

# ***Relational Algebra***

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*Content from Database system concepts-Korth and Dr. Bruns*

# Functions on Relations

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In programming we define objects and methods on those objects.

For example, for strings we have operations like

- `substring(string, first, last)`
- `concat(string1, string2)`

In relational databases the objects are relations, and we define functions on relations.

# Select $\sigma$ : filter tuples of a relation

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Patient No.	Last name	First name	Sex
454	Smith	John	M
223	Jones	Peter	M
597	Brown	Brenda	F
234	Jenkins	Alan	M
244	Wells	Christy	F

$\sigma_{\text{Sex} = \text{F}} (\text{PATIENT})$



Patient No.	Last name	First name	Sex
597	Brown	Brenda	F
244	Wells	Christy	F

# Project $\Pi$ : slice attributes of relation

Patient No.	Last name	First name	Ward No.
454	Smith	John	6
223	Jones	Peter	8
597	Brown	Brenda	3
234	Jenkins	Alan	7
244	Wells	Chris	6

The operation takes a table as input and produces a table as output

$\Pi$  "patient no.", "first name" (PATIENT)

Patient No.	First name
454	John
223	Peter
597	Brenda
234	Alan
244	Chris

# Union: add the tuples of two relations

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Class_101	Student_1
Class_205	Student_4
Class_205	Student_5

Class_101	Student_2
Class_205	Student_3

**Relation1 U Relation2**  
**Expression U Expression**

Class_101	Student_1
Class_205	Student_4
Class_205	Student_5
Class_101	Student_2
Class_205	Student_3

1. Duplicate tuples eliminated
2. Schema of relations must match exactly

# Question about project

Patient No.	Last name	First name	Ward No.
454	Smith	John	6
223	Jones	Peter	8
597	Brown	Brenda	3
234	Jenkins	Alan	7
244	Brown	Chris	6

What happens if we perform a project on this relation using attribute 'Last name'?

$\pi_{\text{last name}}(\text{PATIENT})$

Last name
Smith
Jones
Brown
Jenkins
Brown

Problem?

# Removing duplicate rows

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As part of applying an operation of relation algebra, duplicate rows should be removed.

Duplicate rows aren't technically allowed in a valid table.

In practice a DB system may allow duplicate rows.

$\Pi_{\text{lastname}}$  (PATIENT) will remove duplicate last names from the resultant relation.

Relation Algebra: Sets (duplicates eliminated)

SQL: Multisets, Bags (duplicates allowed)

# Answering questions with Project/Select

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What are the names of students with more than 100 credits?

To get the answer:

1. select the rows of the 'student' table with `tot_cred > 100`
2. get the 'name' attribute

## STUDENT

ID	name	dept_name	tot_cred
128	Zhang	Comp. Sci.	102
12345	Shankar	Comp. Sci.	32
19991	Brandt	History	80
23121	Chavez	Finance	110
44553	Peltier	Physics	56
45678	Levy	Physics	46

$\Pi_{\text{Name}} (\sigma_{\text{tot\_cred} > 100} (\text{STUDENT}))$



# Another example

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What courses are offered in Spring '09, and what buildings are they in?

## SECTION

course_id	sec_id	semester	year	building	room_number	time_slot_id
BIO-101	1	Summer	2009	Painter	514	B
BIO-301	1	Summer	2010	Painter	514	A
CS-101	1	Fall	2009	Packard	101	H
CS-101	1	Spring	2010	Packard	101	F
CS-190	1	Spring	2009	Taylor	3128	E
CS-190	2	Spring	2009	Taylor	3128	A
CS-315	1	Spring	2010	Watson	120	D
CS-319	1	Spring	2010	Watson	100	B

$R = \sigma_{\text{semester}=\text{Spring} \wedge \text{year}=2009}(\text{SECTION})$

$\Pi_{\text{course\_id}, \text{building}}(R)$

# Exercise

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Write the operations to get the IDs of students who took course CS-101 in Fall 2009.

Remember the structure of the 'takes' table:

## TAKES

ID	course_id	sec_id	semester	year	grade
128	CS-101	1	Fall	2009	A
128	CS-347	1	Fall	2009	A-
12345	CS-101	1	Fall	2009	C
12345	CS-190	2	Spring	2009	A
12345	CS-315	1	Spring	2010	A
12345	CS-347	1	Fall	2009	A

$\Pi_{ID} (\sigma_{\text{course\_id}=\text{"CS-101"} \wedge \text{semester}=\text{"Fall"} \wedge \text{year}=2009} (\text{TAKES}))$

# A 'union' example

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Get the ID of courses  
taught in Fall 2015 or  
Fall 2016

