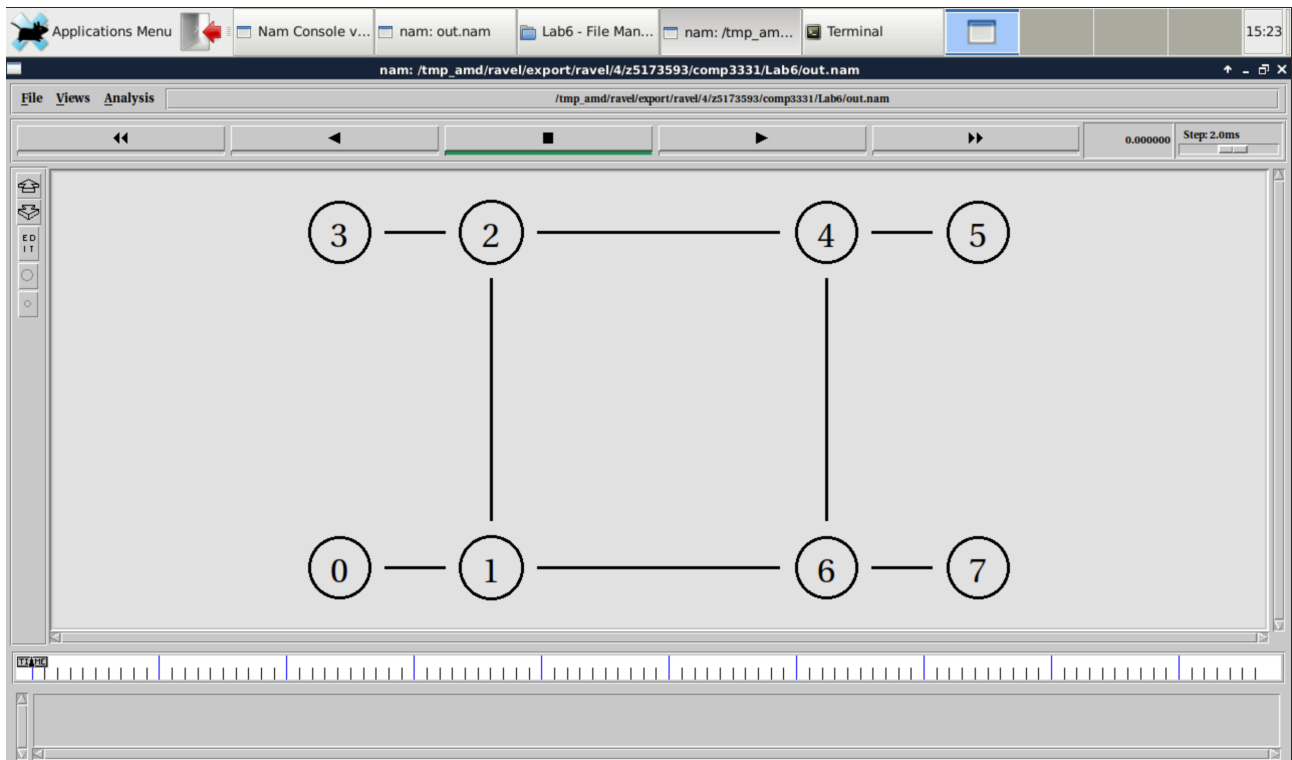
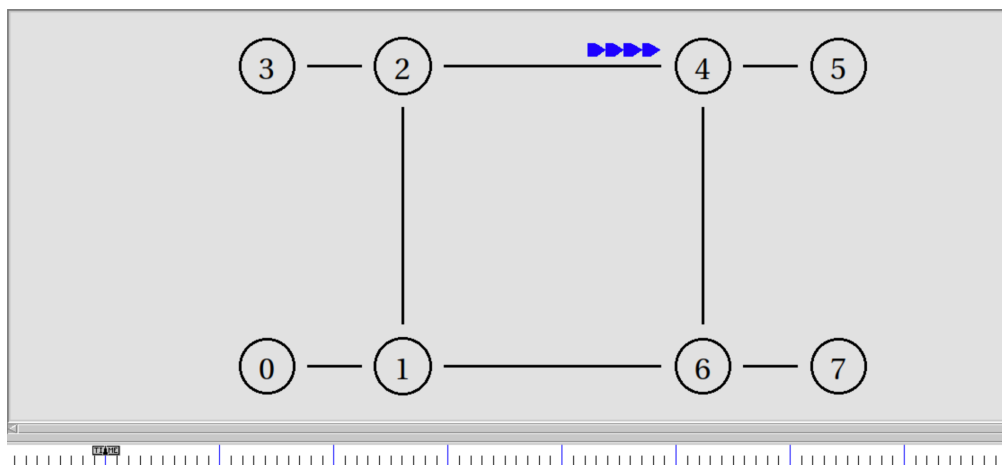


