3. Recursion

- · What is Recursion
- · What it is good for and what it is not good for
- · What are the characteristics of recursion
- · How is a recursive function executed
- · How is recursive function call implemented.

What is Recursion

- Self referential (defined) in terms of itself)
- The laughing-cow (la vache qui rit) package shows a cow wearing laughing-cow packages as earrings, which show a cow wearing laughing-cow packages as earrings which ...



Other examples

A linked list is:

- a) empty, or
- b) has a head (first element) and a tail, which is a linked list

A tree is:

- a) empty, or
- b) has a root, and left and/or right (sub-) trees

Factorial function

Factorial 5, written 5!, is: $5\times4\times3\times2\times1$

and 6! is $6 \times 5 \times 4 \times 3 \times 2 \times 1$, so $6 \times 5!$

Factorial function, for non-negative integers is:

- a) 0! = 1
- b) if n > 0, then $n! = n \times (n 1)!$

```
if (n == 0) return 1;
  else return (n * factorial(n - 1));
 }
```

Caution: inefficient

Useful recursion

- ◆To be useful the recursion must terminate, so there must be at least one non-recursive case such as: 0!
- as well as recursive cases.

such as: n * (n - 1)!

Infinite recursion

```
void TellStory(){
  printf("%s", "It was a dark and stormy night ");
  printf("%s", "and the captain said to the mate ");
  printf("%s", ": `Tell us a story mate' ");
  printf("%s", " and this is the story he told ...");
  TellStory();
}
```

Recursive Programming

- Consider the problem of computing the sum of all the integers between 1 and any positive integer N
- This problem can be recursively defined as:

$$\sum_{i=1}^{N} i = N + \sum_{i=1}^{N-1} i = N + N-1 + \sum_{i=1}^{N-2} i$$

$$= N + N-1 + N-2 + \sum_{i=1}^{N-3} i$$

$$\vdots$$

Recursive Programming

Recursive Programming

- Note that just because we can use recursion to solve a problem, doesn't mean we should
- For instance, we usually would not use recursion to solve the sum of 1 to N problem, because the iterative version is easier to understand and more efficient
- However, for some problems, recursion provides an elegant solution, often cleaner than an iterative version
- You must carefully decide whether recursion is the correct technique for any problem

Recursion 1

Indirect Recursion

- A function invoking itself is considered to be direct recursion
- A function could invoke another function, which invokes another, etc., until eventually the original function is invoked again
- For example, function f1 could invoke f2, which invokes f3, which in turn invokes f1 again
- This is called indirect recursion, and requires all the same care as direct recursion
- It is often more difficult to trace and debug

Recursion 11

Length of a list

- a) the length of an empty list is 0
- b) the length of a (non-empty) list is:
 - 1 + the length of the tail of the list

Length of a list in C

```
int length_v1 (node* p){ /* iteration */
    int countNodes = 0;
    while (p) do {
        countNodes++;
        p = p->next
    }
    return countNodes;
}

int length_v2(node* p){ /* recursion */
    if (p) return (1 + length_v2(p->next));
    else return 0;
}
```

Traversing a list: iterative

```
Traversing a (singly) linked list iteratively in the forward direction is easy:
```

```
void traverse (node* p){
while (p) {
  process(p->data); /* assume a process function */
  p = p->next;
}
```

Traversing iteratively in the backward direction is **hard** (no pointers, so need to *stack* return pointers)

Traversing a list: recursive, forward

Traversing a (singly) linked list *recursively* in the forward direction is easy:

```
void traverse (node* p){
    if (p){
        process(p->data);
        traverse(p->next);
    }
}
```

tursion 15

Traversing a list: recursive, backward

Traversing a (singly) linked list recursively in the **backward** direction is also easy:

```
void reverseTraverse (node* p){
   if (p){
      reverseTraverse(p->next);
      process(p->data);
   }
}
```

How recursion works

- When a function is called, its parameters, local variables and return address are stacked on the function-call stack.
- Nested calls lead to deeper stacking.
- A call of a function to itself is just another nested call.

