

Olympiad Notebook

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Abstract

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1 Combinatorics

1.1 Bijections

1.2 Invariants and Monovariants

1.3 Pigeonhole Principle

Theorem 1.1 (Pigeonhole Principle). *Let m, n be positive integers with $m \geq n$. If $m + 1$ pigeons fly to n pigeonholes, then at least one pigeonhole contains at least $\lfloor \frac{m}{n} \rfloor + 1$ pigeons.*

1.4 Extremal Principle

1.5 Combinatorial Games

The main strategies for analyzing combinatorial games are:

- Play the game: try to find some forced moves.
- Reduce the game to a simpler game.
- Start at the end of the game: find endgame positions which are winning and losing and work backwards.
- Find an invariant or monovariant that a player can control.

Problem 1.2. Four heaps contain 38, 45, 61, and 70 matches respectively. Two players take turns choosing any two of the heaps and removing a non-zero number of matches from each heap. The player who cannot make a move loses. Which one of the players has a winning strategy?

Proof. Denote the heaps with a 4-tuple (w, x, y, z) with $w \leq x \leq y \leq z$. We claim the winning positions are of the form (w, x, y, z) with $w < y$. It is clear that $(0, 0, y, z)$ leads to a win by removing y and z and $(0, x, y, z)$ leads to a win by reducing to $(0, 1, 1, z)$ which is forced to leave either 1 or 2 heaps.

Since we remove tiles on each move, the game must terminate. If we have (w, x, y, z) with $w < y$, we can reduce to (w, w, w, x) by sending y and z to w .

We show that (w, w, w, z) is a losing position. We have three cases:

1. If we remove from two of the w -heaps, we are left with (w', w'', w, z) .
2. If we remove from a w -heap and the z -heap, we are left with either (w', z', w, w) or (w', w, z', w) or (w', w, w, z') .
3. If we remove any number of heaps entirely, the resulting position is clearly winning.

It follows that (w, x, y, z) with $w < y$ is a winning position as desired. □

Problem 1.3. The number 10^{2015} is written on a blackboard. Alice and Bob play a game where each player can do one of the following on each turn:

- replace an integer x on the board with integers $a, b > 1$ so that $x = ab$
- erase one or both of two equal integers on the blackboard.

The player who is not able to make a move loses the game. Who has a winning strategy?

Proof. We claim Alice has a winning strategy. First, it is clear that the game must eventually terminate. On the first turn, Alice can replace 10^{2015} with 2^{2015} and 5^{2015} . We claim that after any of Bob's turns, Alice can move the board into the state

$$2^{\alpha_1} 2^{\alpha_2} \dots 2^{\alpha_k} 5^{\alpha_1} 5^{\alpha_2} \dots 5^{\alpha_k}.$$

If Bob sends 2^{α_j} to $2^{\beta_1}, 2^{\beta_2}$, then Alice can send 5^{α_j} to $5^{\beta_1}, 5^{\beta_2}$ and vice versa. Otherwise, if Bob removes one or two integers $2^{\alpha_j}, 2^{\alpha_k}$, then we have $\alpha_j = \alpha_k$ so Alice can remove one or two of $5^{\alpha_j}, 5^{\alpha_k}$ or vice versa. Since Alice can always follow the copycat strategy and the game eventually terminates, we must have that Bob is unable to make a move at some point, which implies that Alice wins the game as desired. \square

Problem 1.4.

1.6 Algorithms

1.7 Generating Functions

Problem 1.5 (Putnam 2020 A2). Let k be a non-negative integer. Evaluate

$$\sum_{j=0}^k 2^{k-j} \binom{k+j}{j}.$$

Proof. We claim the sum evaluates to 4^k . Note that $\binom{k+j}{j} = \binom{k+j}{k}$. It follows that the sum is the coefficient of x^k in the power series $\sum_{j=0}^n 2^{k-j} (1+x)^{k+j}$. Evaluating this, we find

$$\begin{aligned} \sum_{j=0}^n 2^{k-j} (1+x)^{k+j} &= 2^k (1+x)^k \sum_{j=0}^k 2^{-j} (1+x)^j \\ &= 2^k (1+x)^k \frac{1 - (1+x)^{k+1}/2^{k+1}}{1 - (1+x)/2} \\ &= \frac{2^{k+1} (1+x)^k - (1+x)^{2k+1}}{1-x} \\ &= 2^{k+1} (1+x)^k - (1+x)^{2k+1} \sum_{n \geq 0} x^n. \end{aligned}$$

