

Algorithmic Methods for Mathematical Models – COURSE PROJECT –

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Problem Formal Statement

Given:

- The set of D departments of the university indexed by $\{1, \dots, D\}$.
- The set of N faculty members indexed by $\{1, \dots, N\}$. For each member i its department d_i is specified.
- The compatibility m_{ij} between every pair of members i and j .

Find the commission of members of the faculty subject to the following constraints:

- Each department p has exactly n_p members in the commission.
- There cannot be two members i and j in the commission with compatibility $m_{ij} = 0$.
- For each pair of members i and j in the commission with a compatibility lower than 0.15 there must be a member k in the commission with compatibility higher than 0.85 with both members i and j .

with the *objective* to maximize the average compatibility among all pairs of participants in the commission.

ILP Formulation

- N Total number of faculty members, index i .
- D Total number of faculty departments, index p .
- n_p Number of members of department p required in the commission.
- d_i Department to which member i belongs.
- m_{ij} Compatibility between members i and j .

The following decision variables are also defined:

- c_i Binary variable representing if faculty member i is selected to be part of the committee.
- x_{ij} Binary variable representing if both faculty member i and j are selected to be part of the committee.

ILP Formulation

maximize

$$\sum_{i=1}^N \sum_{j=i+1}^N m_{ij} \cdot x_{ij}$$

subject to:

$$x_{ij} \leq c_i \quad \forall i, j \in \{1, \dots, N\} \quad (1)$$

$$x_{ij} \leq c_j \quad \forall i, j \in \{1, \dots, N\} \quad (2)$$

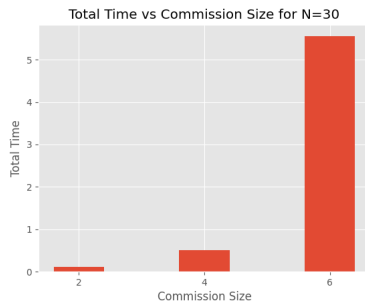
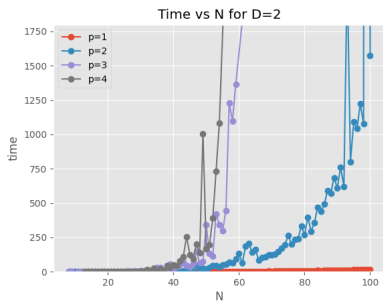
$$x_{ij} \geq c_i + c_j - 1 \quad \forall i, j \in \{1, \dots, N\} \quad (3)$$

$$c_i + c_j \leq 1 \quad \forall i, j \in \{1, \dots, N\} \text{ s.t. } m_{ij} = 0 \quad (4)$$

$$1 - x_{ij} + \sum_{\substack{1 \leq k \leq N, \\ m_{ik}, m_{jk} > 0.85}} c_k \geq 1 \quad \forall i, j \in \{1, \dots, N\} \text{ s.t. } m_{ij} < 0.15 \quad (5)$$

$$\sum_{\substack{1 \leq i \leq N \\ \bar{d}_i = p}} c_i = n_p \quad \forall p \in \{1, \dots, D\} \quad (6)$$

ILP Formulation



N	D	p	size	time
30	2	4	8	3.47
30	4	2	8	0.56
54	2	4	8	1083.16
54	4	2	8	309.06

Algorithm 1: $\text{isFeasible}(C)$

Input: A set C of faculty members.

```
foreach department  $p \in \{1, \dots, D\}$  do  
    if number of members from department  $p$  in  $C \neq n_p$  then  
        return false  
  
foreach pair of members  $i, j \in C$  do  
    if  $m_{ij} = 0$  then  
        return false  
  
foreach pair of members  $i, j \in C$  do  
    if  $m_{ij} < 0.15$  then  
        if there is no member  $k \in C$  s.t.  $m_{ik} > 0.85$  and  $m_{jk} > 0.85$  then  
            return false
```