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When not to use recursion

- Don't use a recursive approach when a simple iterative approach is available
- Examples: searching, traversing and inserting in a list is easy to do iteratively
- Traversing a list backwards ('backtracking') is easy to do recursively but hard to do iteratively.

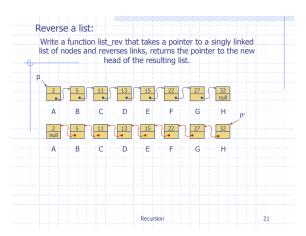
Recursion 18

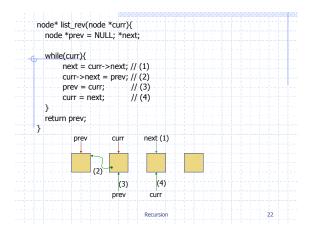
When not to use recursion: example Fibonacci Numbers: fib₀ = 0 fib₁ = 1 fib_n = fib_{n-1} + fib_{n-2}, for n > 0 int fib(n: integer){ /* doubly recursive */ if (n == 0) return 0; else if (n ==1) return 1; return (fib(n - 1) + fib (n - 2)); } Very inefficient: values repeatedly calculated, then 'forgotten'

```
A Better way:

int fib (n: integer){ /*iterative */
    int i, x, y, z;

    i = 1; x = 1; y = 0;
    while (i != n) {
        z = x;
        i++;
        x = x + y;
        y = z;
    }
    return x;
}
```





```
node *list_rev_recursion(node *curr, node *prev) {
    node *revHead;
    if (curr == NULL)
        revHead = prev;
    else {
        revHead = list_rev_recursion(curr->next, curr);
        curr->next = prev;
    }
    return revHead;
}

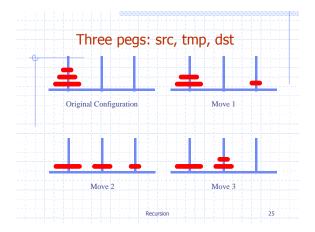
Initial method call should be
    head = list_rev_recursion(head, NULL)

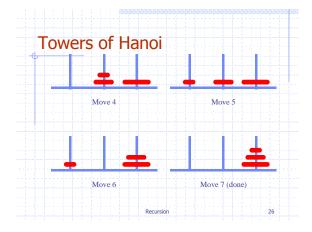
Recursion 23
```

Towers of Hanoi

- The Towers of Hanoi is a puzzle made up of three vertical pegs and several disks that slide on the pegs
- The disks are of varying size, initially placed on one peg with the largest disk on the bottom with increasingly smaller ones on top
- The goal is to move all of the disks from one peg to another under the following rules:
 - Only one disk can be moved at a time
 - A bigger disk can never be placed on top of a smaller one

Recursion 2





Towers of Hanoi An iterative solution to the Towers of Hanoi is quite complex A recursive solution is much shorter and more elegant if (n == 1) { (move one disk directly from src to dst) } else { (move a tower of n-1 disks from src to tmp) (move one disk directly from src to dst) (move a tower of n-1 disk from tmp to dst) } Recursion 27

```
Three Characteristics of
Recursion

Calls itself recursively

Has some terminating condition

Moves "closer" to the terminating condition.
```

```
Two Flavors of Recursion

if (terminating condition) {
    do final actions
} else {
    move one step closer to terminating condition
    recursive call(s)
}

-or-
if (!(terminating condition)) {
    move one step closer to terminating condition
    recursive call(s)
}
```

Analysis of a Recursive Algorithm

- To evaluate the computing costs of an algorithm, we need to analyze
 i+
- The first step in algorithm analysis is to find out the number of times the "primitive operation" is executed.
- In doing so, we express the number of times as a function of n, which is the input size.

ecursion 31

Analysis of a Recursive Algorithm

Steps in analyzing a recursive algorithm:

- 1.Decide on the input size n.
- 2. Identify the primitive operation.
- 3. Set up a recurrence relation, with an appropriate initial condition, for the number of times the primitive operation is executed.
- 4. Solve the recurrence.

Recursion 3

Analysis of a Recursive Algorithm

+Example: recursive computing of factorial

```
int factorial (int n){
  if (n == 0) return 1;
  else return (n * factorial(n - 1));
}
```

Input size: n

Primitive operation: multiplication

on .

