

COMP 2611, DATA STRUCTURES

LECTURE 21

SORTING

- **Performance Analysis of Selection Sort, Bubble Sort, and Insertion Sort**
- **Mergesort**
- **Quicksort**

Mergesort Algorithm

- Uses a “divide-and-conquer” approach.

```
Mergesort (A, lengthA) {  
  
    sort the first half of A  
    sort the second half of A  
    merge the sorted halves  
  
}
```

Mergesort

- Suppose we need to sort the following array:

0	1	2	3	4	5	6	7	8	9
56	34	38	71	28	10	91	18	22	75



First half

Second half

- We sort the first half and the second half separately:

0	1	2	3	4	5	6	7	8	9
28	34	38	56	71	10	18	22	75	91



First half

Second half

Merge

```
void mergeArray (int A[], int p, int q, int r):
```

- This version of *mergeArray* accepts a single array, *A*, as a parameter. The values in *A* from location *p* to location *q* are in sorted order. The values from location *q+1* to location *r* are also in sorted order. However, the values from location *p* to location *r* are not sorted.

To be done
in Lab #10

0	1	2	3	4	5	6	7	8	9
28	34	38	56	71	10	18	22	75	91

p First half *q* *q+1* Second half *r*

Mergesort

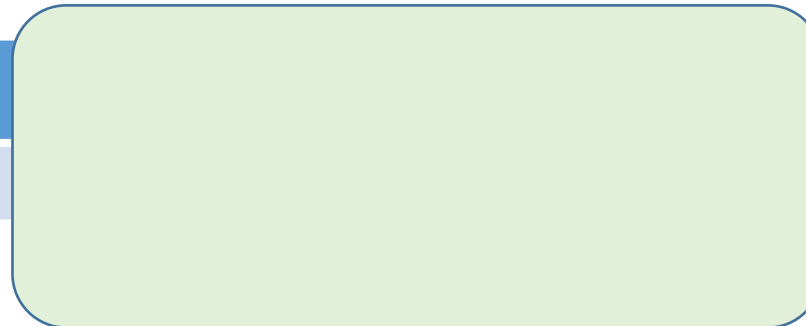
- Suppose we need to sort the following array:

0	1	2	3	4	5	6	7	8	9
56	34	38	71	28	10	91	18	22	75



- How to sort the first half and the second half?

