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Nondeterministic Finite Automata

(Part 1)

Lecture 09
Day 09/31

CS 154
Formal Languages and Computability
Spring 2019

Agenda of Day 09

- Summary of Lecture 08
- Quiz 3
- A few slides from lecture 08
- JFLAP Demo
- Lecture 09: Teaching ...
 - Nondeterministic Finite Automata (Part 1)

Summary of Lecture 08: We learned ...

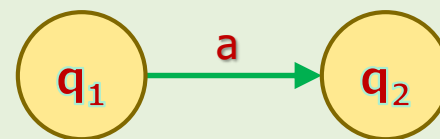
DFAs

- A DFA M is defined by a quintuple:

$$M = (Q, \Sigma, \delta, q_0, F)$$

- Q is ...
 - ... a finite and nonempty set of states.
- Σ is ...
 - ... a finite and nonempty set of symbols called input alphabet.
- δ is ...
 - ... called transition function and is defined as:
 $\delta: Q \times \Sigma \rightarrow Q$
 δ is total function.

- $q_0 \in Q$ is ...
 - ... the initial state.
- $F \subseteq Q$ is ...
 - ... the set of accepting states.
- Every sub-rule like $\delta(q_1, a) = q_2$ represents a transition in transition graph.



Any question?

Summary of Lecture 08: We learned ...

DFAs

- Why total function?
 - ... because if not, in some situations, the DFA does not know where to go!
- DFAs constraint: ...
 - ... DFAs transition function **must** be total function.
- The **consequence** of this constraint is ...
 - ... every state must have an **outgoing transition** for **every symbol** of alphabet.

Any question?

NAME	Alan M. Turing		
SUBJECT	CS 154	TEST NO.	3
DATE	02/21/2019	PERIOD	1 / 2 / 3

TEST RECORD	
PART 1	123
PART 2	
TOTAL	



Quiz 3

Use Scantron

A Few Slides from Lecture 08



Prepare Your **Development Environment**

JFLAP (Java Formal Language and Automata Package)

- We'll use **JFLAP** tool in this course to **develop** and **test** our **automata**.
- **Download it from Canvas:** Files/Misc/JFLAP7.1.jar
 - For uniformity, please use the copy that I provided!
- Tutorial: <http://www.jflap.org/tutorial/>

JFLAP Demo



- **Official website:** <http://www.jflap.org/>
- Official download site: <http://www.jflap.org/jflaptmp/> (Don't !)
- The **stable version 7.1** (Jul 27, 2018)

JFLAP Basic Features



Basic Features

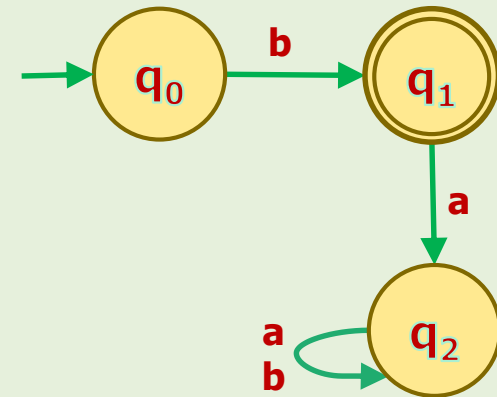
- Creating **states**
- Creating **transitions**
- Defining a state as **initial state** or **final state**
- Deleting
- Shift-Enter to create **multiple transitions**
- **Multiple running** (testing your design)
- **Debugging**: step-by-state
- **Saving** machines (xml file)
- **Selecting** multiple objects and moving
- Changing **state's name** or **label**
- Adding **comment**
- Changing **edges shape**
- **Zoom-in** and **zoom-out**

Nondeterministic Finite Automata (NFA)

DFAs Constraint Violations

DFAs Constraint Violation #1

- What is the problem of the following DFA over $\Sigma = \{a, b\}$?

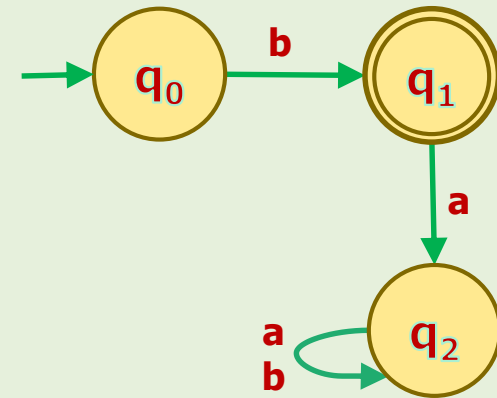


Violation

- The machine has no (zero) transition when it is in state q_0 and the input is a !
 - There is more like this in this graph, what are they?
- In other words, there are some timeframes that the machine does NOT KNOW WHERE TO GO?

DFAs Constraint Violation #1

- What is the value of $\delta(q_0, a)$?
 - $\delta(q_0, a) = \text{"Undefined"}$



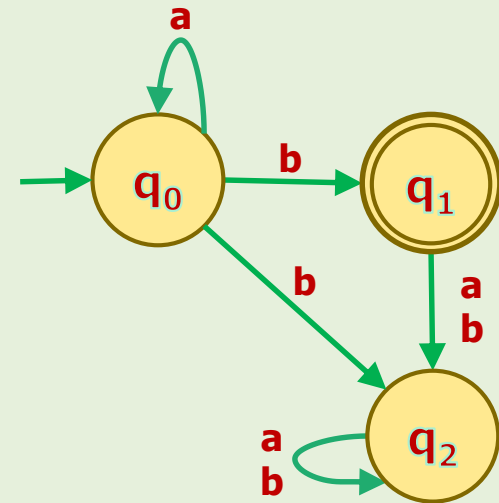
- What type of function is the transition function of this machine?
- Partial function
- ❗ So, the machine is NOT a DFA because it violates DFAs constraint!

