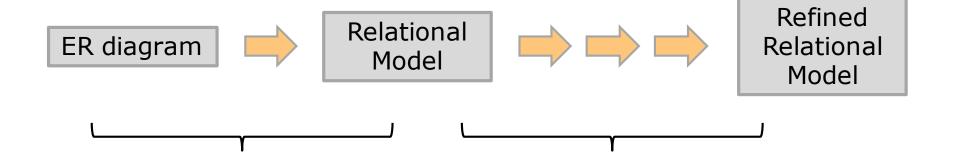
# Normalization 1: Functional dependencies

## Database design process



we discussed how to create an ER design, and how to map it to a relational design

now we discuss how to refine the initial relational design to a final one

# Is this a good schema?

## 

ID name		salary	dept_name	building	budget	
22222	Einstein	95000	Physics	Watson	70000	
12121	Wu	90000	Finance	Painter	120000	
32343	El Said	60000	History	Painter	50000	
45565	45565 Katz		Comp. Sci. Taylor		100000	
98345	Kim	80000	Elec. Eng.	Taylor	85000	
76766	Crick	72000	Biology	Watson	90000	
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000	
58583	Califieri	62000	History	Painter	50000	
83821	Brandt	92000	Comp. Sci.	Taylor	100000	
15151	Mozart	40000	Music	Packard	80000	
33456	Gold	87000	Physics	Watson	70000	
76543	Singh	80000	Finance	Painter	120000	

This schema obtained by joining 'instructor' and 'department'.

## Pros and Cons

ID name		salary	dept_name	building	budget	
22222	Einstein	95000	Physics	Watson	70000	
12121	Wu	90000	Finance	Painter	120000	
32343	El Said	60000	History	Painter	50000	
45565	45565 Katz		Comp. Sci.	Taylor	100000	
98345	Kim	80000	Elec. Eng.	Taylor	85000	
76766	Crick	72000	Biology	Watson	90000	
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000	
58583	Califieri	62000	History	Painter	50000	
83821	Brandt	92000	Comp. Sci.	Taylor	100000	
15151	Mozart	40000	Music	Packard	80000	
33456	Gold	87000	Physics	Watson	70000	
76543	Singh	80000	Finance	Painter	120000	

#### Good news:

Fewer joins needed in queries.

#### Bad news 1:

If Physics budget changes, it must be changed in multiple places.

#### Bad news 2:

How to create a new department if no instructors yet for the department?

#### Bad news 3:

If the last instructor in a department leaves, other dept. info is lost

## Improving a schema

If we were given a schema like this, would we know how to improve it?

ID	name	salary	dept_name	building	budget		
22222	Einstein	95000	Physics	Watson	70000		
12121	Wu	90000	Finance	Painter	120000		
32343	El Said	60000	History	Painter	50000		T
45565	Katz	75000	Comp. Sci.	Taylor	100000	<u></u>	It might b
98345	Kim	80000	Elec. Eng.	Taylor	85000		hard to sp
76766	Crick	72000	Biology	Watson	90000		these kind
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000	<u>,                                      </u>	
58583	Califieri	62000	History	Painter	50000		things
83821	Brandt	92000	Comp. Sci.	Taylor	100000		
15151	Mozart	40000	Music	Packard	80000		
33456	Gold	87000	Physics	Watson	70000		
76543	Singh	80000	Finance	Painter	120000		

## Improving a schema

To improve the schema, it would be helpful if we knew this:

dept\_name → budget

ID	name	salary	dept_name	building	budget		
22222	Einstein	95000	Physics	Watson	70000		
12121	Wu	90000	Finance	Painter	120000		With this
32343	El Said	60000	History	Painter	50000		
45565	Katz	75000	Comp. Sci.	Taylor	100000	_	<u>functional</u>
98345	Kim	80000	Elec. Eng.	Taylor	85000		dependency
76766	Crick	72000	Biology	Watson	90000		above, we can
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000 <		•
58583	Califieri	62000	History	Painter	50000		identify some
83821	Brandt	92000	Comp. Sci.	Taylor	100000		redundancy.
15151	Mozart	40000	Music	Packard	80000		,
33456	Gold	87000	Physics	Watson	70000		
76543	Singh	80000	Finance	Painter	120000		

(table from Database System Concepts, Silbershatz et al)

## Decomposing a schema

#### If these functional dependencies were given:

dept\_name → budget

dept\_name → building

ID name		salary	dept_name	building	budget	
22222	Einstein	95000	Physics	Watson	70000	
12121	Wu	90000	Finance	Painter	120000	
32343	El Said	60000	History	Painter	50000	
45565	Katz	75000	Comp. Sci.	Taylor	100000	
98345	Kim	80000	Elec. Eng.	Taylor	85000	
76766	Crick	72000	Biology	Watson	90000	
10101	Srinivasan	65000	Comp. Sci.	Taylor	100000	
58583	Califieri	62000	History	Painter	50000	
83821	Brandt	92000	Comp. Sci.	Taylor	100000	
15151	Mozart	40000	Music	Packard	80000	
33456	Gold	87000	Physics	Watson	70000	
76543	Singh	80000	Finance	Painter	120000	

we could see the problem, and could fix it by splitting the schema in two

How to do this systematically?

(table from Database System Concepts, Silbershatz et al)

## Summarizing

- Redundancy in relational schemas is bad
  - storage waste
  - leads to inconsistencies
- It's hard to spot redundancy by looking at a table, especially if it's big
- ☐ If data designer supplies functional dependencies, we can spot redundancies
- It can be fixed by splitting schemas but how exactly?

## Functional dependencies

- a kind of integrity constraint
- examples:
  - students are uniquely identified by their ID
  - each instructor has only one name
  - each instructor is associated with just one department
  - each department has only one budget value
- □ dept\_name → budget
- □ a table satisfies dept\_name → budget if:
  - if two rows have the same dept\_name, they have the same budget

# Functional dependencies: definition

- suppose r(R) is a relational schema (R is the set of attributes in schema r)
- a functional dependency has the form

$$X \rightarrow Y$$

where X, Y are subsets of R

- $\square$  an instance of r(R) satisfies  $X \to Y$  if:
  - □ all pairs of rows with the same values for the *X* attributes have the same values for the *Y* attributes

## Examples

- ☐ students are uniquely identified by their ID

  ID → name
- □ each instructor has only one nameID → name
- each instructor is associated with just one department

?

each department has only one budget value

## Functional dependencies and keys

Question: Are keys just a kind of functional dependency?

### Examples:

- department(<u>dept name</u>, building, budget)
- section(<u>course id</u>, <u>sec id</u>, <u>semester</u>, <u>year</u>, building, room\_number, time\_slot\_id)

#### Remember:

 you can't find functional dependencies by looking at a table

## All the FDs of a schema

When we talk about the functional dependencies of a schema, we mean:

- the FDs we may have explicitly identified
- +
- $\square$  "trivial" FDs, like  $X \to Y$  and  $X, Y \to Y$

+

FDs derivable from other FDs of the schema

"derivable" means by one of these rules:

- 1) if  $X \to Y$  and  $Y \to Z$  then  $X \to Z$
- 2) if  $X \rightarrow Y$  then  $X \cup Z \rightarrow Y \cup Z$

So if ID → name, and name → dept\_name, then ID → dept\_name

## Summary

- We want relational schemas that:
  - don't have redundancy
  - can be queried efficiently
- Redundancy is addressed by splitting a relation into smaller relations
- To do this in a systematic way, we use functional dependencies
- Functional dependencies are integrity constraints
- Keys are a special kind of functional dependency