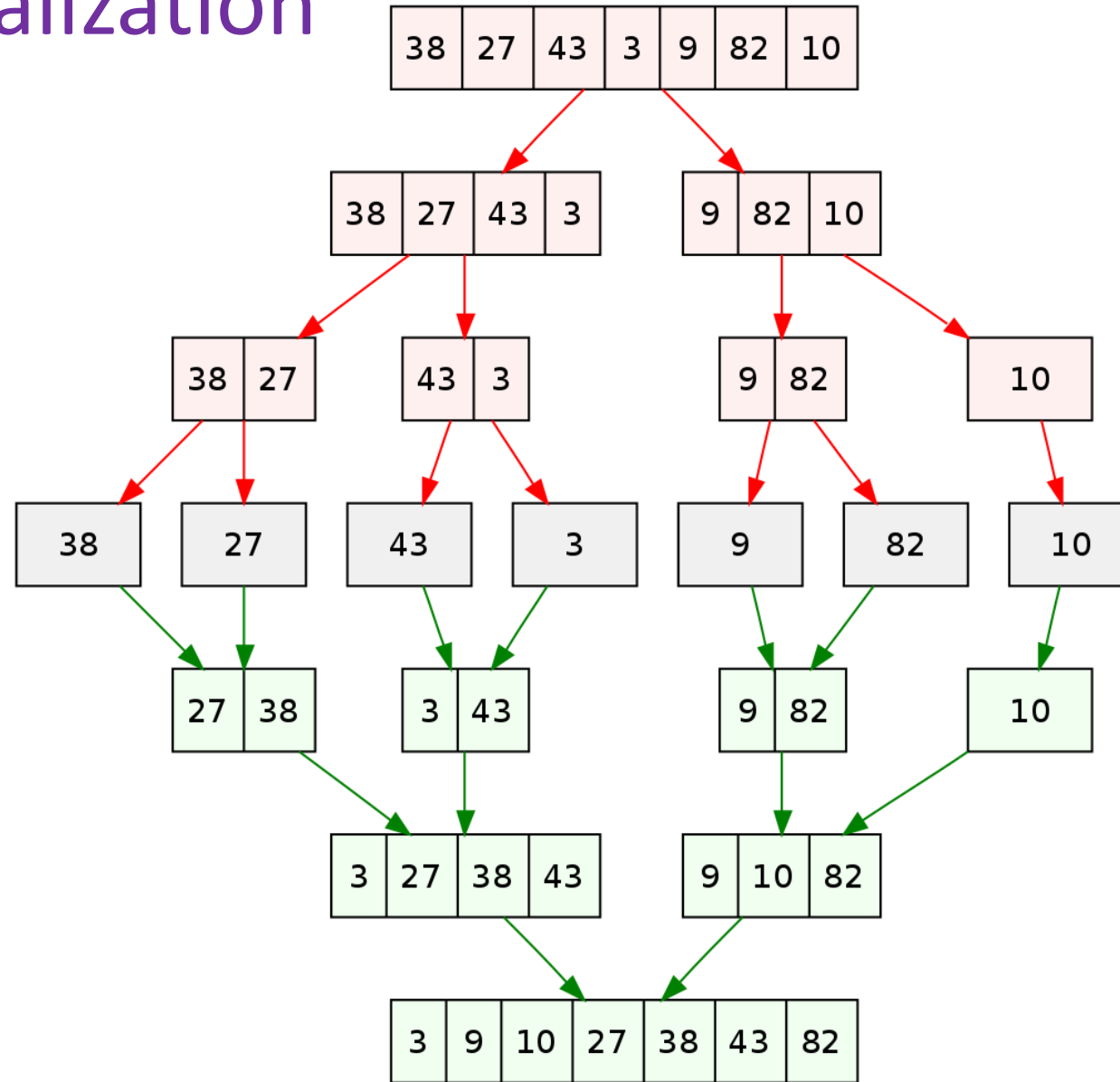
0	1	2	3	4
56	34	38	71	28



Mergesort Function

```
void mergeSort (int A[], int start, int end) {  
    int mid;  
  
    if (start < end) {  
        mid = (start + end) / 2;  
        mergeSort (A, start, mid);  
        mergeSort (A, mid+1, end);  
        merge (A, start, mid, end);  
    }  
}
```

Mergesort Visualization



Mergesort Animation

0	1	2	3	4	5	6
38	27	43	3	9	82	10

0	1	2	3
38	27	43	3

0	1
38	27

0	1
38	27

0	1
27	38

mergeSort (A, 0, 6):

$\text{mid} = (0 + 6) / 2 = 3$

mergeSort (A, 0, 3):

$\text{mid} = (0 + 3) / 2 = 1$

mergeSort (A, 0, 1):

$\text{mid} = (0 + 1) / 2 = 0$

mergeSort (A, 0, 0):

terminates!

mergeSort (A, 1, 1):

terminates!

merge (A, 0, 0, 1)

0	1	2	3	4	5	6
38	27	43	3	9	82	10

0	1	2	3
38	27	43	3

0	1	2	3
38	27	43	3

0	1	2	3
38	27	43	3

0	1	2	3
27	38	3	43

mergeSort (A, 0, 6):

$\text{mid} = (0 + 6) / 2 = 3$

mergeSort (A, 0, 3):

$\text{mid} = (0 + 3) / 2 = 1$

mergeSort (A, 0, 1)

mergeSort (A, 2, 3):

$\text{mid} = (2 + 3) / 2 = 2$

mergeSort (A, 2, 2):

terminates!

mergeSort (A, 3, 3):

terminates!

merge (A, 2, 2, 3)

0	1	2	3	4	5	6
38	27	43	3	9	82	10

0	1	2	3
38	27	43	3

0	1	2	3
38	27	43	3

0	1	2	3
38	27	43	3

0	1	2	3
27	38	3	43

0	1	2	3
3	27	38	43

mergeSort (A, 0, 6):

$\text{mid} = (0 + 6) / 2 = 3$

mergeSort (A, 0, 3):

$\text{mid} = (0 + 3) / 2 = 1$

mergeSort (A, 0, 1)

mergeSort (A, 2, 3)

merge (A, 0, 1, 3)

0	1	2	3	4	5	6
38	27	43	3	9	82	10

0	1	2	3
38	27	43	3

4	5	6
9	82	10

0	1
38	27

2	3
43	3

4	5
9	82

0
38

1
27

2
43

3
3

4
9

5
82

0	1
27	38

2	3
3	43

4	5
9	82

0	1	2	3
3	27	38	43

mergeSort (A, 0, 6):

$\text{mid} = (0 + 6) / 2 = 3$

mergeSort (A, 0, 3)

mergeSort (A, 4, 6):

