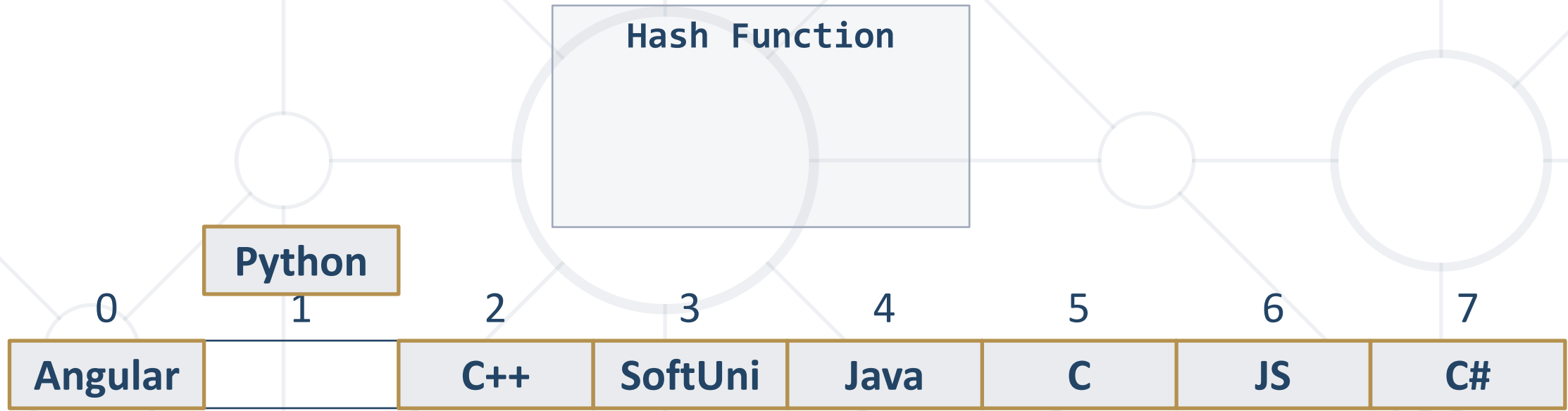
Collision Resolution: Linear Probing



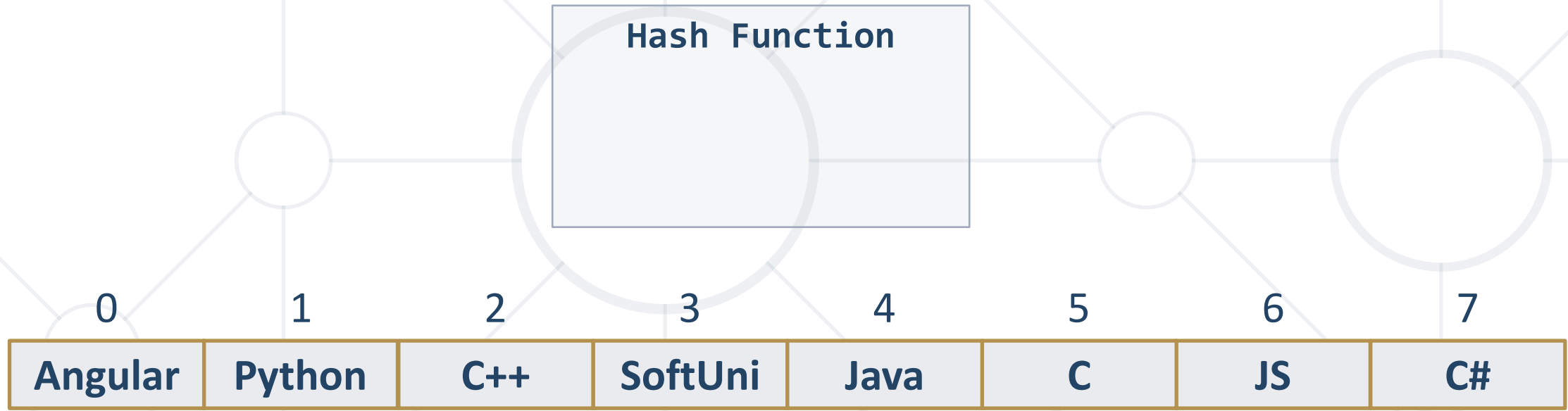
Collision Resolution: Linear Probing



Collision Resolution: Linear Probing



Collision Resolution: Linear Probing



TIME'S

- What is the average running time of delete in linear-probing hash table? Your hash function satisfies the uniform hashing assumption and that the hash table is at most 50% full.
 - $O(1)$
 - $O(\log N)$
 - $O(N)$
 - $O(N \log N)$

- What is the average running time of delete in linear-probing hash table? Your hash function satisfies the uniform hashing assumption and that the hash table is at most 50% full.
 - $O(1)$ ✓
 - $O(\log N)$
 - $O(N)$
 - $O(N \log N)$

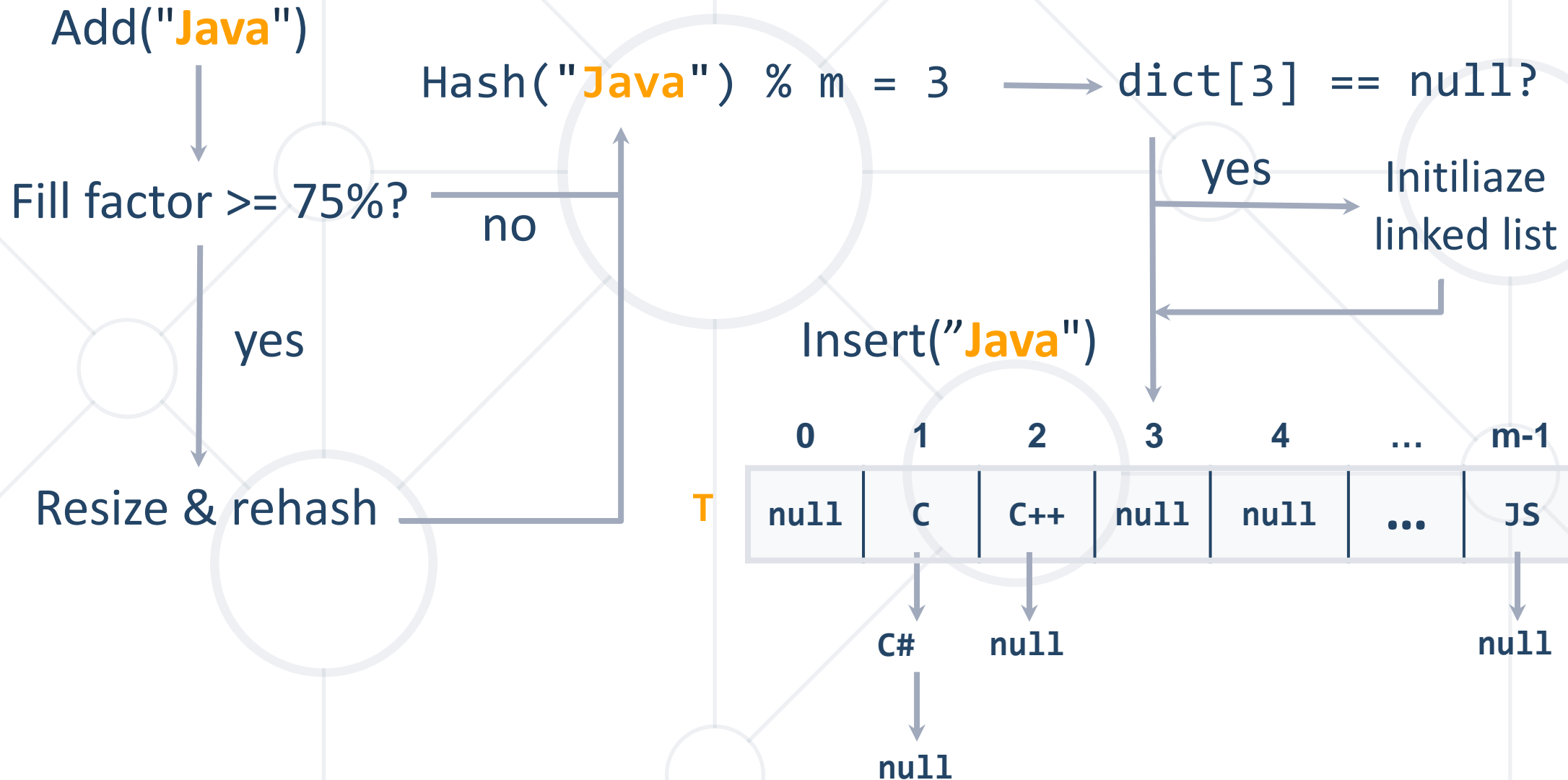
- The hash-table performance depends on the probability of collisions - Less collisions = faster add / find / delete operations
 - **Collisions resolution** algorithm
 - **Fill factor** (used buckets / all buckets)