It follows that the coefficient of x^k is given by

$$2^{k+1} \sum_{j=0}^k \binom{k}{j} - \sum_{j=0}^k \binom{2k+1}{j} = 2^{2k+1} - 2^{2k} = 4^k.$$

\square

Problem 1.6. (CJMO 2020/1) Let N be a positive integer, and let S be the set of all tuples with positive integer elements and a sum of N . For all tuples t , let $p(t)$ denote the product of all the elements of t . Evaluate

$$\sum_{t \in S} p(t).$$

Proof. We claim the sum evaluates to F_{2N} , where F_k denotes the k -th Fibonacci number. Note that the sum can be represented as the coefficient of x^N in $\sum_{k=1}^N \left(\sum_{n \geq 0} nx^n \right)^k$. Evaluating this, we find

$$\begin{aligned} \sum_{k=1}^N \left(\sum_{n \geq 0} nx^n \right)^k &= \sum_{k=1}^N \left(\frac{x}{(1-x)^2} \right)^k \\ &= \sum_{k=1}^N \frac{x^k}{(1-x)^{2k}} \\ &= \sum_{k=1}^N \sum_{j \geq 0} \binom{2k-1+j}{2k-1} x^{j+k}. \end{aligned}$$

The coefficient of x^N is given by

$$\sum_{k=1}^N \binom{N+k-1}{2k-1} = \sum_{k=1}^N \binom{N+k-1}{N-k} = \sum_{j \geq 0} \binom{2N-1-j}{j} = F_{2N}.$$

□

Problem 1.7 (IMO 1995/6). Let p be an odd prime number. How many p -element subsets A of $\{1, 2, \dots, 2p\}$ are there, the sum of whose elements is divisible by p ?

Proof. Define $f(x, y) = \prod_{k=1}^{2p} (1 + x^k y)$. We wish to find the sum of the coefficients of terms of the form $x^{p\ell} y^p$. We do this by first considering f as a generating function in x using the root of unity filter associated to $\omega = e^{\frac{2\pi i}{p}}$. Then, we read off the coefficient of y^p to find the desired expression.

Note that for $1 \leq k \leq p-1$,

$$f(\omega^k, y) = \prod_{k=1}^{2p} (1 + \omega^k y) = \prod_{k=1}^p (1 + \omega^k y)^2 = (1 + y^p)^2.$$

It follows that

$$\begin{aligned} \frac{1}{p} \sum_{i=0}^{p-1} f(\omega^i, y) &= \frac{1}{p} \left((1 + y)^{2p} + \sum_{i=1}^{p-1} f(\omega^i, y) \right) \\ &= \frac{(1 + y)^{2p} + (p-1)(1 + y^p)^2}{p}. \end{aligned}$$

Finally, the coefficient of y^p is given by

$$\frac{\binom{2p}{p} + 2(p-1)}{2}.$$

□

1.8 Enumerative Combinatorics

1.9 Probabilistic Method

1.10 Algebraic Combinatorics

1.11 Combinatorial Geometry

1.11.1 Convex Hull

Problem 1.8 (Happy-Ending Problem). Suppose we have five points in the plane with no three collinear. Show that we can find four points whose convex hull is a quadrilateral.

Proof. Take the convex hull of the five points. If it is a quadrilateral or pentagon, we are done (choose any 4 points in the latter case). Suppose the convex hull is a triangle. Label the points with A through E and without loss of generality, let the points A, B, C form the triangle and D, E , be the points inside the hull.