# Lab Exercise 6: Throughput, IP Fragmentation and Routing

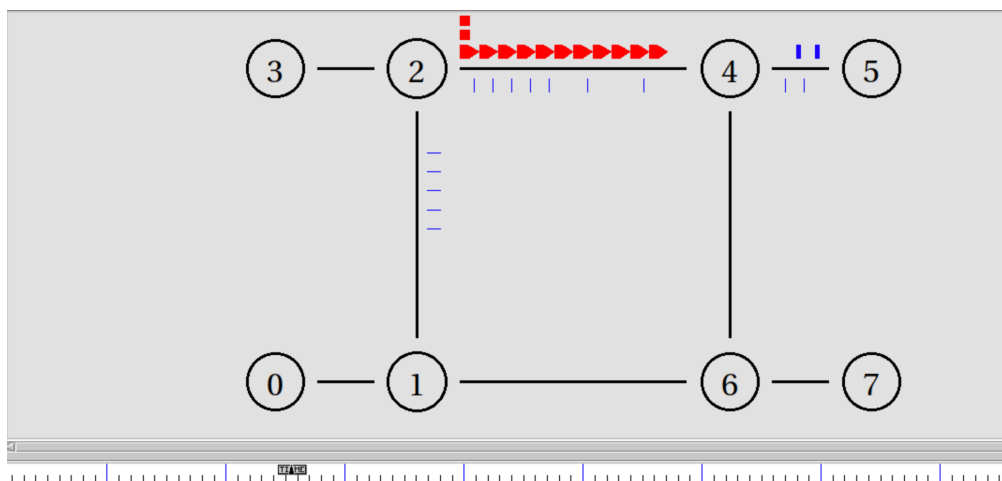
## Exercise 1: Setting up NS2 simulation for measuring TCP throughput