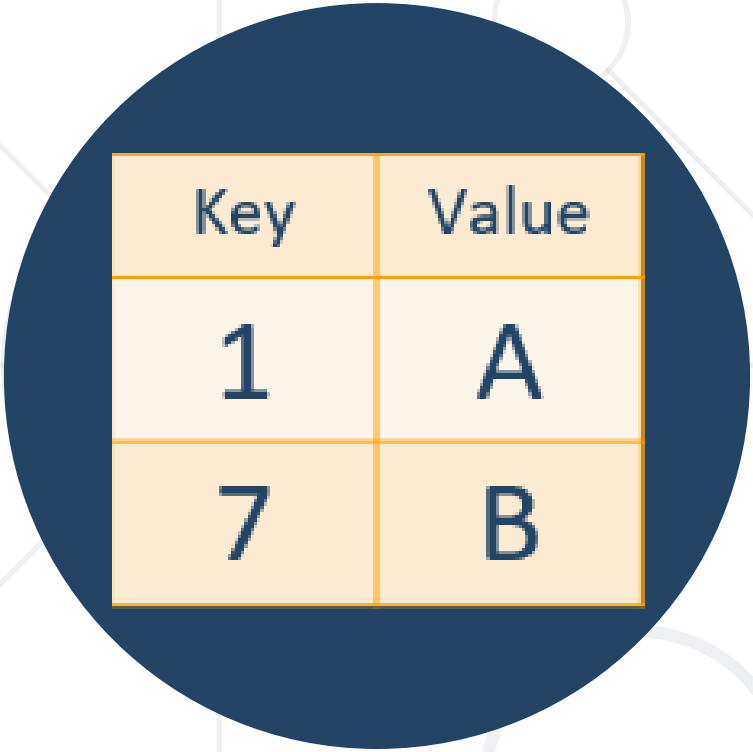
- **Add / Find / Delete** take just few primitive operations
 - Speed does not depend on the size of the hash-table
 - Amortized complexity **$O(1)$** – constant time
- Example:
 - Finding an element in a **hash-table** holding **1 000 000 elements** takes average just **1-2 steps**
 - Finding an element in an **array** holding **1 000 000 elements** takes average **500 000 steps**

How Big the Hash-Table Should Be?

- The **load factor** (fill factor) = **used cells / all cells**
 - How much the hash table is filled, e.g. 65%
- Smaller fill factor leads to less collisions (faster average seek time)
- Recommended fill factors:
 - When **chaining** is used as collision resolution less than **75%**
 - When **open addressing** is used less than **50%**