Analysis of a Recursive Algorithm

Example: recursive computing of factorial

Recurrence relation for the number of times multiplication is executed and initial condition:

$$M(n) = M(n-1) + 1 \text{ for } n > 0$$

 $M(0) = 0$

ursion

Analysis of a Recursive Algorithm

Solve the recurrence relation using the method of backward substitution:

```
\begin{split} &M(n) = M(n-1) + 1 & \text{substitute } M(n-1) = M(n-2) + 1 \\ &= [M(n-2) + 1] + 1 = M(n-2) + 2 & \text{substitute } M(n-2) = M(n-3) + 1 \\ &= [M(n-3) + 1] + 2 = M(n-3) + 3 & \text{substitute } M(n-3) = M(n-4) + 1 \end{split}
```

```
M(n) = M(n-1) + 1 = ... = M(n-i) + i = ...

M(n) = M(n-i) + i, M(0) = 0
```

To eliminate the M term from the right hand side, let i = n and use the initial condition, we have

```
M(n) = n
Recursion 35
```

Analysis of a Recursive Algorithm

Example: Hanoi tower

Analysis of a Recursive Algorithm

Example: Hanoi tower

```
1. Input size: n (the number of disks)
2. Primitive operation: Moving a disk
3. Recurrence and initial condition:
M(n) = M(n-1) + 1 + M(n-1)
   = 2M(n-1) + 1
                             n>1
M(1) = 1
```

Analysis of a Recursive Algorithm

Example: Hanoi tower

```
M(n) = 2M(n-1) + 1
                               sub. M(n-1)=2M(n-2)+1
=2[2M(n-2)+1]+1
=2^{2}M(n-2)+2+1
                               sub. M(n-2)=2M(n-3)+1
=2^{2}[2M(n-3)+1]+2+1
=2^3M(n-3)+2^2+2+1
M(n) = 2^{i}M(n-i) + 2^{i-1} + 2^{i-2} + \cdots + 2 + 1
     = 2^{i}M(n-i)+2^{i}-1
```

Using
$$\sum_{j=0}^{j=i-1} 2^j = 2^{i-1}$$

Analysis of a Recursive Algorithm

```
Example: Hanoi tower
M(n) = 2M(n-1)+1 = 2^{i}M(n-1)+2^{i}-1 n>1
M(1) = 1
To use the initial condition, let i=n-1.
We have
M(n) = 2^{n-1}M[n-(n-1)] + 2^{n-1} - 1
    = 2^{n-1}M(1) + 2^{n-1} - 1
    = 2^{n-1} + 2^{n-1} - 1
    = 2^n - 1
```

That is, to move n disks, the number of movements is 2^n-1 .

Tracing The Recursion

To keep track of recursive execution, do what a computer does: maintain information on an activation stack.

Each stack frame contains:

- Module identifier and variables
- Any unfinished business

ModuleID: **Data values** Unfinished business

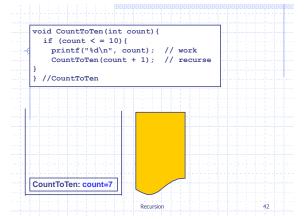
Work and Recursion

Problem: Count from N to 10.

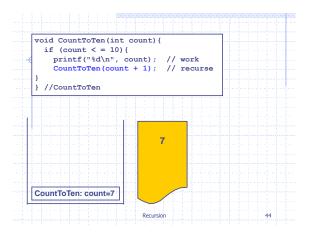
void CountToTen(int count) { if (count <= 10) { printf("%d\n", count); // work CountToTen(count + 1); // recurse

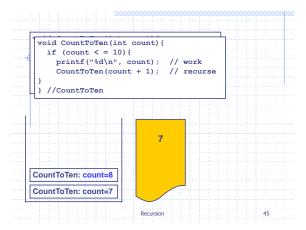
} //CountToTen

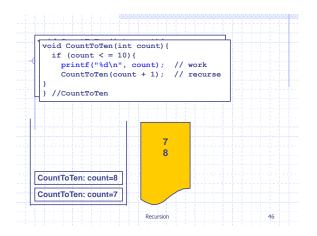
First do the work and then the recursive call!

