Proof Theory in Philosophy

Carlo Nicolai



Slides available at https://carlonicolai.github.io

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What I'm not considering

I don't consider proof-theoretic semantics.

I only briefly touch upon **reductive proof-theory** in the philosophy of mathematics.

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'There never were any set-theoretic paradoxes, but the property theoretic paradoxes are still unresolved' (Gödel to Myhill)

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Abstraction and Truth

'There never were any set-theoretic paradoxes, but the property theoretic paradoxes are still unresolved' (Gödel to Myhill)

Naïve abstraction

$$\forall x \big(x \in \{ v \mid \varphi(v) \} \leftrightarrow \varphi(x) \big)$$

Naïve Truth

$$\operatorname{Tr} \lceil A \rceil \leftrightarrow A$$

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Abstraction and Truth

"There never were any set-theoretic paradoxes, but the property theoretic paradoxes are still unresolved" (Gödel to Myhill)

Naïve abstraction

$$\forall x \big(x \in \{ v \mid \varphi(v) \} \leftrightarrow \varphi(x) \big)$$

Naïve Truth

$$\operatorname{Tr} \lceil A \rceil \leftrightarrow A$$

Here I assume that for any φ in the language there is a term $\{v \mid \varphi(v)\}$ with $FV(\{v \mid \varphi(v)\}) = FV(\varphi) \setminus \{v\}$. If φ is a sentence, I write $\ ^rA \ ^r$ for 'the proposition expressed by A'.

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 $\Gamma, \Delta, \Theta, \Lambda, \dots$ are multisets of formulas.

Truth rules

$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \text{Tr} \lceil A \rceil} \qquad \frac{A, \Gamma \Rightarrow D}{\text{Tr} \lceil A \rceil, \Gamma \Rightarrow D}$$

$$\lambda \Leftrightarrow \neg \text{Tr} \lceil \lambda \rceil \qquad \neg \lambda \Leftrightarrow \text{Tr} \lceil \lambda \rceil$$

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Modal Logic Modal Predicate

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$$\frac{\lambda \Rightarrow \lambda}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow \lambda}$$

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$$\lambda \Leftrightarrow \neg \text{Tr} \, \Gamma \lambda^{\dagger} \qquad \neg \lambda \Leftrightarrow \text{Tr} \, \Gamma \lambda^{\dagger}$$

$$\frac{\frac{\lambda \Rightarrow \lambda}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow \lambda}}{\operatorname{Tr} \lceil \lambda \rceil, \neg \lambda \Rightarrow}$$

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$$\frac{\frac{\lambda \Rightarrow \lambda}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow \lambda}}{\frac{\operatorname{Tr} \lceil \lambda \rceil, \neg \lambda \Rightarrow}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow}}$$

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$$\lambda \Leftrightarrow \neg \text{Tr} \Gamma \lambda^{\dagger}, \qquad \neg \lambda \Leftrightarrow \text{Tr} \Gamma \lambda^{\dagger}$$

$$\frac{\lambda \Rightarrow \lambda}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow \lambda}$$

$$\frac{\operatorname{Tr} \lceil \lambda \rceil, \neg \lambda \Rightarrow}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow}$$

$$\Rightarrow \neg \operatorname{Tr} \lceil \lambda \rceil$$

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$$\lambda \Leftrightarrow \neg \text{Tr} \Gamma \lambda^{\gamma} \qquad \neg \lambda \Leftrightarrow \text{Tr} \Gamma \lambda^{\gamma}$$

$$\frac{\lambda \Rightarrow \lambda}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow \lambda}$$

$$\frac{\operatorname{Tr} \lceil \lambda \rceil, \neg \lambda \Rightarrow}{\operatorname{Tr} \lceil \lambda \rceil \Rightarrow}$$

$$\frac{\lambda \Rightarrow \lambda}{\lambda \Rightarrow \operatorname{Tr} \lceil \lambda \rceil}$$

$$\Rightarrow \neg \operatorname{Tr} \lceil \lambda \rceil \Rightarrow$$

$$\frac{\lambda, \neg \operatorname{Tr} \lceil \lambda \rceil \Rightarrow}{\lambda, \lambda \Rightarrow}$$

$$\Rightarrow \lambda$$

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Truth rules

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$$\lambda \Leftrightarrow \neg \text{Tr} \, \Gamma \lambda^{\text{T}} \qquad \neg \lambda \Leftrightarrow \text{Tr} \, \Gamma \lambda^{\text{T}}$$

$$\frac{\lambda \Rightarrow \lambda}{\operatorname{Tr}'\lambda' \Rightarrow \lambda}$$

$$\frac{\operatorname{Tr}'\lambda', \neg\lambda \Rightarrow}{\operatorname{Tr}'\lambda' \Rightarrow}$$

$$\frac{\neg\operatorname{Tr}'\lambda'}{\Rightarrow \lambda}$$

$$\frac{\lambda \Rightarrow \lambda}{\lambda \Rightarrow \operatorname{Tr}'\lambda'}$$

$$\frac{\lambda, \neg\operatorname{Tr}'\lambda' \Rightarrow}{\lambda, \lambda \Rightarrow}$$

$$\frac{\lambda, \lambda \Rightarrow}{\lambda \Rightarrow}$$

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$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \mathsf{Tr}^{\mathsf{r}} A^{\mathsf{r}}}$$

$$\frac{A, \Gamma \Rightarrow D}{\operatorname{Tr}^{r} A^{r}, \Gamma \Rightarrow D}$$

$$\kappa \Leftrightarrow \operatorname{Tr} \lceil \kappa \rceil \to \bot$$

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$$\kappa \Leftrightarrow \operatorname{Tr} \lceil \kappa \rceil \to \bot$$

$$\frac{\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \mathsf{Tr}^{\lceil \kappa \rceil}} \qquad \bot \Rightarrow \bot}{\kappa, \mathsf{Tr}^{\lceil \kappa \rceil} \rightarrow \bot \Rightarrow \bot}$$

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$$\frac{A, \Gamma \Rightarrow D}{\operatorname{Tr}^{\mathsf{r}} A^{\mathsf{l}}, \Gamma \Rightarrow D}$$

$$\kappa \Leftrightarrow \operatorname{Tr}^{\lceil} \kappa^{\rceil} \to \bot$$

$$\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \operatorname{Tr}^{\lceil \kappa \rceil}} \qquad \perp \Rightarrow \perp$$

$$\frac{\kappa, \operatorname{Tr}^{\lceil \kappa \rceil} \to \perp \Rightarrow \perp}{\kappa, \kappa \Rightarrow \perp}$$

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$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \mathsf{Tr} \, \lceil A \rceil}$$

$$\frac{A, \Gamma \Rightarrow D}{\operatorname{Tr}^{\mathsf{r}} A^{\mathsf{r}}, \Gamma \Rightarrow D}$$

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$$\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \operatorname{Tr} \lceil \kappa \rceil} \qquad \perp \Rightarrow \perp$$

$$\frac{\kappa, \operatorname{Tr} \lceil \kappa \rceil \to \perp \Rightarrow \perp}{\kappa, \kappa \Rightarrow \perp}$$

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$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \operatorname{Tr} \Gamma A}$$

$$\frac{A, \Gamma \Rightarrow D}{\operatorname{Tr}^{\mathsf{r}} A^{\mathsf{r}}, \Gamma \Rightarrow D}$$