Extend the line DE . Note that two points must lie on one side of the line - if not then we have three collinear points. It is easy to show that these four points form a convex quadrilateral. \square

Problem 1.9. There are $n > 3$ coplanar points, no three collinear and every four of them are the vertices of a convex quadrilateral. Prove that the n points are the vertices of a convex n -sided polygon.

Proof. Suppose that some point P is inside the convex hull of the n points. Let Q be some vertex of the convex hull. The diagonals from Q to the other vertices divide the convex hull into triangles and since no three points are collinear, P must lie inside some triangle $\triangle QRS$. But this is a contradiction since P, Q, R, S do not form a convex quadrilateral. \square

Problem 1.10 (1985 IMO Longlist). Let A, B be finite disjoint sets of points in the plane such that any three distinct points in $A \cup B$ are not collinear. Assume that at least one of the sets A, B contains at least five points. Show that there exists a triangle all of whose vertices are contained in A or in B that does not contain in its interior any point from the other set.

Proof. Suppose A has at least five points. Take $A_1 A_2$ on the boundary of the convex hull of A . For any other $A_i \in A$, define $\theta_i = \angle A_1 A_2 A_i$. Without loss of generality, $\theta_3 < \theta_4 < \dots < 180^\circ$. It follows that $\text{conv}(\{A_1, A_2, A_3, A_4, A_5\})$ contains no other points of A . \square

Problem 1.11 (Putnam 2001 B6). Assume that $(a_n)_{n \geq 1}$ is an increasing sequence of positive real numbers such that $\lim_{n \rightarrow \infty} \frac{a_n}{n} = 0$. Must there exist infinitely many positive integers n such that

$$a_{n-i} + a_{n+i} < 2a_n$$

for $i = 1, \dots, n-1$?

Proof. We claim such a subsequence exists. Let $A = \text{conv}\{(n, a_n) : n \in \mathbb{N}\}$ and let ∂A denote the set of points on the boundary of the convex hull.

We claim that ∂A contains infinitely many elements. Suppose not. Then, ∂A has a last point (N, a_N) . If we let $m = \sup_{n > N} \frac{a_n - a_N}{n - N}$, the slope of the line between (N, a_N) and (n, a_n) , then the line through (N, a_N) with slope m lies above (or contains) each point (n, a_n) for $n > N$. However, since $a_n/n \rightarrow 0$ and a_N, N are fixed, we have that

$$\frac{a_n - a_N}{n - N} \rightarrow 0.$$

This implies that the set of slopes attains a maximum, i. e. there is some point (M, a_M) with $M > N$ so that $m = \frac{a_M - a_N}{M - N}$. But then, we must also have that $(M, a_M) \in \partial A$, contradicting the fact that (N, a_N) is the last point in ∂A .

For each point on the boundary $(n, a_n) \in \partial A$, we must have that midpoint of the line through $(n - i, a_{n-i})$ and $(n + i, a_{n+i})$ for $i \in [n - 1]$ must lie below (n, a_n) . From this, it follows that $a_n > \frac{a_{n-i} + a_{n+i}}{2}$, which implies the result. \square

2 Algebra

2.1 Linear Algebra

Problem 2.1. Let $A \in M_n(\mathbb{R})$ be skew-symmetric. Show that $\det(A) \geq 0$.

Proof. If n is odd, note that

$$\det(A) = \det(A^T) = \det(-A) = (-1)^n \det(A) = -\det(A).$$

It follows that $\det(A) = 0$.

Otherwise, suppose n is even and let $p(\lambda) = \det(A - I_n \lambda)$. If $\lambda \neq 0$ is an eigenvalue, note that $p(\lambda) = 0$ by the Cayley-Hamilton Theorem. Moreover,

$$p(-\lambda) = \det(A + I_n \lambda) = \det(A^T + I_n^T \lambda) = \det(-A + I_n \lambda) = 0.$$

Moreover, let v be an eigenvector with corresponding eigenvalue λ . Note that

$$\langle Av, v \rangle = \lambda \langle v, v \rangle = \lambda \|v\|^2,$$

$$\langle Av, v \rangle = \langle v, A^T v \rangle = \langle v, -Av \rangle = -\bar{\lambda} \langle v, v \rangle = -\bar{\lambda} \|v\|^2.$$

It follows that $\lambda = -\bar{\lambda}$, which implies that $\lambda = ri$ for $r \in \mathbb{R}$. Hence,

$$\det(A) = \prod_{j=1}^{n/2} (i\lambda_j)(-i\lambda_j) = \prod_{j=1}^n \lambda_j^2 \geq 0.$$

□

Problem 2.2. Let $A \in M_n(\mathbb{R})$ with $A^3 = A + I_n$. Show that $\det(A) > 0$.