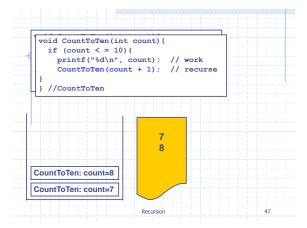


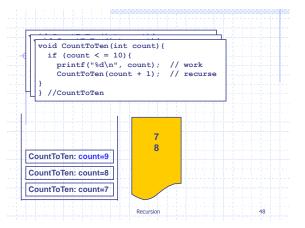
```
void CountToTen(int count){
  if (count < = 10) {
    printf("%d\n", count); // work
    CountToTen(count + 1); // recurse
}
} //CountToTen</pre>
7
CountToTen: count=7
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```

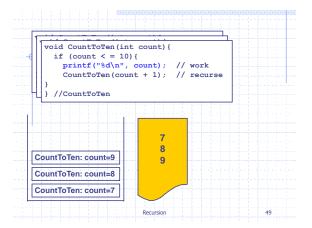


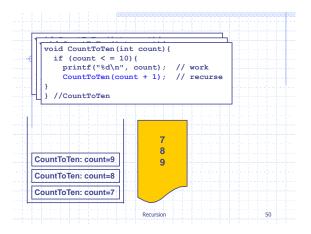


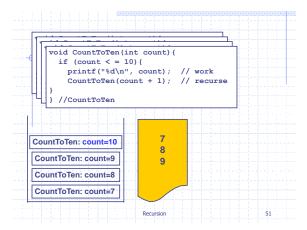


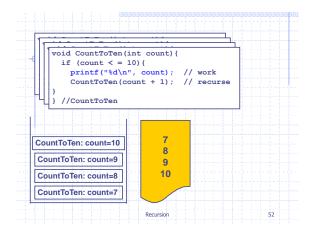


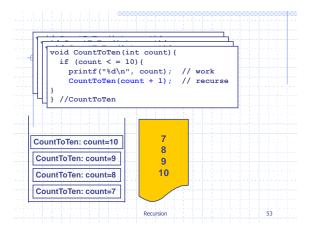


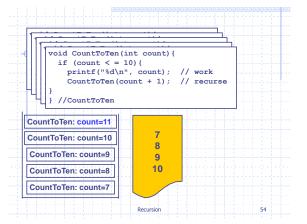


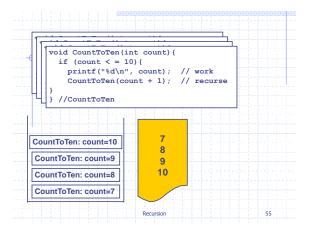


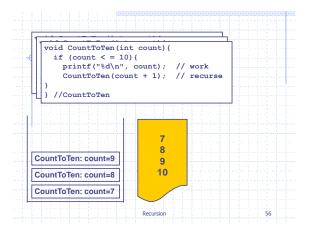


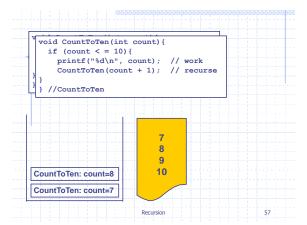


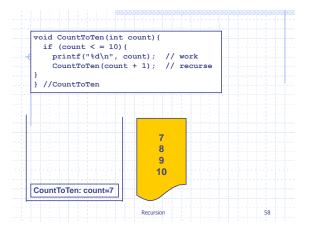


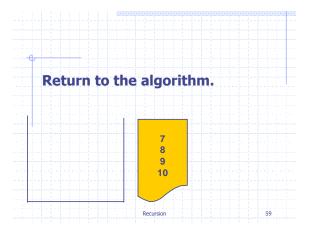












```
Reversing the Work and Recursion

Problem: Count from N to 10.

void CountToTen(int count) {
   if (count <= 10) {
        CountToTen(count + 1); // recurse
        printf("%d\n", count); // work
   }
} //CountToTen

Now the work will happen as the
   frames pop off the stack!
```