$\text{mid} = (4 + 6) / 2 = 5$

mergeSort (A, 4, 5):

$\text{mid} = (4 + 5) / 2 = 4$

mergeSort (A, 4, 4):

terminates!

mergeSort (A, 5, 5):

terminates!

merge (A, 4, 4, 5)

0	1	2	3	4	5	6
38	27	43	3	9	82	10

0	1	2	3
38	27	43	3

4	5	6
9	82	10

0	1
38	27

2	3
43	3

4	5
9	82

6
10

0
38

1
27

2
43

3
3

4
9

5
82

0	1
27	38

2	3
3	43

4	5
9	82

0	1	2	3
3	27	38	43

4	5	6
9	10	82

mergeSort (A, 0, 6):

$\text{mid} = (0 + 6) / 2 = 3$

mergeSort (A, 0, 3)

mergeSort (A, 4, 6):

$\text{mid} = (4 + 6) / 2 = 5$

mergeSort (A, 4, 5)

mergeSort (A, 6, 6):

terminates!

merge (A, 4, 5, 6)

0	1	2	3	4	5	6
38	27	43	3	9	82	10

0	1	2	3
38	27	43	3

4	5	6
9	82	10

0	1
38	27

2	3
43	3

4	5
9	82

6
10

0
38

1
27

2
43

3
3

4
9

5
82

0	1
27	38

2	3
3	43

4	5
9	82

0	1	2	3
3	27	38	43

4	5	6
9	10	82

0	1	2	3	4	5	6
3	9	10	27	38	43	82

mergeSort (A, 0, 6):

$mid = (0 + 6) / 2 = 3$

mergeSort (A, 0, 3)

mergeSort (A, 4, 6)

merge (A, 0, 3, 6)

0	1	2	3	4	5	6
38	27	43	3	9	82	10

mergeSort (A, 0, 6)

0	1	2	3
38	27	43	3

4	5	6
9	82	10

0	1
38	27

2	3
43	3

4	5
9	82

6
10

0
38

1
27

2
43

3
3

4
9

5
82

0	1
27	38

2	3
3	43

4	5
9	82

0	1	2	3
3	27	38	43

4	5	6
9	10	82

0	1	2	3	4	5	6
3	9	10	27	38	43	82

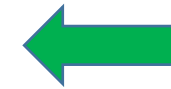
Quicksort Algorithm

- Suppose the portion of the array A between p and r needs to be sorted:

p						r
30	65	10	75	55	90	60

- Find an index q and reorganize elements such that:

- All elements to the left of q are smaller than $A[q]$
- All elements to the right of q are greater than $A[q]$