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$$\frac{\kappa, \operatorname{Tr}^{\lceil \kappa \rceil} \to \perp \Rightarrow \perp}{\kappa, \kappa \Rightarrow \perp}$$

$$\frac{\kappa, \kappa \Rightarrow \perp}{\kappa \Rightarrow \perp}$$

$$\Rightarrow \operatorname{Tr}^{\lceil \kappa \rceil} \to \perp$$

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Truth rules

$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \mathsf{Tr} \, \lceil A \rceil}$$

$$\frac{A, \Gamma \Rightarrow D}{\operatorname{Tr} \lceil A \rceil, \Gamma \Rightarrow D}$$

$$\kappa \Leftrightarrow \operatorname{Tr} \lceil \kappa \rceil \to \bot$$

$$\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \operatorname{Tr}^{\mathsf{r}} \kappa^{\mathsf{r}}} \qquad \bot \Rightarrow \bot$$

$$\frac{\kappa, \operatorname{Tr}^{\mathsf{r}} \kappa^{\mathsf{r}} \to \bot \Rightarrow \bot}{\kappa, \kappa \Rightarrow \bot}$$

$$\frac{\kappa, \kappa \Rightarrow \bot}{\kappa \Rightarrow \bot}$$

$$\Rightarrow \operatorname{Tr}^{\mathsf{r}} \kappa^{\mathsf{r}} \to \bot$$

$$\Rightarrow \kappa$$

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Truth rules

$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \operatorname{Tr} \Gamma A}$$

$$A, \Gamma \Rightarrow D$$

$$Tr'A', \Gamma \Rightarrow D$$

$$\kappa \Leftrightarrow \operatorname{Tr} \lceil \kappa \rceil \to \bot$$

$$\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \operatorname{Tr} \lceil \kappa \rceil} \qquad 1 \Rightarrow 1$$

$$\frac{\kappa, \operatorname{Tr} \lceil \kappa \rceil \Rightarrow 1 \Rightarrow 1}{\kappa, \kappa \Rightarrow 1}$$

$$\Rightarrow \operatorname{Tr} \lceil \kappa \rceil \Rightarrow 1$$

$$\Rightarrow \kappa$$

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$$\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \operatorname{Tr} \lceil \kappa \rceil} \qquad 1 \Rightarrow 1$$

$$\frac{\kappa, \operatorname{Tr} \lceil \kappa \rceil \Rightarrow 1 \Rightarrow 1}{\kappa \Rightarrow 1}$$

$$\Rightarrow \kappa$$

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$$\frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow \text{Tr}^{\Gamma} A^{\Gamma}} \qquad \frac{A, \Gamma \Rightarrow D}{\text{Tr}^{\Gamma} A^{\Gamma}, \Gamma \Rightarrow D}$$

$$\kappa \Leftrightarrow \text{Tr}^{\Gamma} \kappa^{\Gamma} \Rightarrow \Gamma$$

$$\frac{\kappa \Rightarrow \kappa}{\kappa \Rightarrow \operatorname{Tr}[\kappa]} \qquad 1 \Rightarrow 1$$

$$\frac{\kappa, \operatorname{Tr}[\kappa] \to 1 \Rightarrow 1}{\kappa, \kappa \Rightarrow 1}$$

$$\Rightarrow \operatorname{Tr}[\kappa] \to 1$$

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$$\frac{\kappa, \operatorname{Tr}[\kappa] \to 1 \Rightarrow 1}{\kappa \Rightarrow 1}$$

$$\Rightarrow \kappa$$

$$\Rightarrow 1$$

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Internal Curry

'Consequence' predicate

$$\frac{\Gamma, A \Rightarrow B \qquad \Gamma, C \Rightarrow D}{\Gamma, A, C(\lceil B \rceil, \lceil C \rceil) \Rightarrow D}$$

$$v \Leftrightarrow C(\lceil v \rceil, \lceil \bot \rceil)$$

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 $\frac{\Gamma, A \Rightarrow B}{\Gamma \Rightarrow C(\lceil A \rceil, \lceil B \rceil)}$

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'Consequence' predicate

$$\frac{\Gamma, A \Rightarrow B \qquad \Gamma, C \Rightarrow D}{\Gamma, A, C(\lceil B \rceil, \lceil C \rceil) \Rightarrow D} \qquad \frac{\Gamma, A \Rightarrow B}{\Gamma \Rightarrow C(\lceil A \rceil, \lceil B \rceil)}$$

$$\nu \Leftrightarrow C(\lceil \nu \rceil, \lceil \bot \rceil)$$

$$\frac{\begin{array}{c}
v \Rightarrow v & \bot \Rightarrow \bot \\
v, C(\lceil v\rceil, \lceil \bot \rceil) \Rightarrow \bot \\
\hline
v \Rightarrow \bot \\
\hline
\Rightarrow C(\lceil v\rceil, \lceil \bot \rceil) \\
\hline
\Rightarrow v & v & \bot \Rightarrow \bot
\end{array}$$

$$\frac{v \Rightarrow v & \bot \Rightarrow \bot}{v, C(\lceil v\rceil, \lceil \bot \rceil) \Rightarrow \bot}$$

$$\Rightarrow \bot$$

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$$\frac{\Gamma, A \Rightarrow B \qquad \Gamma, C \Rightarrow D}{\Gamma, A, C(\lceil B \rceil, \lceil C \rceil) \Rightarrow D} \qquad \frac{\Gamma, A \Rightarrow B}{\Gamma \Rightarrow C(\lceil A \rceil, \lceil B \rceil)}$$

$$\nu \Leftrightarrow C(\lceil \nu \rceil, \lceil \bot \rceil)$$

$$\frac{\begin{matrix} \mathbf{v} \Rightarrow \mathbf{v} & \mathbf{1} \Rightarrow \mathbf{1} \\ \\ \mathbf{v}, \mathbf{C}(\lceil \mathbf{v}\rceil, \lceil \mathbf{1}\rceil) \Rightarrow \mathbf{1} \end{matrix}}{\begin{matrix} \mathbf{v} \Rightarrow \mathbf{1} \\ \\ \Rightarrow \mathbf{C}(\lceil \mathbf{v}\rceil, \lceil \mathbf{1}\rceil) \end{matrix}} \qquad \frac{\begin{matrix} \mathbf{v} \Rightarrow \mathbf{v} & \mathbf{1} \Rightarrow \mathbf{1} \\ \\ \\ \mathbf{v}, \mathbf{C}(\lceil \mathbf{v}\rceil, \lceil \mathbf{1}\rceil) \Rightarrow \mathbf{1} \end{matrix}}{\begin{matrix} \mathbf{v} \Rightarrow \mathbf{v} \end{matrix}}$$

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Cut-elimination for truth and abstraction

The main extension of the standard inductive strategy consists in the reduction of cuts of the following form:

Tr -rules principal in the last inferences

$$\frac{\Gamma \Rightarrow \Delta, A}{\Gamma \Rightarrow \Delta, \text{Tr}'A'} \qquad \frac{A, \Gamma \Rightarrow \Delta}{\text{Tr}'A', \Gamma \Rightarrow \Delta}$$

$$\frac{\Gamma \Rightarrow \Delta}{\Gamma \Rightarrow \Delta}$$

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$$\frac{\Gamma \Rightarrow \Delta, A}{\Gamma \Rightarrow \Delta, \operatorname{Tr} \Gamma A} \qquad \frac{A, \Gamma \Rightarrow \Delta}{\operatorname{Tr} \Gamma A^{"}, \Gamma \Rightarrow \Delta}$$

$$\Gamma \Rightarrow \Delta$$

... which we would like to reduce to:

$$\begin{array}{ccc} \mathcal{D}_{o} & \mathcal{D}_{1} \\ \Gamma \Rightarrow \Delta, A & A, \Gamma \Rightarrow \Delta \\ \hline \Gamma \Rightarrow \Delta & \end{array}$$

This creates a problem because $\operatorname{Tr} \lceil A \rceil$ is atomic whereas A is of arbitrary (logical) complexity.