Adding Item to Hash Table with Chaining

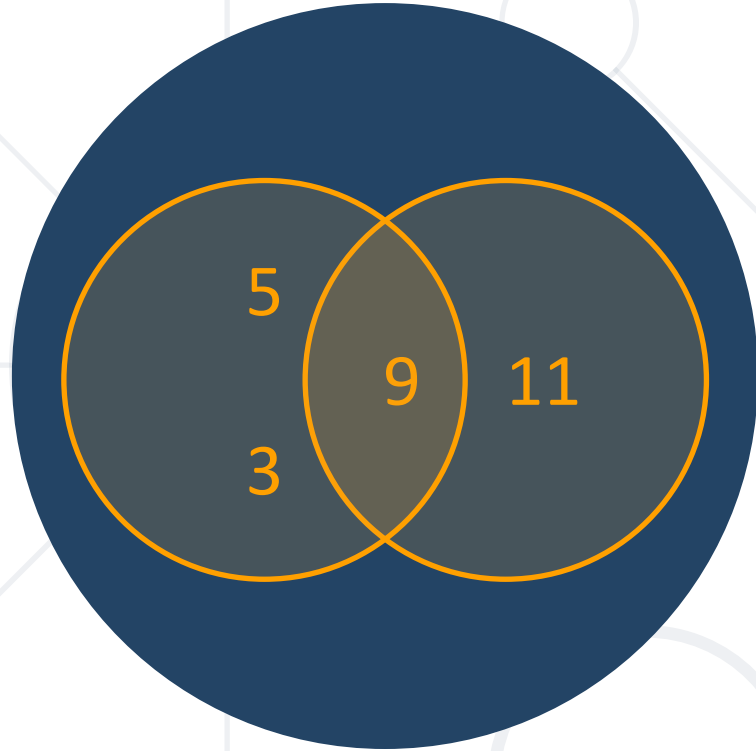




Key	Value
1	A
7	B

Lab Exercise

Implement a Hash-Table with Chaining



Sets and Bags

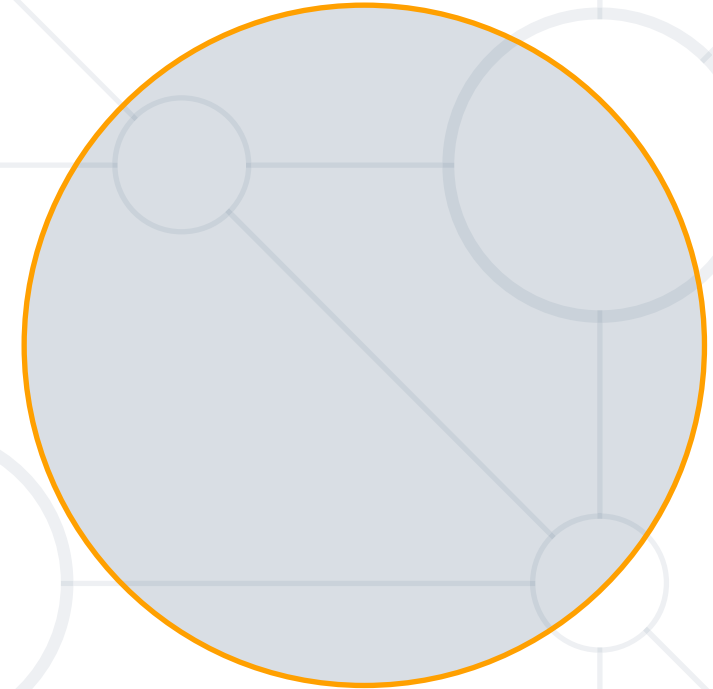
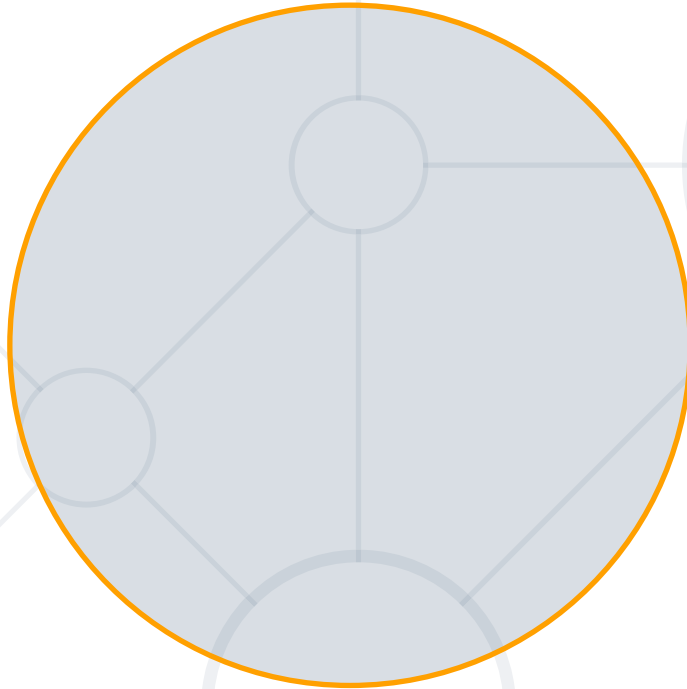
Set Operations

- The abstract data type (ADT) "**Set**" keeps a set of elements with no duplicates
- Sets with duplicates are also known as ADT "**Bag**"
- Set specific operations:
 - **UnionWith(set)**
 - **IntersectWith(set)**
 - **ExceptWith(set)**
 - **SymmetricExceptWith(set)**

