

IP Unicast Routing

Redes de Comunicações II

**Licenciatura em
Engenharia de Computadores e Informática
DETI-UA**



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IP Routing Overview

- Routers forward packets toward destination networks.
- Routers must be aware of destination networks to be able to forward packets to them.
- A router knows about the networks directly attached to its interfaces
- For networks not directly connected to one of its interfaces, however, the router must rely on outside information.
- A router can be made aware of remote networks by:
 - ◆ **Static routing:** An administrator manually configure the information.
 - ◆ **Dynamic routing:** Learns from other routers.
 - ◆ **Policy based routing:** Manually routing rules that outweigh static/dynamic routing and may depend on parameters other than the destination address.



Default Routes

- In some circumstances, a router does not need to recognize the details of remote networks.
- The router can be configured to send all traffic (or all traffic for which there is not a more specific entry in the routing table) to a specific neighbor router.
- This is known as a default route.
- Default routes are either dynamically advertised using routing protocols or statically configured.
- IPv4 default route - 0.0.0.0/0
- IPv6 default route - ::/0

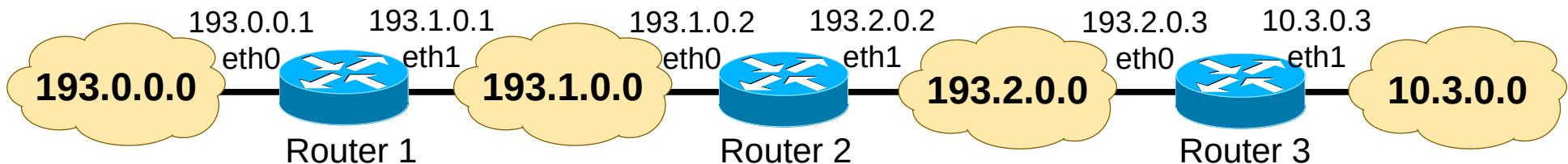


Static Routing

- Stating routing do not react to network topology changes.
 - ◆ If a link fails, the static route is no longer valid if it is configured to use that failed link, so a new static route must be configured.
 - ◆ Connectivity may be lost until intervention of an administrator.
- Static routing does not scale well when network grows.
 - ◆ Administrative burden to maintain routes may can become excessive.
- Static routes can be used in the following circumstances:
 - ◆ When the administrator needs total control over the routes used by the router.
 - ◆ When a backup to a dynamically recognized route is necessary.
 - ◆ When it is used to reach a network accessible by only one path (a stub network).
 - There is no backup link, so dynamic routing has no advantage.
 - ◆ When a router connects to its ISP and needs to have only a default route pointing toward the ISP router, rather than learning many routes from the ISP.
 - Again, a single path of access without backup.
 - ◆ When a router is underpowered and does not have the CPU or memory resources necessary to handle a dynamic routing protocol.
 - ◆ When it is undesirable to have dynamic routing updates forwarded across low bandwidth links.



Static Routing Examples



- Example 1

- Router2 do not know networks 193.0.0.0/24 and 10.3.0.0/24
- Necessary static routes:
 - 193.0.0.0/24 accessible through 193.1.0.1 (eth1, Router1)
 - 10.3.0.0/24 accessible through 193.2.0.3 (eth0, Router3)

- Example 2

- Router1 do not know networks 193.2.0.0/24 and 10.3.0.0/24
- Necessary static routes:
 - 193.2.0.0/24 accessible through 193.1.0.2 (eth0, Router2)
 - 10.3.0.0/24 accessible through 193.1.0.2 (eth0, Router2)
- OR
- Using default route: 0.0.0.0/0 accessible through 193.1.0.2 (eth0, Router2)



Dynamic Routing

- Dynamic routing allows the network to adjust to changes in the topology automatically, without administrator involvement.
- Routers exchange information about the reachable networks and the state of each network/link.
 - ◆ Routers exchange information only with other routers running the same routing protocol.
 - ◆ When the network topology changes, the new information is dynamically propagated throughout the network, and each router updates its routing table to reflect the changes.



(Complex) Routing Tables

- An IP address may have multiple matches on a Routing Table:
 - ◆ Example: 192.168.1.12
 - ◆ Will match:
 - 192.168.1.0/25 via ...
 - 192.168.1.0/24 via ...
 - 192.168.0.0/23 via ...
 - 192.168.0.0/16 via ...
 - ...
 - ◆ Router will choose entry with the largest network prefix (most specific network).
 - i.e., 192.168.1.0/25 via ...
- Load balancing
 - ◆ Routing tables may have more than one path for each network
 - ◆ Traffic will be divided by all entries.
 - ◆ By packet, flow (TCP session, UDP IPs/port), etc...
 - E.g, packet 1 path 1, packet 2 path 2, packet 3 path 1, ...
 - Flow 1 path 1, flow 2 path 2, flow 3 path 3, flow 4 path 1, flow 5 path 2, ...



Administrative Distance

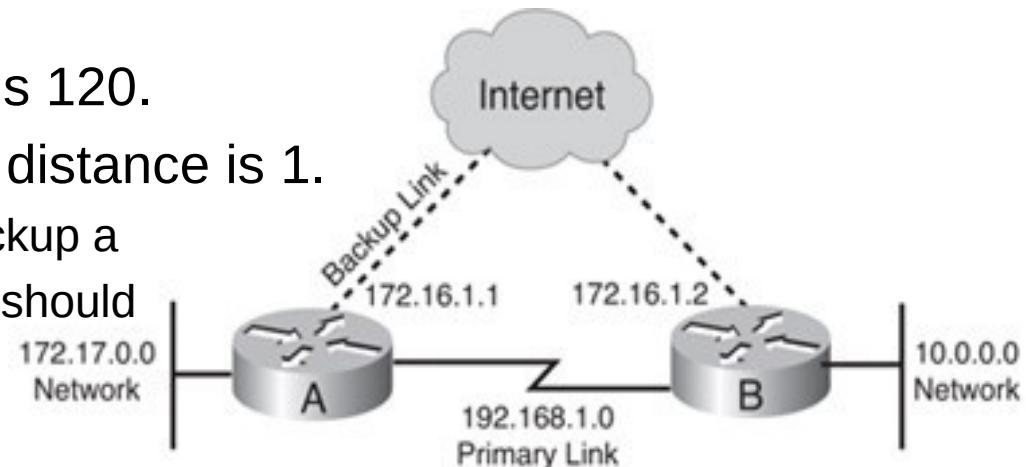
- Most routing protocols have metric structures and algorithms that are incompatible with other protocols.
- It is critical that a network using multiple routing protocols be able to seamlessly exchange route information and be able to select the best path across multiple protocols.
- Routers use a value called administrative distance to select the best path when they learn from different routing protocols the same destination (same network prefix and mask length).
- The Protocol/Method with the lowest Administrative Distance is preferred
 - ◆ The Administrative Distance value is configurable.
- Example:
 - ◆ Static [1/1] 192.168.1.0/24 via ... ← Chosen!
 - ◆ RIP [120/1] 192.168.1.0/24 via ...
 - ◆ OSPF [110/1] 192.168.1.0/24 via ...



Floating Static Routes

Calculam o melhor caminho como nos Spanning Trees usando o Teorema de Bellman. Faz-se uma espécie de lista dos melhores caminhos

- Based on the default administrative distances, routers use static routes over any dynamically learned route.
 - ◆ However, this default behavior might not be the desired behavior.
 - ◆ For example, when you configure a static route as a backup to a dynamically learned route, you do not want the static route to be used as long as the dynamic route is available.
- A static route that appears in the routing table only when the primary route goes away is called a floating static route.
 - ◆ The administrative distance of the static route is configured to be higher than the administrative distance of the primary route and it “floats” above the primary route, until the primary route is no longer available.
- RIP default administrative distance is 120.
- Static Routes default administrative distance is 1.
 - ◆ To create a floating static route (to backup a RIP route) the administrative distance should be greater than 120.
 - ◆ In example: 200.



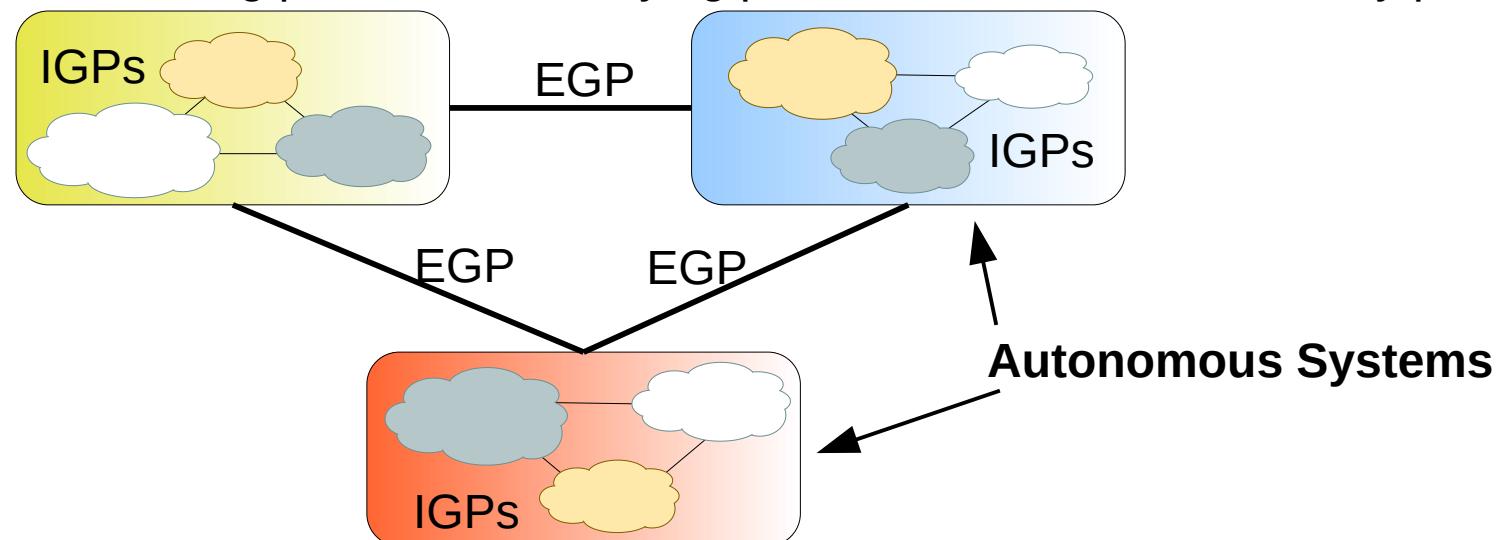
Analisa-se a informação TODA

```
ip route 10.0.0.0 255.0.0.0 172.16.1.2 200
router rip
  network 192.168.1.0
  network 172.17.0.0
```

```
ip route 172.17.0.0 255.255.0.0 172.16.1.1 200
router rip
  network 192.168.1.0
  network 10.0.0.0
```

Autonomous Systems

- AS (Autonomous System) – set of routers/networks with a common routing policy and under the same administration.
- Routing inside an AS is performed by IGPs (Interior Gateway Protocols) such as RIPv1, RIPv2, OSPF, IS-IS and EIGRP.
 - ◆ Called Internal Routing
- Routing between AS is performed by EGPs (Exterior Gateway Protocols) such as BGP.
- IGPs and EGPs have different objectives:
 - ◆ IGPs: optimize routing performance
 - ◆ EGPs: optimize routing performance obeying political, economic and security policies.



Type of Networks

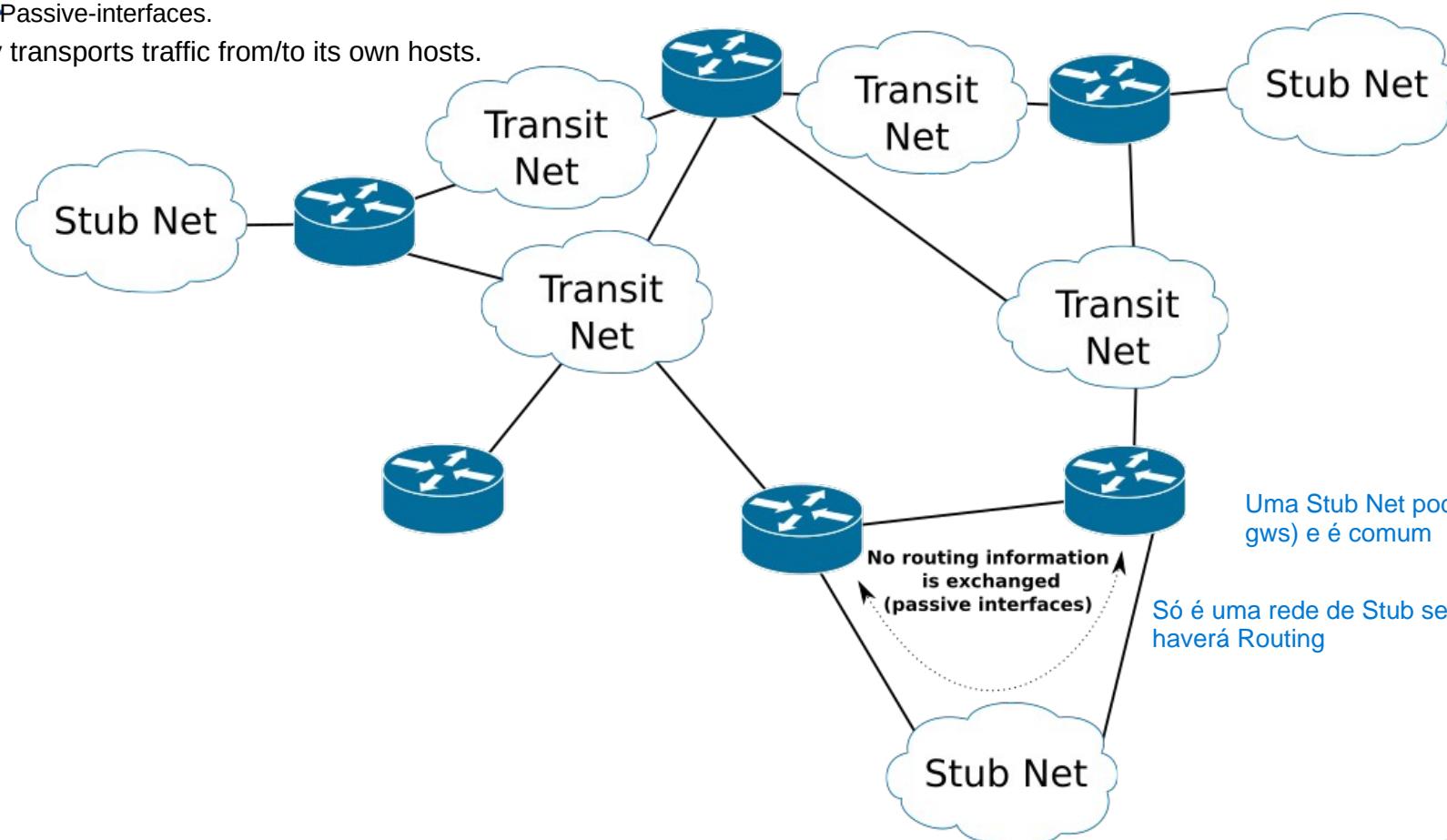
● Transit/Transport

- ◆ Used to interconnect networks.
 - Routers exchange routing information using it.
- ◆ Transports traffic from/to other network hosts and from/to its own hosts.

Transit Net - pode ser um fio ou uma nuvem

● Stub

- ◆ Single router network.
- ◆ or multiple routers network, if routers do not exchange routing information.
 - Passive-interfaces.
- ◆ Only transports traffic from/to its own hosts.



Distance Vector versus Link State Protocols

- Distance vector
 - ◆ Each routers learns networks and best path based on the information sent periodically by its neighbors.
 - Network and cost (distance) to that network.
 - ◆ Each router determines the shortest paths to all know networks based on a distributed and asynchronous version of the Bellman-Ford algorithm.
 - ◆ Examples: RIPv1, RIPv2, IGRP, EIGRP.
- Link state
 - ◆ Routers learn the complete network topology and use a centralized algorithm to determine the shortest paths to all known networks.
 - ◆ The information necessary to construct and maintain in each router a data base with the network topology is obtain by a flooding process.
 - ◆ Network information is only exchanged on bootstrap and after any topology change.
 - ◆ Examples: OSPF, IS-IS.

Mais comum para Redes empresariais

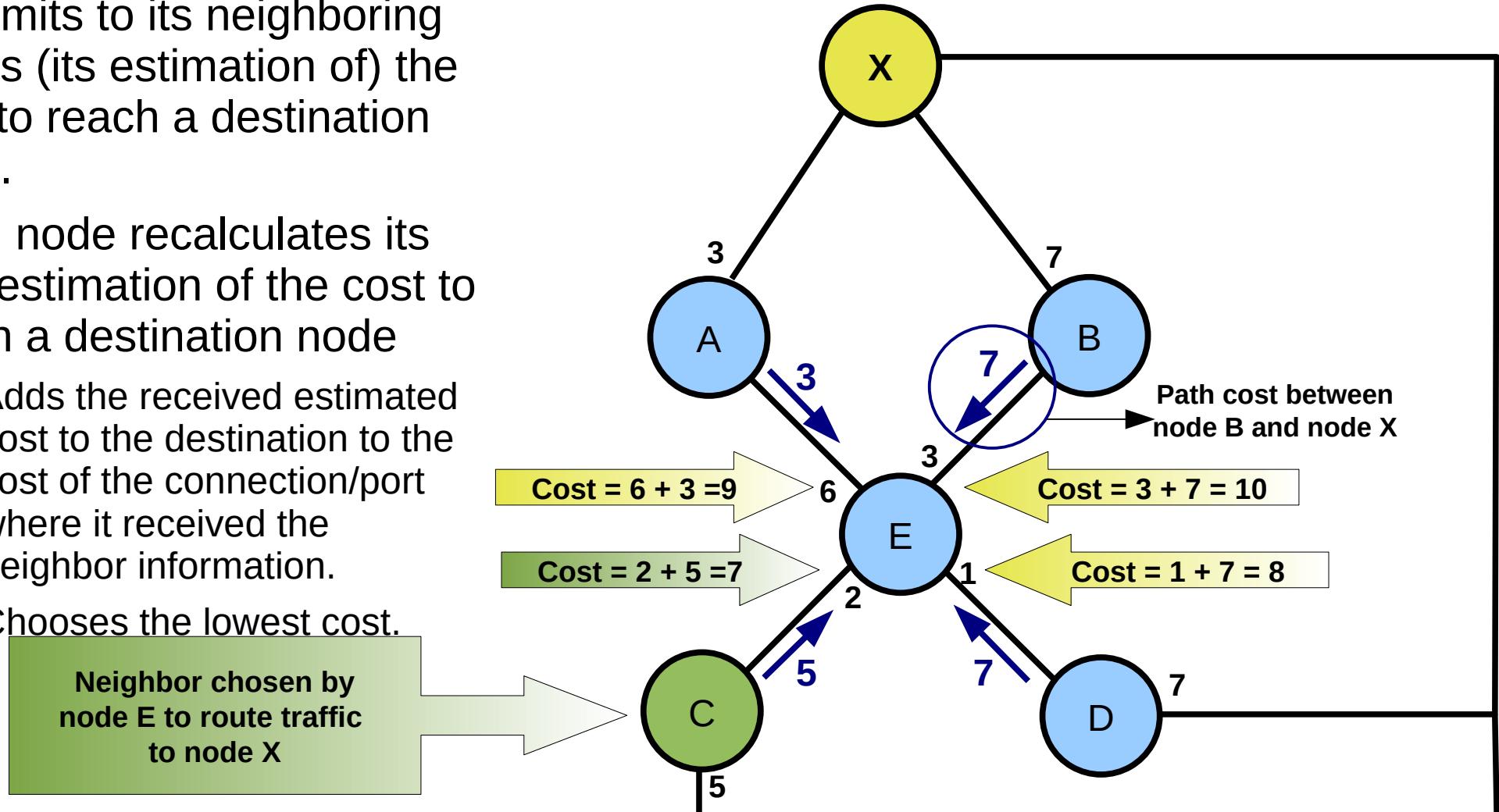
Para grandes operadores, com mais dinheiro -> Mais flexível, mais escalável, funciona melhor em redes grandes



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Distributed and Asynchronous Bellman-Ford Algorithm

- Each node periodically transmits to its neighboring nodes (its estimation of) the cost to reach a destination node.
- Each node recalculates its own estimation of the cost to reach a destination node
 - ◆ Adds the received estimated cost to the destination to the cost of the connection/port where it received the neighbor information.
 - ◆ Chooses the lowest cost.



RIP (Routing Information Protocol)

- Is a *distance vector* protocol
 - ◆ Each router maintains a list of known networks and, for each network, an estimation of the cost to reach it – this is called a distance vector.
 - ◆ Each router periodically send to its neighboring routers its own distance vector (partially or complete) – announcement/update.
 - ◆ Each router uses the distance vector sent by its neighbors to update its own distance vector.
- The path cost to a destination is given by the number of routers/hops in the path.
 - ◆ Maximum cost is 15.
 - ◆ A cost of 16 is considered infinite (or unattainable destination).
- Each router determines the entries in its own routing table, based on the constructed distance vector.
 - ◆ For each destination (network) learned, it adds an entry to that network that uses the path (or paths) with the lowest cost, using as next-hop the neighboring router(s) that announced that network with that lowest cost path.



RIP Version 1

- RIP Version 1 (RIPv1) is a classfull protocol.
 - ◆ Does not announces (sub-)networks masks, only network prefixes.
 - ◆ Network masks are assumed based on the incoming interface mask.
 - ◆ If all networks have the same mask it works perfectly, however, when networks with different masks exist it is problematic.
- RIPv1 uses the broadcast address 255.255.255.255 to send announcements/updates.
 - ◆ All network devices must process the packets.
- Does not support authentication.
 - ◆ Messages may be forged by an attacker.