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Tr -measures

 I will consider two ways of keeping track of applications of the truth rules in derivations: the first applies to nodes in the derivation tree, the second applies to single formulas within derivations.

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Tr -measures

- I will consider two ways of keeping track of applications of the truth rules in derivations: the first applies to nodes in the derivation tree, the second applies to single formulas within derivations.
- ▶ In the first case:

$$\frac{\gamma_{o} \Rightarrow \top \alpha}{\gamma_{o} \Rightarrow \text{Tr} \vdash \neg \alpha + 1} \qquad \gamma_{1} \Rightarrow \text{Tr} \vdash \neg \beta}{\gamma_{o}, \gamma_{1} \Rightarrow \text{Tr} \vdash \neg \wedge \text{Tr} \vdash \neg \max(\alpha, \beta)}$$

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Tr -measures

- I will consider two ways of keeping track of applications of the truth rules in derivations: the first applies to nodes in the derivation tree, the second applies to single formulas within derivations.
- ▶ In the first case:

$$\frac{\frac{y_o \Rightarrow T \alpha}{y_o \Rightarrow Tr'T'\alpha+1} \qquad y_1 \Rightarrow Tr'T'\beta}{y_o, y_1 \Rightarrow Tr'T' \land Tr'T'\max(\alpha, \beta)}$$

▶ In the second case:

$$\frac{\gamma_{o} \Rightarrow {}^{\circ}\mathsf{T}}{\gamma_{o} \Rightarrow {}^{1}\mathsf{Tr}^{\mathsf{r}}\mathsf{T}^{\mathsf{r}} \qquad \gamma_{1} \Rightarrow {}^{o}\mathsf{Tr}^{\mathsf{r}}\mathsf{T}^{\mathsf{r}}}$$

$$\gamma_{o}, \gamma_{1} \Rightarrow {}^{\max(1, n)}\mathsf{Tr}^{\mathsf{r}}\mathsf{T}^{\mathsf{r}} \wedge \mathsf{Tr}^{\mathsf{r}}\mathsf{T}^{\mathsf{r}}$$

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Contraction-Free

Systems of truth and 'set theories' can be proved to be consistent via cut elimination arguments Grišin (1982), Petersen (2000), Cantini (2003).

Truth à la Grišin GT

$$\Gamma, \operatorname{Tr} s \Rightarrow \operatorname{Tr} s, \Delta [o] \qquad \Gamma \Rightarrow \tau, \Delta [o] \qquad \Gamma \Rightarrow \tau, \Delta [o]$$

$$\frac{A, \Gamma \Rightarrow \Delta [\alpha]}{\operatorname{Tr} \Gamma A^{\gamma}, \Gamma \Rightarrow \Delta [\alpha + 1]} \qquad \frac{\Gamma \Rightarrow \Delta, A [\alpha]}{\Gamma \Rightarrow \Delta, A_{i} [\alpha]} \qquad \frac{\Gamma \Rightarrow \Delta, A [\alpha]}{\Gamma \Rightarrow \Delta, A [\alpha]} \qquad \frac{\Gamma \Rightarrow \Delta, B [\beta]}{\Gamma \Rightarrow \Delta, B [\max(\alpha, \beta)]}$$

$$\frac{A, B, \Gamma \Rightarrow \Delta [\alpha]}{A \star B, \Gamma \Rightarrow \Delta [\alpha]} \qquad \frac{\Gamma \Rightarrow \Delta, A [\alpha]}{\Gamma \Rightarrow \Delta, A \star B [\alpha + \beta]} \qquad \frac{\Gamma \Rightarrow \Delta, A \star B [\alpha + \beta]}{\Gamma \Rightarrow \Delta, A \star B [\alpha + \beta]}$$

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Contraction-free

Systems of truth and 'set theories' can be proved to be consistent via cut elimination arguments Grišin (1982), Petersen (2000), Cantini (2003).

Lemma

Given cut-free derivations $\mathcal{D}_o \vdash_{\mathrm{GT}} \Gamma \Rightarrow \Delta$, A and $\mathcal{D}_1 \vdash_{\mathrm{GT}} A, \Theta \Rightarrow \Lambda$, there is a $\mathcal{D} \vdash_{\mathrm{GT}} \Gamma, \Theta \Rightarrow \Delta$, Λ with the Tr-rank ρ of \mathcal{D} is $\leq \rho(\mathcal{D}_o) + \rho(\mathcal{D}_1)$.

Proof Idea.

The induction is on $(\rho(\mathcal{D}_0) + \rho(\mathcal{D}_1), |A|, |\mathcal{D}_0| + |\mathcal{D}_1|)$.

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Contraction-free

Systems of truth and 'set theories' can be proved to be consistent via cut elimination arguments Grišin (1982), Petersen (2000), Cantini (2003).

Two problems of the contraction-free approach:

Viewed as a set theory, GS is inconsistent with extensionality, e.g defined as:

$$s\subseteq t\star t\subseteq s, t\in r\Rightarrow s\in r$$

This is often called **Grišin's paradox**.

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This is often called **Grišin's paradox**.

Viewed as a property theory or a truth theory, there is no known, plausible semantics. Basics

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Contraction-free

Systems of truth and 'set theories' can be proved to be consistent via cut elimination arguments Grišin (1982), Petersen (2000), Cantini (2003).

Two problems of the contraction-free approach:

Viewed as a set theory, GS is inconsistent with extensionality, e.g defined as:

$$s \subseteq t \star t \subseteq s, t \in r \Rightarrow s \in r$$

This is often called **Grišin's paradox**.

Viewed as a property theory or a truth theory, there is no known, plausible semantics.

However, it needs to be added that it also features a 'decent' conditional (compared, e.g. to Field (2008)).

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Given our language $\mathcal{L}_{Tr} := \mathcal{L}_{\cup} \{ Tr \}$, we start with a (classical) model \mathcal{M} of \mathcal{L} such that ${}^{r}\varphi^{\gamma \mathcal{M}} = \varphi$, and set, for $X \subset |\mathcal{M}|$:

$$a \in \Phi(X) \Leftrightarrow a = \ulcorner \mathsf{T} \urcorner$$
, or
$$a = \ulcorner \mathsf{Tr} \ulcorner \varphi \urcorner \text{ and } \ulcorner \varphi \urcorner \in X, \text{ or}$$

$$a = \ulcorner \neg \mathsf{Tr} \ulcorner \varphi \urcorner \text{ and } \ulcorner \neg \varphi \urcorner \in X, \text{ or}$$

$$a = \ulcorner \varphi \land \psi \urcorner \text{ and } \ulcorner \varphi \urcorner \in X \text{ and } \ulcorner \psi \urcorner \in X, \text{ or}$$

$$a = \ulcorner \neg (\varphi \land \psi) \urcorner \text{ and } \ulcorner \neg \varphi \urcorner \in X \text{ or } \lnot \neg \psi \urcorner \in X, \text{ or } ...$$

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Given our language $\mathcal{L}_{Tr} := \mathcal{L}_{\cup} \{ Tr \}$, we start with a (classical) model \mathcal{M} of \mathcal{L} such that ${}^{\mathsf{r}} \varphi^{\mathsf{T} \mathcal{M}} = \varphi$, and set, for $X \subset |\mathcal{M}|$:

$$a \in \Phi(X) \Leftrightarrow a = \lceil \top \rceil$$
, or
$$a = \lceil \top \top \lceil \varphi \rceil \rceil \text{ and } \lceil \varphi \rceil \in X, \text{ or }$$

$$a = \lceil \neg \top \top \lceil \varphi \rceil \rceil \text{ and } \lceil \neg \varphi \rceil \in X, \text{ or }$$

$$a = \lceil \varphi \wedge \psi \rceil \text{ and } \lceil \varphi \rceil \in X \text{ and } \lceil \psi \rceil \in X, \text{ or }$$

$$a = \lceil \neg (\varphi \wedge \psi) \rceil \text{ and } \lceil \neg \varphi \rceil \in X \text{ or } \lceil \neg \psi \rceil \in X, \text{ or } ...$$

Let then
$$\Phi^{\circ}(X) = X$$
, $\Phi^{\alpha+1}(X) = \Phi(\Phi^{\alpha}(X))$, $\Phi^{\lambda}(X) = \bigcup_{\beta < \lambda} \Phi^{\beta}(X)$.