Proof. Let $p(x) = x^3 - x - 1$. Note that $p(0) = -1$, $p(2) = 5$, so the polynomial has a root in the interval $(0, 2)$ by the intermediate value theorem. Furthermore, $p'(x) = 3x^2 - 1$ so the polynomial has critical points at $\pm \frac{1}{\sqrt{3}}$. It is easy to see that at both of these values, $p(x) < 0$ so it follows that the other roots of $p(x)$ are conjugate complex numbers. Let the roots be $\lambda_1, \lambda_2, \lambda_3$ with λ_1 being the positive real root and λ_2, λ_3 the conjugate complex ones. If A satisfies $A^3 = A + I_n$, then we must have the eigenvalues of A are λ_1, λ_2 and λ_3 , with multiplicity $\alpha_1, \alpha_2, \alpha_3$ respectively. Since λ_2, λ_3 are complex conjugates, we must have $\alpha_2 = \alpha_3$, so it follows that

$$\det(A) = \lambda_1^{\alpha_1} (\lambda_2 \lambda_3)^{\alpha_2} = \lambda_1^{\alpha_1} |\lambda_2|^{\alpha_2} > 0.$$

□

Problem 2.3. If $A, B \in M_n(\mathbb{R})$ such that $AB = BA$, then $\det(A^2 + B^2) \geq 0$.

Proof.

$$\det(A^2 + B^2) = \det(A + iB) \det(A - iB) = \det(A + iB) \overline{\det(A + iB)} = |\det(A + iB)|^2 \geq 0.$$

□

Problem 2.4. Let $A, B \in M_2(\mathbb{R})$ such that $AB = BA$ and $\det(A^2 + B^2) = 0$. Show that $\det(A) = \det(B)$.

Proof. Let $p_{A,B}(\lambda) = \det(A + \lambda B) = \det(B)\lambda^2 + (\operatorname{tr} A + \operatorname{tr} B - \operatorname{tr}(AB))\lambda + \det(A)$. By Problem 1.3, we have $\det(A + iB)$ and $\det(A - iB) = 0$, which implies that $p_{A,B}(\lambda) = c(\lambda - i)(\lambda + i) = c(\lambda^2 + 1)$. It follows that $c = \det B = \det A$. \square

Problem 2.5. Let $A \in M_2(\mathbb{R})$ with $\det A = -1$. Show that $\det(A^2 + I_2) \geq 4$. When does equality hold?

Proof. First, note the identity

$$\det(X + Y) + \det(X - Y) = 2(\det X + \det Y).$$

This follows from writing $p(z) = \det(X + zY) = \det(Y)z^2 + (\operatorname{tr} X + \operatorname{tr} Y - \operatorname{tr}(XY))z + \det(X)$ and taking

$$p(1) + p(-1) = \det(X + Y) + \det(X - Y) = 2\det Y + 2\det X.$$

Then, taking $X = A^2 + I$ and $Y = 2A$, we have

$$0 \leq \det(A + I)^2 + \det(A - I)^2 = 2(\det(A^2 + I) + \det(2A)) = 2(\det(A^2 + I) - 4).$$

It follows that $\det(A^2 + I) \geq 4$ as desired. We have equality when the eigenvalues of A are 1 and -1 . \square

Problem 2.6. Let $A, B \in M_3(\mathbb{C})$ with $\det(A) = \det(B) = 1$. Show that $\det(A + \sqrt{2}B) \neq 0$.

Proof. \square