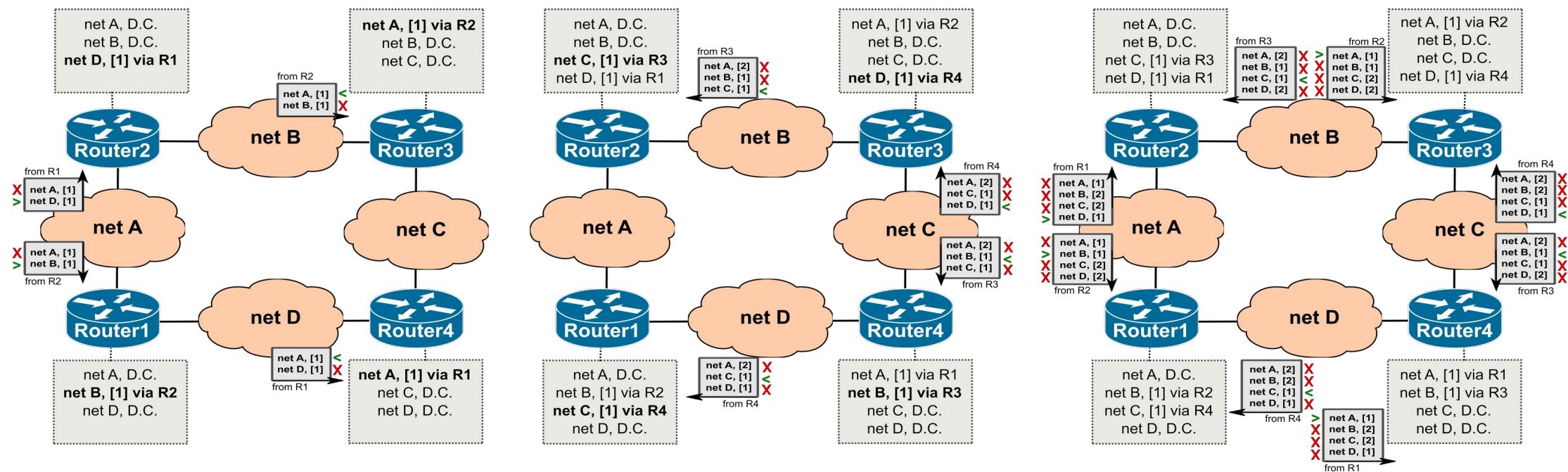
RIP Version 2

- RIP Version 2 (RIPv2) is a *classless* protocol.
 - ◆ RIPv2 announcements include network prefix and mask.
 - ◆ Supports variable length masks.
- RIPv2 used the multicast address 224.0.0.9 to send announcements/updates only to routers running RIPv2.
- RIPv2 supports authentication using message-digest and clear text password.
 - ◆ Clear text password authentication should not be used!



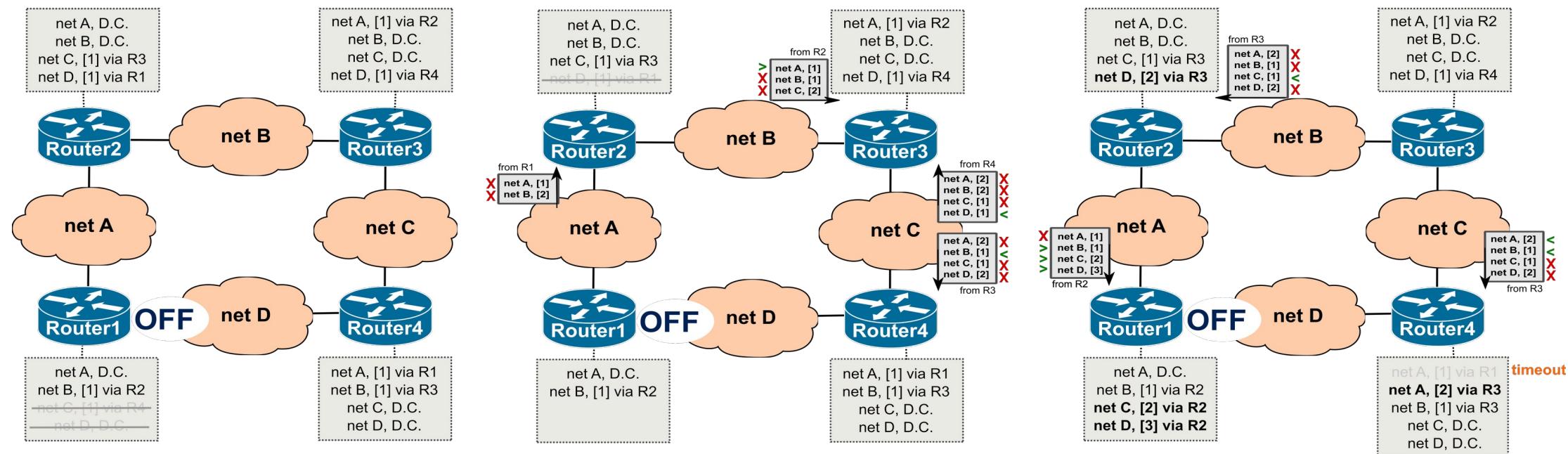
RIP Algorithm (1)

- Assuming that Router1 and Router2 send announcements first.
 - With split-horizon disabled.



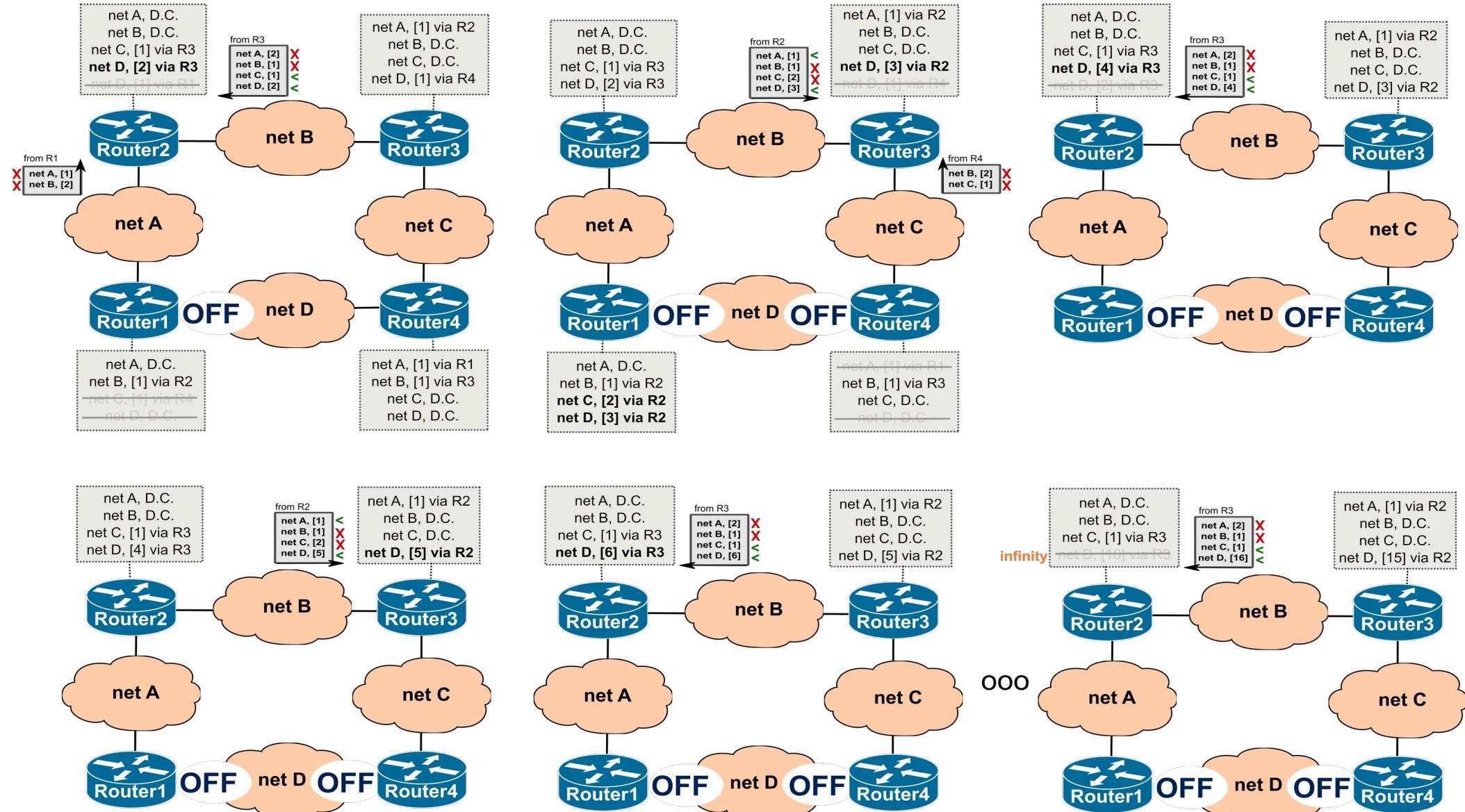
RIP Algorithm (2)

- Assuming Router1 connection to network D goes down.
 - No triggered updates.



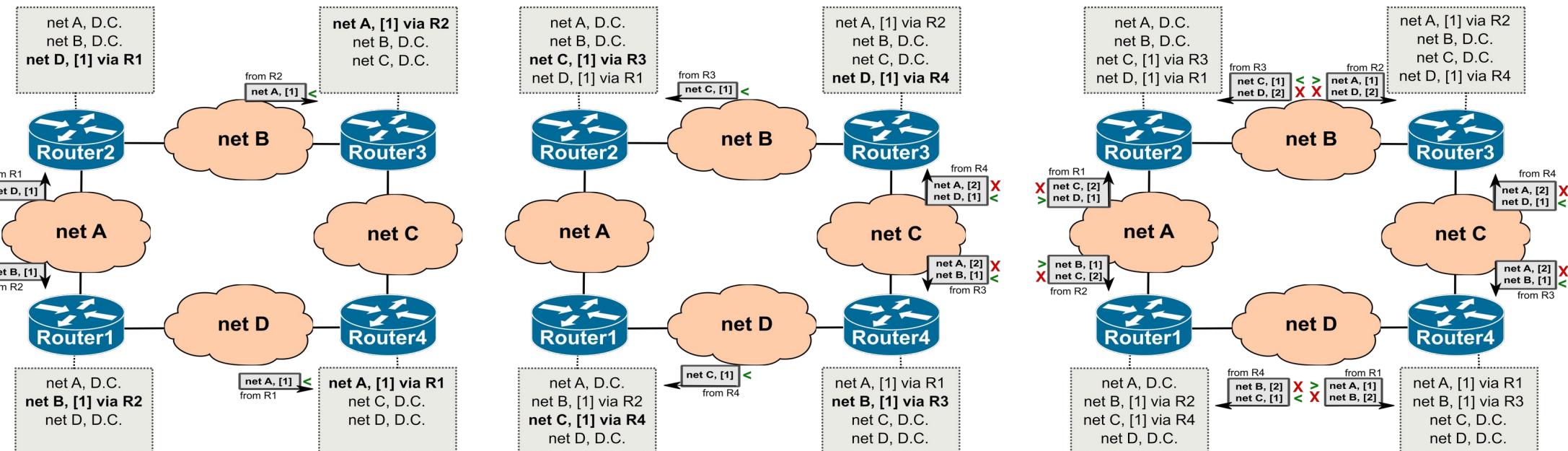
Count to Infinity Problem

- When multiple failures occur before algorithm convergence!



Split-Horizon (1)

- Solution for the count to infinity problem.
- Each Router, in each interface, announces only the networks in which that interface is not used to provide the best path to that destination.
- Split horizon lowers the convergence time of the routing tables when there is a topology change.
 - ◆ RIPv1 e RIPv2 supports it.
- Assuming Router1 and Router2 start sending announcements first:

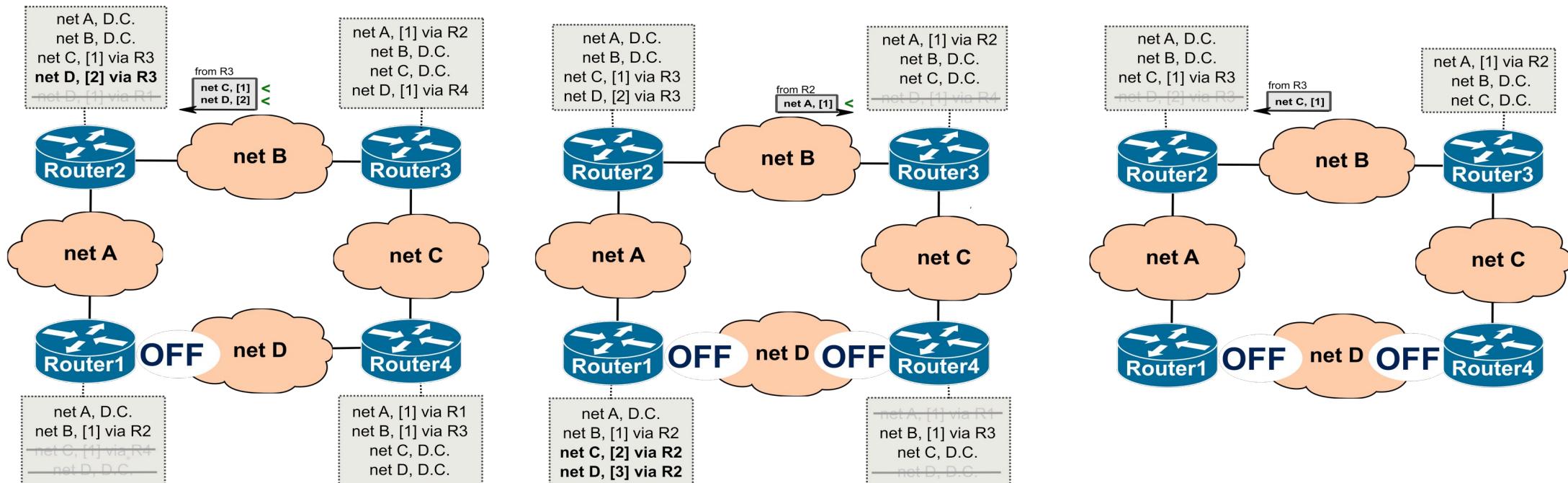


- In Split horizon with Poisoned Reverse, routers announce all networks but set metric to infinity (16) for networks learned by the interface by which they are sending the announcement.
 - ◆ Larger update messages.



Split-Horizon (2)

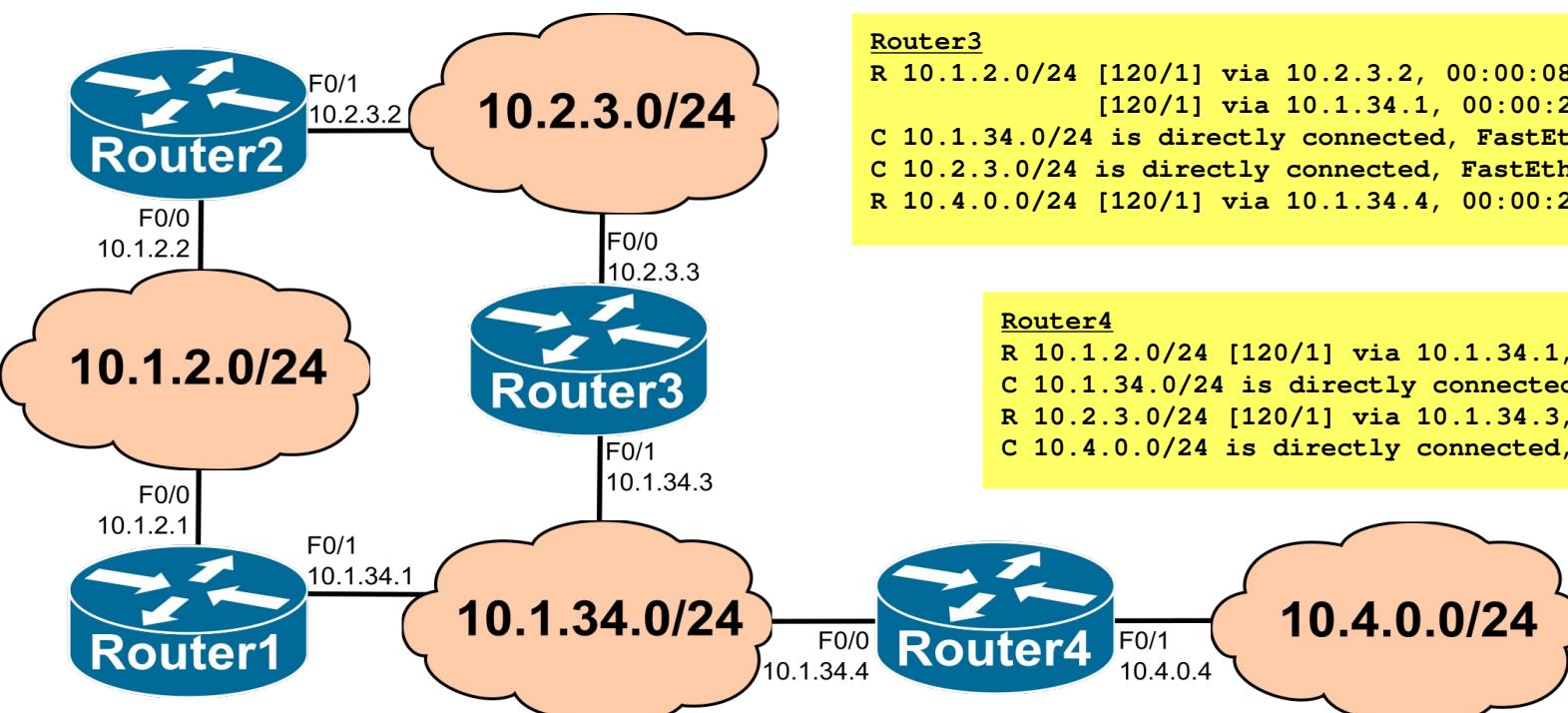
- Solution for the count to infinity problem.
- Prevents any routing loops that involve two routers.
 - ◆ It is possible to end up with patterns in which three or more routers are engaged in mutual deception.
- Assuming Router1 and Router4 loose connection to network D almost simultaneously:



Routing Tables with RIP

Router2

```
C 10.1.2.0/24 is directly connected, FastEthernet0/0
R 10.1.34.0/24 [120/1] via 10.2.3.3, 00:00:21, FastEthernet0/1
          [120/1] via 10.1.2.1, 00:00:11, FastEthernet0/0
C 10.2.3.0/24 is directly connected, FastEthernet0/1
R 10.4.0.0/24 [120/2] via 10.2.3.3, 00:00:21, FastEthernet0/1
          [120/2] via 10.1.2.1, 00:00:11, FastEthernet0/0
```



Router3

```
R 10.1.2.0/24 [120/1] via 10.2.3.2, 00:00:08, FastEthernet0/0
          [120/1] via 10.1.34.1, 00:00:28, FastEthernet0/1
C 10.1.34.0/24 is directly connected, FastEthernet0/1
C 10.2.3.0/24 is directly connected, FastEthernet0/0
R 10.4.0.0/24 [120/1] via 10.1.34.4, 00:00:24, FastEthernet0/1
```

Router4

```
R 10.1.2.0/24 [120/1] via 10.1.34.1, 00:00:18, FastEthernet0/0
C 10.1.34.0/24 is directly connected, FastEthernet0/0
R 10.2.3.0/24 [120/1] via 10.1.34.3, 00:00:29, FastEthernet0/0
C 10.4.0.0/24 is directly connected, FastEthernet0/1
```

Router1

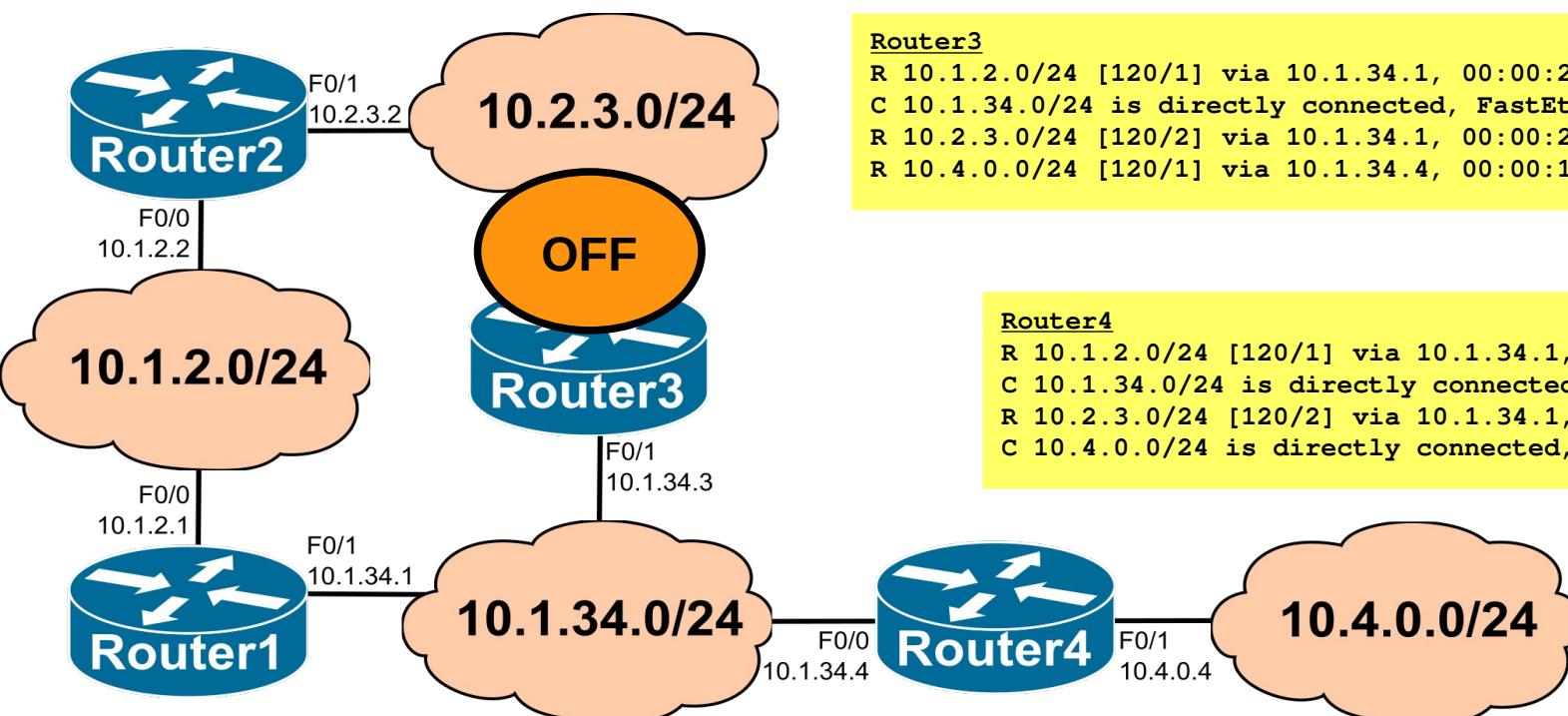
```
C 10.1.2.0/24 is directly connected, FastEthernet0/0
C 10.1.34.0/24 is directly connected, FastEthernet0/1
R 10.2.3.0/24 [120/1] via 10.1.34.3, 00:00:11, FastEthernet0/1
          [120/1] via 10.1.2.2, 00:00:01, FastEthernet0/0
R 10.4.0.0/24 [120/1] via 10.1.34.4, 00:00:24, FastEthernet0/1
```