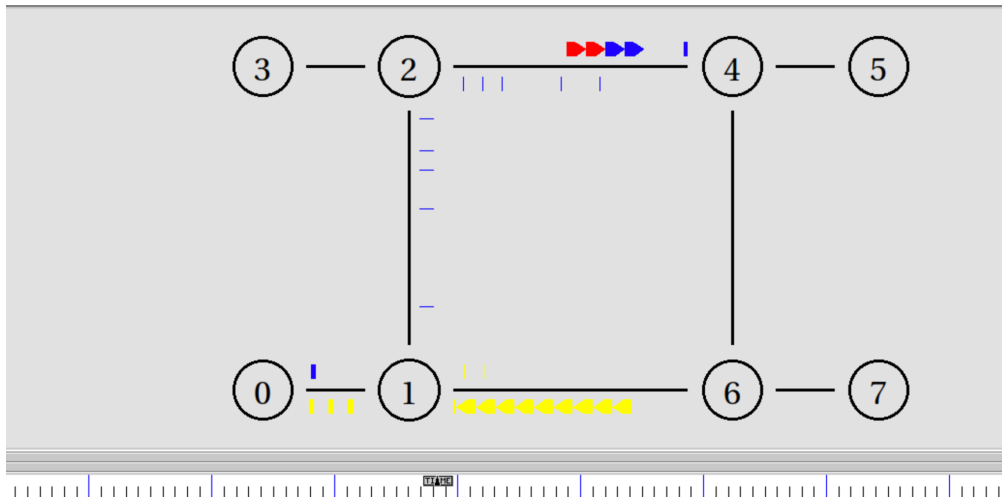
n0 -> n5: 0.5s~8.5s, n0->n1->n2->n4->n5



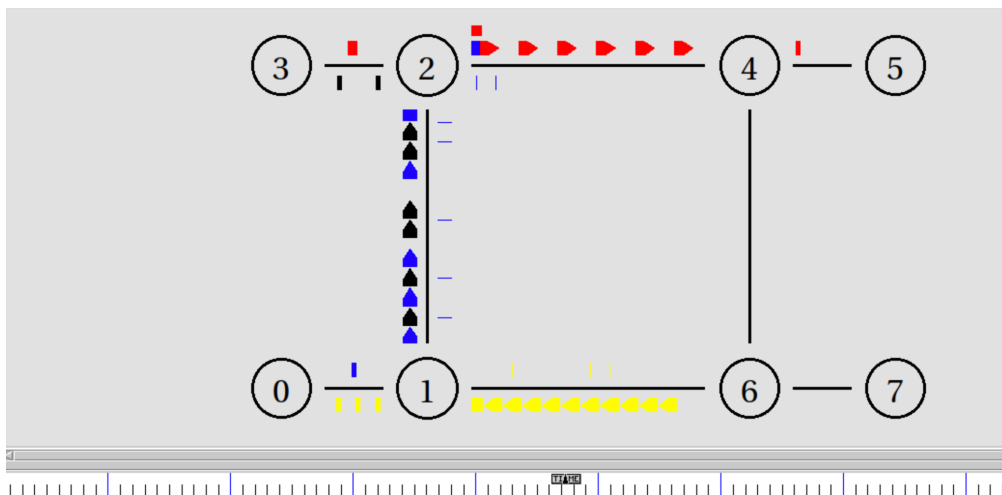
n3 -> n5: 2.0s~9.5s, n3->n2->n4->n5



n7 -> n0: 3s~9.5s, n7->n6->n1->n0



n7 -> n3: 4s~7s, n7->n6->n1->n2->n3



Question 1:

n3->n2 has larger bandwidth than n0->n1->n2.

At 6s, n3->n2 has more packages.

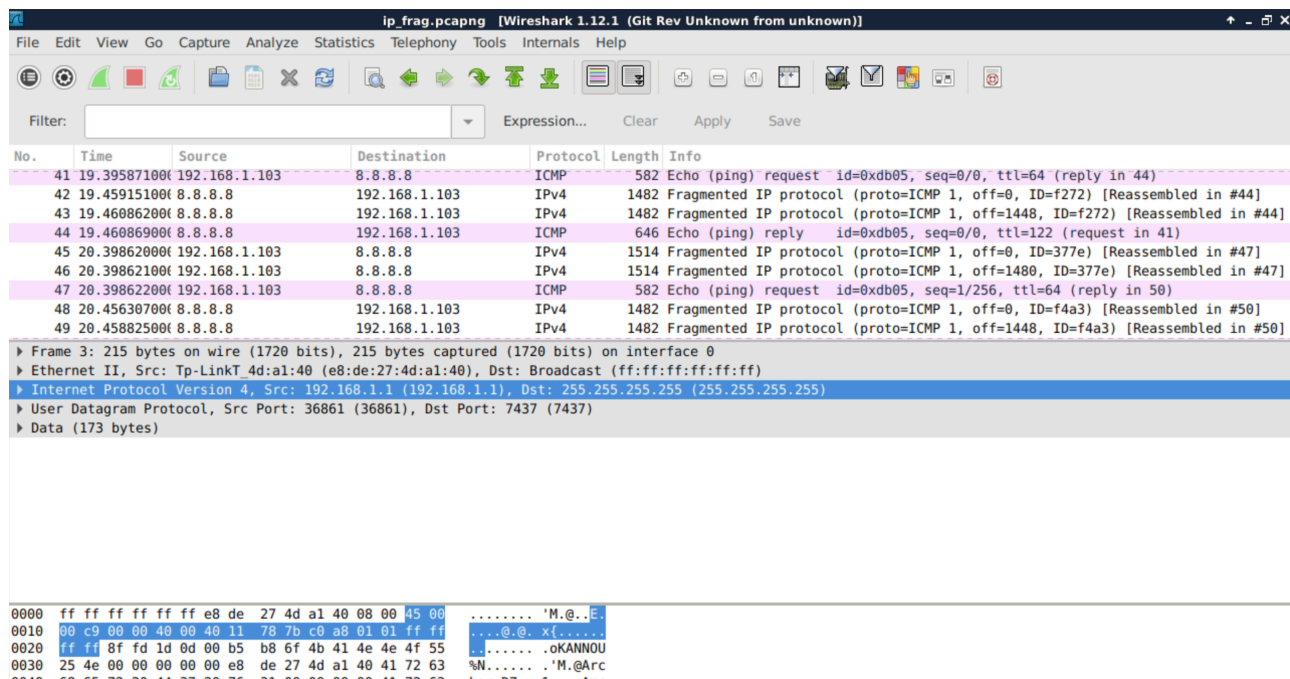
Question 2:

Between 0.5s~2s, tcp1 is on slow start phase, tcp will try to find all the available bandwidth it can use, so the throughput will fluctuating.

