


POLIMORFISMO:

1) C++ MONOMORFIZATIONE LINGUAGGI
TIPATI

SWAP<T>(T *p, *q) { ... }

SWAP<INT>(&x, &y);

SWAP<COUPLE<INT, CHAR>>(--, --); }
} COPICE
} COPIASO
2 VOLTE

SWAP -- INT

SWAP -- COUPLE -- INT -- CHAR

→ EFFICIENZA IN TEMPO

OPTIMIZATION TIPI -
- NO COPICE GENERICO DEPRECATI

- TEMPO DI COMPILAZIONE
- DIMENSIONE CODICE OCCORRE \Rightarrow CACHE MISSES

2) C STRATEGY (In PARTE RUST)

SWAP (void *x, void *y, SITE-k SITE);

- UNICHE OPERAZIONI CONCASSR

COPIA, SPOSTAMENTO, PLACCA MEMORIA,
ALLOCATION MEMORIA

DIMENSIONE DEI
DATI

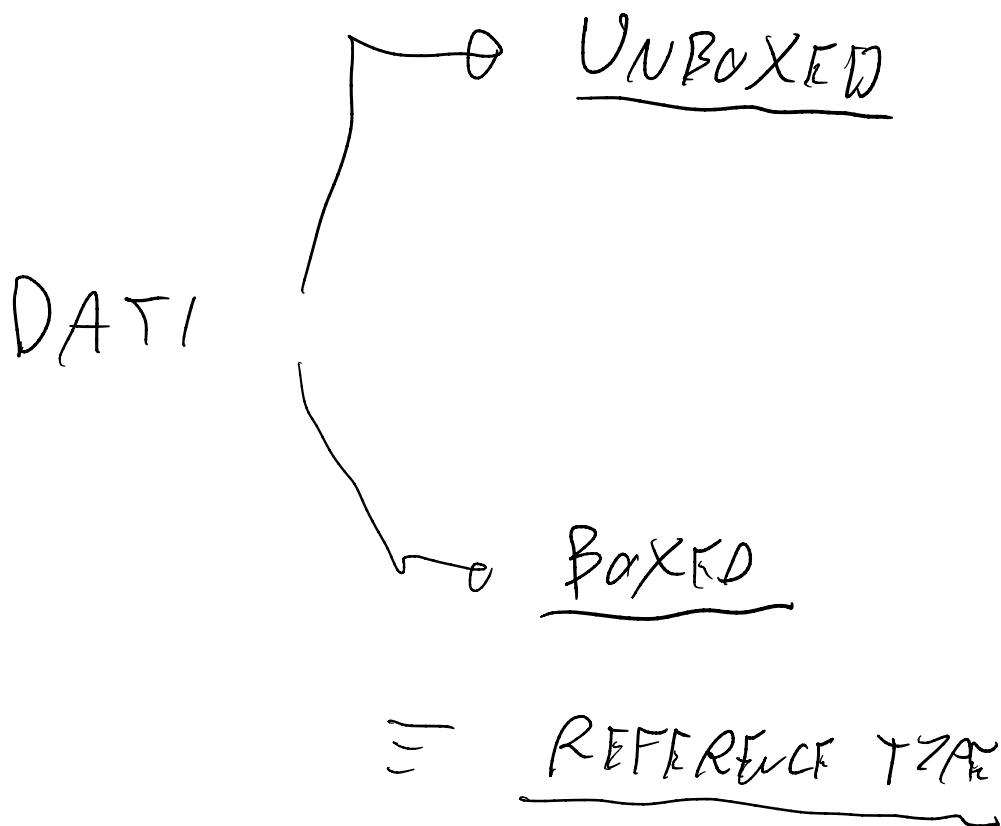
- MENO EFFICIENTE: CICLO FOR SUO SITE

- PRO: NO AUMENTO DIM. CODICE

- NO NECESSITA' DI CONOSCERE TUTTI GLI USI (PROG-SV)