Routing Tables with RIP

Router2 (AFTER 3 minutes TIMEOUT)

```
C 10.1.2.0/24 is directly connected, FastEthernet0/0
R 10.1.34.0/24 [120/1] via 10.1.2.1, 00:00:25, FastEthernet0/0
C 10.2.3.0/24 is directly connected, FastEthernet0/1
R 10.4.0.0/24 [120/2] via 10.1.2.1, 00:00:25, FastEthernet0/0
```



Router3

```
R 10.1.2.0/24 [120/1] via 10.1.34.1, 00:00:22, FastEthernet0/1
C 10.1.34.0/24 is directly connected, FastEthernet0/1
R 10.2.3.0/24 [120/2] via 10.1.34.1, 00:00:22, FastEthernet0/1
R 10.4.0.0/24 [120/1] via 10.1.34.4, 00:00:19, FastEthernet0/11
```

Router4

```
R 10.1.2.0/24 [120/1] via 10.1.34.1, 00:00:18, FastEthernet0/0
C 10.1.34.0/24 is directly connected, FastEthernet0/0
R 10.2.3.0/24 [120/2] via 10.1.34.1, 00:00:29, FastEthernet0/0
C 10.4.0.0/24 is directly connected, FastEthernet0/1
```

Router1

```
C 10.1.2.0/24 is directly connected, FastEthernet0/0
C 10.1.34.0/24 is directly connected, FastEthernet0/1
R 10.2.3.0/24 [120/1] via 10.1.2.2, 00:00:01, FastEthernet0/0
R 10.4.0.0/24 [120/1] via 10.1.34.4, 00:00:24, FastEthernet0/1
```



RIP Message Types

- RIP Response

- RIP Response
 - ◆ Distance vector announcement/update message.
 - ◆ Contains the distance vector.
 - ◆ It is sent:
 - ◆ 1 – Periodically (~30 seconds by default, there is a random component).
 - ◆ 2 – Optionally, when some information changes (triggered updates).
 - ◆ 3 – In response to a RIP Request.
 - ◆ In cases 1 and 2:
 - In RIPv1, is sent to the broadcast address.
 - In RIPv2, is sent to the multicast address 224.0.0.9 (Routers com RIP).
 - ◆ In case 3, it is sent only (unicast) to the router that sent the RIP Request.

- RIP Request (Optional)

- RIP Request (Optional)
 - ◆ Sent by a router that was recently started (bootstrap) or, when the validity of some of the distance vector information has expired (default timeout = 180 seconds)
 - ◆ It may request specific information (a specific network) or, the complete neighbor distance vector.



RIPv1 vs. RIPv2 Responses (1)

- New RIPv2 message fields in Response packets:
 - ◆ Subnet mask
 - ✚ Supports variable length masks.
 - ✚ Makes RIPv2 *classless* protocol.
 - ◆ Route tag
 - ✚ Attribute assigned to a specific network that must be reserved a re-announced.
 - ✚ Provides a method to separate internal (to the RIP domain) and external networks.
 - ◆ Next hop
 - ✚ Address to which the packets must be routed.
 - ✚ 0.0.0.0 indicates that the packets must be routed to the router that sent the RIP message.

RIPv1

| Bit | 0 | 7 | 8 | 15 | 16 | 23 | 24 | 31 |
|-------------|------------------|-------------|---|------------------------------------|---------------|----|--------------|---|
| | Command = 1 or 2 | Version = 1 | | | Must be zero | | | |
| Route Entry | | | | Address family identifier (2 = IP) | | | Must be zero | |
| | | | | | | | | |
| | | | | IP Address (Network Address) | | | | |
| | | | | | Must be zero | | | |
| | | | | | Must be zero | | | |
| | | | | | Metric (Hops) | | | |
| | | | | | | | | Multiple Route Entries, up to a maximum of 25 |

RIPv2

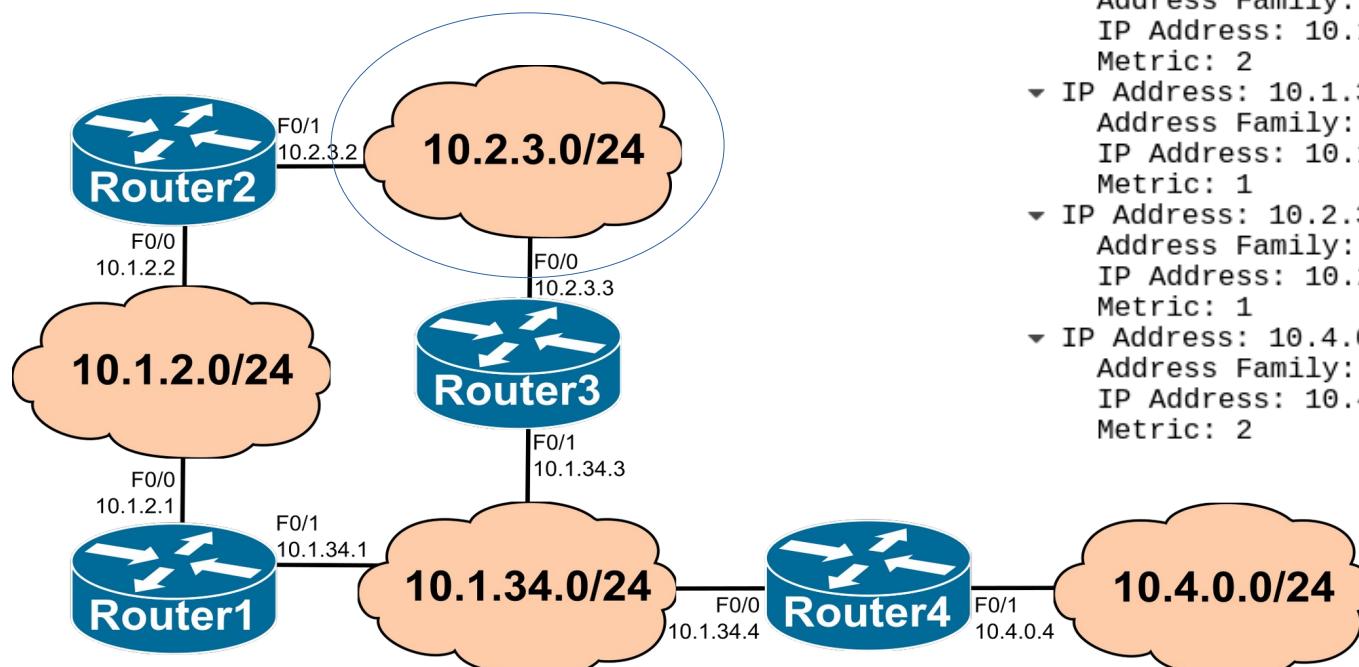
| Bit | 0 | 7 | 8 | 15 | 16 | 23 | 24 | 31 |
|-------------|------------------|-------------|---|------------------------------------|---------------|----|-----------|---|
| | Command = 1 or 2 | Version = 2 | | | Must be zero | | | |
| Route Entry | | | | Address family identifier (2 = IP) | | | Route Tag | |
| | | | | | | | | |
| | | | | IP Address (Network Address) | | | | |
| | | | | | Subnet Mask | | | |
| | | | | | Next Hop | | | |
| | | | | | Metric (Hops) | | | |
| | | | | | | | | Multiple Route Entries, up to a maximum of 25 |



RIPv1 Messages (Example)

Sent by Router3 with Split-Horizon

- ▶ Internet Protocol Version 4, Src: 10.2.3.3, Dst: 255.255.255.255
- ▶ User Datagram Protocol, Src Port: 520, Dst Port: 520
- ▶ Routing Information Protocol
 - Command: Response (2)
 - Version: RIPv1 (1)
 - IP Address: 10.1.34.0, Metric: 1
 - Address Family: IP (2)
 - IP Address: 10.1.34.0
 - Metric: 1
 - IP Address: 10.4.0.0, Metric: 2
 - Address Family: IP (2)
 - IP Address: 10.4.0.0
 - Metric: 2



Sent by Router3 without Split-Horizon

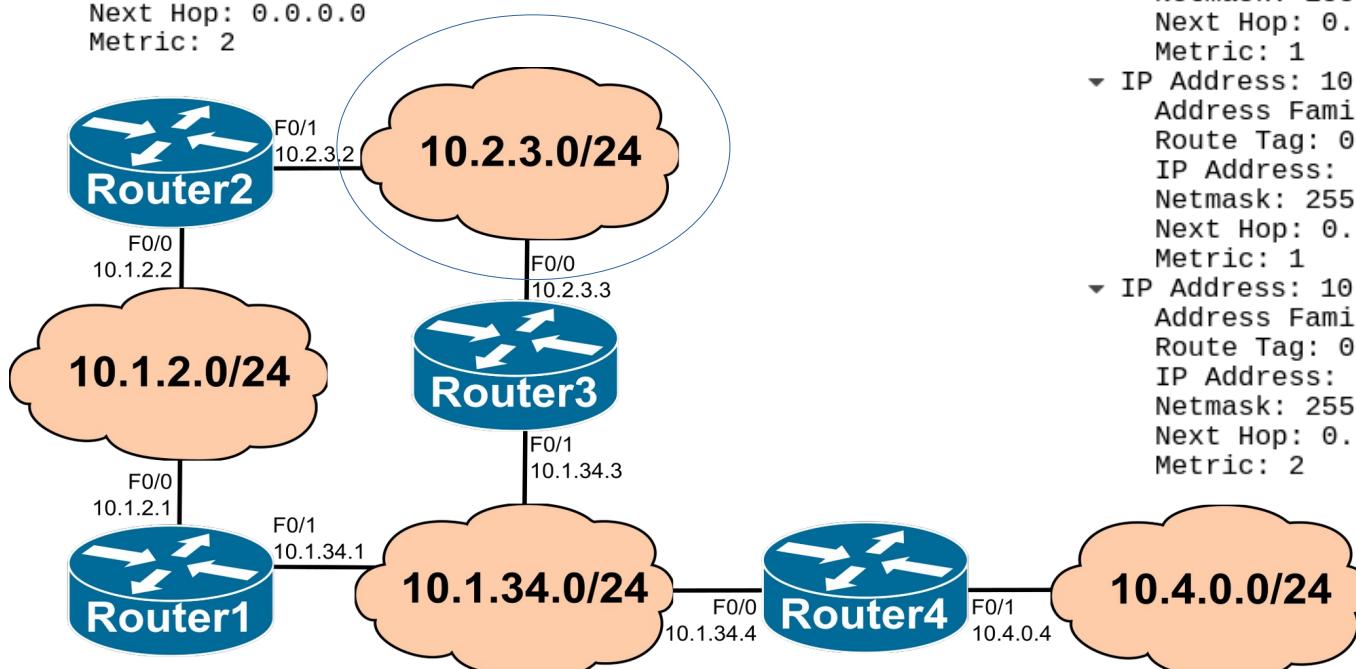
- ▶ Internet Protocol Version 4, Src: 10.2.3.3, Dst: 255.255.255.255
- ▶ User Datagram Protocol, Src Port: 520, Dst Port: 520
- ▶ Routing Information Protocol
 - Command: Response (2)
 - Version: RIPv1 (1)
 - IP Address: 10.1.2.0, Metric: 2
 - Address Family: IP (2)
 - IP Address: 10.1.2.0
 - Metric: 2
 - IP Address: 10.1.34.0, Metric: 1
 - Address Family: IP (2)
 - IP Address: 10.1.34.0
 - Metric: 1
 - IP Address: 10.2.3.0, Metric: 1
 - Address Family: IP (2)
 - IP Address: 10.2.3.0
 - Metric: 1
 - IP Address: 10.4.0.0, Metric: 2
 - Address Family: IP (2)
 - IP Address: 10.4.0.0
 - Metric: 2



RIPv2 Messages (Example)

Sent by Router3 with Split-Horizon

```
► Internet Protocol Version 4, Src: 10.2.3.3, Dst: 224.0.0.9
► User Datagram Protocol, Src Port: 520, Dst Port: 520
▼ Routing Information Protocol
    Command: Response (2)
    Version: RIPv2 (2)
    ▼ IP Address: 10.1.34.0, Metric: 1
        Address Family: IP (2)
        Route Tag: 0
        IP Address: 10.1.34.0
        Netmask: 255.255.255.0
        Next Hop: 0.0.0.0
        Metric: 1
    ▼ IP Address: 10.4.0.0, Metric: 2
        Address Family: IP (2)
        Route Tag: 0
        IP Address: 10.4.0.0
        Netmask: 255.255.255.0
        Next Hop: 0.0.0.0
        Metric: 2
```



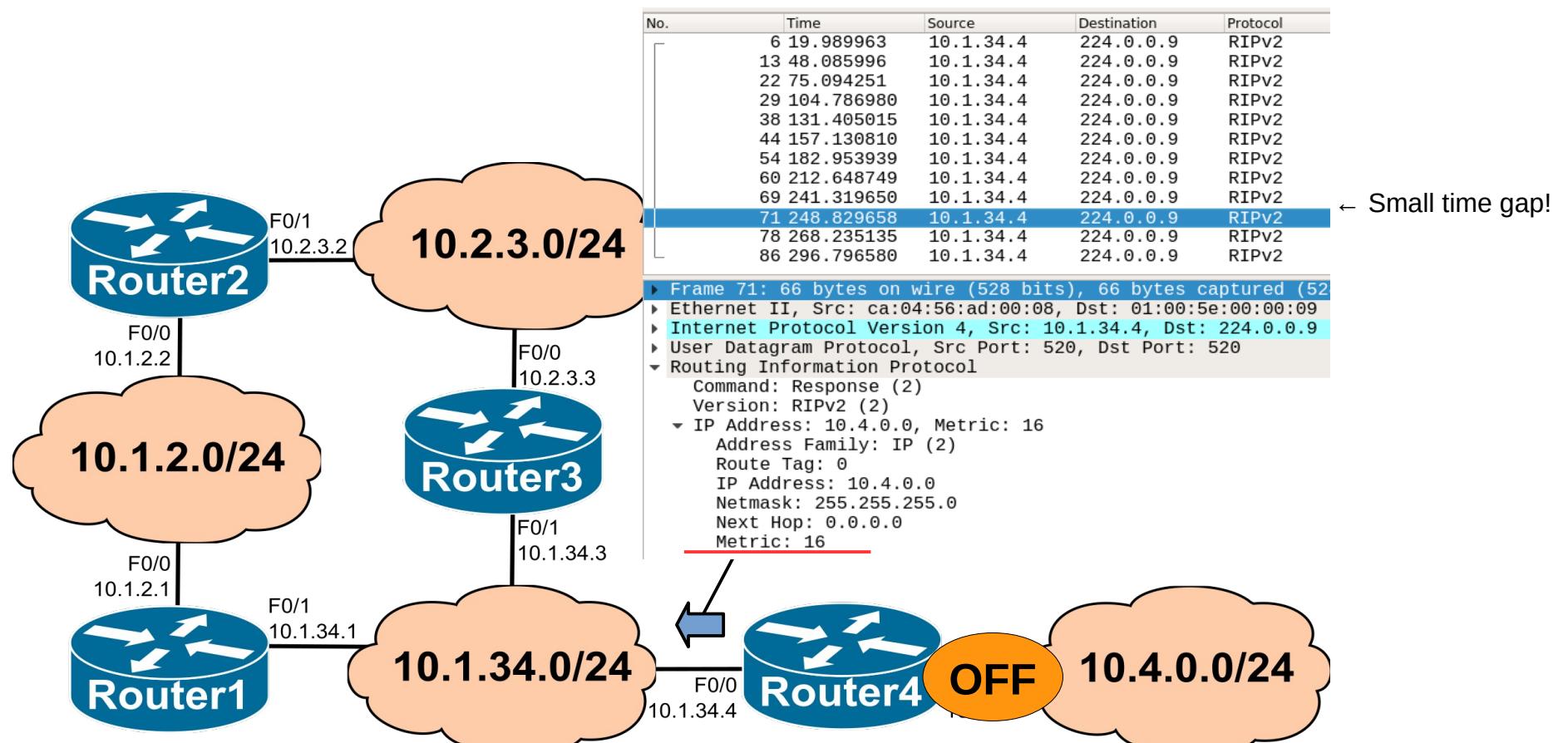
Sent by Router3 without Split-Horizon

```
► Internet Protocol Version 4, Src: 10.2.3.3, Dst: 224.0.0.9
► User Datagram Protocol, Src Port: 520, Dst Port: 520
▼ Routing Information Protocol
    Command: Response (2)
    Version: RIPv2 (2)
    ▼ IP Address: 10.1.2.0, Metric: 2
        Address Family: IP (2)
        Route Tag: 0
        IP Address: 10.1.2.0
        Netmask: 255.255.255.0
        Next Hop: 10.2.3.2
        Metric: 2
    ▼ IP Address: 10.1.34.0, Metric: 1
        Address Family: IP (2)
        Route Tag: 0
        IP Address: 10.1.34.0
        Netmask: 255.255.255.0
        Next Hop: 0.0.0.0
        Metric: 1
    ▼ IP Address: 10.2.3.0, Metric: 1
        Address Family: IP (2)
        Route Tag: 0
        IP Address: 10.2.3.0
        Netmask: 255.255.255.0
        Next Hop: 0.0.0.0
        Metric: 1
    ▼ IP Address: 10.4.0.0, Metric: 2
        Address Family: IP (2)
        Route Tag: 0
        IP Address: 10.4.0.0
        Netmask: 255.255.255.0
        Next Hop: 0.0.0.0
        Metric: 2
```



Triggered Updates

- Prevents any routing loops that involve more than two routers.
- Whenever a router changes the metric for a route, it is required to send update messages almost immediately, even if it is not yet time for one of the regular update message.
- Neighboring routers update routing tables faster and overall convergence is faster.
 - ◆ Including entries that were removed by timeout!



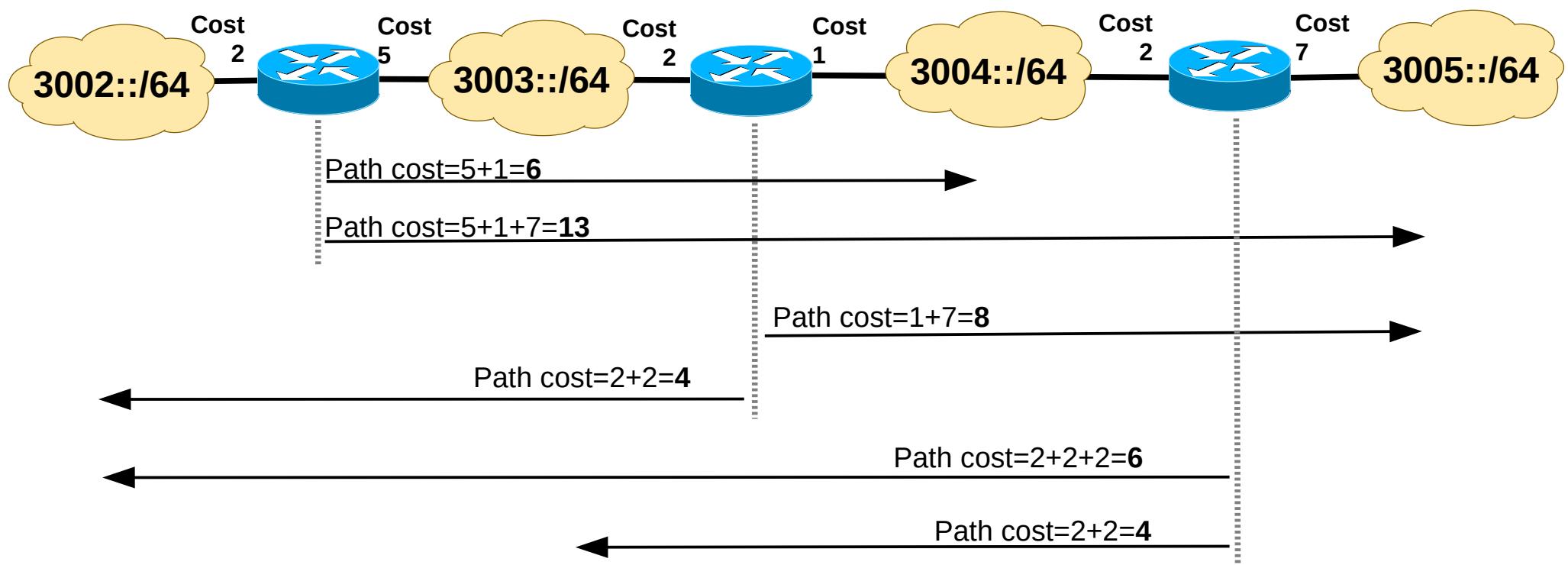
RIPng for IPv6 Routing

- Similar to IPv4 RIPv2:
 - ◆ Distance-vector concept, radius of 15 hops, infinity metric is 16, split-horizon, triggered update.
- Differences between RIPv2 and RIPng
 - ◆ Uses IPv6 for transport.
 - ◆ Uses link-local addresses (not the global ones).
 - ◆ IPv6 prefix, next-hop IPv6 link-local address.
 - ◆ Uses multicast group address FF02::9 (all-RIP-routers) as the destination address for RIP updates.
 - ◆ Routers always add the cost of the interface to the metric received.
 - ◆ Metric is sum of “output interfaces” costs to destination and not number of hops.
 - ◆ If all costs are 1, metric is number of “output interfaces” to destination.
 - ◆ Allows for node/interface costs other than 1.
 - ◆ Cisco calls it “cost offset” per interface (out or in direction).
 - ◆ Cost to network is given by the sum of all output interfaces costs along the path.
 - ◆ With the infinity metric value at 16, this require careful configurations.
 - ◆ Routers always announce a directed connected network to that same network.
 - ◆ Even with split-horizon activated, directed connected networks are not “learned” by RIPng (i.e., are not a RIP entry on the routing table).
 - RFC 2080 states that: “... split horizon algorithm omits routes learned from one neighbor in updates sent to that neighbor.”.
 - ◆ This can be justified by the fact that a router can have multiple IPv6 global network prefixes in one interface and, all routers in that network connection may not have all the same IPv6 global network prefixes configured.
 - ◆ in IOS Cisco
 - ◆ Activation per interface, named process, more than one active process.



