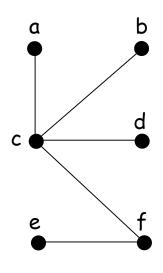
Matemáticas Discretas II

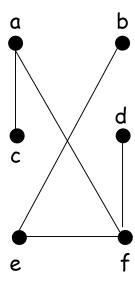
Oscar Bedoya

oscar.bedoya@correounivalle.edu.co

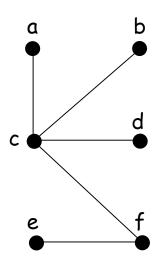
- Árboles
- Aplicaciones de los árboles
- · Recorridos de los árboles
- Algoritmo de Prim
- Algoritmo de Kruskal

Árboles

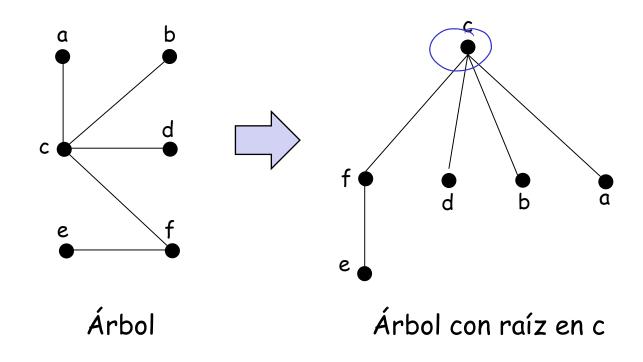


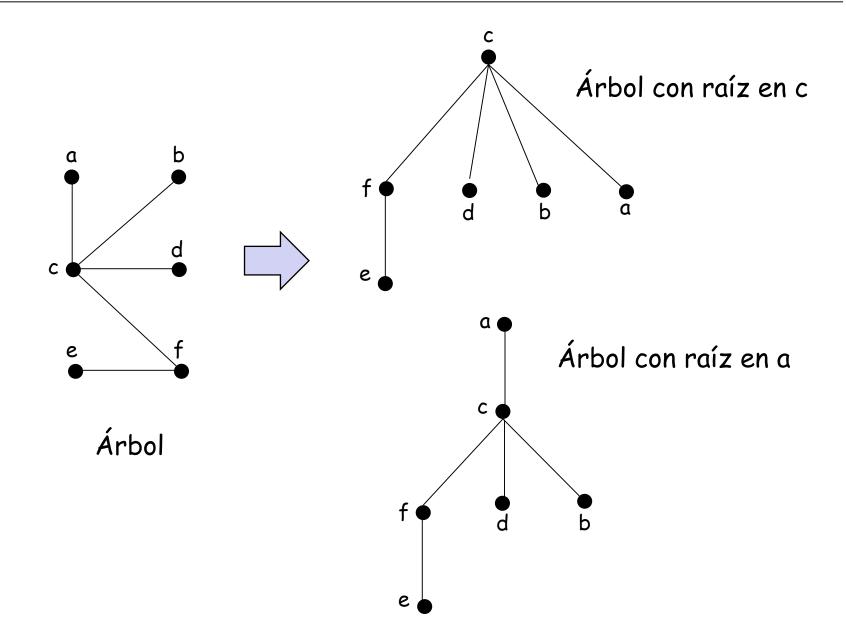


Árboles

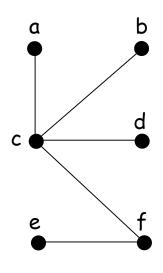


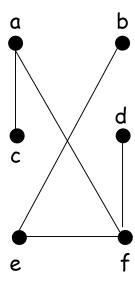
Árboles



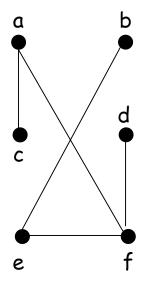


Árboles



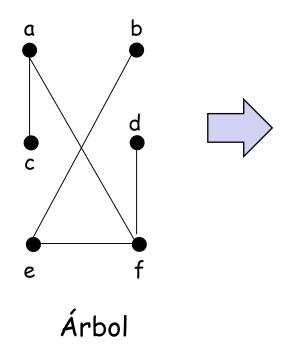


Árboles



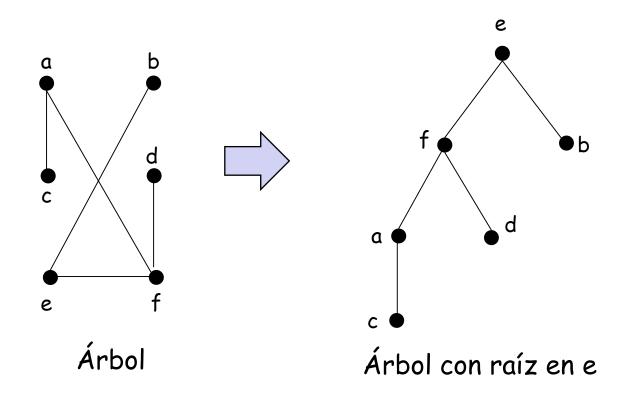
Árboles

Un árbol es un grafo conexo que no tiene circuitos



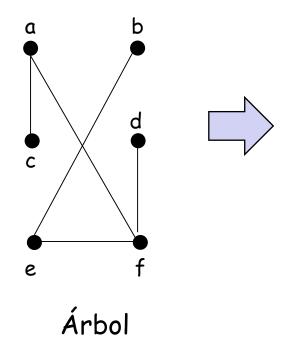
Árbol con raíz en e

Árboles



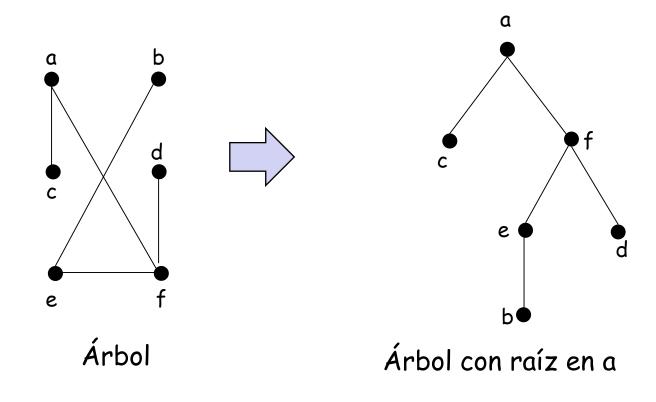
Árboles

Un árbol es un grafo conexo que no tiene circuitos

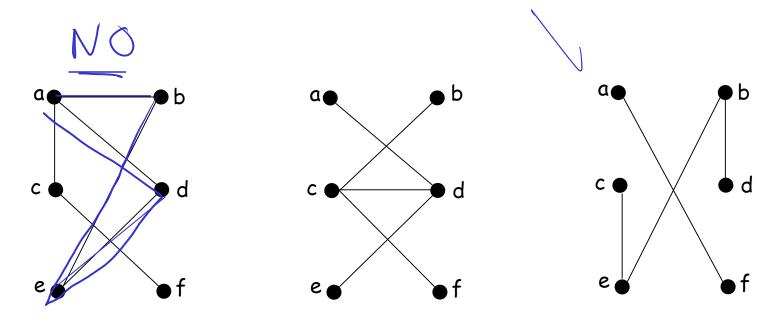


Árbol con raíz en a

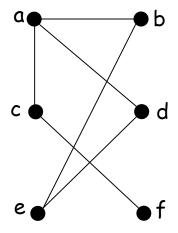
Árboles



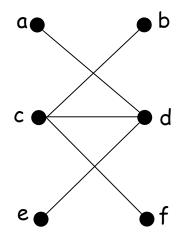
Indicar cuáles de los siguientes grafos son árboles



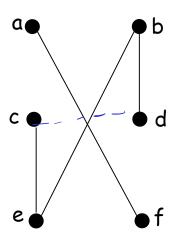
Indicar cuáles de los siguientes grafos son árboles



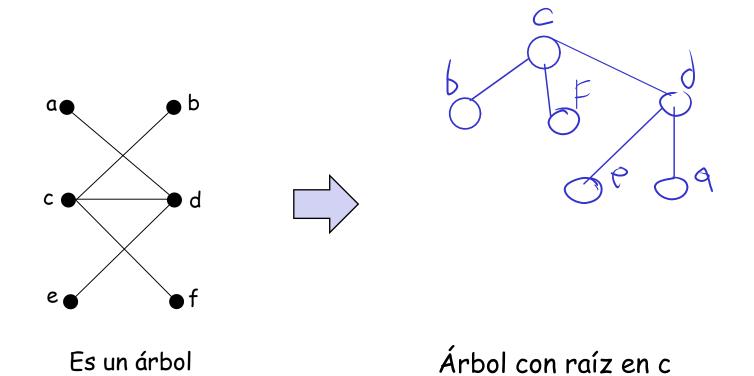
No es árbol porque hay un circuito (a→d→e→b→a)

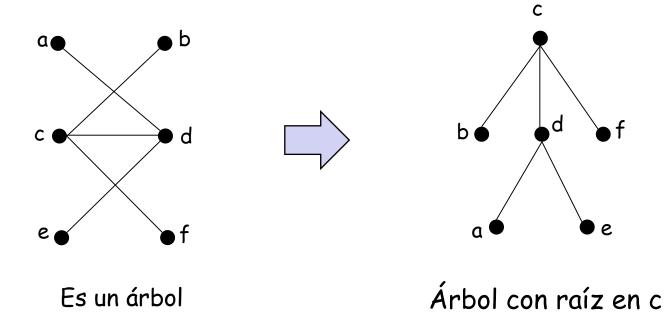


Es un árbol

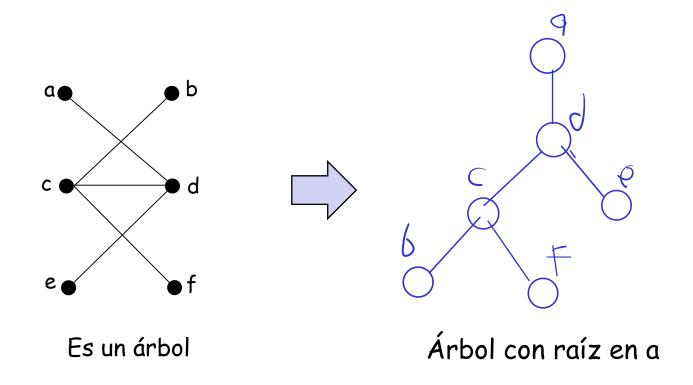


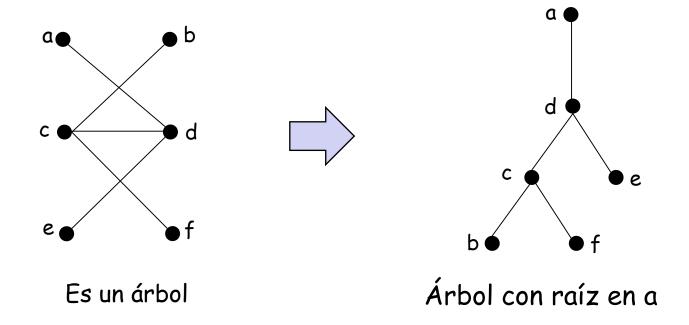
No es árbol porque no es conexo

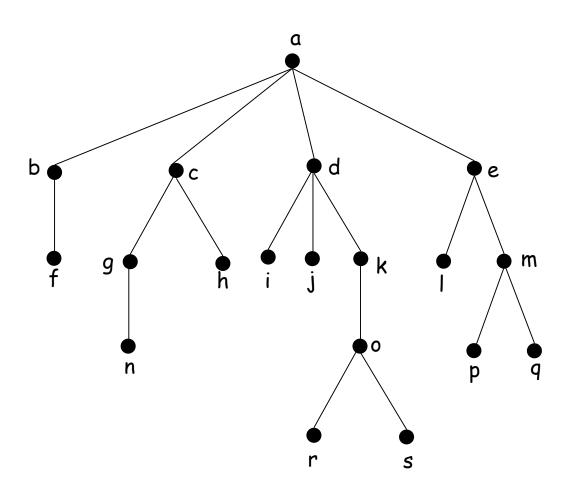


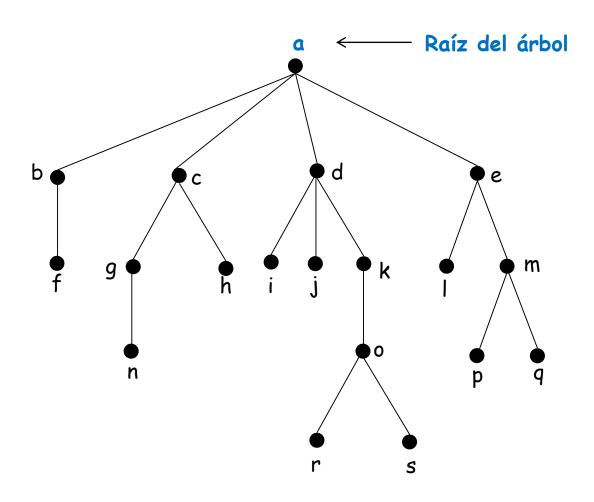


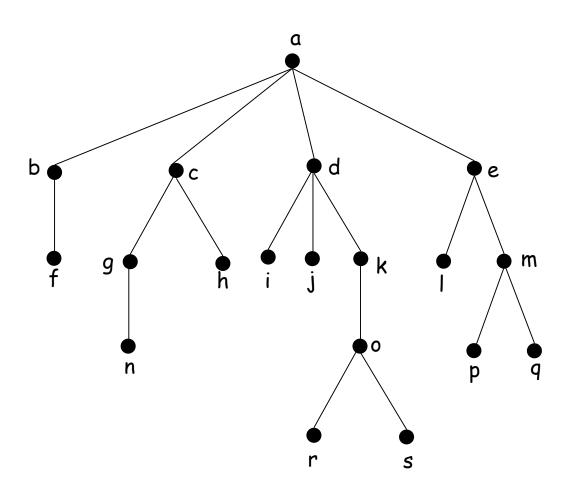
Es un árbol

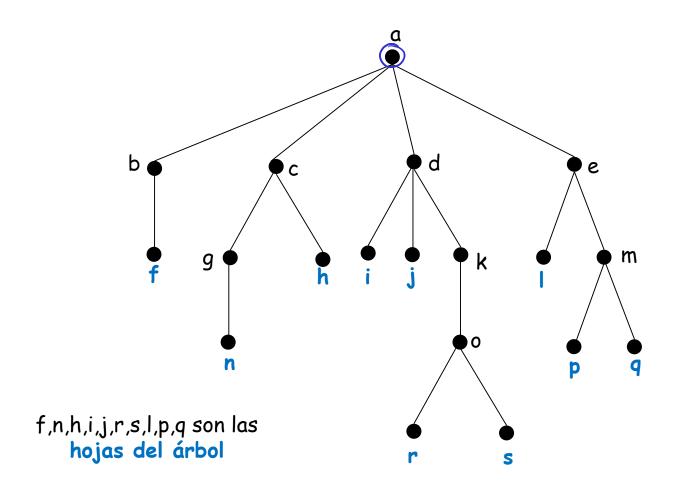


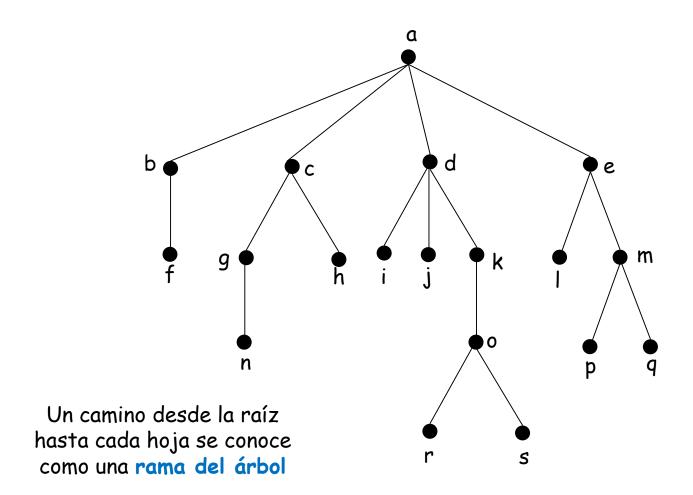


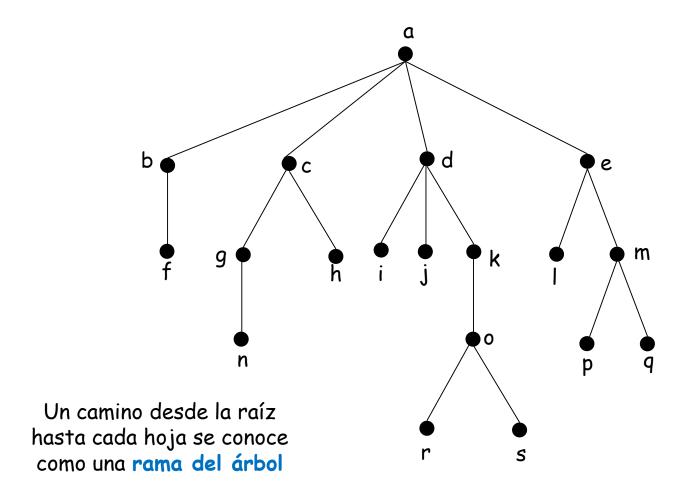




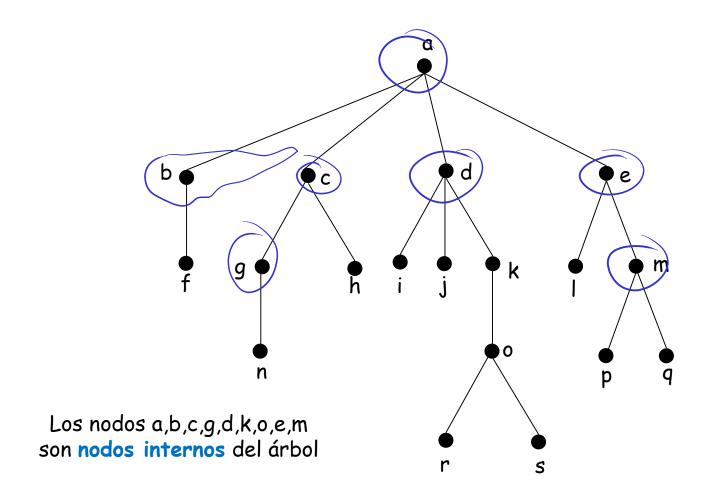


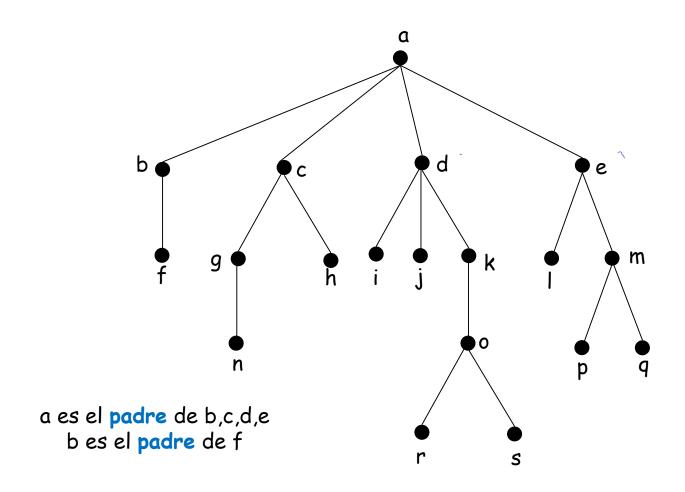


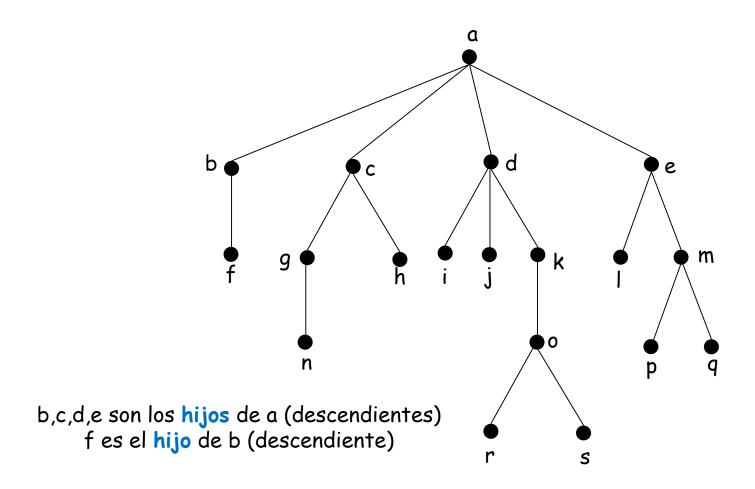


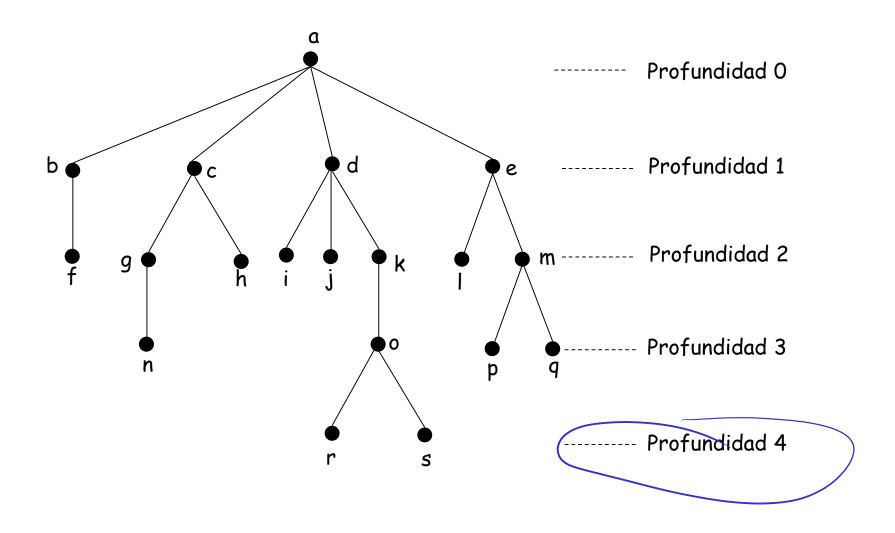


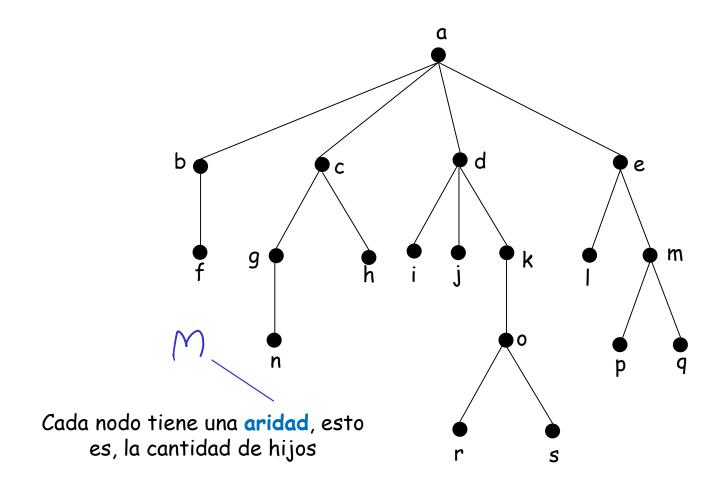
 $a \rightarrow b \rightarrow f$ $a \rightarrow d \rightarrow k \rightarrow o \rightarrow s$

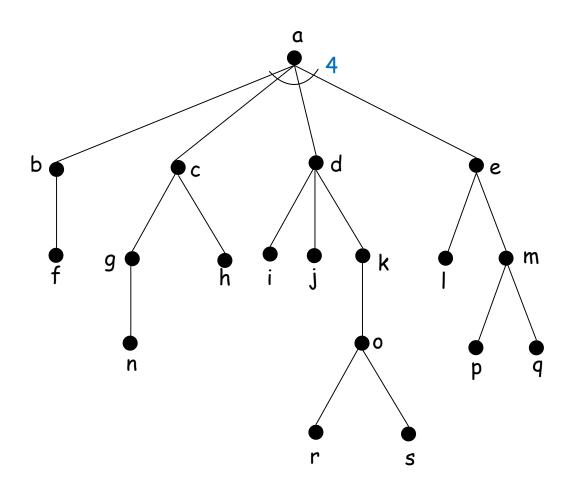


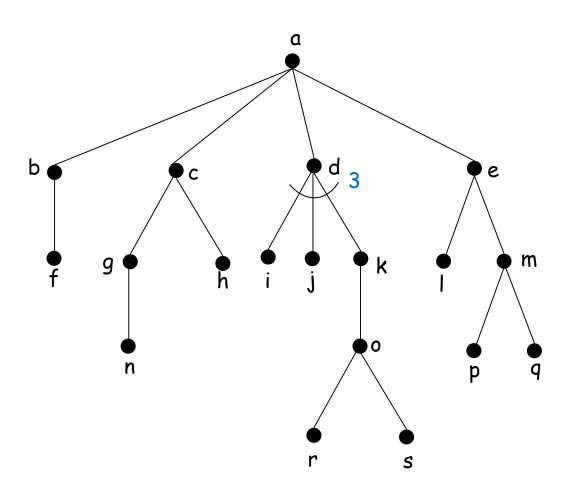


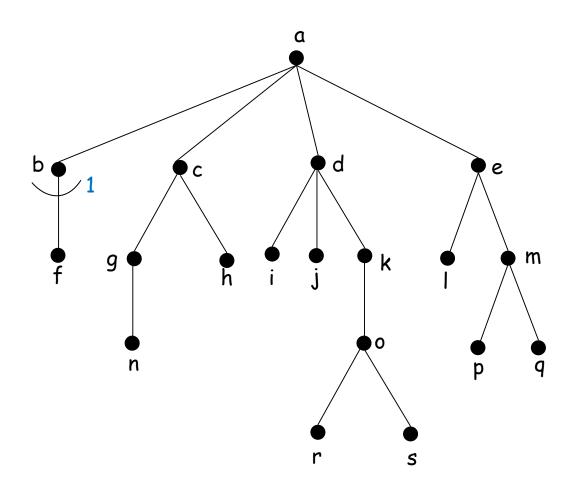


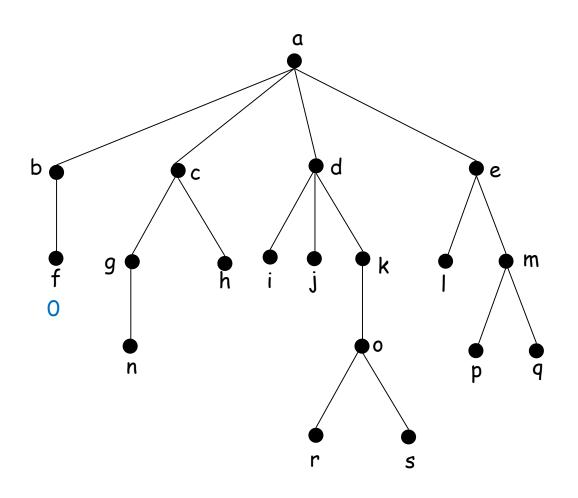






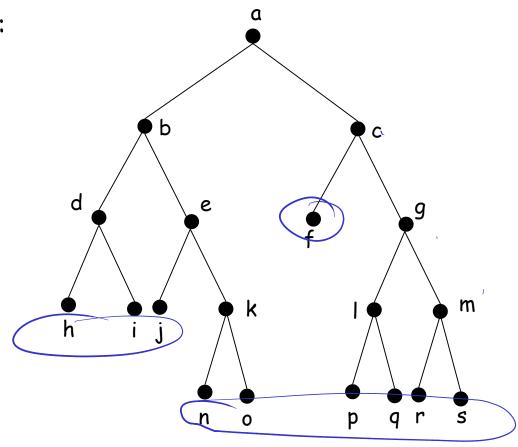






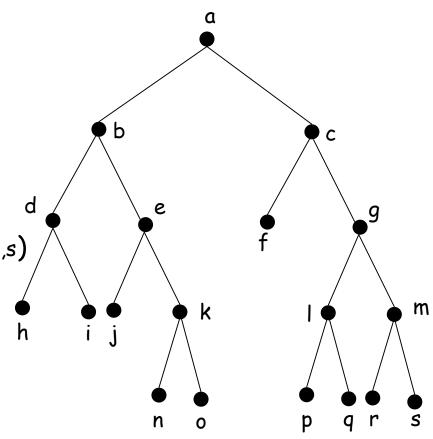
Dado el siguiente árbol indique:

- La raíz del árbol
- El padre de e
- Los hijos de g
- · Las hojas del árbol
- La profundidad de m
- La profundidad de b
- La aridad de c
- La aridad de j



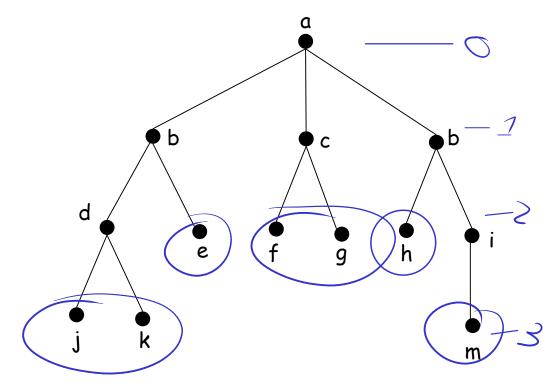
Dado el siguiente árbol indique:

- La raíz del árbol (a)
- El padre de e (b)
- Los hijos de g (l,m)
- Las hojas del árbol (h,i,j,n,o,f,p,q,r,s)
- La profundidad de m (3)
- La profundidad de b (1)
- La aridad de c (2)
- La aridad de j (0)



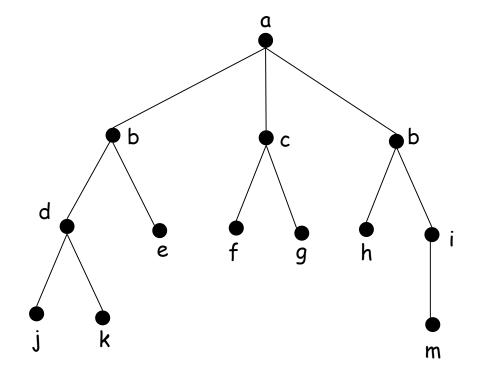
Dado el siguiente árbol indique:

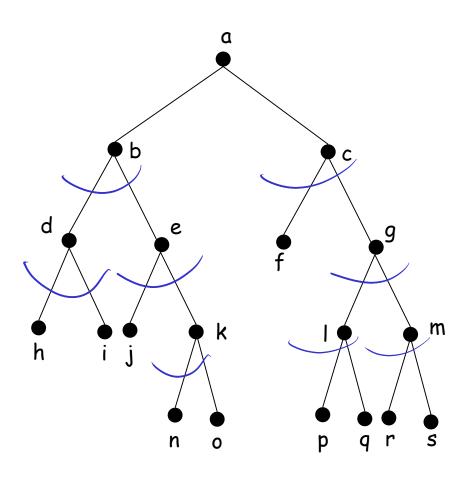
- La raíz del árbol
- El padre de i 👈
- Los hijos de d) K
- · Las hojas del árbol
- La profundidad de e 2
- La profundidad de a
- La aridad de i
- La aridad de a



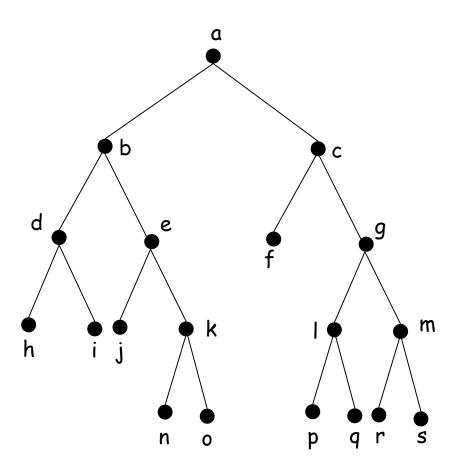
Dado el siguiente árbol indique:

- La raíz del árbol (a)
- El padre de i (b)
- Los hijos de d (j,k)
- Las hojas del árbol (j,k,e,f,g,h,m)
- La profundidad de e (2)
- La profundidad de a (0)
- La aridad de i (1)
- La aridad de a (3)