3) RAPPRESENTAZIONE UNIFORME DEI DATI

1 DATO VIENE RAPP. CON 1 WORD



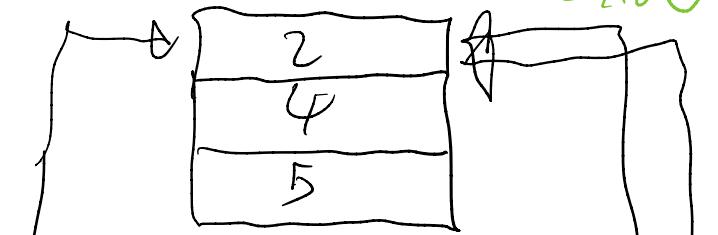
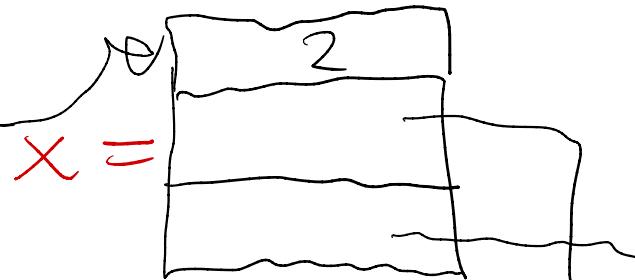
- OCCUPANO MEMO DI UNA PAROLA
- RAPP. SPRECANDO BIT
- DATI ALLOCATI SULLO HEAP COME:
 - [1 CELIA TAG]
 - [1 CELIA DELL'INDIRIZZO]CIOE CONSECUTIVI
- LA WORD CONTIENE UN REF ALLO HEAP

$\{ \{ 3, \text{true} \}, [2, \{ 4, 5 \}, 6] \}$

$\{ x, [y, z, w] \}$

SHARING

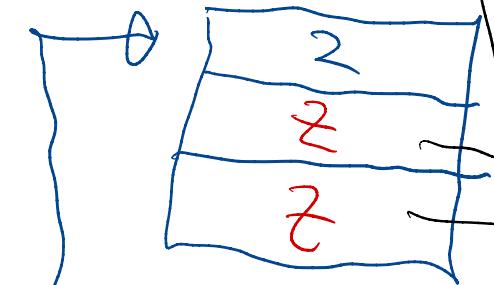
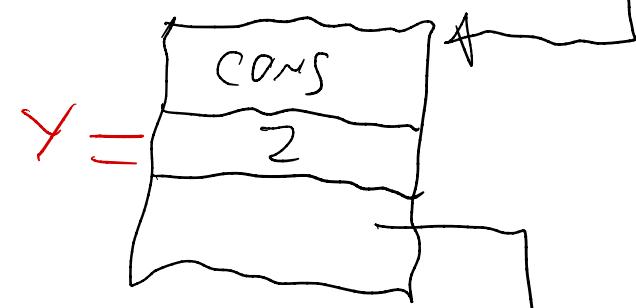
\uparrow



$O(|\text{PATTERN}|)$

OUTPUT:

$\{ 2, 2 \}$



$O(|\text{PATTERN}| \text{ OUTPUT})$

- BASSA COMPLESSITÀ
- CREATIONE DI SHARING AUTOMATICA
- CONCISONE

IF ((INT)*P == 2) {

X = ++P;

Y = ++P;

IF ((TAG)*q == CONS) {

Y = ++q;

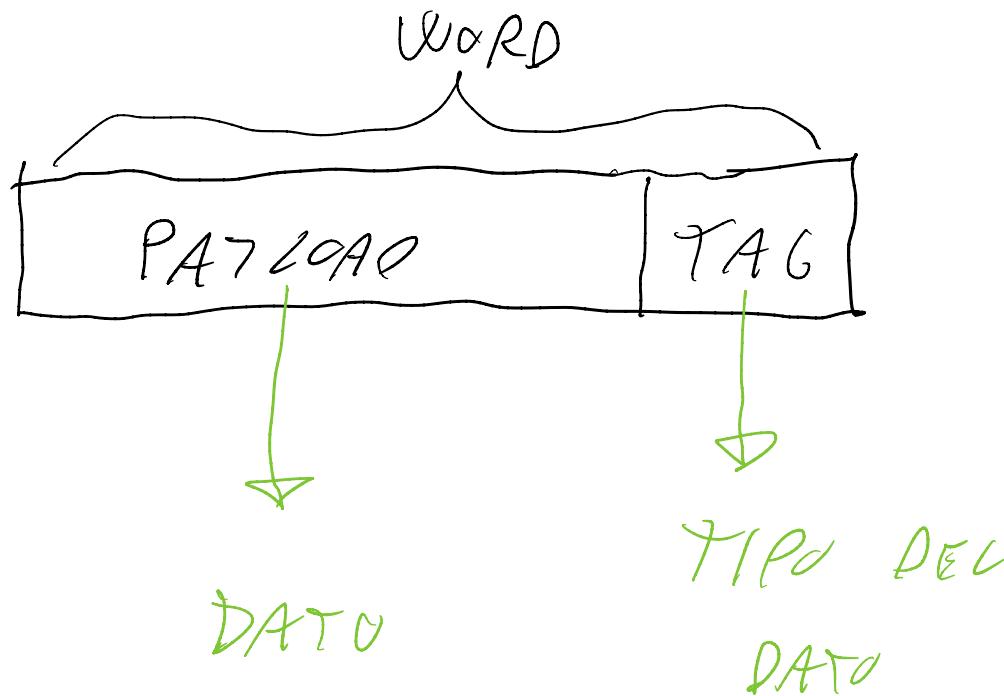
...