Lemma (Kripke (1975), Martin and Woodruff (1975))

If $S \subseteq |\mathcal{M}|$ is a fixed-point of Φ , then for all $\varphi \in \mathcal{L}_{Tr}$:

$$\varphi \in S \text{ iff } Tr \lceil \varphi \rceil \in S$$

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Tr ^{<κ} (⊤)	
Tr <κ(T)	

$$\Phi^{\kappa+1}(\varnothing) = \Phi^{\kappa}(\varnothing)$$

$$\Phi^{\kappa}(\varnothing) = \mathcal{I}_{\Phi}$$

$$\Phi^{\kappa}(\emptyset) = \mathcal{I}_{\mathfrak{A}}$$

	$\operatorname{Tr}(T), T \wedge T \dots$	
	Т	
		$ \mathcal{M} $

 $\Phi^2(\emptyset)$

 $\Phi(\emptyset)$

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:

Tr ^{<κ} (⊺)	$\Phi^{\kappa+1}(\varnothing) = \Phi^{\kappa}(\varnothing)$
Tr ^{<κ} (⊺)	$\Phi^{\kappa}(\varnothing) = \mathcal{I}_{\Phi}$
:	:
·	

	$\operatorname{Tr}(T), T \wedge T \dots$		$\Phi^{2}(\varnothing)$
	Т		$\Phi(\varnothing)$
		$ \mathcal{M} $	Ø

The structure $(\mathcal{M}, \mathcal{I}_{\Phi})$ gives rise to a three-valued model for \mathcal{L}_{Tr} with Tr a 'partial' predicate. Define

$$\mathcal{M} \vDash \Gamma \Rightarrow \Delta : \Leftrightarrow \left(\forall \gamma \in \Gamma \right) |\gamma|_{\mathcal{I}_{\Phi}}^{\mathcal{M}} \neq o \rightarrow \left(\exists \delta \in \Delta \right) |\gamma|_{\mathcal{I}_{\Phi}}^{\mathcal{M}} = 1$$

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Restricting initial sequents

Already known in other contexts Kreuger (1994); Jäger and Stärk (1998); Schroeder-Heister (2016). This is contained in Nicolai (2018). Structural rules are absorbed.

Definition (LPT)

$$\Gamma, \bot \Rightarrow \Delta \qquad \qquad \Gamma \Rightarrow \top, \Delta$$

$$\frac{\Gamma \Rightarrow \Delta, A}{\Gamma \Rightarrow \Delta, \operatorname{Tr} \Gamma A} \qquad \qquad \frac{A, \Gamma \Rightarrow \Delta}{\operatorname{Tr} \Gamma A, \Gamma \Rightarrow \Delta}$$

$$(\neg L) \frac{\Gamma \Rightarrow \Delta, \varphi}{\neg \varphi, \Gamma \Rightarrow \Delta} \qquad \qquad (\neg R) \frac{\varphi, \Gamma \Rightarrow \Delta}{\Gamma \Rightarrow \Delta, \neg \varphi}$$

$$\vdots \qquad \vdots$$

▶ Now $(\mathcal{M}, S) \models LPT$ for S a fixed point of Φ .

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$$(\neg L) \frac{\Gamma \Rightarrow \Delta, \varphi}{\neg \varphi, \Gamma \Rightarrow \Delta} \qquad (\neg R) \frac{\varphi, \Gamma \Rightarrow \Delta}{\Gamma \Rightarrow \Delta, \neg \varphi}$$

$$\vdots \qquad \vdots$$

- ▶ Now $(\mathcal{M}, S) \models LPT$ for S a fixed point of Φ .
- The model $(\mathcal{M}, \mathcal{I}_{\Phi})$ satisfies a fully operational, paracomplete version system of naïve truth based on Strong-Kleene logic (modulo definition of consequence).

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Back to cut elimination

When contraction is around, the notion of Tr -rank is not enough:

$$\begin{array}{c} \mathcal{D}_{\text{oo}} \\ \hline \Gamma \Rightarrow \Delta, \mathsf{Tr}^{\lceil} \psi^{\rceil}, \mathsf{Tr}^{\lceil} \psi^{\rceil} \left[\alpha \right] \\ \hline \hline \Gamma \Rightarrow \Delta, \mathsf{Tr}^{\lceil} \psi^{\rceil} \left[\alpha \right] \\ \hline \Gamma, \Theta \Rightarrow \Delta, \Lambda \left[\alpha + \beta \right] \\ \end{array}$$

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Back to cut elimination

When contraction is around, the notion of Tr -rank is not enough:

$$\frac{\mathcal{D}_{\text{oo}}}{\frac{\Gamma \Rightarrow \Delta, \text{Tr}^{\text{r}}\psi^{\text{l}}, \text{Tr}^{\text{r}}\psi^{\text{l}} \left[\alpha\right]}{\Gamma \Rightarrow \Delta, \text{Tr}^{\text{r}}\psi^{\text{l}} \left[\alpha\right]}} \frac{\mathcal{D}_{\text{l}}}{\text{Tr}^{\text{r}}\psi^{\text{l}}, \Theta \Rightarrow \Lambda \left[\beta\right]}$$

$$\frac{\Gamma \Rightarrow \Delta, \text{Tr}^{\text{r}}\psi^{\text{l}} \left[\alpha\right]}{\Gamma, \Theta \Rightarrow \Delta, \Lambda \left[\alpha + \beta\right]}$$

Now the idea here would be that we transform the derivation in

$$\frac{\mathcal{D}_{\text{oo}}^{*} \qquad \mathcal{D}_{1}^{*}}{\Gamma \Rightarrow \Delta, \psi, \psi \quad [\alpha] \qquad \psi, \Theta \Rightarrow \Lambda \quad [\beta]} \qquad \mathcal{D}_{1}^{*} \\
\frac{\Gamma, \Theta \Rightarrow \Delta, \Lambda, \psi \quad [\alpha + \beta] \qquad \psi, \Theta \Rightarrow \Lambda \quad [\beta]}{\Gamma, \Theta, \Theta \Rightarrow \Delta, \Lambda, \Lambda \quad [2\alpha + \beta]} \\
\frac{\Gamma, \Theta, \Theta \Rightarrow \Delta, \Lambda \quad [2\alpha + \beta]}{\Gamma, \Theta \Rightarrow \Delta, \Lambda \quad [2\alpha + \beta]}$$

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The ordinal Tr -complexity $\kappa_{\mathcal{D}}(\cdot)$ of a formula φ of \mathcal{L}_{Tr} in a derivation \mathcal{D} is defined inductively as follows:

▶ formulas of \mathcal{L} have Tr -complexity o in any \mathcal{D} ;

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- formulas of \mathcal{L} have Tr -complexity o in any \mathcal{D} ;
- If \mathcal{D} is just $\Gamma, \varphi \Rightarrow \varphi, \Delta$ with $\varphi \in \mathcal{L}$, then $\kappa_{\mathcal{D}}(\psi) = \kappa_{\mathcal{D}}(\varphi) = 0$ for all $\psi \in \Gamma, \Delta$. Similarly for $(\top), (\bot)$.