Committee members' weight

The weight of a faculty member i with respect to a set C of faculty members.

$$w(i, C) = \sum_{j \in C} m_{ij} \quad (7)$$

For the committee C , the objective function can be computed as the average compatibility among all pairs of members in the committee.

$$\text{average_compatibility}(C) = \frac{1}{|C|} \sum_{i \in C} \sum_{j \in C, i < j} m_{ij} \quad (8)$$

Committee members' weight (our approach)

$$w(i, C) = w_1(i, C) + w_2(i, C) \quad (9)$$

The following formula takes into account the relationship of the current committee members and the candidate i :

$$w_1(i, C) = \sum_{j \in C} \begin{cases} 2m_{ij}, & \text{if } 0.15 \leq m_{ij} \leq 0.85 \text{ or } \text{intermediate}(i, j, C), \\ -2, & \text{if } m_{ij} < 0.15 \text{ and not } \text{intermediate}(i, j, C) \end{cases} \quad (10)$$

For further refinement, the algorithm will also take into account the future candidates members not in the committee:

$$w_2(i, C) = \sum_{j \notin C} \begin{cases} m_{ij}, & \text{if } 0.15 \leq m_{ij} \text{ or } \text{intermediate}(i, j, C) \text{ and } \text{isValid}(j, C \cup \{i\}), \\ -1, & \text{if } m_{ij} < 0.15 \text{ and not } \text{intermediate}(i, j, C), \text{ and } \text{isValid}(j, C \cup \{i\}). \end{cases} \quad (11)$$

Greedy Constructive Algorithm: Variables

N : Total number of members

D : Total number of departments

n : Array where n_p is the required members capacity of department p ($1 \leq p \leq D$)

d : Array where d_i is the department of member i ($1 \leq i \leq N$)

m : Compatibility matrix where m_{ij} represents compatibility between members i and j

Initialize $C \leftarrow \emptyset$

Initialize $members \leftarrow \{1, \dots, N\}$

Greedy Constructive Algorithm: Loop

```
while  $C$  not a solution do
   $feasible\_members \leftarrow \{\}$ 
   $members\_weights \leftarrow \{\}$ 
  foreach  $member \notin C$  do
    if  $isFeasible(member, C)$  then
       $feasible\_members \leftarrow feasible\_members \cup \{member\}$ 
       $members\_weights[member] \leftarrow w(member, C)$ 
  if  $|feasible\_members| > 0$  then
     $best\_member \leftarrow \arg \max_{m \in feasible\_members} members\_weights[m]$ 
     $C \leftarrow C \cup \{best\_member\}$ 
```

Greedy Constructive Algorithm: Feasible

To check if the solution is feasible, we can check the not computed constraints.

```
foreach pair of members  $(i, j) \in C \times C$  do  
  if  $m_{ij} < 0.15$  then  
    if there does not exist a  $k \in C$  such that  $m_{ik} > 0.85$  and  
       $m_{jk} > 0.85$  then  
        Print "Relaxed solution is not valid for the original problem"  
        return  $C$  // Relaxed solution to be improved later
```

Local Search

```
C ← current_solution
Comp ← average_compatibility(C)
foreach  $i \in \{1, \dots, N\}$  do
    if  $i \in C$  then
        C_neighbor ← C
        C_neighbor ← remove(C_neighbor, i)
        foreach  $j \in \{1, \dots, N\}$  do
            if  $j \notin C\_neighbor$  then
                C_neighbor ← add(C_neighbor, j)
                if isFeasible(C_neighbor) then
                    Comp_neighbor ← average_compatibility(C_neighbor)
                    if  $Comp\_neighbor > Comp$  then
                        Comp ← Comp_neighbor
                        C ← C_neighbor
```

```
C ← ∅
Candidates ← {1, ..., N}
while C not a solution do
    W ← [0, 0, ..., 0] of size N
    foreach n ∈ Candidates do
        |  $W_n \leftarrow w(n, C)$ 
     $W_{min} \leftarrow \min\{W_n \mid n \in \text{Candidates}\}$ 
     $W_{max} \leftarrow \max\{W_n \mid n \in \text{Candidates}\}$ 
     $RCL_{max} \leftarrow \{n \in \text{Candidates} \mid W_n \geq W_{max} - \alpha(W_{max} - W_{min})\}$ 
    Select n ∈ RCL at random
     $C \leftarrow C \cup \{n\}$ 
    Candidates ← Candidates \ {n}
```