## SECTION

course_id	sec_id	semester	year	building	room_number	time_slot_id
BIO-101	1	Summer	2009	Painter	514	B
BIO-301	1	Summer	2010	Painter	514	A
CS-101	1	Fall	2009	Packard	101	H
CS-101	1	Spring	2010	Packard	101	F
CS-190	1	Spring	2009	Taylor	3128	E

$R1 = \sigma_{\text{semester}=\text{"Fall"} \wedge \text{year}=2015}(\text{SECTION})$

$R2 = \sigma_{\text{semester}=\text{"Fall"} \wedge \text{year}=2016}(\text{SECTION})$

$\Pi \text{"course\_ID"} (R1 \cup R2)$

Last line could alternatively be written:

$\Pi \text{"course\_ID"} (R1) \cup \Pi \text{"course\_ID"} (R2)$

# Cartesian Product – combine 2 relations

R1	id	name
	1	Jack
	2	Sally
	3	Trudy

R2	course	id	time
	DB	1	10
	OS	2	4

Schema of resultant relation is  
 $\text{schema(R1)} \cup \text{schema(R2)}$

R1.id	name	course	R2.id	time
1	Jack	DB	1	10
1	Jack	OS	2	4
2	Sally	DB	1	10
2	Sally	OS	2	4
3	Trudy	DB	1	10
3	Trudy	OS	2	4

Question: if you take the product of a table with 5 rows and a table with 8 rows, how many rows do you get?

# Using cross product

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What are the names of CS instructors and what classes do they teach?

**INSTRUCTOR**

ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000

**TEACHES**

ID	course_id	sec_id	semester	year
10101	CS-101	1	Fall	2009
10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-101	1	Spring	2010
15151	MUS-199	1	Spring	2010
22222	PHY-101	1	Fall	2009

$R1 = \text{INSTRUCTOR} \times \text{TEACHES}$

$R2 = \sigma_{\text{instructor.ID} = \text{teaches.ID} \wedge \text{dept\_name} = \text{"Comp. Sci."}} (R1)$

$\Pi_{\text{"name", "course\_id"}} (R2)$

# Natural Join

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- ❑ Enforce equality on all attributes with same name
  - Performs a cross-product
  - Enforces equality on all attributes with the same name
- ❑ Eliminates one copy of duplicate attributes
  - Don't need to keep two copies of the duplicate column because the values are always going to be equal

# Using Natural Join

What are the names of CS instructors and what classes do they teach?

**INSTRUCTOR**

ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000

**TEACHES**

ID	course_id	sec_id	semester	year
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10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-101	1	Spring	2010
15151	MUS-199	1	Spring	2010
22222	PHY-101	1	Fall	2009

$\Pi_{\text{"name", "course\_id"}} (\sigma_{\text{dept\_name}=\text{"Comp. Sci."}} (\text{INSTRUCTOR} \bowtie \text{TEACHES}))$

# Theta Join $\bowtie_{\theta}$

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- $\text{Exp1} \bowtie_{\theta} \text{Exp2} = \sigma_{\theta} (\text{Exp1} \times \text{Exp2})$ 
  - Performs a cross-product (combine all tuples)
  - Keep tuples that pass  $\theta$  the condition



# Lab – Use Case 1

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What are the names of students and instructors?

## INSTRUCTOR

ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000

## STUDENT

ID	name	dept_name	tot_cred
128	Zhang	Comp. Sci.	102
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23121	Chavez	Finance	110
44553	Peltier	Physics	56
45678	Levy	Physics	46

$\Pi \text{ "name"} (\text{INSTRUCTOR}) \cup \Pi \text{ "name"} (\text{STUDENT})$

# Lab – Use Case 2

What are the names of students and instructors?

## INSTRUCTOR

ID	iName	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000

## STUDENT

ID	sName	dept_name	tot_cred
128	Zhang	Comp. Sci.	102
12345	Shankar	Comp. Sci.	32
19991	Brandt	History	80
23121	Chavez	Finance	110
44553	Peltier	Physics	56
45678	Levy	Physics	46

$\Pi$  “iName” (INSTRUCTOR)  $\cup$   $\Pi$  “sName” (STUDENT) ?

# Rename $\rho$

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- Lets say relation  $R1(C1, \dots, Cn)$
- $\rho_{R2(C1, \dots, Cn)}(R1)$
- $\rho_{R2(A1, \dots, An)}(R1)$
- $\rho_{R3(A1, \dots, An)}(Exp)$ 
  - Exp is relation algebra expression
  - Remember: result of expression is relation
- Rename operator reassigns the schema in the result of Exp

# Lab – Use Case 3

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Employees (Employee\_ID, Name, Salary, Manager\_ID)

Find name of each employee's manager?

$$\Pi \text{ "ename", "mname" } ( \sigma_{eMID = mID} ( \rho_{emp(eID, eName, eMID)}(Employees) \times \rho_{mgr(mID, mName, mMID)}(Employees) ) )$$

# Expression Tree

SQL is compiled into an expression tree