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- formulas of \mathcal{L} have Tr -complexity o in any \mathcal{D} ;
- ► If \mathcal{D} is just Γ , $\varphi \Rightarrow \varphi$, Δ with $\varphi \in \mathcal{L}$, then $\kappa_{\mathcal{D}}(\psi) = \kappa_{\mathcal{D}}(\varphi) = 0$ for all $\psi \in \Gamma$, Δ . Similarly for (\top) , (\bot) .
- If \mathcal{D} ends with

$$\frac{\Gamma \Rightarrow \Delta, A}{\Gamma \Rightarrow \Delta, \mathsf{Tr}^{\mathsf{r}} A^{\mathsf{r}}}$$

then the complexity of formulas in Γ , Δ is unchanged and $\kappa_{\mathcal{D}}(\operatorname{Tr} \lceil A \rceil) = \kappa_{\mathcal{D}}(A) + 1$ (similarly for (TR-L)).

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- If \mathcal{D} ends with

$$\frac{\Gamma \Rightarrow \Delta, A}{\Gamma \Rightarrow \Delta, \operatorname{Tr} \lceil A \rceil}$$

then the complexity of formulas in Γ , Δ is unchanged and $\kappa_{\mathcal{D}}(\operatorname{Tr} \, {}^{\Gamma}A^{-}) = \kappa_{\mathcal{D}}(A) + 1$ (similarly for (TR -L)).

• If \mathcal{D} ends with

$$\frac{\gamma_1^1, \dots, \gamma_n^1 \Rightarrow \delta_1^1, \dots, \delta_m^1, \varphi \qquad \gamma_1^2, \dots, \gamma_n^2 \Rightarrow \delta_1^2, \dots, \delta_m^2, \psi}{\gamma_1^3, \dots, \gamma_n^3 \Rightarrow \delta_1^3, \dots, \delta_m^3, \varphi \land \psi}$$

then

$$\begin{split} \kappa_{\mathcal{D}}(\varphi \wedge \psi) &= \max(\kappa_{\mathcal{D}}(\varphi), \kappa_{\mathcal{D}}(\psi)) \\ \kappa_{\mathcal{D}}(\gamma_i^3) &= \max(\kappa_{\mathcal{D}}(\gamma_i^1), \kappa_{\mathcal{D}}(\gamma_i^2)) \ 1 \leq i \leq n \\ \kappa_{\mathcal{D}}(\delta_j^3) &= \max(\kappa_{\mathcal{D}}(\delta_j^1), \kappa_{\mathcal{D}}(\delta_j^2)) \ 1 \leq j \leq m \end{split}$$

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Full cut elimination

Crucially, rules of LPT are κ -invertible, e.g.: If $\mathcal{D} \vdash_{\text{LPT}} \Gamma^1, \text{Tr} \ulcorner A \urcorner \Rightarrow \Delta^1$, then there is $\mathcal{D}' \vdash_{\text{LPT}} A, \Gamma \Rightarrow \Delta$ with $\kappa_{\mathcal{D}'}(\Gamma) \leq \kappa_{\mathcal{D}}(\Gamma^1), \kappa_{\mathcal{D}'}(\Delta) \leq \kappa_{\mathcal{D}}(\Delta^1)$, and $\kappa_{\mathcal{D}'}(A) \leq \kappa_{\mathcal{D}}(\text{Tr} \ulcorner A \urcorner)$, if $\kappa_{\mathcal{D}}(\text{Tr} \ulcorner A \urcorner) = 0$; $\kappa_{\mathcal{D}'}(A) < \kappa_{\mathcal{D}}(\text{Tr} \ulcorner A \urcorner)$, if $\kappa_{\mathcal{D}}(\text{Tr} \ulcorner A \urcorner) \neq 0$.

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Full cut elimination

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Lemma

Contraction is κ -admissible and length-admissible, e.g.: If $\mathcal{D} \vdash_{\mathrm{LPT}}^{n} \Gamma^{1}, \varphi^{1}, \varphi^{2} \Rightarrow \Delta^{1}$, then there is a $\mathcal{D}' \vdash_{\mathrm{LPT}}^{n} \Gamma, \varphi \Rightarrow \Delta$ with $\kappa_{\mathcal{D}'}(\Gamma^{1}) \leq \kappa_{\mathcal{D}'}(\Gamma)$; $\kappa_{\mathcal{D}'}(\Delta^{1}) \leq \kappa_{\mathcal{D}'}(\Delta)$ $\kappa_{\mathcal{D}'}(\varphi) \leq \max(\kappa_{\mathcal{D}'}(\varphi^{1}), \kappa_{\mathcal{D}'}(\varphi^{2}))$.

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Full cut elimination

Crucially, rules of LPT are κ -invertible, e.g.:

If $\mathcal{D} \vdash_{\mathrm{LPT}} \Gamma^1$, $\mathrm{Tr} \lceil A \rceil \Rightarrow \Delta^1$, then there is $\overline{\mathcal{D}'} \vdash_{\mathrm{LPT}} A$, $\Gamma \Rightarrow \Delta$ with $\kappa_{\mathcal{D}'}(\Gamma) \leq \kappa_{\mathcal{D}}(\Gamma^1)$, $\kappa_{\mathcal{D}'}(\Delta) \leq \kappa_{\mathcal{D}}(\Delta^1)$, and

$$\kappa_{\mathcal{D}'}(A) \le \kappa_{\mathcal{D}}(\operatorname{Tr}^{\lceil}A^{\rceil}), \text{ if } \kappa_{\mathcal{D}}(\operatorname{Tr}^{\lceil}A^{\rceil}) = 0;$$

$$\kappa_{\mathcal{D}'}(A) < \kappa_{\mathcal{D}}(\operatorname{Tr}^{\lceil}A^{\rceil}), \text{ if } \kappa_{\mathcal{D}}(\operatorname{Tr}^{\lceil}A^{\rceil}) \neq 0.$$

Lemma

Contraction is κ -admissible and length-admissible, e.g.: If $\mathcal{D} \vdash_{\mathrm{LPT}}^{n} \Gamma^{1}, \varphi^{1}, \varphi^{2} \Rightarrow \Delta^{1}$, then there is a $\mathcal{D}' \vdash_{\mathrm{LPT}}^{n} \Gamma, \varphi \Rightarrow \Delta$ with $\kappa_{\mathcal{D}'}(\Gamma^{1}) \leq \kappa_{\mathcal{D}'}(\Gamma)$; $\kappa_{\mathcal{D}'}(\Delta^{1}) \leq \kappa_{\mathcal{D}'}(\Delta)$ $\kappa_{\mathcal{D}'}(\varphi) \leq \max(\kappa_{\mathcal{D}'}(\varphi^{1}), \kappa_{\mathcal{D}'}(\varphi^{2}))$.

