

CSC 488S/CSC 2107S Lecture Notes

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Reading Assignment

Section 10.1

Chapter 7

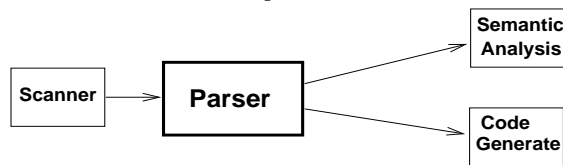
Omit 7.7

Emphasize AST use and structure

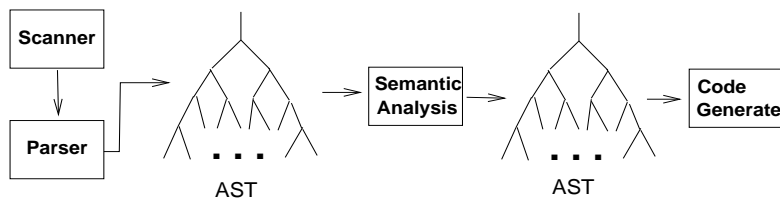
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Single Pass Compiler



Multi Pass Compiler



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Internal Compiler Data Structures

- Information flows between compiler passes
 - Representation(s) of the program
 - Tables
 - Error messages
 - Compiler flags
 - Source program coordinates.
- Form of communication may change as program is processed. A compiler may use multiple representations of a program.
- Use disk resident information for very large programs. Use memory resident information for better compiler speed.
- **Backward information flow should be avoided if possible.**

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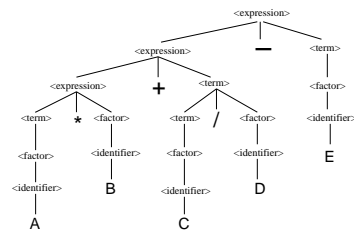
Intermediate Representations

- Represent the structure of the program being compiled
 - Declaration structure.
 - Scope structure.
 - Control flow structure.
 - Source code structure.
- Used to pass information between compiler passes
 - Compact representation desirable.
 - Should be efficient to read and write.
 - Provide utility to print intermediate language for compiler debugging.

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Parse Tree

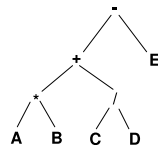
Complete representation of the syntactic structure of the program according to some grammar.



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Abstract Syntax Tree

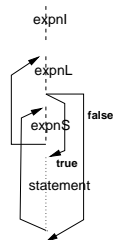
Similar to parse tree but only describes essential program structure.



Directed Acyclic Graphs

Used to represent control structure

for (expnI ; expnL ; expnS) statement



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Intermediate Representation Examples

Condensed Source

A * B + C / D - E

Polish Postfix Notation

A B * C D / + E -

Triples

```
1 ( * , A , B )
2 ( / , C , D )
3 ( + , (1) , (2) )
4 ( - , (3) , E )
```

Quadruples

```
( * , A , B , T1 )
( / , C , D , T2 )
( + , T1 , T2 , T3 )
( - , T3 , E , T4 )
```

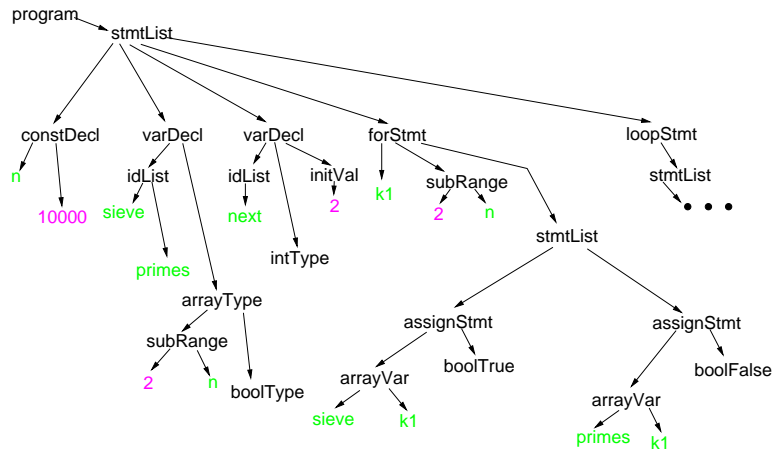
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Abstract Syntax Trees

- Abstract Syntax Trees (ASTs) are designed to capture all of the essential structural information about a program being compiled.
- Bottom Up syntax analysis is a good match with AST building. The child nodes to the tree are always built before the parent node.
- The basic process for *building* an AST is very simple:
 - the leaf nodes are typically constants and identifiers. Build a node to represent each of these entities.
 - Interior nodes are build as needed to represent particular language constructs. Typically interior nodes contain links to one or more child nodes.
- The *processing* of AST nodes for semantic analysis is also simple:
 - Process the AST *depth first*. This guarantees that child nodes are processed before parent nodes.
 - At each parent node do any processing required using information from already processed child nodes.

