## Exam 2

## CSc 75010: Theoretical Computer Science Graduate Center of CUNY 22 November 2002 (Sample Exam)

Do five of the following six problems. Write each answer on a separate piece of paper.

- 1. (a) Define decidable set and give an example of a set that is not decidable.
  - (b) Assume that the alphabet and tape alphabet are:  $\Sigma = \Gamma = \{0, 1\}$ . Give an implementation level description of a Turing machine that decides the language:

 $\{w \mid w \text{ contains an equal number of 0's and 1's}\}$ 

- (c) Define decidable function and give an example of a function that is not decidable.
- (d) Assume that the alphabet and tape alphabet are:  $\Sigma = \Gamma = \{0, 1\}$ . Give an implementation level description of a Turing machine that copies the input string on the tape. That is, if the input to the machine is the string w, the output is ww.
- (e) Define decidable set and give an example of a set that is not decidable.
- (f) Assume that the alphabet and tape alphabet are:  $\Sigma = \Gamma = \{0, 1\}$ . Give an implementation level description of a Turing machine that decides the language:

 $\{w \mid w \text{ is a palindrome}\}$ 

- (g) Define decidable function and give an example of a function that is not decidable.
- (h) Assume that the alphabet and tape alphabet are:  $\Sigma = \Gamma = \{0, 1\}$ . Give an implementation level description of a Turing machine that takes as input a number in binary representation and mulitplies it by 2 (that is, shifts it to the right).
- 2. (a) Show that for any finite set  $\Sigma$ , the set of all finite strings of  $\Sigma$ ,  $\Sigma^*$  is countable.
  - (b) Let  $\Sigma = \{0, 1\}$ . Show that the set of all infinite strings over  $\Sigma$  is uncountable.
  - (c) Show that the set of all positive rational numbers is countable.
  - (d) Show that the set of all positive real numbers is uncountable.
- 3. (a) Show that the Halting Problem is undecidable, using the diagonalization method.
  - (b) Show that the set of all finitely defined computable functions (that is, the set of all functions that diverge for all but a finis not decidable.
  - (c) Show that the Halting Problem is undecidable, using the diagonalization method.
  - (d) Show that the set of all constant computable functions is not decidable.
  - (e) Show that the Halting Problem is undecidable, using the diagonalization method.
  - (f) Show that the set of all totally undefined computable functions (that is the set of all functions that diverge for all inputs) is not decidable.
  - (g) Show that the Halting Problem is undecidable, using the diagonalization method.

- (h) Show that the set of all totally defined computable functions is not decidable.
- 4. (a) Show that the set of decidable languages is closed under concatenation.
  - (b) Show that if A and  $\bar{A}$  are Turing-recognizable, then A is decidable.
  - (c) Show that the set of Turing-recognizable languages is closed under star.
  - (d) Show that if A is decidable, then A and  $\bar{A}$  are Turing-recognizable.
  - (e) Show that the set of decidable languages is closed under intersection.
  - (f) Show that if A and  $\bar{A}$  are Turing-recognizable, then A is decidable.
  - (g) Show that the set of decidable languages is closed under union.
  - (h) Show that if A and  $\bar{A}$  are Turing-recognizable, then A is decidable.
  - (i) Show that the set of Turing-recognizable languages is closed under concatenation.
  - (j) Show that if A is decidable, then A and  $\bar{A}$  are Turing-recognizable.
- 5. (a) State the Post Correspondence Problem (PCP).
  - (b) Show that PCP is decidable over the alphabet  $\Sigma = \{1\}$ .
  - (c) State Rice's Theorem.
  - (d) Prove Rice's Theorem.
- 6. (a) If  $A \leq_m B$  and B regular does that imple A is regular? Why or why not? Justify your answer.
  - (b) Show for all A, B, there exists a set J such that  $A \leq_T J$  and  $B \leq_T J$ .
  - (c) Show that  $\leq_m$  is transitive.
  - (d) Show for all A, B, there exists a set J such that  $A \leq_T J$  and  $B \leq_T J$ .
  - (e) Show that if A Turing-recognizable and  $A \leq_m \bar{A}$ , then A is decidable.
  - (f) Show that if  $A \leq_T B$  and  $B \leq_T C$  implies  $A \leq_T C$ .
  - (g) Show that for all A and B,  $A \leq_m B$  implies  $\bar{A} \leq_m \bar{B}$ .
  - (h) Show for all A, B, there exists a set J such that  $A \leq_T J$  and  $B \leq_T J$ .