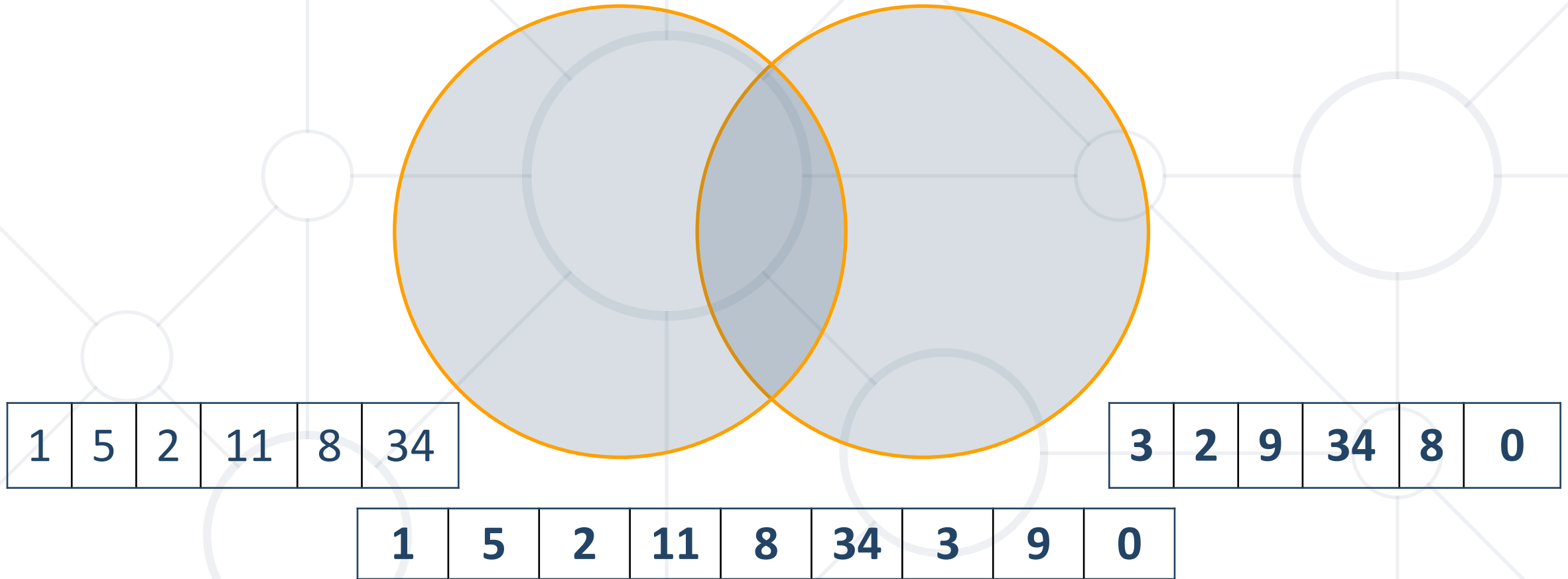
Known as relative complement in math

Known as symmetric difference

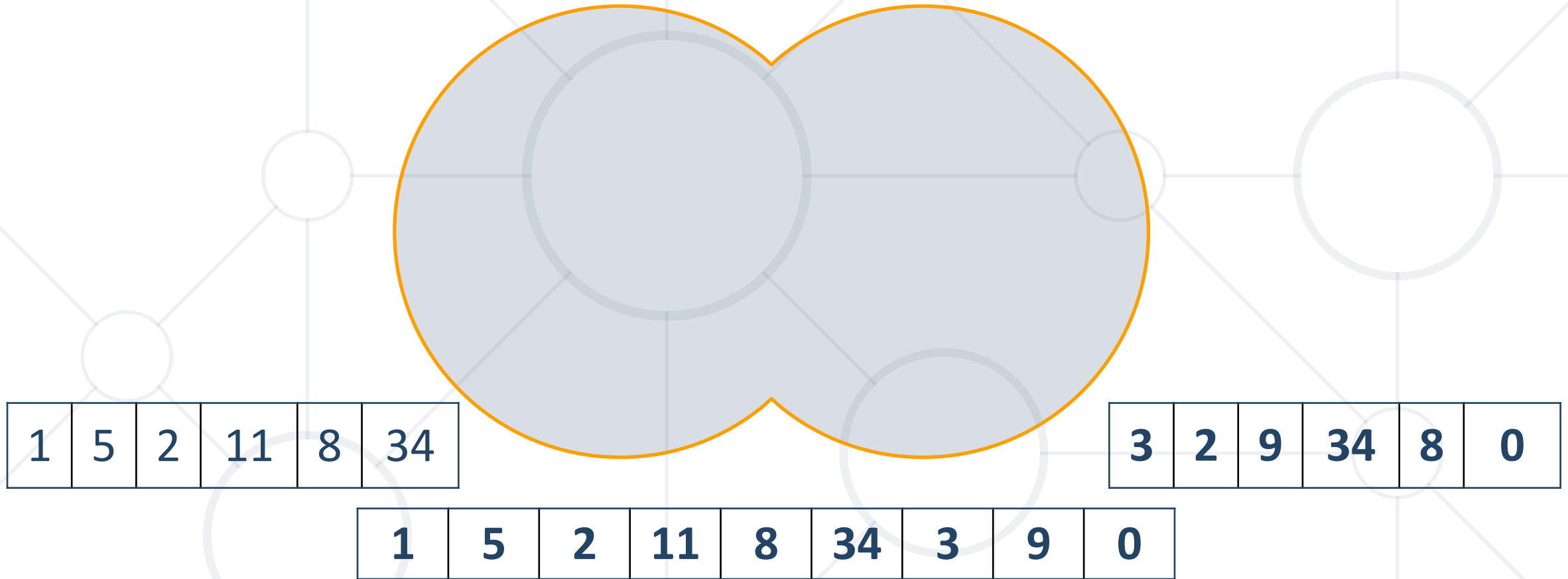
Union



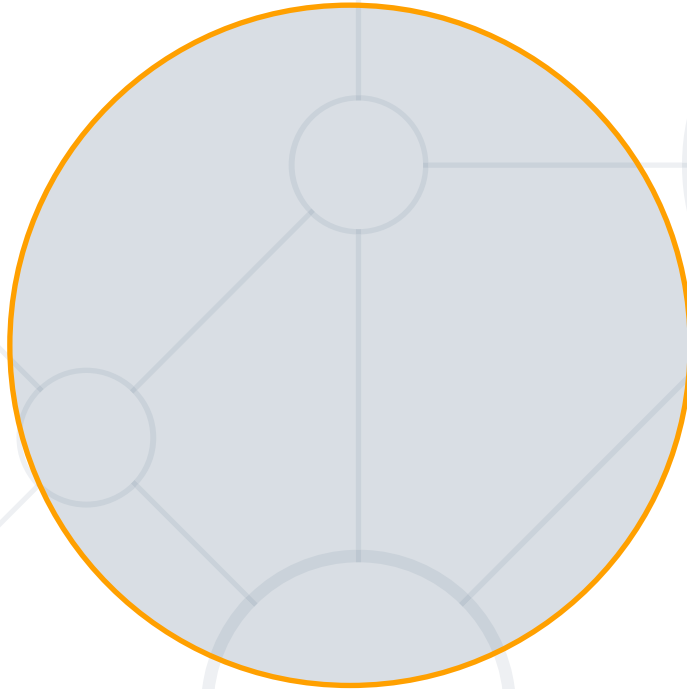
Union



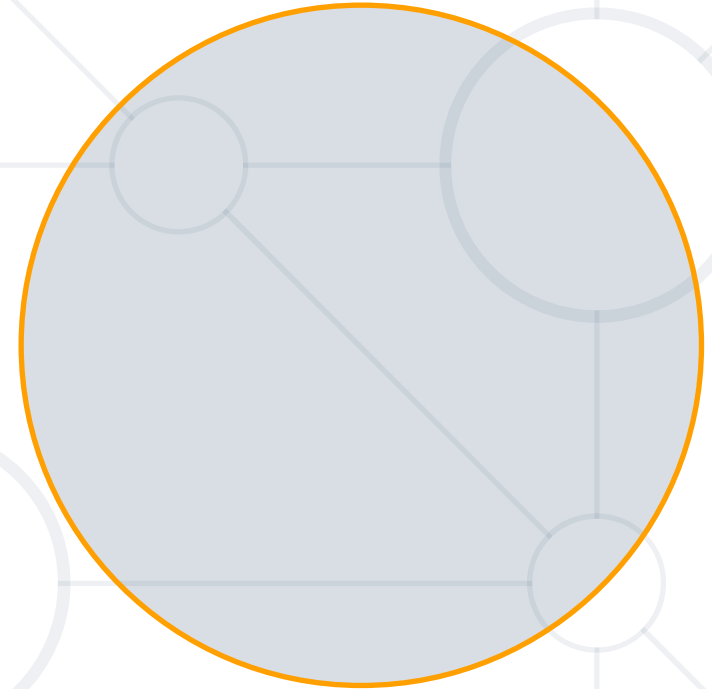
Union



Intersects

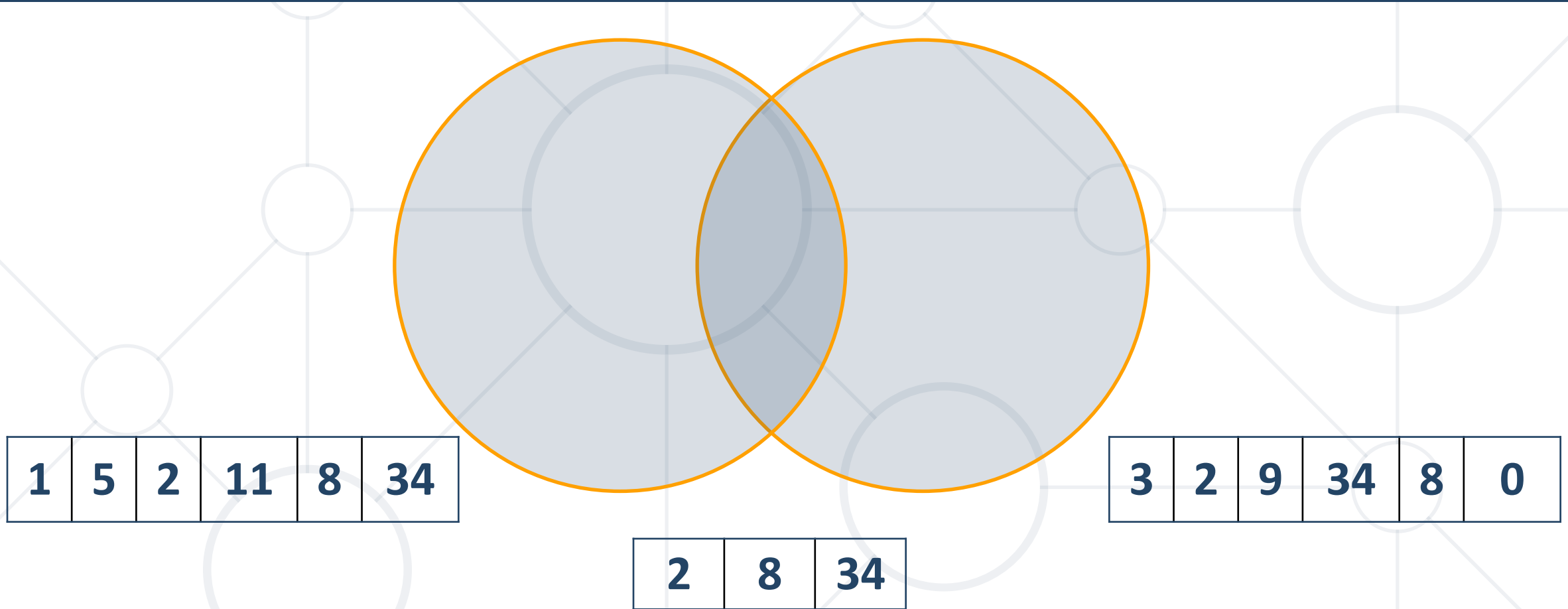


1	5	2	11	8	34
---	---	---	----	---	----



3	2	9	34	8	0
---	---	---	----	---	---

Intersects

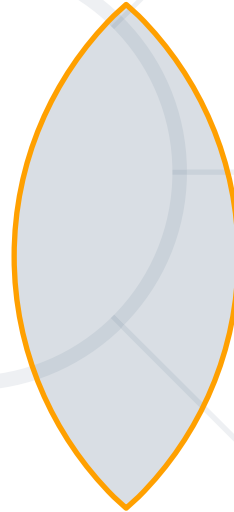


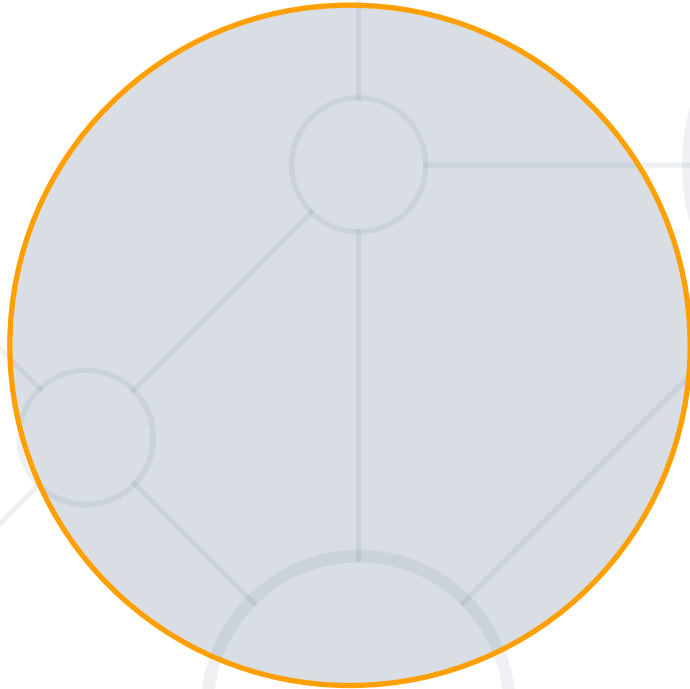
Intersects

1	5	2	11	8	34
---	---	---	----	---	----

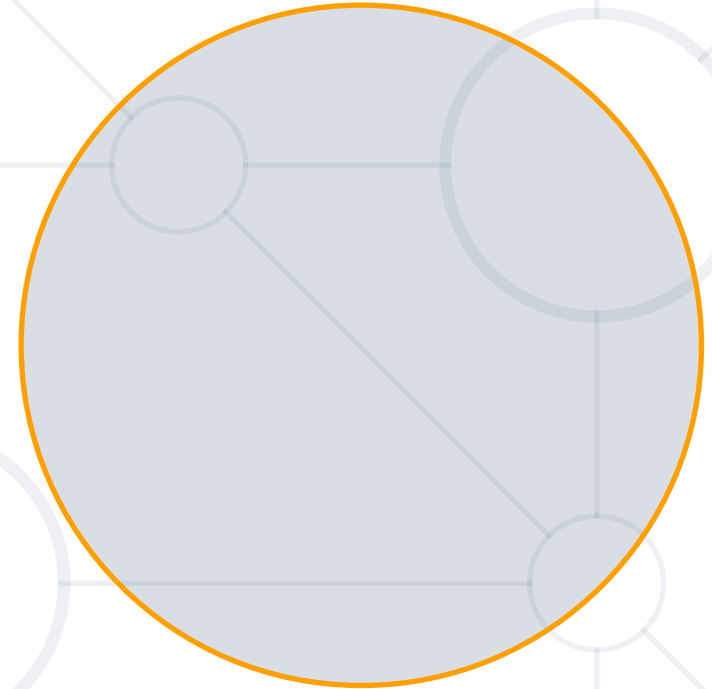
2	8	34
---	---	----

3	2	9	34	8	0
---	---	---	----	---	---