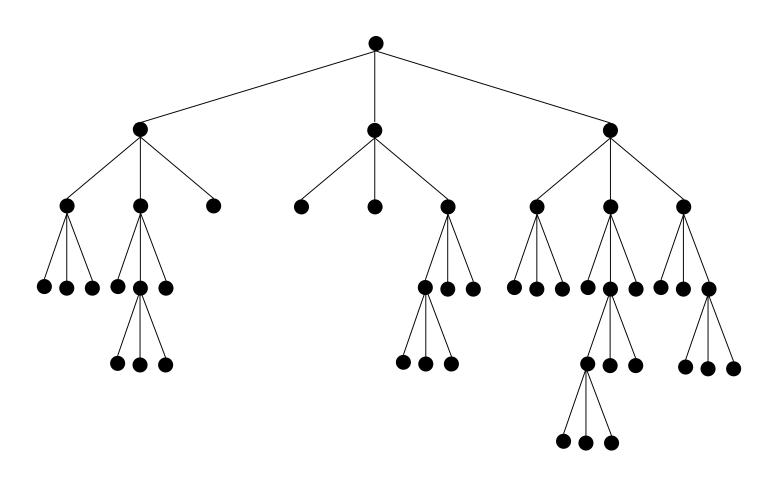


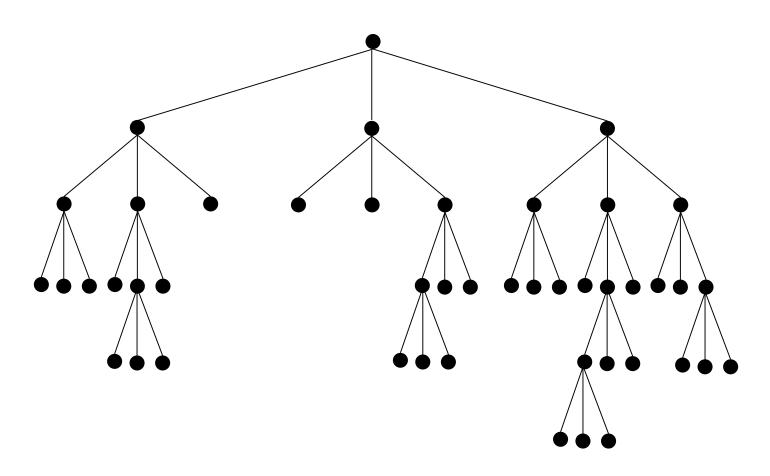


Cuál es la aridad de los nodos internos



Todos los nodos internos tienen la misma aridad

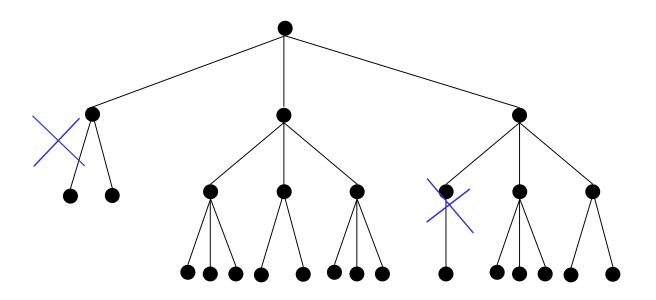


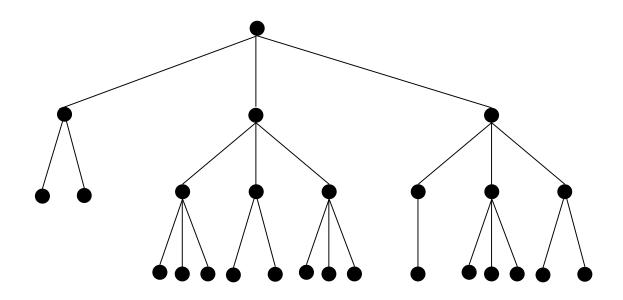


Todos los nodos internos tienen la misma aridad

Árbol n-ario completo

Un árbol es n-ario completo si cada nodo interno tiene exactamente n hijos

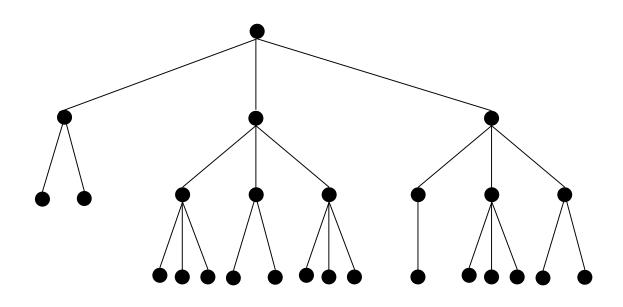




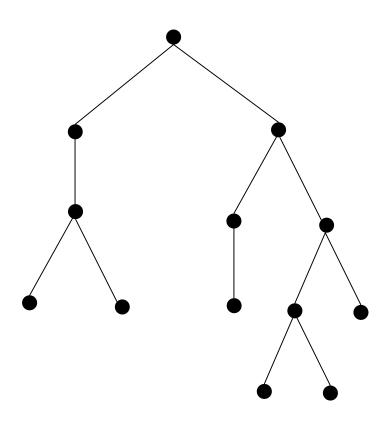
Cada nodo interno tiene aridad 1, 2, o 3

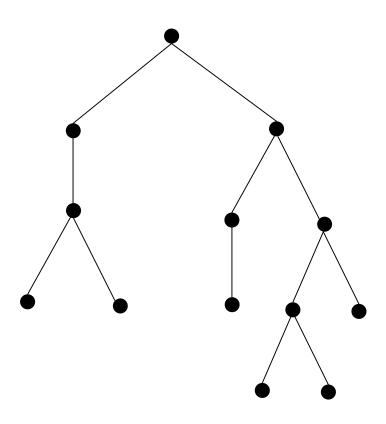
Árbol n-ario

Un árbol es n-ario si cada nodo interno no tiene más de n hijos

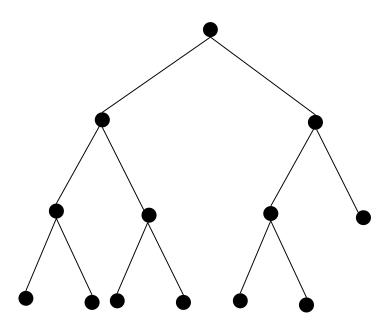


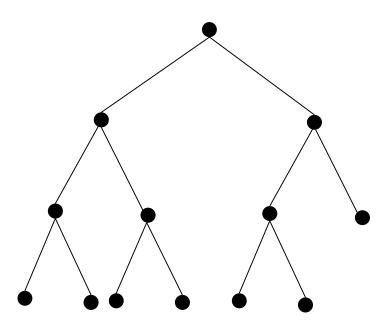
Árbol 3-ario



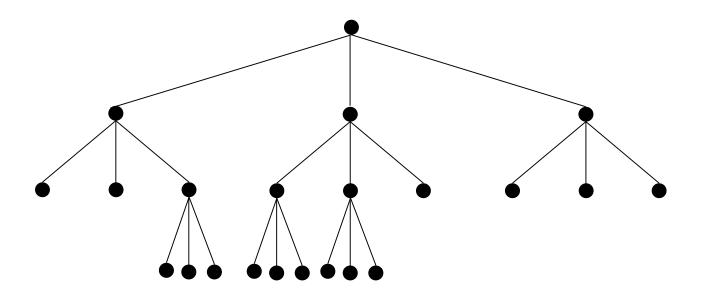


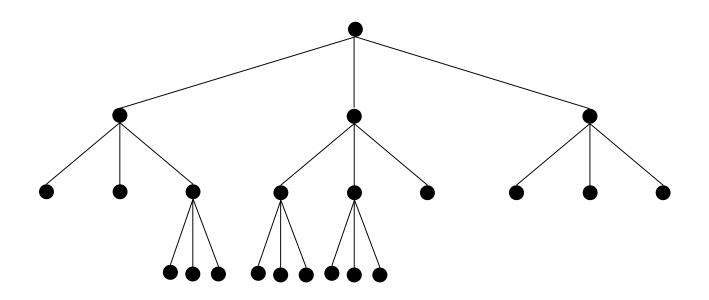
Árbol 2-ario o binario



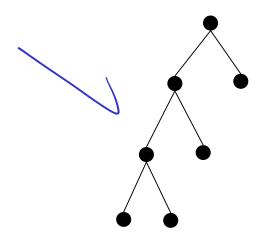


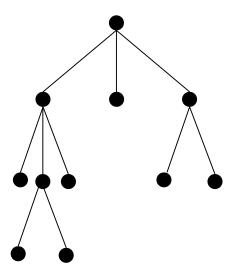
Árbol binario completo

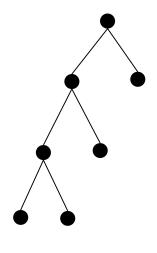




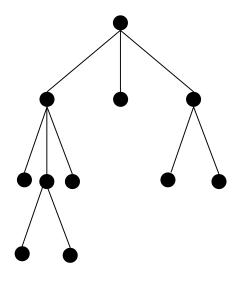
Árbol 3-ario completo



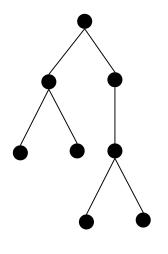


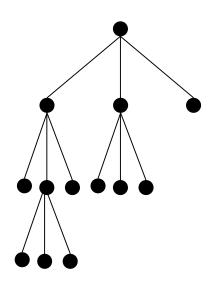


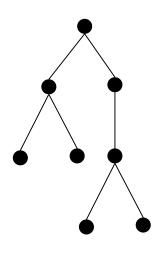
Árbol binario completo



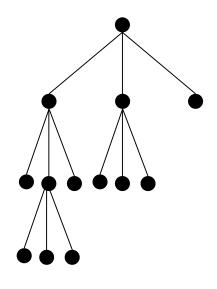
Árbol 3-ario



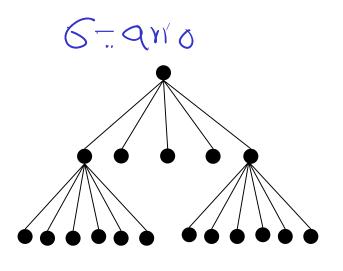


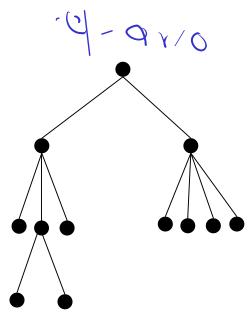


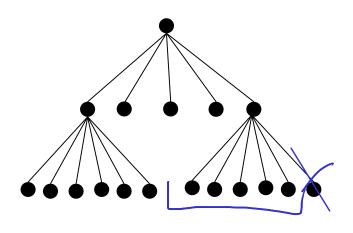
Árbol binario



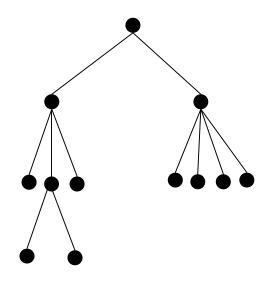
Árbol 3-ario completo





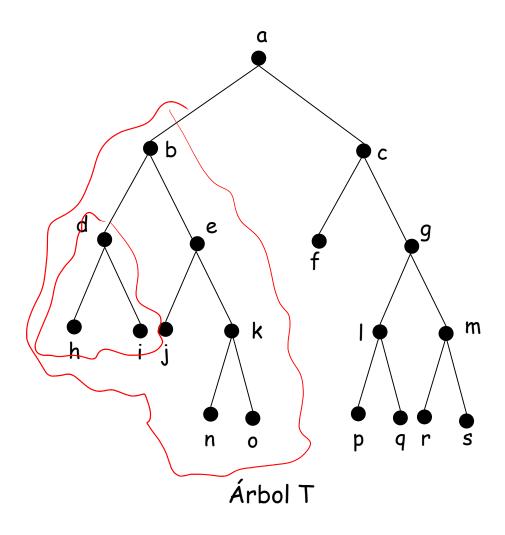


Árbol 5-ario completo

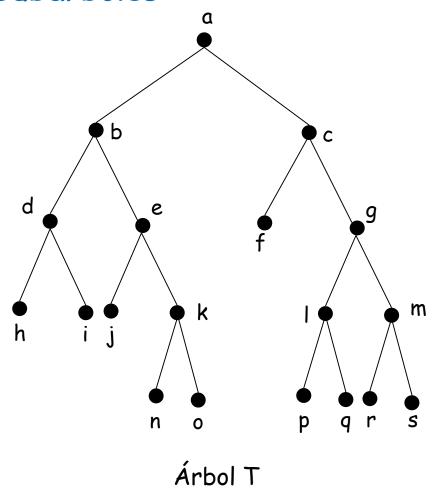


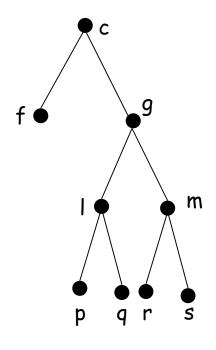
Árbol 4-ario

Subárboles



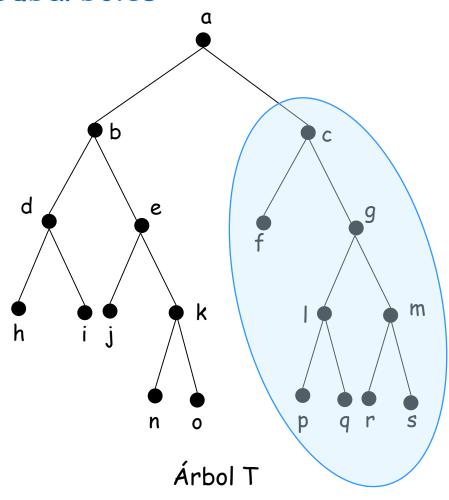
Subárboles

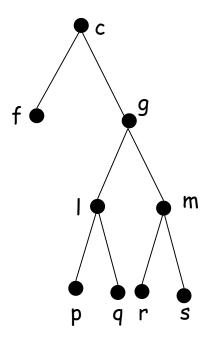




Subárbol de T con raíz en c

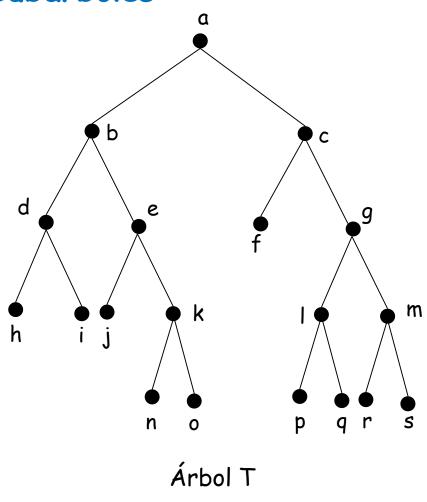
Subárboles

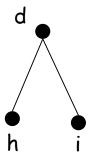




Subárbol de T con raíz en c

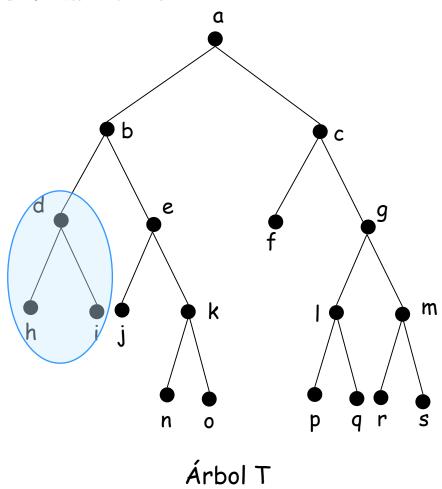
Subárboles

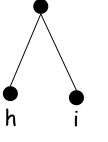




Subárbol de T con raíz en d

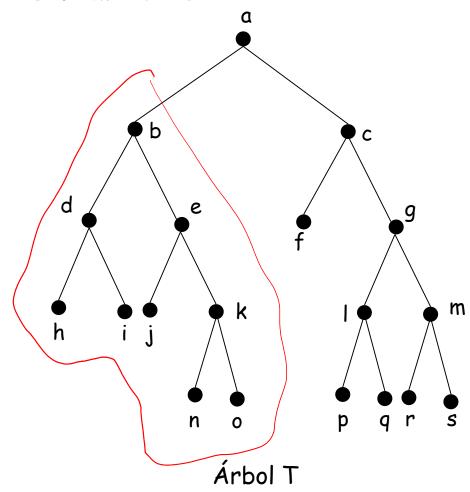
Subárboles





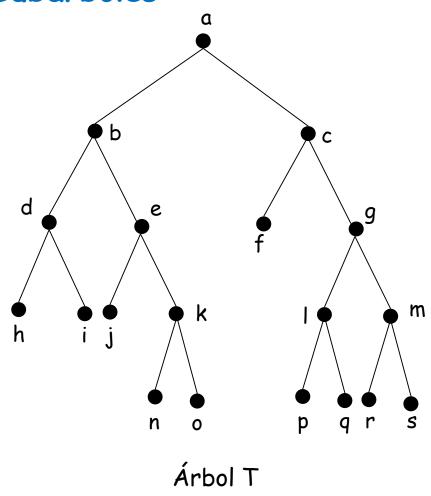
Subárbol de T con raíz en d

Subárboles



Subárbol de T con raíz en b

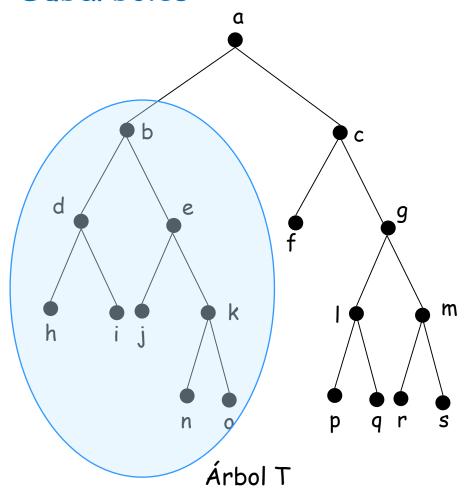
Subárboles

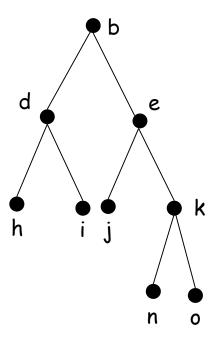


d e k h i j o o

Subárbol de T con raíz en b

Subárboles





Subárbol de T con raíz en b

Aplicaciones de los árboles

- Árboles de juego
- · Árboles binarios de búsqueda
- Árboles de decisión

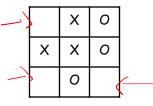
Árboles de juego

- Se modelan los posibles movimientos de cada jugador en un juego con adversario
- Sirve para analizar el efecto de las jugadas

Teoría de juegos

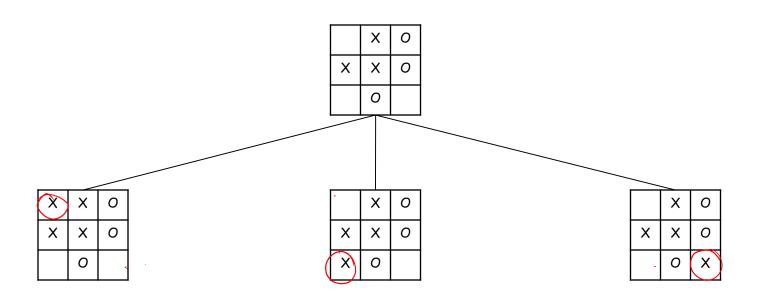
- · Construir el árbol de juego
- · Analizar el árbol

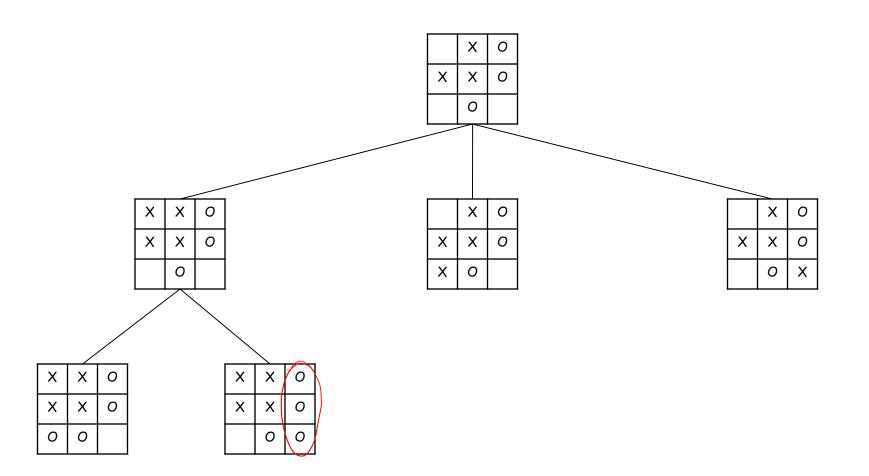
Construir el árbol de juego

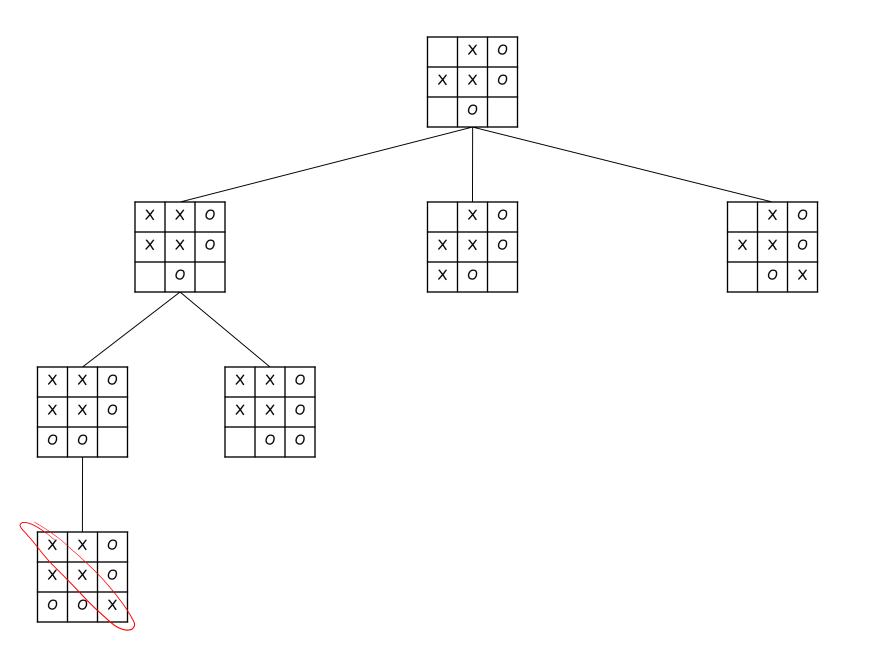


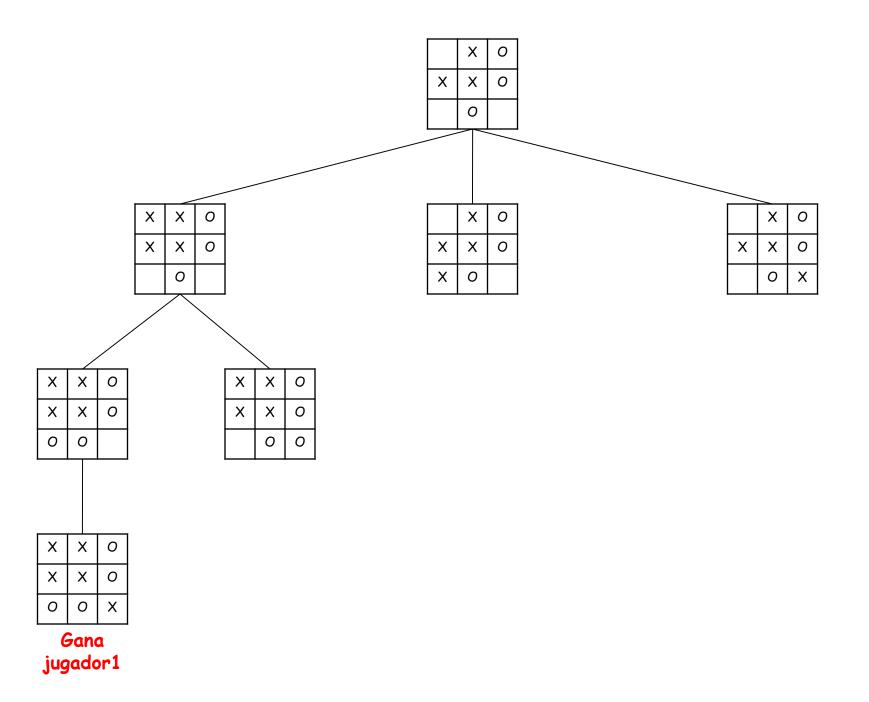
La jugada es de (X)