Proposition

If \mathcal{D}_{\circ} is a cut-free proof of $\Gamma^{1} \Rightarrow \Delta^{1}$, φ^{1} in LPT, and \mathcal{D}_{1} is a cut-free LPT-proof of φ^{2} , $\Gamma^{2} \Rightarrow \Delta^{2}$, then there is a cut-free proof \mathcal{D} of $\Gamma^{3} \Rightarrow \Delta^{3}$ with $\kappa_{\mathcal{D}}(\Gamma^{3}) \leq \max(\kappa_{\mathcal{D}_{\circ}}(\Gamma^{1}), \kappa_{\mathcal{D}_{1}}(\Gamma^{2}))$ and $\kappa_{\mathcal{D}}(\Delta^{3}) \leq \max(\kappa_{\mathcal{D}_{\circ}}(\Delta^{1}), \kappa_{\mathcal{D}_{1}}(\Delta^{2}))$.

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 The ideal of semantic closure is at odds with resources that outstrip the ones available in one's semantic theory.
 Cut-elimination procedures are usually formalizable in weak arithmetical systems. Basics

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- The ideal of semantic closure is at odds with resources that outstrip the ones available in one's semantic theory.
 Cut-elimination procedures are usually formalizable in weak arithmetical systems.
- When nonlogical initial sequents are around, full cut elimination is not in general available: however by eliminating cuts on formulas of the form $\operatorname{Tr} \lceil A \rceil$, one can obtain **conservativity proofs**.

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- The ideal of semantic closure is at odds with resources that outstrip the ones available in one's semantic theory.
 Cut-elimination procedures are usually formalizable in weak arithmetical systems.
- When nonlogical initial sequents are around, full cut elimination is not in general available: however by eliminating cuts on formulas of the form $\operatorname{Tr}^{r}A^{r}$, one can obtain **conservativity proofs**.
- Another advantage of the approach with restricted initial sequents is that unlike the contraction-free approaches there are natural **infinitary systems** that arise from the logic and that give succinct presentations of Π_1^1 -sets.

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Truth Bearers

There are good reasons to require an ontology of bearers of truth prior to discussing principles of truth. We want to prove in the object language things like:

$$\forall \varphi, \psi \,\exists \chi \, (\chi = (\varphi \dot{\wedge} \psi) \wedge \varphi \neq \chi)$$

This is usually guaranteed by assuming a theory of finite objects (as we shall see in a moment).

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Truth Bearers

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$$\forall \varphi, \psi \,\exists \chi \, (\chi = (\varphi \dot{\wedge} \psi) \wedge \varphi \neq \chi)$$

This is usually guaranteed by assuming a theory of finite objects (as we shall see in a moment).

Notice that this is imposing non-trivial constraints. More 'philosophical' theories of truth are often formulated in terms of propositions, and not sentence types (Horwich, 1998; Soames, 1998; Jago, 2018). This rules out that propositions are coarse-grained, e.g. sets of possible worlds. Basics

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Arithmetic

▶ Peano arithmetic (PA) is the preferred base theory for systems of truth. It is usually formulated in $\mathcal{L}_{\mathbb{N}} = \{o, S, +, \times\}$ and features equations for its primitives, e.g.

$$(x + o) = x x + Sy = S(x + y)$$

and the first-order induction schema

$$\varphi(o) \land \forall x (\varphi(x) \to \varphi(Sx)) \to \forall x \varphi(x)$$
 for $\varphi \in \mathcal{L}_{\mathbb{N}}$

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Arithmetic

▶ Peano arithmetic (PA) is the preferred base theory for systems of truth. It is usually formulated in $\mathcal{L}_{\mathbb{N}} = \{o, S, +, \times\}$ and features equations for its primitives, e.g.

$$(x + o) = x$$
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$$\varphi(o) \land \forall x (\varphi(x) \to \varphi(Sx)) \to \forall x \varphi(x)$$
 for $\varphi \in \mathcal{L}_{\mathbb{N}}$

Alternatively, one can employ a theory of strings and concatenation $\hat{}$ with two atoms a, b based on **Tarski's axiom** $a\hat{} y = u\hat{} v \leftrightarrow \exists w \big((x = u\hat{} w \land v = w\hat{} y) \lor (u = x\hat{} w \land y = w\hat{} v) \big)$ With first-order string induction, the two theories are mutually interpretable. An accessible source is Ganea (2009).

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With first-order string induction, the two theories are mutually

With first-order string induction, the two theories are mutually interpretable. An accessible source is Ganea (2009).

- Finite set theories are also a convenient choice. For instance
 - Kaye and Wong (2006) show that PA and KF \setminus {Inf}+ 'every set has a transitive closure' are bi-interpretable;
 - similarly, a neat set theory based by Świerczkowski (2003) based on the **adjunction operation** $x \triangleleft y \mapsto x \cup \{y\}$ is bi-interpretable with PA.

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Doing with less

 Ultimately, what we require to establish the basic properties of the truth bearers are a good notion of sequence, and a minimum of induction to handle suitable forms of recursion.

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Doing with less

- Ultimately, what we require to establish the basic properties of the truth bearers are a good notion of sequence, and a minimum of induction to handle suitable forms of recursion.
- For the former the notion of a sequential theory is enough see Visser (2010) for a comprehensive overview. A theory is sequential if it interprets – with no relativization of quantifiers – the theory AS given by the empyt set axiom and adjunction – which is as strong as Robinson's Q.

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Doing with less

- Ultimately, what we require to establish the basic properties of the truth bearers are a good notion of sequence, and a minimum of induction to handle suitable forms of recursion.
- For the former the notion of a sequential theory is enough see Visser (2010) for a comprehensive overview. A theory is sequential if it interprets – with no relativization of quantifiers – the theory AS given by the empyt set axiom and adjunction – which is as strong as Robinson's Q.
- As to induction, since all the relevant syntactic notions (terms, formulas, proofs) are p-time decidable, the theory S_2^1 by Buss (1986) suffices. However, many of the results that I will treat below are specific to PA (or equivalents), and it is object of current research to check which results are stable over $T \supseteq S_2^1$.

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• A good base theory will also provide a satisfactory representation of ordinals. For our purposes it suffices to require a notation up to the Feferman-Schütte ordinal Γ_0 :

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- A good base theory will also provide a satisfactory representation of ordinals. For our purposes it suffices to require a notation up to the Feferman-Schütte ordinal Γ_0 :
 - α is principal if it cannot be expressed as $\zeta + \eta$ for ζ , $\eta < \alpha$. Define:

$$C(o) :=$$
 'the class of principal ordinals'
$$C(\alpha + 1) :=$$
 'the class of fixed points of the function enumerating $C(\alpha)$ '
$$C(\lambda) := \bigcap_{\zeta < \lambda} C(\zeta) \text{ for } \lambda \text{ a limit ordinal}$$

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- A good base theory will also provide a satisfactory representation of ordinals. For our purposes it suffices to require a notation up to the Feferman-Schütte ordinal Γ_0 :
 - α is principal if it cannot be expressed as $\zeta + \eta$ for $\zeta, \eta < \alpha$. Define:

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• The Veblen functions φ_{α} are the enumerating functions of $C(\alpha)$. The class of *strongly critical* ordinals SC contains precisely the ordinals α that are themselves α -critical. Γ_{ζ} indicates the ζ -th strongly critical ordinal.

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- The *Veblen functions* φ_{α} are the enumerating functions of $C(\alpha)$. The class of *strongly critical* ordinals SC contains precisely the ordinals α that are themselves α -critical. Γ_ζ indicates the ζ-th strongly critical ordinal.
- Principal ordinals α that are not themselves strongly critical are such that $\alpha = \varphi_{\zeta} \eta$ for $\eta, \zeta < \alpha$. Therefore, by this fact and Cantor's normal form theorem, ordinals $< \Gamma_o$ can be uniquely determined as words of the alphabet $(o, +, \varphi.\cdot)$.