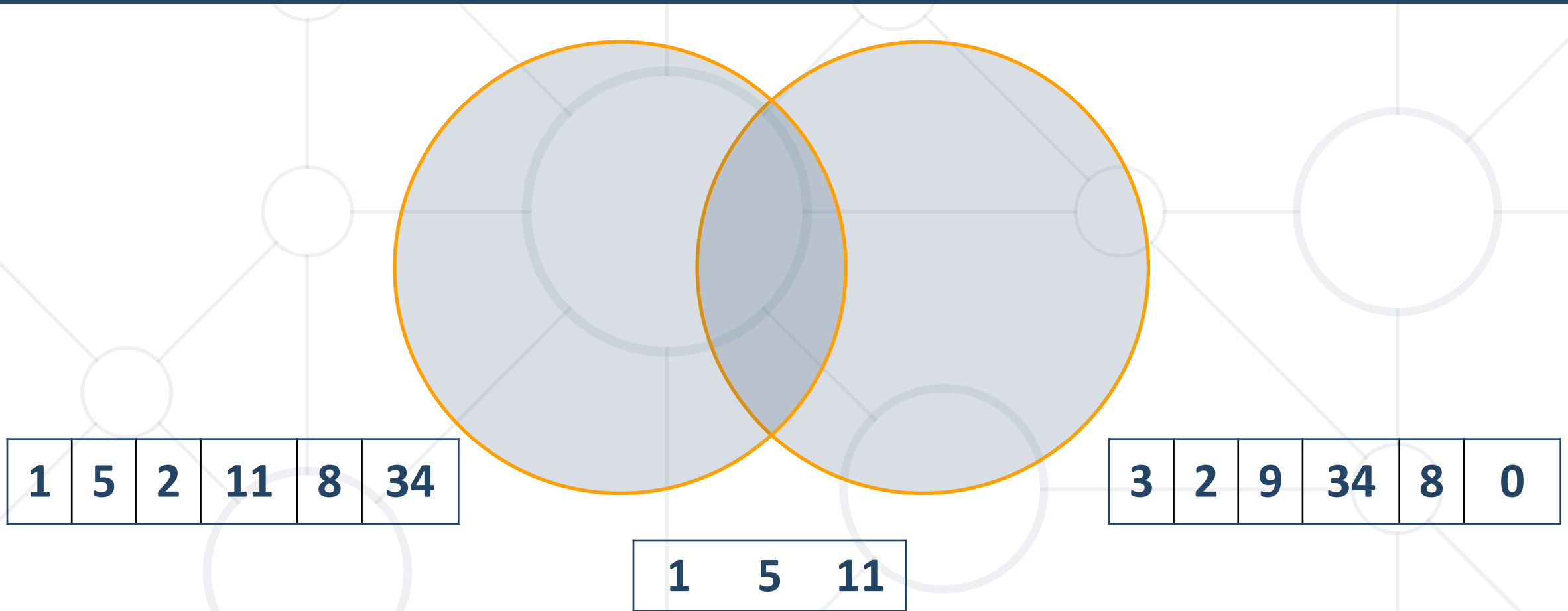




1	5	2	11	8	34
---	---	---	----	---	----



3	2	9	34	8	0
---	---	---	----	---	---

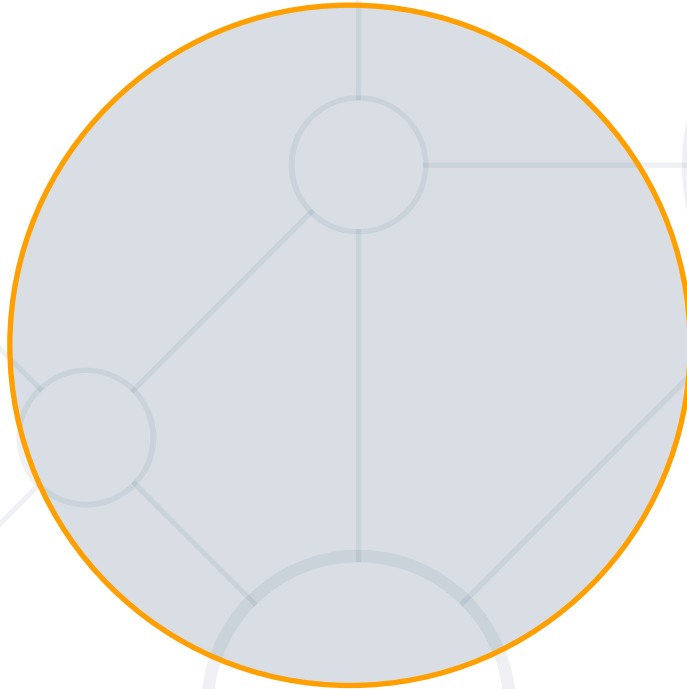


1	5	2	11	8	34
---	---	---	----	---	----

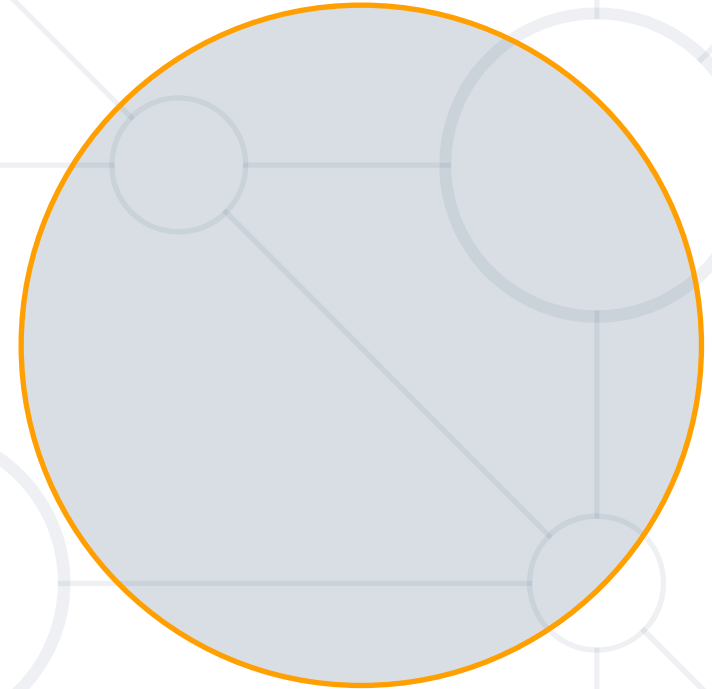
1	5	11
---	---	----

3	2	9	34	8	0
---	---	---	----	---	---

Symmetric Except

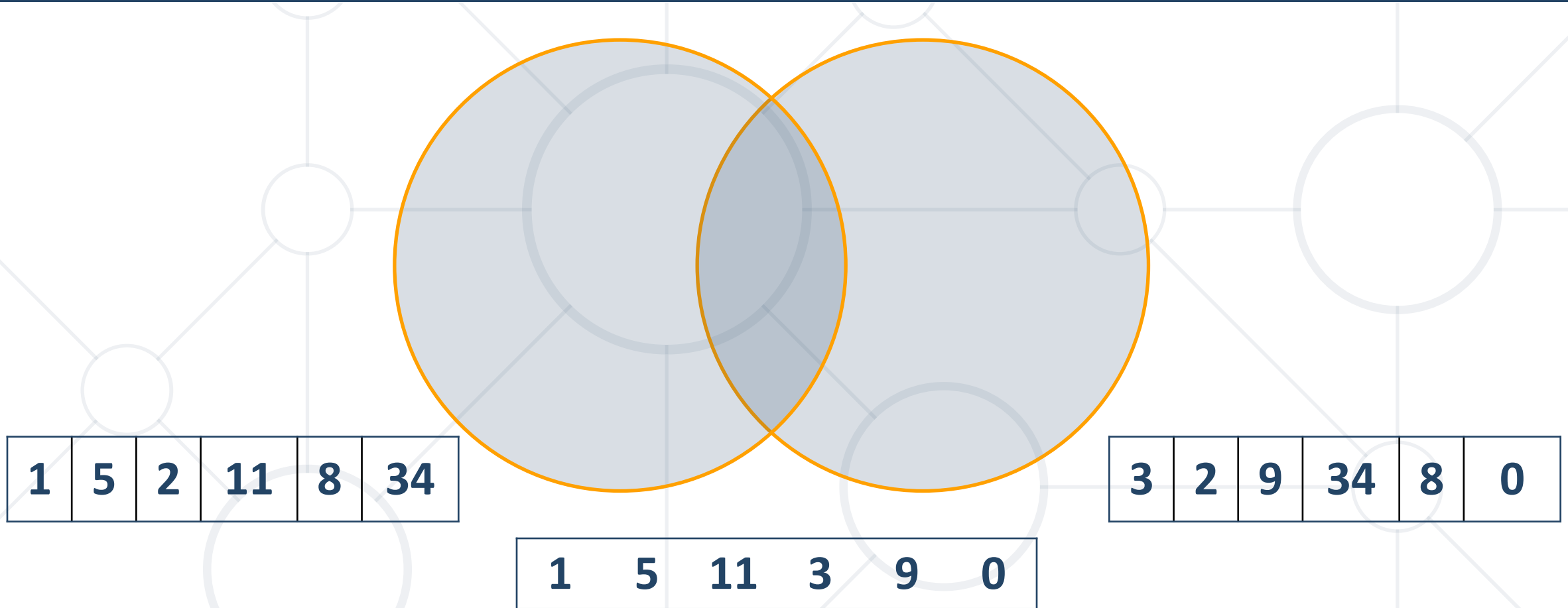


1	5	2	11	8	34
---	---	---	----	---	----

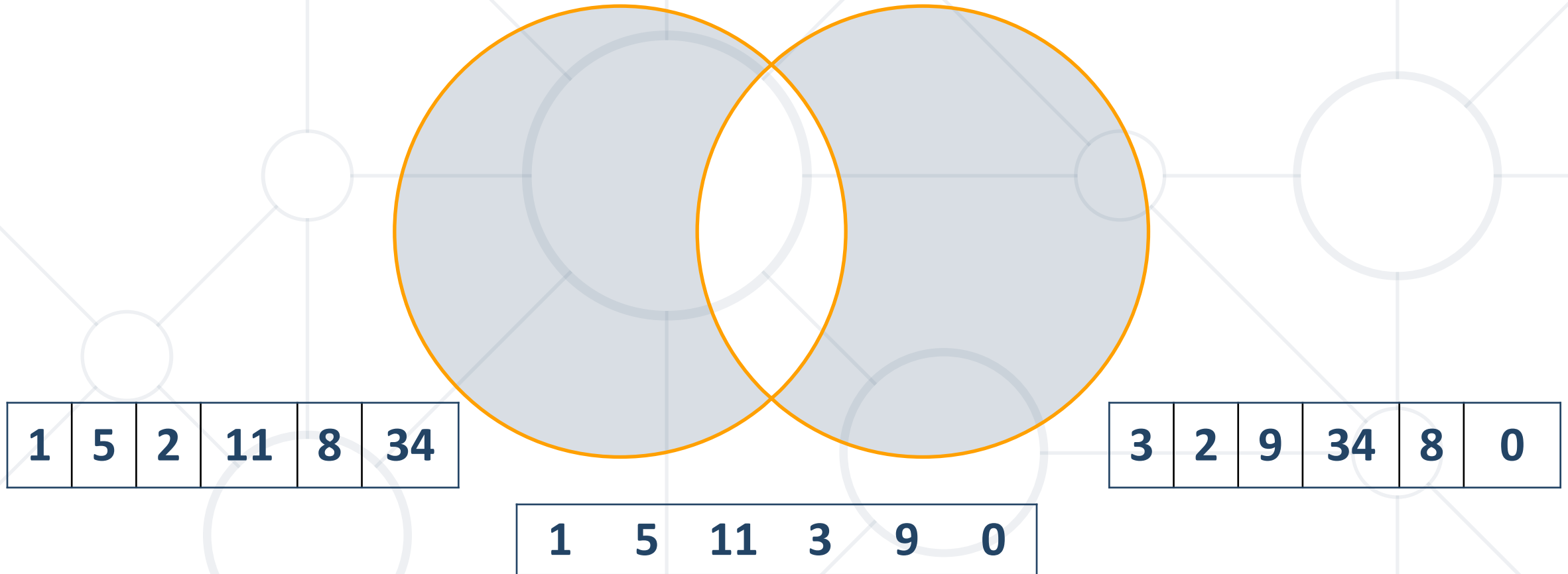


3	2	9	34	8	0
---	---	---	----	---	---

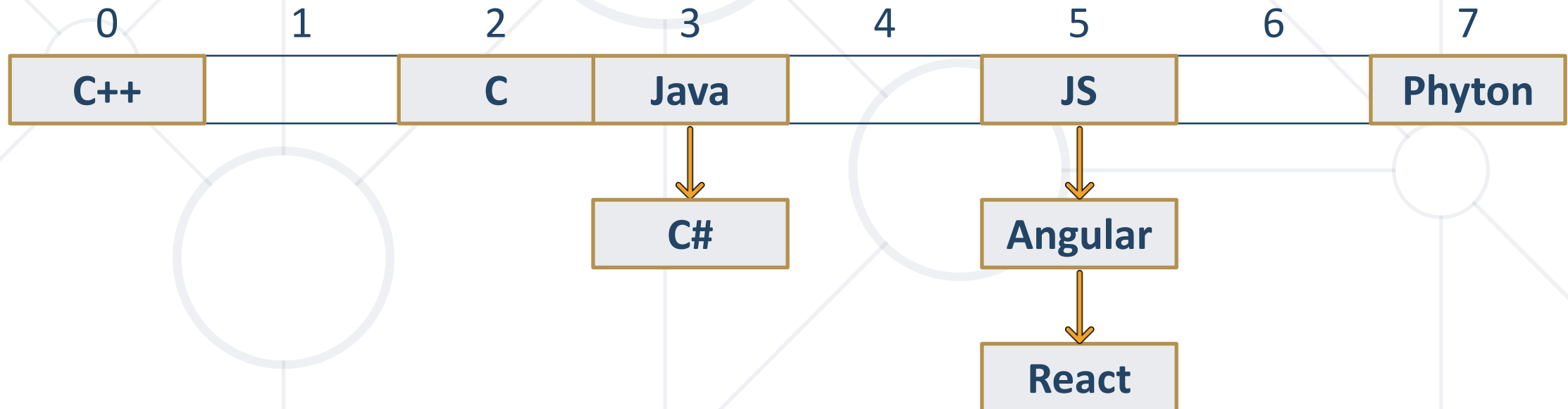
Symmetric Except



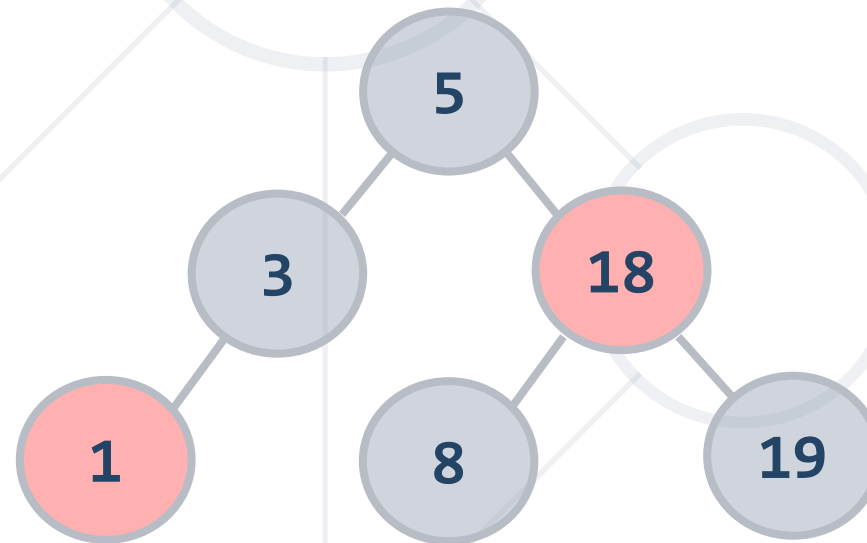
Symmetric Except