Achieved by a
partition algorithm

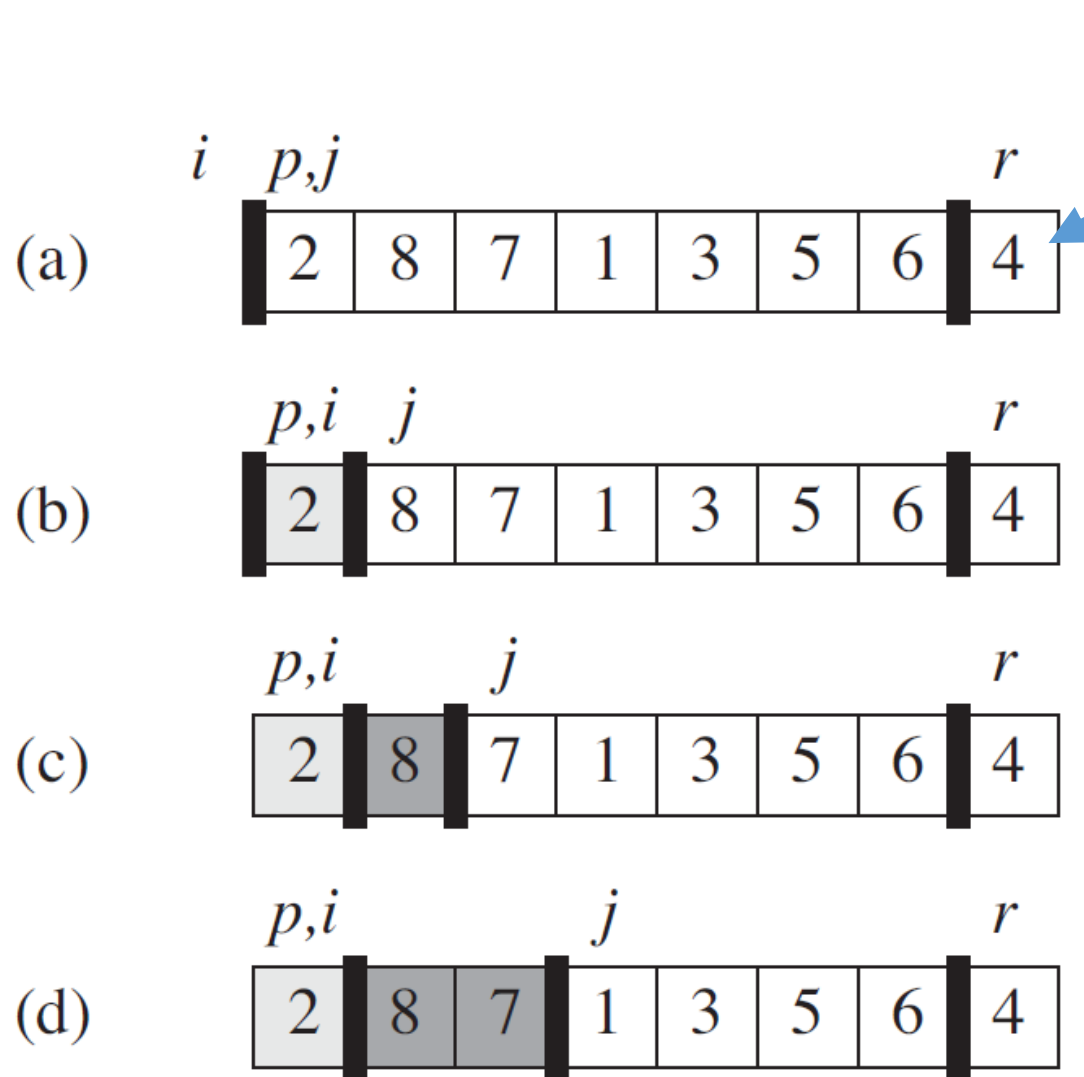
p			q			r
30	10	55	60	65	90	75

Quicksort Function

- Like mergesort, it uses a “divide-and-conquer” approach.

```
void quickSort (int A[], int p, int r) {  
    int q;  
  
    if (p < r) {  
        q = partition (A, p, r);  
        quickSort (A, p, q-1);  
        quickSort (A, q+1, r);  
    }  
}
```

Quicksort: Partition

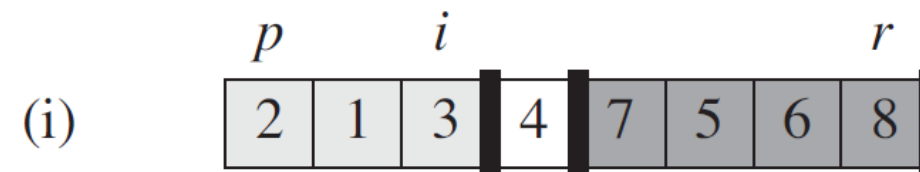
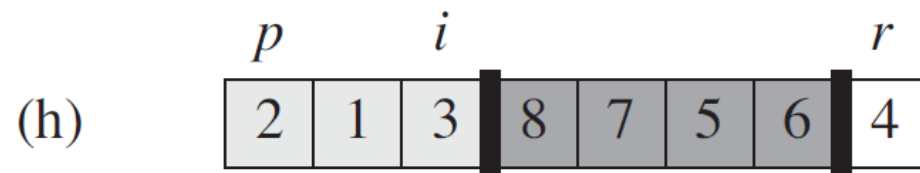
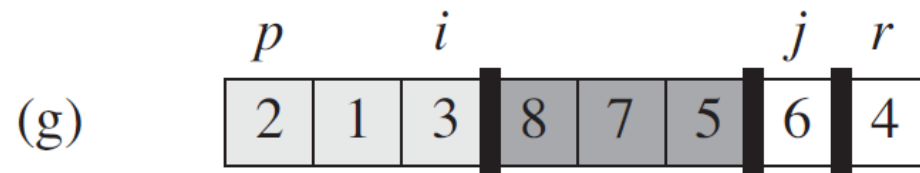
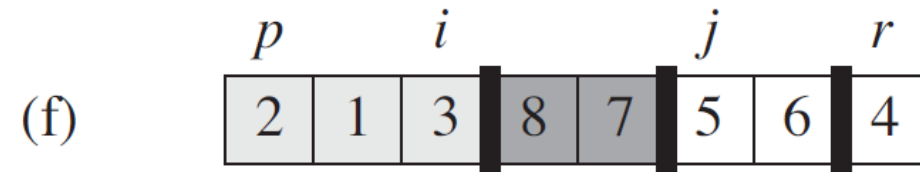
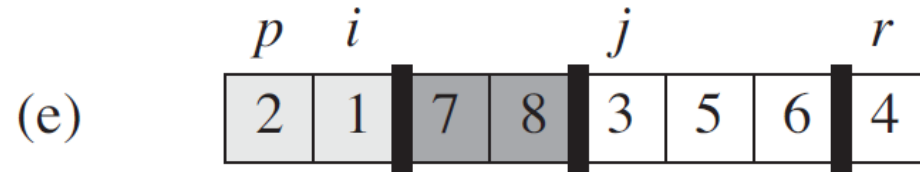