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- A good base theory will also provide a satisfactory representation of ordinals. For our purposes it suffices to require a notation up to the Feferman-Schütte ordinal Γ_0
- A notation system for Γ_0 is of the form (OT, PT, $|\cdot|$, <), with
 - $\blacktriangleright\,$ OT the set of natural number 'codes' for ordinals $<\Gamma_{o}$
 - ▶ $OT \subseteq OT$ the set of codes of principal ordinals
 - $|\cdot|: OT \rightarrow ON$
 - $\qquad n < m : \leftrightarrow n \in \mathtt{OT} \land m \in \mathtt{OT} \land |n| < |m|$

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Ordinals

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- A notation system for Γ_0 is of the form (OT, PT, $|\cdot|$, <), with
 - \blacktriangleright 0T the set of natural number 'codes' for ordinals $<\Gamma_o$
 - OT ⊆ OT the set of codes of principal ordinals
 - $|\cdot|: OT \to ON$
 - $n < m : \leftrightarrow n \in \mathsf{OT} \land m \in \mathsf{OT} \land |n| < |m|$
- Using standard coding techniques one can show that OT, PT, \prec are primitive recursive. Actually, Beckmann et al. (2003) show that they can be showed to be p-time and represented in S_2^1 notice that I **do not** mean that Γ_0 can be well-founded in S_2^1 !

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Schemata

For ordinals $\alpha < \Gamma_0$, we denote with a the corresponding numeral in the representation of OT and we do not distinguish between ordinal functions such as the Veblen functions and their representations.

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Schemata

For ordinals $\alpha < \Gamma_o$, we denote with a the corresponding numeral in the representation of OT and we do not distinguish between ordinal functions such as the Veblen functions and their representations. The system (OT, PT, \prec) enables us to formulate the following principles of transfinite induction:

$$(\operatorname{TI}_{\mathcal{L}_{\operatorname{Tr}}}^{\varepsilon_{\circ}}) \qquad \frac{\forall a < b \, \phi(a), \Gamma \Rightarrow \Delta, \phi(b)}{\Gamma \Rightarrow \Delta, \forall a < \varepsilon_{\circ} \, \phi(a)}$$

$$\forall a < b \, \phi(a), \Gamma \Rightarrow \Delta, \phi(b)$$

$$(\operatorname{TI}_{\mathcal{L}_{\operatorname{Tr}}}^{<\omega^{\omega}}) \qquad \frac{\forall a < b \, \phi(a), \, \Gamma \Rightarrow \Delta, \, \phi(b)}{\Gamma \Rightarrow \Delta, \, \forall a < c \, \phi(a)} \quad \text{for all } \gamma(=|c|) < \omega^{\omega}$$

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Truth as primitive

 Truth-theoretic deflationism holds that truth is not a genuine property and that its function is mainly that of a generalizing device (Quine, 1970; Field, 1994; Horwich, 1998).

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Truth as primitive

- Truth-theoretic deflationism holds that truth is not a genuine property and that its function is mainly that of a generalizing device (Quine, 1970; Field, 1994; Horwich, 1998).
- ► Unlike other notions that have been taken to be primitive for lack of consensus over a definition e.g. knowledge, see Williamson (2000) Tarski's theorem uncontroversially establishes this (Halbach, 2014, Ch. 1).

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Truth as primitive

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- ▶ Unlike other notions that have been taken to be primitive for lack of consensus over a definition e.g. knowledge, see Williamson (2000) Tarski's theorem uncontroversially establishes this (Halbach, 2014, Ch. 1).
- Truth is a fundamental semantic concept. A theory of meaning for natural language expressions is not much more than a (Tarskian) theory of truth for it (Davidson, 1984).

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Tarskian Truth

The theory of truth in \mathcal{L}_{Tr} that Davidson had in mind extends PA with the following:

Definition (CT)

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\forall s, t (\operatorname{Tr}(s=t) \leftrightarrow s^{\circ} = t^{\circ})
\forall \varphi \in \mathcal{L}(\operatorname{Tr}(\neg \varphi) \leftrightarrow \neg \operatorname{Tr} \varphi)
\forall \varphi, \psi \in \mathcal{L}(\operatorname{Tr}(\varphi \land \psi) \leftrightarrow \operatorname{Tr} \varphi \land \operatorname{Tr} \psi)
\forall v, \forall \varphi(v) \in \mathcal{L}(\operatorname{Tr}(\forall v \varphi) \leftrightarrow \forall x \operatorname{Tr} \varphi(\dot{x}))
\varphi(o) \land \forall x (\varphi(x) \rightarrow \varphi(x+1)) \rightarrow \forall x \varphi(x) \qquad \text{with } \varphi(v) \in \mathcal{L}_{\operatorname{Tr}}
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$$\forall v, \forall \varphi(v) \in \mathcal{L}(\operatorname{Tr}(\forall v\varphi) \leftrightarrow \forall x \operatorname{Tr} \varphi(\dot{x}))$$

$$\varphi(o) \land \forall x(\varphi(x) \rightarrow \varphi(x+1)) \rightarrow \forall x \varphi(x) \qquad \text{with } \varphi(v) \in \mathcal{L}_{\operatorname{Tr}}$$

Important variations are obtained by tweaking the induction schema:

- ▶ CT \ (a.k.a. CT) is obtained by restricting induction to \mathcal{L}
- ► CT_{int} is obtained by adopting the internal induction schema

$$\forall \varphi(v) (\operatorname{Tr} \varphi(o/v) \wedge \forall y (\operatorname{Tr} \varphi(\dot{y}/v) \to \operatorname{Tr} \varphi(\dot{S}y/v)) \to \forall x \operatorname{Tr} \varphi(\dot{x}/v))$$

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Thesis (Shapiro, 1998; Ketland, 1999)

CT proves Con(PA), therefore deflationism is untenable.

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This contrasts with:

Proposition (Kotlarski et al. (1981); Visser and Enayat (2015); Leigh (2015))

 $\ensuremath{\text{CT}_{\text{int}}}$ is a conservative extension of PA.

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CT proves Con(PA), therefore deflationism is untenable.

This contrasts with:

Proposition (Kotlarski et al. (1981); Visser and Enayat (2015); Leigh (2015))

CT_{int} is a conservative extension of PA.

The discussion took a strong technical turn, brilliantly summarized in Cieliski (2017) – with many original contributions. It's worth mentioning:

Proposition (Enayat and Pakhomov (2018))

CT ↑ plus 'disjunctive correctness', i.e.

$$\forall s \big(\mathit{Tr} \big(\bigvee_{i < s} \varphi_i \big) \leftrightarrow \exists i < s \; \mathit{Tr} \, \varphi_i \big)$$

is the same theory as $CT[I\Delta_o]$, and therefore proves Con(PA).

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Despite the technical interest, the debate seems to be built on shaky foundations. Virtually no deflationist has thoroughly defended the claim that truth has to be conservative over the base theory.

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- Despite the technical interest, the debate seems to be built on shaky foundations. Virtually no deflationist has thoroughly defended the claim that truth has to be conservative over the base theory.
- By contrast, it has repeatedly been argued that truth has to be nonconservative, but in a way that is distinctively metalinguistic, i.e. it does not interfere with the subject matter of the base theory over which truth is built.

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- Despite the technical interest, the debate seems to be built on shaky foundations. Virtually no deflationist has thoroughly defended the claim that truth has to be conservative over the base theory.
- By contrast, it has repeatedly been argued that truth has to be nonconservative, but in a way that is distinctively metalinguistic, i.e. it does not interfere with the subject matter of the base theory over which truth is built.
- This has led to the programme of 'disentangling' syntactic quantifiers from quantifiers over natural numbers:

Proposition (Nicolai (2015, 2016))

If one formulates CT \ as a two-sorted theory, with 'syntactic' quantifiers and 'number-theoretic' quantifiers, and truth applying only over syntactic objects, then:

- ► The theory of truth becomes *trivially* conservative over PA;
- This version of CT ↑ plus 'all axioms of PA are true' is mutually interpretable with PA + Con(PA).