RIPng Path Costs

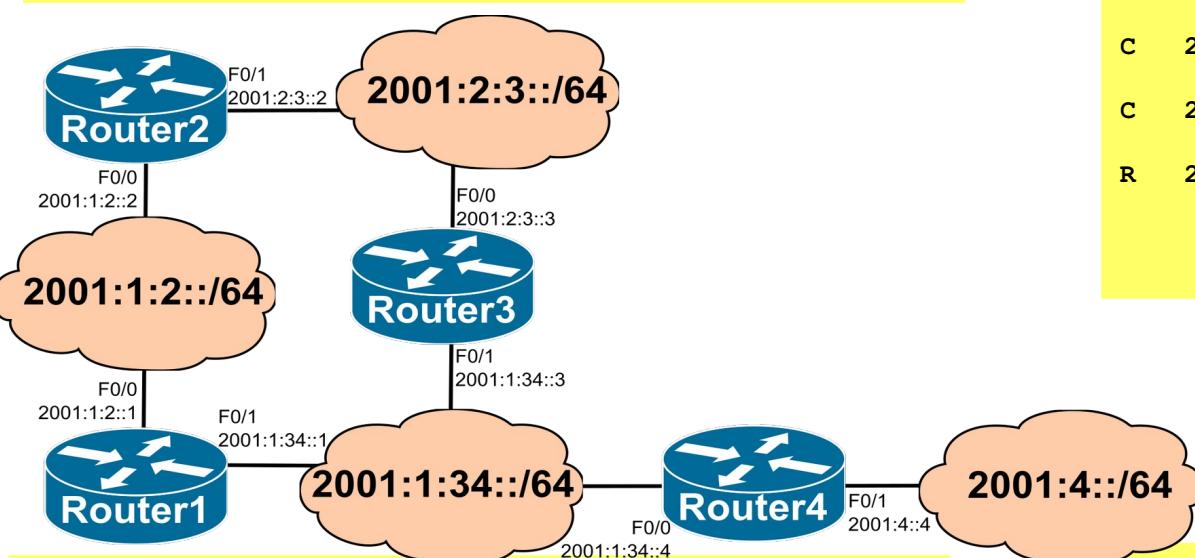
- Each router link/interface has an associated RIPng cost.
- The total cost between a router and a network is given by the sum of all RIPng costs of the (routers) output interfaces along the path.
 - ◆ Routers to access directly connect networks never use RIPng paths.



IPv6 Routing Tables with RIPng

Router2

```
C 2001:1:2::/64 [0/0]
  via FastEthernet0/0, directly connected
R 2001:1:34::/64 [120/2]
  via FE80::C801:54FF:FE41:8, FastEthernet0/0
  via FE80::C803:56FF:FE0A:8, FastEthernet0/1
C 2001:2:3::/64 [0/0]
  via FastEthernet0/1, directly connected
R 2001:4::/64 [120/3]
  via FE80::C801:54FF:FE41:8, FastEthernet0/0
  via FE80::C803:56FF:FE0A:8, FastEthernet0/1
```



Router1

```
C 2001:1:2::/64 [0/0]
  via FastEthernet0/0, directly connected
C 2001:1:34::/64 [0/0]
  via FastEthernet0/1, directly connected
R 2001:2:3::/64 [120/2]
  via FE80::C802:54FF:FEF5:8, FastEthernet0/0
  via FE80::C803:56FF:FE0A:6, FastEthernet0/1
R 2001:4::/64 [120/2]
  via FE80::C804:56FF:FEAD:8, FastEthernet0/1
```

Assuming all interfaces with cost 1.

Router3

```
R 2001:1:2::/64 [120/2]
  via FE80::C802:54FF:FEF5:6, FastEthernet0/0
  via FE80::C801:54FF:FE41:6, FastEthernet0/1
C 2001:1:34::/64 [0/0]
  via FastEthernet0/1, directly connected
C 2001:2:3::/64 [0/0]
  via FastEthernet0/0, directly connected
R 2001:4::/64 [120/2]
  via FE80::C804:56FF:FEAD:8, FastEthernet0/1
```

Router4

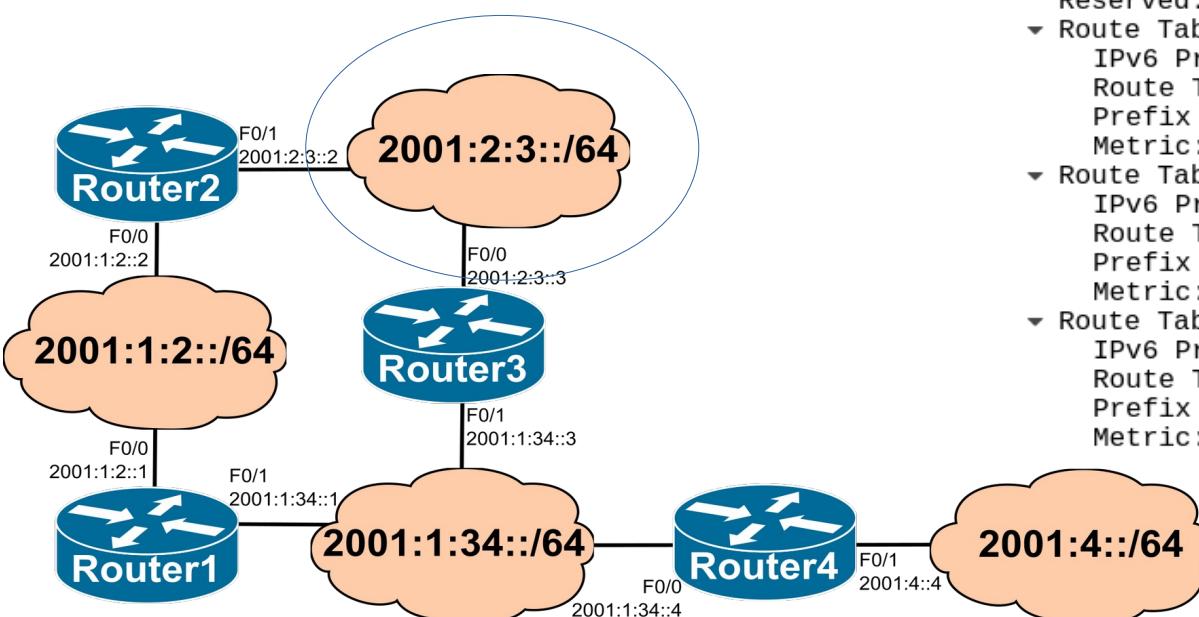
```
R 2001:1:2::/64 [120/2]
  via FE80::C801:54FF:FE41:6, FastEthernet0/0
C 2001:1:34::/64 [0/0]
  via FastEthernet0/0, directly connected
R 2001:2:3::/64 [120/2]
  via FE80::C803:56FF:FE0A:6, FastEthernet0/0
C 2001:4::/64 [0/0]
  via FastEthernet0/1, directly connected
```



RIPng Messages (Example)

Sent by Router2 with Split-Horizon

```
► Internet Protocol Version 6, Src: fe80::c802:54ff:fef5:6, Dst: ff02::9
► User Datagram Protocol, Src Port: 521, Dst Port: 521
▼ RIPng
  Command: Response (2)
  Version: 1
  Reserved: 0000
▼ Route Table Entry: IPv6 Prefix: 2001:1:2::/64 Metric: 1
  IPv6 Prefix: 2001:1:2::
  Route Tag: 0x0000
  Prefix Length: 64
  Metric: 1
▼ Route Table Entry: IPv6 Prefix: 2001:2:3::/64 Metric: 1
  IPv6 Prefix: 2001:2:3::
  Route Tag: 0x0000
  Prefix Length: 64
  Metric: 1
```



- Since network 2001:2:3::/64 does not appear in Routers 2 and 3 routing tables as a RIPng (R) entry, the split-horizon algorithm does not omit it from updates.
- Both routers announce network 2001:2:3::/64 to that same network.

Sent by Router3 with Split-Horizon

```
► Internet Protocol Version 6, Src: fe80::c803:56ff:fe0a:8, Dst: ff02::9
► User Datagram Protocol, Src Port: 521, Dst Port: 521
▼ RIPng
  Command: Response (2)
  Version: 1
  Reserved: 0000
▼ Route Table Entry: IPv6 Prefix: 2001:2:3::/64 Metric: 1
  IPv6 Prefix: 2001:2:3::
  Route Tag: 0x0000
  Prefix Length: 64
  Metric: 1
▼ Route Table Entry: IPv6 Prefix: 2001:1:34::/64 Metric: 1
  IPv6 Prefix: 2001:1:34::
  Route Tag: 0x0000
  Prefix Length: 64
  Metric: 1
▼ Route Table Entry: IPv6 Prefix: 2001:4::/64 Metric: 2
  IPv6 Prefix: 2001:4::
  Route Tag: 0x0000
  Prefix Length: 64
  Metric: 2
```

Perguntas sobre isto não saem no teste

Os RIPvns não saem mt no teste, é mais uma coisa histórica



Open Shortest Path First (OSPF) Protocol

- OSPF is an open-standard protocol based primarily on RFC 2328.
- OSPF is a link-state routing protocol
 - ◆ Respond quickly to network changes,
 - ◆ Send triggered updates when a network change occurs,
 - ◆ Send periodic updates, known as link-state refresh, at long time intervals, such as every 30 minutes.
- Routers running OSPF collect routing information from all other routers in the network (or from within a defined area of the network)
- And then each router independently calculates its best paths to all destinations in the network, using Dijkstra's (SPF) algorithm.



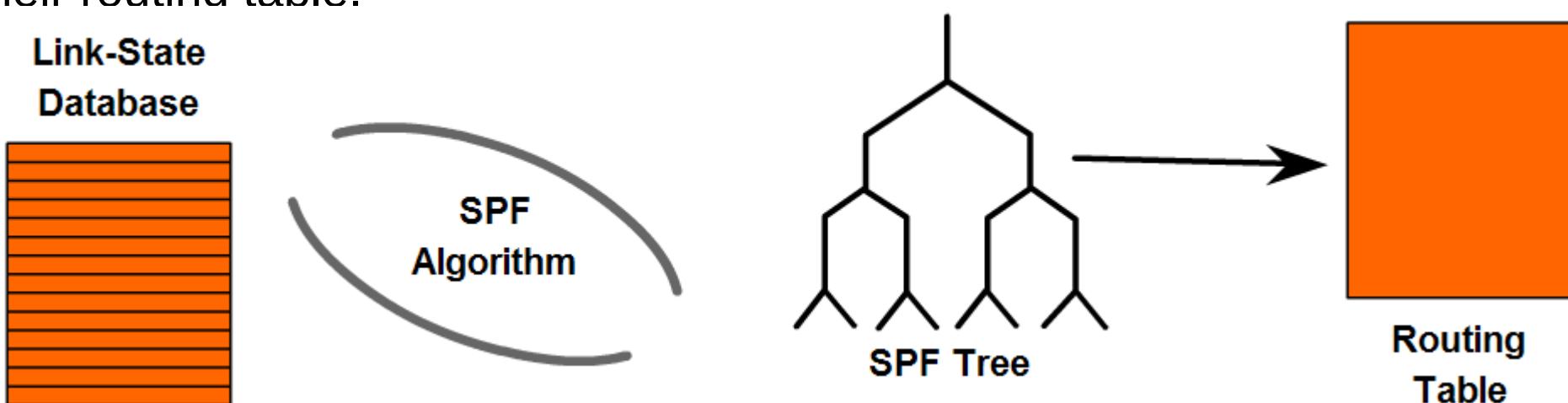
OSPF Necessary Routing Information

- For all the routers in the network to make consistent routing decisions, each link-state router must keep a record of the following information:
 - ◆ Its immediate neighbor routers
 - If the router loses contact with a neighbor router, within a few seconds it invalidates all paths through that router and recalculates its paths through the network.
 - For OSPF, adjacency information about neighbors is stored in the OSPF neighbor table, also known as an adjacency database.
 - ◆ All the other routers in the network, or in its area of the network, and their attached networks
 - The router recognizes other routers and networks through LSAs, which are flooded through the network.
 - LSAs are stored in a topology table or database (which is also called an LSDB).
 - ◆ The best paths to each destination
 - Each router independently calculates the best paths to each destination in the network using Dijkstra's (SPF) algorithm.
 - All paths are kept in the LSDB.
 - The best paths are then offered to the routing table (also called the forwarding database).
 - Packets arriving at the router are forwarded based on the information held in the routing table.



Link-State Protocol Operation

- Link-state routing protocols generate routing updates only when a change occurs in the network topology.
- When a link changes state, the device that detected the change creates a Link-State Advertisement (LSA) concerning that link.
 - ◆ LSA propagates to neighbor devices using a special multicast address.
- Each router stores the LSA, forwards the LSA to neighboring devices and updates its Link-State DataBase (LSDB).
- Link-state routers find the best paths to a destination by applying Dijkstra's algorithm, also known as SPF, against the LSDB to build the SPF tree.
- Each router selects the best paths from their SPF tree and places them in their routing table.



Link-State Advertisement (LSA)

- LSAs report the state of routers and the links between routers.
- Link-state information must be synchronized between routers.
- LSAs have the following characteristics:
 - ◆ LSAs are reliable. There is a method for acknowledging their delivery.
 - ◆ LSAs are flooded throughout the area (or throughout the domain if there is only one area).
 - ◆ LSAs have a sequence number and a set lifetime, so each router recognizes that it has the most current version of the LSA.
 - ◆ LSAs are periodically refreshed to confirm topology information before they age out of the LSDB.



OSPF Router ID (RID)

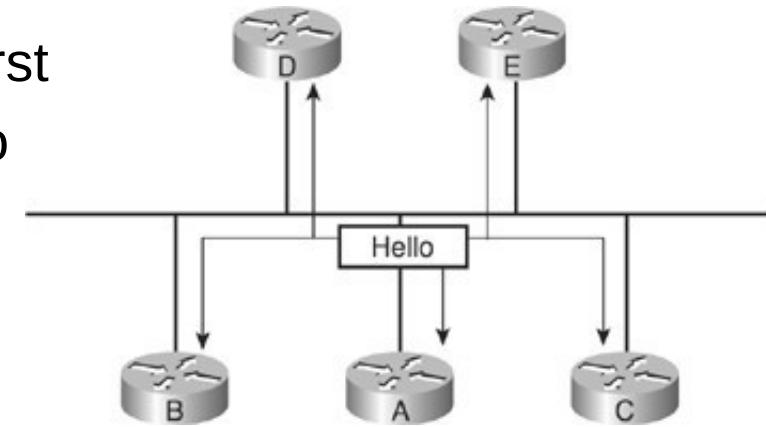
- The Router ID identifies the router and is:
 - ◆ The highest IPv4 address of all router interfaces at the moment of the OSPF process activation.
 - ◆ A value administratively defined.
- If a physical interface address is being used as the router ID, and that physical interface fails, and the router (or OSPF process) is restarted, the router ID will change.
 - ◆ This change in router ID makes it more difficult for network administrators to troubleshoot and manage OSPF.
- Administratively defining the RID or using loopback interfaces for the router ID forces the router ID to stay the same, regardless of the state of the physical interfaces.



OSPF Adjacencies

(Vizinhanças dos Routers, eles têm de descobrir os vizinhos)

- A router running a link-state routing protocol must first establish neighbor adjacencies, by exchanging hello packets with the neighboring routers
- The router sends and receives Hello packets to and from its neighboring routers.
 - ◆ The destination address is typically a multicast address.
 - ◆ It is possible to define unicast OSPF relations.
- The routers exchange hello packets subject to protocol-specific parameters, such as checking whether the neighbor is in the same area, using the same hello interval, and so on.
 - ◆ Routers declare the neighbor up when the exchange is complete.
- Two OSPF routers on a point-to-point serial link, usually encapsulated in High-Level Data Link Control (HDLC) or Point-to-Point Protocol (PPP), form a full adjacency with each other.
- However, OSPF routers on broadcast networks, such as LAN links, elect one router as the designated router (DR) and another as the backup designated router (BDR).
 - ◆ All other routers on the LAN form full adjacencies with these two routers and pass LSAs only to them.



Não temos de saber exatamente os pacotes, mas as funcionalidades dos mesmos



DR and BDR Election

- The first OSPF router to boot becomes the Designated Router (DR).
- The second router to boot becomes the Backup Designated Router (BDR).
- If multiple routers boot simultaneously,
 - ◆ The DR it will be the router with the highest priority. The BDR the second.
 - ◆ The OSPF priority is a administratively defined parameter.
 - ◆ In case of tie, it will be chosen the router with the highest Router ID (RID).
- When the DR fails, the BDR assumes the role of DR.
 - ◆ The BDR does not perform any DR functions when the DR is operating.
 - ◆ The choice of the new BDR is done according to some criteria of the initial election.
- After the election, the DR and BDR maintain that role, independently of which routers join the OSPF process.
- The ID of an OSPF Network is the IP address of the network's Designated Router (DR) interface.