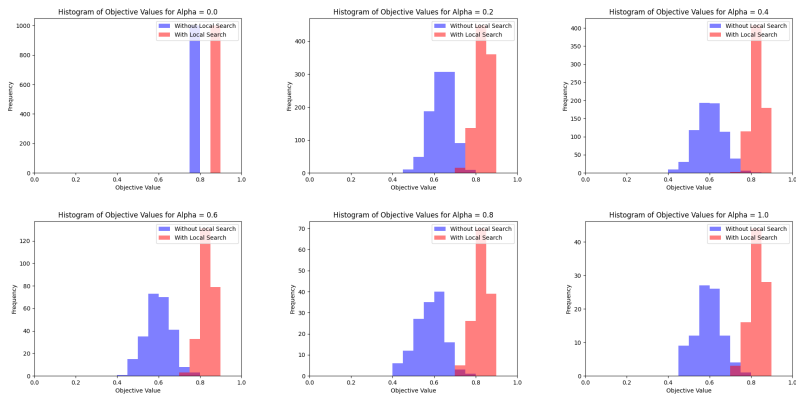
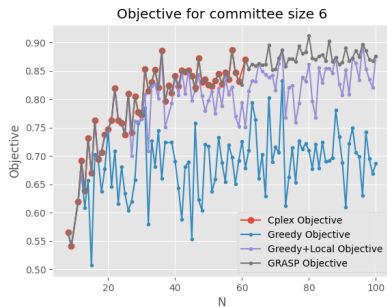
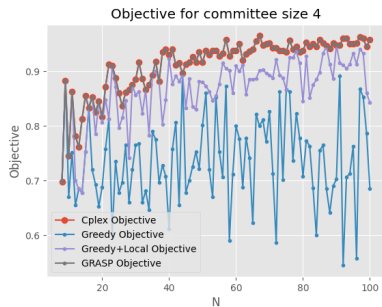
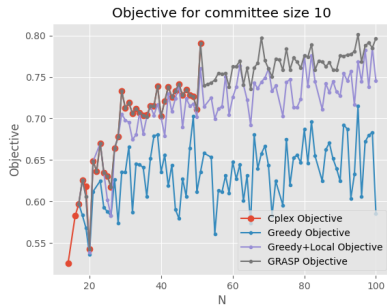
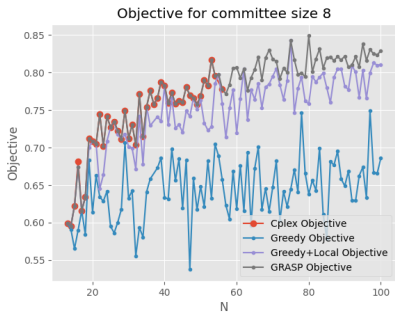
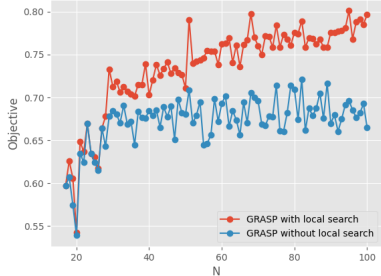


Figure: Distribution of construction phase solution values as a function of the RCL parameter α (1000 repetitions were recorded for each value of α)





Objective for GRASP with and without local search



Time for GRASP with and without local search

