A one phase IPM for non-convex optimization



Hey, my name is Oliver and today I will be talking about work in progress with Yinyu Ye on developing a one phase method for non-convex optimization.

There are three main goals for this work:

- Is to develop a simple one phase interior point method for non-convex optimization inspired by the highly successful infeasible start IPMs for convex optimization
- 2. Improve and clarify the convergence theory for non-convex interior point methods.
- 3. Improve reliability of non-convex interior point codes, particularly on infeasible or degenerate problems.

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intline

Introduction

Introdu

—Outline

This talk is split into three sections.

- 1. In first section, we review current non-convex IPM and explain why one phase algorithms are currently not used for non-convex optimization. We then outline our algorithm.
- In the second section, we present our convergence results notably, we do not require constraint qualifications for either global or local superlinear convergence.
- 3. In the third sections, we briefly discuss implementation details and present preliminary CUTEst results.

—Problem we wish to solve:

So the problem we want to solve has general non-linear objective and inequality constraints, we can, of course, represent any equality constraint by two inequalities.

We assume that:

- 1. The functions have 1st and 2nd derivatives everywhere
- The feasible region is bounded for any perturbation of the right hand side. If necessary, one can add dummy constraints to guarantee this.

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—One phase IPMs for convex optimization

One phase IPMs for convex optimization

Either declare problem infastible (primal or dual) or optimal in one phase

o Infastible start IPM (Mahorta, 1922)

or Homogeness aggloratin (Anderson and Vs. 1958, 1999)

State of the art for convex optimization

One phase algorithms have been successful in linear programming and convex optimization. Starting from an infeasible point they declare a problem optimal or infeasible. At each iteration one fixed newton system is used and the right hand side to decide the target reduction in primal, dual and complementary to get quick convergence.

The two types most successful one phase algorithms are the infeasible start IPM that proves infeasibility through the divergence of dual iterates and the homogenous algorithm that proves infeasibility by convergence.

A one phase IPM for non-convex optimization Introduction

Why one phase non-convex IPM do not exist (yet)

© Some infeasible start algorithms have been tested, but no theoretical guarantees (i.e. LOQO)

 Ideally our algorithm would either find a local optimum or a local minimizer the violation of the constraints

 Infeasible start IPMs may fail this criterion (Wachter and Biegler in 2000).

—Why one phase non-convex IPM do not exist (yet)

- 1. Ideally our algorithm would either find a local optimum or a local minimizer of the constraint violation
- In 2000, Wachter and Biegler gave an example of a problem for which infeasible start algorithm may fail to converge. I think this paper influenced many IPMs including IPOPT and KNITRO.

A one phase IPM for non-convex optimization
Introduction
Current non-convex IPMs
Solutions in literature to this problem

Solutions in literature to this problem

Too places (POPT)
Gody place has a first excision
Gody place has a first excision
Once preference on infeatile publics
O compute two directions (NMTRO)
O Pleasity (or big M) method (Curtis, 2012)

- 1. Two phases (IPOPT): main phase and feasibility restoration phase
 - 1.1 Each phase has a different variables: each time initialize FRP start from scratch, the dual variables = zero, new introduced primal variables have ad hoc values.
 - 1.2 Awkward convergence assumptions: to prove convergence it is assumed constraints are linearly independent in a neighborhood of the feasible region.
 - 1.3 Poor performance on infeasible problems.
- 2. (KNITRO) using two different newton systems
 - 2.1 One direction aims to minimize L2 norm of constraint violation, other step aims to improve optimality
 - 2.2 Not 'pure' IPM in feasibility search, requires factorizing two different newton systems at each step.

1. Curtis use L1 minimization

... Now, as you can see there are a lot of complicated modifications to the infeasible start IPM to guarantee convergence.

It doesn't need to be this complicated!

In this talk, I will show despite the fact an infeasible start algorithm fails with equality constraints, if we:

- 1. Use inequality constraints instead equality constraints
- 2. Adjust the target rate of reduction of primal and complementary

Then an infeasible start algorithm can be proved to converge to a farkas certificate of infeasibility or local optimality certificate. Furthermore, if we add dummy constraints, this farkas certificate, becomes a certificate of L1 minimization of the constraints.

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Introduction

Local infeasibility certificates

Local infeasibility certificates



In convex optimization infeasibility detection is simple. We produce a farkas certificate, that proves the problem is infeasible. In non-convex optimization we cannot guarantee global infeasibility. However, there are many possible types of local infeasibility certificates.

For example, we could produce a farkas certificate of infeasibility, which is a non-negative weight vector y and decision variable x^* such that x^* is a stationary point of $y^T a(x)$ with $y^T a(x^*)$ positive.

A stronger infeasibility criterion would be if x^* was a stationary point of the sum of the constraint violation.

While there are other possible many other possible criterion we could use, in this talk we will focus on these two.



Now we are ready to start describing our algorithm.

Recall the original problem we wish to solve with a general non-linear objective and inequality constraints.