OSPF LS Database

- The OSPF database (LSDB) is organized in two tables.
 - Router Link States – Routers related information table.
 - The routers are identified by theirs RID.
 - Net Link States – Networks/Links related information table.
 - Networks are identified by their ID.

| OSPF Router with ID (20.20.20.1) (Process ID 1) | | | | | |
|---|------------|-----|------------|----------|------------|
| Router Link States (Area 0) | | | | | |
| Link ID | ADV Router | Age | Seq# | Checksum | Link count |
| 20.20.20.1 | 20.20.20.1 | 40 | 0x8000000A | 0x00E7FB | 2 |
| 30.30.30.2 | 30.30.30.2 | 69 | 0x80000006 | 0x002906 | 2 |
| 30.30.30.3 | 30.30.30.3 | 41 | 0x80000007 | 0x00283D | 2 |
| Net Link States (Area 0) | | | | | |
| Link ID | ADV Router | Age | Seq# | Checksum | |
| 10.10.10.3 | 30.30.30.3 | 41 | 0x80000001 | 0x00051C | |
| 20.20.20.2 | 30.30.30.2 | 70 | 0x80000001 | 0x00A164 | |
| 30.30.30.3 | 30.30.30.3 | 154 | 0x80000001 | 0x00A91C | |



OSPF LS Database Tables (1)

- Router Link States

- For each router, it contains the information about the networks directly connected to that router.

```
LS age: 321
Options: (No TOS-capability, DC)
LS Type: Router Links
Link State ID: 20.20.20.1 ← Router ID
Advertising Router: 20.20.20.1
LS Seq Number: 8000000A
Checksum: 0xE7FB
Length: 48
Number of Links: 2 ← Number of Links → Nº de interfaces (nº de ligações a redes)

Link connected to: a Transit Network ← Network Type
(Link ID) Designated Router address: 20.20.20.2 ← Network ID
(Link Data) Router Interface address: 20.20.20.1 ← Interface IP Address
Number of TOS metrics: 0
TOS 0 Metrics: 1 ← Interface Cost

Link connected to: a Transit Network
(Link ID) Designated Router address: 10.10.10.3
(Link Data) Router Interface address: 10.10.10.1
Number of TOS metrics: 0
TOS 0 Metrics: 1
```



OSPF LS Database Tables (2)

- Network Link States

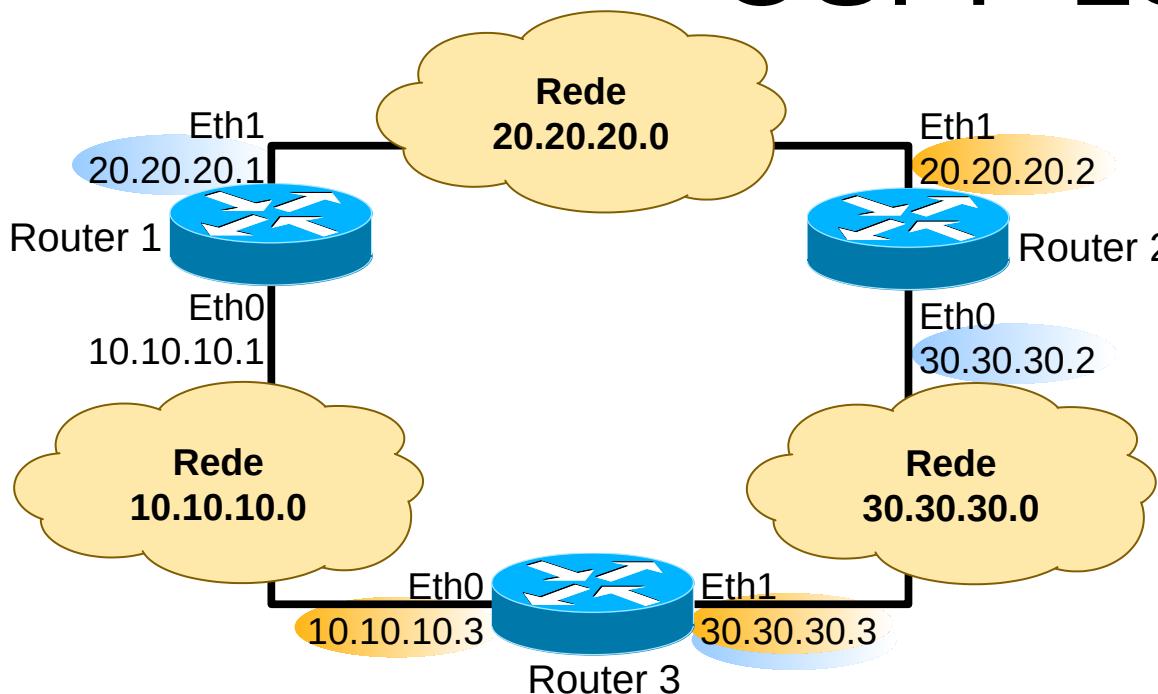
- For each network, it contains the information about the routers directly attached to that network.

```
Routing Bit Set on this LSA
LS age: 483
Options: (No TOS-capability, DC)
LS Type: Network Links
Link State ID: 10.10.10.3 (address of Designated Router)
Advertising Router: 30.30.30.3
LS Seq Number: 80000001
Checksum: 0x51C
Length: 32
Network Mask: /24
Attached Router: 30.30.30.3
Attached Router: 20.20.20.1 }
```

Network ID
Attached routers (RID)



OSPF LSDatabase Example



Routing Bit Set on this LSA

LS age: 208
Options: (No TOS-capability, DC)
LS Type: Network Links
Link State ID: 20.20.20.2 (address of Designated Router)
Advertising Router: 30.30.30.2
LS Seq Number: 80000001
Checksum: 0xA164
Length: 32
Network Mask: /24
Attached Router: 30.30.30.2
Attached Router: 20.20.20.1

Network 20.20.20.0's Network Link State

LS age: 321
Options: (No TOS-capability, DC)
LS Type: Router Links
Link State ID: 20.20.20.1
Advertising Router: 20.20.20.1
LS Seq Number: 8000000A
Checksum: 0xE7FB
Length: 48
Number of Links: 2

Link connected to: a Transit Network
(Link ID) Designated Router address: 20.20.20.2
(Link Data) Router Interface address: 20.20.20.1
Number of TOS metrics: 0
TOS 0 Metrics: 1

Link connected to: a Transit Network
(Link ID) Designated Router address: 10.10.10.3
(Link Data) Router Interface address: 10.10.10.1
Number of TOS metrics: 0
TOS 0 Metrics: 1

Router 1's Router Link State



OSPF Packets

- Hello - Discovers neighbors and builds adjacencies between them.
- Database Description (DBD) - Checks for database synchronization between routers.
- Link-State Request (LSR) - Requests specific link-state records from another router.
- Link-State Update (LSU) - Sends specifically requested link-state records.
- LSAck - Acknowledges the other packet types.



OSPF Packet Format

- Version Number

- Set to 2 for OSPF Version 2, the IPv4 version of OSPF.
- Set to 3 for OSPF Version 3, the IPv6 version of OSPF.

- Type

- Differentiates the five OSPF packet types.

- Packet Length

- The length of the OSPF packet in bytes.

- Router ID

- Defines which router is the packet's source.

- Area ID

- Defines the area in which the packet originated.

- Checksum

- Used for packet header error detection to ensure that the OSPF packet was not corrupted during transmission.

- Authentication Type

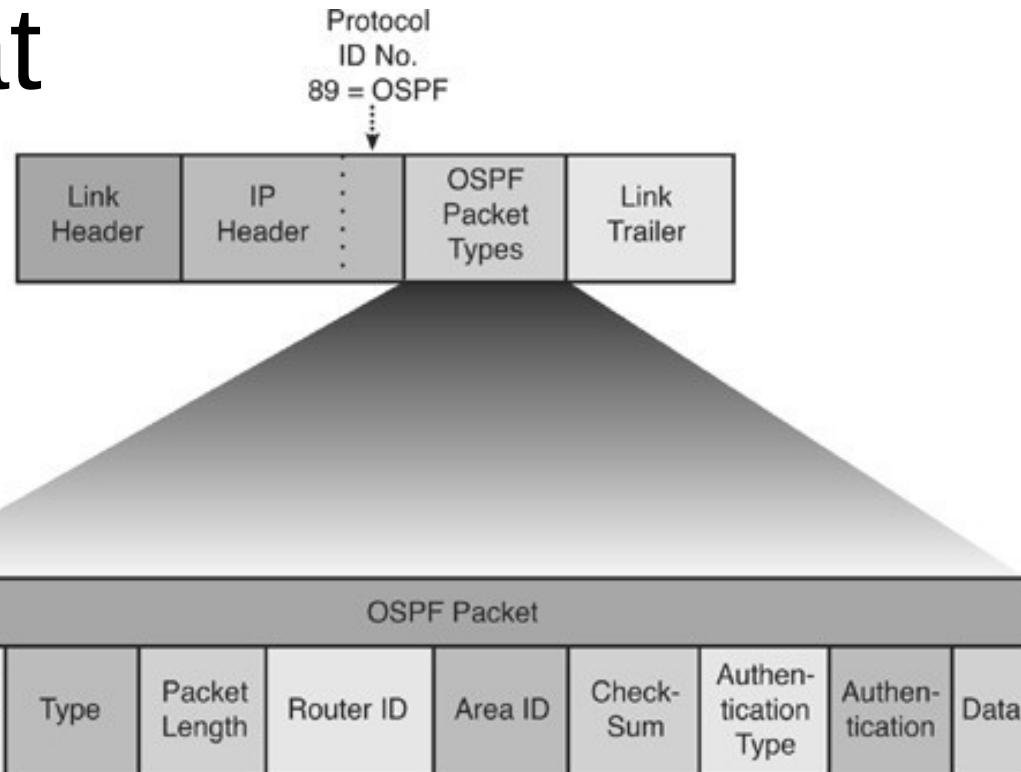
- An option in OSPF that describes either no authentication, clear-text passwords, or encrypted message digest 5 (MD5) for router authentication.

- Authentication

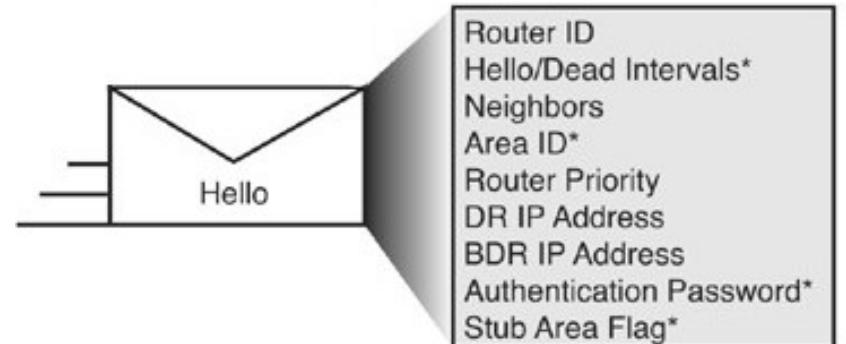
- Used with authentication type.

- Data, contains different information, depending on the OSPF packet type:

- For the Hello packet - Contains a list of known neighbors.
- For the DBD packet - Contains a summary of the LSDB, which includes all known router IDs and their last sequence number, among several other fields.
- For the LSR packet - Contains the type of LSU needed and the router ID of the router that has the needed LSU.
- For the LSU packet - Contains the full LSA entries. Multiple LSA entries can fit in one OSPF update packet.
- For the LSAck packet - This data field is empty.



OSPF Hello Packets



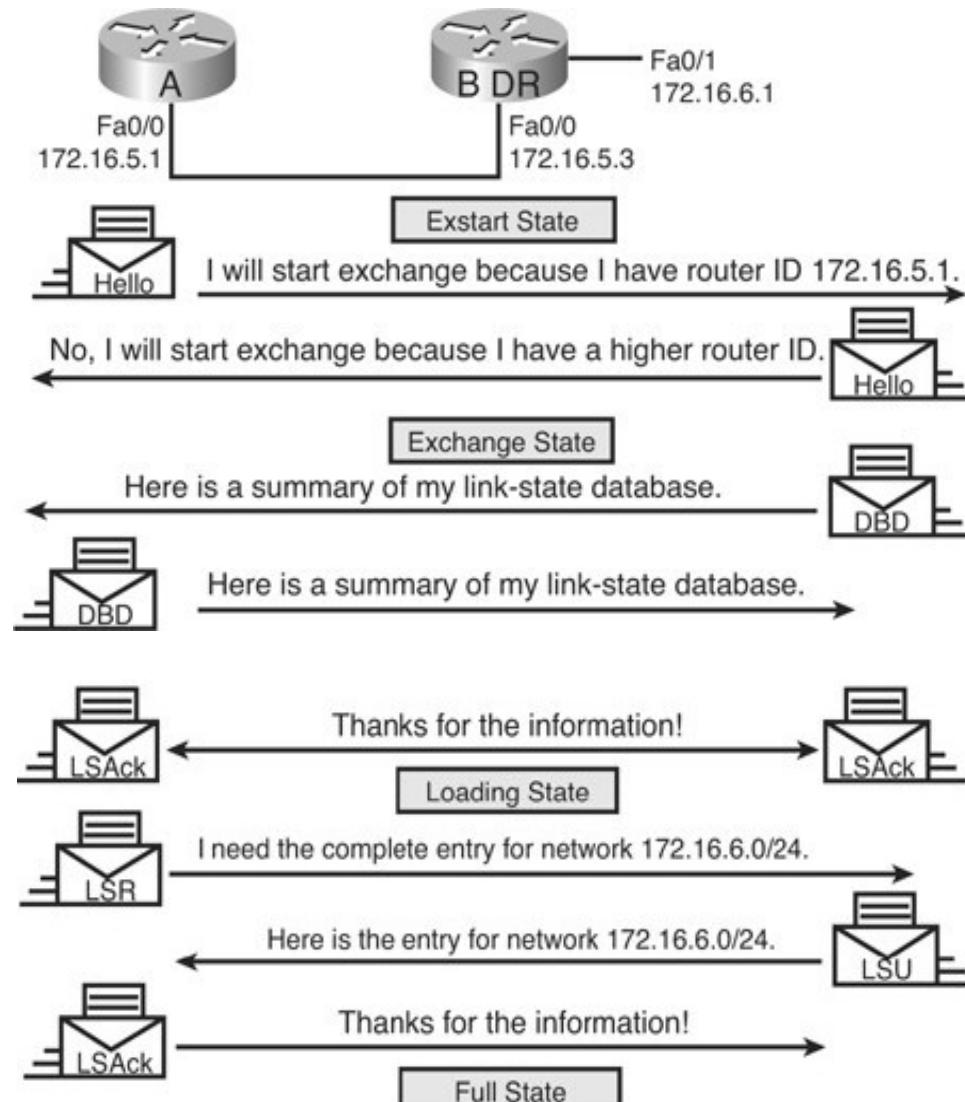
- An hello packet contains the following information:

- Router ID
 - A 32-bit number that uniquely identifies the router.
- Hello and dead intervals
 - The hello interval specifies how often, in seconds, a router sends hello packets (10 seconds is the default on multiaccess networks).
 - The dead interval is the amount of time in seconds that a router waits to hear from a neighbor before declaring the neighbor router out of service (the dead interval is four times the hello interval by default).
 - These timers must be the same on neighboring routers; otherwise an adjacency will not be established.
- Neighbors
 - The Neighbors field lists the adjacent routers with which this router has established bidirectional communication.
 - Bidirectional communication is indicated when the router sees itself listed in the Neighbors field of the hello packet from the neighbor.
- Area ID
 - To communicate, two routers must share a common segment, and their interfaces must belong to the same OSPF area on that segment.
 - These routers will all have the same link-state information for that area.
- Router priority
 - An 8-bit number that indicates a router's priority. Priority is used when electing a DR and BDR.
- DR and BDR IP addresses
 - If known, the IP addresses of the DR and BDR for the specific multiaccess network.
- Authentication password
 - If router authentication is enabled, two routers must exchange the same password.
 - Authentication is not required, but if it is enabled, all peer routers must have the same password.
- Stub area flag
 - A stub area is a special area.
 - The stub area technique reduces routing updates by replacing them with a default route.
 - Two neighboring routers must agree on the stub area flag in the hello packets.

- Hello Interval, Dead Interval, Area ID, Authentication Password and Stub Area Flag fields must match on neighboring routers for them to establish an adjacency.



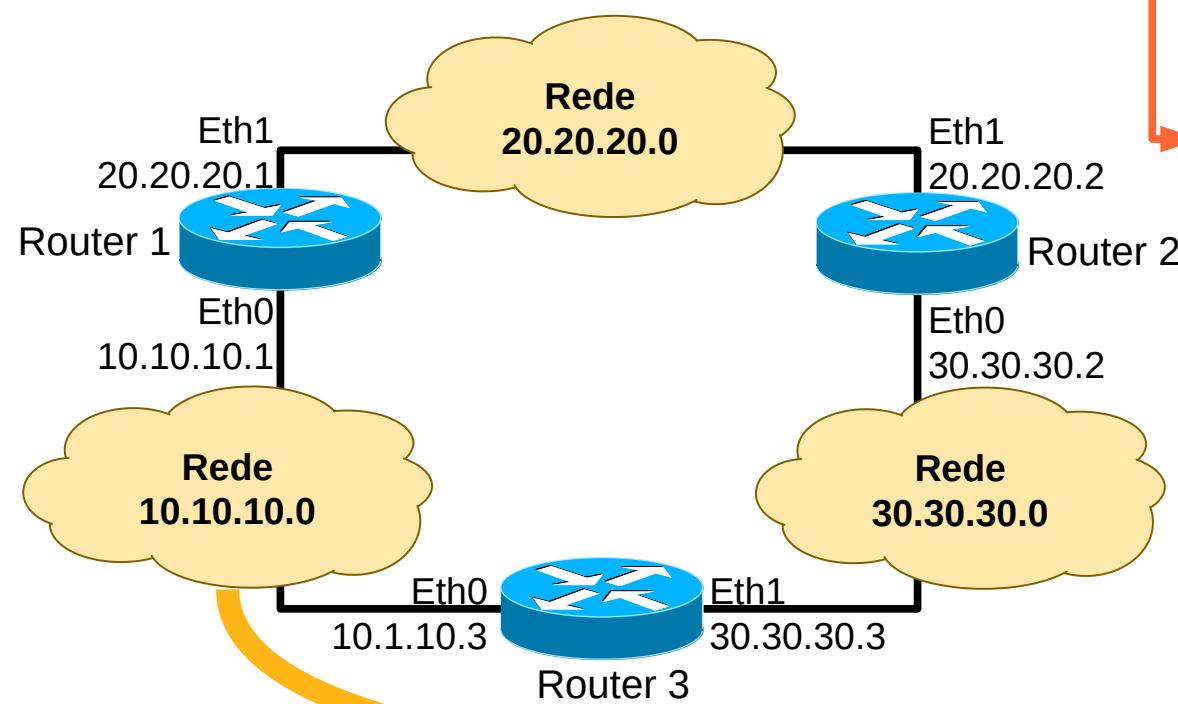
Discovering the Network Routes



- A master and slave relationship is created between each router and its adjacent DR and BDR.
 - ◆ Only the DR exchanges and synchronizes link-state information with the routers to which it has established adjacencies.
- The master and slave routers exchange one or more DBD packets.
 - ◆ A DBD includes information about the LSA entry header that appears in the router's LSDB.
 - ◆ The entries can be about a link or about a network.
 - ◆ Each LSA entry header includes information about the link-state type, the address of the advertising router, the link's cost, and the sequence number.
 - ◆ The router uses the sequence number to determine the "newness" of the received link-state information.
- It acknowledges the receipt of the DBD using the LSAck packet.
 - ◆ It compares the information it received with the information it has in its own LSDB.
- If the DBD has a more current link-state entry, the router sends an LSR to the other router.
- The other router responds with the complete information about the requested entry in an LSU packet.
- Again, when the router receives an LSU, it sends an LSAck.
- The router adds the new link-state entries to its LSDB.



OSPF Example



OSPF activated on Router 1

OSPF activated on Router 3

OSPF activated on Router 2

| Time | Source | Destination | Protocol Info |
|------------|------------|-------------|---------------------|
| 0.000000 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| 10.002318 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| 20.003116 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| ----- | | | |
| 80.000000 | 10.10.10.3 | 224.0.0.5 | OSPF Hello Packet |
| 83.683033 | 10.10.10.3 | 224.0.0.5 | OSPF LS Update |
| 83.715683 | 10.10.10.3 | 224.0.0.5 | OSPF Hello Packet |
| 83.717864 | 10.10.10.1 | 10.10.10.3 | OSPF Hello Packet |
| 83.726166 | 10.10.10.3 | 10.10.10.1 | OSPF DB Descr. |
| 83.726258 | 10.10.10.3 | 10.10.10.1 | OSPF Hello Packet |
| 83.728433 | 10.10.10.1 | 10.10.10.3 | OSPF DB Descr. |
| 83.732590 | 10.10.10.3 | 10.10.10.1 | OSPF DB Descr. |
| 83.734733 | 10.10.10.1 | 10.10.10.3 | OSPF DB Descr. |
| 83.738942 | 10.10.10.3 | 10.10.10.1 | OSPF LS Request |
| 83.741083 | 10.10.10.1 | 10.10.10.3 | OSPF LS Update |
| 84.240362 | 10.10.10.3 | 224.0.0.5 | OSPF LS Update |
| 86.245792 | 10.10.10.3 | 224.0.0.5 | OSPF LS Acknowledge |
| 86.380876 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| 86.741036 | 10.10.10.1 | 224.0.0.5 | OSPF LS Acknowledge |
| 93.721376 | 10.10.10.3 | 224.0.0.5 | OSPF Hello Packet |
| 96.380005 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| ----- | | | |
| 213.780338 | 10.10.10.3 | 224.0.0.5 | OSPF Hello Packet |
| 216.542473 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| 216.568852 | 10.10.10.1 | 224.0.0.5 | OSPF LS Update |
| 217.048427 | 10.10.10.1 | 224.0.0.5 | OSPF LS Update |
| 217.084909 | 10.10.10.1 | 224.0.0.5 | OSPF LS Update |
| 219.067748 | 10.10.10.3 | 224.0.0.5 | OSPF LS Acknowledge |
| 219.650308 | 10.10.10.1 | 224.0.0.5 | OSPF LS Update |
| 222.150349 | 10.10.10.3 | 224.0.0.5 | OSPF LS Acknowledge |
| 223.779492 | 10.10.10.3 | 224.0.0.5 | OSPF Hello Packet |
| 224.284149 | 10.10.10.3 | 224.0.0.5 | OSPF LS Update |
| 224.789598 | 10.10.10.1 | 224.0.0.5 | OSPF LS Update |
| 224.789775 | 10.10.10.3 | 224.0.0.5 | OSPF LS Update |
| 226.545718 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |
| 226.785254 | 10.10.10.1 | 224.0.0.5 | OSPF LS Acknowledge |
| 227.294756 | 10.10.10.3 | 224.0.0.5 | OSPF LS Acknowledge |
| 233.779863 | 10.10.10.3 | 224.0.0.5 | OSPF Hello Packet |
| 236.544658 | 10.10.10.1 | 224.0.0.5 | OSPF Hello Packet |