}

}

DATI

UNBOXED:



PROBLEMA: USARE

SEQUENZA DI BIT

DISGENTE PER
GARANTIRE

$\emptyset \neq []$

TAG A LUNGHEZZA VARIABILE:

es. ~~01, 011, 111~~ — NESSUNA SECUANZA E'
SUFFISSO DI UN'ALTRA
— Non USARE LO \emptyset

DATI

BOXED:

WORD



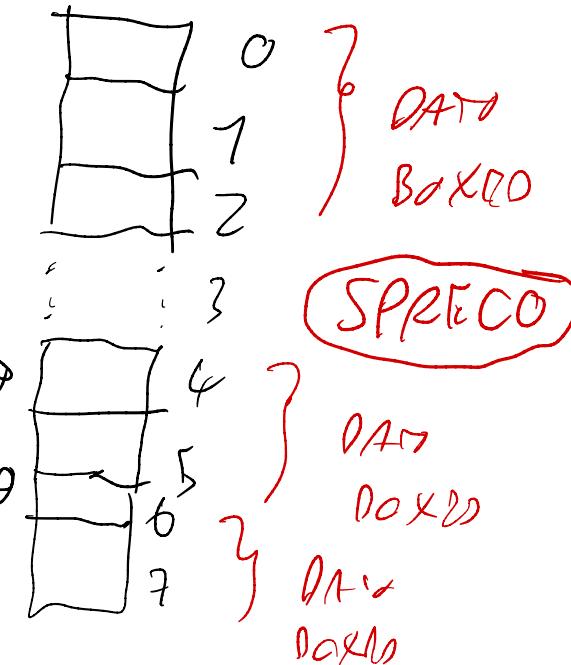
BIT

MENO

SIGNIFICATIVO

⇒ Solo Indiritti
PARI

⇒ Solo Indiritti
MEMORY ALIGNED



Poco Specie
di RAM

CONS: ARITMETICA

$$n \xrightarrow{+} \boxed{n \quad |1|} = 2 \cdot n + 1$$

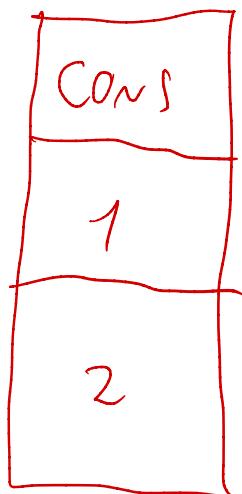
$$\begin{aligned} n+m &\xrightarrow{+} \boxed{n+m \quad |1|} = 2 \cdot (n+m) + 1 \\ &\neq (2 \cdot n + 1) + (2 \cdot m + 1) \end{aligned}$$

$$\begin{aligned} n \times n &\xrightarrow{+} \boxed{n \times n \quad |1|} = 2 \cdot n \cdot n + 1 \\ &\neq (2 \cdot n + 1) (2n + 1) \end{aligned}$$

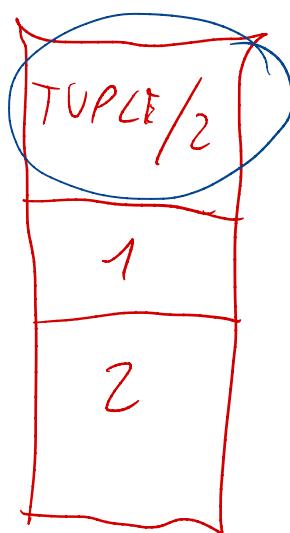
FIX: 1 OPERATION ARITMETICA \Rightarrow n OPERATIONS
MACCHINA