Question 3:

All links competing together, they need to share bandwidth, so they must be lower than max bandwidth — 2.5Mbps.

## Exercise 2: Understanding IP Fragmentation



### Question 1:

MTU= 1500 bytes

IP = 2000 + 8 = 2008 bytes > 1500 bytes, caused fragmentation.

host: 192.169.1.103

2008 - 20 = 1988 bytes, p1 = 1480+20=1500 bytes, p2 = 528+20=548 bytes

2 fragments.

### Question 2:

Yes, 3500 bytes > 1500 bytes. So it also get fragmented.

### Question 3:

ID	Length	Flag	Offset
7a7b	1514	0x2000, more fragments	0
7a7b	1514	0x20b9, more fragments	185
7a7b	582	172	370

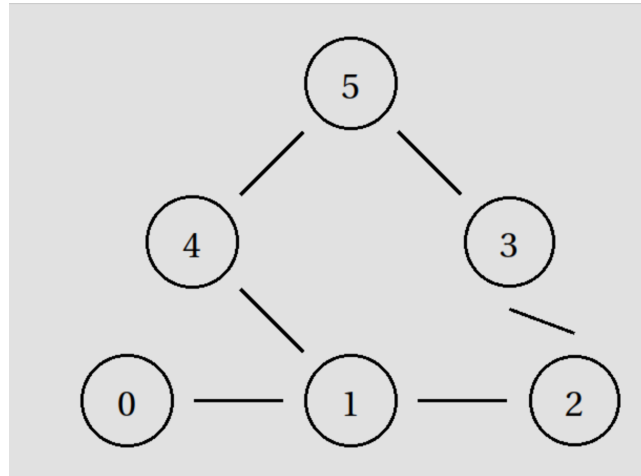
### Question 4:

Yes, 3500 bytes > 1480 bytes.

### Question 5:

When one fragment lost, it will retransmit the whole package, so 3 fragments will be retransmit for 192.168.1.103.

### Exercise 3: Understanding the Impact of Network Dynamics on Routing



Question 1:

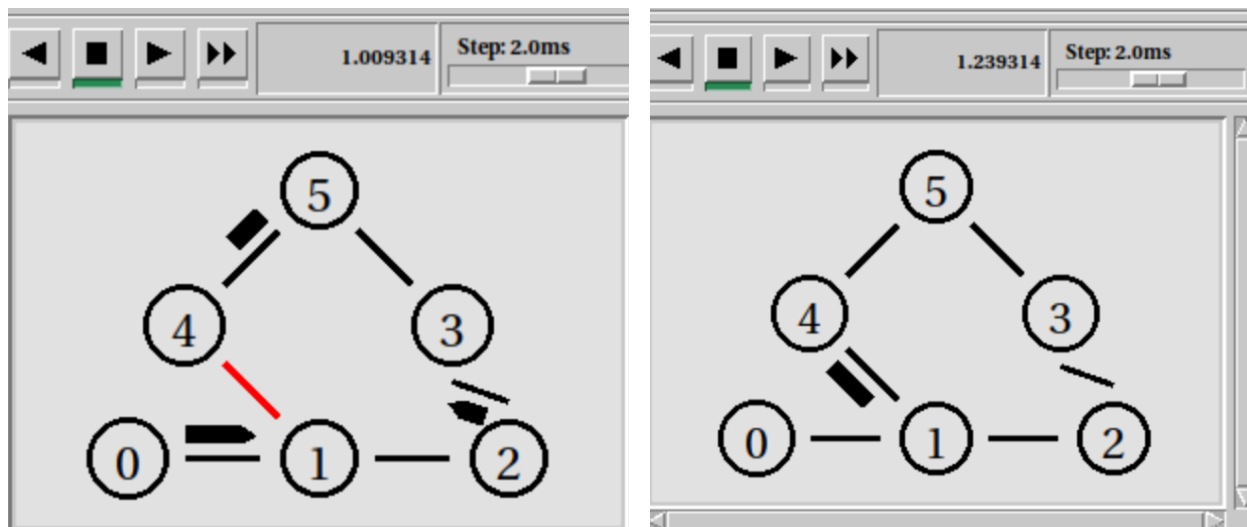
n0 communicate with n5, n2 communicate with n5