- **HashSet<T>** implements ADT **Set** by hash table
 - Elements are in no particular order
- All major operations are fast: **Add / Delete / Contains**




- **SortedSet<T>** implements ADT **Set** by balanced search tree (red-black tree)
 - Elements are sorted in increasing order



TIME'S

- For given sets - $\{1, 2, 3, 4, 5\}$ and $\{3, 4, 5, 6, 7\}$, what is the operation that will give us the following result: $\{1, 2, 6, 7\}$
 - Union
 - Intersects
 - Except
 - SymmetricExcept

- For given sets - $\{1, 2, 3, 4, 5\}$ and $\{3, 4, 5, 6, 7\}$, what is the operation that will give us the following result: $\{1, 2, 6, 7\}$
 - Union
 - Intersects
 - Except
 - SymmetricExcept 



Dictionaries

Definition and Operations

The Dictionary (Map) ADT

- The abstract data type (ADT) "**dictionary**" maps key to values
 - Also known as "**map**" or "**associative array**"
 - Holds a set of **{key, value} pairs**
- Many implementations
 - Hash table, balanced tree, list, array, ...

key	value
John Smith	+1-555-8976
Sam Doe	+1-555-5030

ADT Map – Example

- Sample dictionary:

Key	Value
C#	Modern general-purpose object-oriented programming language
PHP	Popular server-side scripting language for Web development
compiler	Software that transforms a computer program to executable machine code
...	...