Our idea is to modify this problem in two ways. Firstly, we add in a log barrier term to penalize getting too close to the boundary of the constraints, this is a very typical strategy. Secondly, we set a target for reducing the constraint violation. It is this second feature, which is common in convex codes, but uncommon in non-convex codes, which is critical to our result.

At each iteration we reduce the penalty parameter μ^k and the primal feasibility by the same rate η^k .

Next, we take the barrier problem and derive a typical primal-dual interior point direction. We modify the inertia of the hessian matrix to ensure it is sufficiently positive definite. Observe that by changing η^k which only appears in the right hand side we can change the target reduction in primal feasibility and complementary.

On this slide we present an outline of our algorithm. At each iteration, we either take an aggressive step or a stabilization step. If the dual feasibility is sufficiently small, then we take an aggressive step. In an aggressive step we reduce μ^k and the primal feasibility by η^k . We then take the largest step in that direction while maintaining complementary.

Otherwise, we take a stabilization step where we keep the primal feasibility and complementary the same. We perform a backtracking search on the direction using a merit function.

Given a direction, and a step size α^k this slide describes how we update the iterates. We update the x and y variables in a typical fashion, however, the s variables and barrier parameter μ are updated to ensure that the constraint violation and μ are reduced by **exactly** α^k times η^k .

Furthermore, we reject any iterates that do not approximately satisfy complementary.

During the stabilization step, we use the following, typical, barrier merit function to ensure convergence. We use a backtracking line search to take the largest step that makes sufficient progress on this merit function. During the aggressive step, we do not use a merit function.

Finally, there are two ways our algorithm can terminate.

- 1. It satisfies the KKT conditions to sufficient degree of accuracy
- 2. It finds a Farkas certificate of infeasibility

Now, assuming δ^k and η^k are carefully chosen then it possible to show that our algorithm either converges is a first-order ϵ -local optimum or a first-order ϵ -local farkas infeasibility certificate. It is worth noting that this result does not say anything about whether the sequence of dual multipliers is bounded.

Our one phase algorithm has nice convergence properties. If the algorithm is converging towards a KKT point that satisfies strict complementary and the sufficient conditions for local optimality are satisfied at this point, then the algorithm converges Q-quadratically to a local optimality certificate. Observe that this result is somewhat unusual since we require minimal assumptions and no constraint qualifications.

So far we have shown that if our algorithm does not find a local optimum it finds a farkas certificate of local infeasibility.

What if we would like a different type of infeasibility certificate? For example, if we want to minimize the sum of constraint violation then it is sufficient to embed the original problem inside a slightly larger problem.

Assuming the slack variables are appropriately initialized then the farkas certificate of local infeasibility for the larger problem will imply that we have minimized the sum of constraint violation.

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Practice

Implementation details

Implementation details

mplementation details

- Use fraction to boundary rule i.e. y^{k+1} ≥ 0.01y^k and s^{k+1} ≥ 0.01s^k
 Choose η^k on aggressive steps using predictor-corrector
- Search on negative eigenvector to guarantee second-order necessary conditions are satisfied

Our implemented algorithm, is more sophisticated than the simple algorithm described at the beginning of the talk. We are in the early stages and the design decisions are still fluid. I am just going to briefly discuss a few implementation details:

- 1. To prevent iterates use converging too quickly to the boundary, we only accept iterates that satisfy the fraction to the boundary criterion.
- 2. For the theoretical proofs we never give a practical choice for η^k , in practice, we use a predictor-corrector technique. First, we compute take an aggressive step with $\eta^k=1$, and on aggressive steps using predictor-corrector technique
- 3. We use the factorized matrix as a pre-conditioner to generate negative eigenvectors. We search on this negative eigenvector to guarantee second-order necessary conditions are satisfied.

Using a filter to prevent cycling

We have found that the practical performance of the algorithm can be improved substantially by adding a second merit function during the stabilization step.

Preliminary empirical results with CUTEst
Empirical results

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On this slide we compare the iteration count for an old version of our algorithm compared with IPOPT and KNITRO. This algorithm is based on similar ideas, but has significant differences from the current implementation.

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—Summary

Immunity

Introduction

Introd

Finally, before the presentation ends, I wish to give a quick summary. Currently, in optimization literature it is thought that infeasible start IPM needed to be significantly modified in order to guarantee convergence for non-convex optimization problems. In this talk, we have demonstrated by avoiding using equality constraints and gradually reducing the constraint violation it is possible to give convergence guarantees for a pure infeasible start algorithm. Furthermore, preliminary computational tests indicate this could become efficient algorithm in practice.

A one phase IPM for non-convex optimization $% \left(1\right) =\left(1\right) \left(1\right) \left($

—References

References

Ital Ashlagi, Mark Braverman, and Avinatan Haesidim (2014) Stability in Large Matching Markets with Complementarities Operations Research 62:4, 173-732. —Conjectures

Conjectures

Assume (x^*, s^*) is a limit point of the algorithm and there exists finite dual multipliers at (x^*, s^*) then for some subsequence of the dual multipliers $y^{e(k)} \rightarrow y^*$.

After a polynomial number of steps a simple variant of the algorithm has found either a first-order «-local optimum or a first-order «-local infeasibility certificate.