Part 25: Mutex 15 AUGUST 2017

Welcome to tutorial no. 25 in Golang tutorial series.

In this tutorial we will learn about mutexes. We will also

learn how to solve race conditions using mutexes and channels. **Critical section**

Before jumping to mutex, it is important to understand the concept of <u>critical section</u> in concurrent

programming. When a program runs concurrently, the parts of code which modify shared resources should not be accessed by multiple Goroutines at the same time. This section of code which modifies shared resources is called critical section. For example lets assume that we have some piece of code which increments a variable x by 1. x = x + 1

As long as the above piece of code is accessed by a single

```
Let's see why this code will fail when there are multiple
Goroutines running concurrently. For the sake of
```

Goroutine, there shouldn't be any problem.

simplicity lets assume that we have 2 Goroutines running the above line of code concurrently.

Internally the above line of code will be executed by the system in the following steps(there are more technical details involving registers, how addition works and so on but for the sake of this tutorial lets assume that these are the three steps),

1. get the current value of x 2. compute x + 13. assign the computed value in step 2 to x When these three steps are carried out by only one Goroutine, all is well.

code concurrently. The picture below depicts one scenario of what could happen when two Goroutines

Goroutine 1 x = 01. Current value of x is 0 2. compute x + 1

Goroutine 1 3. Assign computed value to x x = 1Now x becomes 1 Goroutine 2 3. Assign computed value to x x = 1Now x becomes 1 x = 1 after both Goroutines finish x = 1executing

Assume x = 0 initially Value of x after each step Goroutine 1 1. Current value of x is 0 2. compute x + 1 x = 13. Assign computed value to x Now x becomes 1 Goroutine 2 1. Current value of x is 1 2. compute x + 1x = 23. Assign computed value to x Now x becomes 2

So from the two cases you can see that the final value of x is 1 or 2 depending on how context switching happens.

This type of undesirable situation where the output of

the program depends on the sequence of execution of

In the above scenario, the race condition could have been

avoided if only one Goroutine was allowed to access the

critical section of the code at any point of time. This is

Goroutines is called **race condition**.

made possible by using Mutex.

mutex.Lock()

mutex.Unlock()

x = x + 1

race condition.

package main

"fmt"

"sync"

x = x + 1

var x = 0

import (

Mutex

A Mutex is used to provide a locking mechanism to ensure that only one Goroutine is running the critical section of code at any point of time to prevent race condition from happening. Mutex is available in the sync package. There are two methods defined on Mutex namely Lock and Unlock. Any

In the above code, x = x + 1 will be executed by only one Goroutine at any point of time thus preventing race condition. If one Goroutine already holds the lock and if a new Goroutine is trying to acquire a lock, the new Goroutine

wg.Done() 10 func main() { var w sync.WaitGroup for i := 0; i < 1000; i++ { 13 W.Add(1)go increment(&w)

func increment(wg *sync.WaitGroup) {

for each time because of race condition. Some of the outputs which I encountered are final value of x 941, final value of x 928, final value of x 922 and so on.

Solving the race condition using

In the program above, we spawn 1000 Goroutines. If each

increments the value of x by 1, the final desired value of x

should be 1000. In this section, we will fix the race

condition in the program above using mutex.

package main

"fmt"

"sync"

var x = 0

Run in playground

import (

var w sync.WaitGroup 14 var m sync.Mutex 15 for i := 0; i < 1000; i++ { w.Add(1) 17 go increment(&w, &m) 18 19 w.Wait() fmt.Println("final value of x", x) 21 22

Mutex is a struct type and we create a zero valued

variable m of type Mutex in line no. 15. In the above

and m. Unlock(). Now this code is void of any race

this piece of code at any point of time.