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Partial^a Abstract Syntax Tree for Running Example (Slide 25)

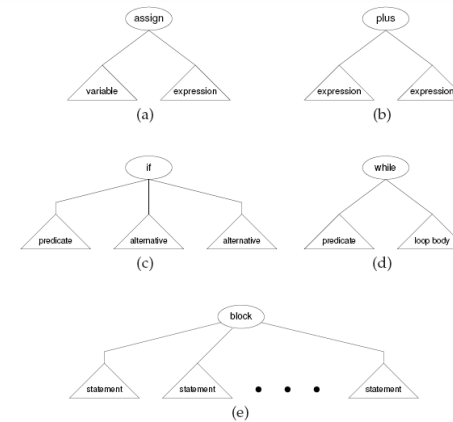


^aCompare this to the full parse tree in Slide 58

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ASTs Are Language Specific

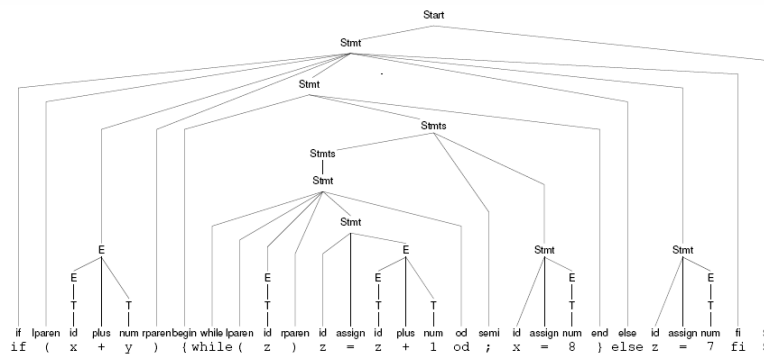
Typically AST nodes are custom designed for each of the major constructs in the language. Some Examples (Fischer, Cytron, LeBlanc Figure 7-15)^a:



^aFor CSC488S Winter 2015/2016 examples see the AST classes: AssignStmt, BinaryExpn, IfStmt, WhileDoStmt, and Scope

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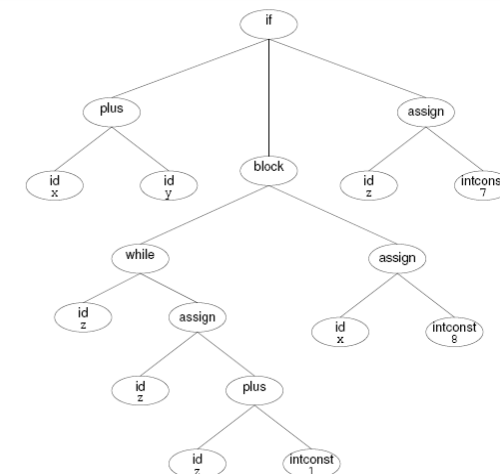
Complete Syntax Tree for Program Fragment^a



^a Fischer, Cytron, LeBlanc Figure 7-18

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Abstract Syntax Tree for Program Fragment^a



^a Fischer, Cytron, LeBlanc Figure 7-19

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A Very General Abstract Syntax Tree Node^a

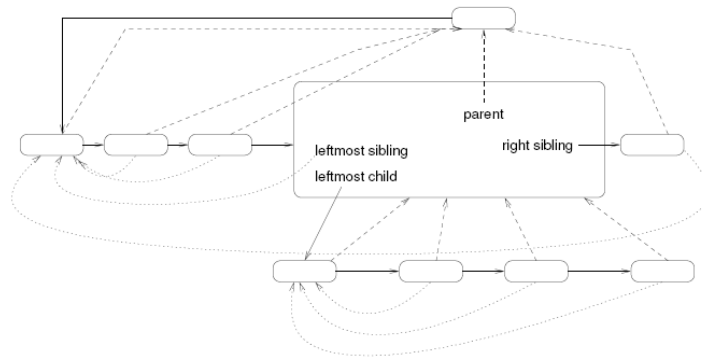


Figure 7.12: Internal format of an AST node. A dashed line connects a node with its parent; a dotted line connects a node with its leftmost sibling. Each node also has a solid connection to its leftmost child and right sibling.

^a Fischer, Cytron, LeBlanc Figure 7-12