DISTINGUERI DATI BOXED:

$$(\text{Ex. } [1|2] \neq \{1, 2\})$$



VS

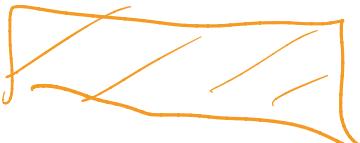


TAG + NURATO DI

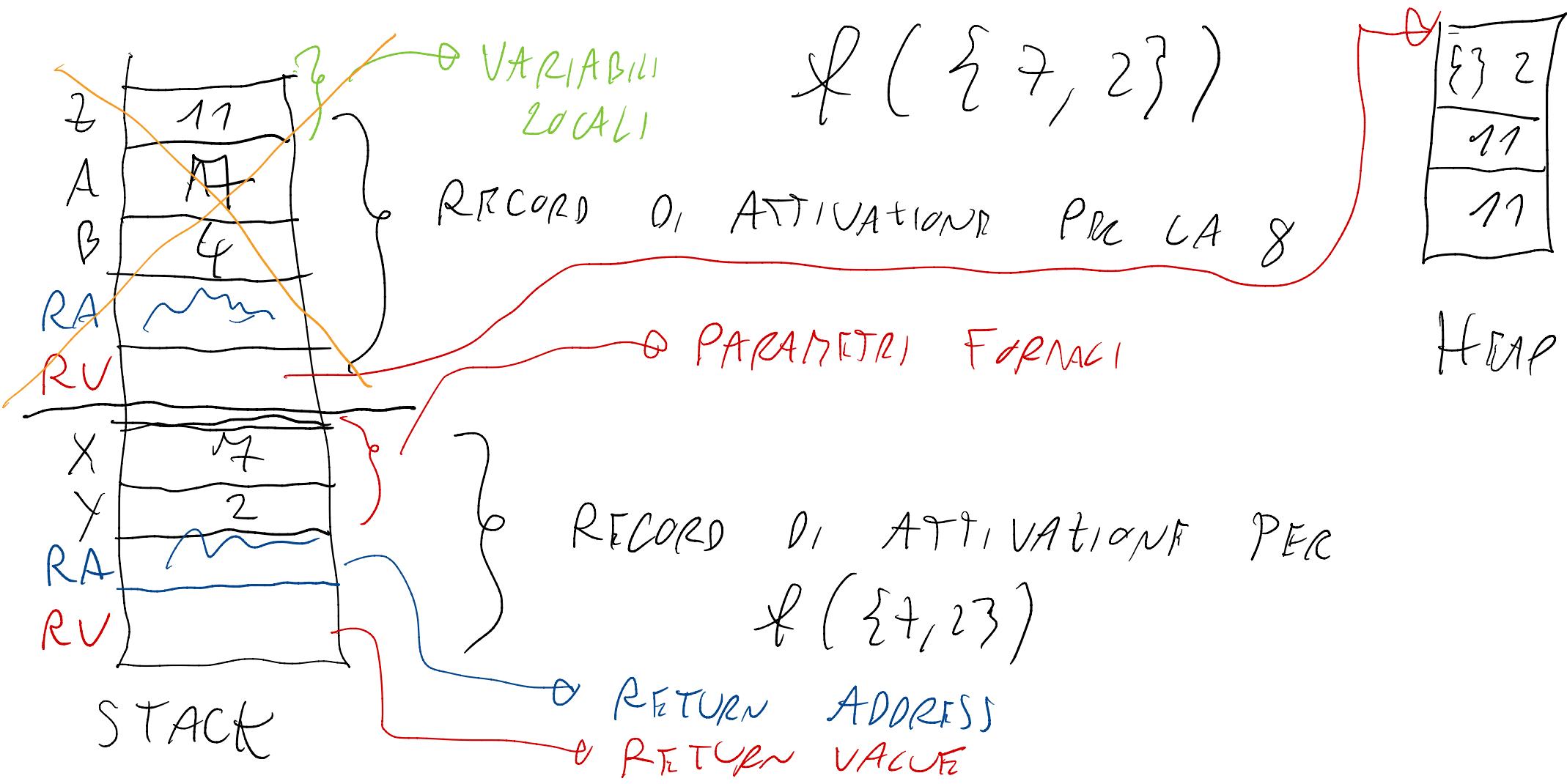
CARPI COPIFICATI

NELLA STESSA

VOCO

$f(\{x, y\}) \rightarrow g(x, y+2)$. 

