

1. Let $\text{DOUBLE-SAT} = \{\langle \emptyset \rangle \mid \emptyset \text{ has at least two satisfying assignments}\}$ Show DOUBLE-SAT is NP-complete

Proof:

Show that DOUBLE-SAT is NP:

On input \emptyset a nondeterministic polynomial machine can guess two different assignments, and accept if \emptyset returns true for both.

Show that DOUBLE-SAT is NP-complete by mapping $\text{SAT} \leq_p \text{DOUBLE-SAT}$.

TM $M =$ "On input $\langle \emptyset \rangle$ where $\langle \emptyset \rangle$ is the encoding of a boolean formula.

1. Create a new variable x .
2. Let $\emptyset_2 = \emptyset \wedge (x \vee \bar{x})$
3. Output \emptyset_2

Validation: If $\emptyset \in \text{SAT}$, then \emptyset has a least 1 satisfying assignment. Due to this we know that \emptyset_2 has two satisfying assignments:

The assignment that satisfied $\emptyset \wedge x = \text{True}$

The assignment that satisfied $\emptyset \wedge x = \text{False}$

Therefore If $\emptyset \in \text{SAT}$, then $\emptyset_2 \in \text{DOUBLE-SAT}$.

Otherwise $\emptyset \notin \text{SAT}$, then \emptyset has no satisfying assignment, therefore $\emptyset \wedge (x \vee \bar{x})$ will always evaluate to false, because false and (anything) is always false.

Therefore If $\emptyset \notin \text{SAT}$, then $\emptyset_2 \notin \text{DOUBLE-SAT}$.

Thus, $\text{SAT} \leq_p \text{DOUBLE-SAT}$ and DOUBLE-SAT is NP-complete.