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To give a nice presentation of the **reflective closure** of PA – and possibly of further 'natural' theories: i.e. the (truth-)theory that makes explicit all that is implicit in the acceptance of PA.

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To give a nice presentation of the **reflective closure** of PA – and possibly of further 'natural' theories: i.e. the (truth-)theory that makes explicit all that is implicit in the acceptance of PA.

First attempt: $CT_{< a}$, $a \le \Gamma_0$

$$\mathcal{L}_{o} \coloneqq \mathcal{L}_{\mathsf{Tr}}$$
 $\mathcal{L}_{<\gamma} \coloneqq \mathcal{L} \cup \{\mathsf{Tr}_{b} \mid b < \gamma\}$

With $b < \gamma$:

$$\forall s, t(\operatorname{Tr}_{b}(s=t) \leftrightarrow s^{\circ} = t^{\circ})$$

$$\forall \varphi \in \mathcal{L}_{

$$\forall \varphi, \psi \in \mathcal{L}_{< b} (\operatorname{Tr}_{b}(\varphi \wedge \psi) \leftrightarrow \operatorname{Tr}_{b}\varphi \wedge \operatorname{Tr}_{b}\psi)$$

$$\forall v, \forall \varphi(v) \in \mathcal{L}(\operatorname{Tr}_{b}(\forall v\varphi) \leftrightarrow \forall x \operatorname{Tr}_{b}\varphi(\dot{x}))$$

$$\forall \varphi \in \mathcal{L}_{a< b}(\operatorname{Tr}_{b}\operatorname{Tr}_{a}\varphi \leftrightarrow \operatorname{Tr}_{b}\varphi)$$

$$\forall d < b, \forall \varphi \in \mathcal{L}_{< d}(\operatorname{Tr}_{b}\operatorname{Tr}_{d}\varphi \leftrightarrow \operatorname{Tr}_{b}\varphi)$$$$

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▶ The project of isolating $CT_{<\epsilon_o}$ or $CT_{<\Gamma_o}$ as **natural stopping points**, that was congenial to Feferman's project, depended essentially on other results, such as Feferman's and Schütte's independent characterization of Γ_o , or the provable well-orderings of PA.

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- The project of isolating $CT_{<\epsilon_o}$ or $CT_{<\Gamma_o}$ as **natural stopping points**, that was congenial to Feferman's project, depended essentially on other results, such as Feferman's and Schütte's independent characterization of Γ_o, or the provable well-orderings of PA.
- The next step was to find an independent characterization of such theories:

Definition (KF)

```
\forall s, t(\operatorname{Tr}(s=t) \leftrightarrow s^{\circ} = t^{\circ})
\forall s, t(\operatorname{Tr}(s\neq t) \leftrightarrow s^{\circ} \neq t^{\circ})
\forall t(\operatorname{Tr}\operatorname{Tr} t \leftrightarrow \operatorname{Tr} t^{\circ})
\forall t(\operatorname{Tr}\operatorname{Tr} \tau \leftrightarrow \operatorname{Tr} \tau^{\circ})
\forall \varphi \in \mathcal{L}_{\operatorname{Tr}}(\operatorname{Tr}(\neg \neg \varphi) \leftrightarrow \operatorname{Tr} \varphi)
\forall \varphi, \psi \in \mathcal{L}_{\operatorname{Tr}}(\operatorname{Tr}(\varphi \wedge \psi) \leftrightarrow \operatorname{Tr} \varphi \wedge \operatorname{Tr} \psi)
\forall \varphi, \psi \in \mathcal{L}_{\operatorname{Tr}}(\operatorname{Tr} \neg (\varphi \wedge \psi) \leftrightarrow \operatorname{Tr} \neg \varphi \vee \operatorname{Tr} \neg \psi)
\forall v, \forall \varphi(v) \in \mathcal{L}_{\operatorname{Tr}}(\operatorname{Tr}(\forall v\varphi) \leftrightarrow \forall x \operatorname{Tr} \varphi(\dot{x}))
\forall v, \forall \varphi(v) \in \mathcal{L}_{\operatorname{Tr}}(\operatorname{Tr}(\neg \forall v\varphi) \leftrightarrow \exists x \operatorname{Tr} \neg \varphi(\dot{x}))
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• Semantically KF fits nicely with Kripke's fixed-point semantics (Kripke, 1975),

 $(\mathbb{N}, S) \models KF$ iff S is a fixed point of Kripke's theory of truth

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$$(\mathbb{N}, S) \models KF \text{ iff } S \text{ is a fixed point of Kripke's theory of truth}$$

► The full Tr -schema is available for meaningful predicates satisfying $D(x) : \leftrightarrow \text{Tr } x \lor \text{Tr } \neg x$, i.e. for all $A \in \mathcal{L}_{\text{Tr}}$:

$$KF \vdash D(\lceil A \rceil) \rightarrow (Tr \lceil A \rceil \leftrightarrow A)$$

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► The full Tr -schema is available for meaningful predicates satisfying $D(x) :\leftrightarrow Tr x \lor Tr \neg x$, i.e. for all $A \in \mathcal{L}_{Tr}$:

$$KF \vdash D(\lceil A \rceil) \rightarrow (Tr \lceil A \rceil \leftrightarrow A)$$

Proposition (Feferman (1991); Cantini (1989))

KF is proof-theoretically equivalent to $(\Pi_1^{\circ}-CA)_{<\epsilon_{\circ}}$.

Proof Idea.

Lower bound: PA in \mathcal{L}_{Tr} proves $TI_{\mathcal{L}_{Tr}}^{\langle \epsilon_o}$. Now KF proves:

$$\varphi \in \mathcal{L}_{< a} \to D(\varphi) \Rightarrow \varphi \in \mathcal{L}_a \to D(\varphi)$$

An application of $TI_{\mathcal{L}_{Tr}}^{<\epsilon_o}$ yields an embedding of $CT_{<\epsilon_o}$, which suffices.

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Proposition (Feferman (1991); Cantini (1989))

KF is proof-theoretically equivalent to $(\Pi_1^{\circ}-CA)_{\leq \varepsilon_0}$.

Proof Idea.

Upper bound: One formulates KF in a Tait (one-sided) infinitary calculus, and analyzes quai-normal derivations, i.e. derivations with only cuts on Tr t and \neg Tr t and proves in CT_{< ϵ} that

if
$$KF^{\infty} \vdash^{\alpha} Tr \lceil A \rceil$$
, then $Tr_{2^{\alpha}} \lceil A \rceil$

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Symmetries

• One important drawback of KF is that its internal theory $\{\varphi \in \mathcal{L}_{Tr} \mid KF \vdash \varphi\}$ is different from its theorems: for instance $KF \vdash \lambda \lor \neg \lambda$ but $KF \nvdash Tr \ulcorner \lambda \lor \neg \lambda \urcorner$.

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- ▶ To overcome this:

Reinhardt's thesis

One should adopt an **instrumental** reading of KF. Its conceptual core is given by its **internal theory**.

Lemma (Halbach and Horsten (2006))

There are A's in \mathcal{L}_{Tr} such that $KF \vdash Tr^rA^r$ but the proof essentially employs B's such that $KF \nvdash Tr^rB^r$.

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