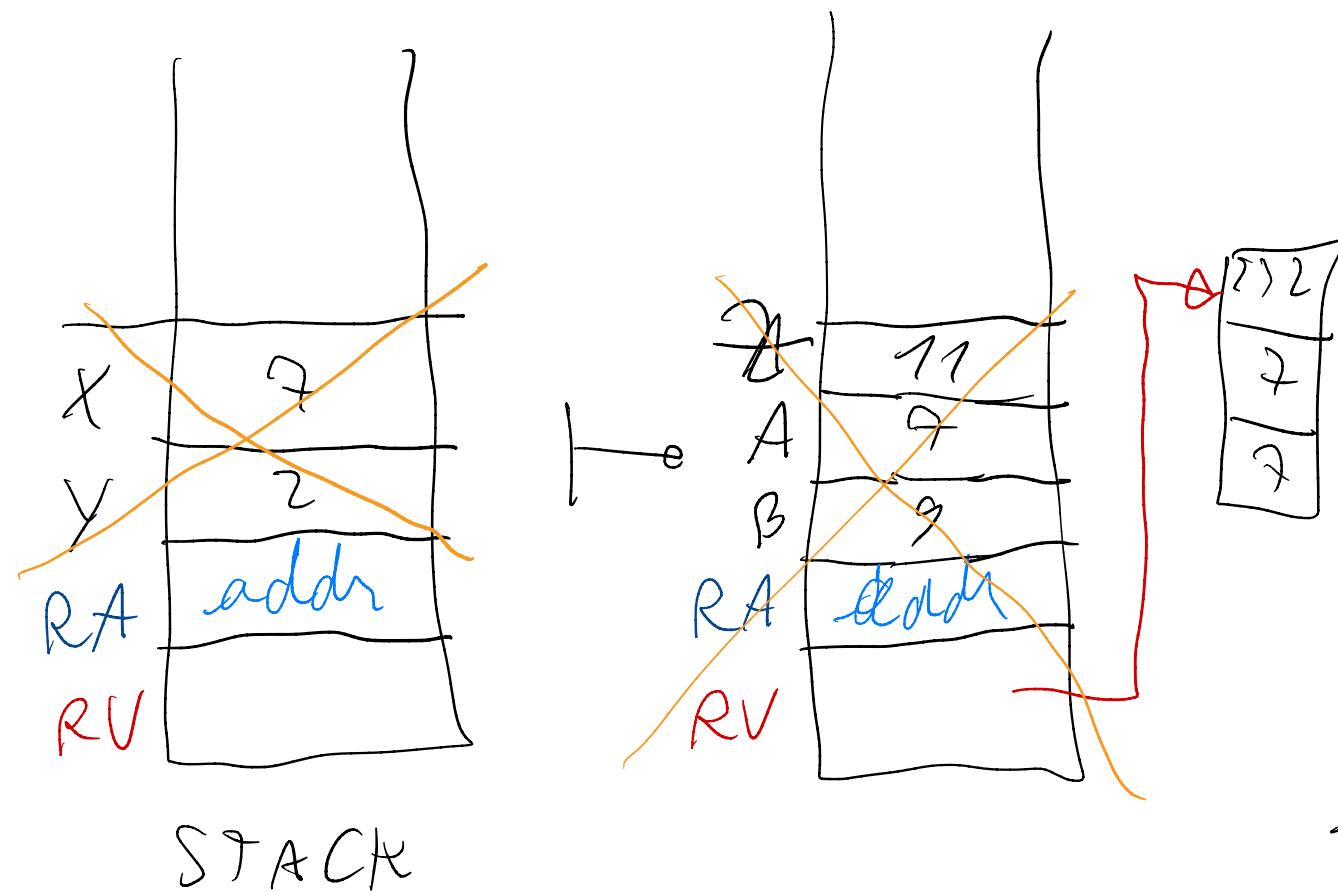
$g(A, B) \rightarrow z = A + B, \{z, z\}$.



$\star(\{x, y\}) \rightarrow \underline{g}(x, y+2).$

$g(A, B) \rightarrow \begin{cases} z = A + B, & \{z, z\}. \end{cases}$

Y CHIAMATA DI CODA:



- LA \star NON FA nulla
Dopo

- LA \star RESTITUISCE
QUANTO RESTITUITO DA
g

$\star(\{z, y\})$