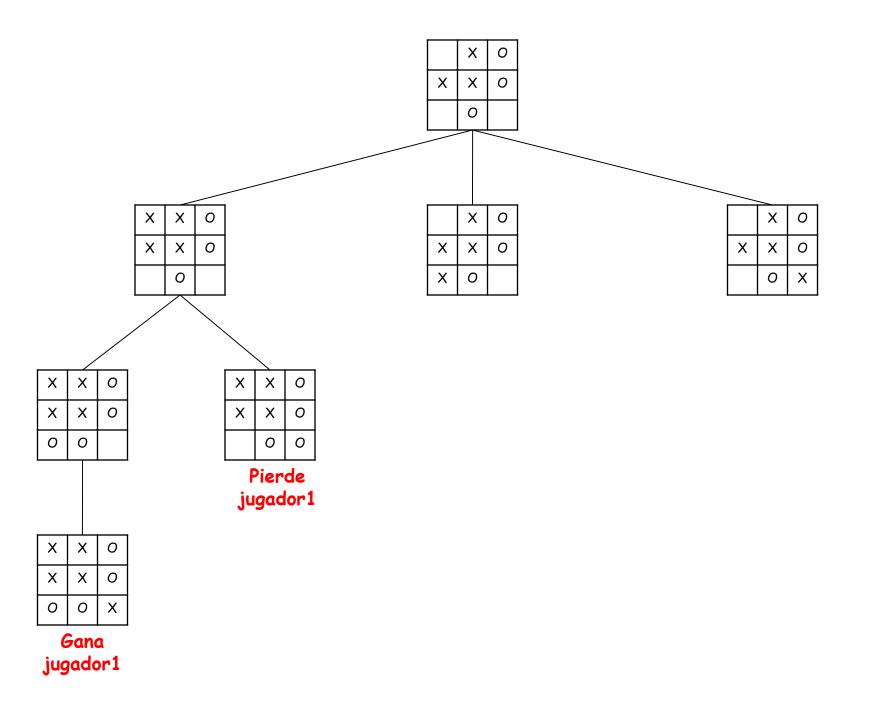
	×	0
X	×	0
	0	

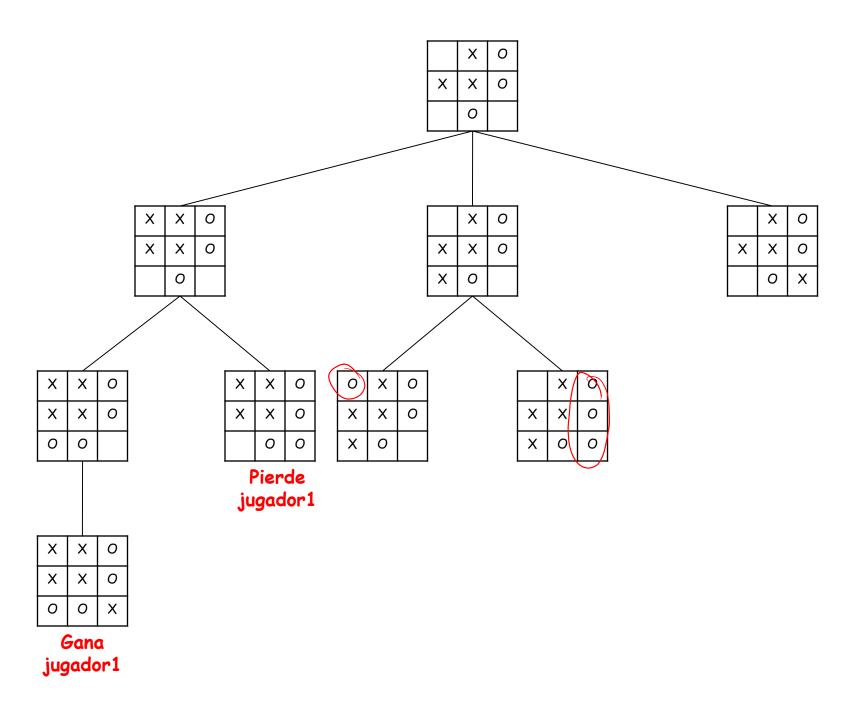


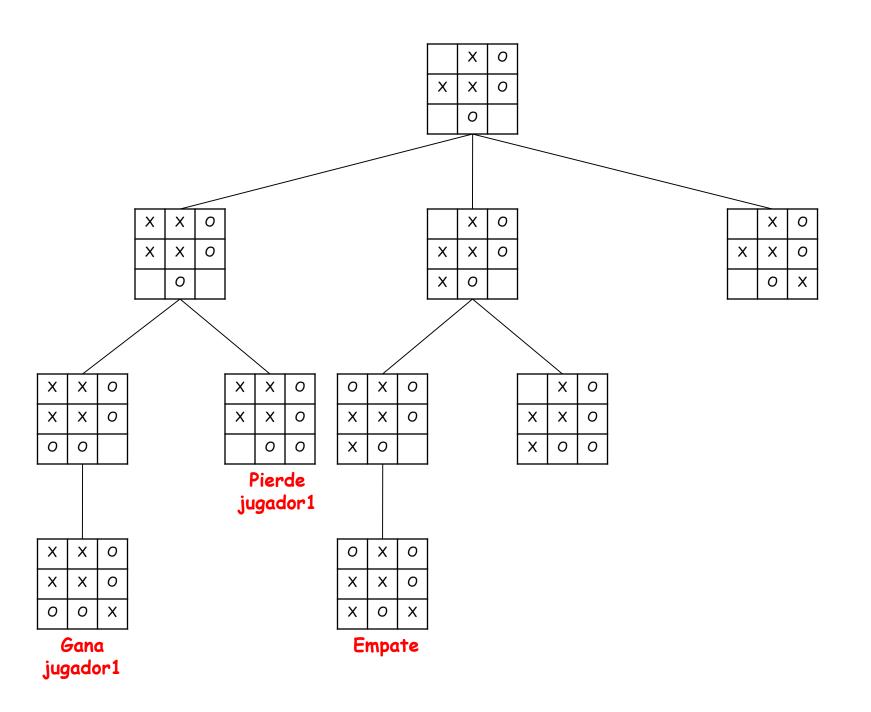


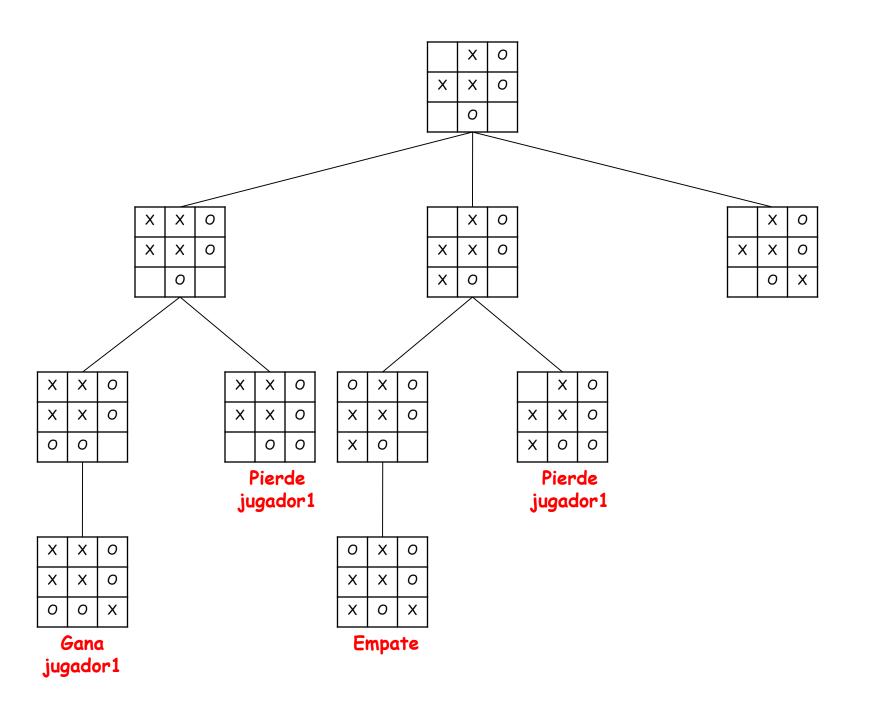


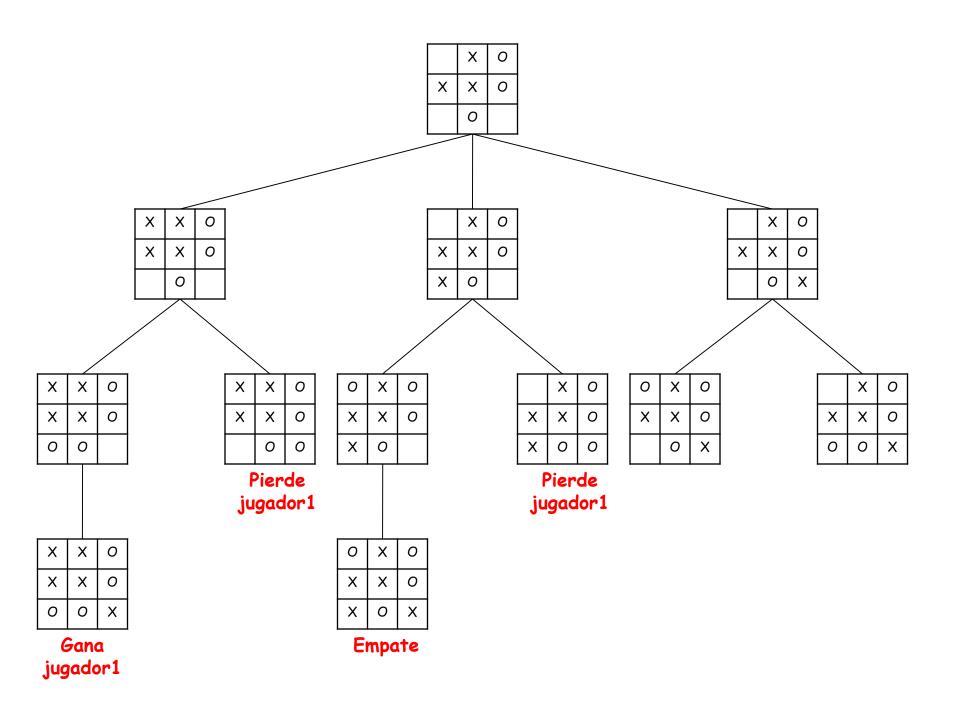


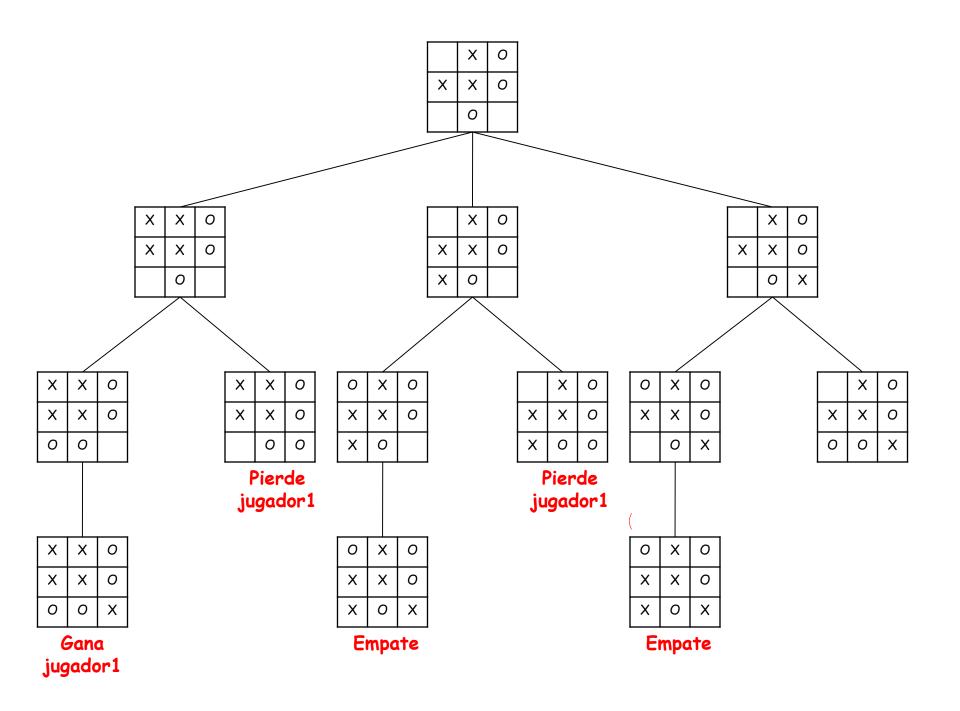


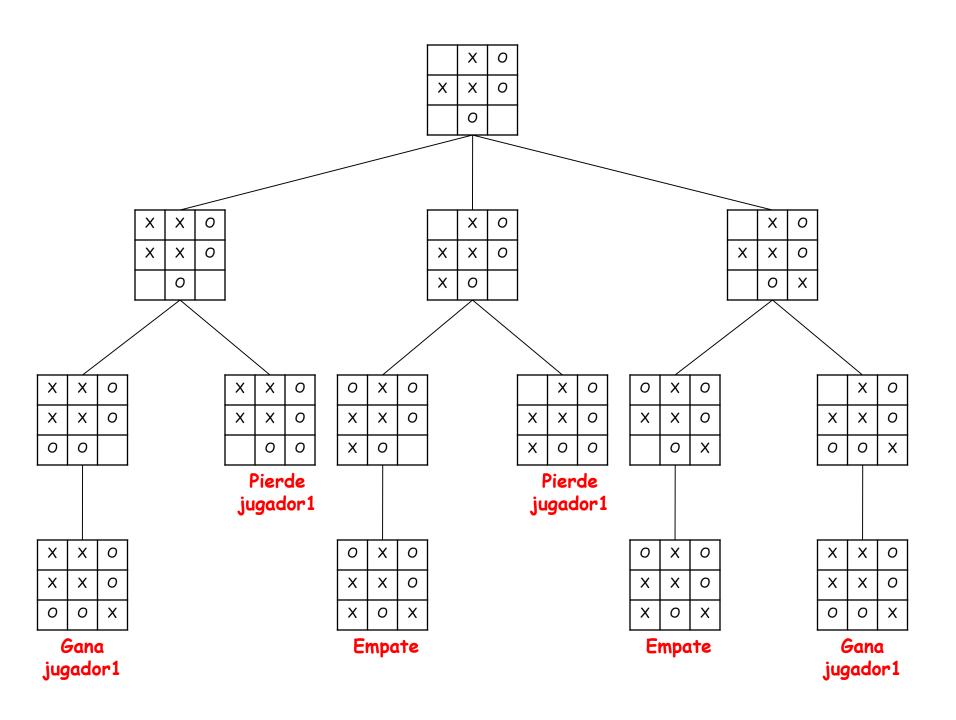


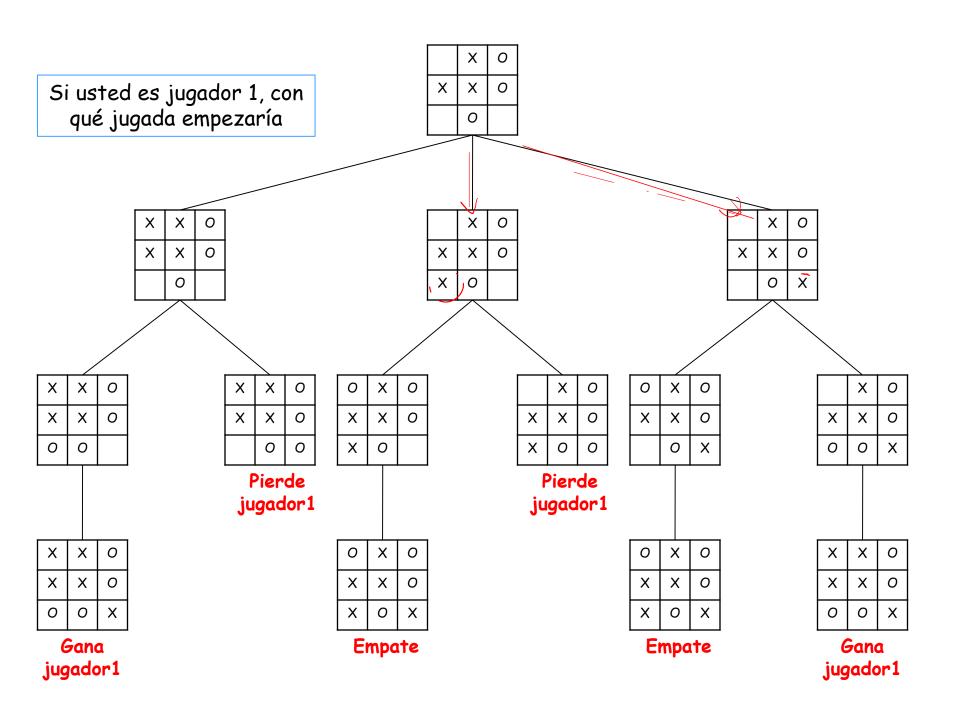


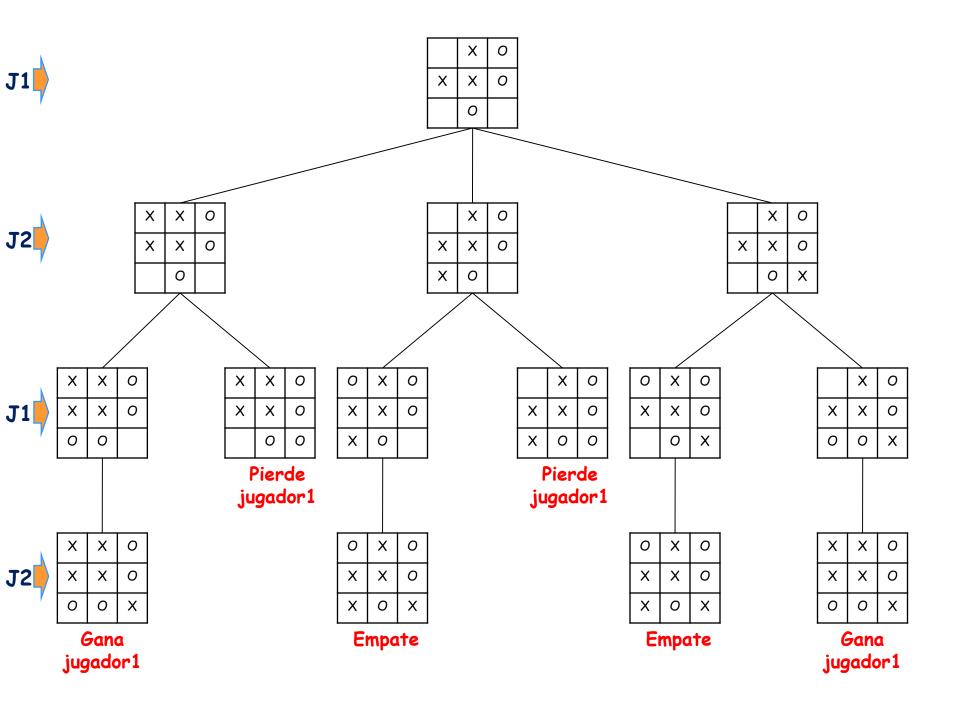
































Jugador2



. .

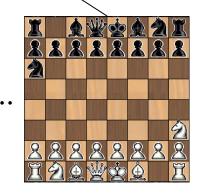


Jugador1

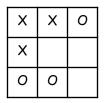




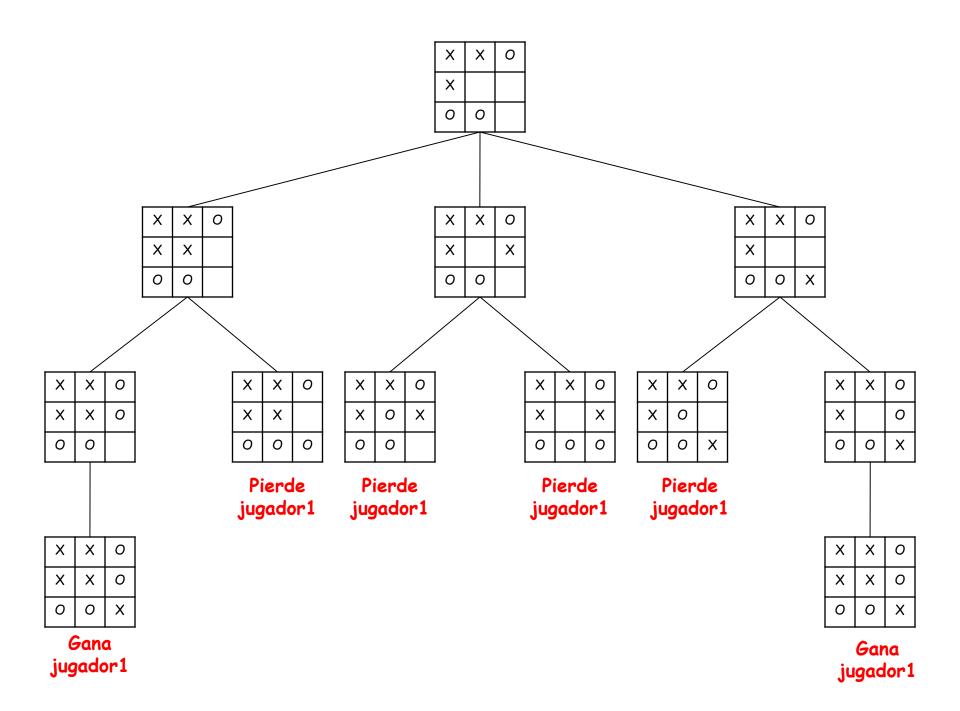


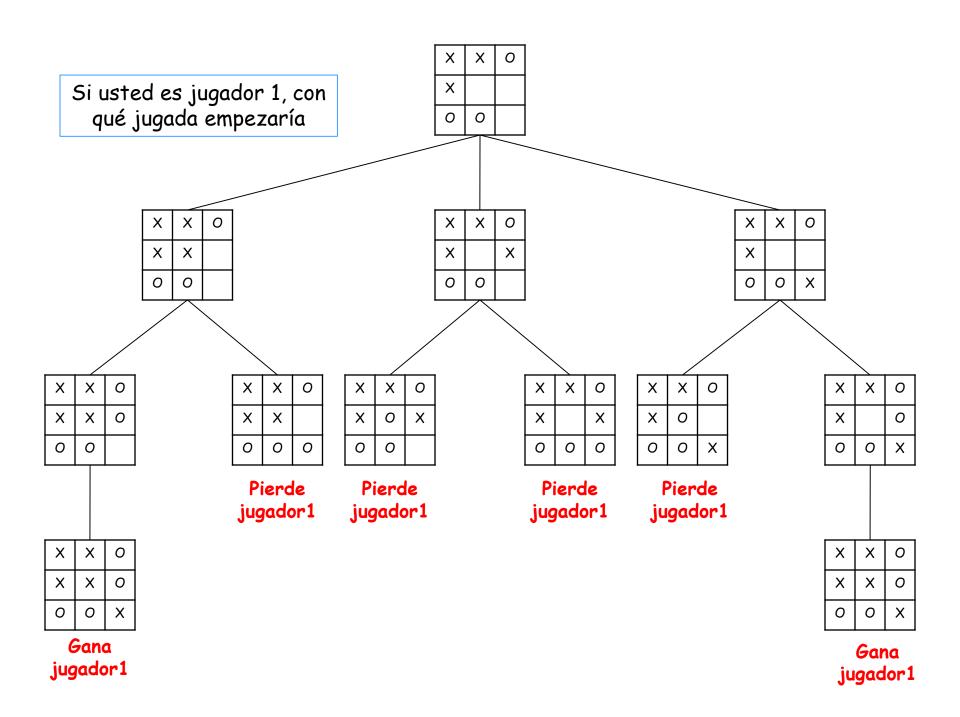


Construir el árbol de juego

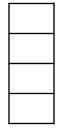


• La jugada es de (X)

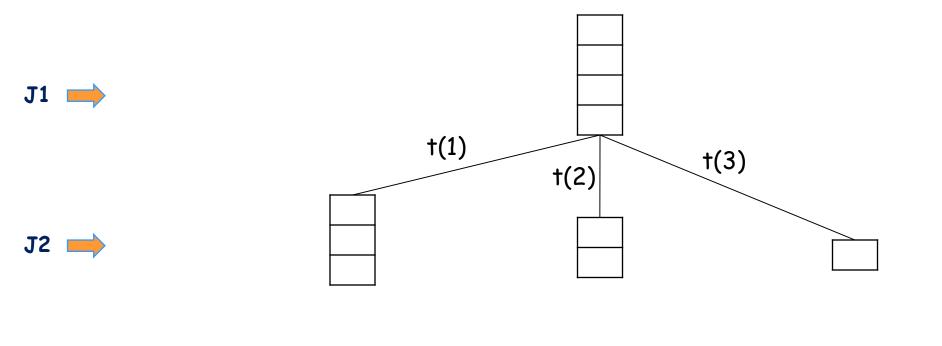


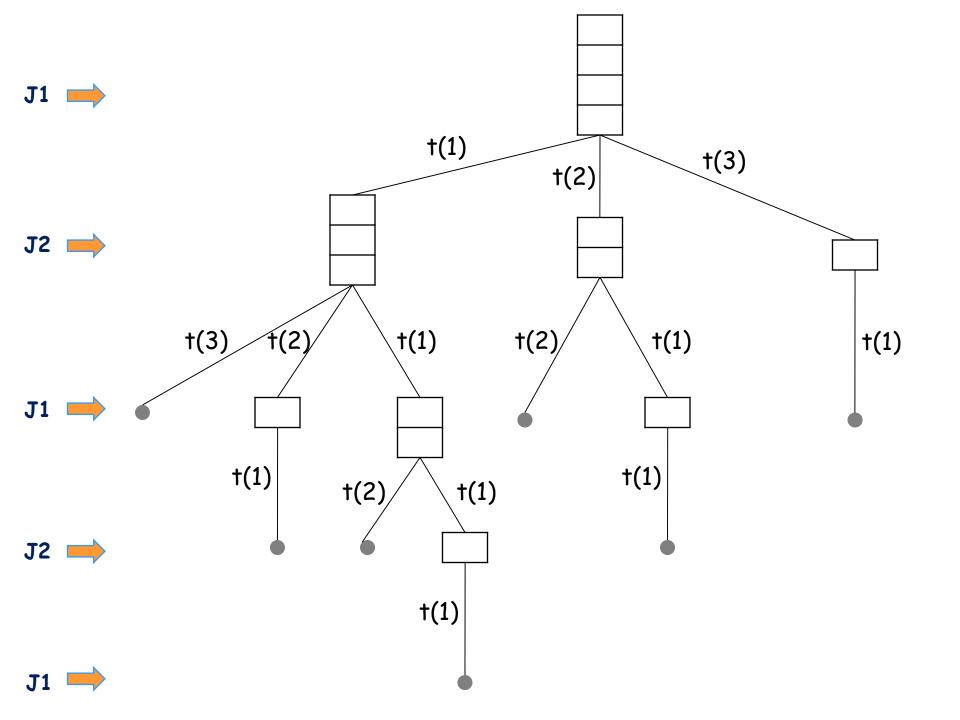


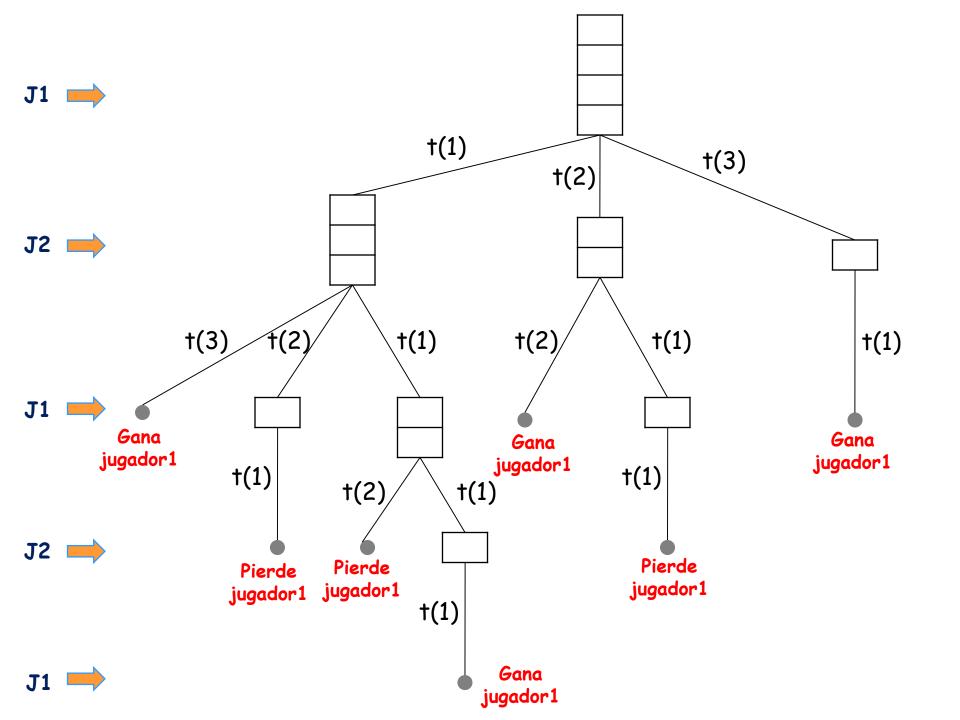
Construir el árbol de juego



• El juego del NIM. Se tiene una pila de 4 fichas de la cual cada jugador puede tomar 1, 2 ó 3. El objetivo de cada jugador es obligar a su adversario a tomar la última ficha. Como los elementos están apilados, solo se pueden tomar fichas de su tope





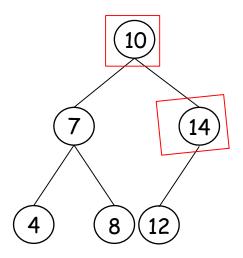


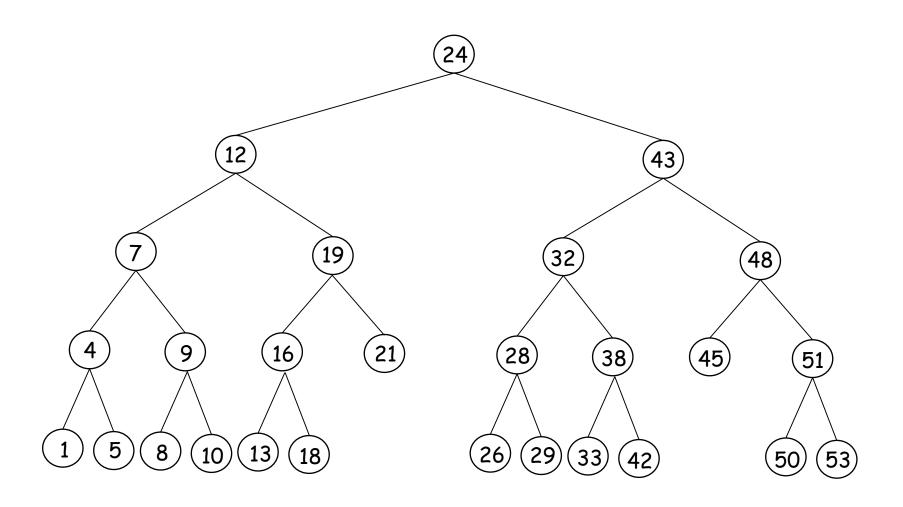
Aplicaciones de los árboles

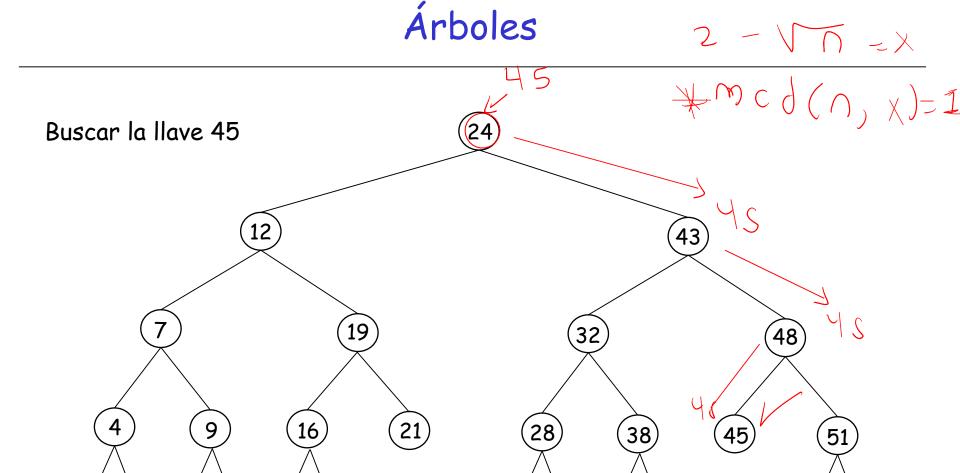
- Árboles de juego
- · Árboles binarios de búsqueda
- Árboles de decisión

Árboles binarios de búsqueda

Es un árbol binario en el que cada vértice tiene una llave. La llave de un vértice es mayor que las llaves de los vértices del subárbol izquierdo y menor que las llaves de los vértices del subárbol derecho



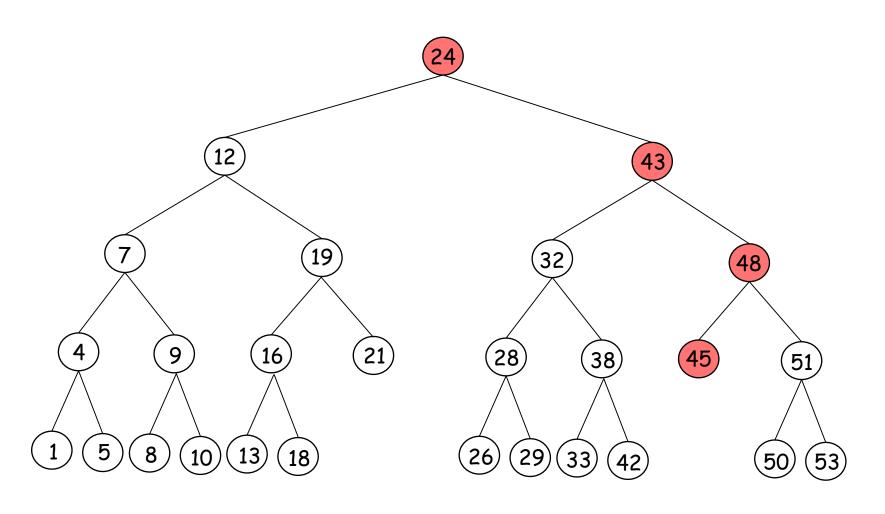




10

(29)

(33)



Se necesitan 4 comparaciones

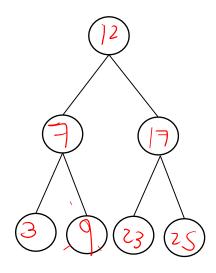
Buscar la llave 45



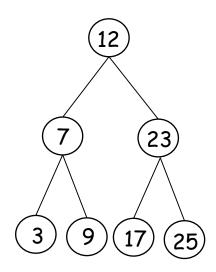


Se necesitan 14 comparaciones

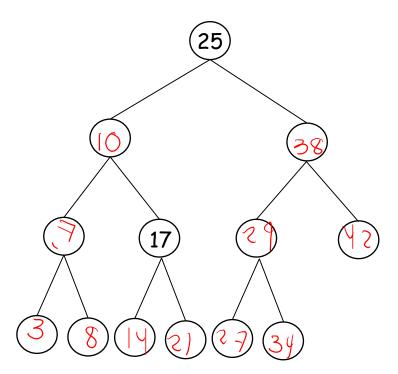
Crear un árbol de búsqueda binaria con las llaves 3,7,9,12,17,23,25 que tenga la siguiente estructura



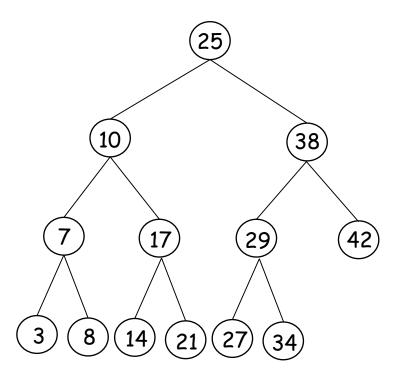
Crear un árbol de búsqueda binaria con las llaves 3,7,9,12,17,23,25 que tenga la siguiente estructura



Crear un árbol de búsqueda binaria con las llaves 3,7,8,10,14,17,21,25,27,29,34,38,42 que tenga la siguiente estructura



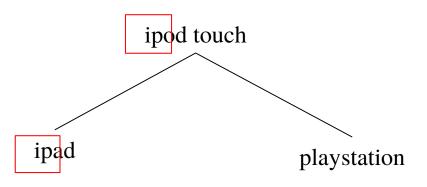
Crear un árbol de búsqueda binaria con las llaves 3,7,8,10,14,17,21,25,27,29,34,38,42 que tenga la siguiente estructura



Crear un árbol de búsqueda binaria con las siguientes llaves:

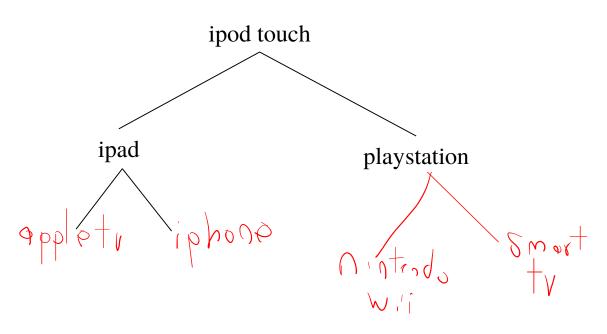
- apple tv
- ipad
- iphone
- ipod touch
- nintendo wii
- playstation
- smart tv

apple tv
ipad
iphone
ipod touch
nintendo wii
playstation
smart tv

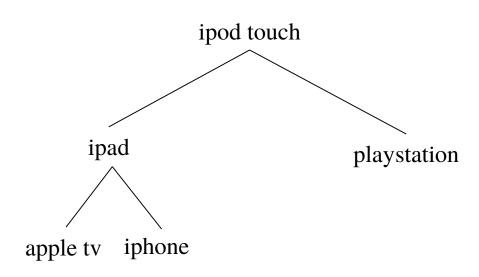


9x bolos B-trer

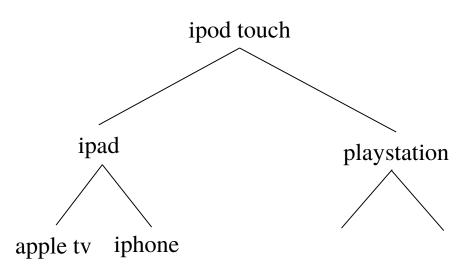
apple tv
ipad
iphone
ipod touch
nintendo wii
playstation
smart tv



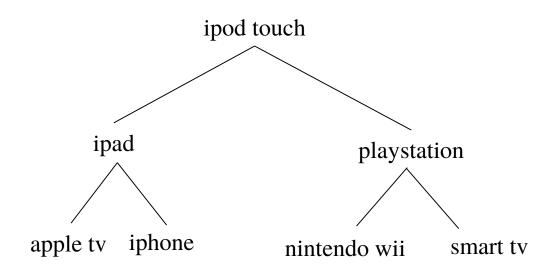
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ipad
iphone
ipod touch
nintendo wii
playstation
smart tv