Pivot

PARTITION(A, p, r)

```
1   $x = A[r]$ 
2   $i = p - 1$ 
3  for  $j = p$  to  $r - 1$ 
4      if  $A[j] \leq x$ 
5           $i = i + 1$ 
6          exchange  $A[i]$  with  $A[j]$ 
7  exchange  $A[i + 1]$  with  $A[r]$ 
8  return  $i + 1$ 
```

Quicksort: Partition



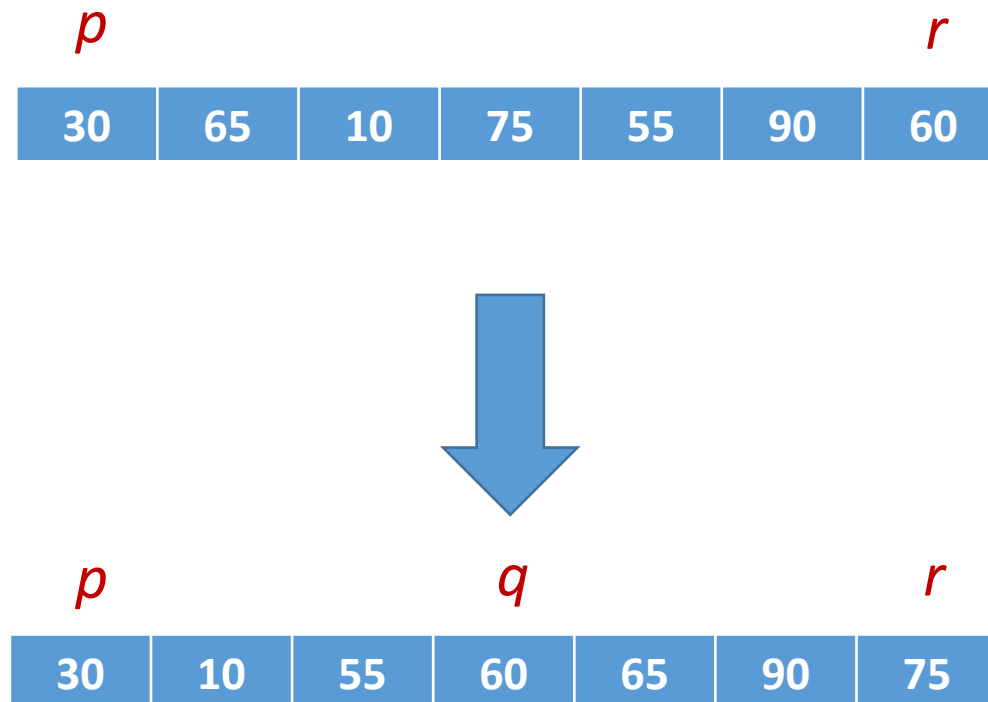
PARTITION(A, p, r)

```

1   $x = A[r]$ 
2   $i = p - 1$ 
3  for  $j = p$  to  $r - 1$ 
4      if  $A[j] \leq x$ 
5           $i = i + 1$ 
6          exchange  $A[i]$  with  $A[j]$ 
7  exchange  $A[i + 1]$  with  $A[r]$ 
8  return  $i + 1$ 
    
```

Quicksort: Partition Exercise

- What is the effect of partition (A, p, r) on the following array, A ?



PARTITION(A, p, r)

```
1   $x = A[r]$ 
2   $i = p - 1$ 
3  for  $j = p$  to  $r - 1$ 
4      if  $A[j] \leq x$ 
5           $i = i + 1$ 
6          exchange  $A[i]$  with  $A[j]$ 
7  exchange  $A[i + 1]$  with  $A[r]$ 
8  return  $i + 1$ 
```