route 1: n0->n1->n4->n5

route 2: n2->n3->n5

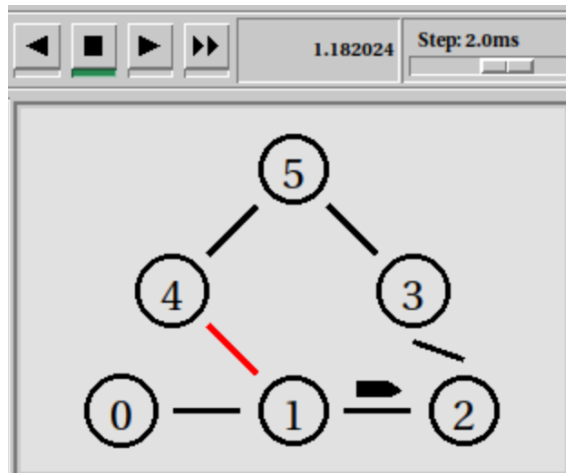
not change.

Question 2:



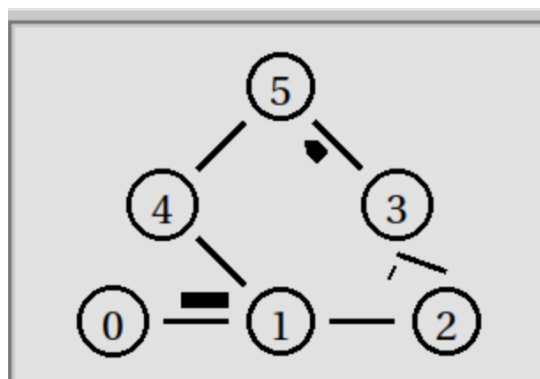
At time 1.0s~1.2s, because node 1 and node 4 down, the link from n1->n4 can not use.  
After 1.2s, node 1 and node 4 up again, the link can be use again.

Question 3:



Between 1.0s~1.2s, when link between  $n1 \rightarrow n4$  down, the package transfer from  $n0 \rightarrow n1 \rightarrow n2 \rightarrow n3 \rightarrow n5$ .

Question 4:



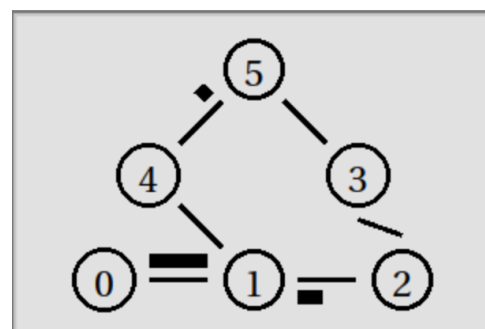
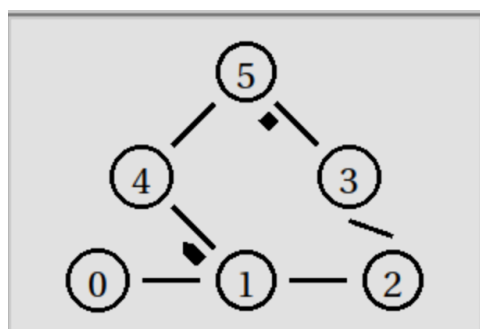
No package transfer using link  $n1 \rightarrow n4 \rightarrow n5$ , they all use  $n2 \rightarrow n3 \rightarrow n5$

route 1:  $n0 \rightarrow n1 \rightarrow n2 \rightarrow n3 \rightarrow n5$

route 2:  $n2 \rightarrow n3 \rightarrow n5$

Because the cost from  $n1 \rightarrow n4$  is 3, which means from  $n0 \rightarrow n1 \rightarrow n2 \rightarrow n3 \rightarrow n5$  will cheaper than  $n0 \rightarrow n1 \rightarrow n4 \rightarrow n5$ .

Question 5:



Cost for  $n0 \rightarrow n1 \rightarrow n4 \rightarrow n5 = 4$ ,  $n0 \rightarrow n1 \rightarrow n2 \rightarrow n3 \rightarrow n5 = 5$ ,  $n2 \rightarrow n3 \rightarrow n5 = 4$ ,  $n2 \rightarrow n1 \rightarrow n4 \rightarrow n5 = 4$ . So cheapest cost was chosen for route 1:  $n0 \rightarrow n1 \rightarrow n4 \rightarrow n5$ . The cost for two ways are same for route 2, and we 'set multiPath', so the path for route 2 will vary.