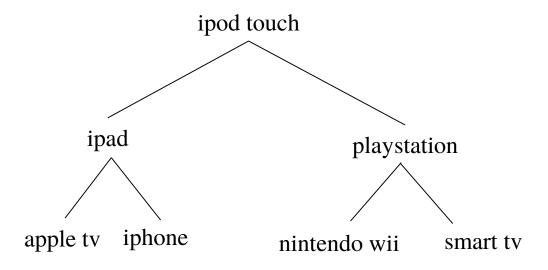


apple tv
ipad
iphone
ipod touch
nintendo wii
playstation
smart tv

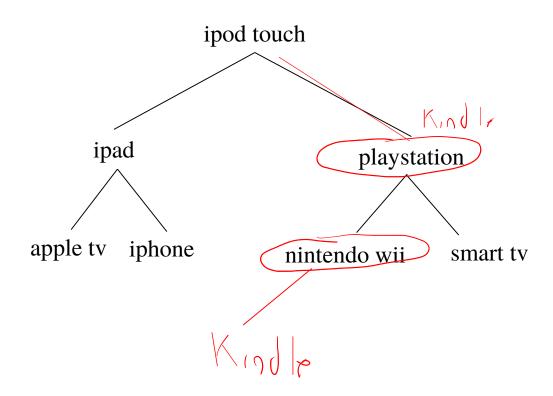


apple tv
ipad
iphone
ipod touch
nintendo wii
playstation
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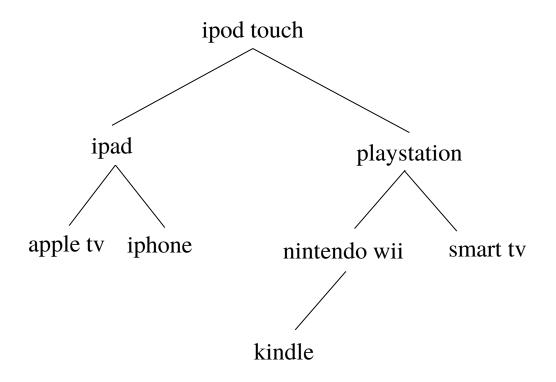




Insertar la llave kindle



Insertar la llave kindle

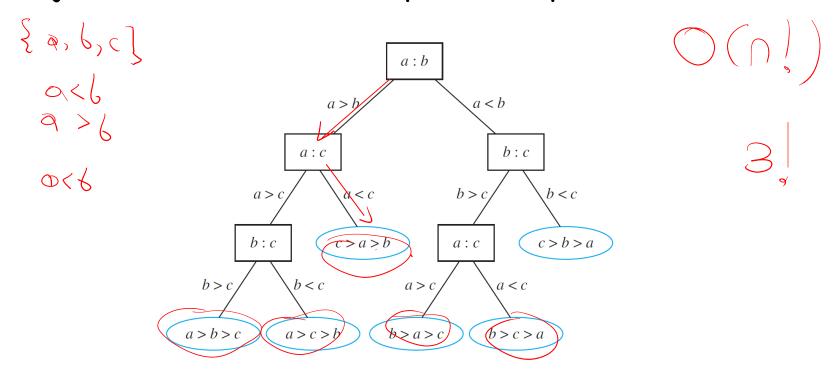


Aplicaciones de los árboles

- Árboles de juego
- · Árboles binarios de búsqueda
- Árboles de decisión

Árboles de decisión

Es un árbol en cuyos vértices se tienen condiciones y en las hojas decisiones sobre un problema particular



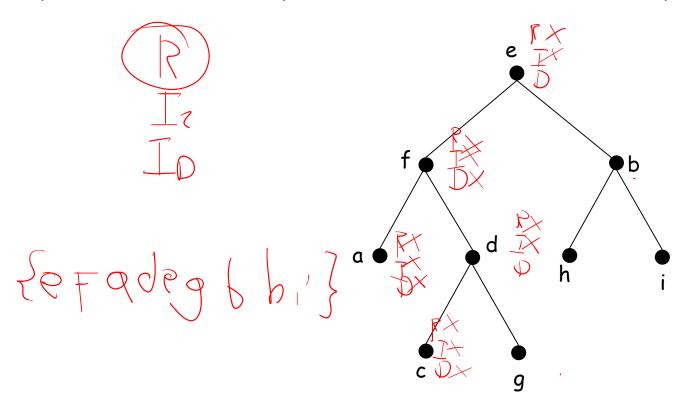
Árbol de decisión para ordenar 3 elementos

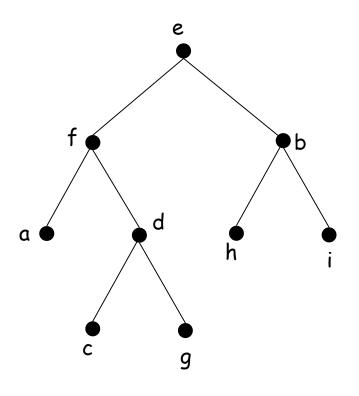
Recorridos de los árboles

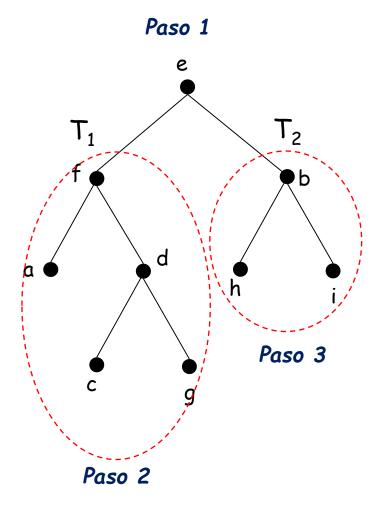
- Preorden
- Inorden
- Postorden

Recorridos en preorden

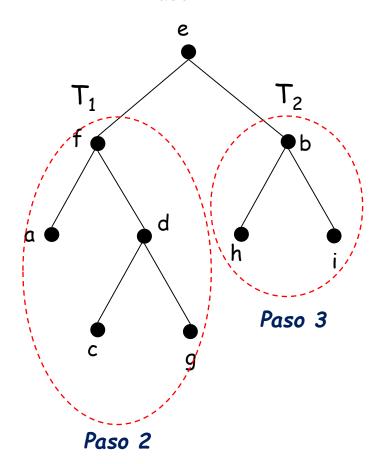
Sea T un árbol con raíz r y subárboles T_1 , T_2 , ..., T_n . El recorrido en preorden se hace visitando r, luego T_1 en preorden, T_2 en preorden, así hasta T_n en preorden





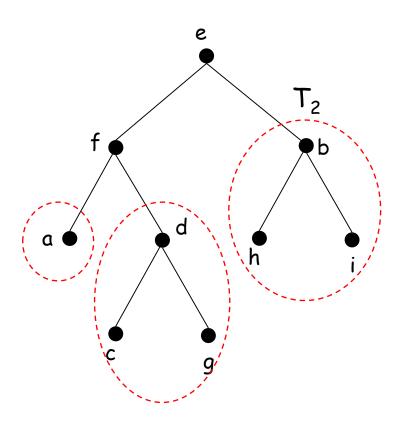


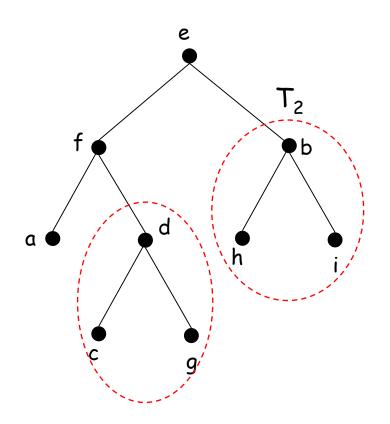


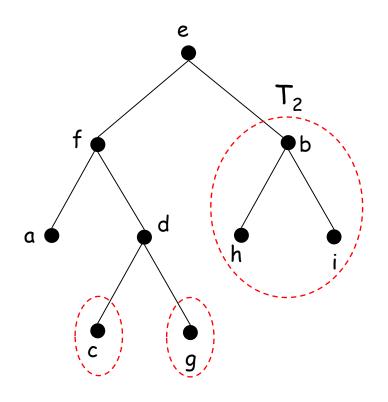


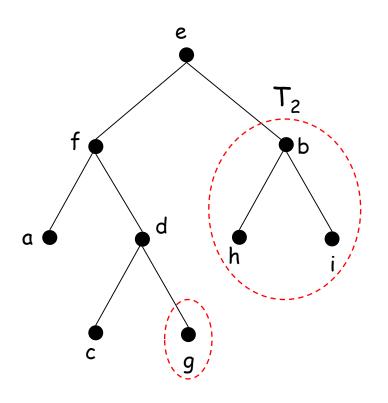
Recorrido en preorden:

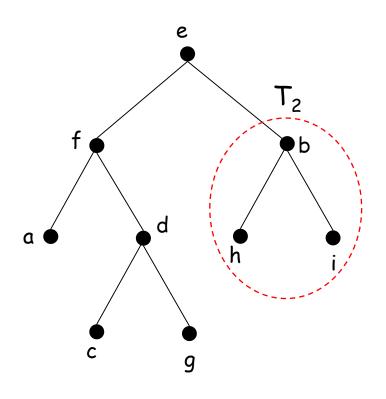
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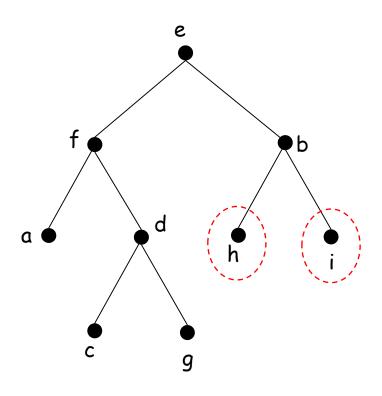


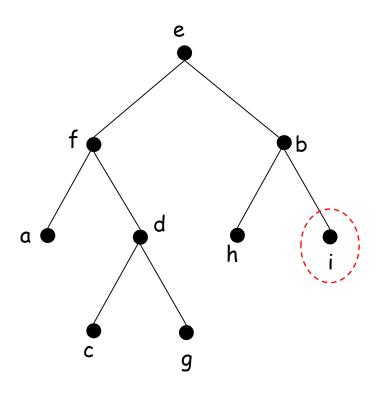


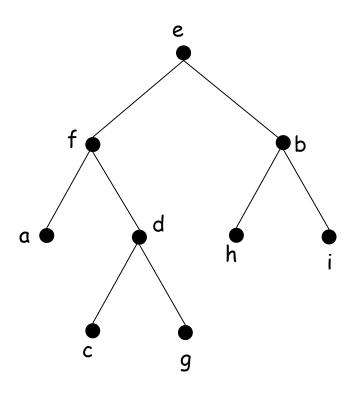


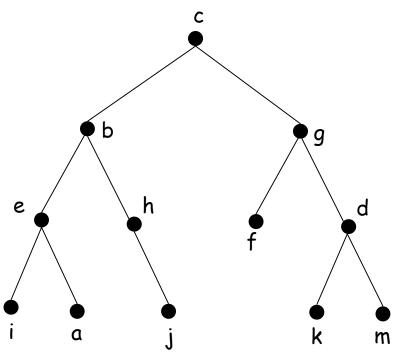


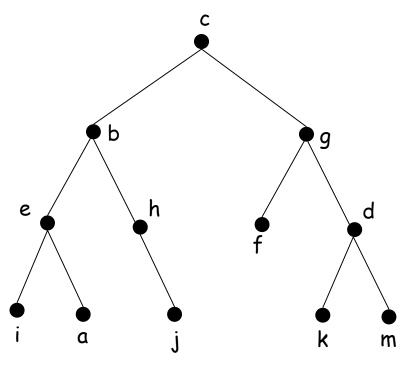


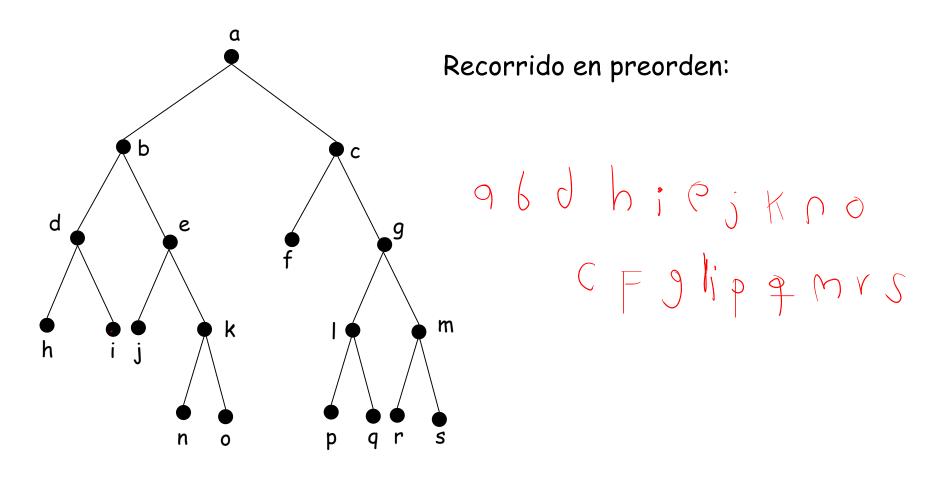


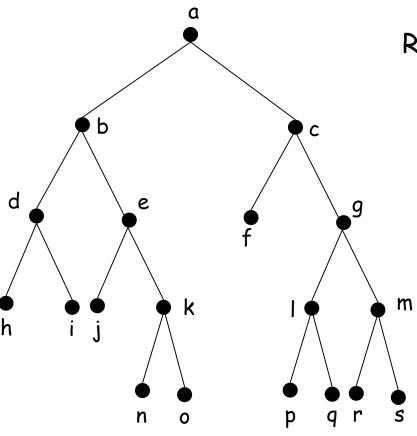






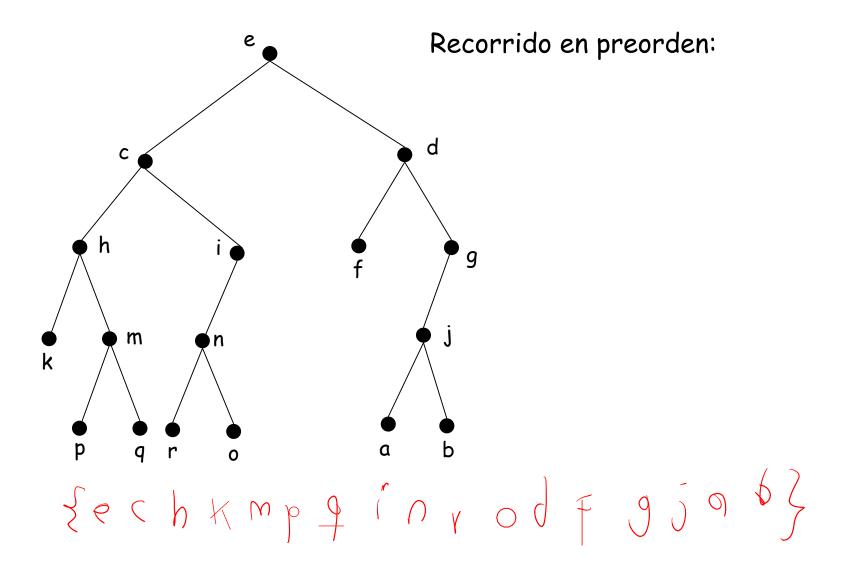


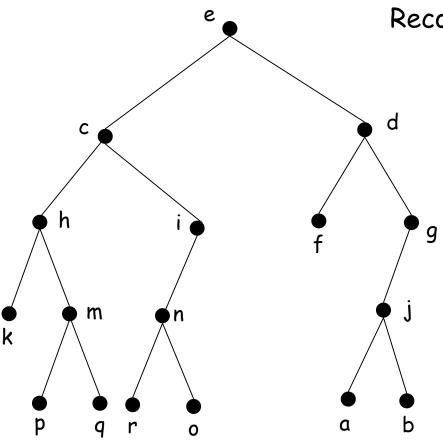




Recorrido en preorden:

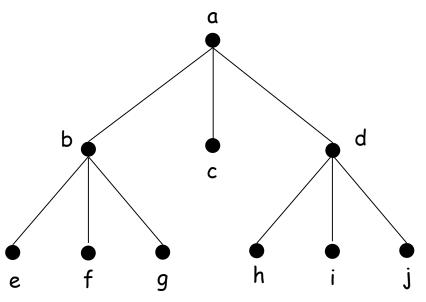
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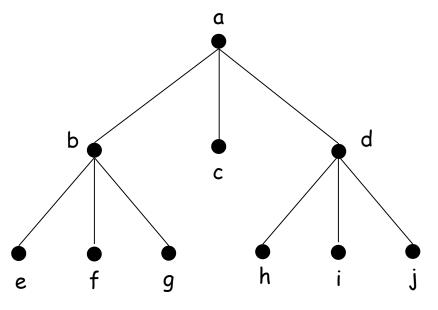


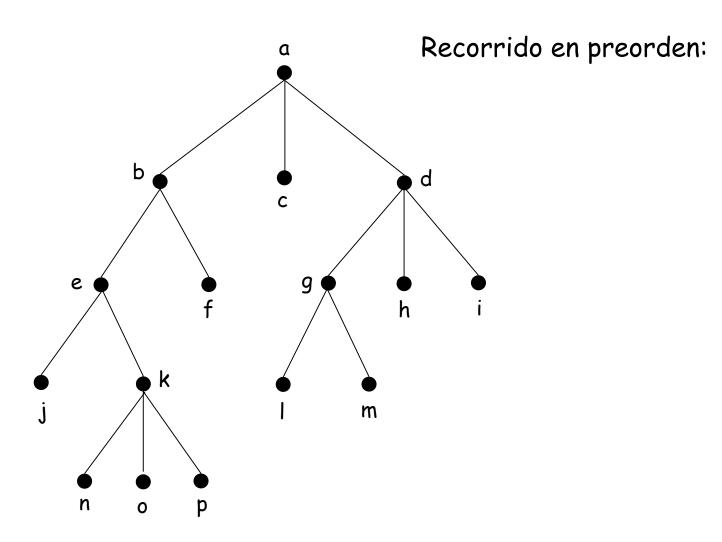


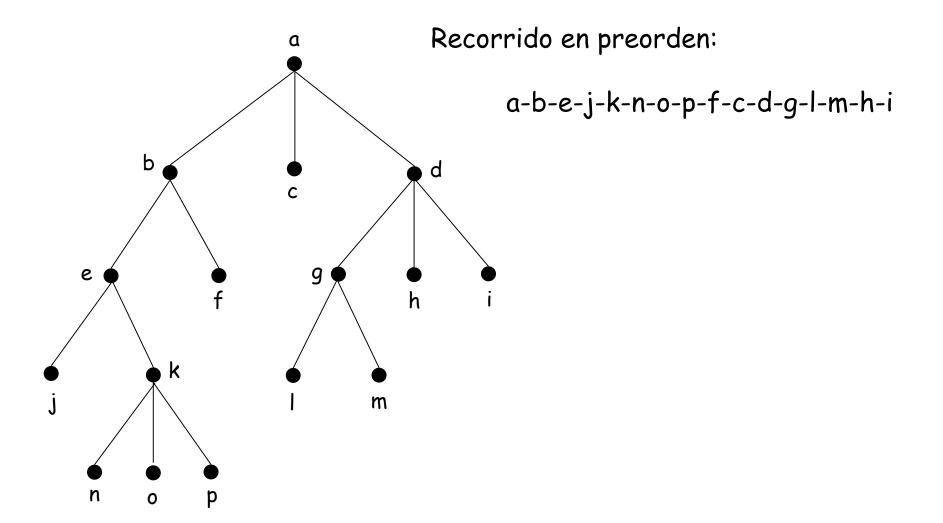
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e-c-h-k-m-p-q-i-n-r-o-d-f-g-j-a-b









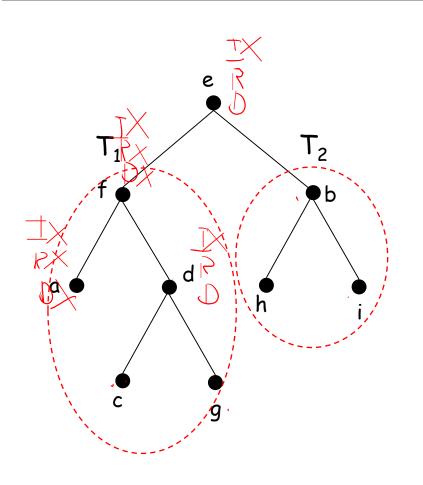
Recorridos de los árboles

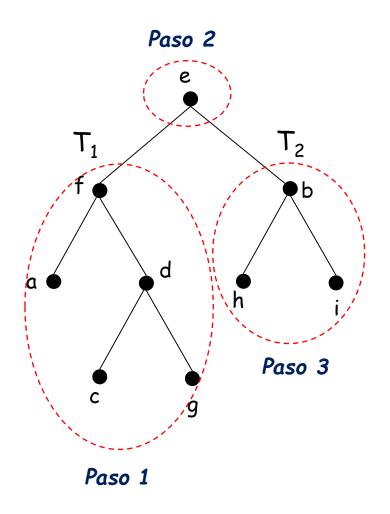
- Preorden

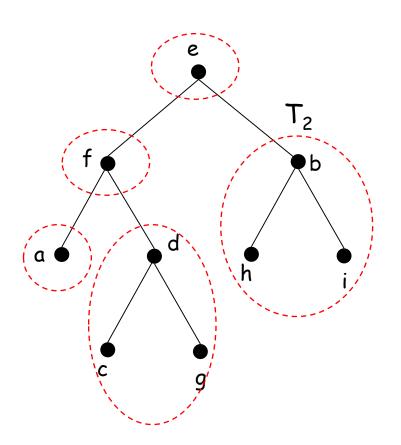
- Postorden

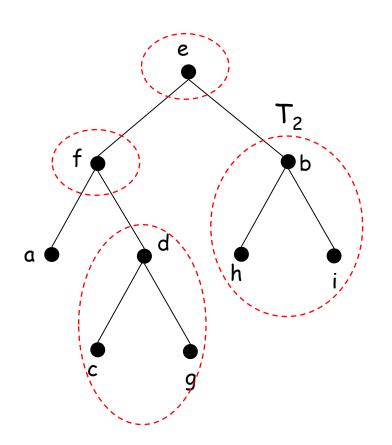
Recorridos en inorden

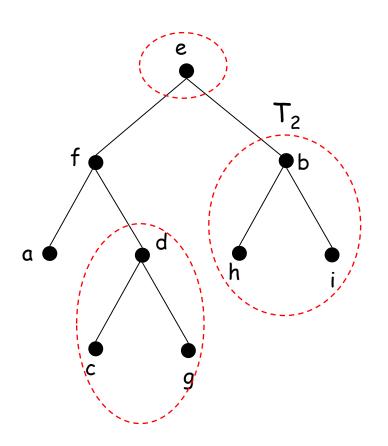
Sea T un árbol con raíz r y subárboles T_1 , T_2 , ..., T_n . El recorrido en inorden se hace realizando el recorrido de T_1 en inorden, luego visitando r, luego se hace el recorrido de T_2 en inorden, así hasta T_n en inorden

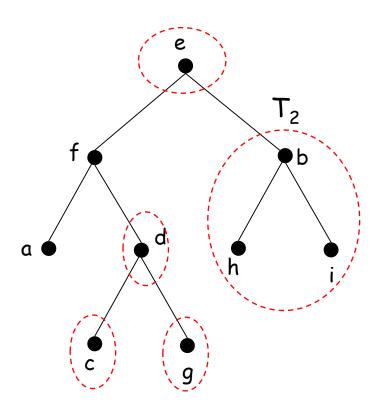


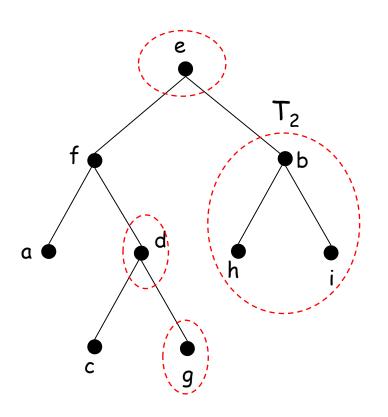


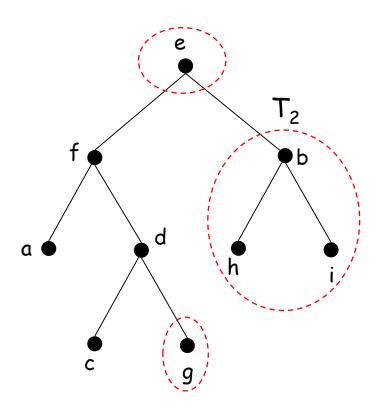


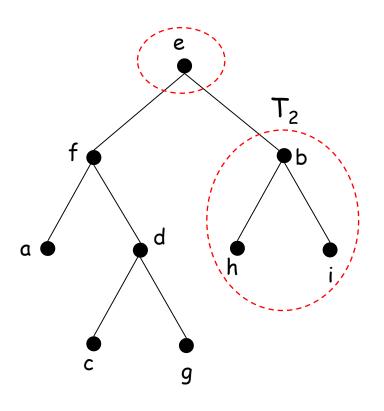


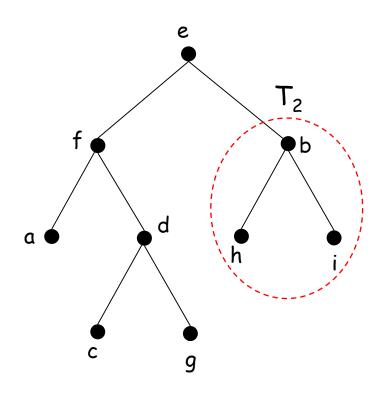


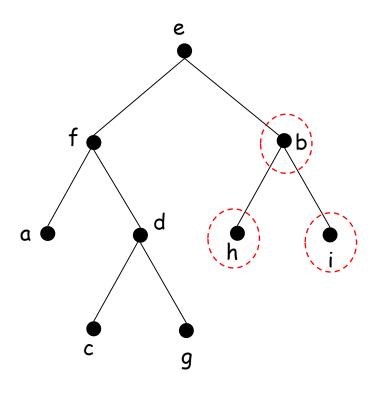


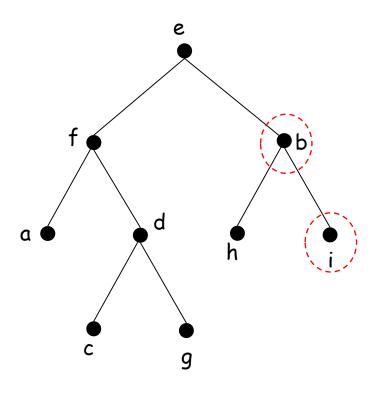


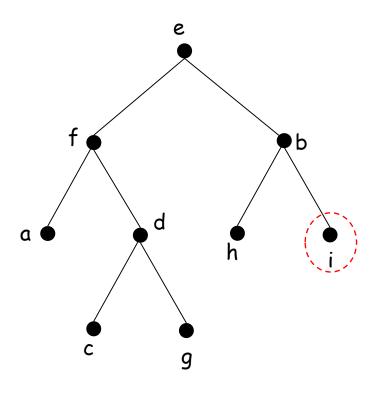


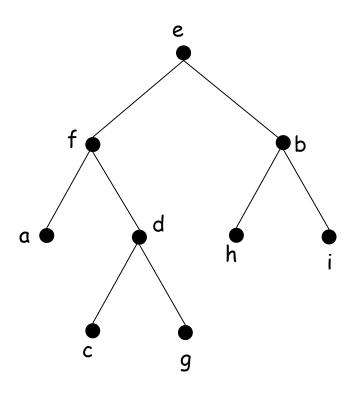


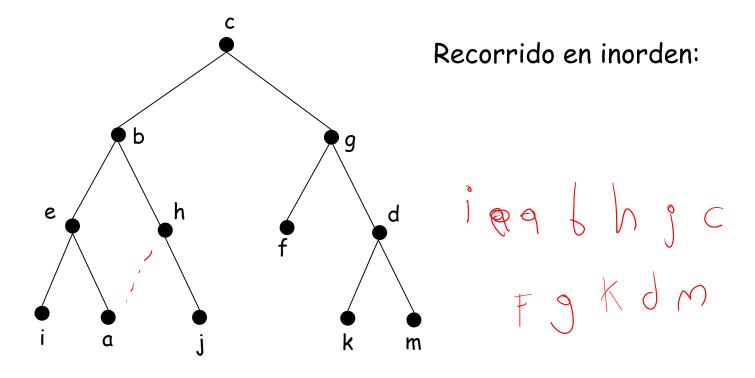


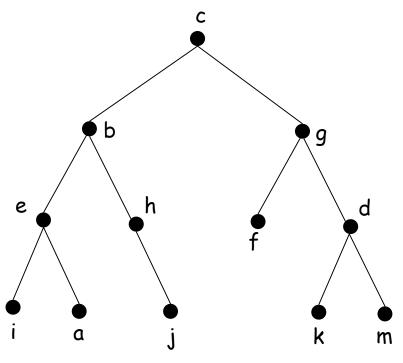


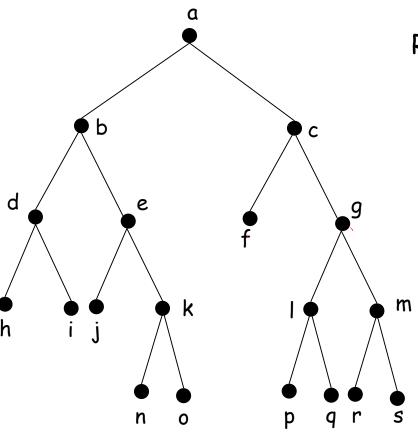


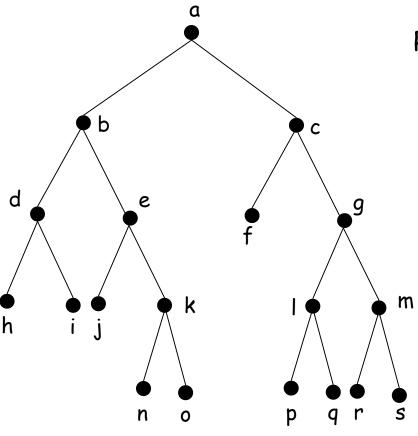






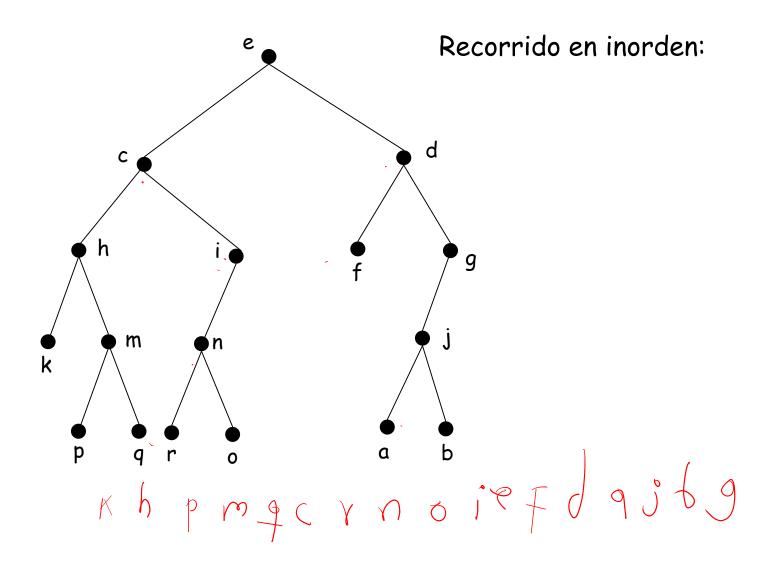


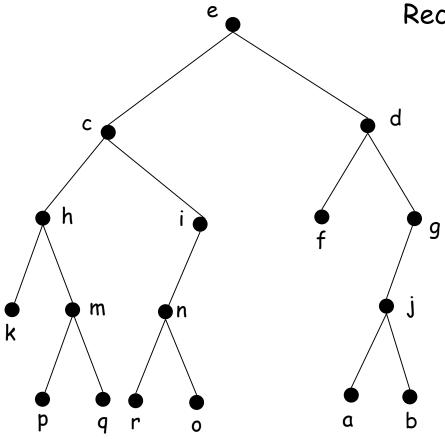




Recorrido en inorden:

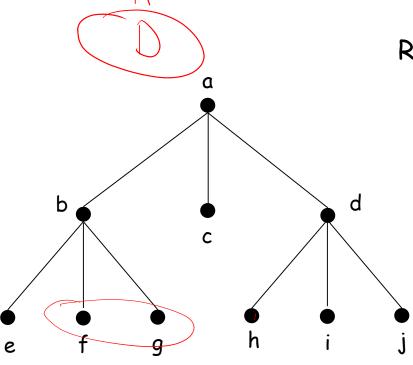
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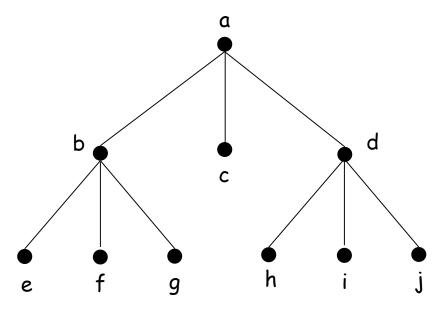


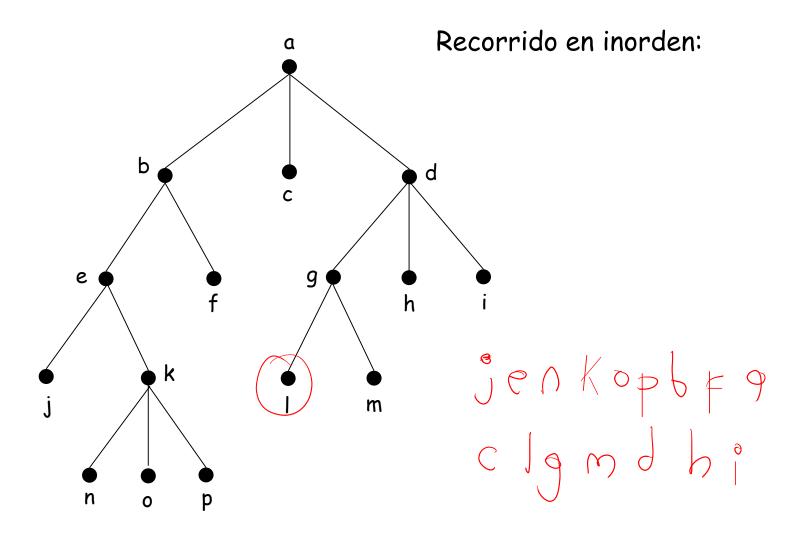


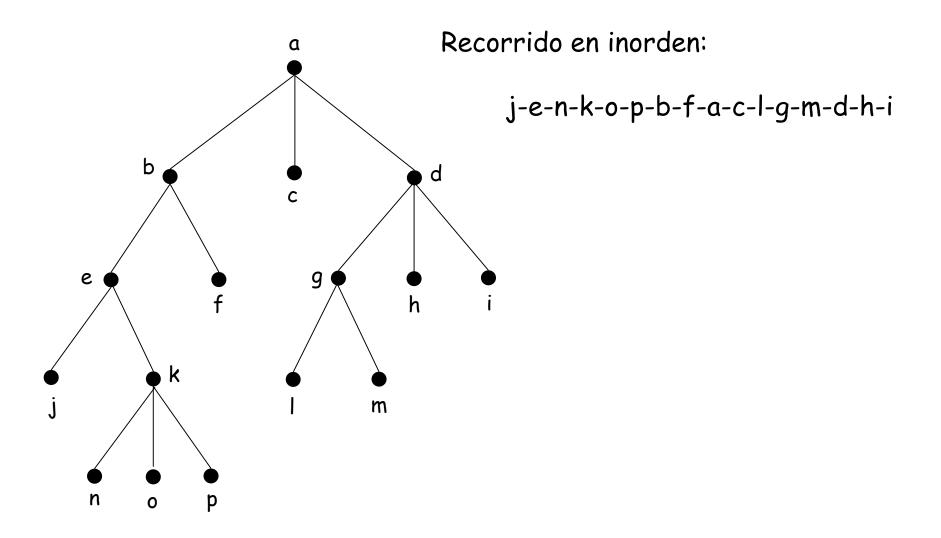
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k-h-p-m-q-c-r-n-o-i-e-f-d-a-j-b-g









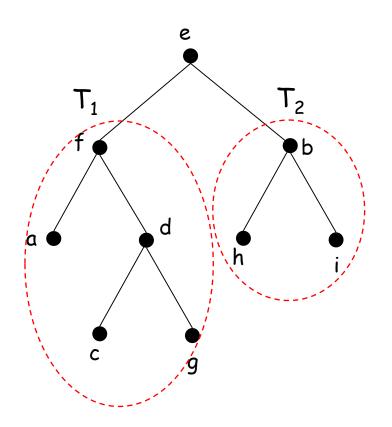
Recorridos de los árboles

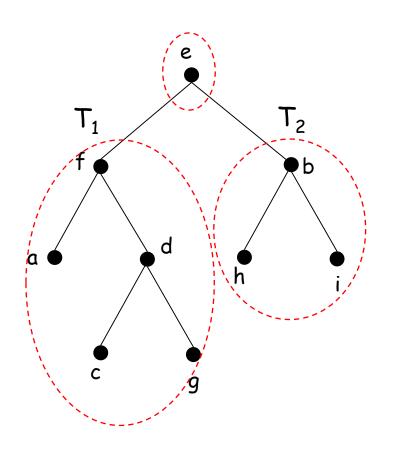
- Preorden
- Inorden
- Postorden —

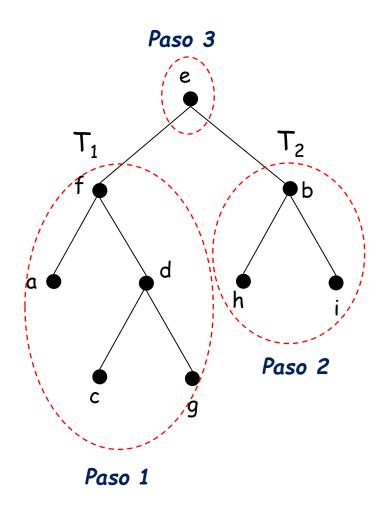


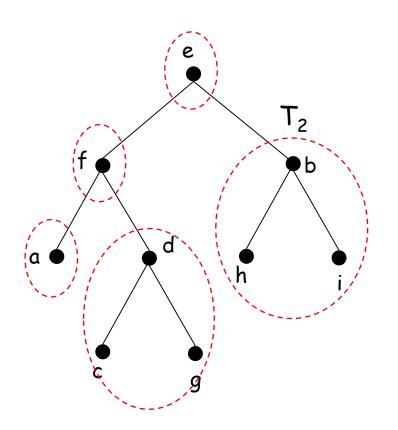
Recorridos en postorden

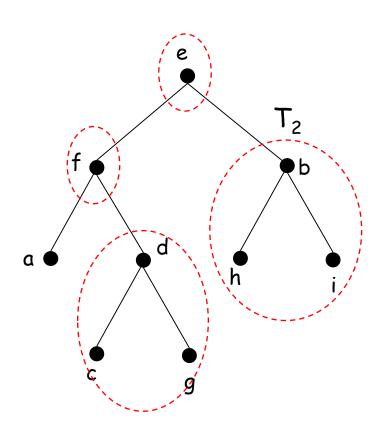
Sea T un árbol con raíz r y subárboles T_1 , T_2 , ..., T_n . El recorrido en postorden se hace realizando el recorrido de T_1 en postorden, T_2 en postorden, hasta T_n en postorden, y luego visitando r





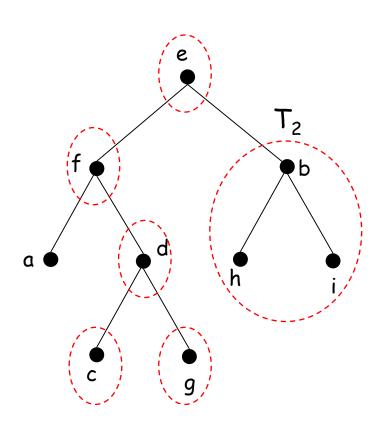






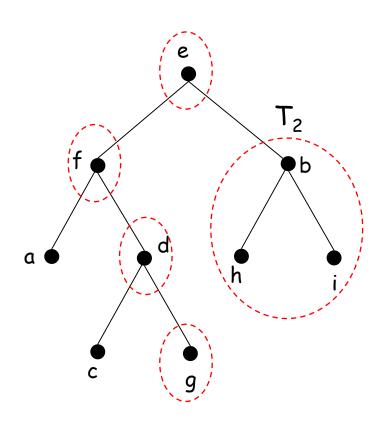
Recorrido en postorden:

a



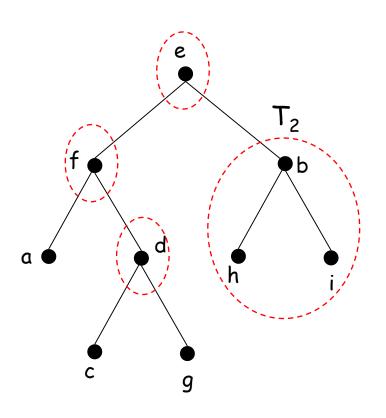
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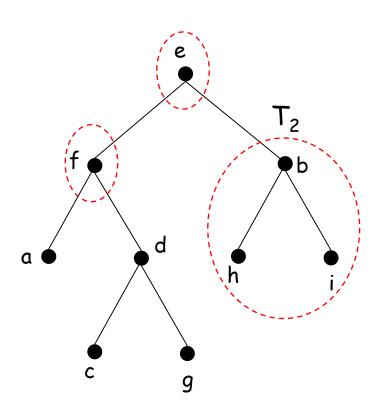
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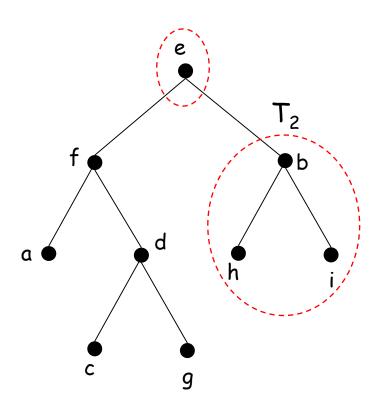


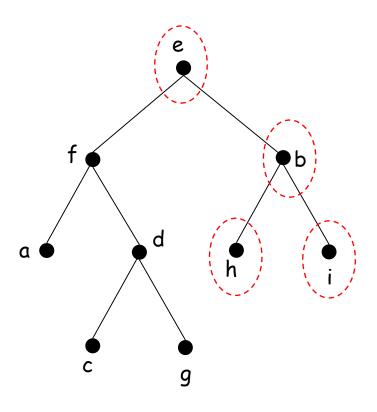
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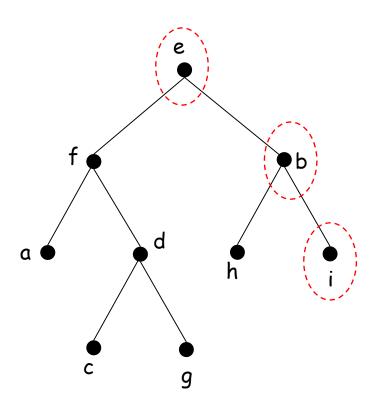
a - c

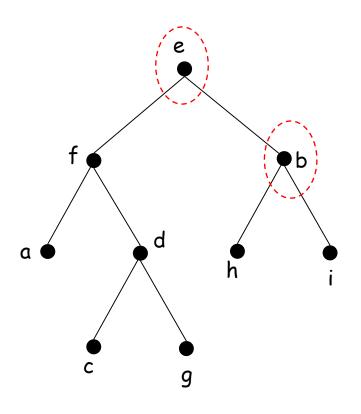


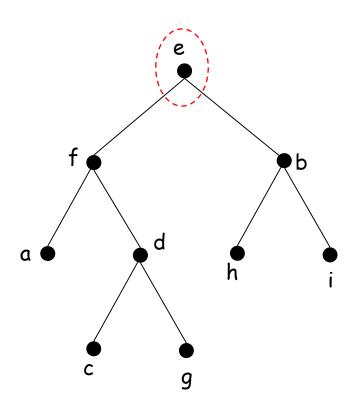


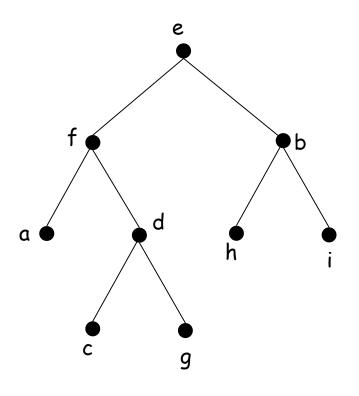




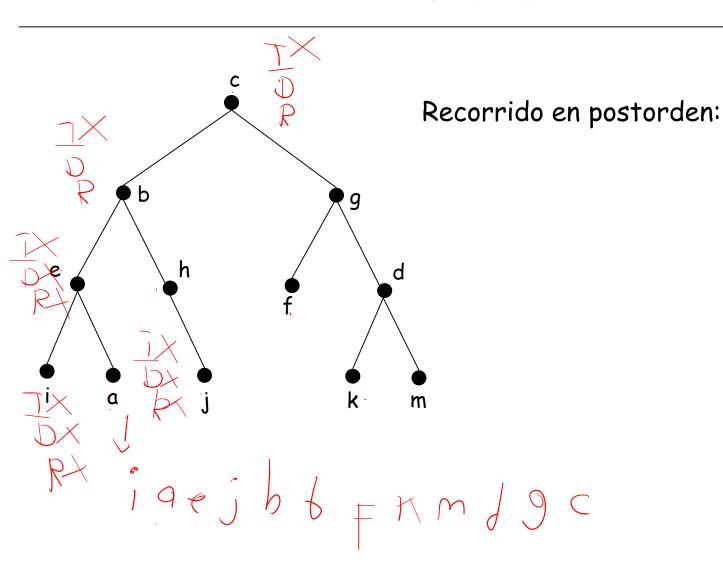


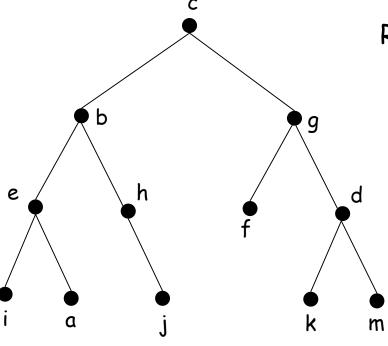




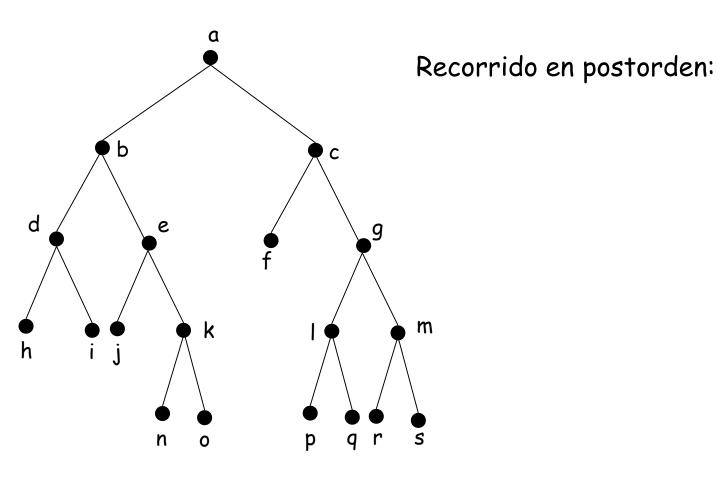


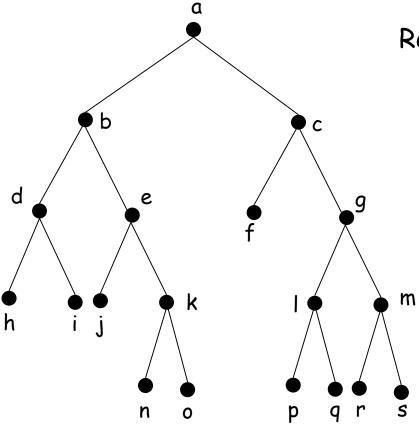
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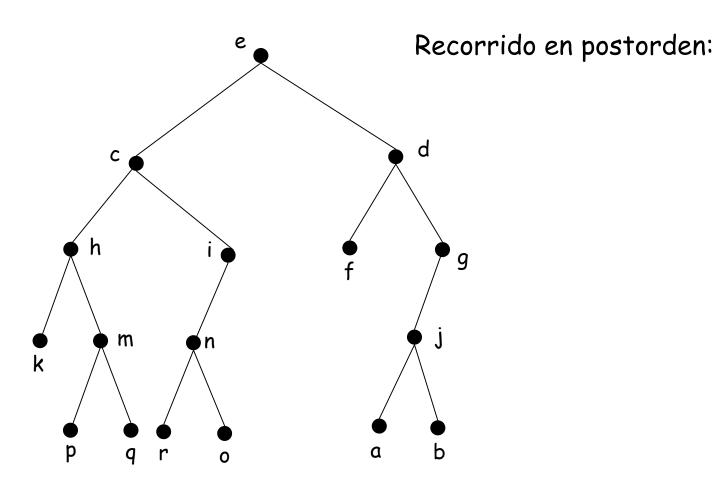
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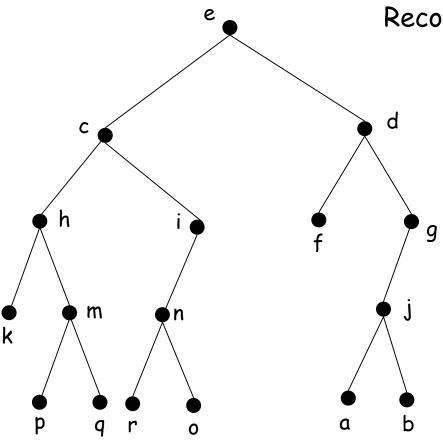




Recorrido en postorden:

h-i-d-j-n-o-k-e-b-f-p-q-l-r-s-m-g-c-a

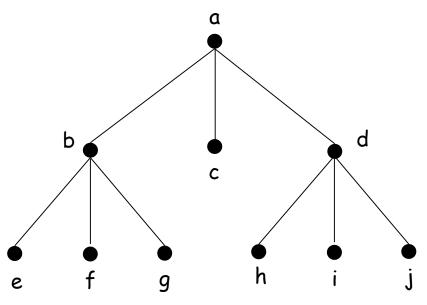




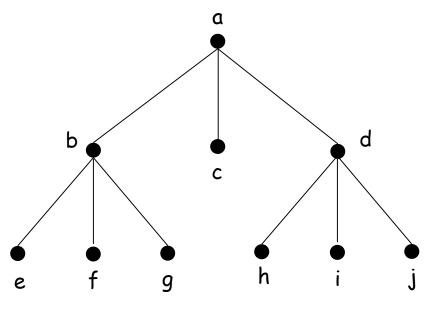
Recorrido en postorden:

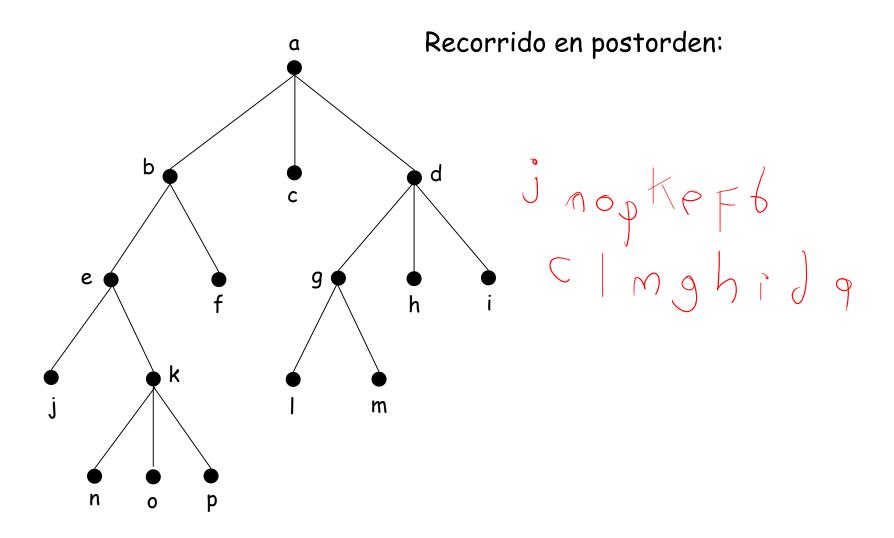
k-p-q-m-h-r-o-n-i-c-f-a-b-j-g-d-e

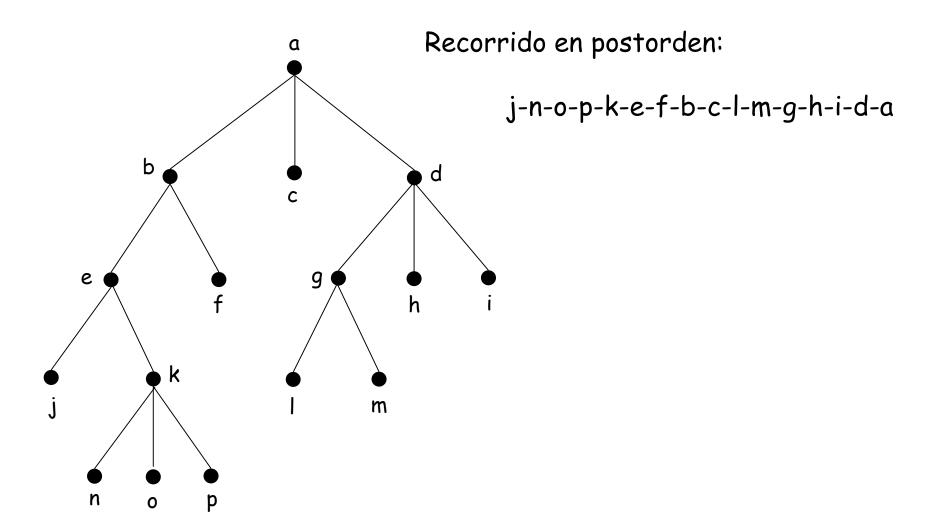
Recorrido en postorden:



Recorrido en postorden:





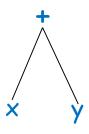


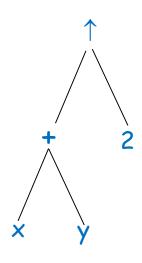
Notación infija, prefija y postfija

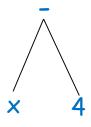
Permite representar expresiones complejas como proposiciones compuestas, combinaciones de conjuntos y expresiones aritméticas

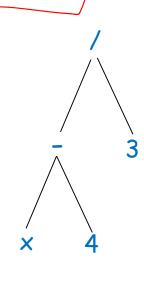
Expresiones aritméticas

- Suma (+)
- Resta (-)
- Multiplicación (*)
- División (/)
- Potencia (↑)



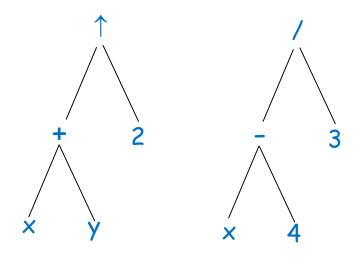


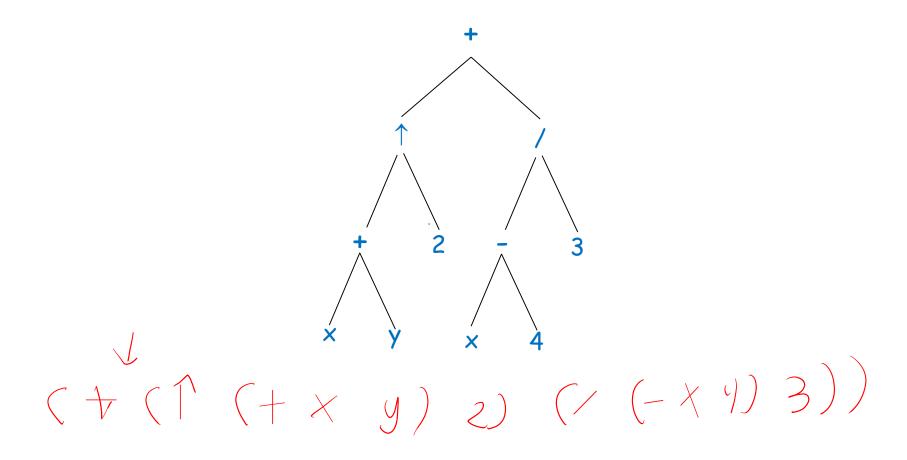




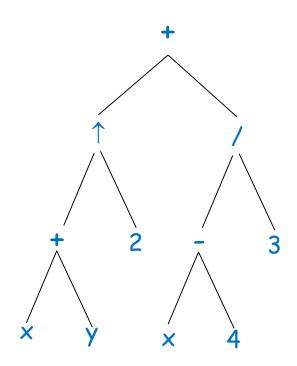
$$(/(= \times 4) 3$$

 $(\times - 4)/3$



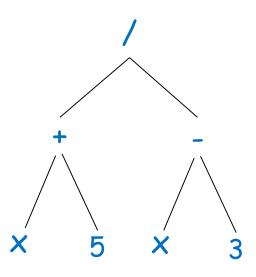


Muestre un árbol binario que represente la expresión matemática $((x+y)^2)+((x-4)/3)$



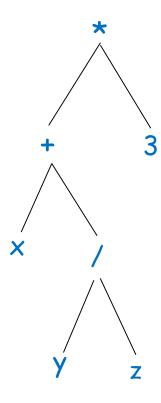
La expresión $((x+y)^2)+((x-4)/3)$ se obtiene al hacer el recorrido en inorden

$$(x + s) / (x - 3)$$
 $((x + x s) (-x 3))$
 $((x + x s) (x 3 - x 3))$



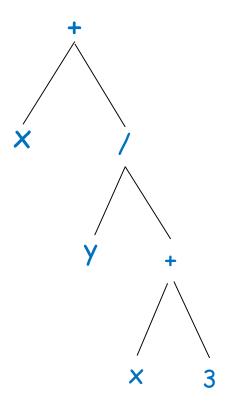
Árbol binario que representa la expresión matemática (x+5)/(x-3)

Muestre un árbol binario que represente la expresión matemática (x+(y/z))*3

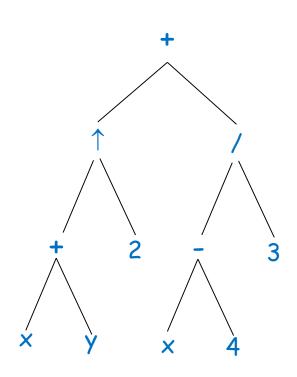


Árbol binario que representa la expresión matemática (x+(y/z))*3

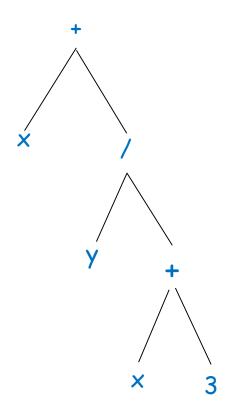
Muestre un árbol binario que represente la expresión matemática x+(y/(x+3))



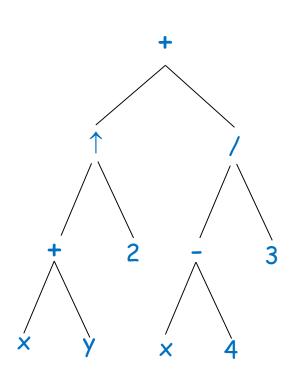
Árbol binario que representa la expresión matemática x+(y/(x+3))



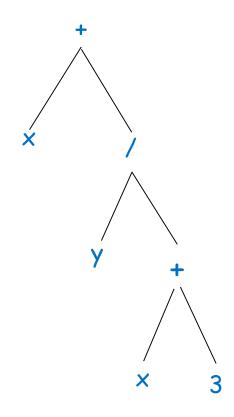
Árbol binario que representa la expresión matemática $((x+y)^{2})+((x-4)/3)$



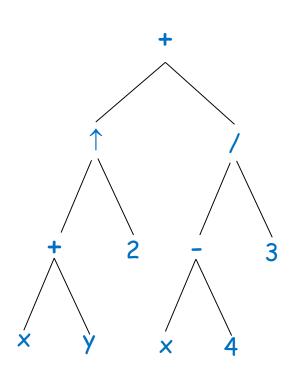
Árbol binario que representa la expresión matemática x+(y/(x+3))



Recorrido en inorden:

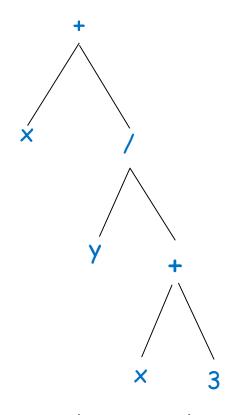


Recorrido en inorden:



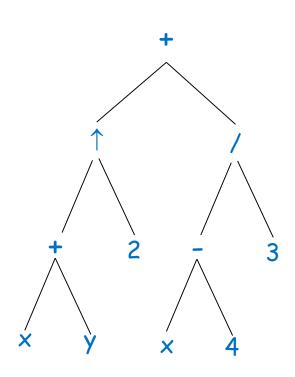
Recorrido en inorden:

$$((x+y)^2)+((x-4)/3)$$



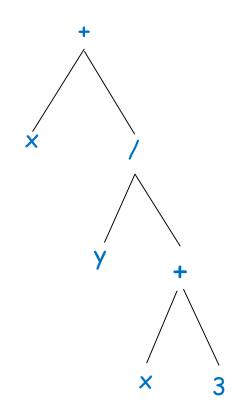
Recorrido en inorden:

$$x+(y/(x+3))$$



Recorrido en inorden:

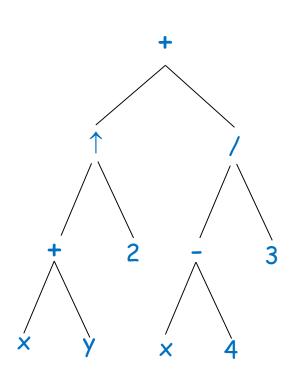
$$((x+y)^2)+((x-4)/3)$$



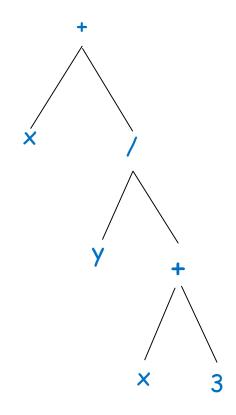
Recorrido en inorden:

$$x+(y/(x+3))$$

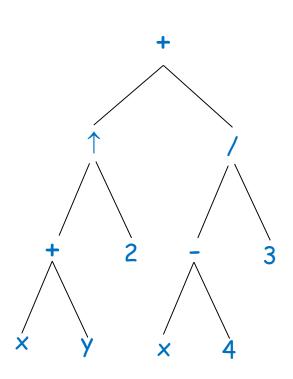
Notación infija



Recorrido en preorden:

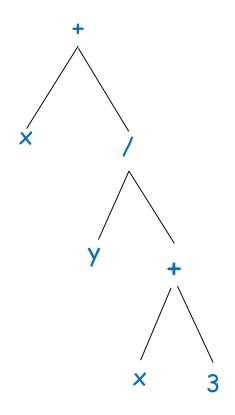


Recorrido en preorden:



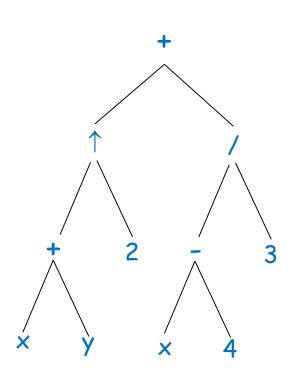
Recorrido en preorden:

$$+ \uparrow + x y 2 / - x 4 3$$



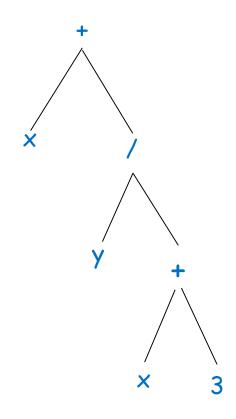
Recorrido en preorden:

$$+x/y+x$$
 3



Recorrido en preorden:

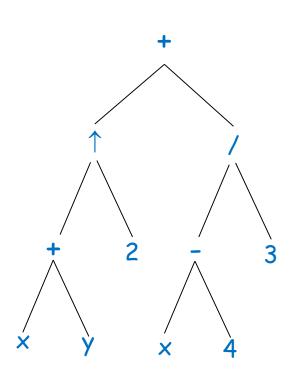
$$+ \uparrow + x y 2 / - x 4 3$$



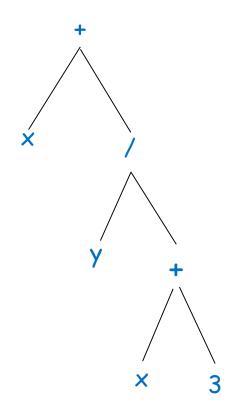
Recorrido en preorden:

$$+x/y+x$$
 3

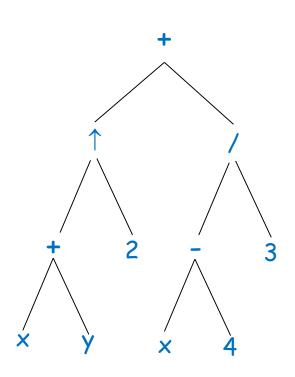
Notación prefija



Recorrido en postorden:

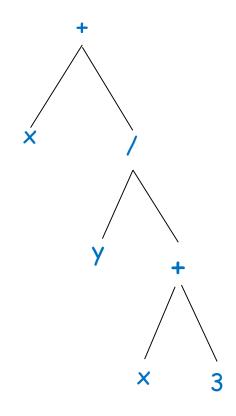


Recorrido en postorden:

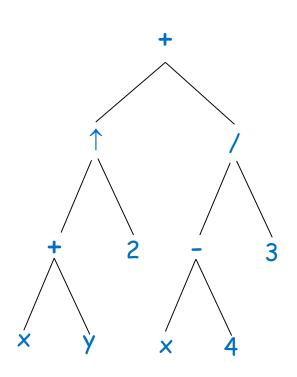


Recorrido en postorden:

$$xy + 2 \uparrow x4 - 3 / +$$

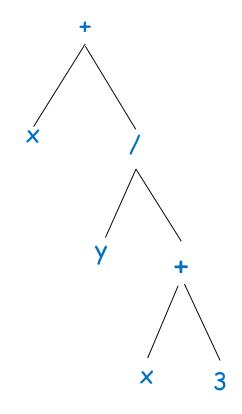


Recorrido en postorden:



Recorrido en postorden:

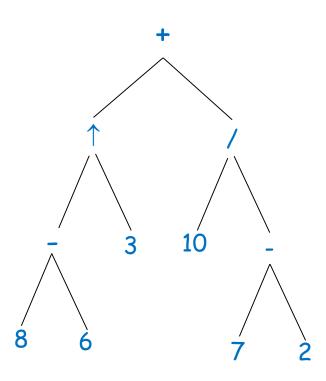
$$xy + 2 \uparrow x4 - 3 / +$$

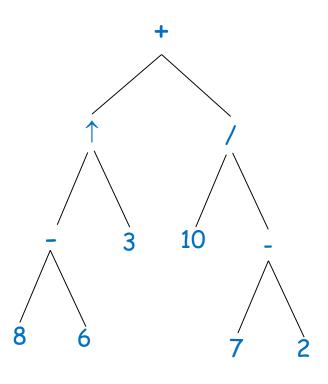


Recorrido en postorden:

Notación postfija

Representar en notación infija, prefija y postfija, la expresión matemática dada por el siguiente árbol:

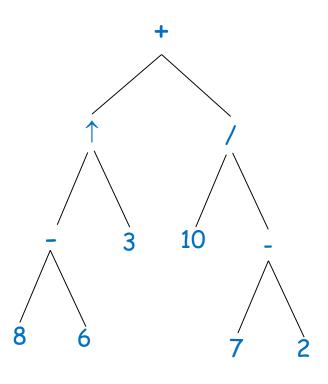




Notación **infija**: $((8-6)^{\uparrow}3)+(10/(7-2))$

Notación **prefija**: + 1 - 8 6 3 / 10 - 7 2

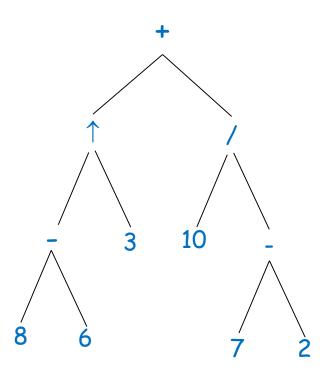
Notación **postfija**: 8 6 - 3 ↑ 10 7 2 - / +



Notación **infija**: $((8-6)^{\uparrow}3)+(10/(7-2)) = ?$

Notación **prefija**: + 1 - 8 6 3 / 10 - 7 2

Notación **postfija**: 8 6 − 3 ↑ 10 7 2 − / +



Notación **infija**: $((8-6)^{\uparrow}3)+(10/(7-2))=10$

Notación **prefija**: + 1 - 8 6 3 / 10 - 7 2

Notación **postfija**: 8 6 − 3 ↑ 10 7 2 − / +

Notación **prefija**: + 1 - 8 6 3 / 10 - 7 2

Notación **prefija**: + 1 - 8 6 3 / 10 - 7 2

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$

 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$

 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$

 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$
 $+ \uparrow 232$

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$
 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$
 $+ \uparrow 232$

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$
 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$
 $+ \uparrow 232$
 $+ 82$

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$
 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$
 $+ \uparrow 232$
 $+ 82$

Notación **prefija**:
$$+ \uparrow - 863 / 10 - 72$$

 $+ \uparrow - 863 / 105$
 $+ \uparrow - 8632$
 $+ \uparrow 232$
 $+ 82$

Notación **postfija**: 86 - 3 ↑ 1072 - / +

Notación **postfija**: 86 - 3 ↑ 1072 - / +

Notación **postfija**: 86 - 3 ↑ 1072 - / +

23 1072-/+

Notación **postfija**:
$$86 - 3 \uparrow 1072 - / +$$

$$23 \uparrow 1072 - / +$$

$$81072 - / +$$

Notación **postfija**:
$$86 - 3 \uparrow 1072 - / +$$
 $23 \uparrow 1072 - / +$ $81072 - / +$

Notación **postfija**:
$$86 - 3 \uparrow 1072 - / +$$

$$23 \uparrow 1072 - / +$$

$$81072 - / +$$

$$8105/ +$$

$$82 +$$

Indique el valor de la siguiente expresión que está en notación prefija

Notación **prefija**: + - *235/↑234

Indique el valor de la siguiente expresión que está en notación prefija

```
Notación prefija: + - * 2 3 5 / ↑ 2 3 4

+ - * 2 3 5 / 8 4

+ - * 2 3 5 2

+ - 6 5 2

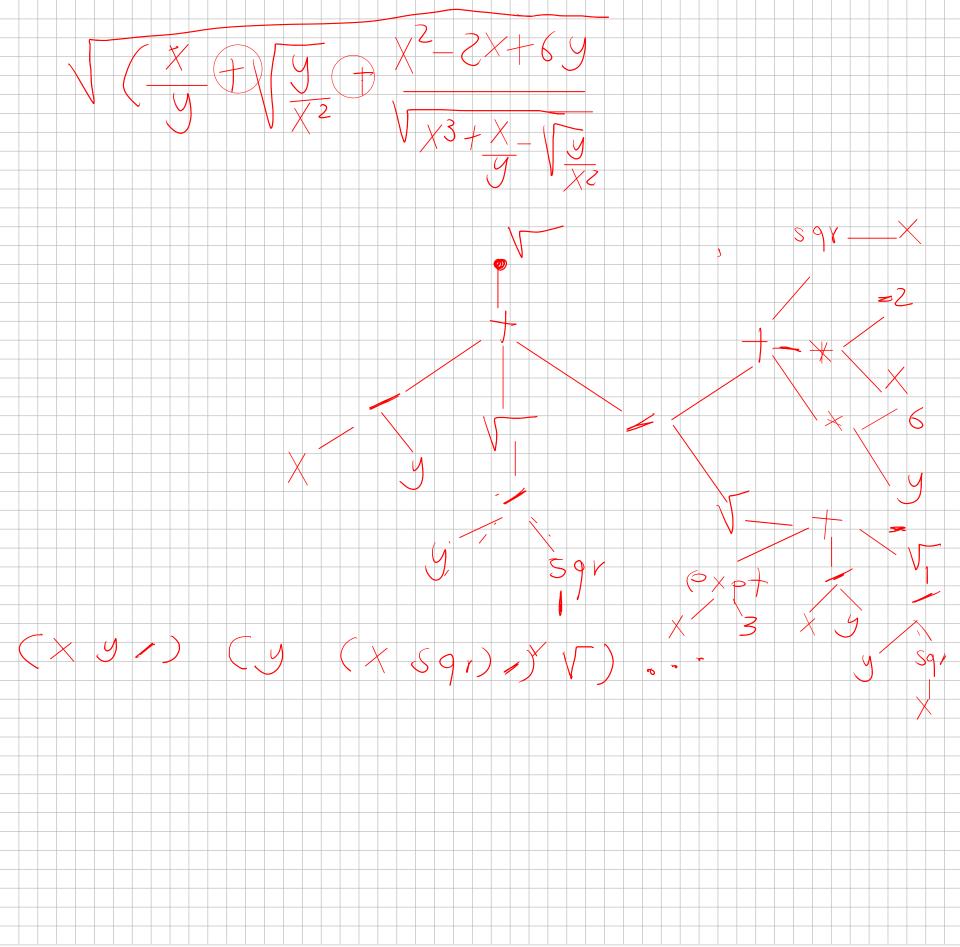
+ 1 2

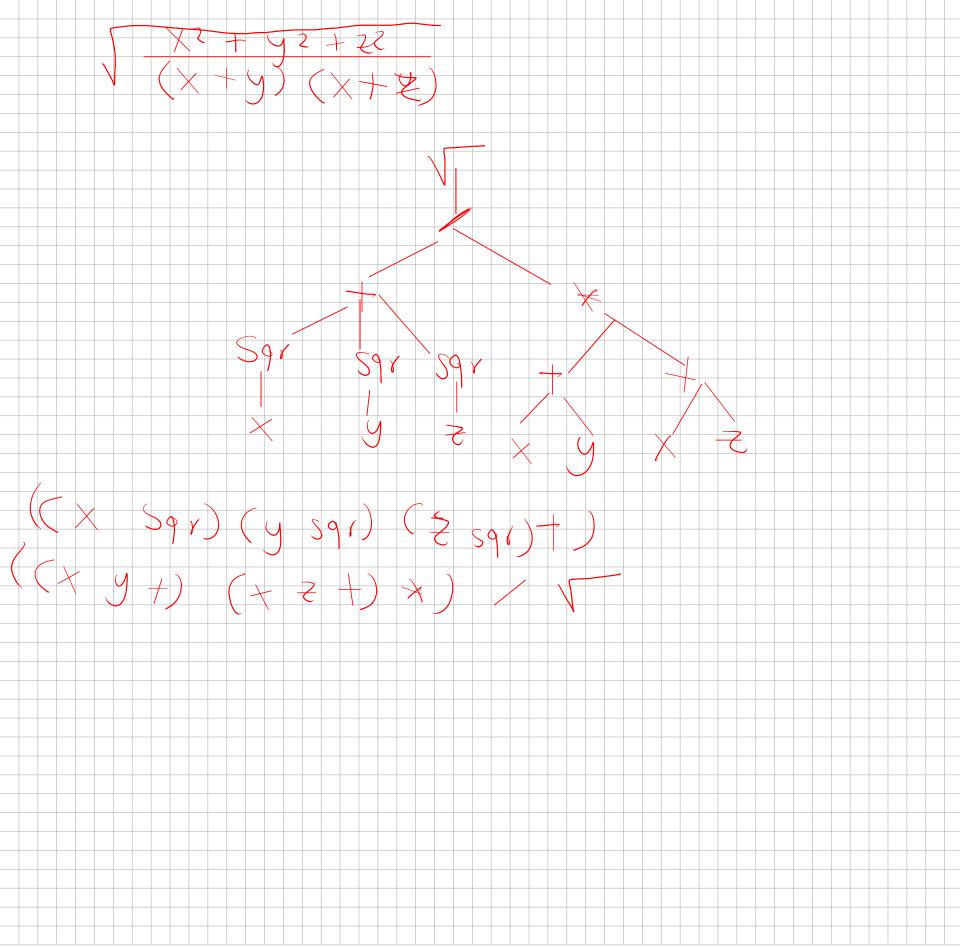
3
```

Indique el valor de la siguiente expresión que está en notación postfija

Notación **postfija**: 723 * - 4 ↑ 93 / +

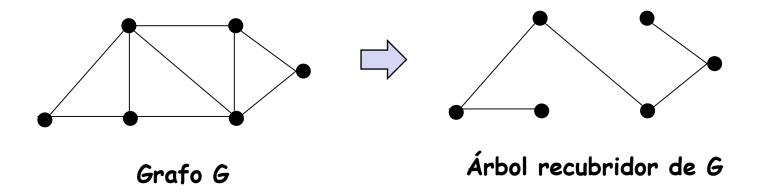
Indique el valor de la siguiente expresión que está en notación postfija





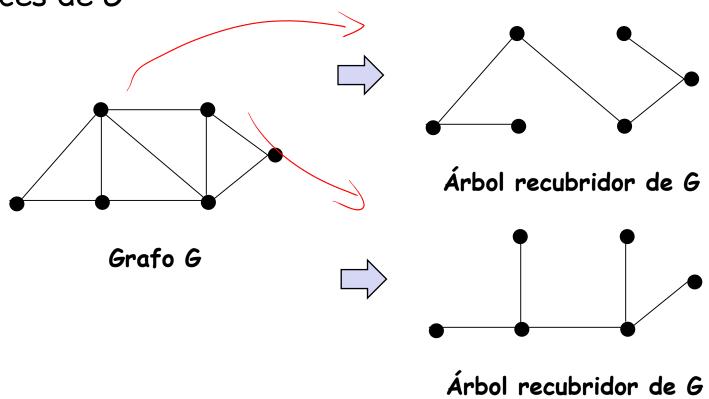
Árbol recubridor

Sea G un grafo simple, un árbol recubridor de G es un subgrafo de G que es un árbol y contiene todos los vértices de G

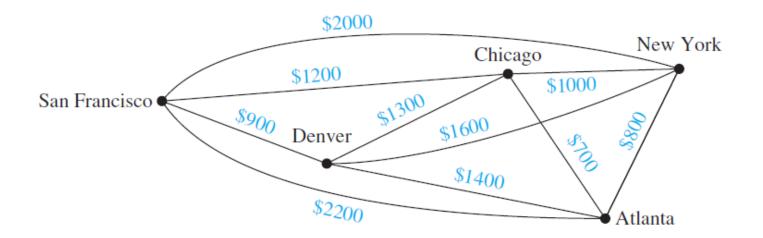


Árbol recubridor

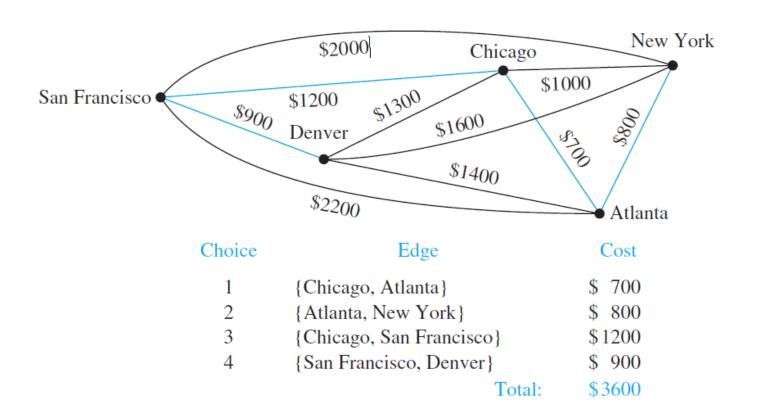
Sea G un grafo simple, un árbol recubridor de G es un subgrafo de G que es un árbol y contiene todos los vértices de G



El siguiente grafo indica los costos de una red de comunicaciones. ¿Qué enlaces se deben mantener para asegurar que hay una forma de comunicar cada dos ciudades a un costo mínimo?c



El siguiente grafo indica los costos de una red de comunicaciones. ¿Qué enlaces se deben mantener para asegurar que hay una forma de comunicar cada dos ciudades a un costo mínimo?



Algoritmo Prim

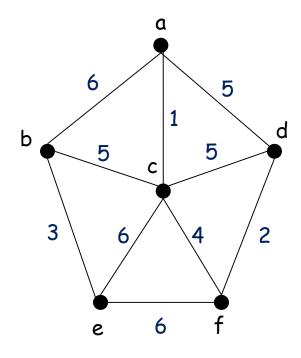
- Escoga la arista con menor peso y adiciónela al árbol recubridor
- Seleccione la arista con menor peso que sea incidente con el árbol recubridor y que no cree un circuito.
 Adiciónela al árbol.
- Repita el proceso hasta cuando el árbol tenga n-1 aristas (n es el número de vértices)

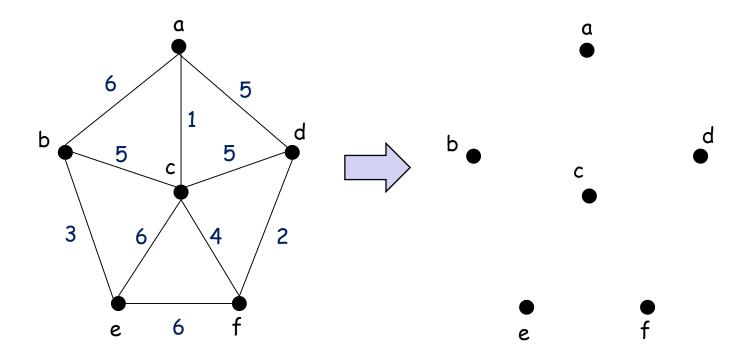
Algoritmo Prim

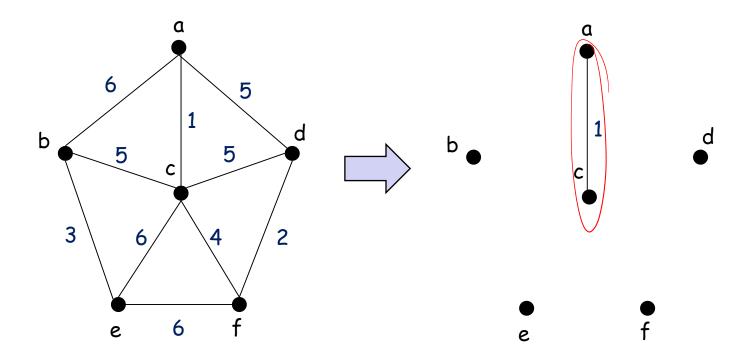
- Escoga la arista con menor peso y adiciónela al árbol recubridor
- Seleccione la arista con menor peso que sea incidente con el árbol recubridor y que no cree un circuito.
 Adiciónela al árbol.
- Repita el proceso hasta cuando el árbol tenga n-1 aristas (n es el número de vértices)

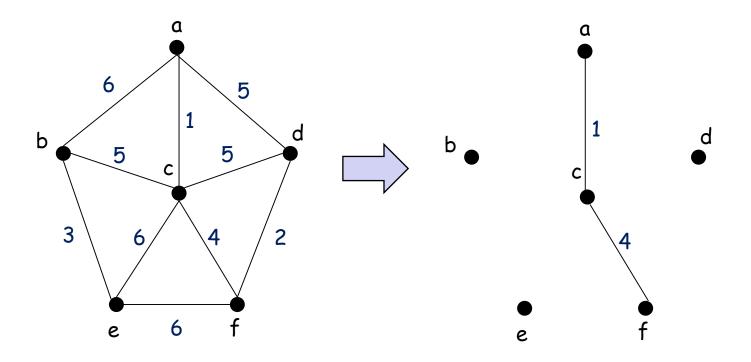
Al usar el algoritmo de Prim se pueden obtener árboles recubridores diferentes para un mismo grafo

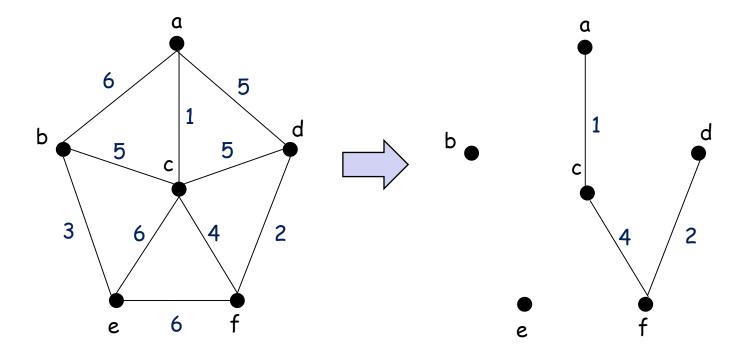
Encontrar un árbol recubridor mínimo usando el algoritmo de Prim

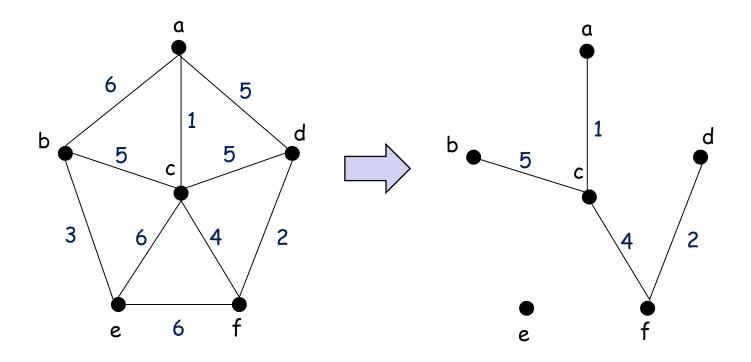


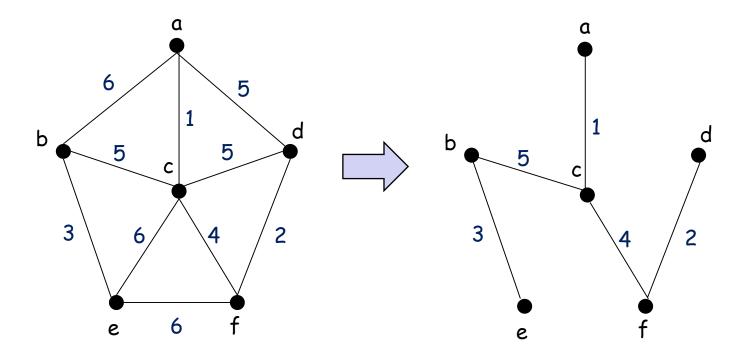




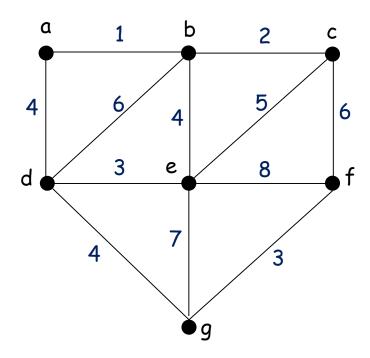


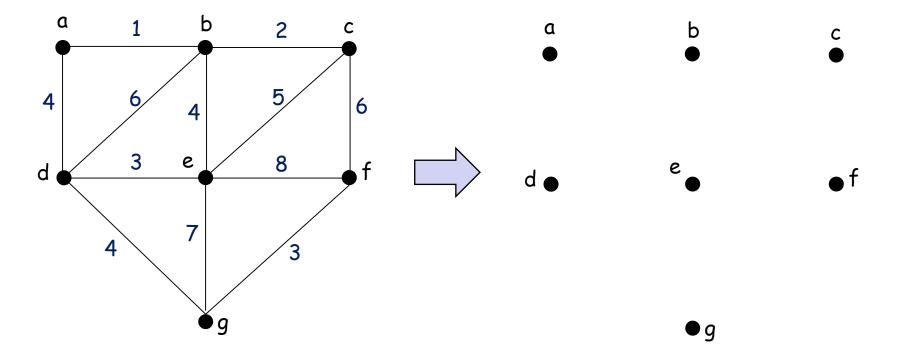


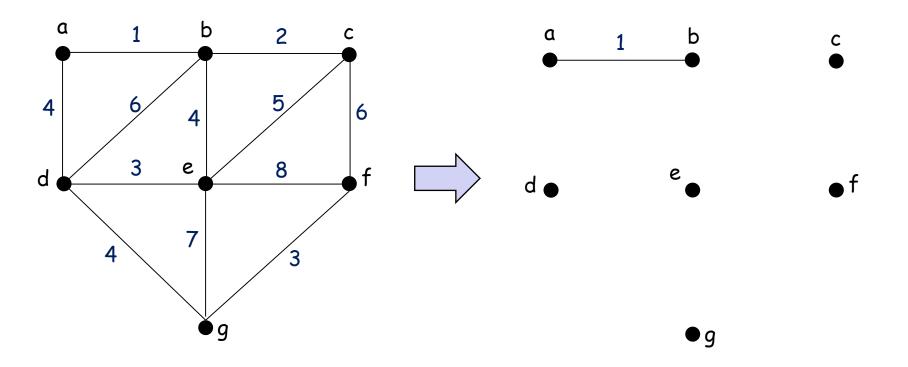


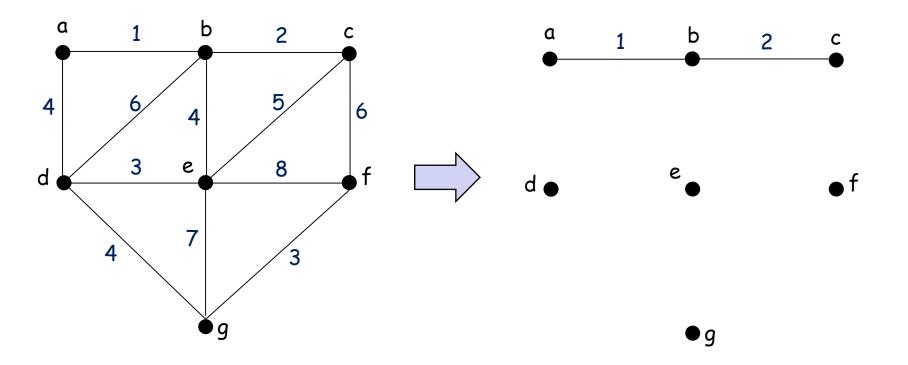


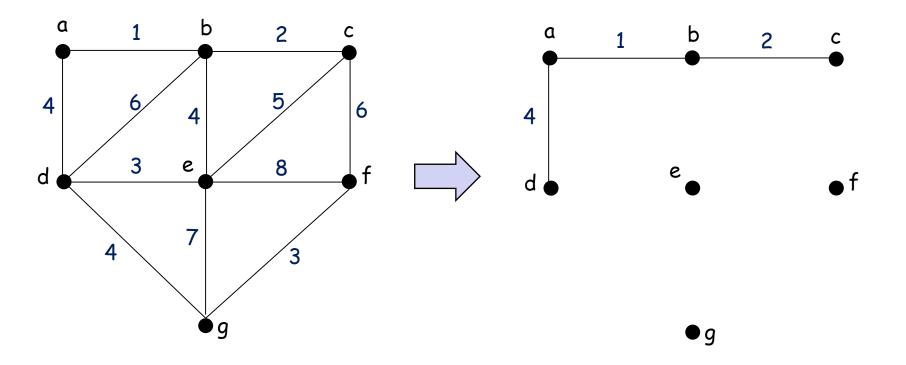
Encontrar un árbol recubridor mínimo usando el algoritmo de Prim

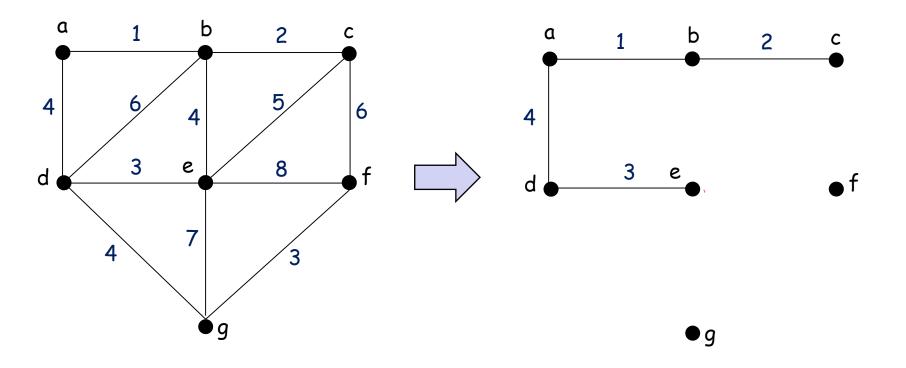


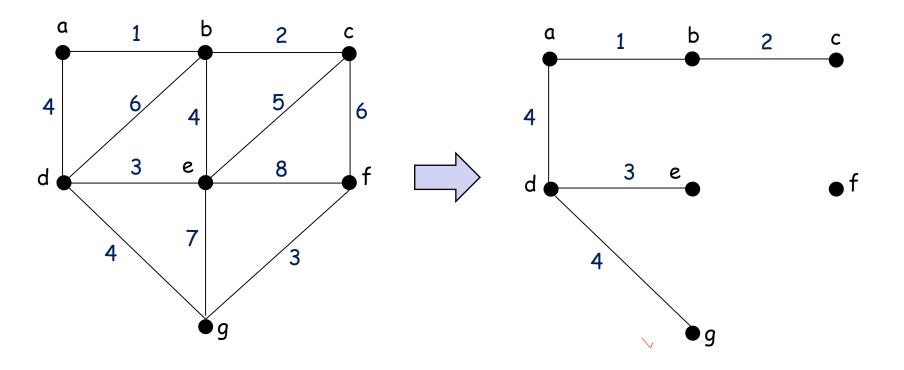


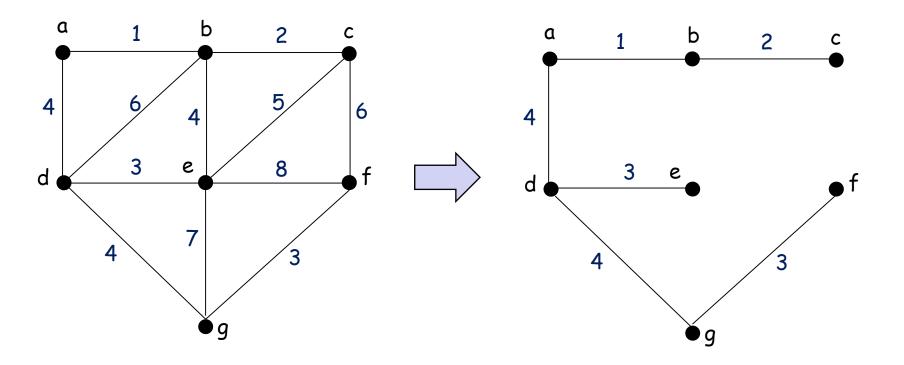




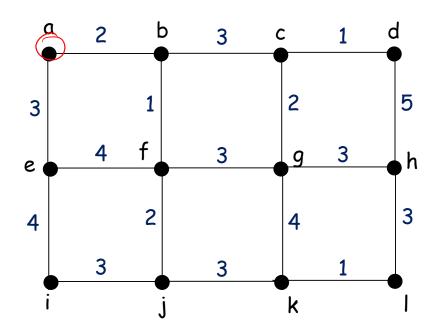


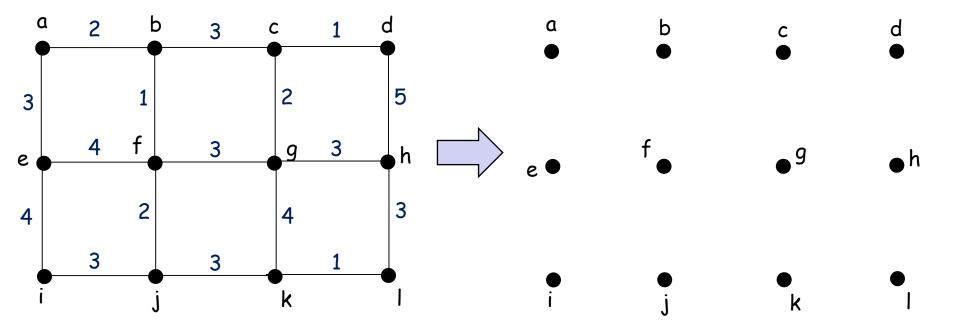


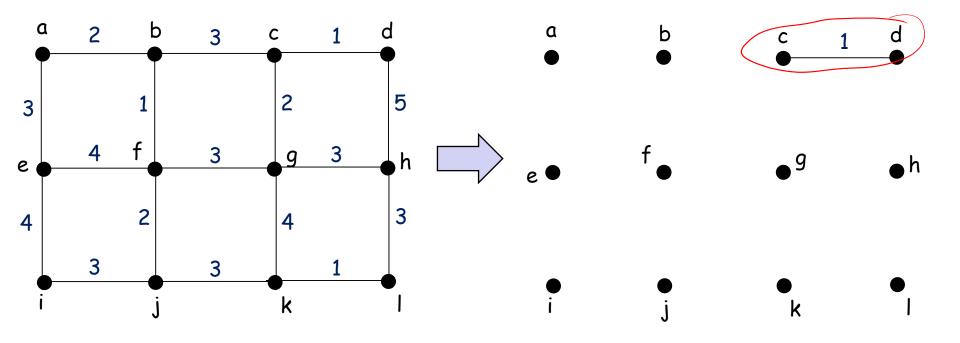


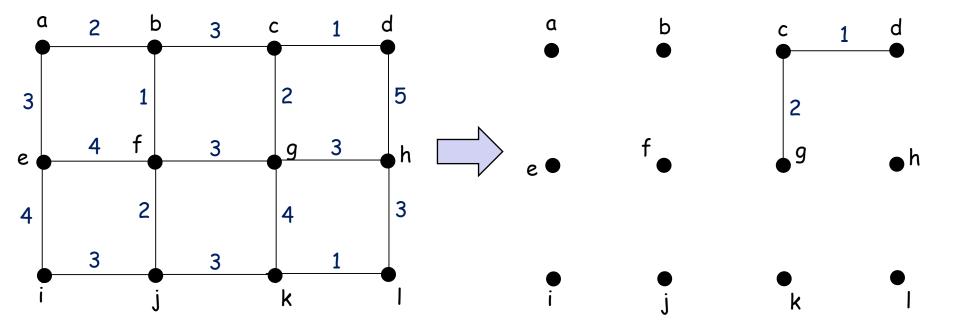


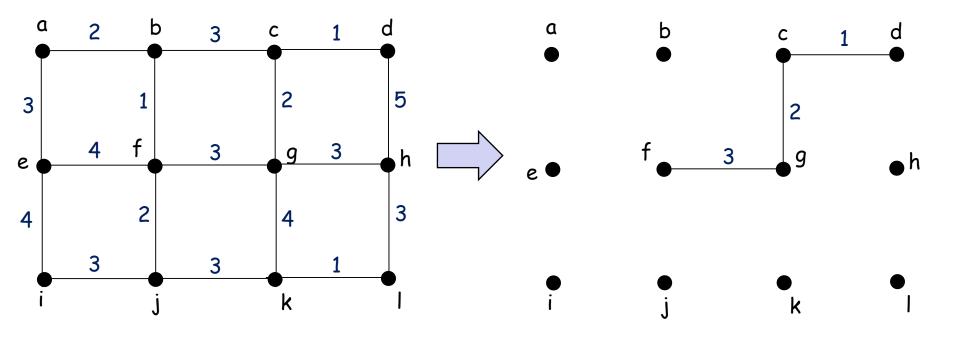
Encontrar un árbol recubridor mínimo usando el algoritmo de Prim