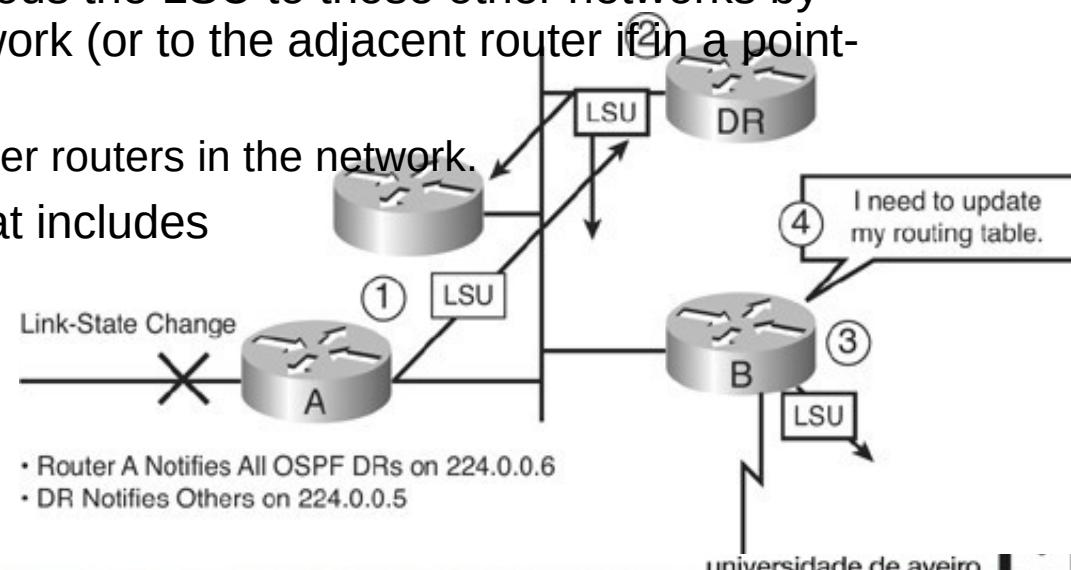


Maintaining Routing Information

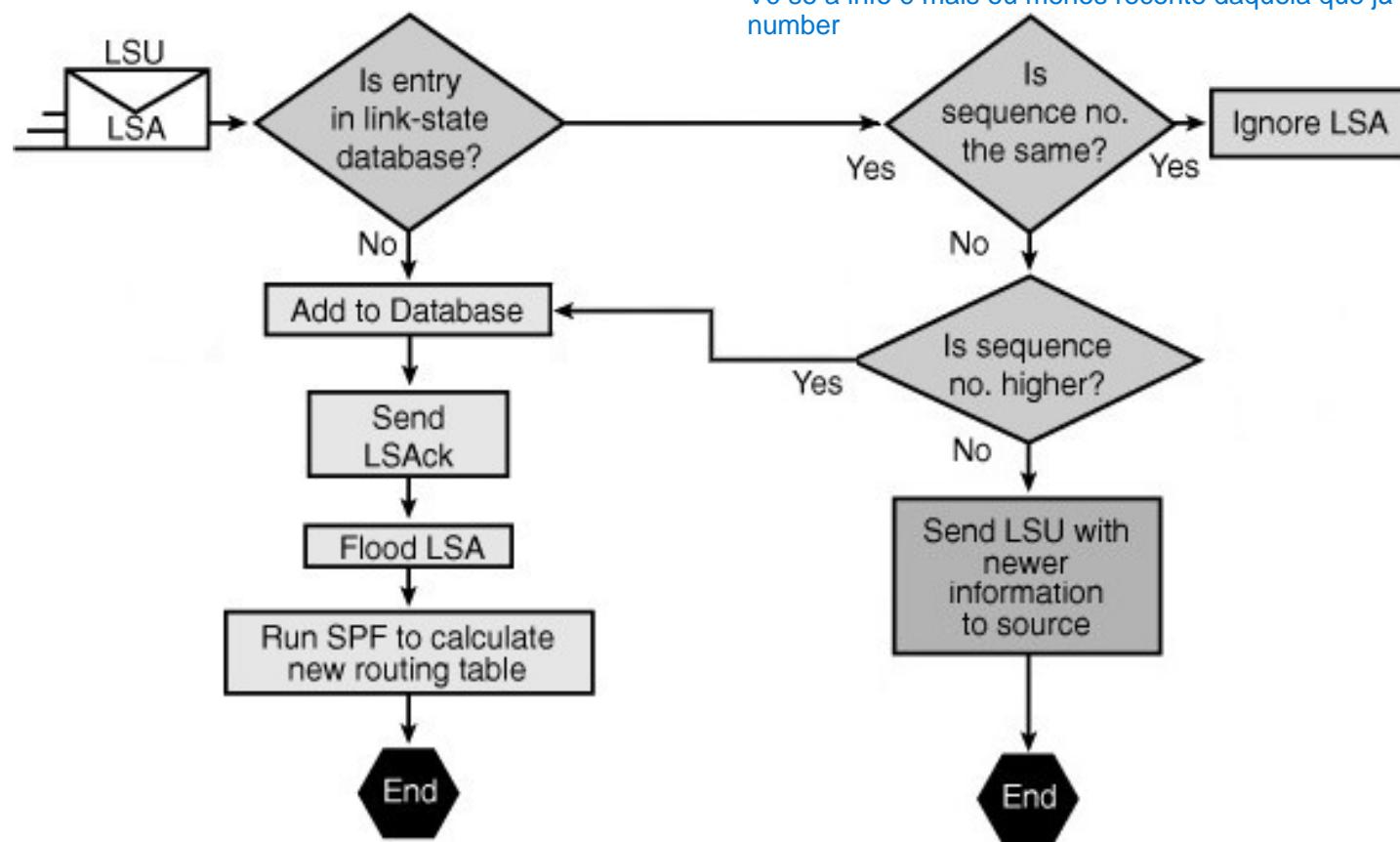
Há dois grupos de Multicast

- Flooding process:

- A router notices a change in a link state and multicasts an LSU packet, which includes the updated LSA entry with the sequence number incremented, to 224.0.0.6.
 - This address goes to all OSPF DRs and BDRs.
 - On point-to-point links, the LSU is multicast to 224.0.0.5.)
 - An LSU packet might contain several distinct LSAs.
- The DR receives the LSU, processes it, acknowledges the receipt of the change and floods the LSU to other routers on the network using the OSPF multicast address 224.0.0.5.
 - After receiving the LSU, each router responds to the DR with an LSAck.
 - To make the flooding procedure reliable, each LSA must be acknowledged separately.
- If a router is connected to other networks, it floods the LSU to those other networks by forwarding the LSU to the DR of the other network (or to the adjacent router if in a point-to-point network).
 - That DR, in turn, multicasts the LSU to the other routers in the network.
- The router updates its LSDB using the LSU that includes the changed LSA.
- It then recomputes the SPF algorithm against the updated database after a short delay and updates the routing table as necessary.



LSA Operation



A diferença entre OSPF e RIP é que no RIP, ao desligar uma interface, tem de calcular os caminhos tds outra vez, no OSPF ao mudar pequenas coisas não tem de calcular tdo outra vez, apenas são trocados alguns pacotes e atualiza-se a topologia da rede

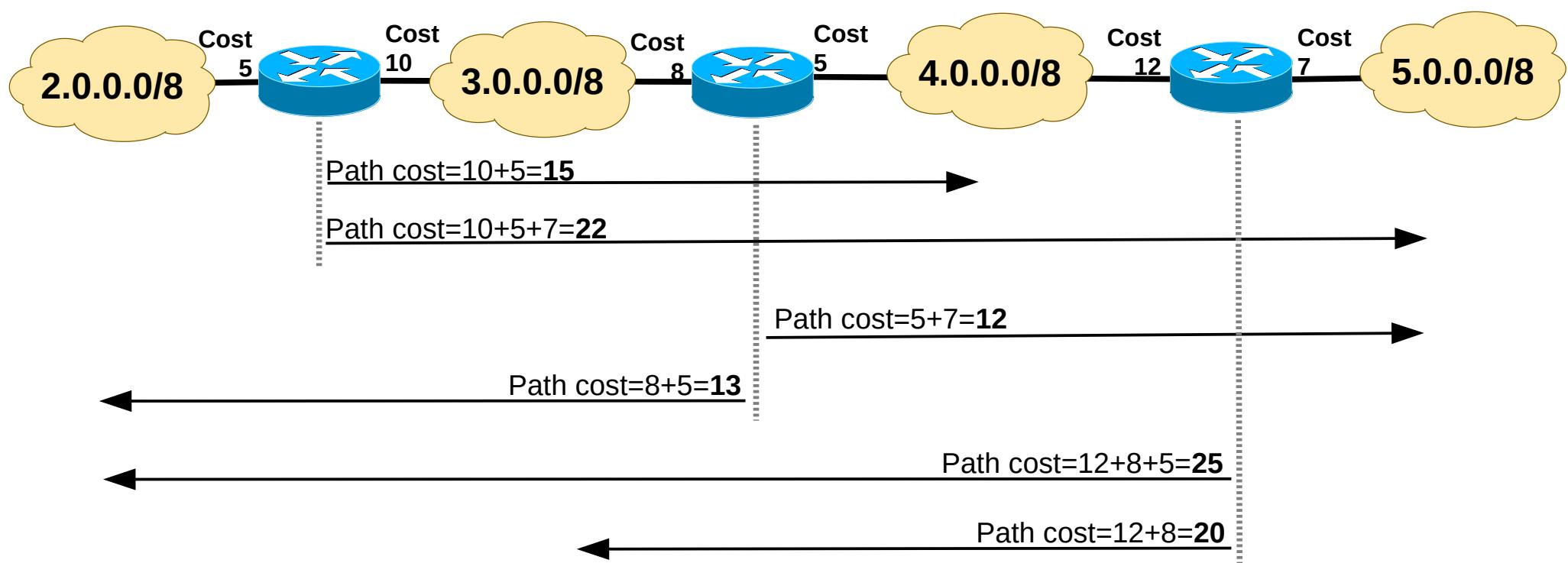
- When each router receives the LSU:

- If the LSA entry does not already exist, the router adds the entry to its LSDB, sends back a link-state acknowledgment (LSAck), floods the information to other routers, runs SPF, and updates its routing table.
- If the entry already exists and the received LSA has the same sequence number, the router ignores the LSA entry.
- If the entry already exists but the LSA includes newer information (it has a higher sequence number), the router adds the entry to its LSDB, sends back an LSAck, floods the information to other routers, runs SPF, and updates its routing table.
- If the entry already exists but the LSA includes older information, it sends an LSU to the sender with its newer information.

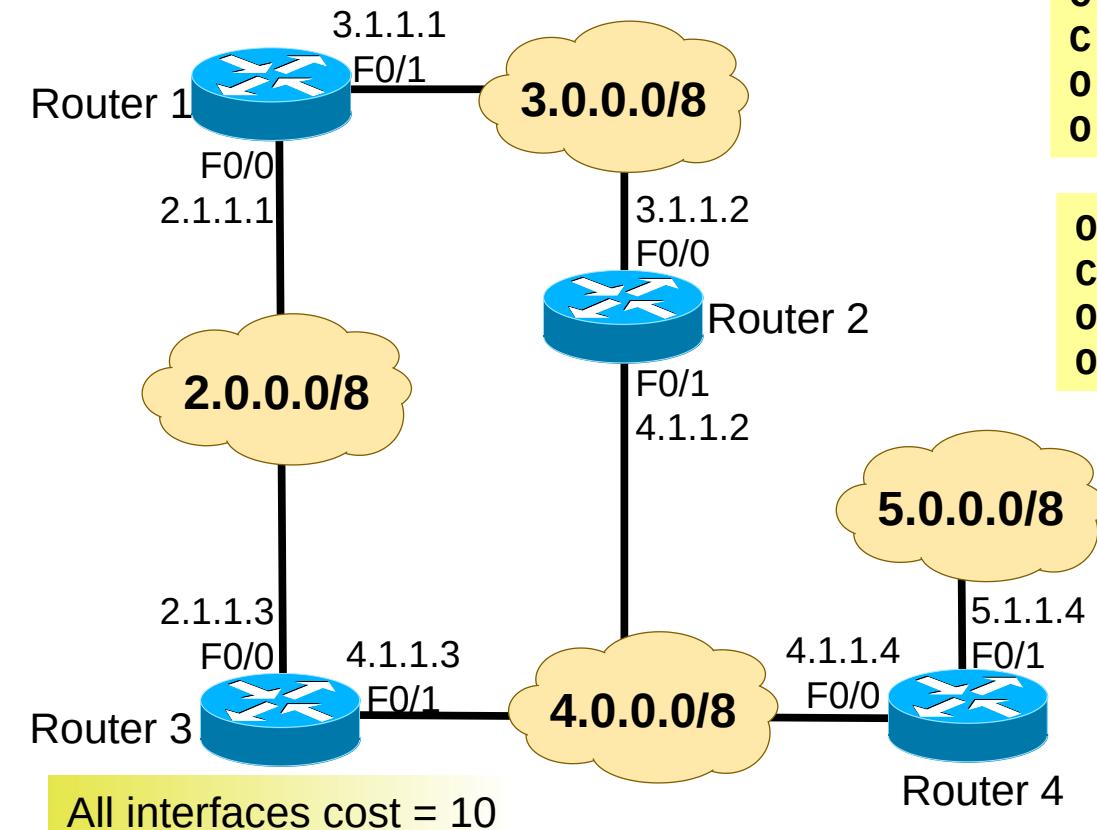


OSPF Path Costs

- Each router link/interface has an associated OSPF cost.
- The total cost between a router and a network is given by the sum of all OSPF costs of the (routers) output interfaces along the path.
 - ◆ Routers to access directly connect networks never use OSPF paths.



OSPF Example



```
C 2.0.0.0/8 is directly connected, F0/0
C 3.0.0.0/8 is directly connected, F0/1
O 4.0.0.0/8 [110/20] via 2.1.1.3, 00:01:18, F0/0
O 5.0.0.0/8 [110/30] via 2.1.1.3, 00:01:00, F0/0
```

```
O 2.0.0.0/8 [110/20] via 3.1.1.1, 00:01:13, F0/0
C 3.0.0.0/8 is directly connected, F0/0
O 4.0.0.0/8 [110/30] via 3.1.1.1, 00:01:13, F0/0
O 5.0.0.0/8 [110/40] via 3.1.1.1, 00:01:10, F0/0
```

Router 1 and Router 2 after disconnecting the F0/1 at Router2

```
C 2.0.0.0/8 is directly connected, F0/0
C 3.0.0.0/8 is directly connected, F0/1
O 4.0.0.0/8 [110/15] via 3.1.1.2, 00:01:13, F0/1
O 5.0.0.0/8 [110/25] via 3.1.1.2, 00:01:10, F0/1
```

Router1, now with the cost of Router2 F0/1 interface equal to 5

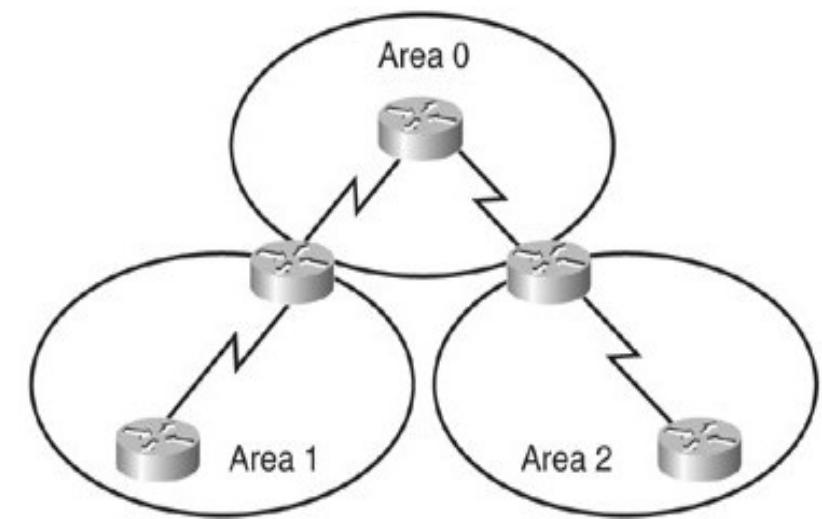
```
C 2.0.0.0/8 is directly connected, F0/0
C 3.0.0.0/8 is directly connected, F0/1
O 4.0.0.0/8 [110/20] via 3.1.1.2, 00:01:13, F0/1
[110/20] via 2.1.1.3, 00:01:31, F0/0
O 5.0.0.0/8 [110/30] via 3.1.1.2, 00:01:10, F0/1
[110/30] via 2.1.1.3, 00:01:10, F0/0
```

Router 1 initial table



OSPF Hierarchical Routing (1)

- In small networks, the web of router links is not complex, and paths to individual destinations are easily deduced.
- In large networks, the resulting web is highly complex, and the number of potential paths to each destination is large.
 - ◆ Dijkstra calculations comparing all of these possible routes can be very complex and can take significant time.
 - ✚ Large LSDB. Because the LSDB covers the topology of the entire network, each router must maintain an entry for every network in the area, even if not every route is selected for the routing table.
 - ✚ Frequent SPF algorithm calculations. In a large network, changes are inevitable, so the routers spend many CPU cycles recalculating the SPF algorithm and updating the routing table.
 - ✚ Large routing table. OSPF does not perform route summarization by default. If the routes are not summarized, the routing tables can become very large, depending on the size of the network.
- Link-state routing protocols usually reduce the size of the Dijkstra calculations by partitioning the network into areas.



OSPF Hierarchical Routing (2)

- Using multiple OSPF areas has several important advantages:
 - ◆ Reduced frequency of SPF calculations.
 - ✚ Detailed route information only exists within each area
 - ✚ It is not necessary to flood all link-state changes to all other areas.
 - ✚ Only routers that are affected by the change need to recalculate the SPF algorithm and the impact of the change is localized within the area.
 - ◆ Reduced updates overhead.
 - ✚ Rather than send an update about each network within an area, a router can advertise a single summarized route or a small number of routes between areas, thereby reducing the overhead associated with updates when they cross areas.
 - ◆ Smaller routing tables.
 - ✚ Detailed route entries for specific networks within an area can remain in the area.
 - ✚ Routers can be configured to summarize the routes into one or more summary addresses.
 - ✚ Advertising these summaries reduces the number of messages propagated between areas but keeps all networks reachable.



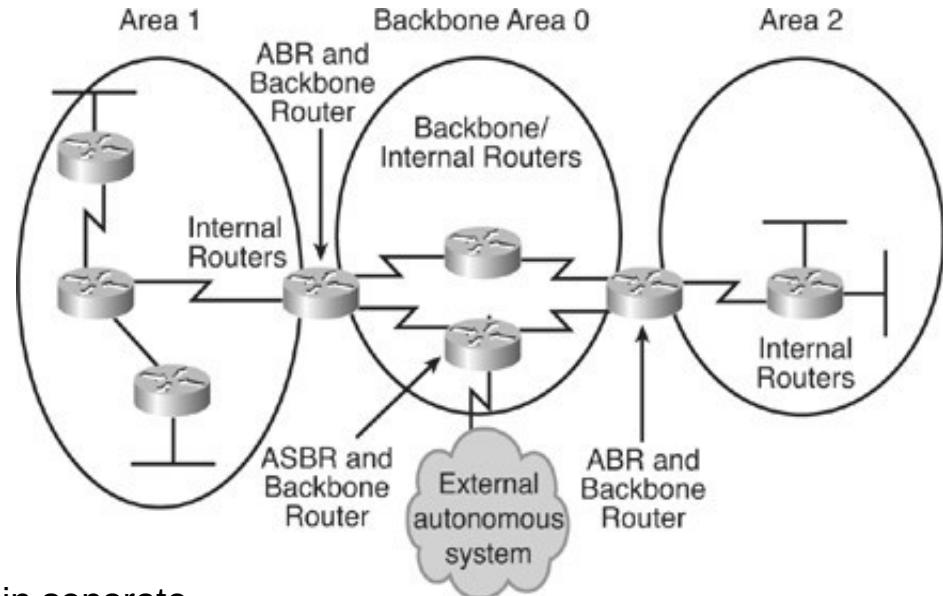
OSPF Two-Layer Area Hierarchy

- OSPF uses a two-layer area hierarchy:
- Backbone area
 - ◆ An OSPF area whose primary function is the fast and efficient movement of IP packets.
 - ◆ The backbone area interconnect with all other OSPF areas.
 - ◆ Generally, end users are not found within a backbone area.
 - ◆ The backbone area is also called OSPF area 0.
 - ◆ Hierarchical networking defines area 0 as the core to which all other areas connect (directly or virtually).
- Regular (non backbone) area
 - ◆ An OSPF area whose primary function is to connect users and resources.
 - ◆ Regular areas (also called normal areas) are usually set up along functional or geographic groupings.
 - ◆ By default, a regular area does not allow traffic from another area to use its links to reach other areas.
 - ◆ By default, all traffic from other areas must cross backbone area 0.
 - ◆ Regular areas can have several subtypes, including standard area, stub area, totally stubby area, not-so-stubby area (NSSA), and totally stubby NSSA.

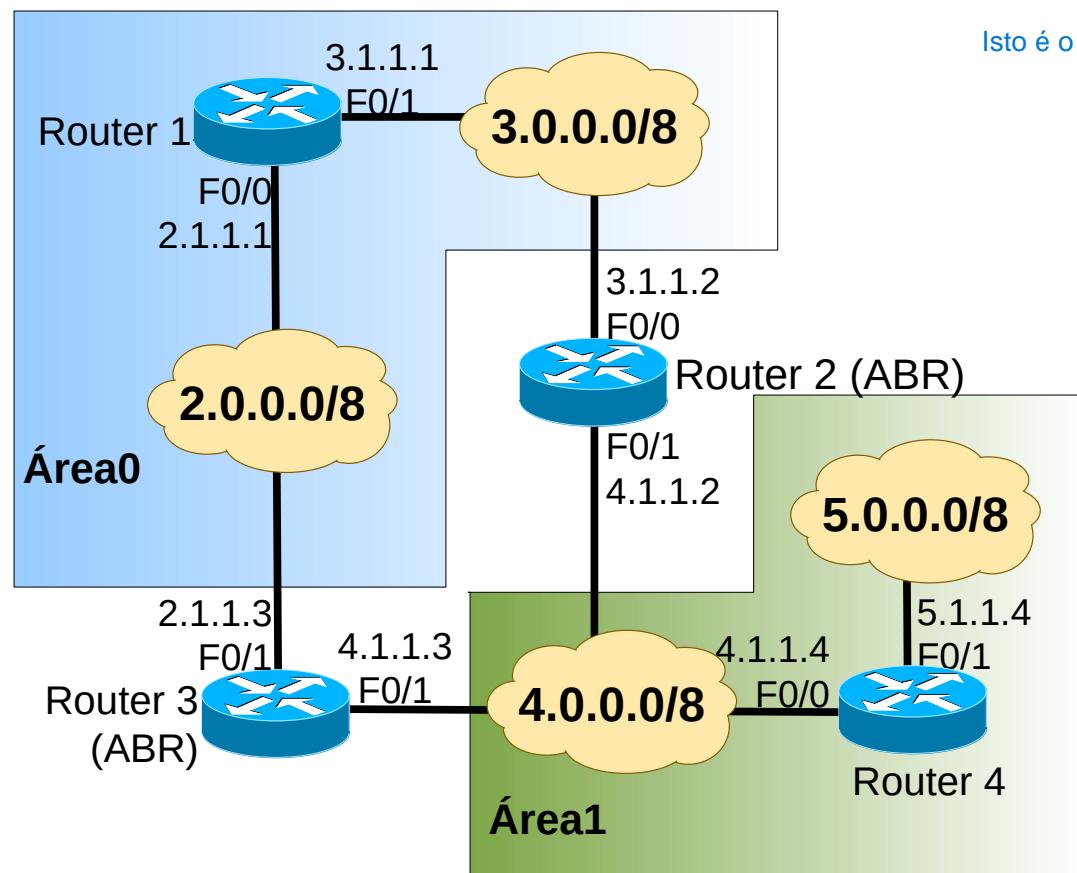


OSPF Routers Types

- Internal router
 - ◆ Routers that have all of their interfaces in the same area.
 - ◆ All routers within the same area have identical LSDBs.
- Backbone router
 - ◆ Routers that sit in the perimeter of the backbone area 0 and that have at least one interface connected to area 0.
 - ◆ Backbone routers maintain OSPF routing information using the same procedures and algorithms as internal routers.
- Area Border Router (ABR)
 - ◆ Routers that have interfaces attached to multiple areas, maintain separate LSDBs for each area to which they connect, and route traffic destined for or arriving from other areas.
 - ◆ Connect area 0 to a non backbone area and are exit points for the area
 - ◆ Routing information destined for another area can get there only via the ABR of the local area.
 - ◆ The ideal design is to have each ABR connected to two areas only, the backbone and another area.
 - ◆ The recommended upper limit is three areas.
- Autonomous System Boundary Router (ASBR)
 - ◆ Routers that have at least one interface attached to a different routing domain (such as another OSPF autonomous system or a domain using other routing protocol).
 - ◆ An OSPF autonomous system consists of all the OSPF areas and the routers within them.
 - ◆ ASBRs can redistribute external routes into the OSPF domain and vice versa.
- A router can be more than one router type.
 - ◆ For example, if a router interconnects to area 0 and area 1, and to a non-OSPF network, it is both an ABR and an ASBR.



