

## Practice Set 7 (Breadth First Search)

Data Structure (CS 223)

**Q1:** Answer the following questions for an undirected graph with  $N$  vertices and  $M$  edges:

- What is complexity of the BFS algorithm?

**Answer:**  $O(N + M)$

- How many times is each edge inspected (i.e., the vertex on either end is looked at) during the BFS algorithm on an undirected graph?

**Answer:** Twice (once each when the vertex on either end is dequeued).

- What is the maximum number of times a vertex is enqueued?

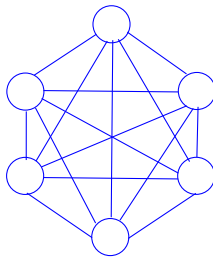
**Answer:** Once (when the first edge incident on the vertex is relaxed).

- While executing BFS, other than the start vertex, how many times is each vertex looked at?

**Answer:** As many times as the number of edges incident on the vertex.

- At any point, what is the maximum number of vertices on the queue? Draw a graph on 6 vertices for which this maximum number is realized, irrespective of the starting node.

**Answer:**  $N - 1$ , where  $N$  is the number of vertices.



**Q2:** Draw an undirected graph on 6 nodes such that at any point the queue contains at most 2 vertices, irrespective of where BFS starts.

**Answer:**  $v_0 - v_1 - v_2 - v_3 - v_4 - v_5$

**Q3:** Describe how you can use BFS to determine whether or not an undirected graph is connected.

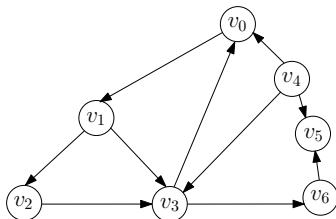
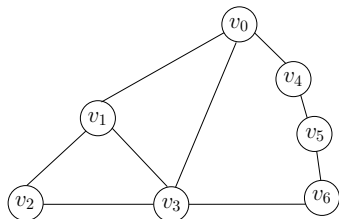
**Answer:** Start BFS at an arbitrary vertex. Upon termination, if there exists a vertex  $v$  for which  $level(v) = \infty$ , then the graph is not connected, otherwise it is connected.

**Q4:** Describe how you can use BFS to determine whether or not  $v$  is reachable from  $u$ .

**Answer:** Start BFS at  $u$ . Then a vertex  $v$  is reachable if and only if  $level(v) \neq \infty$  upon termination.

**Q5:** Starting from node  $v_1$ , illustrate the BFS algorithm on the following graphs. Show:

- the queue, and the level array at each stage
- the final BFS tree.



### BFS Undirected

	level array							queue
	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$	
<b>At Start:</b>	$\infty$	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$[v_1]$
Dequeue (Relax $v_1$ )	1	0	1	1	$\infty$	$\infty$	$\infty$	$[v_0, v_2, v_3]$
Dequeue (Relax $v_0$ )	1	0	1	1	2	$\infty$	$\infty$	$[v_2, v_3, v_4]$
Dequeue (Relax $v_2$ )	1	0	1	1	2	$\infty$	$\infty$	$[v_3, v_4]$
Dequeue (Relax $v_3$ )	1	0	1	1	2	$\infty$	2	$[v_4, v_6]$
Dequeue (Relax $v_4$ )	1	0	1	1	2	3	2	$[v_6, v_5]$
Dequeue (Relax $v_6$ )	1	0	1	1	2	3	2	$[v_5]$
Dequeue (Relax $v_5$ )	1	0	1	1	2	3	2	$[ ]$

### BFS Directed

	level array							queue
	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$	
<b>At Start:</b>	$\infty$	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$[v_1]$
Dequeue (Relax $v_1$ )	$\infty$	0	1	1	$\infty$	$\infty$	$\infty$	$[v_2, v_3]$
Dequeue (Relax $v_2$ )	$\infty$	0	1	1	$\infty$	$\infty$	$\infty$	$[v_3]$
Dequeue (Relax $v_3$ )	2	0	1	1	$\infty$	$\infty$	2	$[v_0, v_6]$
Dequeue (Relax $v_0$ )	2	0	1	1	$\infty$	$\infty$	2	$[v_6]$
Dequeue (Relax $v_6$ )	2	0	1	1	$\infty$	3	2	$[v_5]$
Dequeue (Relax $v_5$ )	2	0	1	1	$\infty$	3	2	$[ ]$

### BFS Trees