- Major operations:
 - **Add(key, value)** – adds an element by key + value
 - **Remove(key)** – removes a value by key
 - **this[key] = value** – add / replace element by key
 - **this[key]** – returns the value by key
 - **Keys** – returns a collection of all keys (in order of entry)
 - **Values** – returns a collection of all values (in order of entry)

- Major operations:
 - **ContainsKey(key)** – checks if given key exists in the dictionary
 - **ContainsValue(value)** – checks whether the dictionary contains given value
 - Warning: slow operation – **$O(n)$**
 - **TryGetValue(key, outvalue)**
 - if the key is found, returns it in the value
 - otherwise returns false

- **SortedDictionary<Key, Value>** implements the ADT "dictionary" as self-balancing search tree
 - Elements are arranged in the tree ordered by key
 - Traversing the tree returns the elements in increasing order
 - **Add / Find / Delete** perform **log N** operations
- Use **SortedDictionary<Key, Value>** when you need the elements sorted by key
 - Otherwise use **Dictionary<Key, Value>** – it has better performance


TIME'S

- Which built-in implementation of **IDictionary<Key, Value>** sorts the items by value?
 - Dictionary<Key, Value>
 - SortedDictionary<Key, Value>
 - None

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 - Dictionary<Key, Value>
 - SortedDictionary<Key, Value>
 - **None** 

TIME'S

- Which is the main reason to use a hash table instead of a red-black BST?
 - Supports more operations efficiently
 - Better worst-case performance guarantee
 - Better performance in practice on typical inputs

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Comparing Keys

Using Custom Key Classes

- **Dictionary<Key, Value>** relies on
 - **Object.Equals()** – for comparing the keys
 - **Object.GetHashCode()** – for calculating the hash codes of the keys
- **SortedDictionary<Key, Value>** relies on **IComparable<Key>** for ordering the keys

Implementing Equals() and GetHashCode()

```
public class Point {  
    public int x, y;  
    public override bool Equals(Object obj) {  
        if (!obj is Point) || (obj == null) return false;  
        Point p = (Point) obj;  
        return (x == p.x) && (y == p.y);  
    }  
  
    public int GetHashCode() {  
        return (x << 16 | y >> 16) ^ y;  
    }  
}
```

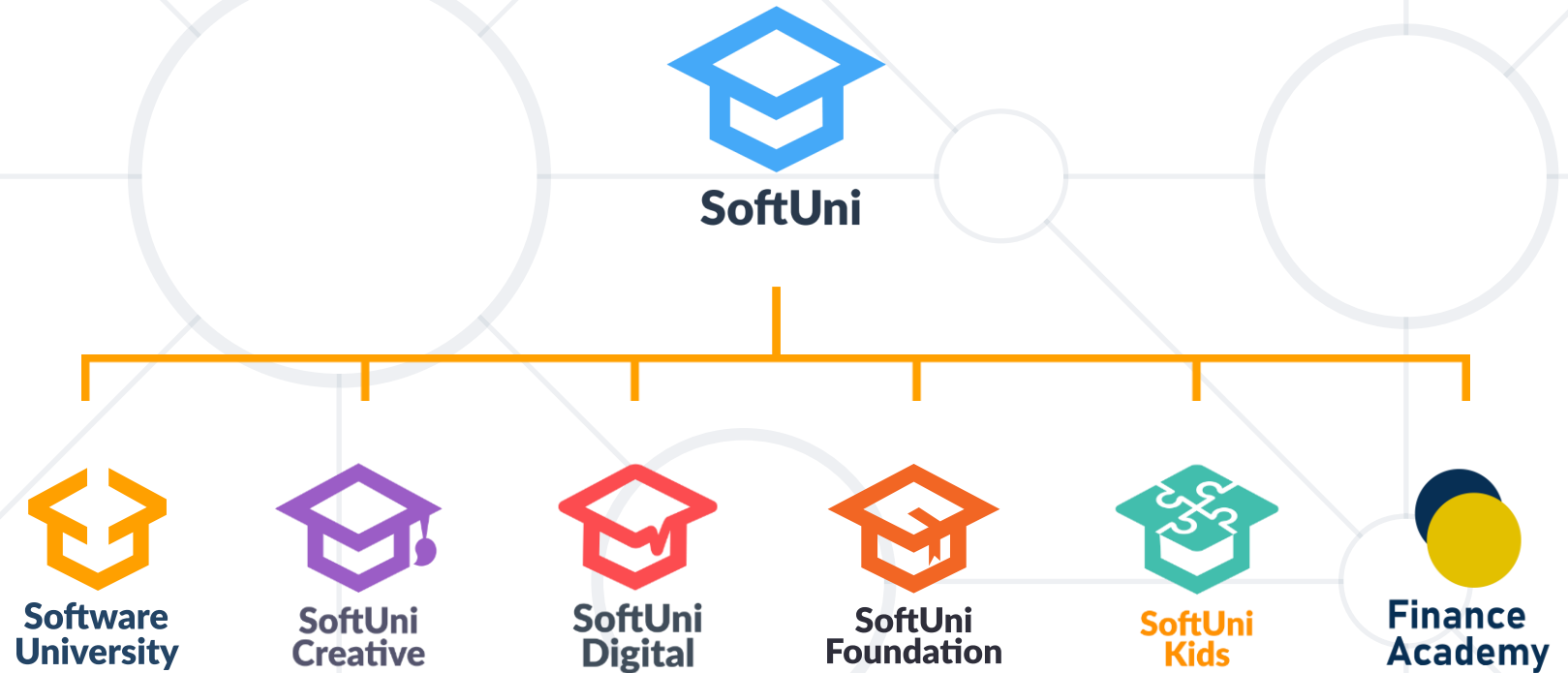
Implementing IComparable<T>

```
public class Point : IComparable<Point> {  
    public int x, y;  
  
    public int CompareTo(Point other) {  
        if (x != other.x) {  
            return this.X.CompareTo(other.x);  
        }  
        else {  
            return this.y.CompareTo(other.y);  
        }  
    }  
}
```


- **Hash-tables** map keys to values
 - Rely on hash-functions to distribute the keys in the table
 - Collisions needs resolution algorithm (e.g., chaining)
 - Very fast add / find / delete – **$O(1)$**
- **Sets** hold a group of elements
- **Dictionaries** map key to value



Questions?



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