OSPF Hierarchical Routing Example



Net Link States from Router 1

Isto é o que torna possível diminuir o tamanho da base de dados

| |
|-----------------------------------|
| Advertising Router: 4.1.1.2 |
| Number of Links: 1 |
| Router Interface address: 3.1.1.2 |
| TOS 0 Metrics: 10 |
| Advertising Router: 4.1.1.3 |
| Number of Links: 1 |
| Router Interface address: 2.1.1.1 |
| TOS 0 Metrics: 10 |
| Advertising Router: 4.1.1.3 |
| Number of Links: 1 |
| Router Interface address: 2.1.1.3 |
| TOS 0 Metrics: 10 |

Router Link States from Router 1

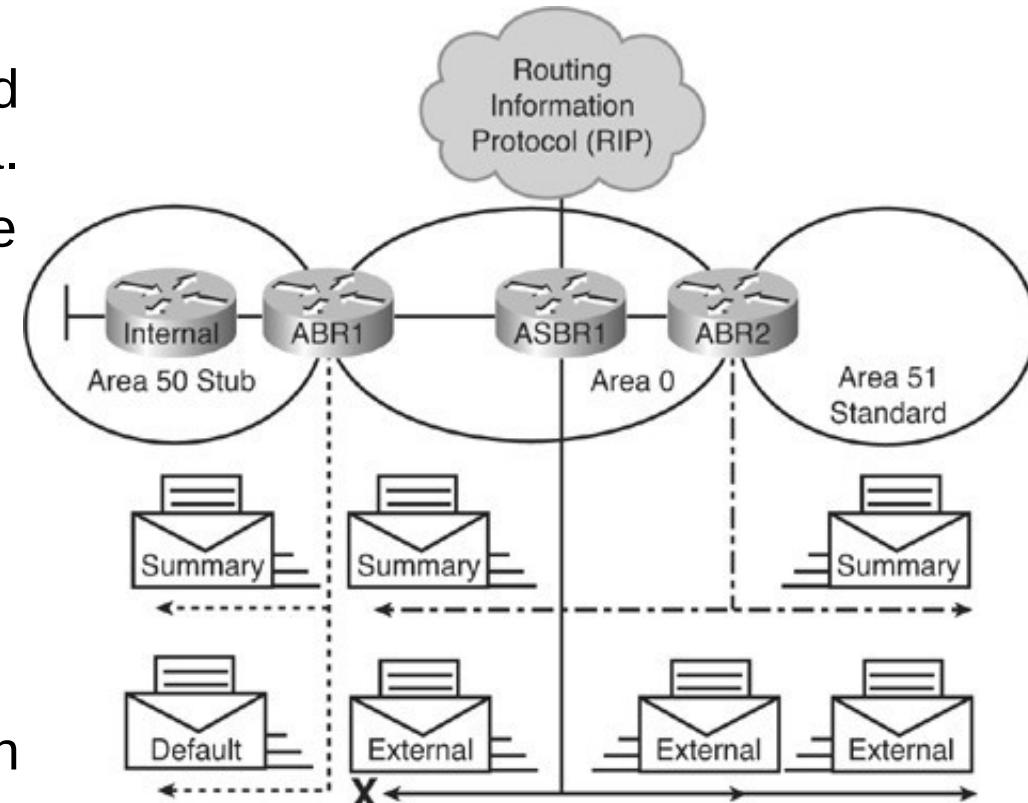
| | |
|-----------------------------|-----------------------------|
| Link State ID: 4.0.0.0 | Link State ID: 5.0.0.0 |
| Advertising Router: 4.1.1.2 | Advertising Router: 4.1.1.2 |
| Network Mask: /8 | Network Mask: /8 |
| TOS: 0 Metric: 10 | TOS: 0 Metric: 20 |
| Link State ID: 4.0.0.0 | Link State ID: 5.0.0.0 |
| Advertising Router: 4.1.1.3 | Advertising Router: 4.1.1.3 |
| Network Mask: /8 | Network Mask: /8 |
| TOS: 0 Metric: 10 | TOS: 0 Metric: 20 |

Summary Net Link States from Router 1



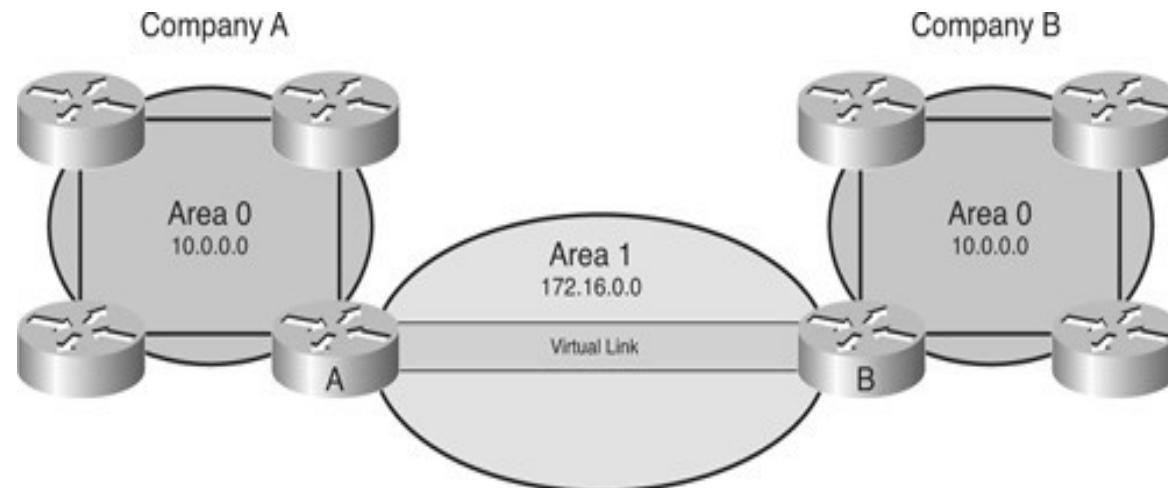
Stub Areas

- Configuring a stub area reduces the size of the LSDB inside an area, resulting in reduced memory requirements for routers in that area.
- Routers within the stub area also do not have to run the SPF algorithm as often because they will receive fewer routing updates.
- External network LSAs (type 5), such as those redistributed from other routing protocols into OSPF, are not permitted to flood into a stub area.
- Routing from these areas to a route external to the OSPF autonomous system is based on a default route (0.0.0.0).
 - Stub area ABR when receives an external LSA, sends a 0.0.0.0 LSA to the stub area.
 - If a packet is addressed to a network that is not in the routing table of an internal router, the router automatically forwards the packet to the ABR that originates a 0.0.0.0 LSA.

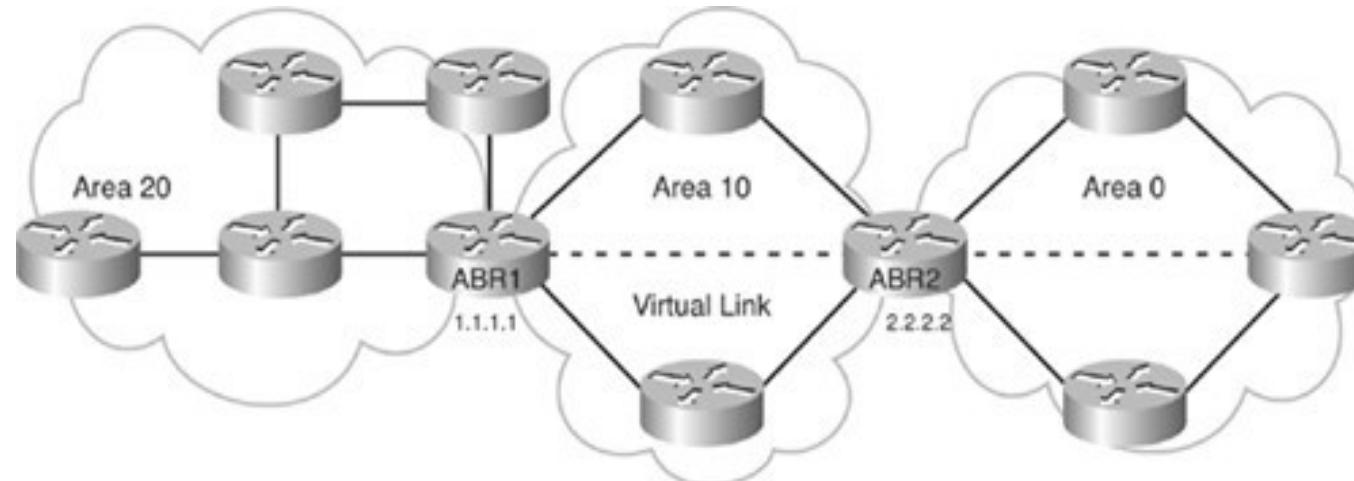


Areas Virtual Links

- Virtual Links can be used to connect a discontiguous Area 0.

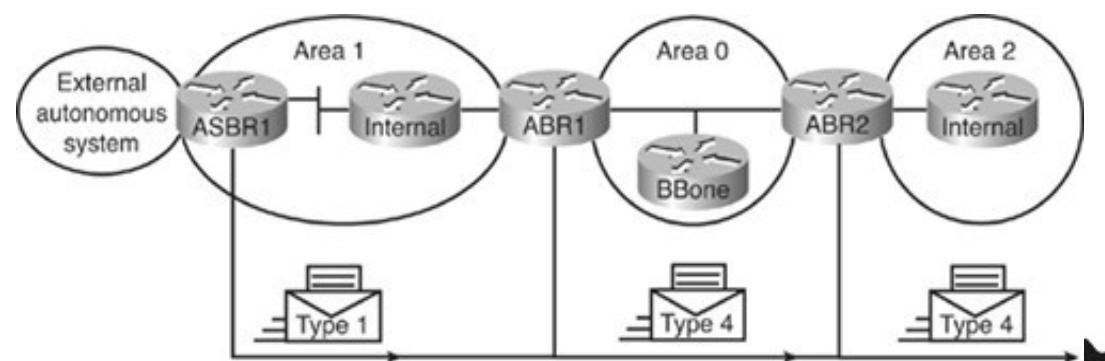
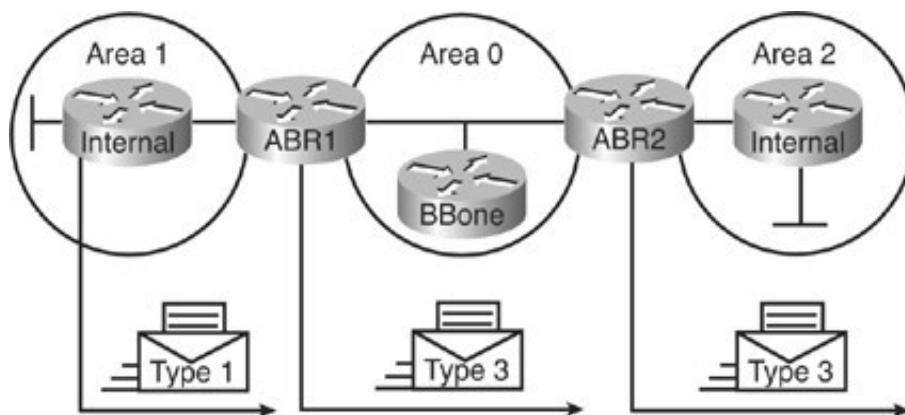


- Virtual Links can be used to connect an area to the backbone Area.



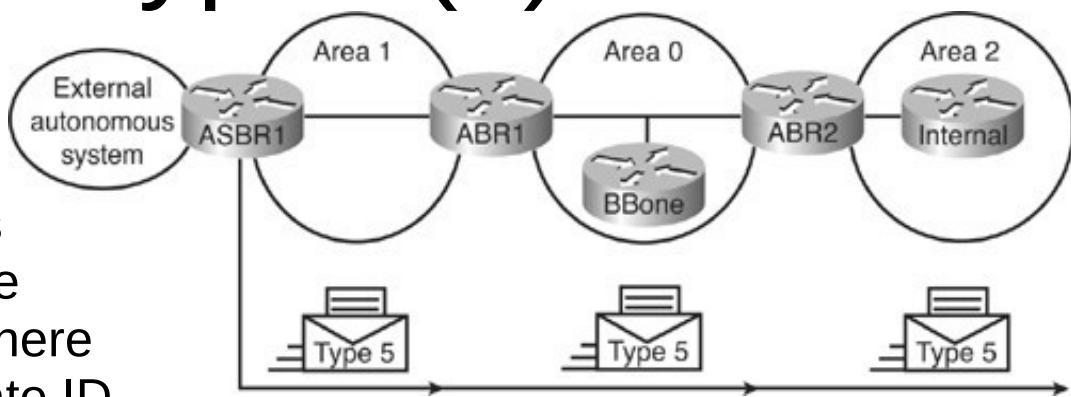
OSPF LSA Types (1)

- Type 1 (Router LSA) - Every router generates router-link advertisements for each area to which it belongs. Router-link advertisements describe the states of the router's links to the area and are flooded only within a particular area. All types of LSAs have 20-byte LSA headers. One of the fields of the LSA header is the link-state ID. The link-state ID of the type 1 LSA is the originating router's ID.
- Type 2 (Network LSA) - DRs generate network link advertisements for multiaccess networks, which describe the set of routers attached to a particular multiaccess network. Network link advertisements are flooded in the area that contains the network. The link-state ID of the type 2 LSA is the DR's IP interface address.
- Types 3 and 4 (Summary LSA) - ABRs generate summary link advertisements. Summary link advertisements describe the following inter-area routes
 - ◆ Type 3 describes routes to the area's networks (and may include aggregate routes) [Inter-Area Prefix LSA].
 - ◆ Type 4 describes routes to ASBRs [Inter-Area Router LSA].



OSPF LSA Types (2)

- Type 5 (AS external LSA) - ASBRs generate autonomous system external link advertisements. External link advertisements describe routes to destinations external to the autonomous system and are flooded everywhere except to any type of stub areas. The link-state ID of the type 5 LSA is the external network number.
- Type 6 (Multicast OSPF LSA) - These LSAs are used in multicast OSPF applications.
- Type 7 (LSAs for NSSAs) - These LSAs are used in NSSAs.
- Type 8 (External attributes LSA for BGP) - These LSAs are used to internetwork OSPF and BGP.
- Types 9, 10, or 11 (Opaque LSAs) - These LSA types are designated for future upgrades to OSPF for distributing application-specific information through an OSPF domain. Standard LSDB flooding mechanisms are used to distribute opaque LSAs. Each of the three types has a different flooding scope.
 - ◆ Type 9 LSAs are not flooded beyond the local network or subnetwork.
 - ◆ Type 10 LSAs are not flooded beyond the borders of their associated area.
 - ◆ Type 11 LSAs are flooded throughout the autonomous system (the same as for Type 5 LSAs). (Opaque LSAs are defined in RFC 5250, The OSPF Opaque LSA Option.)



Types of OSPF Routes

- OSPF intra-area (router LSA) and network LSA

- ◆ Networks from within the router's area, advertised by way of router LSAs and network LSAs.

- Inter-area (summary LSA)

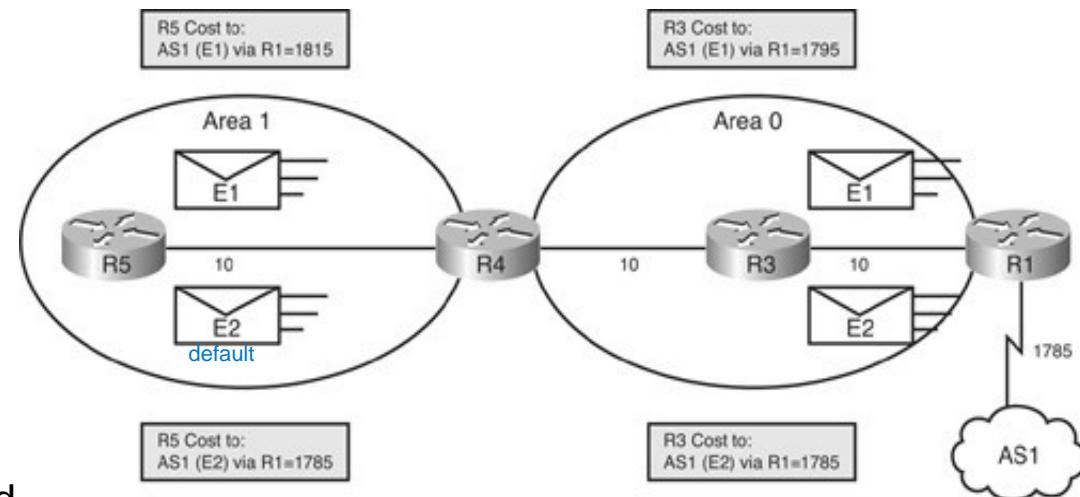
- ◆ Networks from outside the router's area but within the OSPF autonomous system, advertised by way of summary LSAs.

- Type 2 external routes (E2)

- ◆ Networks from outside the OSPF domain, advertised by way of external LSAs.
 - ◆ The cost of OSPF E2 routes is always the external cost only.
 - ◆ Use this type if only one ASBR is advertising an external route to the autonomous system.
 - ◆ This is usually the default for external routes.

- Type 1 external routes (E1) Este é preferível

- ◆ Networks from outside the OSPF domain, advertised by way of external LSAs.
 - ◆ Calculate the cost by adding the external cost to the internal cost of each link the packet crosses.
 - ◆ Use this type when multiple ASBRs are advertising an external route to the same autonomous system, to avoid suboptimal routing.
 - ◆ Always preferred over Type 2 external routes (E2). Even for higher metrics!



Falta a noção da distância física. Exemplo: Pôr um Router na UA e mandar os pacotes ir pra o servidor da Vodafone em Águeda pra depois voltar, quando temos o Router com ligação à internet ao lado



OSPF Area Types

Standard area

- ◆ This default area type accepts link updates, route summaries, and external routes.

Backbone area

- ◆ The backbone area is labeled area 0, and all other areas connect to this area to exchange and route information.
- ◆ The OSPF backbone has all the properties of a standard OSPF area.

Stub area

- ◆ Cannot contain ASBRs (except ABRs that may also be ASBRs).
- ◆ From Area 0 ABR, receives summary routes (LSA Type 3) and automatic default route. External routes (LSA Type 5) are blocked.

Totally stubby area

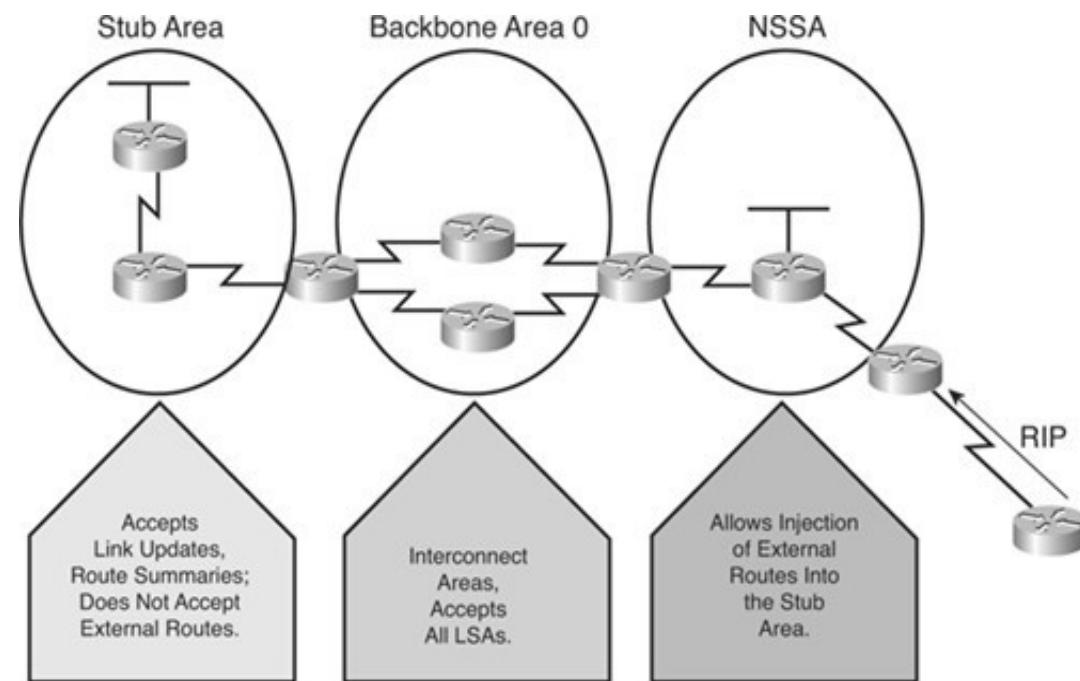
- ◆ Cisco proprietary area type.
- ◆ Cannot contain ASBRs (except ABRs that may also be ASBRs).
- ◆ From Area 0 ABRs, receives automatic default route. Summary routes (LSA Type 3) and external routes (LSA Type 5) are blocked.