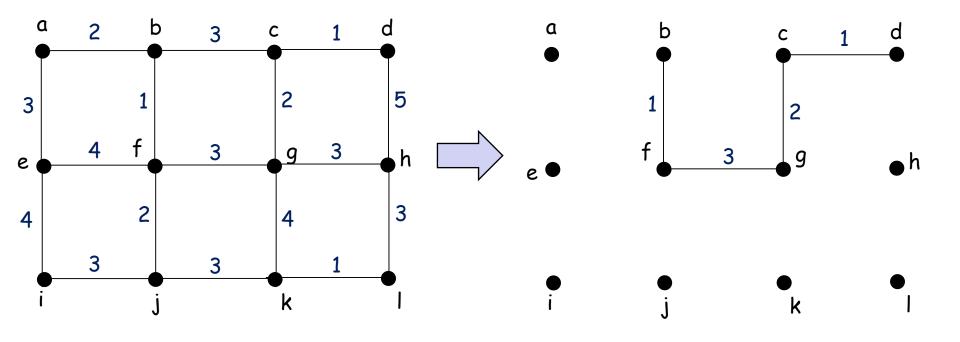


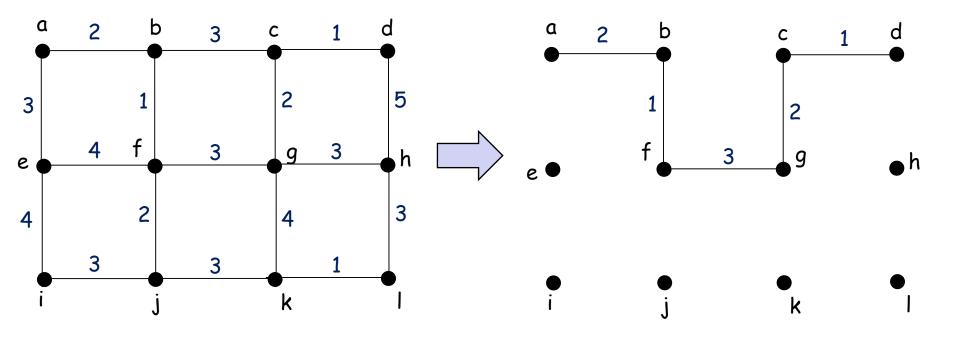


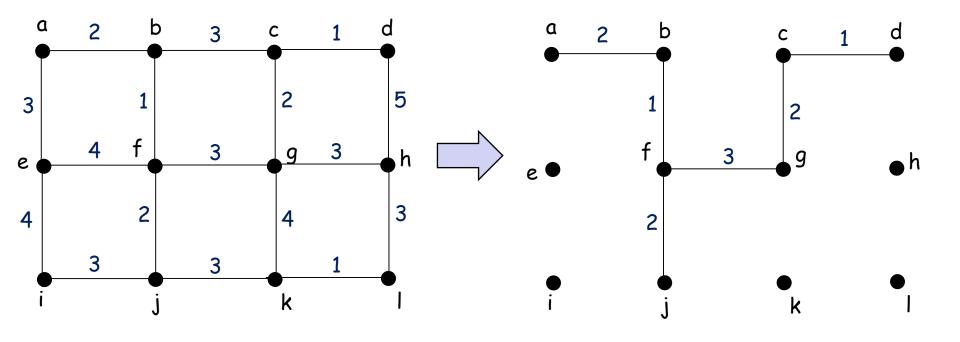


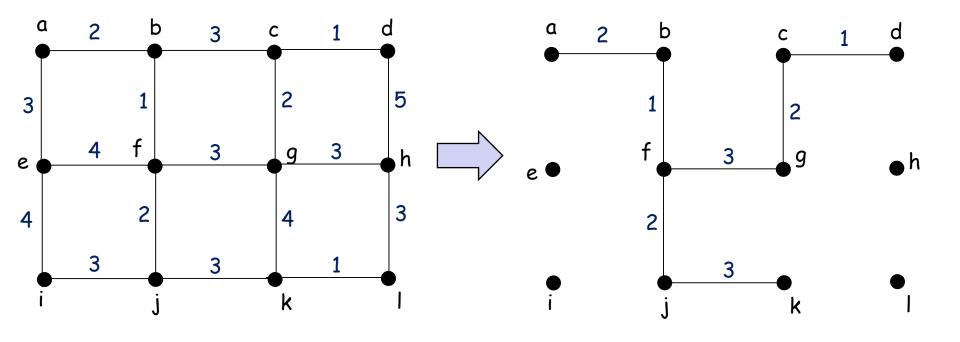


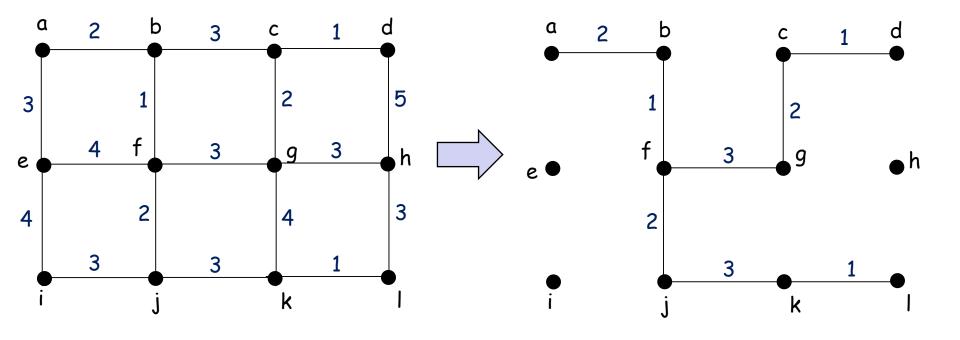


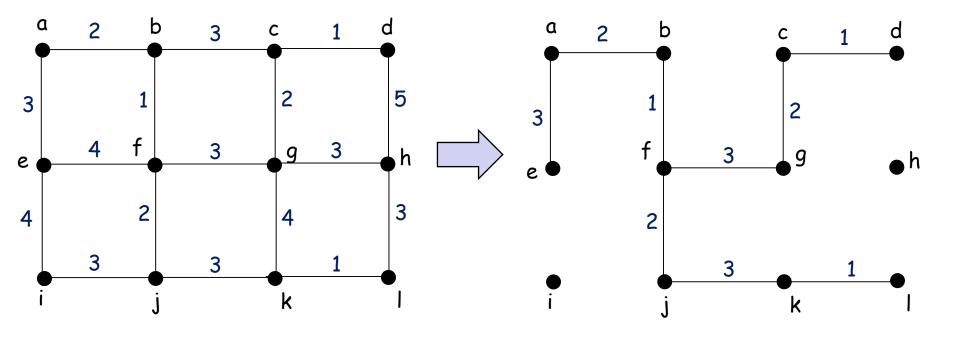


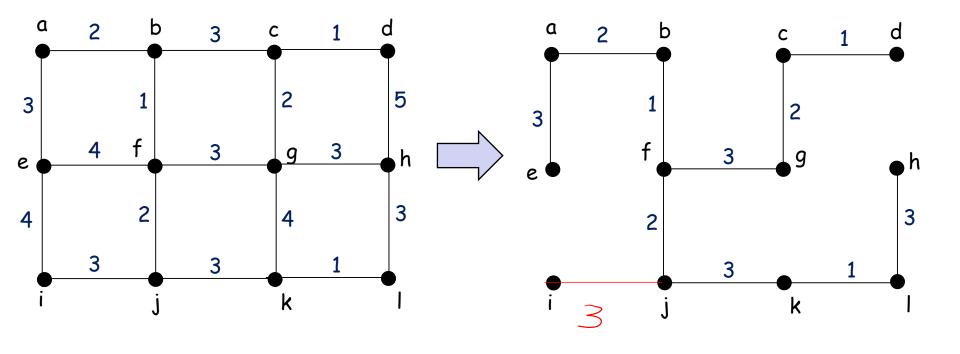




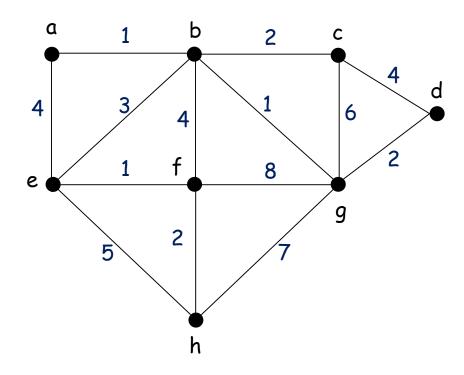


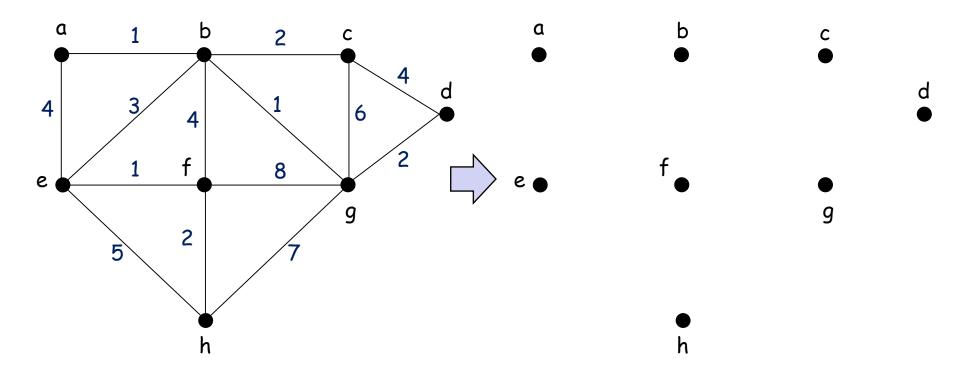


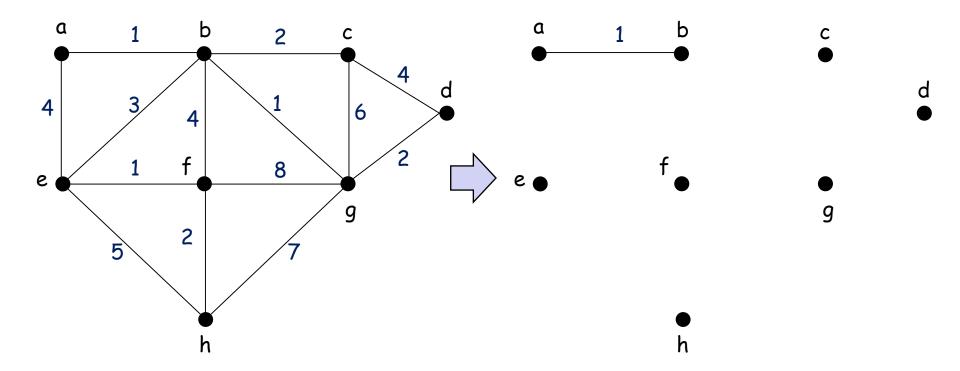


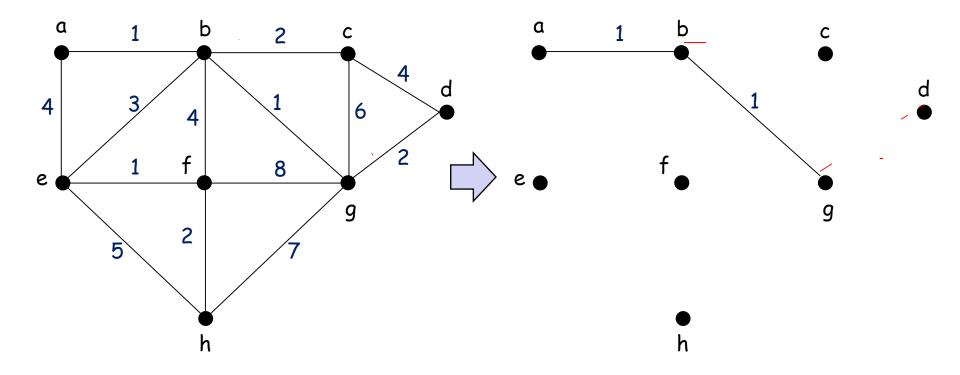


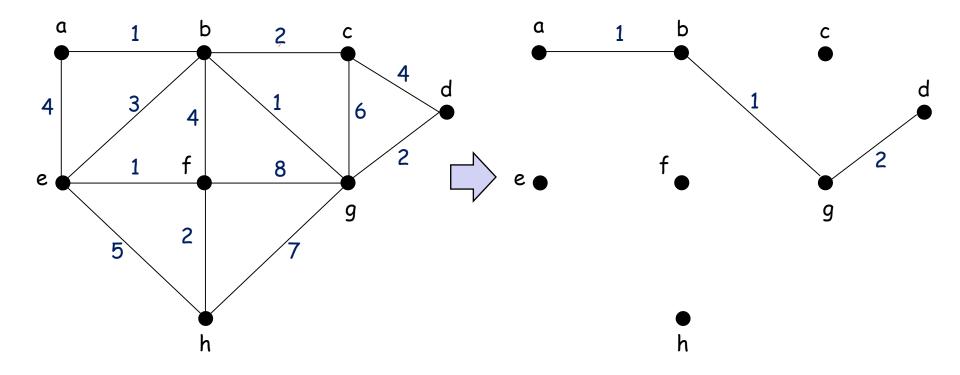
Encontrar un árbol recubridor mínimo usando el algoritmo de Prim

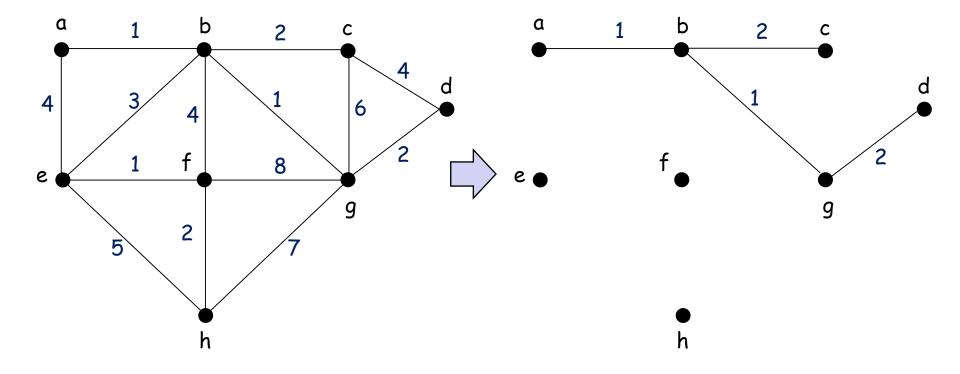


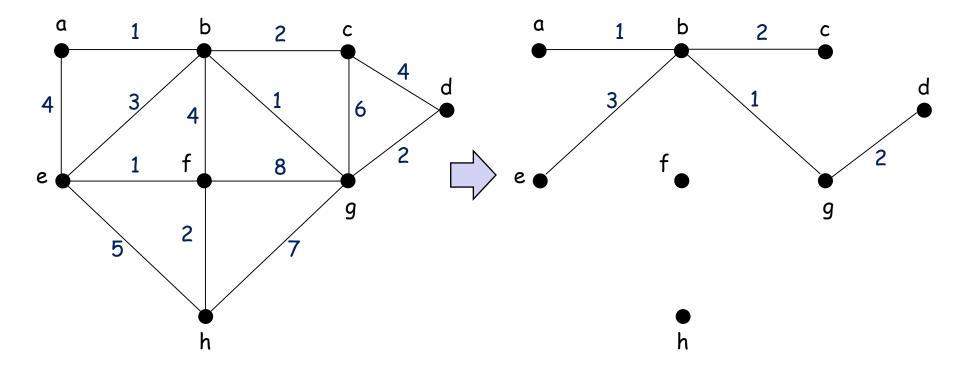


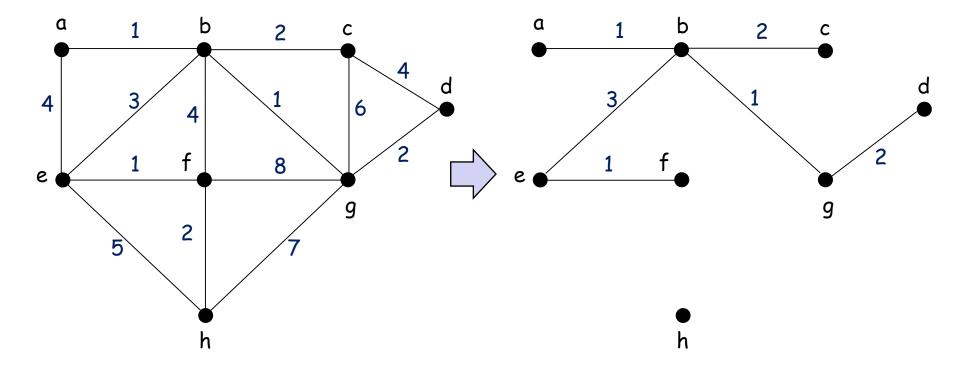


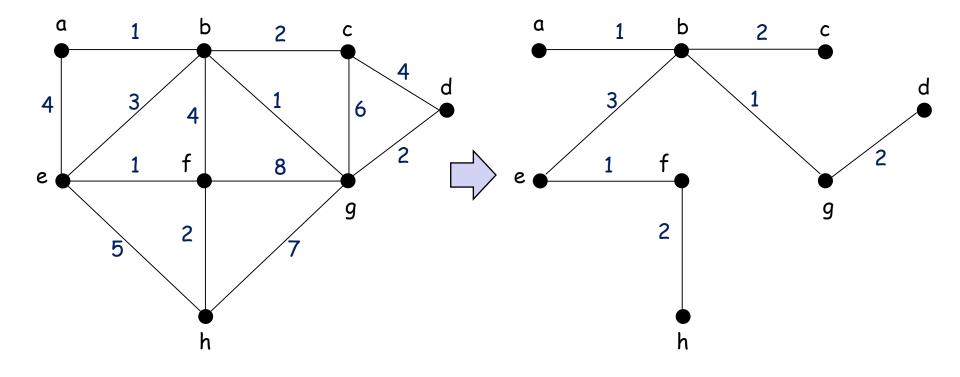












Algoritmo de Kruskal

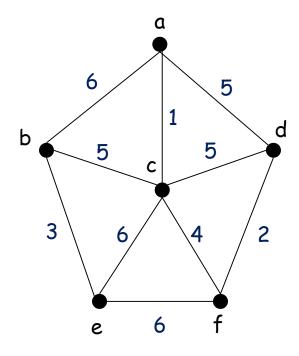
- Seleccione la arista con menor peso y adiciónela al árbol recubridor
- Adicione al árbol la arista con menor peso que no cree un circuito
- Repita el proceso cuando el árbol tenga n-1 aristas (n es el número de vértices)

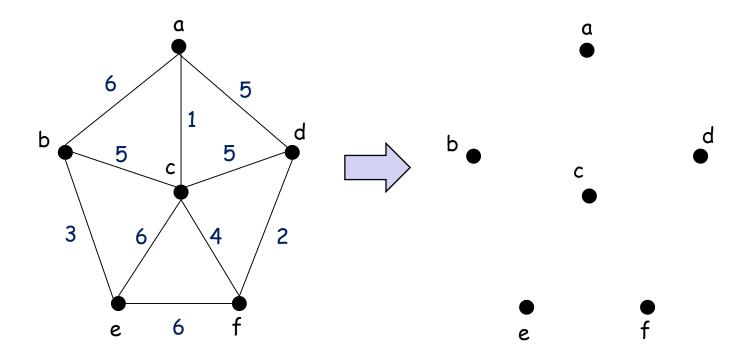
Algoritmo de Kruskal

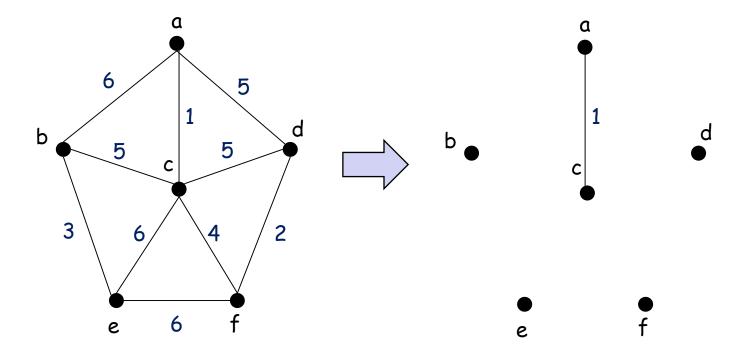
- Seleccione la arista con menor peso y adiciónela al árbol recubridor
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- Repita el proceso cuando el árbol tenga n-1 aristas (n es el número de vértices)

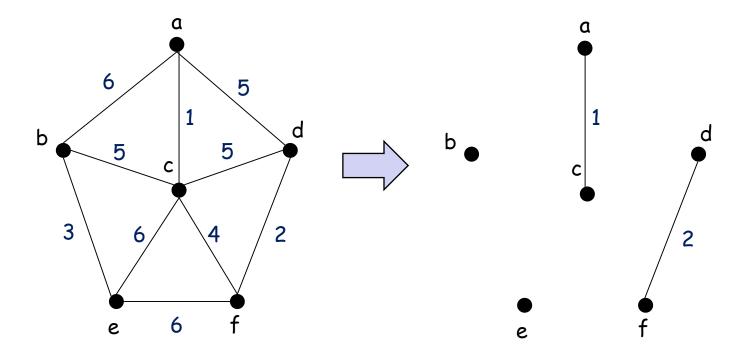
La arista que se selecciona con el algoritmo de Prim debe ser incidente en el árbol recubridor, mientras que en el algoritmo de Kruskal no

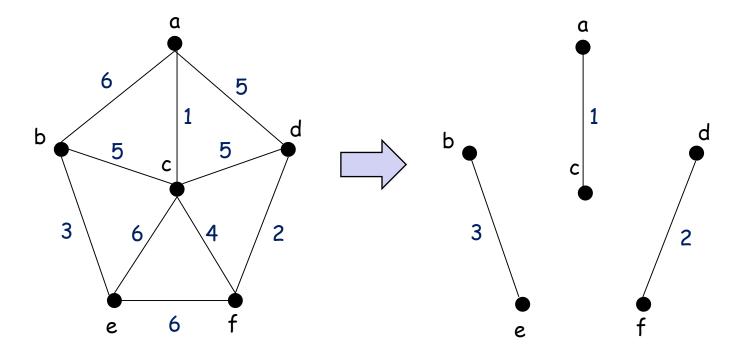
Encontrar un árbol recubridor mínimo usando el algoritmo de Kruskal

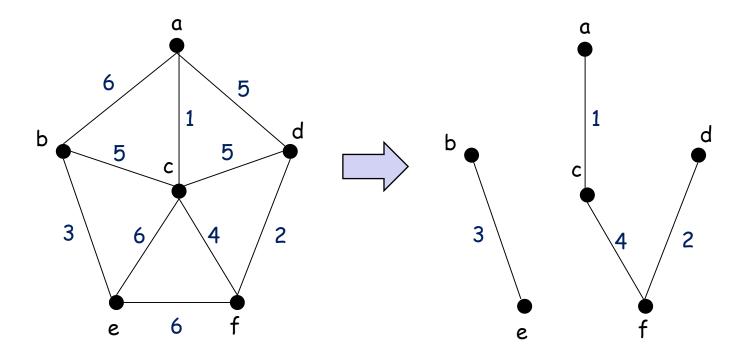


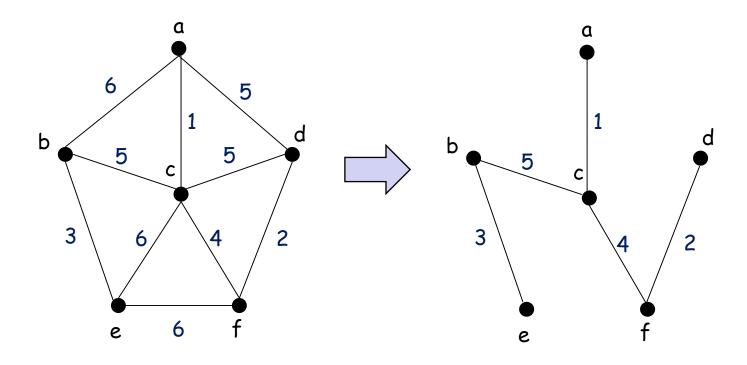


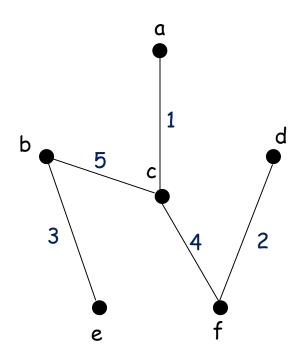




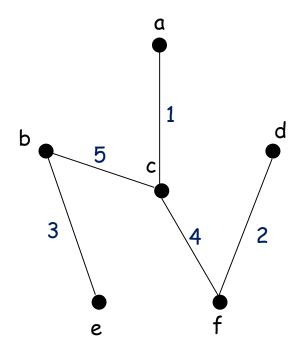






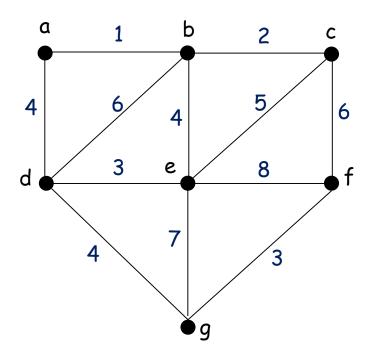


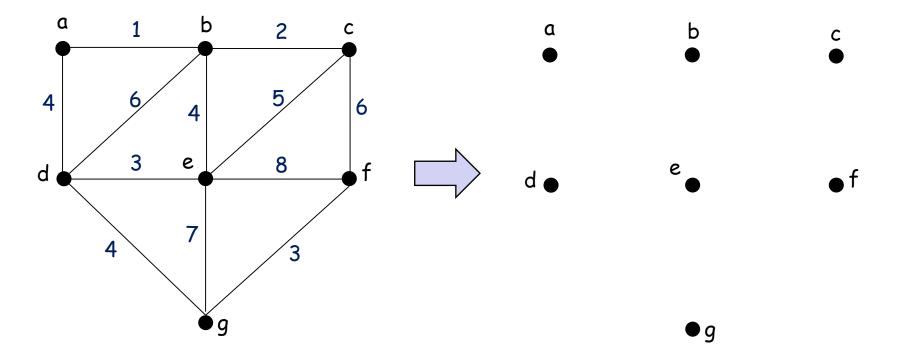
Árbol recubridor mínimo obtenido con el algoritmo de Prim

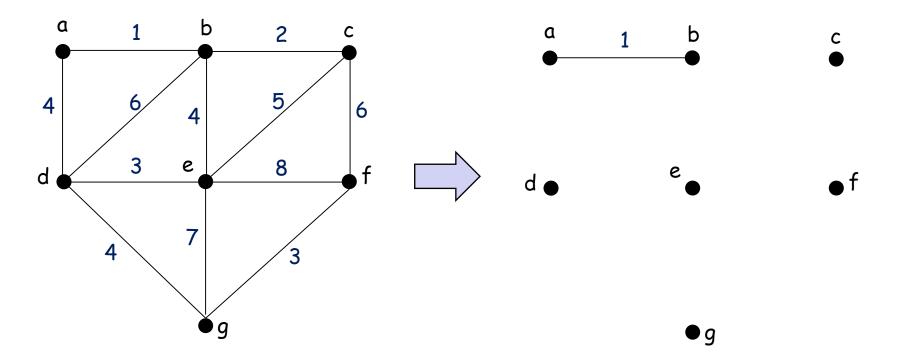


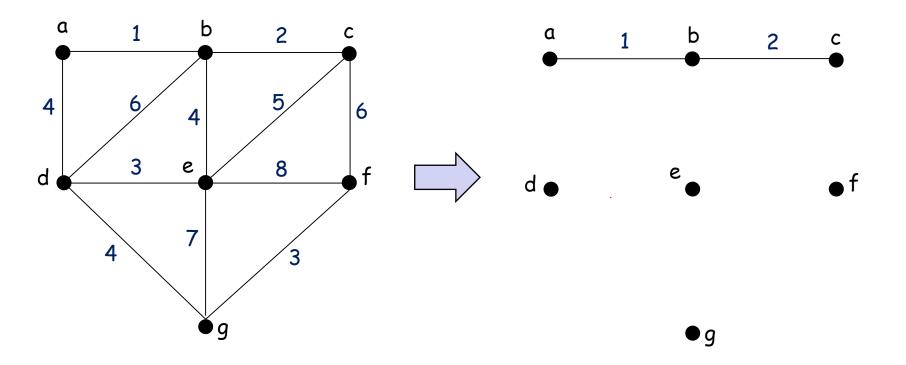
Árbol recubridor mínimo obtenido con el algoritmo de Kruskal

