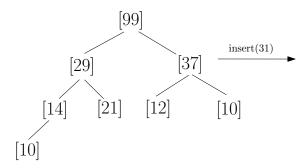
Homework 6 (COMPSCI 223)

Q1: Answer the following:

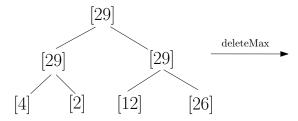
- Starting with the heap array [10, 15, 28, 16], show the final heap array after we insert 1 Answer: [1, 10, 28, 16, 15]
- Starting with the heap array [-1, 2, 7, 6, 18, 92], show the final heap array after we delete the minimum.

Answer: [2, 6, 7, 92, 18]

Q2: For each of the following MAX heaps, show the final heap array when the adjoining operation is performed.



Answer: [99, 31, 37, 29, 21, 12, 10, 10, 14]



Answer: [29, 26, 29, 4, 2, 12] or [29, 29, 26, 4, 2, 12]

Q3: Write true or false for the following:

- In the worst-case, the complexity of inserting in a heap having N nodes is $O(\log N)$. True.
- For sorting N numbers, the worst case complexity of heap sort is $O(N^2)$, whereas the best case complexity is $O(N \log N)$.

False. Both $O(N \log N)$

• Given a min-heap containing only distinct numbers, the second minimum number in a min-heap is a child of the root node.

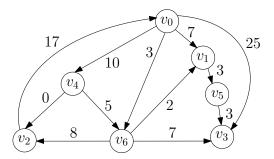
True.

• Suppose you have a max-heap, where all numbers are distinct. Then, the third highest number is definitely a child of the root node.

False. Depends on the structure. The third highest can be a child of the second highest or a child of the maximum.

Q4: The best-case complexity of inserting a number into a max-heap is O(1); for example, when you insert a new minimum.

Q5: Starting from node v_0 , illustrate Dijkstra's algorithm on the following graph. At each stage, show the relaxed vertex, the open and closed sets, and the dist and parent arrays.



	open	closed	dist array							parent array						
			v_0	v_1	v_2	v_3	v_4	v_5	v_6	v_0	v_1	v_2	v_3	v_4	v_5	v_6
Start:	v_0		0	∞	∞	∞	∞	∞	∞	Ø	Ø	Ø	Ø	Ø	Ø	Ø
relax v_0	v_1, v_3, v_4, v_6	v_0	0	7	∞	25	10	∞	3	Ø	v_0	Ø	v_0	v_0	Ø	v_0
relax v_6	v_1, v_2, v_3, v_4	v_0, v_6	0	5	11	10	10	∞	3	Ø	v_6	v_6	v_6	v_0	Ø	v_0
relax v_1	v_2, v_3, v_4, v_5	v_0, v_6, v_1	0	5	11	10	10	8	3	Ø	v_6	v_6	v_6	v_0	v_1	v_0
relax v_5	v_2, v_3, v_4	v_0, v_6, v_1, v_5	0	5	11	10	10	8	3	Ø	v_6	v_6	v_6	v_0	v_1	v_0
relax v_3	v_2, v_4	v_0, v_6, v_1, v_5, v_3	0	5	11	10	10	8	3	Ø	v_6	v_6	v_6	v_0	v_1	v_0
relax v_4	v_2	$v_0, v_6, v_1, v_5, v_3, v_4$	0	5	10	10	10	8	3	Ø	v_6	v_4	v_6	v_0	v_1	v_0
relax v_2		$v_0, v_6, v_1, v_5, v_3, v_4, v_2$	0	5	10	10	10	8	3	Ø	v_6	v_4	v_6	v_0	v_1	v_0

Q6: How many times is a vertex relaxed in Dijkstra's algorithm?

- At most once \leftarrow **Answer**
- As many times as the number of vertices in the graph
- As many times as the number of edges in the graph
- Depends on the professor's mood

Q7: Dijkstra's algorithm finds (choose the most appropriate):

- a shortest path from a source to a destination
- shortest paths from a source to every vertex ← Answer
- shortest paths from a source to all vertices other than source
- shortest paths from every vertex to every vertex

Q8: Given a directed weighted graph, is there a path from a vertex s to a vertex t whose weight is at most W? We can use Dijkstra's algorithm to answer this question as follows:

- Run Dijkstra from t. The answer is YES if dist[s] > W, else NO.
- Run Dijkstra from t. The answer is NO if $dist[s] \leq W$, else YES.
- Run Dijkstra from s. The answer is NO if dist[s] > W, else YES.
- Run Dijkstra from s. The answer is NO if dist[t] > W, else YES. \leftarrow Answer

Q9: A vertex is closed more than once in Dijkstra's algorithm. True or False?

Answer: False

Q10: You are given a pre-built max-heap having N numbers. Suppose the maximum number in the heap occurs k times. Give an $O(k \log N)$ time algorithm to compute k. Write pseudo-code.

Answer: Just keep deleting the maximum until you get a new number. More specifically,

- max = deleteMax()
- *k* = 1
- while $(size() \neq 0 \text{ and } max \text{ equals } deleteMax())$
 - -k++
- \bullet return k