DFAs Constraint Violation #2

- What is the **problem** of the following DFA over $\Sigma = \{a, b\}$?

Violation

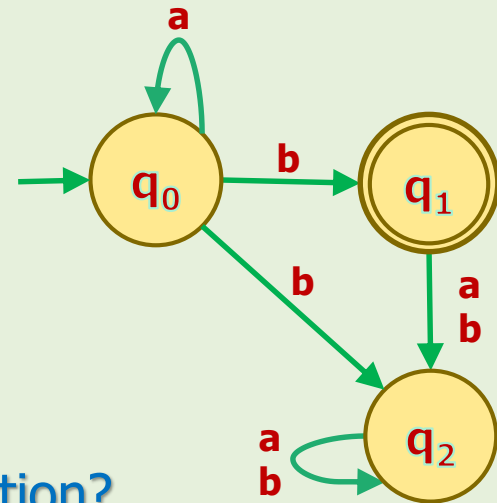
- The machine has **more than one transition** when it is in **state q_0** and the input is **b** !



- In other words, there are **some timeframes** that the machine does **NOT KNOW WHERE TO GO?**
 - Because there are **more than one choice!**

DFAs Constraint Violation #2

- What is the value of $\delta(q_0, b)$?
 - $\delta(q_0, b) = \{q_1, q_2\}$
 - The range has more than one value!
 - So, we have to put them in a set of Qs .



- What type of function is the transition function?
 - It is NOT a regular function because it violates the definition of function.
 - It is called "multifunction" (aka multivalued function) in math.



- So, the machine is not a DFA because it violates the DFAs constraint!

DFAs Constraint Violations Summary

- **Violation #1:** There are some timeframes that the machine has no (zero) transition.
 - The transition function is NOT total function.
- **Violation #2:** There are some timeframes that the machine has more than one transition.
 - The transition function is NOT a regular function.
 - It is multifunction (aka multivalued function).
- Let's relax the DFAs constraint and define a new class of machines!
- We are going to have a new class of machines that these violations are legal.

Let's Relax the DFAs Constraint!

- Recall that DFAs' transition function is defined as:

$$\delta: Q \times \Sigma \rightarrow Q$$

δ is total function.

- To accommodate those two violations, we change the RANGE of the function to a set.
- In this way, the range can have zero, one, or more states.
- In other words, the range of this function is a set of Qs.
- We already know that 2^Q is the power set of Q and it contains all subsets of Q.
- Therefore, we change the range from Q to 2^Q .

$$\delta: Q \times \Sigma \rightarrow 2^Q$$

- Let's take some examples.

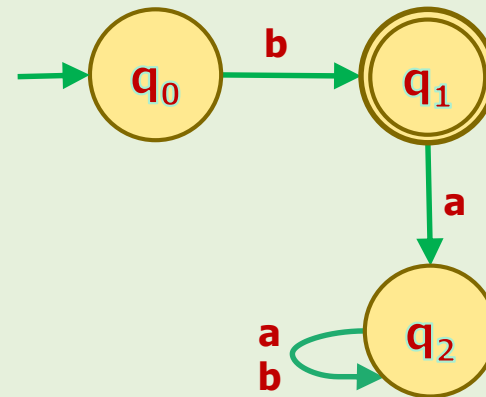
Relaxed Transition Function Examples

Example 1

- Write the **rule** of the following transition graph over $\Sigma = \{a, b\}$.

Solution

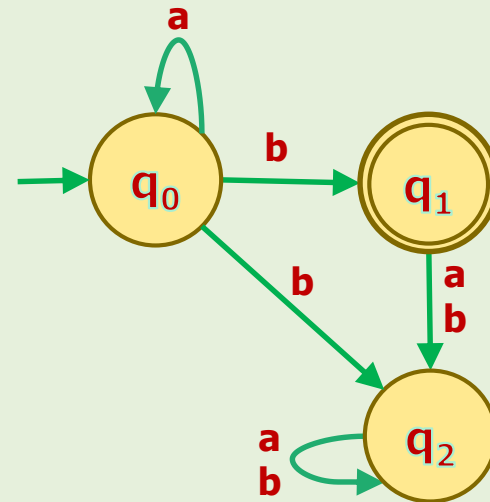
$$\begin{cases} \delta(q_0, a) = \{\} \\ \delta(q_0, b) = \{q_1\} \\ \delta(q_1, a) = \{q_2\} \\ \delta(q_1, b) = \{\} \\ \delta(q_2, a) = \{q_2\} \\ \delta(q_2, b) = \{q_2\} \end{cases}$$



Relaxed Transition Function Examples

Example 2

- Write the **rule** of the following transition graph over $\Sigma = \{a, b\}$.



Solution

$$\begin{cases} \delta(q_0, a) = \{q_0\} \\ \delta(q_0, b) = \{q_1, q_2\} \\ \delta(q_1, a) = \{q_2\} \\ \delta(q_1, b) = \{q_2\} \\ \delta(q_2, a) = \{q_2\} \\ \delta(q_2, b) = \{q_2\} \end{cases}$$

References

1. Linz, Peter, "An Introduction to Formal Languages and Automata, 5th ed.," Jones & Bartlett Learning, LLC, Canada, 2012
2. Michael Sipser, "Introduction to the Theory of Computation, 3rd ed.," CENGAGE Learning, United States, 2013
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