NSSA (Not So Stubby Area)

- ◆ Is an addendum to the OSPF RFC.
- ◆ Contain ASBRs (receives external routes).
- ◆ From Area 0 ABRs, receives summary routes (LSA Type 3). External routes (LSA Type 5) are blocked.
- ◆ No automatic default route is sent to NSSA Area by ABR.
- ◆ Uses a LSA type 7 to announce external routes to Area 0 ABR, ABR transforms the LSA Type 7 into a LSA type 5 and sends it to Area 0.

Totally stubby NSSA

- ◆ Contain ASBRs (receives external routes).
- ◆ From Area 0 ABRs, receives automatic default route. Summary routes (LSA Type 3) and external routes (LSA Type 5) are blocked.
- ◆ Uses a LSA type 7 to announce external routes to Area 0 ABR, ABR transforms the LSA Type 7 into a LSA type 5 and sends it to Area 0.



Routing - OSPFv3

- Based on OSPFv2, with enhancements:
 - ◆ Uses IPv6 for transport
 - ◆ Distributes IPv6 prefixes
 - ◆ Uses multicast group addresses FF02::5 (OSPF IGP) and FF02::6 (OSPF IGP Designated Routers)
 - ◆ Runs over a link rather than a subnet
 - ◆ Multiple instances per link
 - ◆ Topology not IPv6-specific
 - ✚ Router ID, Area ID, Link ID remain a 4 bytes number
 - ✚ Neighbors are always identified by Router ID (4 bytes)
 - ✚ With an additional table with mapping between IPv6 prefixes and Link IDs
 - ◆ Uses link-local addresses as IPv6 source addresses



OSPFv3 - LSA Types

- Link LSA (Type 8)
 - ◆ Informs neighbors of link local address
 - ◆ Informs neighbors of IPv6 prefixes on link
- Intra-Area Prefix LSA (Type 9)
 - ◆ Associates IPv6 prefixes with a network or router
- Flooding scope for LSAs has been generalized
 - ◆ Three flooding scopes for LSAs
 - ◆ Link-local
 - ◆ Area
 - ◆ AS
- LSA Type encoding expanded to 16 bits
 - ◆ Includes flooding scope



Integrated System-Integrated System (IS-IS) Protocol

é mais flexível

- IS-IS was defined in 1992 in the ISO/IEC recommendation 10589.
- IS-IS is a link-state routing protocol.
 - ◆ Provides fast convergence and excellent scalability.
 - ◆ Very efficient in its use of network bandwidth.
- Uses Dijkstra's Shortest Path First algorithm (SPF).
- Types of packets
 - ◆ IS-IS Hello packet (IIH), Link State Packet (LSP), Partial Sequence Number Packet (PSNP) and Complete Sequence Number Packet (CSNP).
- Link States are called LSPs
 - ◆ Contain all information about one router adjacencies, connected IP prefixes, OSI end systems, area addresses, etc.
 - ◆ One LSP per router (plus fragments).
 - ◆ One LSP per LAN network.
- IS-IS has 2 layers of hierarchy
 - ◆ The backbone is called level-2.
 - ◆ Areas are called level-1.
 - ◆ A router can take part in L1 and L2 inter-area routing (or inter-level routing).

é subjetivo se este protocolo é melhor do que OSPF.
O OSPF é mais usado, exceto quando se trata de
operadoras em grandes cidades, onde é preciso
mudar o core, este processo é mais fácil/pragmático

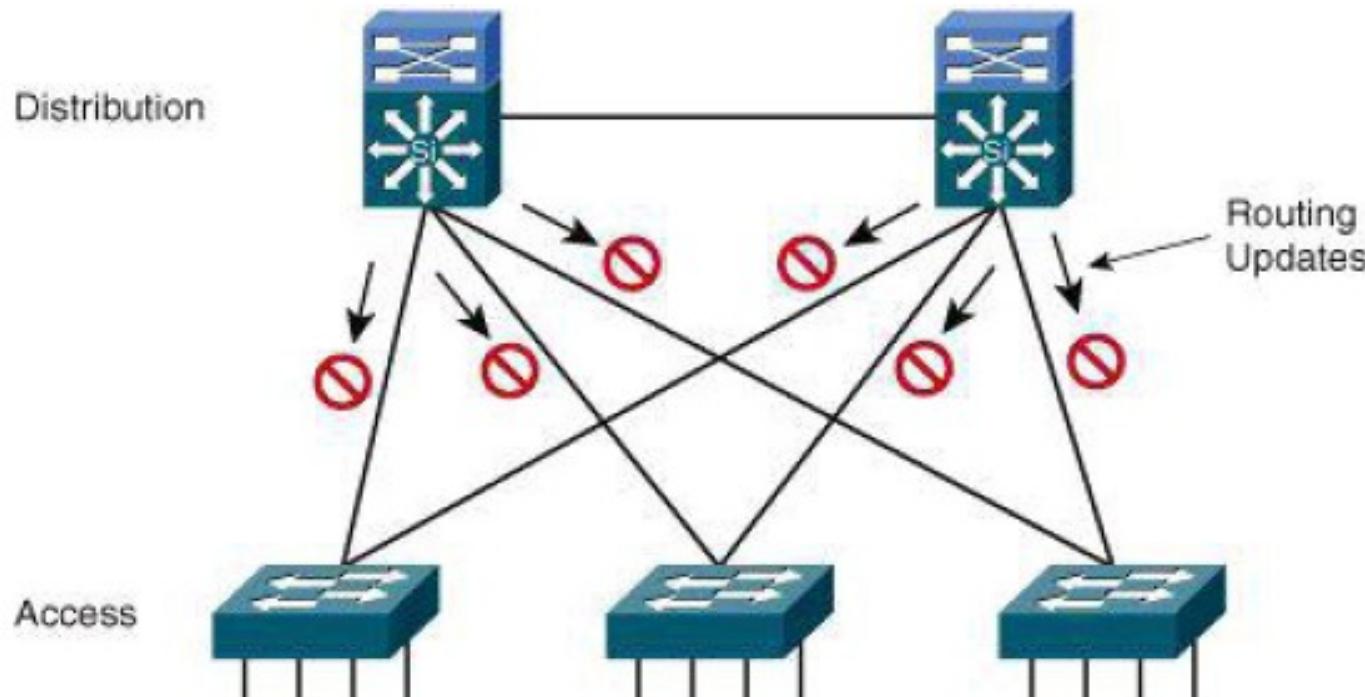


Enhanced Interior Gateway Routing Protocol (EIGRP) Protocol

- EIGRP is a Cisco-proprietary protocol that combines the advantages of link-state and distance vector routing protocols.
- EIGRP has its roots as a distance vector routing protocol and is predictable in its behavior.
- What makes EIGRP an advanced distance vector protocol is the addition of several link-state features, such as dynamic neighbor discovery.
 - ◆ EIGRP Maintains a Neighbor Table, a Topology Table, and a Routing Table.
- EIGRP has Variable-length subnet masking (VLSM) support.
- Has a sophisticated metric that considers five criteria:
 - ◆ Two by default:
 - ◆ Bandwidth - The smallest (slowest) bandwidth between the source and destination.
 - ◆ Delay - The cumulative interface delay along the path.
 - ◆ Available, are not commonly used, because they typically result in frequent recalculation of the topology table:
 - ◆ Reliability - The worst reliability between the source and destination, based on keepalives.
 - ◆ Loading - The worst load on a link between the source and destination based on the packet rate and the interface's configured bandwidth.
 - ◆ Maximum transmission unit (MTU) - The smallest MTU in the path.
- A significant advantage of EIGRP (and IGRP) over other protocols is its support for unequal metric load balancing that allows administrators to better distribute traffic flow in their networks.



Passive Interfaces on Access Layer

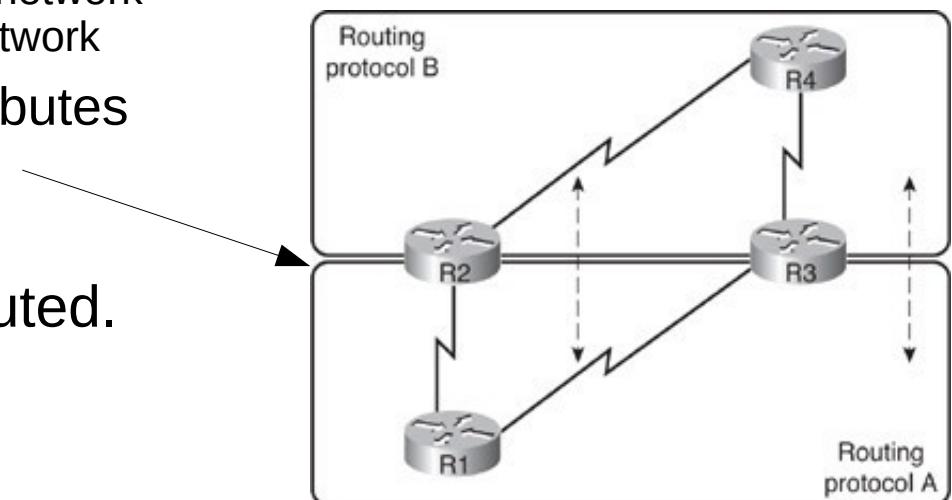
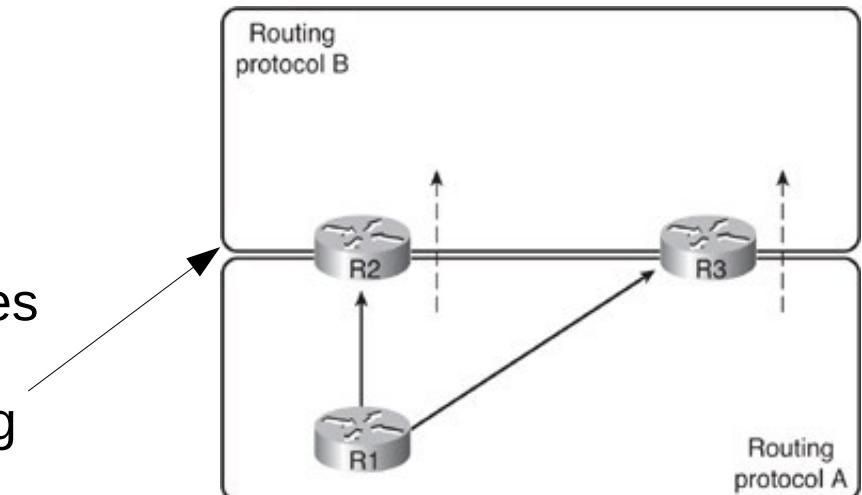


- As a recommended practice, limit unnecessary L3 routing peer adjacencies by configuring the ports toward Layer 2 access switches as passive.
 - ◆ Suppress the advertising of routing updates.
 - ◆ If a distribution switch does not receive L3 routing updates from a potential peer on a specific interface, it does not form a neighbor adjacency with the potential peer across that interface.



Route Redistribution

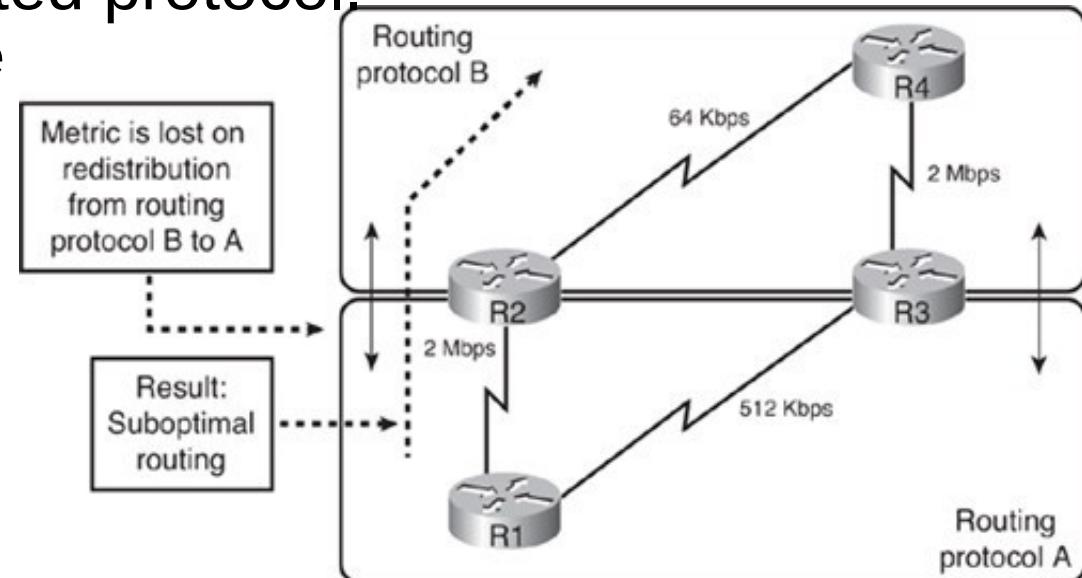
- Domains with different routing protocols can exchange routes.
 - ◆ This is called route redistribution.
 - ◆ One-way redistribution - Redistributions only the networks learned from one routing protocol into the other routing protocol.
 - Uses a default or static route so that devices in that other part of the network can reach the first part of the network
 - ◆ Two-way redistribution - Redistributions routes between the two routing processes in both directions
 - ◆ Static routes can also be redistributed.



Redistribution Issues

- Lost metric from redistributed protocol.

- ◆ It is not possible to achieve an optimal overall routing.



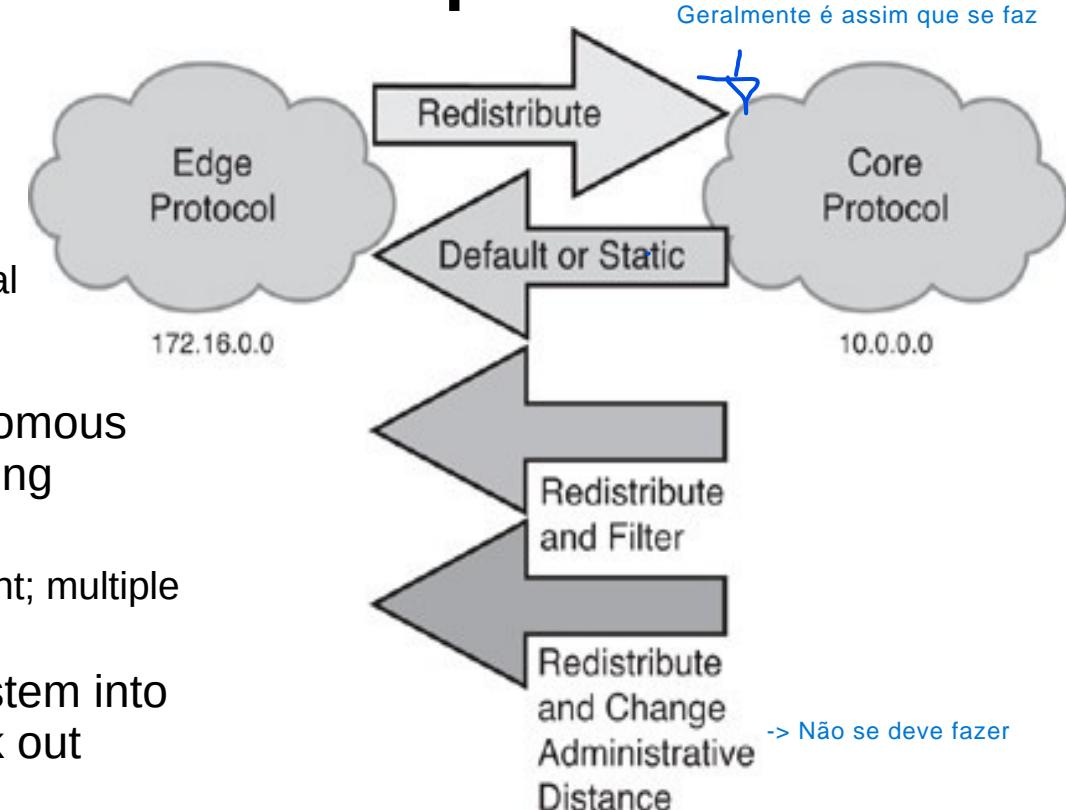
- Preventing Routing Loops in a Redistribution Environment.

- ◆ Safest way to perform redistribution is to redistribute routes in only one direction, on only one boundary router within the network.
 - ◆ However, that this results in a single point of failure in the network.
 - ◆ If redistribution must be done in both directions or on multiple boundary routers, the redistribution should be tuned to avoid problems such as suboptimal routing and routing loops.



Redistribution Techniques

- Redistribute a default route from the core autonomous system into the edge autonomous system, and redistribute routes from the edge routing protocols into the core routing protocol.
 - ◆ This technique helps prevent route feedback, suboptimal routing, and routing loops.
- Redistribute multiple static routes about the core autonomous system networks into the edge autonomous system, and redistribute routes from the edge routing protocols into the core routing protocol.
 - ◆ This method works if there is only one redistribution point; multiple redistribution points might cause route feedback.
- Redistribute routes from the core autonomous system into the edge autonomous system with filtering to block out inappropriate routes.
 - ◆ For example, when there are multiple boundary routers, routes redistributed from the edge autonomous system at one boundary router should not be redistributed back into the edge autonomous system from the core at another redistribution point.
- Redistribute all routes from the core autonomous system into the edge autonomous system, and from the edge autonomous system into the core autonomous system, and then modify the administrative distance associated with redistributed routes so that they are not the selected routes when multiple routes exist for the same destination.

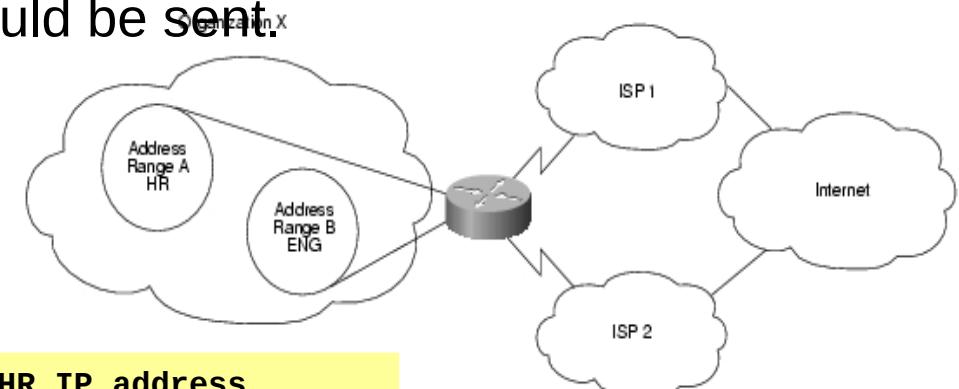


Do que está mais na ponta da rede, manda para o mais central (core)



Policy-Based Routing (PBR)

- PBR allows the operator to define routing policy other than basic destination-based routing using the routing table.
- PBR rules can be used to match source and destination addresses, protocol types, and end-user applications.
- When a match occurs, a set command can be used to define the interface or next-hop address to which the packet should be sent.



```
access-list 1 permit 209.165.200.225
access-list 2 permit 209.165.200.226
!
interface ethernet 1
  ip policy route-map ChooseISP
!
route-map ChooseISP permit 10
  match ip address 1
  set ip next-hop 209.165.200.227
!
route-map ChooseISP permit 20
  match ip address 2
  set ip next-hop 209.165.200.228
```

!From HR IP address
!From ENG IP address

Defines order of the rules

!To ISP2 next-hop

!To ISP1 next-hop

O encaminhamento depende dos endereços.
Exemplo: ter um conjunto específico de servidores em Aveiro e só esses endereços devem ir a uns servidores específicos no Brasil

Como se fosse um if, se se cumprir a condição, manda-se/força-se os pacotes a ir por um determinado tunel



Podemos atribuir duas Gateways a um PC, a maioria dos PCs dão mais peso a uma gw em detrimento da outra. O ideal seria que mande para essa gw principal, e se falhar, mandar para a outra, mas mts vezes isso não acontece, o Router continua a mandar para o primeiro, independentemente de estar a funcionar ou não

Virtual Router Redundancy Protocol

- VRRP is a standard protocol defined by the IETF in RFC 3768 to create a virtual gateway.
 - ◆ Cisco has HSRP (Hot Standby Routing Protocol) which is very similar.
- A cluster of routers can be handled as a single router using a
 - ◆ Virtual IP address and virtual MAC address.
 - ◆ Default gateway to the clients.
 - ◆ One of the individual routers acts as master (working router), however, upon a failure one of the other individual routers becomes the working router.
 - ◆ Router states are maintained and verified using a multicast group.

Estes 3 Routers vão partilhar um endereço virtual (10.2.2.1). Eles vão comunicar entre si. Este serviço é para: Caso o PC vai abaixo, o outro assume o IP e o MAC.

Do ponto de vista do terminal nada muda