Now if this program is run, it will output

program we have changed the increment function so that

the code which increments x = x + 1 is between m. Lock ()

conditions since only one Goroutine is allowed to execute

no. 18. If the mutex is passed by value instead of passing the address, each Goroutine will have its own copy of the mutex and the race condition will still occur. Solving the race condition using channel

It is important to pass the address of the mutex in line

In the program above, we have created a buffered channel of capacity 1 and this is passed to the increment Goroutine in line no. 18. This buffered channel is used to ensure that only one Goroutine access the critical section

final value of x 1000 **Mutex vs Channels** We have solved the race condition problem using both mutexes and channels. So how do we decide what to use when. The answer lies in the problem you are trying to solve. If the problem you are trying to solve is a better fit for mutexes then go ahead and use mutex. Do not hesitate to use mutex if needed. If the problem seems to be a better fit for channels, then use it :).

In general use channels when Goroutines need to communicate with each other and mutexes when only one Goroutine should access the critical section of code. In the case of the problem which we solved above, I would prefer to use mutex since this problem does not

My advice would be to choose the tool for the problem and do not try to fit the problem for the tool:) This brings us to an end of this tutorial. Have a great day.

Next tutorial - <u>Structs Instead of Classes</u>

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About

Lets discuss what happens when 2 Goroutines run this

access the line of code x = x + 1 concurrently. Assume x = 0 initially Value of x after each step

Goroutine 2

1. Current value of x is 0 x = 02. compute x + 1We have assumed the initial value of x to be 0. Goroutine 1 gets the initial value of x, computes x + 1 and before it could assign the computed value to x, the system context switches to Goroutine 2. Now Goroutine 2 gets the initial value of x which is still 0, computes x + 1. After this the system context switches again to Goroutine 1. Now

Goroutine 1 assigns it's computed value 1 to x and hence x

becomes 1. Then Goroutine 2 starts execution again and

then assigns it's computed value, which is again 1 to x

Now lets see a different scenario of what could happen.

and hence x is 1 after both Goroutines execute.

In the above scenario, Goroutine 1 starts execution and finishes all its three steps and hence the value of x

becomes 1. Then Goroutine 2 starts execution. Now the

value of x is 1 and when Goroutine 2 finishes execution,

x = 2

x = 2 after both Goroutines finish

executing

the value of x is 2.

code that is present between a call to Lock and Unlock will be executed by only one Goroutine, thus avoiding race condition.

will be blocked until the mutex is unlocked. **Program with race condition** In this section, we will write a program which has race

condition and in the upcoming sections we will fix the

16 w.Wait() 17 fmt.Println("final value of x", x) 18 19

In the program above, the increment function in line no. 7

increments the value of x by 1 and then calls Done () on

We spawn 1000 increment Goroutines from line no. 15 of

concurrently and race condition occurs when trying to

increment x is line no. 8 as multiple Goroutines try to

Please run this program in your local as the playground is

the program above. Each of these Goroutines run

the WaitGroup to notify its completion.

access the value of x concurrently.

mutex

deterministic and the race condition will not occur in the playground. Run this program multiple times in your local machine and you can see that the output will be different

m.Lock() x = x + 1m.Unlock() 10 wg.Done() 11 12 func main() { 13

func increment(wg *sync.WaitGroup, m *sync.Mutex)

We can solve the race condition using channels too. Lets see how this is done.

package main

"fmt"

"sync"

ch <- true

x = x + 1

<- ch

21

22

Run in playground

import (

final value of x 1000

wg.Done() 11 12 func main() { 13 var w sync.WaitGroup 14 ch := make(chan bool, 1) 15 for i := 0; i < 1000; i++ { 16 w.Add(1) 17 go increment(&w, ch) 18 19 w.Wait() 20

fmt.Println("final value of x", x)

func increment(wg *sync.WaitGroup, ch chan bool)

of code which increments x. This is done by passing true to the buffered channel in line no. 8 just before x is incremented. Since the buffered channel has a capacity of 1, all other Goroutines trying to write to this channel are blocked until the value is read from this channel after incrementing x in line no. 10. Effectively this allows only one Goroutine to access the critical section. This program also prints

Most Go newbies try to solve every concurrency problem using a channel as it is a cool feature of the language. This is wrong. The language gives us the option to either use Mutex or Channel and there is no wrong in choosing either.

Hence mutex would be a natural fit.

require any communication between the goroutines.

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