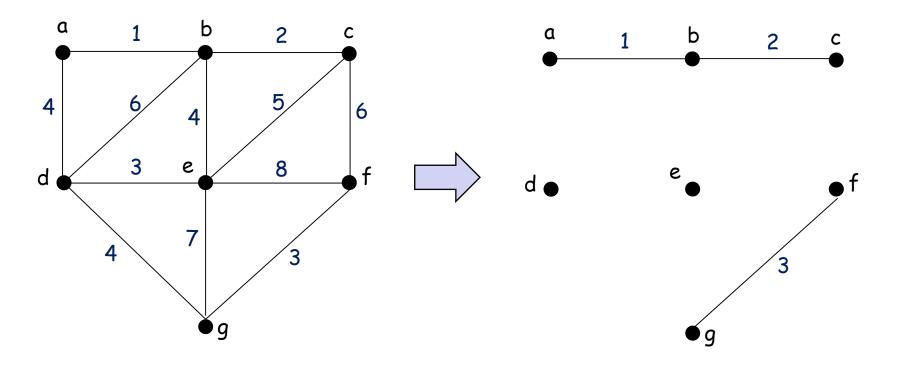
Encontrar un árbol recubridor mínimo usando el algoritmo de Kruskal

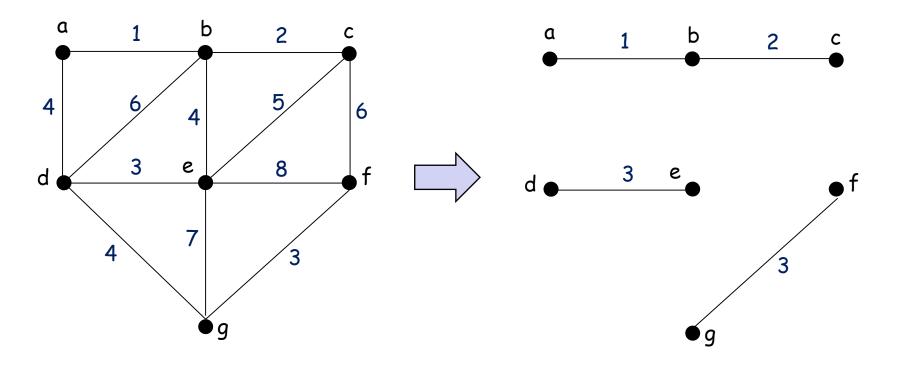


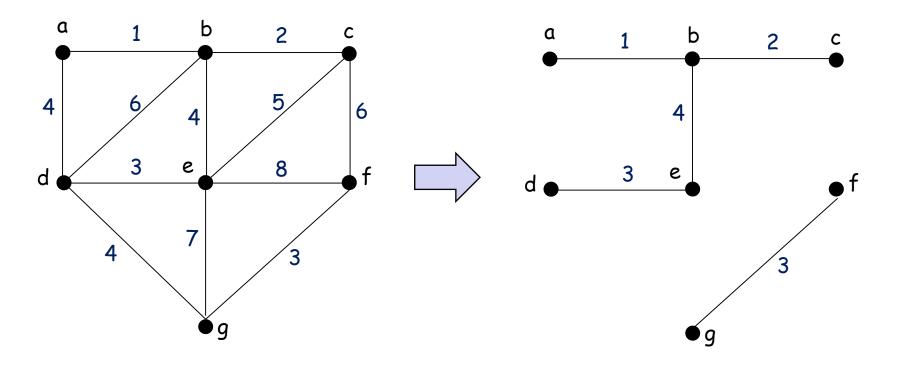


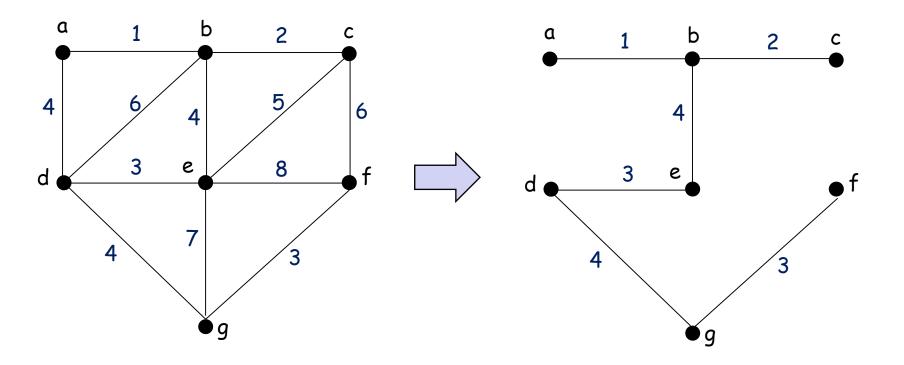


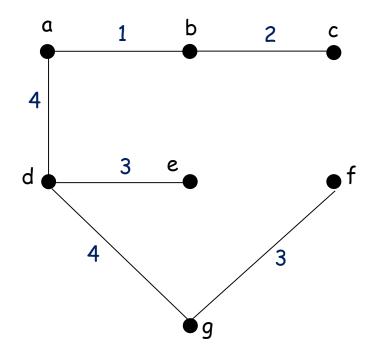




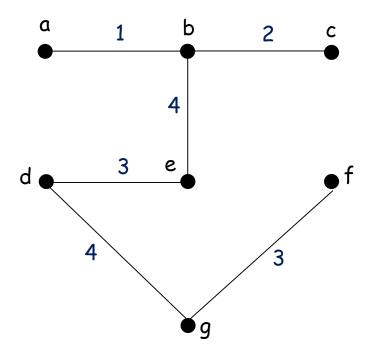






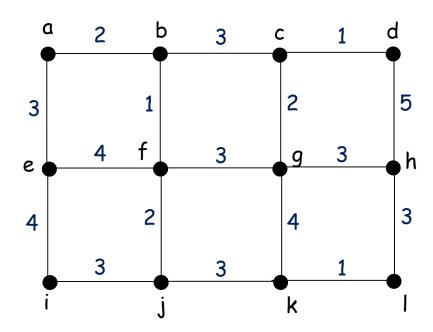


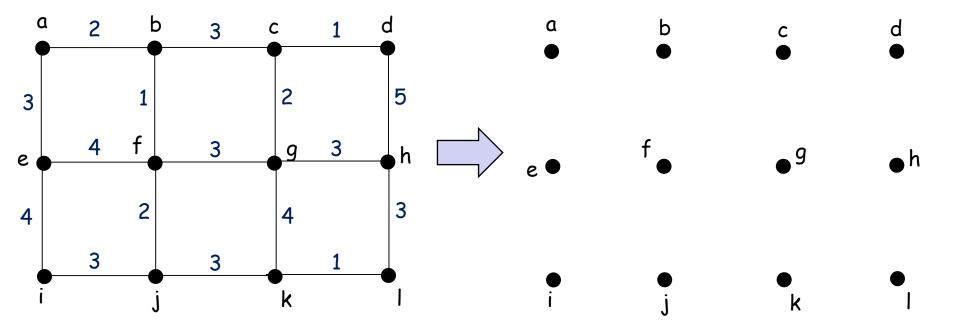
Árbol recubridor mínimo obtenido con el algoritmo de Prim

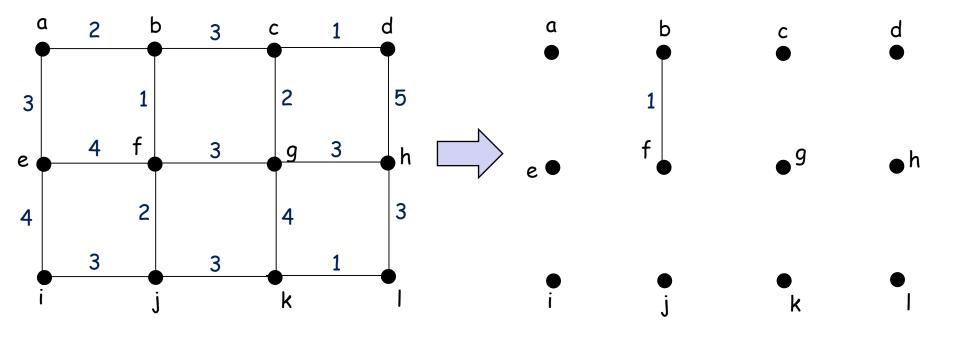


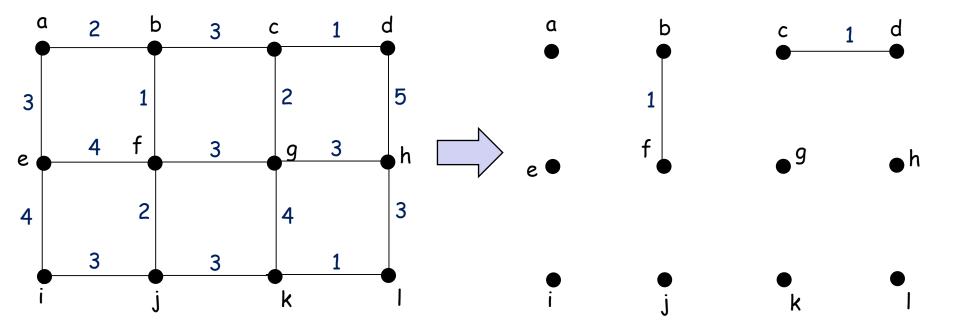
Árbol recubridor mínimo obtenido con el algoritmo de Kruskal

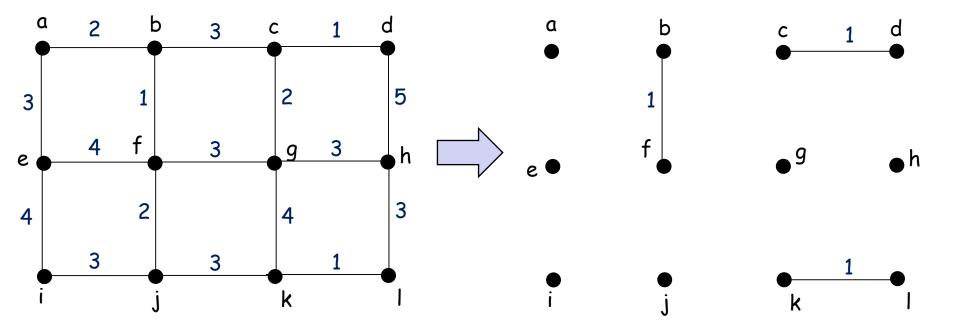
Encontrar un árbol recubridor mínimo usando el algoritmo de Kruskal

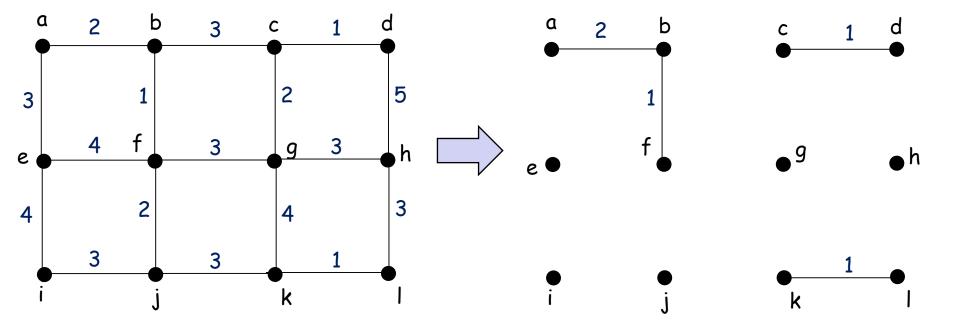


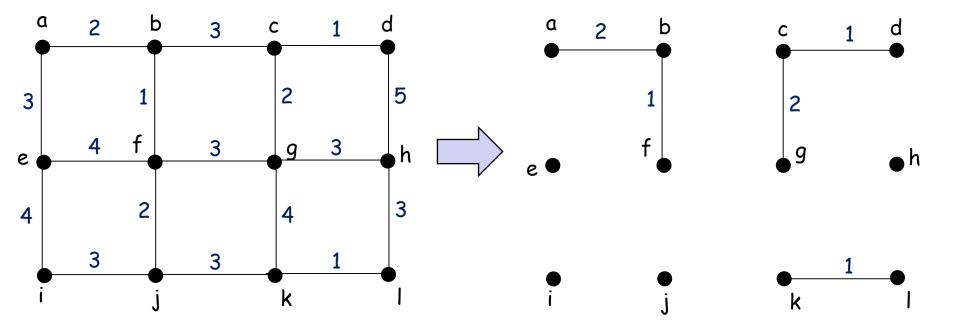


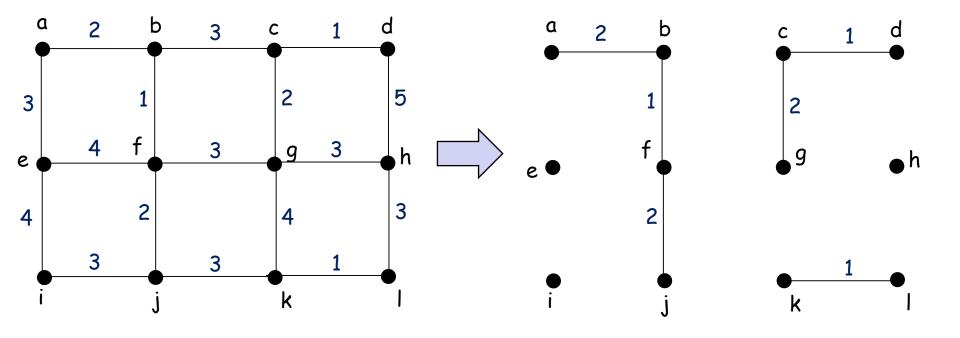


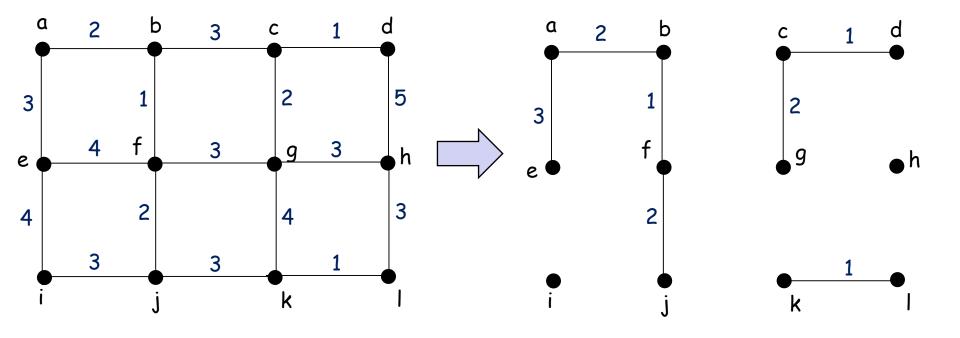


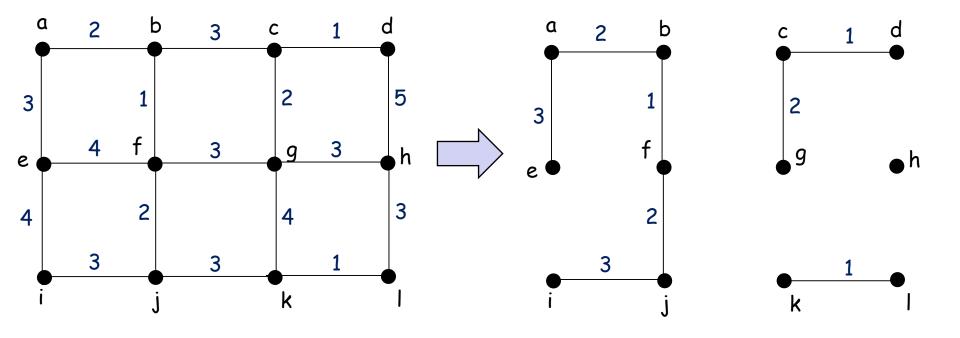


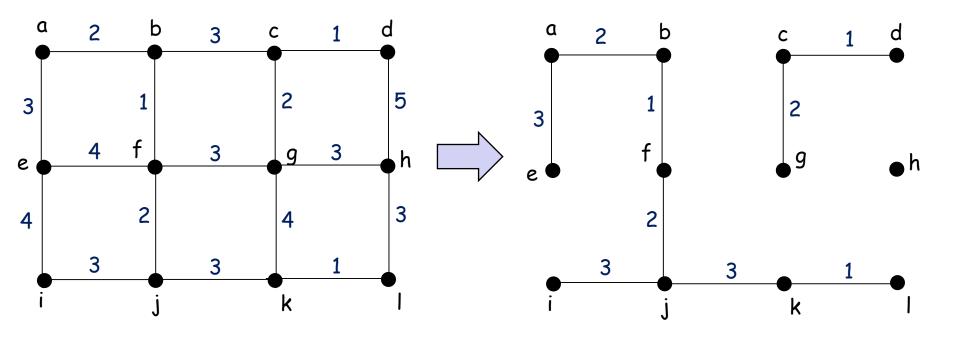


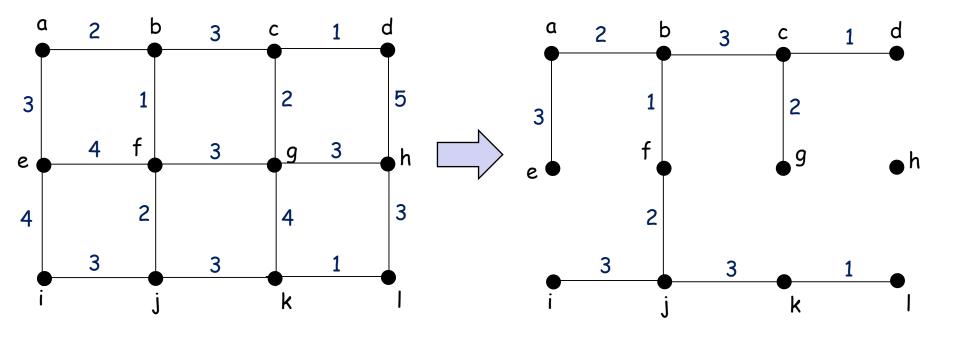


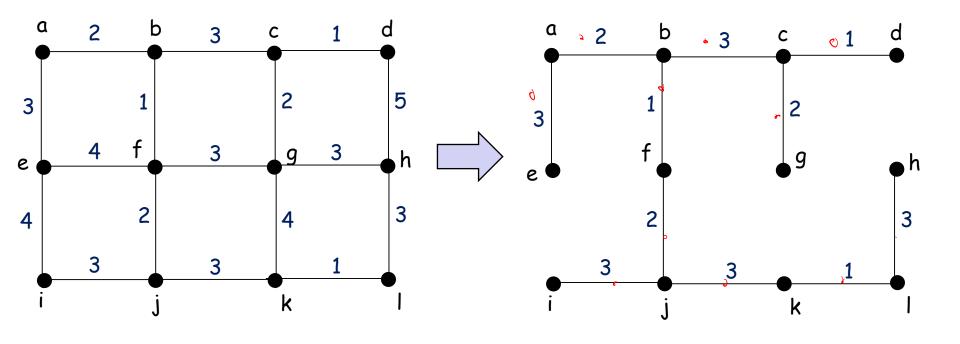


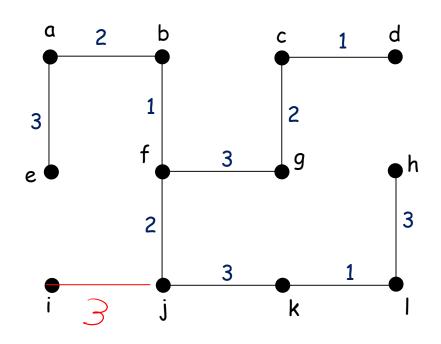


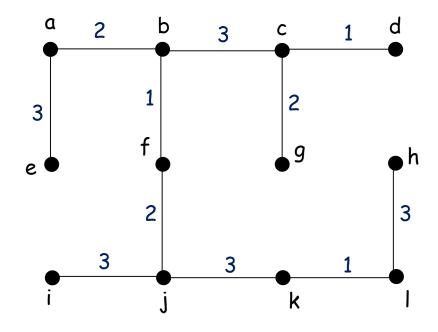








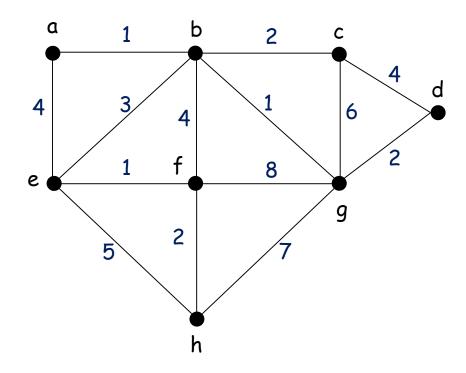


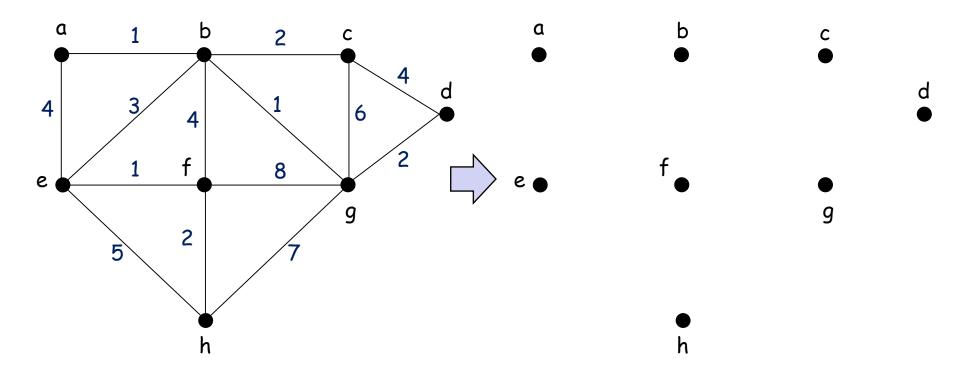


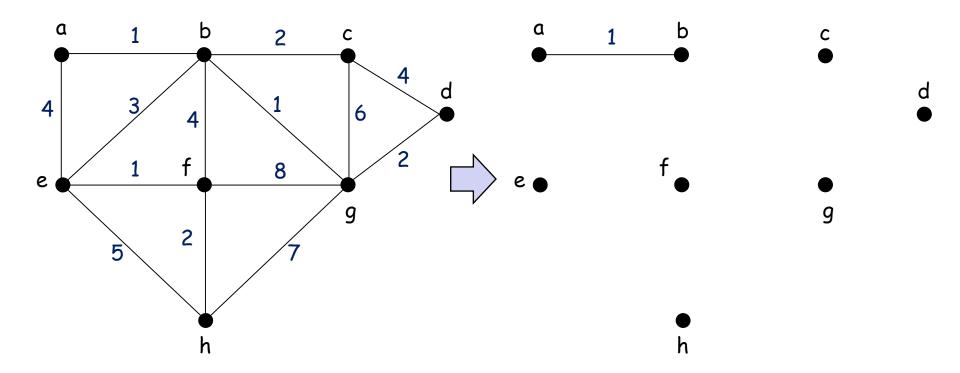
Árbol recubridor mínimo obtenido con el algoritmo de Prim

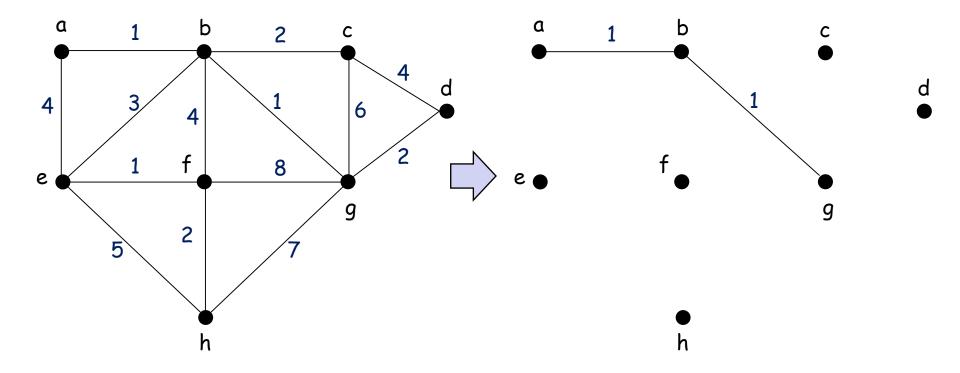
Árbol recubridor mínimo obtenido con el algoritmo de Kruskal

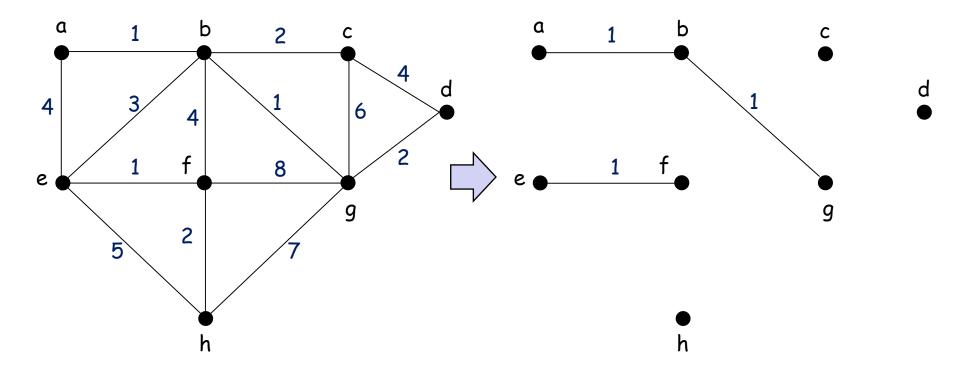
Encontrar un árbol recubridor mínimo usando el algoritmo de Kruskal

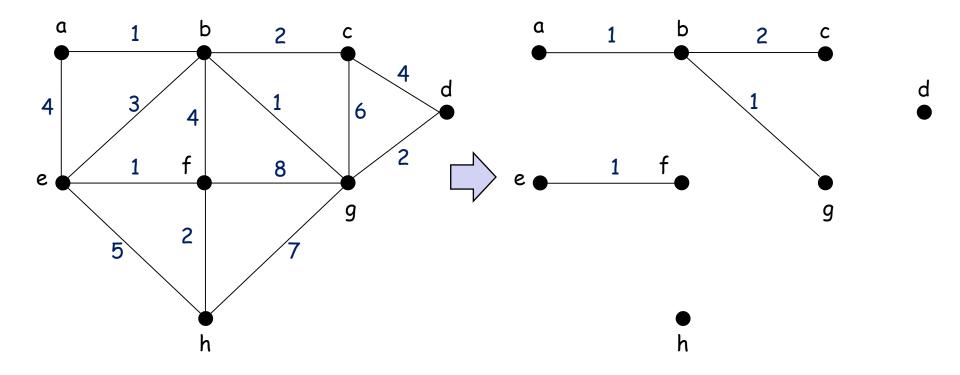


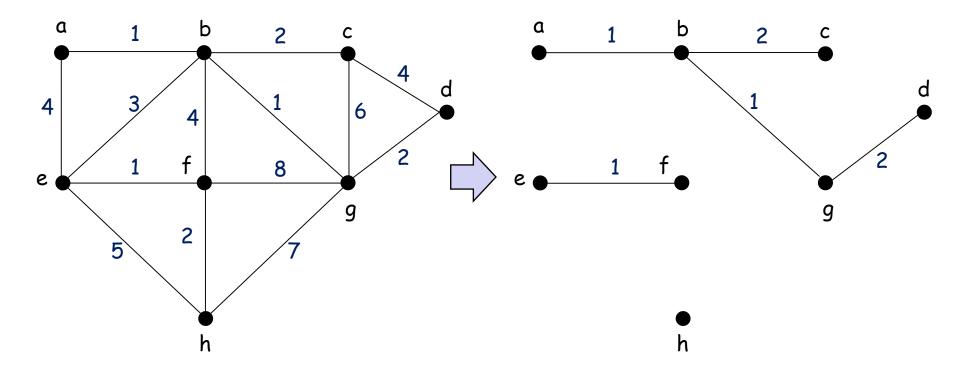


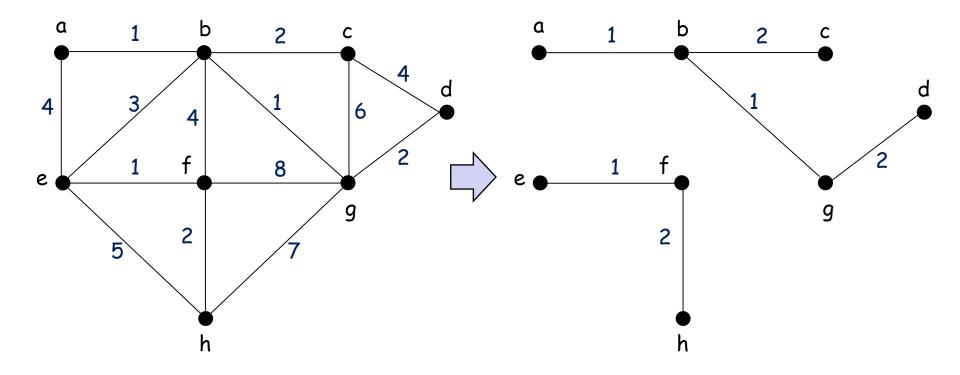


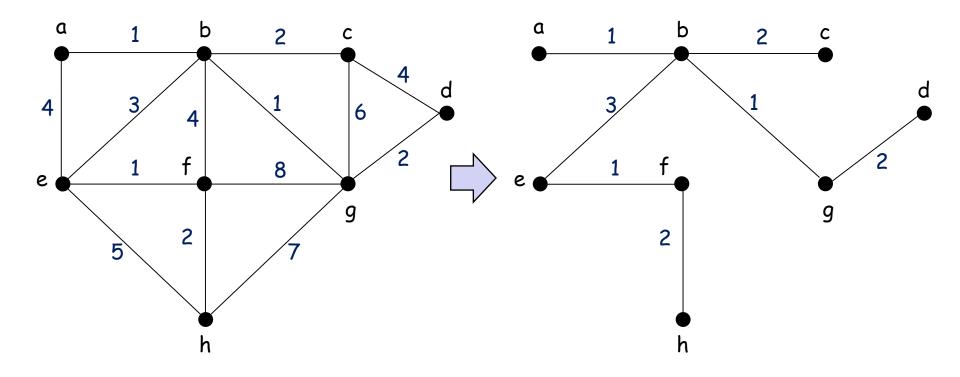


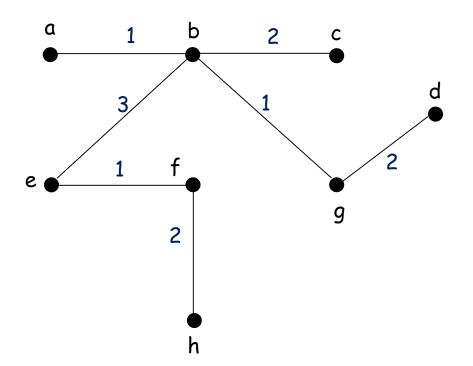




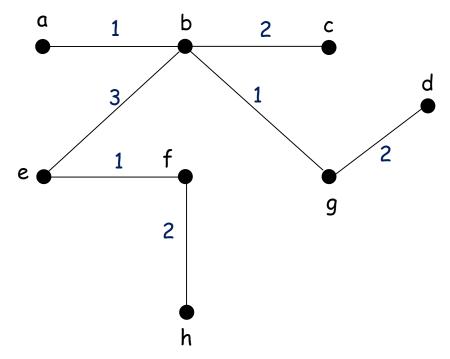








Árbol recubridor mínimo obtenido con el algoritmo de Prim



Árbol recubridor mínimo obtenido con el algoritmo de Kruskal